

Complex Analysis

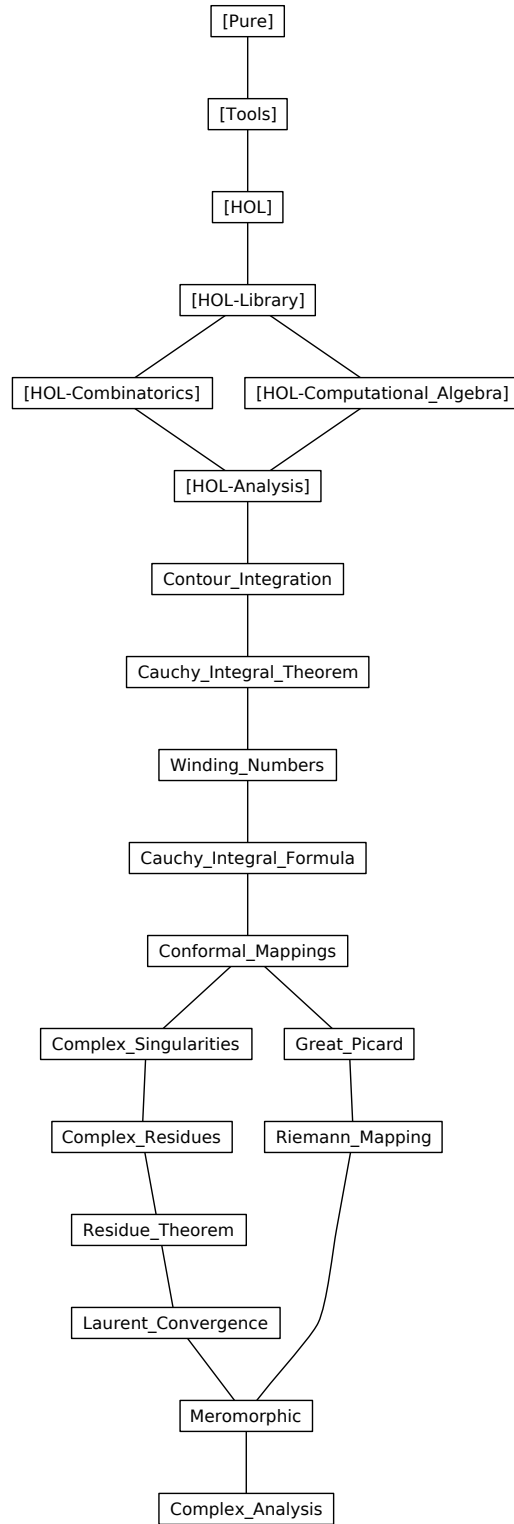
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1 Contour integration

theory *Contour_Integration*

imports *HOL-Analysis.Analysis*

begin

lemma *lhopital_complex_simple*:

assumes (*f has_field_derivative f'*) (*at z*)

assumes (*g has_field_derivative g'*) (*at z*)

assumes $f z = 0$ $g z = 0$ $g' \neq 0$ $f' / g' = c$

shows $((\lambda w. f w / g w) \longrightarrow c)$ (*at z*)

proof –

have *eventually* $(\lambda w. w \neq z)$ (*at z*)

by (*auto simp: eventually_at_filter*)

hence *eventually* $(\lambda w. ((f w - f z) / (w - z)) / ((g w - g z) / (w - z)) = f w / g w)$ (*at z*)

by *eventually_elim* (*simp add: assms field_split_simps*)

moreover **have** $((\lambda w. ((f w - f z) / (w - z)) / ((g w - g z) / (w - z))) \longrightarrow f' / g')$ (*at z*)

by (*intro tendsto_divide has_field_derivativeD assms*)

ultimately **have** $((\lambda w. f w / g w) \longrightarrow f' / g')$ (*at z*)

by (*blast intro: Lim_transform_eventually*)

with *assms* **show** *?thesis* **by** *simp*

qed

1.1 Definition

This definition is for complex numbers only, and does not generalise to line integrals in a vector field

definition *has_contour_integral* :: $(\text{complex} \Rightarrow \text{complex}) \Rightarrow \text{complex} \Rightarrow (\text{real} \Rightarrow \text{complex}) \Rightarrow \text{bool}$

(**infix** *has'_contour'_integral* 50)

where (*f has_contour_integral i*) $g \equiv$

$((\lambda x. f(g x) * \text{vector_derivative } g \text{ (at } x \text{ within } \{0..1\}))$

has_integral i) $\{0..1\}$

definition *contour_integrable_on*

(**infix** *contour'_integrable'_on* 50)

where *f contour_integrable_on g* $\equiv \exists i. (f \text{ has_contour_integral } i) g$

definition *contour_integral*

where *contour_integral g f* $\equiv \text{SOME } i. (f \text{ has_contour_integral } i) g \vee \neg f \text{ contour_integrable_on } g \wedge i=0$

lemma *not_integrable_contour_integral*: $\neg f \text{ contour_integrable_on } g \implies \text{contour_integral } g f = 0$

unfolding *contour_integrable_on_def* *contour_integral_def* **by** *blast*

lemma *contour_integral_unique*: $(f \text{ has_contour_integral } i) g \implies \text{contour_integral } g f$

$g f = i$

apply (*simp add: contour_integral_def has_contour_integral_def contour_integrable_on_def*)
using *has_integral_unique* **by** *blast*

lemma *has_contour_integral_eqpath:*

$\llbracket (f \text{ has_contour_integral } y) p; f \text{ contour_integrable_on } \gamma;$
 $\text{contour_integral } p f = \text{contour_integral } \gamma f \rrbracket$
 $\implies (f \text{ has_contour_integral } y) \gamma$

using *contour_integrable_on_def contour_integral_unique* **by** *auto*

lemma *has_contour_integral_integral:*

$f \text{ contour_integrable_on } i \implies (f \text{ has_contour_integral } (\text{contour_integral } i f))$
 i

by (*metis contour_integral_unique contour_integrable_on_def*)

lemma *has_contour_integral_unique:*

$(f \text{ has_contour_integral } i) g \implies (f \text{ has_contour_integral } j) g \implies i = j$

using *has_integral_unique*

by (*auto simp: has_contour_integral_def*)

lemma *has_contour_integral_integrable:* $(f \text{ has_contour_integral } i) g \implies f \text{ con-}$
 $\text{tour_integrable_on } g$

using *contour_integrable_on_def* **by** *blast*

Show that we can forget about the localized derivative.

lemma *has_integral_localized_vector_derivative:*

$((\lambda x. f (g x) * \text{vector_derivative } p (\text{at } x \text{ within } \{a..b\})) \text{ has_integral } i) \{a..b\}$
 \longleftrightarrow

$((\lambda x. f (g x) * \text{vector_derivative } p (\text{at } x)) \text{ has_integral } i) \{a..b\}$

proof –

have $*$: $\{a..b\} - \{a,b\} = \text{interior } \{a..b\}$

by (*simp add: atLeastAtMost_diff_ends*)

show *?thesis*

by (*rule has_integral_spike_eq [of \{a,b\}] (auto simp: at_within_interior [of*
 $_ \{a..b\}]$)

qed

lemma *integrable_on_localized_vector_derivative:*

$(\lambda x. f (g x) * \text{vector_derivative } p (\text{at } x \text{ within } \{a..b\})) \text{ integrable_on } \{a..b\}$
 \longleftrightarrow

$(\lambda x. f (g x) * \text{vector_derivative } p (\text{at } x)) \text{ integrable_on } \{a..b\}$

by (*simp add: integrable_on_def has_integral_localized_vector_derivative*)

lemma *has_contour_integral:*

$(f \text{ has_contour_integral } i) g \longleftrightarrow$

$((\lambda x. f (g x) * \text{vector_derivative } g (\text{at } x)) \text{ has_integral } i) \{0..1\}$

by (*simp add: has_integral_localized_vector_derivative has_contour_integral_def*)

lemma *contour_integrable_on:*

f *contour_integrable_on* $g \iff$
 $(\lambda t. f(g\ t) * \text{vector_derivative } g \text{ (at } t)) \text{ integrable_on } \{0..1\}$
by (*simp add: has_contour_integral integrable_on_def contour_integrable_on_def*)

lemma *has_contour_integral_mirror_iff*:

assumes *valid_path* g
shows $(f \text{ has_contour_integral } I) (-g) \iff ((\lambda x. -f(-x)) \text{ has_contour_integral } I) g$

proof -

from *assms* **have** g *piecewise_differentiable_on* $\{0..1\}$
by (*auto simp: valid_path_def piecewise_C1_imp_differentiable*)
then obtain S **where** *finite* S **and** $S: \bigwedge x. x \in \{0..1\} - S \implies g$ *differentiable*
at x *within* $\{0..1\}$

unfolding *piecewise_differentiable_on_def* **by** *blast*

have S' : g *differentiable* *at* x **if** $x \in \{0..1\} - (\{0, 1\} \cup S)$ **for** x

proof -

from *that* **have** $x \in \text{interior } \{0..1\}$ **by** *auto*

with $S[\text{of } x]$ **that show** *?thesis* **by** (*auto simp: at_within_interior[of _ \{0..1\}]*)

qed

have $(f \text{ has_contour_integral } I) (-g) \iff$
 $((\lambda x. f(-g\ x) * \text{vector_derivative } (-g) \text{ (at } x)) \text{ has_integral } I) \{0..1\}$
by (*simp add: has_contour_integral*)
also have $\dots \iff ((\lambda x. -f(-g\ x) * \text{vector_derivative } g \text{ (at } x)) \text{ has_integral } I) \{0..1\}$

by (*intro has_integral_spike_finite_eq[of S \cup \{0, 1\}]*)

