

Miscellaneous FOL Examples

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1 Natural numbers

```
theory Natural-Numbers
imports FOL
begin
```

Theory of the natural numbers: Peano's axioms, primitive recursion. (Modernized version of Larry Paulson's theory "Nat".)

```
typedecl nat
instance nat :: <term> ..
```

axiomatization

```
Zero :: <nat>    (<0>) and
Suc  :: <nat => nat> and
rec  :: <[nat, 'a, [nat, 'a] => 'a] => 'a>
where
induct [case-names 0 Suc, induct type: nat]:
  <P(0) ==> (!x. P(x) ==> P(Suc(x))) ==> P(n)> and
Suc-inject: <Suc(m) = Suc(n) ==> m = n> and
Suc-neq-0: <Suc(m) = 0 ==> R> and
rec-0: <rec(0, a, f) = a> and
```

```

rec-Suc: ⟨rec(Suc(m), a, f) = f(m, rec(m, a, f))⟩

lemma Suc-n-not-n: ⟨Suc(k) ≠ k⟩
proof (induct ⟨k⟩)
  show ⟨Suc(0) ≠ 0⟩
  proof
    assume ⟨Suc(0) = 0⟩
    then show ⟨False⟩ by (rule Suc-neq-0)
  qed
next
  fix n assume hyp: ⟨Suc(n) ≠ n⟩
  show ⟨Suc(Suc(n)) ≠ Suc(n)⟩
  proof
    assume ⟨Suc(Suc(n)) = Suc(n)⟩
    then have ⟨Suc(n) = n⟩ by (rule Suc-inject)
    with hyp show ⟨False⟩ by contradiction
  qed
qed

definition add :: ⟨nat => nat => nat⟩ (infixl ⟨+⟩ 60)
  where ⟨m + n = rec(m, n, λx y. Suc(y))⟩

lemma add-0 [simp]: ⟨0 + n = n⟩
  unfolding add-def by (rule rec-0)

lemma add-Suc [simp]: ⟨Suc(m) + n = Suc(m + n)⟩
  unfolding add-def by (rule rec-Suc)

lemma add-assoc: ⟨(k + m) + n = k + (m + n)⟩
  by (induct ⟨k⟩) simp-all

lemma add-0-right: ⟨m + 0 = m⟩
  by (induct ⟨m⟩) simp-all

lemma add-Suc-right: ⟨m + Suc(n) = Suc(m + n)⟩
  by (induct ⟨m⟩) simp-all

lemma
  assumes ⟨!!n. f(Suc(n)) = Suc(f(n))⟩
  shows ⟨f(i + j) = i + f(j)⟩
  using assms by (induct ⟨i⟩) simp-all

end

```

2 Examples for the manual “Introduction to Isabelle”

```
theory Intro
imports FOL
begin
```

2.0.1 Some simple backward proofs

```
lemma mythm:  $\langle P \vee P \longrightarrow P \rangle$ 
apply (rule impI)
apply (rule disjE)
prefer 3 apply (assumption)
prefer 2 apply (assumption)
apply assumption
done
```

```
lemma  $\langle (P \wedge Q) \vee R \longrightarrow (P \vee R) \rangle$ 
apply (rule impI)
apply (erule disjE)
apply (drule conjunct1)
apply (rule disjI1)
apply (rule-tac [2] disjI2)
apply assumption+
done
```

Correct version, delaying use of *spec* until last.

```
lemma  $\langle (\forall x y. P(x,y)) \longrightarrow (\forall z w. P(w,z)) \rangle$ 
apply (rule impI)
apply (rule allI)
apply (rule allI)
apply (drule spec)
apply (drule spec)
apply assumption
done
```

2.0.2 Demonstration of *fast*

```
lemma  $\langle (\exists y. \forall x. J(y,x) \longleftrightarrow \neg J(x,x)) \longrightarrow \neg (\forall x. \exists y. \forall z. J(z,y) \longleftrightarrow \neg J(z,x)) \rangle$ 
apply fast
done
```

```
lemma  $\langle \forall x. P(x,f(x)) \longleftrightarrow (\exists y. (\forall z. P(z,y) \longrightarrow P(z,f(x))) \wedge P(x,y)) \rangle$ 
apply fast
done
```

2.0.3 Derivation of conjunction elimination rule

```
lemma
```

```

assumes major:  $\langle P \wedge Q \rangle$ 
  and minor:  $\langle [P; Q] \implies R \rangle$ 
shows  $\langle R \rangle$ 
apply (rule minor)
apply (rule major [THEN conjunct1])
apply (rule major [THEN conjunct2])
done

```

2.1 Derived rules involving definitions

Derivation of negation introduction

```

lemma
  assumes  $\langle P \implies False \rangle$ 
  shows  $\langle \neg P \rangle$ 
apply (unfold not-def)
apply (rule impI)
apply (rule assms)
apply assumption
done

```

```

lemma
  assumes major:  $\langle \neg P \rangle$ 
  and minor:  $\langle P \rangle$ 
shows  $\langle R \rangle$ 
apply (rule FalseE)
apply (rule mp)
apply (rule major [unfolded not-def])
apply (rule minor)
done

```

Alternative proof of the result above

```

lemma
  assumes major:  $\langle \neg P \rangle$ 
  and minor:  $\langle P \rangle$ 
shows  $\langle R \rangle$ 
apply (rule minor [THEN major [unfolded not-def, THEN mp, THEN FalseE]])
done

```

end

3 Theory of the natural numbers: Peano's axioms, primitive recursion

```

theory Nat
  imports FOL
begin

```

```

typedecl nat
instance nat :: ⟨term⟩ ..

```

axiomatization

```

Zero :: ⟨nat⟩ (⟨0⟩) and
Suc :: ⟨nat ⇒ nat⟩ and
rec :: ⟨[nat, 'a, [nat, 'a] ⇒ 'a] ⇒ 'a⟩
where
induct: ⟨[P(0); ∧x. P(x) ⇒ P(Suc(x))] ⇒ P(n)⟩ and
Suc-inject: ⟨Suc(m)=Suc(n) ⇒ m=n⟩ and
Suc-neq-0: ⟨Suc(m)=0 ⇒ R⟩ and
rec-0: ⟨rec(0,a,f) = a⟩ and
rec-Suc: ⟨rec(Suc(m), a, f) = f(m, rec(m,a,f))⟩

```

```

definition add :: ⟨[nat, nat] ⇒ nat⟩ (infixl ⟨+⟩ 60)
where ⟨m + n ≡ rec(m, n, λx y. Suc(y))⟩

```

3.1 Proofs about the natural numbers

```

lemma Suc-n-not-n: ⟨Suc(k) ≠ k⟩
apply (rule-tac n = ⟨k⟩ in induct)
apply (rule notI)
apply (erule Suc-neq-0)
apply (rule notI)
apply (erule notE)
apply (erule Suc-inject)
done

```

```

lemma ⟨(k+m)+n = k+(m+n)⟩
apply (rule induct)
back
back
back
back
back
back
oops

```

```

lemma add-0 [simp]: ⟨0+n = n⟩
apply (unfold add-def)
apply (rule rec-0)
done

```

```

lemma add-Suc [simp]: ⟨Suc(m)+n = Suc(m+n)⟩
apply (unfold add-def)
apply (rule rec-Suc)
done

```

```

lemma add-assoc: ⟨(k+m)+n = k+(m+n)⟩

```

```

apply (rule-tac n = ⟨k⟩ in induct)
apply simp
apply simp
done

```

```

lemma add-0-right: ⟨m+0 = m⟩
apply (rule-tac n = ⟨m⟩ in induct)
apply simp
apply simp
done

```

```

lemma add-Suc-right: ⟨m+Suc(n) = Suc(m+n)⟩
apply (rule-tac n = ⟨m⟩ in induct)
apply simp-all
done

```

```

lemma
  assumes prem: ⟨ $\bigwedge n. f(\text{Suc}(n)) = \text{Suc}(f(n))$ ⟩
  shows ⟨ $f(i+j) = i+f(j)$ ⟩
apply (rule-tac n = ⟨i⟩ in induct)
apply simp
apply (simp add: prem)
done

```

end

4 Theory of the natural numbers: Peano's axioms, primitive recursion

```

theory Nat-Class
  imports FOL
begin

```

This is an abstract version of `Nat.thy`. Instead of axiomatizing a single type `nat`, it defines the class of all these types (up to isomorphism).

Note: The `rec` operator has been made *monomorphic*, because class axioms cannot contain more than one type variable.

```

class nat =
  fixes Zero :: 'a → ⟨0⟩
    and Suc :: 'a ⇒ 'a
    and rec :: 'a ⇒ 'a ⇒ ('a ⇒ 'a ⇒ 'a) ⇒ 'a
  assumes induct: ⟨ $P(0) \implies (\bigwedge x. P(x) \implies P(\text{Suc}(x))) \implies P(n)$ ⟩
    and Suc-inject: ⟨ $\text{Suc}(m) = \text{Suc}(n) \implies m = n$ ⟩
    and Suc-neq-Zero: ⟨ $\text{Suc}(m) = 0 \implies R$ ⟩
    and rec-Zero: ⟨ $\text{rec}(0, a, f) = a$ ⟩
    and rec-Suc: ⟨ $\text{rec}(\text{Suc}(m), a, f) = f(m, \text{rec}(m, a, f))$ ⟩
begin

```

definition $add :: \langle 'a \Rightarrow 'a \Rightarrow 'a \rangle$ (**infixl** $\langle + \rangle$ 60)
where $\langle m + n = rec(m, n, \lambda x y. Suc(y)) \rangle$

lemma $Suc\text{-}n\text{-not-}n: \langle Suc(k) \neq (k::'a) \rangle$
apply (rule-tac $n = \langle k \rangle$ in induct)
apply (rule notI)
apply (erule Suc-neq-Zero)
apply (rule notI)
apply (erule notE)
apply (erule Suc-inject)
done

lemma $\langle (k + m) + n = k + (m + n) \rangle$
apply (rule induct)
back
back
back
back
back
back
oops

lemma $add\text{-Zero}$ [simp]: $\langle 0 + n = n \rangle$
apply (unfold add-def)
apply (rule rec-Zero)
done

lemma $add\text{-Suc}$ [simp]: $\langle Suc(m) + n = Suc(m + n) \rangle$
apply (unfold add-def)
apply (rule rec-Suc)
done

lemma $add\text{-assoc}$: $\langle (k + m) + n = k + (m + n) \rangle$
apply (rule-tac $n = \langle k \rangle$ in induct)
apply simp
apply simp
done

lemma $add\text{-Zero-right}$: $\langle m + 0 = m \rangle$
apply (rule-tac $n = \langle m \rangle$ in induct)
apply simp
apply simp
done

lemma $add\text{-Suc-right}$: $\langle m + Suc(n) = Suc(m + n) \rangle$
apply (rule-tac $n = \langle m \rangle$ in induct)
apply simp-all
done

lemma


```

assumes prem:  $\langle \bigwedge n. f(\text{Suc}(n)) = \text{Suc}(f(n)) \rangle$ 
shows  $\langle f(i + j) = i + f(j) \rangle$ 
apply (rule-tac  $n = \langle i \rangle$  in induct)
  apply simp
apply (simp add: prem)
done

```

end

end

5 Intuitionistic FOL: Examples from The Foundation of a Generic Theorem Prover

```

theory Foundation
imports IFOL
begin

```

```

lemma  $\langle A \wedge B \longrightarrow (C \longrightarrow A \wedge C) \rangle$ 
apply (rule impI)
apply (rule impI)
apply (rule conjI)
prefer 2 apply assumption
apply (rule conjunct1)
apply assumption
done

```

A form of conj-elimination

```

lemma
  assumes  $\langle A \wedge B \rangle$ 
  and  $\langle A \Longrightarrow B \Longrightarrow C \rangle$ 
  shows  $\langle C \rangle$ 
apply (rule assms)
apply (rule conjunct1)
apply (rule assms)
apply (rule conjunct2)
apply (rule assms)
done

```

```

lemma
  assumes  $\langle \bigwedge A. \neg \neg A \Longrightarrow A \rangle$ 
  shows  $\langle B \vee \neg B \rangle$ 
apply (rule assms)
apply (rule notI)
apply (rule-tac  $P = \langle \neg B \rangle$  in notE)
apply (rule-tac [2] notI)
apply (rule-tac [2]  $P = \langle B \vee \neg B \rangle$  in notE)
prefer 2 apply assumption

```

```

apply (rule-tac [2] disjI1)
prefer 2 apply assumption
apply (rule notI)
apply (rule-tac P = ⟨B ∨ ¬ B⟩ in notE)
apply assumption
apply (rule disjI2)
apply assumption
done

```

```

lemma
  assumes ⟨∧A. ¬ ¬ A ⇒ A⟩
  shows ⟨B ∨ ¬ B⟩
apply (rule assms)
apply (rule notI)
apply (rule notE)
apply (rule-tac [2] notI)
apply (erule-tac [2] notE)
apply (erule-tac [2] disjI1)
apply (rule notI)
apply (erule notE)
apply (erule disjI2)
done

```

```

lemma
  assumes ⟨A ∨ ¬ A⟩
  and ⟨¬ ¬ A⟩
  shows ⟨A⟩
apply (rule disjE)
apply (rule assms)
apply assumption
apply (rule FalseE)
apply (rule-tac P = ⟨¬ A⟩ in notE)
apply (rule assms)
apply assumption
done

```

5.1 Examples with quantifiers

```

lemma
  assumes ⟨∀z. G(z)⟩
  shows ⟨∀z. G(z) ∨ H(z)⟩
apply (rule allI)
apply (rule disjI1)
apply (rule assms [THEN spec])
done

```

```

lemma ⟨∀x. ∃y. x = y⟩
apply (rule allI)

```

```

apply (rule exI)
apply (rule refl)
done

```

```

lemma  $\langle \exists y. \forall x. x = y \rangle$ 
apply (rule exI)
apply (rule allI)
apply (rule refl)?
oops

```

Parallel lifting example.

```

lemma  $\langle \exists u. \forall x. \exists v. \forall y. \exists w. P(u,x,v,y,w) \rangle$ 
apply (rule exI allI)
apply (rule exI allI)
apply (rule exI allI)
apply (rule exI allI)
apply (rule exI allI)
apply (rule exI allI)
oops

```

```

lemma
  assumes  $\langle (\exists z. F(z)) \wedge B \rangle$ 
  shows  $\langle \exists z. F(z) \wedge B \rangle$ 
apply (rule conjE)
apply (rule assms)
apply (rule exE)
apply assumption
apply (rule exI)
apply (rule conjI)
apply assumption
apply assumption
done

```

A bigger demonstration of quantifiers – not in the paper.

```

lemma  $\langle (\exists y. \forall x. Q(x,y)) \longrightarrow (\forall x. \exists y. Q(x,y)) \rangle$ 
apply (rule impI)
apply (rule allI)
apply (rule exE, assumption)
apply (rule exI)
apply (rule allE, assumption)
apply assumption
done

```

end

6 First-Order Logic: PROLOG examples

```

theory Prolog
imports FOL

```

```

begin

typedecl 'a list
instance list :: (<term>) <term> ..

axiomatization
  Nil    :: <'a list> and
  Cons   :: <['a, 'a list]=> 'a list> (infixr <: 60) and
  app    :: <['a list, 'a list, 'a list] => o> and
  rev    :: <['a list, 'a list] => o>
where
  appNil: <app(Nil,ys,ys)> and
  appCons: <app(xs,ys,zs) ==> app(x:xs, ys, x:zs)> and
  revNil: <rev(Nil,Nil)> and
  revCons: <[| rev(xs,ys); app(ys, x:Nil, zs) |] ==> rev(x:xs, zs)>

schematic-goal <app(a:b:c:Nil, d:e:Nil, ?x)>
apply (rule appNil appCons)
apply (rule appNil appCons)
apply (rule appNil appCons)
apply (rule appNil appCons)
done

schematic-goal <app(?x, c:d:Nil, a:b:c:d:Nil)>
apply (rule appNil appCons)+
done

schematic-goal <app(?x, ?y, a:b:c:d:Nil)>
apply (rule appNil appCons)+
back
back
back
back
done

lemmas rules = appNil appCons revNil revCons

schematic-goal <rev(a:b:c:d:Nil, ?x)>
apply (rule rules)+
done

schematic-goal <rev(a:b:c:d:e:f:g:h:i:j:k:l:m:n:Nil, ?w)>
apply (rule rules)+
done

schematic-goal <rev(?x, a:b:c:Nil)>

```

```

apply (rule rules)+ — does not solve it directly!
back
back
done

```

```

ML <
  fun prolog-tac ctxt =
    DEPTH-FIRST (has-fewer-prems 1) (resolve-tac ctxt @{thms rules} 1)
  >

```

```

schematic-goal <rev(?x, a:b:c:Nil)>
apply (tactic <prolog-tac context>)
done

```

```

schematic-goal <rev(a:?x:c:?y:Nil, d:?z:b:?u)>
apply (tactic <prolog-tac context>)
done

```

```

schematic-goal <rev(a:b:c:d:e:f:g:h:i:j:k:l:m:n:o:p:Nil, ?w)>
apply (tactic <
  DEPTH-SOLVE (resolve-tac context ([@{thm refl}], @{thm conjI}]) @ @{thms
  rules}) 1)>)
done

```

```

schematic-goal <a:b:c:d:e:f:g:h:i:j:k:l:m:n:o:p:Nil = ?x ∧ app(?x, ?x, ?y) ∧ rev(?y, ?w)>
apply (tactic <
  DEPTH-SOLVE (resolve-tac context ([@{thm refl}], @{thm conjI}]) @ @{thms
  rules}) 1)>)
done

```

```

end

```

7 Intuitionistic First-Order Logic

```

theory Intuitionistic
imports IFOL
begin

```

Metatheorem (for *propositional* formulae): P is classically provable iff $\neg\neg P$ is intuitionistically provable. Therefore $\neg P$ is classically provable iff it is intuitionistically provable.

Proof: Let Q be the conjunction of the propositions $A \vee \neg A$, one for each atom A in P . Now $\neg\neg Q$ is intuitionistically provable because $\neg\neg(A \vee \neg A)$ is and because double-negation distributes over conjunction. If P is provable classically, then clearly $Q \rightarrow P$ is provable intuitionistically, so $\neg\neg(Q \rightarrow P)$

is also provable intuitionistically. The latter is intuitionistically equivalent to $\neg\neg Q \rightarrow \neg\neg P$, hence to $\neg\neg P$, since $\neg\neg Q$ is intuitionistically provable. Finally, if P is a negation then $\neg\neg P$ is intuitionistically equivalent to P . [Andy Pitts]

lemma $\langle \neg\neg (P \wedge Q) \longleftrightarrow \neg\neg P \wedge \neg\neg Q \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

lemma $\langle \neg\neg ((\neg P \rightarrow Q) \rightarrow (\neg P \rightarrow \neg Q) \rightarrow P) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

Double-negation does NOT distribute over disjunction.

lemma $\langle \neg\neg (P \rightarrow Q) \longleftrightarrow (\neg\neg P \rightarrow \neg\neg Q) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

lemma $\langle \neg\neg\neg P \longleftrightarrow \neg P \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

lemma $\langle \neg\neg ((P \rightarrow Q \vee R) \rightarrow (P \rightarrow Q) \vee (P \rightarrow R)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

lemma $\langle (P \longleftrightarrow Q) \longleftrightarrow (Q \longleftrightarrow P) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

lemma $\langle ((P \rightarrow (Q \vee (Q \rightarrow R))) \rightarrow R) \rightarrow R \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

lemma
 $\langle (((G \rightarrow A) \rightarrow J) \rightarrow D \rightarrow E) \rightarrow (((H \rightarrow B) \rightarrow I) \rightarrow C \rightarrow J)$
 $\rightarrow (A \rightarrow H) \rightarrow F \rightarrow G \rightarrow (((C \rightarrow B) \rightarrow I) \rightarrow D) \rightarrow (A \rightarrow C)$
 $\rightarrow (((F \rightarrow A) \rightarrow B) \rightarrow I) \rightarrow E \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

Admissibility of the excluded middle for negated formulae

lemma $\langle (P \vee \neg P \rightarrow \neg Q) \rightarrow \neg Q \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

The same in a more general form, no ex falso quodlibet

lemma $\langle (P \vee (P \rightarrow R) \rightarrow Q \rightarrow R) \rightarrow Q \rightarrow R \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

7.1 Lemmas for the propositional double-negation translation

lemma $\langle P \rightarrow \neg\neg P \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

lemma $\langle \neg\neg (\neg\neg P \rightarrow P) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

lemma $\langle \neg \neg P \wedge \neg \neg (P \longrightarrow Q) \longrightarrow \neg \neg Q \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

The following are classically but not constructively valid. The attempt to prove them terminates quickly!

lemma $\langle ((P \longrightarrow Q) \longrightarrow P) \longrightarrow P \rangle$
apply (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)?
apply (*rule asm-rl*) — Checks that subgoals remain: proof failed.
oops

lemma $\langle (P \wedge Q \longrightarrow R) \longrightarrow (P \longrightarrow R) \vee (Q \longrightarrow R) \rangle$
apply (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)?
apply (*rule asm-rl*) — Checks that subgoals remain: proof failed.
oops

7.2 de Bruijn formulae

de Bruijn formula with three predicates

lemma
 $\langle ((P \longleftrightarrow Q) \longrightarrow P \wedge Q \wedge R) \wedge$
 $((Q \longleftrightarrow R) \longrightarrow P \wedge Q \wedge R) \wedge$
 $((R \longleftrightarrow P) \longrightarrow P \wedge Q \wedge R) \longrightarrow P \wedge Q \wedge R \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

de Bruijn formula with five predicates

lemma
 $\langle ((P \longleftrightarrow Q) \longrightarrow P \wedge Q \wedge R \wedge S \wedge T) \wedge$
 $((Q \longleftrightarrow R) \longrightarrow P \wedge Q \wedge R \wedge S \wedge T) \wedge$
 $((R \longleftrightarrow S) \longrightarrow P \wedge Q \wedge R \wedge S \wedge T) \wedge$
 $((S \longleftrightarrow T) \longrightarrow P \wedge Q \wedge R \wedge S \wedge T) \wedge$
 $((T \longleftrightarrow P) \longrightarrow P \wedge Q \wedge R \wedge S \wedge T) \longrightarrow P \wedge Q \wedge R \wedge S \wedge T \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

Problems from of Sahlin, Franzen and Haridi, An Intuitionistic Predicate Logic Theorem Prover. J. Logic and Comp. 2 (5), October 1992, 619-656.