(*insert \langle finite S \rangle S'*, *auto simp: o_def fun_Comp_def*)

also have $\dots \iff ((\lambda x. -f(-x)) \text{ has_contour_integral } I) g$

by (*simp add: has_contour_integral*)

finally show *?thesis* .

qed

lemma *contour_integral_on_mirror_iff*:

assumes *valid_path* g
shows f *contour_integrable_on* $(-g) \iff (\lambda x. -f(-x))$ *contour_integrable_on* g
by (*auto simp: contour_integrable_on_def has_contour_integral_mirror_iff assms*)

lemma *contour_integral_mirror*:

assumes *valid_path* g
shows *contour_integral* $(-g) f = \text{contour_integral } g (\lambda x. -f(-x))$
proof (*cases f contour_integrable_on (-g)*)
case *True* **with** *contour_integral_unique* *assms* **show** *?thesis*
by (*auto simp: contour_integrable_on_def has_contour_integral_mirror_iff*)
next
case *False* **then show** *?thesis*
by (*simp add: assms contour_integral_on_mirror_iff not_integrable_contour_integral*)
qed

1.2 Reversing a path

lemma *has_contour_integral_reversepath*:

assumes *valid_path* *g* **and** *f*: (*f* *has_contour_integral* *i*) *g*
shows (*f* *has_contour_integral* $(-i)$) (*reversepath* *g*)

proof –

{ **fix** *S* *x*

assume *xs*: *g* *C1_differentiable_on* ($\{0..1\} - S$) $x \notin (-) 1$ ‘ *S* $0 \leq x \leq 1$

have *vector_derivative* $(\lambda x. g (1 - x))$ (*at* *x* *within* $\{0..1\}$) =
 $-$ *vector_derivative* *g* (*at* $(1 - x)$ *within* $\{0..1\}$)

proof –

obtain *f'* **where** *f'*: (*g* *has_vector_derivative* *f'*) (*at* $(1 - x)$)

using *xs*

by (*force simp: has_vector_derivative_def C1_differentiable_on_def*)

have $(g \circ (\lambda x. 1 - x))$ *has_vector_derivative* $-1 *_{\mathbb{R}}$ *f'*) (*at* *x*)

by (*intro vector_diff_chain_within has_vector_derivative_at_within* [*OF* *f'*] *derivative_eq_intros* | *simp*)+

then have *mf'*: $((\lambda x. g (1 - x))$ *has_vector_derivative* $-f'$) (*at* *x*)

by (*simp add: o_def*)

show *?thesis*

using *xs*

by (*auto simp: vector_derivative_at_within_ivl* [*OF* *mf'*] *vector_derivative_at_within_ivl* [*OF* *f'*])

qed

} **note** $*$ = *this*

obtain *S* **where** *S*: *continuous_on* $\{0..1\}$ *g* *finite* *S* *g* *C1_differentiable_on* $\{0..1\} - S$

using *assms* **by** (*auto simp: valid_path_def piecewise_C1_differentiable_on_def*)

have $((\lambda x. - (f (g (1 - x)) * \text{vector_derivative } g \text{ (at } (1 - x) \text{ within } \{0..1\})))$
has_integral $-i$)

$\{0..1\}$

using *has_integral_affinity01* [**where** $m = -1$ **and** $c = 1$, *OF* *f* [*unfolded* *has_contour_integral_def*]]

by (*simp add: has_integral_neg*)

then show *?thesis*

using *S*

unfolding *reversepath_def* *has_contour_integral_def*

by (*rule_tac* *S* = $(\lambda x. 1 - x)$ ‘ *S* **in** *has_integral_spike_finite*) (*auto simp: **)

qed

lemma *contour_integrable_reversepath*:

valid_path *g* \implies *f* *contour_integrable_on* *g* \implies *f* *contour_integrable_on* (*reversepath* *g*)

using *has_contour_integral_reversepath* *contour_integrable_on_def* **by** *blast*

lemma *contour_integrable_reversepath_eq*:

valid_path *g* \implies (*f* *contour_integrable_on* (*reversepath* *g*) \iff *f* *contour_integrable_on* *g*)

using *contour_integrable_reversepath* *valid_path_reversepath* **by** *fastforce*


```

lemma contour_integral_reversepath:
  assumes valid_path g
  shows contour_integral (reversepath g) f = - (contour_integral g f)
proof (cases f contour_integrable_on g)
  case True then show ?thesis
    by (simp add: assms contour_integral_unique has_contour_integral_integral
has_contour_integral_reversepath)
  next
  case False then have  $\neg f$  contour_integrable_on (reversepath g)
    by (simp add: assms contour_integrable_reversepath_eq)
  with False show ?thesis by (simp add: not_integrable_contour_integral)
qed

```

1.3 Joining two paths together

```

lemma has_contour_integral_join:
  assumes (f has_contour_integral i1) g1 (f has_contour_integral i2) g2
    valid_path g1 valid_path g2
  shows (f has_contour_integral (i1 + i2)) (g1 +++ g2)
proof -
  obtain s1 s2
    where s1: finite s1  $\forall x \in \{0..1\} - s1. g1$  differentiable at  $x$ 
      and s2: finite s2  $\forall x \in \{0..1\} - s2. g2$  differentiable at  $x$ 
    using assms
  by (auto simp: valid_path_def piecewise_C1_differentiable_on_def C1_differentiable_on_eq)
  have 1: (( $\lambda x. f (g1 x) * vector\_derivative g1 (at x)$ ) has_integral i1) {0..1}
    and 2: (( $\lambda x. f (g2 x) * vector\_derivative g2 (at x)$ ) has_integral i2) {0..1}
    using assms
  by (auto simp: has_contour_integral)
  have i1: (( $\lambda x. (2*f (g1 (2*x))) * vector\_derivative g1 (at (2*x))$ ) has_integral
i1) {0..1/2}
    and i2: (( $\lambda x. (2*f (g2 (2*x - 1))) * vector\_derivative g2 (at (2*x - 1))$ )
has_integral i2) {1/2..1}
  using has_integral_affinity01 [OF 1, where m= 2 and c=0, THEN has_integral_cmul
[where c=2]]
    has_integral_affinity01 [OF 2, where m= 2 and c=-1, THEN has_integral_cmul
[where c=2]]
  by (simp_all only: image_affinity_atLeastAtMost_div_diff, simp_all add:
scaleR_conv_of_real mult_ac)
  have g1: vector_derivative ( $\lambda x. if x*2 \leq 1$  then  $g1 (2*x)$  else  $g2 (2*x - 1)$ )
(at z) =
     $2 *_{\mathbb{R}} vector\_derivative g1 (at (z*2))$ 
    if  $0 \leq z$   $z*2 < 1$   $z*2 \notin s1$  for  $z$ 
proof (rule vector_derivative_at [OF has_vector_derivative_transform_within])
  show  $0 < |z - 1/2|$ 
    using that by auto
  have (*) 2 has_vector_derivative 2) (at z)
    by (simp add: has_vector_derivative_def has_derivative_def bounded_linear_mult_left)
  moreover have (g1 has_vector_derivative vector_derivative g1 (at (z * 2)))