Problem 1.1

lemma
 $\langle (\forall x. \exists y. \forall z. p(x) \wedge q(y) \wedge r(z)) \longleftrightarrow$
 $(\forall z. \exists y. \forall x. p(x) \wedge q(y) \wedge r(z)) \rangle$
by (*tactic* $\langle \text{IntPr.best-dup-tac context 1} \rangle$) — SLOW

Problem 3.1

lemma $\langle \neg (\exists x. \forall y. \text{mem}(y,x) \longleftrightarrow \neg \text{mem}(x,x)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

Problem 4.1: hopeless!

lemma

$\langle (\forall x. p(x) \longrightarrow p(h(x)) \vee p(g(x))) \wedge (\exists x. p(x)) \wedge (\forall x. \neg p(h(x)))$
 $\longrightarrow (\exists x. p(g(g(g(g(x)))))) \rangle$

oops

7.3 Intuitionistic FOL: propositional problems based on Pelletier.

$\neg\neg 1$

lemma $\langle \neg\neg ((P \longrightarrow Q) \longleftrightarrow (\neg Q \longrightarrow \neg P)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

$\neg\neg 2$

lemma $\langle \neg\neg (\neg\neg P \longleftrightarrow P) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

3

lemma $\langle \neg (P \longrightarrow Q) \longrightarrow (Q \longrightarrow P) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

$\neg\neg 4$

lemma $\langle \neg\neg ((\neg P \longrightarrow Q) \longleftrightarrow (\neg Q \longrightarrow P)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

$\neg\neg 5$

lemma $\langle \neg\neg ((P \vee Q \longrightarrow P \vee R) \longrightarrow P \vee (Q \longrightarrow R)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

$\neg\neg 6$

lemma $\langle \neg\neg (P \vee \neg P) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

$\neg\neg 7$

lemma $\langle \neg\neg (P \vee \neg\neg\neg P) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

$\neg\neg 8$. Peirce's law

lemma $\langle \neg\neg (((P \longrightarrow Q) \longrightarrow P) \longrightarrow P) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

9

lemma $\langle ((P \vee Q) \wedge (\neg P \vee Q) \wedge (P \vee \neg Q)) \longrightarrow \neg(\neg P \vee \neg Q) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

10

lemma $\langle (Q \longrightarrow R) \longrightarrow (R \longrightarrow P \wedge Q) \longrightarrow (P \longrightarrow (Q \vee R)) \longrightarrow (P \longleftrightarrow Q) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

7.4 11. Proved in each direction (incorrectly, says Pelletier!!)

lemma $\langle P \longleftrightarrow P \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

$\neg\neg$ 12. Dijkstra's law

lemma $\langle \neg\neg ((P \longleftrightarrow Q) \longleftrightarrow R) \longleftrightarrow (P \longleftrightarrow (Q \longleftrightarrow R)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

lemma $\langle ((P \longleftrightarrow Q) \longleftrightarrow R) \longrightarrow \neg\neg (P \longleftrightarrow (Q \longleftrightarrow R)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

13. Distributive law

lemma $\langle P \vee (Q \wedge R) \longleftrightarrow (P \vee Q) \wedge (P \vee R) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

$\neg\neg$ 14

lemma $\langle \neg\neg ((P \longleftrightarrow Q) \longleftrightarrow ((Q \vee \neg P) \wedge (\neg Q \vee P))) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

$\neg\neg$ 15

lemma $\langle \neg\neg ((P \longrightarrow Q) \longleftrightarrow (\neg P \vee Q)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

$\neg\neg$ 16

lemma $\langle \neg\neg ((P \longrightarrow Q) \vee (Q \longrightarrow P)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

$\neg\neg$ 17

lemma $\langle \neg\neg (((P \wedge (Q \longrightarrow R)) \longrightarrow S) \longleftrightarrow ((\neg P \vee Q \vee S) \wedge (\neg P \vee \neg R \vee S))) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

Dijkstra's "Golden Rule"

lemma $\langle (P \wedge Q) \longleftrightarrow P \longleftrightarrow Q \longleftrightarrow (P \vee Q) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

8 Examples with quantifiers

8.1 The converse is classical in the following implications ...

lemma $\langle (\exists x. P(x) \longrightarrow Q) \longrightarrow (\forall x. P(x)) \longrightarrow Q \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

lemma $\langle ((\forall x. P(x)) \longrightarrow Q) \longrightarrow \neg (\forall x. P(x) \wedge \neg Q) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

lemma $\langle ((\forall x. \neg P(x)) \longrightarrow Q) \longrightarrow \neg (\forall x. \neg (P(x) \vee Q)) \rangle$

by (tactic <IntPr.fast-tac **context** 1>)

lemma $\langle (\forall x. P(x)) \vee Q \longrightarrow (\forall x. P(x) \vee Q) \rangle$
by (tactic <IntPr.fast-tac **context** 1>)

lemma $\langle (\exists x. P \longrightarrow Q(x)) \longrightarrow (P \longrightarrow (\exists x. Q(x))) \rangle$
by (tactic <IntPr.fast-tac **context** 1>)

8.2 The following are not constructively valid!

The attempt to prove them terminates quickly!

lemma $\langle ((\forall x. P(x)) \longrightarrow Q) \longrightarrow (\exists x. P(x) \longrightarrow Q) \rangle$
apply (tactic <IntPr.fast-tac **context** 1>)?
apply (rule asm-rl) — Checks that subgoals remain: proof failed.
oops

lemma $\langle (P \longrightarrow (\exists x. Q(x))) \longrightarrow (\exists x. P \longrightarrow Q(x)) \rangle$
apply (tactic <IntPr.fast-tac **context** 1>)?
apply (rule asm-rl) — Checks that subgoals remain: proof failed.
oops

lemma $\langle (\forall x. P(x) \vee Q) \longrightarrow ((\forall x. P(x)) \vee Q) \rangle$
apply (tactic <IntPr.fast-tac **context** 1>)?
apply (rule asm-rl) — Checks that subgoals remain: proof failed.
oops

lemma $\langle (\forall x. \neg \neg P(x)) \longrightarrow \neg \neg (\forall x. P(x)) \rangle$
apply (tactic <IntPr.fast-tac **context** 1>)?
apply (rule asm-rl) — Checks that subgoals remain: proof failed.
oops

Classically but not intuitionistically valid. Proved by a bug in 1986!

lemma $\langle \exists x. Q(x) \longrightarrow (\forall x. Q(x)) \rangle$
apply (tactic <IntPr.fast-tac **context** 1>)?
apply (rule asm-rl) — Checks that subgoals remain: proof failed.
oops

8.3 Hard examples with quantifiers

The ones that have not been proved are not known to be valid! Some will require quantifier duplication – not currently available.

$\neg\neg 18$

lemma $\langle \neg \neg (\exists y. \forall x. P(y) \longrightarrow P(x)) \rangle$
oops — NOT PROVED

$\neg\neg 19$

lemma $\langle \neg \neg (\exists x. \forall y z. (P(y) \longrightarrow Q(z)) \longrightarrow (P(x) \longrightarrow Q(x))) \rangle$

oops — NOT PROVED

20

lemma

$\langle (\forall x y. \exists z. \forall w. (P(x) \wedge Q(y) \longrightarrow R(z) \wedge S(w)))$
 $\longrightarrow (\exists x y. P(x) \wedge Q(y)) \longrightarrow (\exists z. R(z)) \rangle$

by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

21

lemma $\langle (\exists x. P \longrightarrow Q(x)) \wedge (\exists x. Q(x) \longrightarrow P) \longrightarrow \neg \neg (\exists x. P \longleftrightarrow Q(x)) \rangle$

oops — NOT PROVED; needs quantifier duplication

22

lemma $\langle (\forall x. P \longleftrightarrow Q(x)) \longrightarrow (P \longleftrightarrow (\forall x. Q(x))) \rangle$

by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

$\neg\neg$ 23

lemma $\langle \neg \neg ((\forall x. P \vee Q(x)) \longleftrightarrow (P \vee (\forall x. Q(x)))) \rangle$

by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

24

lemma

$\langle \neg (\exists x. S(x) \wedge Q(x)) \wedge (\forall x. P(x) \longrightarrow Q(x) \vee R(x)) \wedge$
 $(\neg (\exists x. P(x)) \longrightarrow (\exists x. Q(x))) \wedge (\forall x. Q(x) \vee R(x) \longrightarrow S(x))$
 $\longrightarrow \neg \neg (\exists x. P(x) \wedge R(x)) \rangle$

Not clear why *fast-tac*, *best-tac*, *ASTAR* and *ITER-DEEPEN* all take forever.

apply (*tactic* $\langle \text{IntPr.safe-tac context} \rangle$)

apply (*erule impE*)

apply (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

apply (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

done

25

lemma

$\langle (\exists x. P(x)) \wedge$
 $(\forall x. L(x) \longrightarrow \neg (M(x) \wedge R(x))) \wedge$
 $(\forall x. P(x) \longrightarrow (M(x) \wedge L(x))) \wedge$
 $((\forall x. P(x) \longrightarrow Q(x)) \vee (\exists x. P(x) \wedge R(x)))$
 $\longrightarrow (\exists x. Q(x) \wedge P(x)) \rangle$

by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

$\neg\neg$ 26

lemma

$\langle (\neg \neg (\exists x. p(x)) \longleftrightarrow \neg \neg (\exists x. q(x))) \wedge$
 $(\forall x. \forall y. p(x) \wedge q(y) \longrightarrow (r(x) \longleftrightarrow s(y))) \rangle$

$\longrightarrow ((\forall x. p(x) \longrightarrow r(x)) \longleftrightarrow (\forall x. q(x) \longrightarrow s(x)))\rangle$
oops — NOT PROVED

27

lemma

$\langle (\exists x. P(x) \wedge \neg Q(x)) \wedge$
 $(\forall x. P(x) \longrightarrow R(x)) \wedge$
 $(\forall x. M(x) \wedge L(x) \longrightarrow P(x)) \wedge$
 $((\exists x. R(x) \wedge \neg Q(x)) \longrightarrow (\forall x. L(x) \longrightarrow \neg R(x)))$
 $\longrightarrow (\forall x. M(x) \longrightarrow \neg L(x))\rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

$\neg\neg$ 28. AMENDED

lemma

$\langle (\forall x. P(x) \longrightarrow (\forall x. Q(x))) \wedge$
 $(\neg\neg (\forall x. Q(x) \vee R(x)) \longrightarrow (\exists x. Q(x) \wedge S(x))) \wedge$
 $(\neg\neg (\exists x. S(x)) \longrightarrow (\forall x. L(x) \longrightarrow M(x)))$
 $\longrightarrow (\forall x. P(x) \wedge L(x) \longrightarrow M(x))\rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

29. Essentially the same as Principia Mathematica *11.71

lemma

$\langle (\exists x. P(x)) \wedge (\exists y. Q(y))$
 $\longrightarrow ((\forall x. P(x) \longrightarrow R(x)) \wedge (\forall y. Q(y) \longrightarrow S(y)) \longleftrightarrow$
 $(\forall x y. P(x) \wedge Q(y) \longrightarrow R(x) \wedge S(y)))\rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

$\neg\neg$ 30

lemma

$\langle (\forall x. (P(x) \vee Q(x)) \longrightarrow \neg R(x)) \wedge$
 $(\forall x. (Q(x) \longrightarrow \neg S(x)) \longrightarrow P(x) \wedge R(x))$
 $\longrightarrow (\forall x. \neg\neg S(x))\rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

31

lemma

$\langle \neg (\exists x. P(x) \wedge (Q(x) \vee R(x))) \wedge$
 $(\exists x. L(x) \wedge P(x)) \wedge$
 $(\forall x. \neg R(x) \longrightarrow M(x))$
 $\longrightarrow (\exists x. L(x) \wedge M(x))\rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

32

lemma

$\langle (\forall x. P(x) \wedge (Q(x) \vee R(x)) \longrightarrow S(x)) \wedge$
 $(\forall x. S(x) \wedge R(x) \longrightarrow L(x)) \wedge$
 $(\forall x. M(x) \longrightarrow R(x))$
 $\longrightarrow (\forall x. P(x) \wedge M(x) \longrightarrow L(x))\rangle$

by (tactic <IntPr.fast-tac context 1>)