```

```

(at (2 * z))
  using s1 that by (auto simp: algebra_simps vector_derivative_works)
  ultimately
  show (( $\lambda x. g1 (2 * x)$ ) has_vector_derivative 2 *R vector_derivative g1 (at
(z * 2))) (at z)
  by (intro vector_diff_chain_at [simplified o_def])
  qed (use that in <simp_all add: dist_real_def abs_if_split: if_split_asm>)

  have g2: vector_derivative ( $\lambda x. \text{if } x \leq 1 \text{ then } g1 (2*x) \text{ else } g2 (2*x - 1)$ )
(at z) =
    2 *R vector_derivative g2 (at (z * 2 - 1))
  if 1 < z * 2 z ≤ 1 z * 2 - 1 ∉ s2 for z
  proof (rule vector_derivative_at [OF has_vector_derivative_transform_within])
    show 0 < |z - 1/2|
      using that by auto
    have (( $\lambda x. 2 * x - 1$ ) has_vector_derivative 2) (at z)
    by (simp add: has_vector_derivative_def has_derivative_def bounded_linear_mult_left)
    moreover have (g2 has_vector_derivative vector_derivative g2 (at (z * 2 -
1))) (at (2 * z - 1))
      using s2 that by (auto simp: algebra_simps vector_derivative_works)
    ultimately
    show (( $\lambda x. g2 (2 * x - 1)$ ) has_vector_derivative 2 *R vector_derivative g2
(at (z * 2 - 1))) (at z)
    by (intro vector_diff_chain_at [simplified o_def])
    qed (use that in <simp_all add: dist_real_def abs_if_split: if_split_asm>)

  have (( $\lambda x. f ((g1 +++ g2) x) * \text{vector\_derivative } (g1 +++ g2) (at x)$ ) has_integral
i1) {0..1/2}
  proof (rule has_integral_spike_finite [OF _ _ i1])
    show finite (insert (1/2) ((* ) 2 - ' s1))
      using s1 by (force intro: finite_vimageI [where h = (* ) 2] inj_onI)
    qed (auto simp add: joinpaths_def scaleR_conv_of_real mult_ac g1)
    moreover have (( $\lambda x. f ((g1 +++ g2) x) * \text{vector\_derivative } (g1 +++ g2) (at
x)$ ) has_integral i2) {1/2..1}
    proof (rule has_integral_spike_finite [OF _ _ i2])
      show finite (insert (1/2) (( $\lambda x. 2 * x - 1$ ) - ' s2))
        using s2 by (force intro: finite_vimageI [where h =  $\lambda x. 2*x-1$ ] inj_onI)
      qed (auto simp add: joinpaths_def scaleR_conv_of_real mult_ac g2)
    ultimately
    show ?thesis
      by (simp add: has_contour_integral has_integral_combine [where c = 1/2])
  qed

```

lemma contour_integrable_joinI:

```

assumes f contour_integrable_on g1 f contour_integrable_on g2
      valid_path g1 valid_path g2
  shows f contour_integrable_on (g1 +++ g2)
  using assms
  by (meson has_contour_integral_join contour_integrable_on_def)

```

```

lemma contour_integrable_joinD1:
  assumes f contour_integrable_on (g1 +++ g2) valid_path g1
  shows f contour_integrable_on g1
proof –
  obtain s1
    where s1: finite s1  $\forall x \in \{0..1\} - s1$ . g1 differentiable at x
    using assms by (auto simp: valid_path_def piecewise_C1_differentiable_on_def
C1_differentiable_on_eq)
    have  $(\lambda x. f ((g1 +++ g2) (x/2)) * \text{vector\_derivative } (g1 +++ g2) \text{ (at } (x/2)))$ 
integrable_on  $\{0..1\}$ 
    using assms integrable_affinity [of _ 0 1/2 1/2 0] integrable_on_subcbx
[where a=0 and b=1/2]
    by (fastforce simp: contour_integrable_on)
    then have *:  $(\lambda x. (f ((g1 +++ g2) (x/2)))/2) * \text{vector\_derivative } (g1 +++$ 
g2) at  $(x/2))$  integrable_on  $\{0..1\}$ 
    by (auto dest: integrable_cmul [where c=1/2] simp: scaleR_conv_of_real)
    have g1: vector_derivative  $(\lambda x. \text{if } x*2 \leq 1 \text{ then } g1 (2*x) \text{ else } g2 (2*x - 1))$ 
at  $(z/2)$  =
       $2 *_{\mathbb{R}} \text{vector\_derivative } g1 \text{ (at } z)$ 
    if  $0 < z < 1$   $z \notin s1$  for z
    proof (rule vector_derivative_at [OF has_vector_derivative_transform_within])
      show  $0 < |(z - 1)/2|$ 
      using that by auto
      have §:  $(\lambda x. x * 2)$  has_vector_derivative 2 at  $(z/2)$ 
      using s1 by (auto simp: has_vector_derivative_def has_derivative_def
bounded_linear_mult_left)
      have  $(g1 \text{ has\_vector\_derivative } \text{vector\_derivative } g1 \text{ (at } z))$  at z
      using s1 that by (auto simp: vector_derivative_works)
      then show  $(\lambda x. g1 (2 * x))$  has_vector_derivative  $2 *_{\mathbb{R}} \text{vector\_derivative } g1$ 
at  $(z)$  at  $(z/2)$ 
      using vector_diff_chain_at [OF §] by (auto simp: field_simps o_def)
    qed (use that in  $\langle \text{simp\_all add: field\_simps dist\_real\_def abs\_if\_split: if\_split\_asm} \rangle$ )
    have fin01: finite  $(\{0, 1\} \cup s1)$ 
    by (simp add: s1)
    show ?thesis
    unfolding contour_integrable_on
    by (intro integrable_spike_finite [OF fin01 _ *]) (auto simp: joinpaths_def
scaleR_conv_of_real g1)
  qed

```

```

lemma contour_integrable_joinD2:
  assumes f contour_integrable_on (g1 +++ g2) valid_path g2
  shows f contour_integrable_on g2
proof –
  obtain s2
    where s2: finite s2  $\forall x \in \{0..1\} - s2$ . g2 differentiable at x
    using assms by (auto simp: valid_path_def piecewise_C1_differentiable_on_def
C1_differentiable_on_eq)

```

```

have ( $\lambda x. f ((g1 +++ g2) (x/2 + 1/2)) * \text{vector\_derivative } (g1 +++ g2) \text{ (at } (x/2 + 1/2))$ ) integrable_on {0..1}
  using assms integrable_affinity [of_ 1/2::real 1 1/2 1/2]
    integrable_on_subcbox [where a=1/2 and b=1]
  by (fastforce simp: contour_integrable_on image_affinity_atLeastAtMost_diff)
then have *: ( $\lambda x. (f ((g1 +++ g2) (x/2 + 1/2))/2) * \text{vector\_derivative } (g1 +++ g2) \text{ (at } (x/2 + 1/2))$ )
  integrable_on {0..1}
  by (auto dest: integrable_cmul [where c=1/2] simp: scaleR_conv_of_real)
have g2: vector_derivative ( $\lambda x. \text{if } x*2 \leq 1 \text{ then } g1 (2*x) \text{ else } g2 (2*x - 1)$ )
  (at ( $z/2+1/2$ )) =
   $2 *_{\mathbb{R}} \text{vector\_derivative } g2 \text{ (at } z)$ 
  if  $0 < z < 1$   $z \notin s2$  for  $z$ 
proof (rule vector_derivative_at [OF has_vector_derivative_transform_within])
  show  $0 < |z/2|$ 
  using that by auto
  have  $\S$ : ( $\lambda x. x * 2 - 1$ ) has_vector_derivative  $2$  (at ( $(1 + z)/2$ ))
  using  $s2$  by (auto simp: has_vector_derivative_def has_derivative_def
bounded_linear_mult_left)
  have ( $g2$  has_vector_derivative vector_derivative  $g2$  (at  $z$ )) (at  $z$ )
  using  $s2$  that by (auto simp: vector_derivative_works)
  then show ( $\lambda x. g2 (2*x - 1)$ ) has_vector_derivative  $2 *_{\mathbb{R}} \text{vector\_derivative } g2$ 
  (at  $z$ ) (at ( $z/2 + 1/2$ ))
  using vector_diff_chain_at [OF  $\S$ ] by (auto simp: field_simps o_def)
qed (use that in <simp_all add: field_simps dist_real_def abs_if_split: if_split_asm>)
have fin01: finite ( $\{0, 1\} \cup s2$ )
  by (simp add: s2)
show ?thesis
  unfolding contour_integrable_on
  by (intro integrable_spike_finite [OF fin01 _ *] (auto simp: joinpaths_def
scaleR_conv_of_real g2))
qed

```

```

lemma contour_integrable_join [simp]:
   $\llbracket \text{valid\_path } g1; \text{valid\_path } g2 \rrbracket$ 
   $\implies f \text{ contour\_integrable\_on } (g1 +++ g2) \iff f \text{ contour\_integrable\_on } g1$ 
 $\wedge f \text{ contour\_integrable\_on } g2$ 
using contour_integrable_joinD1 contour_integrable_joinD2 contour_integrable_joinI
by blast