¬¬33

lemma

⟨(∀ x. ¬ ¬ (P(a) ∧ (P(x) → P(b)) → P(c))) ↔
(∀ x. ¬ ¬ ((¬ P(a) ∨ P(x) ∨ P(c)) ∧ (¬ P(a) ∨ ¬ P(b) ∨ P(c))))⟩
apply (tactic <IntPr.best-tac context 1>)
done

36

lemma

⟨(∀ x. ∃ y. J(x,y)) ∧
(∀ x. ∃ y. G(x,y)) ∧
(∀ x y. J(x,y) ∨ G(x,y) → (∀ z. J(y,z) ∨ G(y,z) → H(x,z)))
→ (∀ x. ∃ y. H(x,y))⟩
by (tactic <IntPr.fast-tac context 1>)

37

lemma

⟨(∀ z. ∃ w. ∀ x. ∃ y.
¬ ¬ (P(x,z) → P(y,w)) ∧ P(y,z) ∧ (P(y,w) → (∃ u. Q(u,w)))) ∧
(∀ x z. ¬ P(x,z) → (∃ y. Q(y,z))) ∧
(¬ ¬ (∃ x y. Q(x,y)) → (∀ x. R(x,x)))
→ ¬ ¬ (∀ x. ∃ y. R(x,y))⟩
oops — NOT PROVED

39

lemma <¬ (∃ x. ∀ y. F(y,x) ↔ ¬ F(y,y))⟩
by (tactic <IntPr.fast-tac context 1>)

40. AMENDED

lemma

⟨(∃ y. ∀ x. F(x,y) ↔ F(x,x)) →
¬ (∀ x. ∃ y. ∀ z. F(z,y) ↔ ¬ F(z,x))⟩
by (tactic <IntPr.fast-tac context 1>)

44

lemma

⟨(∀ x. f(x) →
(∃ y. g(y) ∧ h(x,y) ∧ (∃ y. g(y) ∧ ¬ h(x,y)))) ∧
(∃ x. j(x) ∧ (∀ y. g(y) → h(x,y)))
→ (∃ x. j(x) ∧ ¬ f(x))⟩
by (tactic <IntPr.fast-tac context 1>)

48

lemma <(a = b ∨ c = d) ∧ (a = c ∨ b = d) → a = d ∨ b = c>
by (tactic <IntPr.fast-tac context 1>)

51

lemma

$\langle (\exists z w. \forall x y. P(x,y) \longleftrightarrow (x = z \wedge y = w)) \longrightarrow$
 $(\exists z. \forall x. \exists w. (\forall y. P(x,y) \longleftrightarrow y = w) \longleftrightarrow x = z) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

52

Almost the same as 51.

lemma

$\langle (\exists z w. \forall x y. P(x,y) \longleftrightarrow (x = z \wedge y = w)) \longrightarrow$
 $(\exists w. \forall y. \exists z. (\forall x. P(x,y) \longleftrightarrow x = z) \longleftrightarrow y = w) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

56

lemma $\langle (\forall x. (\exists y. P(y) \wedge x = f(y)) \longrightarrow P(x)) \longleftrightarrow (\forall x. P(x) \longrightarrow P(f(x))) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

57

lemma

$\langle P(f(a,b), f(b,c)) \wedge P(f(b,c), f(a,c)) \wedge$
 $(\forall x y z. P(x,y) \wedge P(y,z) \longrightarrow P(x,z)) \longrightarrow P(f(a,b), f(a,c)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

60

lemma $\langle \forall x. P(x, f(x)) \longleftrightarrow (\exists y. (\forall z. P(z, y) \longrightarrow P(z, f(x))) \wedge P(x, y)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

end

9 First-Order Logic: propositional examples (intuitionistic version)

theory *Propositional-Int*

imports *IFOL*

begin

commutative laws of \wedge and \vee

lemma $\langle P \wedge Q \longrightarrow Q \wedge P \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

lemma $\langle P \vee Q \longrightarrow Q \vee P \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context 1} \rangle$)

associative laws of \wedge and \vee

lemma $\langle (P \wedge Q) \wedge R \longrightarrow P \wedge (Q \wedge R) \rangle$

by (tactic IntPr.fast-tac context 1)

lemma $\langle (P \vee Q) \vee R \longrightarrow P \vee (Q \vee R) \rangle$
by (tactic IntPr.fast-tac context 1)

distributive laws of \wedge and \vee

lemma $\langle (P \wedge Q) \vee R \longrightarrow (P \vee R) \wedge (Q \vee R) \rangle$
by (tactic IntPr.fast-tac context 1)

lemma $\langle (P \vee R) \wedge (Q \vee R) \longrightarrow (P \wedge Q) \vee R \rangle$
by (tactic IntPr.fast-tac context 1)

lemma $\langle (P \vee Q) \wedge R \longrightarrow (P \wedge R) \vee (Q \wedge R) \rangle$
by (tactic IntPr.fast-tac context 1)

lemma $\langle (P \wedge R) \vee (Q \wedge R) \longrightarrow (P \vee Q) \wedge R \rangle$
by (tactic IntPr.fast-tac context 1)

Laws involving implication

lemma $\langle (P \longrightarrow R) \wedge (Q \longrightarrow R) \longleftrightarrow (P \vee Q \longrightarrow R) \rangle$
by (tactic IntPr.fast-tac context 1)

lemma $\langle (P \wedge Q \longrightarrow R) \longleftrightarrow (P \longrightarrow (Q \longrightarrow R)) \rangle$
by (tactic IntPr.fast-tac context 1)

lemma $\langle ((P \longrightarrow R) \longrightarrow R) \longrightarrow ((Q \longrightarrow R) \longrightarrow R) \longrightarrow (P \wedge Q \longrightarrow R) \longrightarrow R \rangle$
by (tactic IntPr.fast-tac context 1)

lemma $\langle \neg (P \longrightarrow R) \longrightarrow \neg (Q \longrightarrow R) \longrightarrow \neg (P \wedge Q \longrightarrow R) \rangle$
by (tactic IntPr.fast-tac context 1)

lemma $\langle (P \longrightarrow Q \wedge R) \longleftrightarrow (P \longrightarrow Q) \wedge (P \longrightarrow R) \rangle$
by (tactic IntPr.fast-tac context 1)

Propositions-as-types

lemma $\langle P \longrightarrow (Q \longrightarrow P) \rangle$
by (tactic IntPr.fast-tac context 1)

— The combinator S

lemma $\langle (P \longrightarrow Q \longrightarrow R) \longrightarrow (P \longrightarrow Q) \longrightarrow (P \longrightarrow R) \rangle$
by (tactic IntPr.fast-tac context 1)

— Converse is classical

lemma $\langle (P \longrightarrow Q) \vee (P \longrightarrow R) \longrightarrow (P \longrightarrow Q \vee R) \rangle$
by (tactic IntPr.fast-tac context 1)

lemma $\langle (P \longrightarrow Q) \longrightarrow (\neg Q \longrightarrow \neg P) \rangle$
by (tactic IntPr.fast-tac context 1)

Schwichtenberg's examples (via T. Nipkow)

lemma *stab-imp*: $\langle (((Q \rightarrow R) \rightarrow R) \rightarrow Q) \rightarrow (((P \rightarrow Q) \rightarrow R) \rightarrow R) \rightarrow P \rightarrow Q \rangle$

by (*tactic IntPr.fast-tac context 1*)

lemma *stab-to-peirce*:

$\langle (((P \rightarrow R) \rightarrow R) \rightarrow P) \rightarrow (((Q \rightarrow R) \rightarrow R) \rightarrow Q) \rightarrow ((P \rightarrow Q) \rightarrow P) \rightarrow P \rangle$

by (*tactic IntPr.fast-tac context 1*)

lemma *peirce-imp1*:

$\langle (((Q \rightarrow R) \rightarrow Q) \rightarrow Q) \rightarrow (((P \rightarrow Q) \rightarrow R) \rightarrow P \rightarrow Q) \rightarrow P \rightarrow Q \rangle$

by (*tactic IntPr.fast-tac context 1*)

lemma *peirce-imp2*: $\langle (((P \rightarrow R) \rightarrow P) \rightarrow P) \rightarrow ((P \rightarrow Q \rightarrow R) \rightarrow P) \rightarrow P \rangle$

by (*tactic IntPr.fast-tac context 1*)

lemma *mints*: $\langle (((P \rightarrow Q) \rightarrow P) \rightarrow P) \rightarrow Q \rightarrow Q \rangle$

by (*tactic IntPr.fast-tac context 1*)

lemma *mints-solovev*: $\langle (P \rightarrow (Q \rightarrow R) \rightarrow Q) \rightarrow ((P \rightarrow Q) \rightarrow R) \rightarrow R \rangle$

by (*tactic IntPr.fast-tac context 1*)

lemma *tatsuta*:

$\langle (((P7 \rightarrow P1) \rightarrow P10) \rightarrow P4 \rightarrow P5) \rightarrow (((P8 \rightarrow P2) \rightarrow P9) \rightarrow P3 \rightarrow P10) \rightarrow (P1 \rightarrow P8) \rightarrow P6 \rightarrow P7 \rightarrow (((P3 \rightarrow P2) \rightarrow P9) \rightarrow P4) \rightarrow (P1 \rightarrow P3) \rightarrow (((P6 \rightarrow P1) \rightarrow P2) \rightarrow P9) \rightarrow P5 \rangle$

by (*tactic IntPr.fast-tac context 1*)

lemma *tatsuta1*:

$\langle (((P8 \rightarrow P2) \rightarrow P9) \rightarrow P3 \rightarrow P10) \rightarrow (((P3 \rightarrow P2) \rightarrow P9) \rightarrow P4) \rightarrow (((P6 \rightarrow P1) \rightarrow P2) \rightarrow P9) \rightarrow (((P7 \rightarrow P1) \rightarrow P10) \rightarrow P4 \rightarrow P5) \rightarrow (P1 \rightarrow P3) \rightarrow (P1 \rightarrow P8) \rightarrow P6 \rightarrow P7 \rightarrow P5 \rangle$

by (*tactic IntPr.fast-tac context 1*)

end

10 First-Order Logic: quantifier examples (intuitionistic version)

theory *Quantifiers-Int*
imports *IFOL*

begin

lemma $\langle (\forall x y. P(x,y)) \longrightarrow (\forall y x. P(x,y)) \rangle$
by (*tactic IntPr.fast-tac context 1*)

lemma $\langle (\exists x y. P(x,y)) \longrightarrow (\exists y x. P(x,y)) \rangle$
by (*tactic IntPr.fast-tac context 1*)

— Converse is false

lemma $\langle (\forall x. P(x)) \vee (\forall x. Q(x)) \longrightarrow (\forall x. P(x) \vee Q(x)) \rangle$
by (*tactic IntPr.fast-tac context 1*)

lemma $\langle (\forall x. P \longrightarrow Q(x)) \longleftrightarrow (P \longrightarrow (\forall x. Q(x))) \rangle$
by (*tactic IntPr.fast-tac context 1*)

lemma $\langle (\forall x. P(x) \longrightarrow Q) \longleftrightarrow ((\exists x. P(x)) \longrightarrow Q) \rangle$
by (*tactic IntPr.fast-tac context 1*)

Some harder ones

lemma $\langle (\exists x. P(x) \vee Q(x)) \longleftrightarrow (\exists x. P(x)) \vee (\exists x. Q(x)) \rangle$
by (*tactic IntPr.fast-tac context 1*)

— Converse is false

lemma $\langle (\exists x. P(x) \wedge Q(x)) \longrightarrow (\exists x. P(x)) \wedge (\exists x. Q(x)) \rangle$
by (*tactic IntPr.fast-tac context 1*)