```

```

lemma contour_integral_join [simp]:
   $\llbracket f \text{ contour\_integrable\_on } g1; f \text{ contour\_integrable\_on } g2; \text{valid\_path } g1; \text{valid\_path } g2 \rrbracket$ 
   $\implies \text{contour\_integral } (g1 +++ g2) f = \text{contour\_integral } g1 f + \text{contour\_integral } g2 f$ 
  by (simp add: has_contour_integral_integral has_contour_integral_join contour_integral_unique)

```

1.4 Shifting the starting point of a (closed) path

```

lemma has_contour_integral_shiftpath:
  assumes f: (f has_contour_integral i) g valid_path g
    and a: a ∈ {0..1}
  shows (f has_contour_integral i) (shiftpath a g)
proof -
  obtain S
  where S: finite S and g: ∀ x∈{0..1} - S. g differentiable at x
  using assms by (auto simp: valid_path_def piecewise_C1_differentiable_on_def
    C1_differentiable_on_eq)
  have *: ((λx. f (g x) * vector_derivative g (at x)) has_integral i) {0..1}
  using assms by (auto simp: has_contour_integral)
  then have i: i = integral {a..1} (λx. f (g x) * vector_derivative g (at x)) +
    integral {0..a} (λx. f (g x) * vector_derivative g (at x))
  apply (rule has_integral_unique)
  apply (subst add_commute)
  apply (subst Henstock_Kurzweil_Integration.integral_combine)
  using assms * integral_unique by auto

  have vd1: vector_derivative (shiftpath a g) (at x) = vector_derivative g (at (x
+ a))
  if 0 ≤ x + a < 1 x ∉ (λx. x - a) ' S for x
  unfolding shiftpath_def
proof (rule vector_derivative_at [OF has_vector_derivative_transform_within])
  have ((λx. g (x + a)) has_vector_derivative vector_derivative g (at (a + x)))
(at x)
  proof (rule vector_diff_chain_at [of _ 1, simplified o_def scaleR_one])
    show ((λx. x + a) has_vector_derivative 1) (at x)
    by (rule derivative_eq_intros | simp)+
    have g differentiable at (x + a)
    using g a that by force
    then show (g has_vector_derivative vector_derivative g (at (a + x))) (at (x
+ a))
    by (metis add_commute vector_derivative_works)
  qed
  then
  show ((λx. g (a + x)) has_vector_derivative vector_derivative g (at (x + a)))
(at x)
  by (auto simp: field_simps)
  show 0 < dist (1 - a) x
  using that by auto
qed (use that in ⟨auto simp: dist_real_def⟩)

  have vd2: vector_derivative (shiftpath a g) (at x) = vector_derivative g (at (x
+ a - 1))
  if x ≤ 1 1 < x + a x ∉ (λx. x - a + 1) ' S for x
  unfolding shiftpath_def
proof (rule vector_derivative_at [OF has_vector_derivative_transform_within])
  have ((λx. g (x + a - 1)) has_vector_derivative vector_derivative g (at

```

```

(a+x-1))) (at x)
  proof (rule vector_diff_chain_at [of _ 1, simplified o_def scaleR_one])
    show (( $\lambda x. x + a - 1$ ) has_vector_derivative 1) (at x)
      by (rule derivative_eq_intros | simp)+
    have g differentiable at (x+a-1)
      using g a that by force
    then show (g has_vector_derivative vector_derivative g (at (a+x-1))) (at
(x + a - 1))
      by (metis add.commute vector_derivative_works)
    qed
    then show (( $\lambda x. g (a + x - 1)$ ) has_vector_derivative vector_derivative g (at
(x + a - 1))) (at x)
      by (auto simp: field_simps)
    show 0 < dist (1 - a) x
      using that by auto
    qed (use that in ‹auto simp: dist_real_def›)

have va1: ( $\lambda x. f (g x) * \text{vector\_derivative } g \text{ (at } x)$ ) integrable_on ({a..1})
  using * a by (fastforce intro: integrable_subinterval_real)
have v0a: ( $\lambda x. f (g x) * \text{vector\_derivative } g \text{ (at } x)$ ) integrable_on ({0..a})
  using * a by (force intro: integrable_subinterval_real)
have finite ({1 - a}  $\cup$  ( $\lambda x. x - a$ ) ‘ S)
  using S by blast
then have (( $\lambda x. f (\text{shiftpath } a \ g \ x) * \text{vector\_derivative } (\text{shiftpath } a \ g) \text{ (at } x)$ )
  has_integral integral {a..1} ( $\lambda x. f (g x) * \text{vector\_derivative } g \text{ (at } x)$ )) {0..1
- a}
  apply (rule has_integral_spike_finite
[where f =  $\lambda x. f (g(a+x)) * \text{vector\_derivative } g \text{ (at}(a+x))$ ])
  subgoal
    using a by (simp add: vd1) (force simp: shiftpath_def add.commute)
  subgoal
    using has_integral_affinity [where m=1 and c=a] integrable_integral [OF
va1]
    by (force simp add: add.commute)
  done
moreover
have finite ({1 - a}  $\cup$  ( $\lambda x. x - a + 1$ ) ‘ S)
  using S by blast
then have (( $\lambda x. f (\text{shiftpath } a \ g \ x) * \text{vector\_derivative } (\text{shiftpath } a \ g) \text{ (at } x)$ )
  has_integral integral {0..a} ( $\lambda x. f (g x) * \text{vector\_derivative } g \text{ (at } x)$ ))
{1 - a..1}
  apply (rule has_integral_spike_finite
[where f =  $\lambda x. f (g(a+x-1)) * \text{vector\_derivative } g \text{ (at}(a+x-1))$ ])
  subgoal
    using a by (simp add: vd2) (force simp: shiftpath_def add.commute)
  subgoal
    using has_integral_affinity [where m=1 and c=a-1, simplified, OF inte-
grable_integral [OF v0a]]
    by (force simp add: algebra_simps)

```

```

done
ultimately show ?thesis
using a
by (auto simp: i has_contour_integral intro: has_integral_combine [where c
= 1-a])
qed

lemma has_contour_integral_shiftpath_D:
  assumes (f has_contour_integral i) (shiftpath a g)
    valid_path g pathfinish g = pathstart g a  $\in \{0..1\}$ 
  shows (f has_contour_integral i) g
proof -
  obtain S
  where S: finite S and g:  $\forall x \in \{0..1\} - S. g$  differentiable at x
  using assms by (auto simp: valid_path_def piecewise_C1_differentiable_on_def
C1_differentiable_on_eq)
  { fix x
  assume x:  $0 < x < 1$   $x \notin S$ 
  then have gx: g differentiable at x
  using g by auto
  have §: shiftpath (1 - a) (shiftpath a g) differentiable at x
  using assms x
  by (intro differentiable_transform_within [OF gx, of min x (1-x)])
  (auto simp: dist_real_def shiftpath_shiftpath abs_if_split: if_split_asm)
  have vector_derivative g (at x within {0..1}) =
  vector_derivative (shiftpath (1 - a) (shiftpath a g)) (at x within {0..1})
  apply (rule vector_derivative_at_within_ivl
[OF has_vector_derivative_transform_within_open
[where f = (shiftpath (1 - a) (shiftpath a g)) and S =
{0 < .. < 1} - S]])
  using S assms x §
  apply (auto simp: finite_imp_closed open_Diff shiftpath_shiftpath
at_within_interior [of _ {0..1}] vector_derivative_works
[symmetric])
  done
} note vd = this
have fi: (f has_contour_integral i) (shiftpath (1 - a) (shiftpath a g))
using assms by (auto intro!: has_contour_integral_shiftpath)
show ?thesis
unfolding has_contour_integral_def
proof (rule has_integral_spike_finite [of {0,1}  $\cup$  S, OF _ _ fi [unfolded
has_contour_integral_def]])
  show finite ({0, 1}  $\cup$  S)
  by (simp add: S)
qed (use S assms vd in <auto simp: shiftpath_shiftpath>)
qed

```

```

lemma has_contour_integral_shiftpath_eq:
  assumes valid_path g pathfinish g = pathstart g a  $\in \{0..1\}$ 

```

shows $(f \text{ has_contour_integral } i) (\text{shiftpath } a \ g) \longleftrightarrow (f \text{ has_contour_integral } i) \ g$
using *assms* *has_contour_integral_shiftpath* *has_contour_integral_shiftpath_D*
by *blast*

lemma *contour_integrable_on_shiftpath_eq*:

assumes *valid_path* *g* *pathfinish* $g = \text{pathstart } g \ a \in \{0..1\}$
shows $f \text{ contour_integrable_on } (\text{shiftpath } a \ g) \longleftrightarrow f \text{ contour_integrable_on } g$
using *assms* *contour_integrable_on_def* *has_contour_integral_shiftpath_eq* **by** *auto*

lemma *contour_integral_shiftpath*:

assumes *valid_path* *g* *pathfinish* $g = \text{pathstart } g \ a \in \{0..1\}$
shows $\text{contour_integral } (\text{shiftpath } a \ g) \ f = \text{contour_integral } g \ f$
using *assms*
by (*simp* *add: contour_integral_def* *contour_integrable_on_def* *has_contour_integral_shiftpath_eq*)