Basic test of quantifier reasoning

lemma $\langle (\exists y. \forall x. Q(x,y)) \longrightarrow (\forall x. \exists y. Q(x,y)) \rangle$
by (*tactic IntPr.fast-tac context 1*)

lemma $\langle (\forall x. Q(x)) \longrightarrow (\exists x. Q(x)) \rangle$
by (*tactic IntPr.fast-tac context 1*)

The following should fail, as they are false!

lemma $\langle (\forall x. \exists y. Q(x,y)) \longrightarrow (\exists y. \forall x. Q(x,y)) \rangle$
apply (*tactic IntPr.fast-tac context 1*)?
oops

lemma $\langle (\exists x. Q(x)) \longrightarrow (\forall x. Q(x)) \rangle$
apply (*tactic IntPr.fast-tac context 1*)?
oops

schematic-goal $\langle P(?a) \longrightarrow (\forall x. P(x)) \rangle$
apply (*tactic IntPr.fast-tac context 1*)?
oops

schematic-goal $\langle (P(?a) \longrightarrow (\forall x. Q(x))) \longrightarrow (\forall x. P(x) \longrightarrow Q(x)) \rangle$

apply (*tactic IntPr.fast-tac context 1*)?
oops

Back to things that are provable ...

lemma $\langle (\forall x. P(x) \longrightarrow Q(x)) \wedge (\exists x. P(x)) \longrightarrow (\exists x. Q(x)) \rangle$
by (*tactic IntPr.fast-tac context 1*)

— An example of why exI should be delayed as long as possible

lemma $\langle (P \longrightarrow (\exists x. Q(x))) \wedge P \longrightarrow (\exists x. Q(x)) \rangle$
by (*tactic IntPr.fast-tac context 1*)

schematic-goal $\langle (\forall x. P(x) \longrightarrow Q(f(x))) \wedge (\forall x. Q(x) \longrightarrow R(g(x))) \wedge P(d) \longrightarrow R(?a) \rangle$
by (*tactic IntPr.fast-tac context 1*)

lemma $\langle (\forall x. Q(x)) \longrightarrow (\exists x. Q(x)) \rangle$
by (*tactic IntPr.fast-tac context 1*)

Some slow ones

lemma $\langle (\forall x y. P(x) \longrightarrow Q(y)) \longleftrightarrow ((\exists x. P(x)) \longrightarrow (\forall y. Q(y))) \rangle$
by (*tactic IntPr.fast-tac context 1*)

lemma $\langle (\exists x y. P(x) \wedge Q(x,y)) \longleftrightarrow (\exists x. P(x) \wedge (\exists y. Q(x,y))) \rangle$
by (*tactic IntPr.fast-tac context 1*)

lemma $\langle (\exists y. \forall x. P(x) \longrightarrow Q(x,y)) \longrightarrow (\forall x. P(x) \longrightarrow (\exists y. Q(x,y))) \rangle$
by (*tactic IntPr.fast-tac context 1*)

end

11 Classical Predicate Calculus Problems

theory *Classical*
imports *FOL*
begin

lemma $\langle (P \longrightarrow Q \vee R) \longrightarrow (P \longrightarrow Q) \vee (P \longrightarrow R) \rangle$
by *blast*

11.0.1 If and only if

lemma $\langle (P \longleftrightarrow Q) \longleftrightarrow (Q \longleftrightarrow P) \rangle$
by *blast*

lemma $\langle \neg (P \longleftrightarrow \neg P) \rangle$
by *blast*

11.1 Pelletier's examples

Sample problems from

- F. J. Pelletier, Seventy-Five Problems for Testing Automatic Theorem Provers, J. Automated Reasoning 2 (1986), 191-216. Errata, JAR 4 (1988), 236-236.

The hardest problems – judging by experience with several theorem provers, including matrix ones – are 34 and 43.

1

lemma $\langle (P \rightarrow Q) \leftrightarrow (\neg Q \rightarrow \neg P) \rangle$
by *blast*

2

lemma $\langle \neg \neg P \leftrightarrow P \rangle$
by *blast*

3

lemma $\langle \neg (P \rightarrow Q) \rightarrow (Q \rightarrow P) \rangle$
by *blast*

4

lemma $\langle (\neg P \rightarrow Q) \leftrightarrow (\neg Q \rightarrow P) \rangle$
by *blast*

5

lemma $\langle ((P \vee Q) \rightarrow (P \vee R)) \rightarrow (P \vee (Q \rightarrow R)) \rangle$
by *blast*

6

lemma $\langle P \vee \neg P \rangle$
by *blast*

7

lemma $\langle P \vee \neg \neg \neg P \rangle$
by *blast*

8. Peirce's law

lemma $\langle ((P \rightarrow Q) \rightarrow P) \rightarrow P \rangle$
by *blast*

9

lemma $\langle ((P \vee Q) \wedge (\neg P \vee Q) \wedge (P \vee \neg Q)) \rightarrow \neg (\neg P \vee \neg Q) \rangle$
by *blast*

10

lemma $\langle (Q \longrightarrow R) \wedge (R \longrightarrow P \wedge Q) \wedge (P \longrightarrow Q \vee R) \longrightarrow (P \longleftrightarrow Q) \rangle$
by *blast*

11. Proved in each direction (incorrectly, says Pelletier!!)

lemma $\langle P \longleftrightarrow P \rangle$
by *blast*

12. "Dijkstra's law"

lemma $\langle ((P \longleftrightarrow Q) \longleftrightarrow R) \longleftrightarrow (P \longleftrightarrow (Q \longleftrightarrow R)) \rangle$
by *blast*

13. Distributive law

lemma $\langle P \vee (Q \wedge R) \longleftrightarrow (P \vee Q) \wedge (P \vee R) \rangle$
by *blast*

14

lemma $\langle (P \longleftrightarrow Q) \longleftrightarrow ((Q \vee \neg P) \wedge (\neg Q \vee P)) \rangle$
by *blast*

15

lemma $\langle (P \longrightarrow Q) \longleftrightarrow (\neg P \vee Q) \rangle$
by *blast*

16

lemma $\langle (P \longrightarrow Q) \vee (Q \longrightarrow P) \rangle$
by *blast*

17

lemma $\langle ((P \wedge (Q \longrightarrow R)) \longrightarrow S) \longleftrightarrow ((\neg P \vee Q \vee S) \wedge (\neg P \vee \neg R \vee S)) \rangle$
by *blast*

11.2 Classical Logic: examples with quantifiers

lemma $\langle (\forall x. P(x) \wedge Q(x)) \longleftrightarrow (\forall x. P(x)) \wedge (\forall x. Q(x)) \rangle$
by *blast*

lemma $\langle (\exists x. P \longrightarrow Q(x)) \longleftrightarrow (P \longrightarrow (\exists x. Q(x))) \rangle$
by *blast*

lemma $\langle (\exists x. P(x) \longrightarrow Q) \longleftrightarrow (\forall x. P(x)) \longrightarrow Q \rangle$
by *blast*

lemma $\langle (\forall x. P(x)) \vee Q \longleftrightarrow (\forall x. P(x) \vee Q) \rangle$
by *blast*

Discussed in Avron, Gentzen-Type Systems, Resolution and Tableaux, JAR 10 (265-281), 1993. Proof is trivial!

lemma $\langle \neg ((\exists x. \neg P(x)) \wedge ((\exists x. P(x)) \vee (\exists x. P(x) \wedge Q(x))) \wedge \neg (\exists x. P(x))) \rangle$
by *blast*

11.3 Problems requiring quantifier duplication

Theorem B of Peter Andrews, Theorem Proving via General Matings, JACM 28 (1981).

lemma $\langle (\exists x. \forall y. P(x) \longleftrightarrow P(y)) \longrightarrow ((\exists x. P(x)) \longleftrightarrow (\forall y. P(y))) \rangle$
by *blast*

Needs multiple instantiation of ALL.

lemma $\langle (\forall x. P(x) \longrightarrow P(f(x))) \wedge P(d) \longrightarrow P(f(f(f(d)))) \rangle$
by *blast*

Needs double instantiation of the quantifier

lemma $\langle \exists x. P(x) \longrightarrow P(a) \wedge P(b) \rangle$
by *blast*

lemma $\langle \exists z. P(z) \longrightarrow (\forall x. P(x)) \rangle$
by *blast*

lemma $\langle \exists x. (\exists y. P(y)) \longrightarrow P(x) \rangle$
by *blast*

V. Lifschitz, What Is the Inverse Method?, JAR 5 (1989), 1–23. NOT PROVED.

lemma
 $\langle \exists x x'. \forall y. \exists z z'.$
 $(\neg P(y,y) \vee P(x,x) \vee \neg S(z,x)) \wedge$
 $(S(x,y) \vee \neg S(y,z) \vee Q(z',z')) \wedge$
 $(Q(x',y) \vee \neg Q(y,z') \vee S(x',x')) \rangle$
oops

11.4 Hard examples with quantifiers

18

lemma $\langle \exists y. \forall x. P(y) \longrightarrow P(x) \rangle$
by *blast*

19

lemma $\langle \exists x. \forall y z. (P(y) \longrightarrow Q(z)) \longrightarrow (P(x) \longrightarrow Q(x)) \rangle$
by *blast*

20

lemma $\langle (\forall x y. \exists z. \forall w. (P(x) \wedge Q(y) \longrightarrow R(z) \wedge S(w)))$
 $\longrightarrow (\exists x y. P(x) \wedge Q(y)) \longrightarrow (\exists z. R(z)) \rangle$

by *blast*

21

lemma $\langle (\exists x. P \longrightarrow Q(x)) \wedge (\exists x. Q(x) \longrightarrow P) \longrightarrow (\exists x. P \longleftrightarrow Q(x)) \rangle$
by *blast*

22

lemma $\langle (\forall x. P \longleftrightarrow Q(x)) \longrightarrow (P \longleftrightarrow (\forall x. Q(x))) \rangle$
by *blast*

23

lemma $\langle (\forall x. P \vee Q(x)) \longleftrightarrow (P \vee (\forall x. Q(x))) \rangle$
by *blast*

24

lemma
 $\langle \neg (\exists x. S(x) \wedge Q(x)) \wedge (\forall x. P(x) \longrightarrow Q(x) \vee R(x)) \wedge$
 $\quad (\neg (\exists x. P(x)) \longrightarrow (\exists x. Q(x))) \wedge (\forall x. Q(x) \vee R(x) \longrightarrow S(x))$
 $\quad \longrightarrow (\exists x. P(x) \wedge R(x)) \rangle$
by *blast*

25

lemma
 $\langle (\exists x. P(x)) \wedge$
 $\quad (\forall x. L(x) \longrightarrow \neg (M(x) \wedge R(x))) \wedge$
 $\quad (\forall x. P(x) \longrightarrow (M(x) \wedge L(x))) \wedge$
 $\quad ((\forall x. P(x) \longrightarrow Q(x)) \vee (\exists x. P(x) \wedge R(x)))$
 $\quad \longrightarrow (\exists x. Q(x) \wedge P(x)) \rangle$
by *blast*

26

lemma
 $\langle ((\exists x. p(x)) \longleftrightarrow (\exists x. q(x))) \wedge$
 $\quad (\forall x. \forall y. p(x) \wedge q(y) \longrightarrow (r(x) \longleftrightarrow s(y)))$
 $\quad \longrightarrow ((\forall x. p(x) \longrightarrow r(x)) \longleftrightarrow (\forall x. q(x) \longrightarrow s(x))) \rangle$
by *blast*

27

lemma
 $\langle (\exists x. P(x) \wedge \neg Q(x)) \wedge$
 $\quad (\forall x. P(x) \longrightarrow R(x)) \wedge$
 $\quad (\forall x. M(x) \wedge L(x) \longrightarrow P(x)) \wedge$
 $\quad ((\exists x. R(x) \wedge \neg Q(x)) \longrightarrow (\forall x. L(x) \longrightarrow \neg R(x)))$
 $\quad \longrightarrow (\forall x. M(x) \longrightarrow \neg L(x)) \rangle$
by *blast*