1.5 More about straight-line paths

lemma *has_contour_integral_linepath*:

shows $(f \text{ has_contour_integral } i) (\text{linepath } a \ b) \longleftrightarrow$
 $((\lambda x. f(\text{linepath } a \ b \ x) * (b - a)) \text{ has_integral } i) \ \{0..1\}$
by (*simp* *add: has_contour_integral*)

lemma *has_contour_integral_trivial [iff]*: $(f \text{ has_contour_integral } 0) (\text{linepath } a \ a)$

by (*simp* *add: has_contour_integral_linepath*)

lemma *has_contour_integral_trivial_iff [simp]*: $(f \text{ has_contour_integral } i) (\text{linepath } a \ a) \longleftrightarrow i=0$

using *has_contour_integral_unique* **by** *blast*

lemma *contour_integral_trivial [simp]*: $\text{contour_integral } (\text{linepath } a \ a) \ f = 0$

using *has_contour_integral_trivial* *contour_integral_unique* **by** *blast*

1.6 Relation to subpath construction

lemma *has_contour_integral_subpath_refl [iff]*: $(f \text{ has_contour_integral } 0) (\text{subpath } u \ u \ g)$

by (*simp* *add: has_contour_integral_subpath_def*)

lemma *contour_integrable_subpath_refl [iff]*: $f \text{ contour_integrable_on } (\text{subpath } u \ u \ g)$

using *has_contour_integral_subpath_refl* *contour_integrable_on_def* **by** *blast*

lemma *contour_integral_subpath_refl [simp]*: $\text{contour_integral } (\text{subpath } u \ u \ g) \ f = 0$

by (*simp* *add: contour_integral_unique*)

lemma *has_contour_integral_subpath*:


```

assumes f: f contour_integrable_on g and g: valid_path g
and uv: u ∈ {0..1} v ∈ {0..1} u ≤ v
shows (f has_contour_integral integral {u..v} (λx. f(g x) * vector_derivative
g (at x)))
      (subpath u v g)
proof (cases v=u)
case True
then show ?thesis
using f by (simp add: contour_integrable_on_def subpath_def has_contour_integral)
next
case False
obtain S where S: ∧x. x ∈ {0..1} - S ⇒ g differentiable at x and fs: finite
S
using g unfolding piecewise_C1_differentiable_on_def C1_differentiable_on_eq
valid_path_def by blast
have §: (λt. f (g t) * vector_derivative g (at t)) integrable_on {u..v}
using contour_integrable_on f integrable_on_subinterval uv by fastforce
then have *: ((λx. f (g ((v - u) * x + u)) * vector_derivative g (at ((v - u)
* x + u)))
      has_integral (1 / (v - u)) * integral {u..v} (λt. f (g t) * vector_derivative
g (at t)))
      {0..1}
using uv False unfolding has_integral_integral
apply simp
apply (drule has_integral_affinity [where m=v-u and c=u, simplified])
apply (simp_all add: image_affinity_atLeastAtMost_div_diff scaleR_conv_of_real)
apply (simp add: divide_simps)
done

have vd: vector_derivative (λx. g ((v-u) * x + u)) (at x) = (v-u) *R vec-
tor_derivative g (at ((v-u) * x + u))
if x ∈ {0..1} x ∉ (λt. (v-u) *R t + u) -‘ S for x
proof (rule vector_derivative_at [OF vector_diff_chain_at [simplified o_def]])
show ((λx. (v - u) * x + u) has_vector_derivative v - u) (at x)
by (intro derivative_eq_intros | simp)+
qed (use S uv mult_left_le [of x v-u] that in ⟨auto simp: vector_derivative_works⟩)

have fin: finite ((λt. (v - u) *R t + u) -‘ S)
using fs by (auto simp: inj_on_def False finite_vimageI)
show ?thesis
unfolding subpath_def has_contour_integral
apply (rule has_integral_spike_finite [OF fin])
using has_integral_cmul [OF *, where c = v-u] fs assms
by (auto simp: False vd scaleR_conv_of_real)
qed

lemma contour_integrable_subpath:
assumes f contour_integrable_on g valid_path g u ∈ {0..1} v ∈ {0..1}
shows f contour_integrable_on (subpath u v g)

```

```

proof (cases u v rule: linorder_class.le_cases)
  case le
  then show ?thesis
  by (metis contour_integrable_on_def has_contour_integral_subpath [OF assms])
next
  case ge
  with assms show ?thesis
  by (metis (no_types, lifting) contour_integrable_on_def contour_integrable_reversepath_eq
  has_contour_integral_subpath reversepath_subpath valid_path_subpath)
qed

```

```

lemma has_integral_contour_integral_subpath:
  assumes f contour_integrable_on g valid_path g u ∈ {0..1} v ∈ {0..1} u ≤ v
  shows (((λx. f(g x) * vector_derivative g (at x)))
  has_integral contour_integral (subpath u v g) f) {u..v}
  using assms
proof -
  have (λr. f (g r) * vector_derivative g (at r)) integrable_on {u..v}
  by (metis (full_types) assms(1) assms(3) assms(4) atLeastAtMost_iff atLeas-
  tatMost_subset_iff contour_integrable_on integrable_on_subinterval)
  then have ((λr. f (g r) * vector_derivative g (at r)) has_integral integral {u..v}
  (λr. f (g r) * vector_derivative g (at r))) {u..v}
  by blast
  then show ?thesis
  by (metis (full_types) assms contour_integral_unique has_contour_integral_subpath)
qed

```

```

lemma contour_integral_subcontour_integral:
  assumes f contour_integrable_on g valid_path g u ∈ {0..1} v ∈ {0..1} u ≤ v
  shows contour_integral (subpath u v g) f =
  integral {u..v} (λx. f(g x) * vector_derivative g (at x))
  using assms has_contour_integral_subpath contour_integral_unique by blast

```

```

lemma contour_integral_subpath_combine_less:
  assumes f contour_integrable_on g valid_path g u ∈ {0..1} v ∈ {0..1} w ∈
  {0..1}
  u < v v < w
  shows contour_integral (subpath u v g) f + contour_integral (subpath v w g)
  f =
  contour_integral (subpath u w g) f

```

```

proof -
  have (λx. f (g x) * vector_derivative g (at x)) integrable_on {u..w}
  using integrable_on_subcbox [where a=u and b=w and S = {0..1}] assms
  by (auto simp: contour_integrable_on)
  with assms show ?thesis
  by (auto simp: contour_integral_subcontour_integral Henstock_Kurzweil_Integration.integral_combine)
qed

```

```

lemma contour_integral_subpath_combine:

```

assumes f *contour_integrable_on* g *valid_path* g $u \in \{0..1\}$ $v \in \{0..1\}$ $w \in \{0..1\}$

shows $\text{contour_integral (subpath } u \ v \ g) f + \text{contour_integral (subpath } v \ w \ g)$
 $f =$

$\text{contour_integral (subpath } u \ w \ g) f$

proof (*cases* $u \neq v \wedge v \neq w \wedge u \neq w$)

case *True*

have $*$: $\text{subpath } v \ u \ g = \text{reversepath (subpath } u \ v \ g) \wedge$
 $\text{subpath } w \ u \ g = \text{reversepath (subpath } u \ w \ g) \wedge$
 $\text{subpath } w \ v \ g = \text{reversepath (subpath } v \ w \ g)$

by (*auto simp: reversepath_subpath*)

have $u < v \wedge v < w \vee$

$u < w \wedge w < v \vee$

$v < u \wedge u < w \vee$

$v < w \wedge w < u \vee$

$w < u \wedge u < v \vee$

$w < v \wedge v < u$

using *True* *assms* **by** *linarith*

with *assms* **show** *?thesis*

using *contour_integral_subpath_combine_less* [*of* $f \ g \ u \ v \ w$]

contour_integral_subpath_combine_less [*of* $f \ g \ u \ w \ v$]

contour_integral_subpath_combine_less [*of* $f \ g \ v \ u \ w$]

contour_integral_subpath_combine_less [*of* $f \ g \ v \ w \ u$]

contour_integral_subpath_combine_less [*of* $f \ g \ w \ u \ v$]

contour_integral_subpath_combine_less [*of* $f \ g \ w \ v \ u$]

by (*elim disjE*) (*auto simp: * contour_integral_reversepath contour_integrable_subpath valid_path_subpath algebra_simps*)

next

case *False*

with *assms* **show** *?thesis*

by (*metis add.right_neutral contour_integral_reversepath contour_integral_subpath_refl diff_0_eq_diff_eq add_0 reversepath_subpath valid_path_subpath*)

qed

lemma *contour_integral_integral*:

$\text{contour_integral } g \ f = \text{integral } \{0..1\} (\lambda x. f (g \ x) * \text{vector_derivative } g \ (\text{at } x))$

by (*simp add: contour_integral_def integral_def has_contour_integral contour_integrable_on*)

lemma *contour_integral_cong*:

assumes $g = g' \wedge x. x \in \text{path_image } g \implies f \ x = f' \ x$

shows $\text{contour_integral } g \ f = \text{contour_integral } g' \ f'$

unfolding *contour_integral_integral* **using** *assms*

by (*intro integral_cong*) (*auto simp: path_image_def*)

lemma *contour_integral_spike_finite_simple_path*:

assumes *finite* A *simple_path* $g \ g = g' \wedge x. x \in \text{path_image } g - A \implies f \ x = f' \ x$

shows $\text{contour_integral } g \ f = \text{contour_integral } g' \ f'$

```

unfolding contour_integral_integral
proof (rule integral_spike)
  have finite (g - ` A ∩ {0<..<1}) using ⟨simple_path g⟩ ⟨finite A⟩
    by (intro finite_vimage_IntI simple_path_inj_on) auto
  hence finite ({0, 1} ∪ g - ` A ∩ {0<..<1}) by auto
  thus negligible ({0, 1} ∪ g - ` A ∩ {0<..<1}) by (rule negligible_finite)
next
  fix x assume x ∈ {0..1} - ({0, 1} ∪ g - ` A ∩ {0<..<1})
  hence g x ∈ path_image g - A by (auto simp: path_image_def)
  from assms(4)[OF this] and assms(3)
    show f' (g' x) * vector_derivative g' (at x) = f (g x) * vector_derivative g (at x) by simp
  qed

```

Contour integral along a segment on the real axis

```

lemma has_contour_integral_linepath_Reals_iff:
  fixes a b :: complex and f :: complex ⇒ complex
  assumes a ∈ Reals b ∈ Reals Re a < Re b
  shows (f has_contour_integral I) (linepath a b) ↔
    ((λx. f (of_real x)) has_integral I) {Re a..Re b}
proof -
  from assms have [simp]: of_real (Re a) = a of_real (Re b) = b
    by (simp_all add: complex_eq_iff)
  from assms have a ≠ b by auto
  have ((λx. f (of_real x)) has_integral I) (cbox (Re a) (Re b)) ↔
    ((λx. f (a + b * of_real x - a * of_real x)) has_integral I /R (Re b - Re a)) {0..1}
    by (subst has_integral_affinity_iff [of Re b - Re a _ Re a, symmetric])
      (insert assms, simp_all add: field_simps scaleR_conv_of_real)
  also have (λx. f (a + b * of_real x - a * of_real x)) =
    (λx. (f (a + b * of_real x - a * of_real x) * (b - a)) /R (Re b - Re a))
    by (simp)
  using ⟨a ≠ b⟩ by (auto simp: field_simps fun_eq_iff scaleR_conv_of_real)
  also have (... has_integral I /R (Re b - Re a)) {0..1} ↔
    ((λx. f (linepath a b x) * (b - a)) has_integral I) {0..1} using assms
    by (subst has_integral_cmul_iff) (auto simp: linepath_def scaleR_conv_of_real algebra_simps)
  also have ... ↔ (f has_contour_integral I) (linepath a b) unfolding has_contour_integral_def
    by (intro has_integral_cong) (simp add: vector_derivative_linepath_within)
  finally show ?thesis by simp
qed

```

```

lemma contour_integrable_linepath_Reals_iff:
  fixes a b :: complex and f :: complex ⇒ complex
  assumes a ∈ Reals b ∈ Reals Re a < Re b
  shows (f contour_integrable_on linepath a b) ↔
    (λx. f (of_real x)) integrable_on {Re a..Re b}
  using has_contour_integral_linepath_Reals_iff [OF assms, of f]
  by (auto simp: contour_integrable_on_def integrable_on_def)

```

```

lemma contour_integral_linepath_Reals_eq:
  fixes a b :: complex and f :: complex  $\Rightarrow$  complex
  assumes a  $\in$  Reals b  $\in$  Reals Re a < Re b
  shows contour_integral (linepath a b) f = integral {Re a..Re b} ( $\lambda$ x. f (of_real x))
proof (cases f contour_integrable_on linepath a b)
  case True
  thus ?thesis using has_contour_integral_linepath_Reals_iff[OF assms, of f]
    using has_contour_integral_integral has_contour_integral_unique by blast
next
  case False
  thus ?thesis using contour_integrable_linepath_Reals_iff[OF assms, of f]
    by (simp add: not_integrable_contour_integral not_integrable_integral)
qed

```

1.7 Cauchy's theorem where there's a primitive

```

lemma contour_integral_primitive_lemma:
  fixes f :: complex  $\Rightarrow$  complex and g :: real  $\Rightarrow$  complex
  assumes a  $\leq$  b
    and  $\bigwedge$ x. x  $\in$  S  $\implies$  (f has_field_derivative f' x) (at x within S)
    and g piecewise_differentiable_on {a..b}  $\bigwedge$ x. x  $\in$  {a..b}  $\implies$  g x  $\in$  S
  shows (( $\lambda$ x. f'(g x) * vector_derivative g (at x within {a..b}))
    has_integral (f(g b) - f(g a))) {a..b}
proof -
  obtain K where finite K and K:  $\forall$  x $\in$ {a..b} - K. g differentiable (at x within {a..b})
  and cg: continuous_on {a..b} g
  using assms by (auto simp: piecewise_differentiable_on_def)
  have continuous_on (g ' {a..b}) f
  using assms
  by (metis field_differentiable_def field_differentiable_imp_continuous_at continuous_on_eq_continuous_within continuous_on_subset_image_subset_iff)
  then have cfg: continuous_on {a..b} ( $\lambda$ x. f (g x))
  by (rule continuous_on_compose [OF cg, unfolded o_def])
  { fix x::real
    assume a < x and b: x < b and xk: x  $\notin$  K
    then have g differentiable at x within {a..b}
      using K by (simp add: differentiable_at_withinI)
    then have (g has_vector_derivative vector_derivative g (at x within {a..b}))
      (at x within {a..b})
      by (simp add: vector_derivative_works has_field_derivative_def scaleR_conv_of_real)
    then have gdiff: (g has_derivative ( $\lambda$ u. u * vector_derivative g (at x within {a..b})))
      (at x within {a..b})
      by (simp add: has_vector_derivative_def scaleR_conv_of_real)
    have (f has_field_derivative (f' (g x))) (at (g x) within g ' {a..b})
      using assms by (metis a atLeastAtMost_iff b DERIV_subset_image_subset_iff less_eq_real_def)
    then have fdiff: (f has_derivative (*) (f' (g x))) (at (g x) within g ' {a..b})

```

```

    by (simp add: has_field_derivative_def)
    have (( $\lambda x. f (g x)$ ) has_vector_derivative  $f' (g x) * \text{vector\_derivative } g$  (at  $x$  within  $\{a..b\}$ )) (at  $x$  within  $\{a..b\}$ )
    using diff_chain_within [OF gdiff fdiff]
    by (simp add: has_vector_derivative_def scaleR_conv_of_real o_def mult_ac)
  } note * = this
  show ?thesis
    using assms cfg *
    by (force simp: at_within_Icc_at intro: fundamental_theorem_of_calculus_interior_strong [OF <finite K>])
qed

```

lemma *contour_integral_primitive:*

```

  assumes  $\bigwedge x. x \in S \implies (f \text{ has\_field\_derivative } f' x)$  (at  $x$  within  $S$ )
    and valid_path  $g$  path_image  $g \subseteq S$ 
    shows ( $f'$  has_contour_integral ( $f(\text{pathfinish } g) - f(\text{pathstart } g)$ ))  $g$ 
  using assms
  apply (simp add: valid_path_def path_image_def pathfinish_def pathstart_def
    has_contour_integral_def)
  apply (auto intro!: piecewise_C1_imp_differentiable contour_integral_primitive_lemma [of 0 1 S])
  done

```

corollary *Cauchy_theorem_primitive:*

```

  assumes  $\bigwedge x. x \in S \implies (f \text{ has\_field\_derivative } f' x)$  (at  $x$  within  $S$ )
    and valid_path  $g$  path_image  $g \subseteq S$  pathfinish  $g = \text{pathstart } g$ 
    shows ( $f'$  has_contour_integral 0)  $g$ 
  using assms by (metis diff_self contour_integral_primitive)