28. AMENDED

lemma

$$\langle (\forall x. P(x) \longrightarrow (\forall x. Q(x))) \wedge \\ ((\forall x. Q(x) \vee R(x)) \longrightarrow (\exists x. Q(x) \wedge S(x))) \wedge \\ ((\exists x. S(x)) \longrightarrow (\forall x. L(x) \longrightarrow M(x))) \\ \longrightarrow (\forall x. P(x) \wedge L(x) \longrightarrow M(x)) \rangle$$

by blast

29. Essentially the same as Principia Mathematica *11.71

lemma

$$\langle (\exists x. P(x)) \wedge (\exists y. Q(y)) \\ \longrightarrow ((\forall x. P(x) \longrightarrow R(x)) \wedge (\forall y. Q(y) \longrightarrow S(y)) \longleftrightarrow \\ (\forall x y. P(x) \wedge Q(y) \longrightarrow R(x) \wedge S(y))) \rangle$$

by blast

30

lemma

$$\langle (\forall x. P(x) \vee Q(x) \longrightarrow \neg R(x)) \wedge \\ (\forall x. (Q(x) \longrightarrow \neg S(x)) \longrightarrow P(x) \wedge R(x)) \\ \longrightarrow (\forall x. S(x)) \rangle$$

by blast

31

lemma

$$\langle \neg (\exists x. P(x) \wedge (Q(x) \vee R(x))) \wedge \\ (\exists x. L(x) \wedge P(x)) \wedge \\ (\forall x. \neg R(x) \longrightarrow M(x)) \\ \longrightarrow (\exists x. L(x) \wedge M(x)) \rangle$$

by blast

32

lemma

$$\langle (\forall x. P(x) \wedge (Q(x) \vee R(x)) \longrightarrow S(x)) \wedge \\ (\forall x. S(x) \wedge R(x) \longrightarrow L(x)) \wedge \\ (\forall x. M(x) \longrightarrow R(x)) \\ \longrightarrow (\forall x. P(x) \wedge M(x) \longrightarrow L(x)) \rangle$$

by blast

33

lemma

$$\langle (\forall x. P(a) \wedge (P(x) \longrightarrow P(b)) \longrightarrow P(c)) \longleftrightarrow \\ (\forall x. (\neg P(a) \vee P(x) \vee P(c)) \wedge (\neg P(a) \vee \neg P(b) \vee P(c))) \rangle$$

by blast

34. AMENDED (TWICE!!). Andrews's challenge.

lemma

$$\langle ((\exists x. \forall y. p(x) \longleftrightarrow p(y)) \longleftrightarrow ((\exists x. q(x)) \longleftrightarrow (\forall y. p(y)))) \longleftrightarrow \\ ((\exists x. \forall y. q(x) \longleftrightarrow q(y)) \longleftrightarrow ((\exists x. p(x)) \longleftrightarrow (\forall y. q(y)))) \rangle$$

by blast

35

lemma $\langle \exists x y. P(x,y) \longrightarrow (\forall u v. P(u,v)) \rangle$
by *blast*

36

lemma
 $\langle (\forall x. \exists y. J(x,y)) \wedge$
 $(\forall x. \exists y. G(x,y)) \wedge$
 $(\forall x y. J(x,y) \vee G(x,y) \longrightarrow (\forall z. J(y,z) \vee G(y,z) \longrightarrow H(x,z)))$
 $\longrightarrow (\forall x. \exists y. H(x,y)) \rangle$
by *blast*

37

lemma
 $\langle (\forall z. \exists w. \forall x. \exists y.$
 $(P(x,z) \longrightarrow P(y,w)) \wedge P(y,z) \wedge (P(y,w) \longrightarrow (\exists u. Q(u,w)))) \wedge$
 $(\forall x z. \neg P(x,z) \longrightarrow (\exists y. Q(y,z))) \wedge$
 $((\exists x y. Q(x,y)) \longrightarrow (\forall x. R(x,x)))$
 $\longrightarrow (\forall x. \exists y. R(x,y)) \rangle$
by *blast*

38

lemma
 $\langle (\forall x. p(a) \wedge (p(x) \longrightarrow (\exists y. p(y) \wedge r(x,y))) \longrightarrow$
 $(\exists z. \exists w. p(z) \wedge r(x,w) \wedge r(w,z))) \longleftrightarrow$
 $(\forall x. (\neg p(a) \vee p(x) \vee (\exists z. \exists w. p(z) \wedge r(x,w) \wedge r(w,z))) \wedge$
 $(\neg p(a) \vee \neg (\exists y. p(y) \wedge r(x,y)) \vee$
 $(\exists z. \exists w. p(z) \wedge r(x,w) \wedge r(w,z)))) \rangle$
by *blast*

39

lemma $\langle \neg (\exists x. \forall y. F(y,x) \longleftrightarrow \neg F(y,y)) \rangle$
by *blast*

40. AMENDED

lemma
 $\langle (\exists y. \forall x. F(x,y) \longleftrightarrow F(x,x)) \longrightarrow$
 $\neg (\forall x. \exists y. \forall z. F(z,y) \longleftrightarrow \neg F(z,x)) \rangle$
by *blast*

41

lemma
 $\langle (\forall z. \exists y. \forall x. f(x,y) \longleftrightarrow f(x,z) \wedge \neg f(x,x))$
 $\longrightarrow \neg (\exists z. \forall x. f(x,z)) \rangle$
by *blast*

42

lemma $\langle \neg (\exists y. \forall x. p(x,y) \longleftrightarrow \neg (\exists z. p(x,z) \wedge p(z,x))) \rangle$
by *blast*

43

lemma
 $\langle (\forall x. \forall y. q(x,y) \longleftrightarrow (\forall z. p(z,x) \longleftrightarrow p(z,y)))$
 $\longrightarrow (\forall x. \forall y. q(x,y) \longleftrightarrow q(y,x)) \rangle$
by *blast*

Other proofs: Can use *auto*, which cheats by using rewriting! *Deepen-tac* alone requires 253 secs. Or *by (mini-tac 1 THEN Deepen-tac 5 1)*.

44

lemma
 $\langle (\forall x. f(x) \longrightarrow (\exists y. g(y) \wedge h(x,y) \wedge (\exists y. g(y) \wedge \neg h(x,y)))) \wedge$
 $(\exists x. j(x) \wedge (\forall y. g(y) \longrightarrow h(x,y)))$
 $\longrightarrow (\exists x. j(x) \wedge \neg f(x)) \rangle$
by *blast*

45

lemma
 $\langle (\forall x. f(x) \wedge (\forall y. g(y) \wedge h(x,y) \longrightarrow j(x,y))$
 $\longrightarrow (\forall y. g(y) \wedge h(x,y) \longrightarrow k(y))) \wedge$
 $\neg (\exists y. l(y) \wedge k(y)) \wedge$
 $(\exists x. f(x) \wedge (\forall y. h(x,y) \longrightarrow l(y)) \wedge (\forall y. g(y) \wedge h(x,y) \longrightarrow j(x,y)))$
 $\longrightarrow (\exists x. f(x) \wedge \neg (\exists y. g(y) \wedge h(x,y))) \rangle$
by *blast*

46

lemma
 $\langle (\forall x. f(x) \wedge (\forall y. f(y) \wedge h(y,x) \longrightarrow g(y)) \longrightarrow g(x)) \wedge$
 $(\exists x. f(x) \wedge \neg g(x)) \longrightarrow$
 $(\exists x. f(x) \wedge \neg g(x) \wedge (\forall y. f(y) \wedge \neg g(y) \longrightarrow j(x,y))) \wedge$
 $(\forall x y. f(x) \wedge f(y) \wedge h(x,y) \longrightarrow \neg j(y,x))$
 $\longrightarrow (\forall x. f(x) \longrightarrow g(x)) \rangle$
by *blast*

11.5 Problems (mainly) involving equality or functions

48

lemma $\langle (a = b \vee c = d) \wedge (a = c \vee b = d) \longrightarrow a = d \vee b = c \rangle$
by *blast*

49. NOT PROVED AUTOMATICALLY. Hard because it involves substitution for Vars; the type constraint ensures that x,y,z have the same type as a,b,u.

lemma

$\langle (\exists x y :: 'a. \forall z. z = x \vee z = y) \wedge P(a) \wedge P(b) \wedge a \neq b \longrightarrow (\forall u :: 'a. P(u)) \rangle$
apply safe
apply (*rule-tac* $x = \langle a \rangle$ **in** *allE, assumption*)
apply (*rule-tac* $x = \langle b \rangle$ **in** *allE, assumption*)
apply fast — blast's treatment of equality can't do it
done

50. (What has this to do with equality?)

lemma $\langle (\forall x. P(a,x) \vee (\forall y. P(x,y))) \longrightarrow (\exists x. \forall y. P(x,y)) \rangle$
by blast

51

lemma
 $\langle (\exists z w. \forall x y. P(x,y) \longleftrightarrow (x = z \wedge y = w)) \longrightarrow$
 $\langle (\exists z. \forall x. \exists w. (\forall y. P(x,y) \longleftrightarrow y=w) \longleftrightarrow x = z) \rangle$
by blast

52

Almost the same as 51.

lemma
 $\langle (\exists z w. \forall x y. P(x,y) \longleftrightarrow (x = z \wedge y = w)) \longrightarrow$
 $\langle (\exists w. \forall y. \exists z. (\forall x. P(x,y) \longleftrightarrow x = z) \longleftrightarrow y = w) \rangle$
by blast

55

Non-equational version, from Manthey and Bry, CADE-9 (Springer, 1988).
fast DISCOVERS who killed Agatha.

schematic-goal

$\langle \text{lives}(\text{agatha}) \wedge \text{lives}(\text{butler}) \wedge \text{lives}(\text{charles}) \wedge$
 $(\text{killed}(\text{agatha}, \text{agatha}) \vee \text{killed}(\text{butler}, \text{agatha}) \vee \text{killed}(\text{charles}, \text{agatha})) \wedge$
 $(\forall x y. \text{killed}(x,y) \longrightarrow \text{hates}(x,y) \wedge \neg \text{richer}(x,y)) \wedge$
 $(\forall x. \text{hates}(\text{agatha}, x) \longrightarrow \neg \text{hates}(\text{charles}, x)) \wedge$
 $(\text{hates}(\text{agatha}, \text{agatha}) \wedge \text{hates}(\text{agatha}, \text{charles})) \wedge$
 $(\forall x. \text{lives}(x) \wedge \neg \text{richer}(x, \text{agatha}) \longrightarrow \text{hates}(\text{butler}, x)) \wedge$
 $(\forall x. \text{hates}(\text{agatha}, x) \longrightarrow \text{hates}(\text{butler}, x)) \wedge$
 $(\forall x. \neg \text{hates}(x, \text{agatha}) \vee \neg \text{hates}(x, \text{butler}) \vee \neg \text{hates}(x, \text{charles})) \longrightarrow$
 $\text{killed}(\text{?who}, \text{agatha}) \rangle$
by fast — MUCH faster than blast

56

lemma $\langle (\forall x. (\exists y. P(y) \wedge x = f(y)) \longrightarrow P(x)) \longleftrightarrow (\forall x. P(x) \longrightarrow P(f(x))) \rangle$
by blast

57

lemma
 $\langle P(f(a,b), f(b,c)) \wedge P(f(b,c), f(a,c)) \wedge$

$\langle (\forall x y z. P(x,y) \wedge P(y,z) \longrightarrow P(x,z)) \longrightarrow P(f(a,b), f(a,c)) \rangle$
by blast

58 NOT PROVED AUTOMATICALLY

lemma $\langle (\forall x y. f(x) = g(y)) \longrightarrow (\forall x y. f(f(x)) = f(g(y))) \rangle$
by (*slow elim: subst-context*)

59

lemma $\langle (\forall x. P(x) \longleftrightarrow \neg P(f(x))) \longrightarrow (\exists x. P(x) \wedge \neg P(f(x))) \rangle$
by blast

60

lemma $\langle \forall x. P(x, f(x)) \longleftrightarrow (\exists y. (\forall z. P(z, y) \longrightarrow P(z, f(x))) \wedge P(x, y)) \rangle$
by blast

62 as corrected in JAR 18 (1997), page 135

lemma
 $\langle (\forall x. p(a) \wedge (p(x) \longrightarrow p(f(x))) \longrightarrow p(f(f(x)))) \longleftrightarrow$
 $(\forall x. (\neg p(a) \vee p(x) \vee p(f(f(x)))) \wedge$
 $(\neg p(a) \vee \neg p(f(x)) \vee p(f(f(x)))) \rangle$
by blast

From Davis, Obvious Logical Inferences, IJCAI-81, 530-531 fast indeed copes!

lemma
 $\langle (\forall x. F(x) \wedge \neg G(x) \longrightarrow (\exists y. H(x,y) \wedge J(y))) \wedge$
 $(\exists x. K(x) \wedge F(x) \wedge (\forall y. H(x,y) \longrightarrow K(y))) \wedge$
 $(\forall x. K(x) \longrightarrow \neg G(x)) \longrightarrow (\exists x. K(x) \wedge J(x)) \rangle$
by fast

From Rudnicki, Obvious Inferences, JAR 3 (1987), 383-393. It does seem obvious!

lemma
 $\langle (\forall x. F(x) \wedge \neg G(x) \longrightarrow (\exists y. H(x,y) \wedge J(y))) \wedge$
 $(\exists x. K(x) \wedge F(x) \wedge (\forall y. H(x,y) \longrightarrow K(y))) \wedge$
 $(\forall x. K(x) \longrightarrow \neg G(x)) \longrightarrow (\exists x. K(x) \longrightarrow \neg G(x)) \rangle$
by fast

Halting problem: Formulation of Li Dafa (AAR Newsletter 27, Oct 1994.)
author U. Egly.

lemma
 $\langle ((\exists x. A(x) \wedge (\forall y. C(y) \longrightarrow (\forall z. D(x,y,z)))) \longrightarrow$
 $(\exists w. C(w) \wedge (\forall y. C(y) \longrightarrow (\forall z. D(w,y,z)))))) \wedge$
 $(\forall w. C(w) \wedge (\forall u. C(u) \longrightarrow (\forall v. D(w,u,v))) \longrightarrow$
 $(\forall y z.$
 $(C(y) \wedge P(y,z) \longrightarrow Q(w,y,z) \wedge OO(w,g)) \wedge$
 $(C(y) \wedge \neg P(y,z) \longrightarrow Q(w,y,z) \wedge OO(w,b)))) \rangle$

$$\begin{aligned}
& \wedge \\
& (\forall w. C(w) \wedge \\
& \quad (\forall y z. \\
& \quad \quad (C(y) \wedge P(y,z) \longrightarrow Q(w,y,z) \wedge OO(w,g)) \wedge \\
& \quad \quad (C(y) \wedge \neg P(y,z) \longrightarrow Q(w,y,z) \wedge OO(w,b))) \longrightarrow \\
& \quad (\exists v. C(v) \wedge \\
& \quad \quad (\forall y. ((C(y) \wedge Q(w,y,y)) \wedge OO(w,g) \longrightarrow \neg P(v,y)) \wedge \\
& \quad \quad \quad ((C(y) \wedge Q(w,y,y)) \wedge OO(w,b) \longrightarrow P(v,y) \wedge OO(v,b)))))) \\
& \longrightarrow \neg (\exists x. A(x) \wedge (\forall y. C(y) \longrightarrow (\forall z. D(x,y,z)))) \rangle
\end{aligned}$$

by (*blast 12*)

— Needed because the search for depths below 12 is very slow.

Halting problem II: credited to M. Bruschi by Li Dafa in JAR 18(1), p. 105.

lemma

$$\begin{aligned}
& \langle ((\exists x. A(x) \wedge (\forall y. C(y) \longrightarrow (\forall z. D(x,y,z)))) \longrightarrow \\
& \quad (\exists w. C(w) \wedge (\forall y. C(y) \longrightarrow (\forall z. D(w,y,z)))))) \\
& \wedge \\
& (\forall w. C(w) \wedge (\forall u. C(u) \longrightarrow (\forall v. D(w,u,v))) \longrightarrow \\
& \quad (\forall y z. \\
& \quad \quad (C(y) \wedge P(y,z) \longrightarrow Q(w,y,z) \wedge OO(w,g)) \wedge \\
& \quad \quad (C(y) \wedge \neg P(y,z) \longrightarrow Q(w,y,z) \wedge OO(w,b)))))) \\
& \wedge \\
& ((\exists w. C(w) \wedge (\forall y. (C(y) \wedge P(y,y) \longrightarrow Q(w,y,y) \wedge OO(w,g)) \wedge \\
& \quad (C(y) \wedge \neg P(y,y) \longrightarrow Q(w,y,y) \wedge OO(w,b)))))) \\
& \longrightarrow \\
& (\exists v. C(v) \wedge (\forall y. (C(y) \wedge P(y,y) \longrightarrow P(v,y) \wedge OO(v,g)) \wedge \\
& \quad (C(y) \wedge \neg P(y,y) \longrightarrow P(v,y) \wedge OO(v,b)))))) \\
& \longrightarrow \\
& ((\exists v. C(v) \wedge (\forall y. (C(y) \wedge P(y,y) \longrightarrow P(v,y) \wedge OO(v,g)) \wedge \\
& \quad (C(y) \wedge \neg P(y,y) \longrightarrow P(v,y) \wedge OO(v,b)))))) \\
& \longrightarrow \\
& (\exists u. C(u) \wedge (\forall y. (C(y) \wedge P(y,y) \longrightarrow \neg P(u,y)) \wedge \\
& \quad (C(y) \wedge \neg P(y,y) \longrightarrow P(u,y) \wedge OO(u,b)))))) \\
& \longrightarrow \neg (\exists x. A(x) \wedge (\forall y. C(y) \longrightarrow (\forall z. D(x,y,z)))) \rangle
\end{aligned}$$

by *blast*

Challenge found on info-hol.

lemma $\langle \forall x. \exists v w. \forall y z. P(x) \wedge Q(y) \longrightarrow (P(v) \vee R(w)) \wedge (R(z) \longrightarrow Q(v)) \rangle$

by *blast*

Attributed to Lewis Carroll by S. G. Pulman. The first or last assumption can be deleted.

lemma

$$\begin{aligned}
& \langle (\forall x. \text{honest}(x) \wedge \text{industrious}(x) \longrightarrow \text{healthy}(x)) \wedge \\
& \quad \neg (\exists x. \text{grocer}(x) \wedge \text{healthy}(x)) \wedge \\
& \quad (\forall x. \text{industrious}(x) \wedge \text{grocer}(x) \longrightarrow \text{honest}(x)) \wedge \\
& \quad (\forall x. \text{cyclist}(x) \longrightarrow \text{industrious}(x)) \wedge \\
& \quad (\forall x. \neg \text{healthy}(x) \wedge \text{cyclist}(x) \longrightarrow \neg \text{honest}(x)) \\
& \quad \longrightarrow (\forall x. \text{grocer}(x) \longrightarrow \neg \text{cyclist}(x)) \rangle
\end{aligned}$$

by *blast*

end

12 First-Order Logic: propositional examples (classical version)

theory *Propositional-Cla*
imports *FOL*
begin

commutative laws of \wedge and \vee

lemma $\langle P \wedge Q \longrightarrow Q \wedge P \rangle$
by (*tactic IntPr.fast-tac context 1*)

lemma $\langle P \vee Q \longrightarrow Q \vee P \rangle$
by *fast*

associative laws of \wedge and \vee

lemma $\langle (P \wedge Q) \wedge R \longrightarrow P \wedge (Q \wedge R) \rangle$
by *fast*

lemma $\langle (P \vee Q) \vee R \longrightarrow P \vee (Q \vee R) \rangle$
by *fast*

distributive laws of \wedge and \vee

lemma $\langle (P \wedge Q) \vee R \longrightarrow (P \vee R) \wedge (Q \vee R) \rangle$
by *fast*

lemma $\langle (P \vee R) \wedge (Q \vee R) \longrightarrow (P \wedge Q) \vee R \rangle$
by *fast*

lemma $\langle (P \vee Q) \wedge R \longrightarrow (P \wedge R) \vee (Q \wedge R) \rangle$
by *fast*

lemma $\langle (P \wedge R) \vee (Q \wedge R) \longrightarrow (P \vee Q) \wedge R \rangle$
by *fast*

Laws involving implication

lemma $\langle (P \longrightarrow R) \wedge (Q \longrightarrow R) \longleftrightarrow (P \vee Q \longrightarrow R) \rangle$
by *fast*

lemma $\langle (P \wedge Q \longrightarrow R) \longleftrightarrow (P \longrightarrow (Q \longrightarrow R)) \rangle$
by *fast*

lemma $\langle ((P \rightarrow R) \rightarrow R) \rightarrow ((Q \rightarrow R) \rightarrow R) \rightarrow (P \wedge Q \rightarrow R) \rightarrow R \rangle$
by fast

lemma $\langle \neg (P \rightarrow R) \rightarrow \neg (Q \rightarrow R) \rightarrow \neg (P \wedge Q \rightarrow R) \rangle$
by fast

lemma $\langle (P \rightarrow Q \wedge R) \leftrightarrow (P \rightarrow Q) \wedge (P \rightarrow R) \rangle$
by fast

Propositions-as-types

lemma $\langle P \rightarrow (Q \rightarrow P) \rangle$
by fast

— The combinator S

lemma $\langle (P \rightarrow Q \rightarrow R) \rightarrow (P \rightarrow Q) \rightarrow (P \rightarrow R) \rangle$
by fast

— Converse is classical

lemma $\langle (P \rightarrow Q) \vee (P \rightarrow R) \rightarrow (P \rightarrow Q \vee R) \rangle$
by fast

lemma $\langle (P \rightarrow Q) \rightarrow (\neg Q \rightarrow \neg P) \rangle$
by fast

Schwichtenberg's examples (via T. Nipkow)

lemma *stab-imp*: $\langle (((Q \rightarrow R) \rightarrow R) \rightarrow Q) \rightarrow (((P \rightarrow Q) \rightarrow R) \rightarrow R) \rightarrow P \rightarrow Q \rangle$
by fast

lemma *stab-to-peirce*:

$\langle (((P \rightarrow R) \rightarrow R) \rightarrow P) \rightarrow (((Q \rightarrow R) \rightarrow R) \rightarrow Q) \rightarrow ((P \rightarrow Q) \rightarrow P) \rightarrow P \rangle$
by fast

lemma *peirce-imp1*:

$\langle (((Q \rightarrow R) \rightarrow Q) \rightarrow Q) \rightarrow (((P \rightarrow Q) \rightarrow R) \rightarrow P \rightarrow Q) \rightarrow P \rightarrow Q \rangle$
by fast

lemma *peirce-imp2*: $\langle (((P \rightarrow R) \rightarrow P) \rightarrow P) \rightarrow ((P \rightarrow Q \rightarrow R) \rightarrow P) \rightarrow P \rangle$
by fast

lemma *mits*: $\langle (((P \rightarrow Q) \rightarrow P) \rightarrow P) \rightarrow Q \rightarrow Q \rangle$
by fast

lemma *mits-solovev*: $\langle (P \rightarrow (Q \rightarrow R) \rightarrow Q) \rightarrow ((P \rightarrow Q) \rightarrow R) \rightarrow R \rangle$

by fast

lemma tatsuta:

$\langle (((P7 \rightarrow P1) \rightarrow P10) \rightarrow P4 \rightarrow P5)$
 $\rightarrow (((P8 \rightarrow P2) \rightarrow P9) \rightarrow P3 \rightarrow P10)$
 $\rightarrow (P1 \rightarrow P8) \rightarrow P6 \rightarrow P7$
 $\rightarrow (((P3 \rightarrow P2) \rightarrow P9) \rightarrow P4)$
 $\rightarrow (P1 \rightarrow P3) \rightarrow (((P6 \rightarrow P1) \rightarrow P2) \rightarrow P9) \rightarrow P5 \rangle$
by fast

lemma tatsuta1:

$\langle (((P8 \rightarrow P2) \rightarrow P9) \rightarrow P3 \rightarrow P10)$
 $\rightarrow (((P3 \rightarrow P2) \rightarrow P9) \rightarrow P4)$
 $\rightarrow (((P6 \rightarrow P1) \rightarrow P2) \rightarrow P9)$
 $\rightarrow (((P7 \rightarrow P1) \rightarrow P10) \rightarrow P4 \rightarrow P5)$
 $\rightarrow (P1 \rightarrow P3) \rightarrow (P1 \rightarrow P8) \rightarrow P6 \rightarrow P7 \rightarrow P5 \rangle$
by fast

end

13 First-Order Logic: quantifier examples (classical version)

theory *Quantifiers-Cla*

imports *FOL*

begin

lemma $\langle (\forall x y. P(x,y)) \rightarrow (\forall y x. P(x,y)) \rangle$
by fast

lemma $\langle (\exists x y. P(x,y)) \rightarrow (\exists y x. P(x,y)) \rangle$
by fast

Converse is false.

lemma $\langle (\forall x. P(x)) \vee (\forall x. Q(x)) \rightarrow (\forall x. P(x) \vee Q(x)) \rangle$
by fast

lemma $\langle (\forall x. P \rightarrow Q(x)) \leftrightarrow (P \rightarrow (\forall x. Q(x))) \rangle$
by fast

lemma $\langle (\forall x. P(x) \rightarrow Q) \leftrightarrow ((\exists x. P(x)) \rightarrow Q) \rangle$
by fast

Some harder ones.

lemma $\langle (\exists x. P(x) \vee Q(x)) \leftrightarrow (\exists x. P(x)) \vee (\exists x. Q(x)) \rangle$
by fast

— Converse is false.

lemma $\langle (\exists x. P(x) \wedge Q(x)) \longrightarrow (\exists x. P(x)) \wedge (\exists x. Q(x)) \rangle$
by *fast*

Basic test of quantifier reasoning.

lemma $\langle (\exists y. \forall x. Q(x,y)) \longrightarrow (\forall x. \exists y. Q(x,y)) \rangle$
by *fast*

lemma $\langle (\forall x. Q(x)) \longrightarrow (\exists x. Q(x)) \rangle$
by *fast*

The following should fail, as they are false!

lemma $\langle (\forall x. \exists y. Q(x,y)) \longrightarrow (\exists y. \forall x. Q(x,y)) \rangle$
apply *fast?*
oops

lemma $\langle (\exists x. Q(x)) \longrightarrow (\forall x. Q(x)) \rangle$
apply *fast?*
oops

schematic-goal $\langle P(?a) \longrightarrow (\forall x. P(x)) \rangle$
apply *fast?*
oops

schematic-goal $\langle (P(?a) \longrightarrow (\forall x. Q(x))) \longrightarrow (\forall x. P(x) \longrightarrow Q(x)) \rangle$
apply *fast?*
oops

Back to things that are provable ...

lemma $\langle (\forall x. P(x) \longrightarrow Q(x)) \wedge (\exists x. P(x)) \longrightarrow (\exists x. Q(x)) \rangle$
by *fast*

An example of why *exI* should be delayed as long as possible.

lemma $\langle (P \longrightarrow (\exists x. Q(x))) \wedge P \longrightarrow (\exists x. Q(x)) \rangle$
by *fast*

schematic-goal $\langle (\forall x. P(x) \longrightarrow Q(f(x))) \wedge (\forall x. Q(x) \longrightarrow R(g(x))) \wedge P(d) \longrightarrow R(?a) \rangle$
by *fast*

lemma $\langle (\forall x. Q(x)) \longrightarrow (\exists x. Q(x)) \rangle$
by *fast*

Some slow ones

Principia Mathematica *11.53

lemma $\langle (\forall x y. P(x) \longrightarrow Q(y)) \longleftrightarrow ((\exists x. P(x)) \longrightarrow (\forall y. Q(y))) \rangle$
by *fast*

lemma $\langle (\exists x y. P(x) \wedge Q(x,y)) \longleftrightarrow (\exists x. P(x) \wedge (\exists y. Q(x,y))) \rangle$
by *fast*

lemma $\langle (\exists y. \forall x. P(x) \longrightarrow Q(x,y)) \longrightarrow (\forall x. P(x) \longrightarrow (\exists y. Q(x,y))) \rangle$
by *fast*

end

theory *Miniscope*
imports *FOL*
begin

lemmas *ccontr* = *FalseE* [*THEN classical*]

13.1 Negation Normal Form

13.1.1 de Morgan laws

lemma *demorgans1*:
 $\langle \neg (P \wedge Q) \longleftrightarrow \neg P \vee \neg Q \rangle$
 $\langle \neg (P \vee Q) \longleftrightarrow \neg P \wedge \neg Q \rangle$
 $\langle \neg \neg P \longleftrightarrow P \rangle$
by *blast+*

lemma *demorgans2*:
 $\langle \bigwedge P. \neg (\forall x. P(x)) \longleftrightarrow (\exists x. \neg P(x)) \rangle$
 $\langle \bigwedge P. \neg (\exists x. P(x)) \longleftrightarrow (\forall x. \neg P(x)) \rangle$
by *blast+*

lemmas *demorgans* = *demorgans1 demorgans2*

lemma *nnf-simps*:
 $\langle (P \longrightarrow Q) \longleftrightarrow (\neg P \vee Q) \rangle$
 $\langle \neg (P \longrightarrow Q) \longleftrightarrow (P \wedge \neg Q) \rangle$
 $\langle (P \longleftrightarrow Q) \longleftrightarrow (\neg P \vee Q) \wedge (\neg Q \vee P) \rangle$
 $\langle \neg (P \longleftrightarrow Q) \longleftrightarrow (P \vee Q) \wedge (\neg P \vee \neg Q) \rangle$
by *blast+*

13.1.2 Pushing in the existential quantifiers

lemma *ex-simps*:
 $\langle (\exists x. P) \longleftrightarrow P \rangle$
 $\langle \bigwedge P Q. (\exists x. P(x) \wedge Q) \longleftrightarrow (\exists x. P(x)) \wedge Q \rangle$
 $\langle \bigwedge P Q. (\exists x. P \wedge Q(x)) \longleftrightarrow P \wedge (\exists x. Q(x)) \rangle$

```

⟨ $\bigwedge P Q. (\exists x. P(x) \vee Q(x)) \longleftrightarrow (\exists x. P(x)) \vee (\exists x. Q(x))$ ⟩
⟨ $\bigwedge P Q. (\exists x. P(x) \vee Q) \longleftrightarrow (\exists x. P(x)) \vee Q$ ⟩
⟨ $\bigwedge P Q. (\exists x. P \vee Q(x)) \longleftrightarrow P \vee (\exists x. Q(x))$ ⟩
by blast+

```

13.1.3 Pushing in the universal quantifiers

lemma *all-simps*:

```

⟨ $(\forall x. P) \longleftrightarrow P$ ⟩
⟨ $\bigwedge P Q. (\forall x. P(x) \wedge Q(x)) \longleftrightarrow (\forall x. P(x)) \wedge (\forall x. Q(x))$ ⟩
⟨ $\bigwedge P Q. (\forall x. P(x) \wedge Q) \longleftrightarrow (\forall x. P(x)) \wedge Q$ ⟩
⟨ $\bigwedge P Q. (\forall x. P \wedge Q(x)) \longleftrightarrow P \wedge (\forall x. Q(x))$ ⟩
⟨ $\bigwedge P Q. (\forall x. P(x) \vee Q) \longleftrightarrow (\forall x. P(x)) \vee Q$ ⟩
⟨ $\bigwedge P Q. (\forall x. P \vee Q(x)) \longleftrightarrow P \vee (\forall x. Q(x))$ ⟩
by blast+

```

lemmas *mini-simps = demorgans nnf-simps ex-simps all-simps*

ML <

```

val mini-ss = simpset-of (context addsimps @ {thms mini-simps});
fun mini-tac ctxt =
  resolve-tac ctxt @ {thms ccontr} THEN' asm-full-simp-tac (put-simpset mini-ss
  ctxt);
>

```

end

14 First-Order Logic: the 'if' example

theory *If*

imports *FOL*

begin

definition *if* :: $\langle [o, o, o] \Rightarrow o \rangle$
where $\langle if(P, Q, R) \equiv P \wedge Q \vee \neg P \wedge R \rangle$

lemma *ifI*: $\langle \llbracket P \Longrightarrow Q; \neg P \Longrightarrow R \rrbracket \Longrightarrow if(P, Q, R) \rangle$
unfolding *if-def* by blast

lemma *ifE*: $\langle \llbracket if(P, Q, R); \llbracket P; Q \rrbracket \Longrightarrow S; \llbracket \neg P; R \rrbracket \Longrightarrow S \rrbracket \Longrightarrow S \rangle$
unfolding *if-def* by blast

lemma *if-commute*: $\langle if(P, if(Q, A, B), if(Q, C, D)) \longleftrightarrow if(Q, if(P, A, C), if(P, B, D)) \rangle$
apply (*rule iffI*)
apply (*erule ifE*)
apply (*erule ifE*)
apply (*rule ifI*)
apply (*rule ifI*)
oops

Trying again from the beginning in order to use *blast*

declare *ifI* [*intro!*]

declare *ifE* [*elim!*]

lemma *if-commute*: $\langle \text{if}(P, \text{if}(Q, A, B), \text{if}(Q, C, D)) \longleftrightarrow \text{if}(Q, \text{if}(P, A, C), \text{if}(P, B, D)) \rangle$
by *blast*

lemma $\langle \text{if}(\text{if}(P, Q, R), A, B) \longleftrightarrow \text{if}(P, \text{if}(Q, A, B), \text{if}(R, A, B)) \rangle$
by *blast*

Trying again from the beginning in order to prove from the definitions

lemma $\langle \text{if}(\text{if}(P, Q, R), A, B) \longleftrightarrow \text{if}(P, \text{if}(Q, A, B), \text{if}(R, A, B)) \rangle$
unfolding *if-def* **by** *blast*

An invalid formula. High-level rules permit a simpler diagnosis.

lemma $\langle \text{if}(\text{if}(P, Q, R), A, B) \longleftrightarrow \text{if}(P, \text{if}(Q, A, B), \text{if}(R, B, A)) \rangle$
apply *auto*
— The next step will fail unless subgoals remain
apply (*tactic all-tac*)
oops

Trying again from the beginning in order to prove from the definitions.

lemma $\langle \text{if}(\text{if}(P, Q, R), A, B) \longleftrightarrow \text{if}(P, \text{if}(Q, A, B), \text{if}(R, B, A)) \rangle$
unfolding *if-def*
apply *auto*
— The next step will fail unless subgoals remain
apply (*tactic all-tac*)
oops

end