```

Existence of path integral for continuous function

lemma *contour_integrable_continuous_linepath:*

```

  assumes continuous_on (closed_segment  $a$   $b$ )  $f$ 
  shows  $f$  contour_integrable_on (linepath  $a$   $b$ )
  proof -
    have continuous_on (closed_segment  $a$   $b$ ) ( $\lambda x. f x * (b - a)$ )
      by (rule continuous_intros | simp add: assms)+
    then have continuous_on {0..1} ( $\lambda x. f (\text{linepath } a \ b \ x) * (b - a)$ )
      by (metis (no_types, lifting) continuous_on_compose continuous_on_cong
        continuous_on_linepath linepath_image_01 o_apply)
    then have ( $\lambda x. f (\text{linepath } a \ b \ x) * \text{vector\_derivative } (\text{linepath } a \ b)$ 
      (at  $x$  within {0..1})) integrable_on {0..1}
      by (metis (no_types, lifting) continuous_on_cong integrable_continuous_real
        vector_derivative_linepath_within)
    then show ?thesis
      by (simp add: contour_integrable_on_def has_contour_integral_def integrable_on_def [symmetric])
  qed

```

lemma *has_field_der_id*: $((\lambda x. x^2/2) \text{ has_field_derivative } x) \text{ (at } x)$
by (*rule has_derivative_imp_has_field_derivative*)
(rule derivative_intros | simp)+

lemma *contour_integral_id* [*simp*]: $\text{contour_integral (linepath } a \ b) (\lambda y. y) = (b^2 - a^2)/2$
using *contour_integral_primitive* [*of UNIV* $\lambda x. x^2/2$ $\lambda x. x$ *linepath* *a b*] *contour_integral_unique*
by (*simp add: has_field_der_id*)

lemma *contour_integrable_on_const* [*iff*]: $(\lambda x. c) \text{ contour_integrable_on (linepath } a \ b)$
by (*simp add: contour_integrable_continuous_linepath*)

lemma *contour_integrable_on_id* [*iff*]: $(\lambda x. x) \text{ contour_integrable_on (linepath } a \ b)$
by (*simp add: contour_integrable_continuous_linepath*)

1.8 Arithmetical combining theorems

lemma *has_contour_integral_neg*:
 $(f \text{ has_contour_integral } i) \ g \implies ((\lambda x. -(f \ x)) \text{ has_contour_integral } (-i)) \ g$
by (*simp add: has_integral_neg has_contour_integral_def*)

lemma *has_contour_integral_add*:
 $\llbracket (f1 \text{ has_contour_integral } i1) \ g; (f2 \text{ has_contour_integral } i2) \ g \rrbracket$
 $\implies ((\lambda x. f1 \ x + f2 \ x) \text{ has_contour_integral } (i1 + i2)) \ g$
by (*simp add: has_integral_add has_contour_integral_def algebra_simps*)

lemma *has_contour_integral_diff*:
 $\llbracket (f1 \text{ has_contour_integral } i1) \ g; (f2 \text{ has_contour_integral } i2) \ g \rrbracket$
 $\implies ((\lambda x. f1 \ x - f2 \ x) \text{ has_contour_integral } (i1 - i2)) \ g$
by (*simp add: has_integral_diff has_contour_integral_def algebra_simps*)

lemma *has_contour_integral_lmul*:
 $(f \text{ has_contour_integral } i) \ g \implies ((\lambda x. c * (f \ x)) \text{ has_contour_integral } (c*i)) \ g$
by (*simp add: has_contour_integral_def algebra_simps has_integral_mult_right*)

lemma *has_contour_integral_rmul*:
 $(f \text{ has_contour_integral } i) \ g \implies ((\lambda x. (f \ x) * c) \text{ has_contour_integral } (i*c)) \ g$
by (*simp add: mult.commute has_contour_integral_lmul*)

lemma *has_contour_integral_div*:
 $(f \text{ has_contour_integral } i) \ g \implies ((\lambda x. f \ x / c) \text{ has_contour_integral } (i/c)) \ g$
by (*simp add: field_class.field_divide_inverse*) (*metis has_contour_integral_rmul*)

lemma *has_contour_integral_eq*:
 $\llbracket (f \text{ has_contour_integral } y) \ p; \bigwedge x. x \in \text{path_image } p \implies f \ x = g \ x \rrbracket \implies (g$

has_contour_integral y) p

by (metis (mono_tags, lifting) has_contour_integral_def has_integral_eq image_eqI path_image_def)

lemma *has_contour_integral_bound_linepath*:

assumes (f *has_contour_integral* i) (linepath a b)

$0 \leq B$ and $B: \bigwedge x. x \in \text{closed_segment } a \ b \implies \text{norm}(f \ x) \leq B$

shows $\text{norm } i \leq B * \text{norm}(b - a)$

proof –

have $\text{norm } i \leq (B * \text{norm}(b - a)) * \text{content } (\text{cbox } 0 \ (1::\text{real}))$

proof (rule *has_integral_bound*)

[of $\lambda x. f(\text{linepath } a \ b \ x) * \text{vector_derivative } (\text{linepath } a \ b)$ (at x within $\{0..1\}$)]]

show $\text{cmod}(f(\text{linepath } a \ b \ x) * \text{vector_derivative } (\text{linepath } a \ b)$ (at x within $\{0..1\}$))

$\leq B * \text{cmod}(b - a)$

if $x \in \text{cbox } 0 \ 1$ for $x::\text{real}$

using that *box_real(2)* *norm_mult*

by (metis *B linepath_in_path mult_right_mono norm_ge_zero vector_derivative_linepath_within*)

qed (use *assms has_contour_integral_def* in *auto*)

then show *?thesis*

by (*auto simp: content_real*)

qed

lemma *has_contour_integral_const_linepath*: $((\lambda x. c) \text{ has_contour_integral } c*(b - a))(\text{linepath } a \ b)$

unfolding *has_contour_integral_linepath*

by (metis *content_real diff_0_right has_integral_const_real lambda_one_of_real_1 scaleR_conv_of_real zero_le_one*)

lemma *has_contour_integral_0*: $((\lambda x. 0) \text{ has_contour_integral } 0) \ g$

by (*simp add: has_contour_integral_def*)

lemma *has_contour_integral_is_0*:

$(\bigwedge z. z \in \text{path_image } g \implies f \ z = 0) \implies (f \text{ has_contour_integral } 0) \ g$

by (rule *has_contour_integral_eq [OF has_contour_integral_0]*) *auto*

lemma *has_contour_integral_sum*:

$[\text{finite } s; \bigwedge a. a \in s \implies (f \ a \text{ has_contour_integral } i \ a) \ p]$

$\implies ((\lambda x. \text{sum } (\lambda a. f \ a \ x) \ s) \text{ has_contour_integral } \text{sum } i \ s) \ p$

by (*induction s rule: finite_induct*) (*auto simp: has_contour_integral_0 has_contour_integral_add*)

1.9 Operations on path integrals

lemma *contour_integral_const_linepath [simp]*: $\text{contour_integral } (\text{linepath } a \ b) (\lambda x. c) = c*(b - a)$

by (rule *contour_integral_unique [OF has_contour_integral_const_linepath]*)

lemma *contour_integral_neg*: $\text{contour_integral } g (\lambda z. -f \ z) = -\text{contour_integral } g \ f$

$g f$
by (*simp add: contour_integral_integral*)

lemma *contour_integral_add*:

$f1 \text{ contour_integrable_on } g \implies f2 \text{ contour_integrable_on } g \implies \text{contour_integral } g (\lambda x. f1 x + f2 x) =$
 $\text{contour_integral } g f1 + \text{contour_integral } g f2$

by (*simp add: contour_integral_unique has_contour_integral_integral has_contour_integral_add*)

lemma *contour_integral_diff*:

$f1 \text{ contour_integrable_on } g \implies f2 \text{ contour_integrable_on } g \implies \text{contour_integral } g (\lambda x. f1 x - f2 x) =$
 $\text{contour_integral } g f1 - \text{contour_integral } g f2$

by (*simp add: contour_integral_unique has_contour_integral_integral has_contour_integral_diff*)

lemma *contour_integral_lmul*:

shows $f \text{ contour_integrable_on } g$
 $\implies \text{contour_integral } g (\lambda x. c * f x) = c * \text{contour_integral } g f$

by (*simp add: contour_integral_unique has_contour_integral_integral has_contour_integral_lmul*)

lemma *contour_integral_rmul*:

shows $f \text{ contour_integrable_on } g$
 $\implies \text{contour_integral } g (\lambda x. f x * c) = \text{contour_integral } g f * c$

by (*simp add: contour_integral_unique has_contour_integral_integral has_contour_integral_rmul*)

lemma *contour_integral_div*:

shows $f \text{ contour_integrable_on } g$
 $\implies \text{contour_integral } g (\lambda x. f x / c) = \text{contour_integral } g f / c$

by (*simp add: contour_integral_unique has_contour_integral_integral has_contour_integral_div*)

lemma *contour_integral_eq*:

$(\bigwedge x. x \in \text{path_image } p \implies f x = g x) \implies \text{contour_integral } p f = \text{contour_integral } p g$

using *contour_integral_cong contour_integral_def* **by** *fastforce*

lemma *contour_integral_eq_0*:

$(\bigwedge z. z \in \text{path_image } g \implies f z = 0) \implies \text{contour_integral } g f = 0$

by (*simp add: has_contour_integral_is_0 contour_integral_unique*)

lemma *contour_integral_bound_linepath*:

shows
 $\llbracket f \text{ contour_integrable_on } (\text{linepath } a b);$
 $0 \leq B; \bigwedge x. x \in \text{closed_segment } a b \implies \text{norm}(f x) \leq B \rrbracket$
 $\implies \text{norm}(\text{contour_integral } (\text{linepath } a b) f) \leq B * \text{norm}(b - a)$

by (*meson has_contour_integral_bound_linepath has_contour_integral_integral*)

lemma *contour_integral_0 [simp]*: $\text{contour_integral } g (\lambda x. 0) = 0$

by (*simp add: contour_integral_unique has_contour_integral_0*)

lemma *contour_integral_sum*:

$\llbracket \text{finite } s; \bigwedge a. a \in s \implies (f\ a)\ \text{contour_integrable_on } p \rrbracket$
 $\implies \text{contour_integral } p\ (\lambda x. \text{sum } (\lambda a. f\ a\ x)\ s) = \text{sum } (\lambda a. \text{contour_integral } p\ (f\ a))\ s$
by (*auto simp: contour_integral_unique has_contour_integral_sum has_contour_integral_integral*)

lemma *contour_integrable_eq*:

$\llbracket f\ \text{contour_integrable_on } p; \bigwedge x. x \in \text{path_image } p \implies f\ x = g\ x \rrbracket \implies g\ \text{contour_integrable_on } p$
unfolding *contour_integrable_on_def*
by (*metis has_contour_integral_eq*)

1.10 Arithmetic theorems for path integrability

lemma *contour_integrable_neg*:

$f\ \text{contour_integrable_on } g \implies (\lambda x. -(f\ x))\ \text{contour_integrable_on } g$
using *has_contour_integral_neg contour_integrable_on_def* **by** *blast*

lemma *contour_integrable_add*:

$\llbracket f1\ \text{contour_integrable_on } g; f2\ \text{contour_integrable_on } g \rrbracket \implies (\lambda x. f1\ x + f2\ x)\ \text{contour_integrable_on } g$
using *has_contour_integral_add contour_integrable_on_def*
by *fastforce*

lemma *contour_integrable_diff*:

$\llbracket f1\ \text{contour_integrable_on } g; f2\ \text{contour_integrable_on } g \rrbracket \implies (\lambda x. f1\ x - f2\ x)\ \text{contour_integrable_on } g$
using *has_contour_integral_diff contour_integrable_on_def*
by *fastforce*

lemma *contour_integrable_lmul*:

$f\ \text{contour_integrable_on } g \implies (\lambda x. c * f\ x)\ \text{contour_integrable_on } g$
using *has_contour_integral_lmul contour_integrable_on_def*
by *fastforce*

lemma *contour_integrable_rmul*:

$f\ \text{contour_integrable_on } g \implies (\lambda x. f\ x * c)\ \text{contour_integrable_on } g$
using *has_contour_integral_rmul contour_integrable_on_def*
by *fastforce*

lemma *contour_integrable_div*:

$f\ \text{contour_integrable_on } g \implies (\lambda x. f\ x / c)\ \text{contour_integrable_on } g$
using *has_contour_integral_div contour_integrable_on_def*
by *fastforce*

lemma *contour_integrable_sum*:

$\llbracket \text{finite } s; \bigwedge a. a \in s \implies (f\ a)\ \text{contour_integrable_on } p \rrbracket$
 $\implies (\lambda x. \text{sum } (\lambda a. f\ a\ x)\ s)\ \text{contour_integrable_on } p$
unfolding *contour_integrable_on_def*

by (*metis has_contour_integral_sum*)

lemma *contour_integrable_neg_iff*:

$(\lambda x. -f x) \text{ contour_integrable_on } g \longleftrightarrow f \text{ contour_integrable_on } g$
using *contour_integrable_neg[of f g] contour_integrable_neg[of $\lambda x. -f x g$]* **by**
auto

lemma *contour_integrable_lmul_iff*:

$c \neq 0 \implies (\lambda x. c * f x) \text{ contour_integrable_on } g \longleftrightarrow f \text{ contour_integrable_on } g$
using *contour_integrable_lmul[of f g c] contour_integrable_lmul[of $\lambda x. c * f x g$ inverse c]*
by (*auto simp: field_simps*)

lemma *contour_integrable_rmul_iff*:

$c \neq 0 \implies (\lambda x. f x * c) \text{ contour_integrable_on } g \longleftrightarrow f \text{ contour_integrable_on } g$
using *contour_integrable_rmul[of f g c] contour_integrable_rmul[of $\lambda x. c * f x g$ inverse c]*
by (*auto simp: field_simps*)

lemma *contour_integrable_div_iff*:

$c \neq 0 \implies (\lambda x. f x / c) \text{ contour_integrable_on } g \longleftrightarrow f \text{ contour_integrable_on } g$
using *contour_integrable_rmul_iff[of inverse c]* **by** (*simp add: field_simps*)

1.11 Reversing a path integral

lemma *has_contour_integral_reverse_linepath*:

$(f \text{ has_contour_integral } i) (\text{linepath } a \ b) \implies (f \text{ has_contour_integral } (-i)) (\text{linepath } b \ a)$
using *has_contour_integral_reversepath valid_path_linepath* **by** *fastforce*

lemma *contour_integral_reverse_linepath*:

$\text{continuous_on } (\text{closed_segment } a \ b) \ f \implies \text{contour_integral } (\text{linepath } a \ b) \ f = - (\text{contour_integral } (\text{linepath } b \ a) \ f)$
using *contour_integral_reversepath* **by** *fastforce*

Splitting a path integral in a flat way.*)

lemma *has_contour_integral_split*:

assumes *f*: $(f \text{ has_contour_integral } i) (\text{linepath } a \ c) (f \text{ has_contour_integral } j) (\text{linepath } c \ b)$

and *k*: $0 \leq k \leq 1$

and *c*: $c - a = k *_R (b - a)$

shows $(f \text{ has_contour_integral } (i + j)) (\text{linepath } a \ b)$

proof (*cases k = 0 \vee k = 1*)

case *True*

then show *?thesis*

using *assms* **by** *auto*

```

next
case False
then have k: 0 < k k < 1 complex_of_real k ≠ 1
  using assms by auto
have c': c = k *R (b - a) + a
  by (metis diff_add_cancel c)
have bc: (b - c) = (1 - k) *R (b - a)
  by (simp add: algebra_simps c')
{ assume *: ((λx. f ((1 - x) *R a + x *R c) * (c - a)) has_integral i) {0..1}
  have ∧x. (x / k) *R a + ((k - x) / k) *R a = a
  using False by (simp add: field_split_simps flip: real_vector.scale_left_distrib)
  then have ∧x. ((k - x) / k) *R a + (x / k) *R c = (1 - x) *R a + x *R b
  using False by (simp add: c' algebra_simps)
  then have ((λx. f ((1 - x) *R a + x *R b) * (b - a)) has_integral i) {0..k}
  using k has_integral_affinity01 [OF *, of inverse k 0]
  by (force dest: has_integral_cmul [where c = inverse k]
      simp add: divide_simps mult.commute [of _ k] image_affinity_atLeastAtMost
  c)
} note fi = this
{ assume *: ((λx. f ((1 - x) *R c + x *R b) * (b - c)) has_integral j) {0..1}
  have **: ∧x. (((1 - x) / (1 - k)) *R c + ((x - k) / (1 - k)) *R b) = ((1 -
x) *R a + x *R b)
  using k unfolding c' scaleR_conv_of_real
  apply (simp add: divide_simps)
  apply (simp add: distrib_right distrib_left right_diff_distrib left_diff_distrib)
  done
  have ((λx. f ((1 - x) *R a + x *R b) * (b - a)) has_integral j) {k..1}
  using k has_integral_affinity01 [OF *, of inverse(1 - k) -(k/(1 - k))]
  apply (simp add: divide_simps mult.commute [of _ 1-k] image_affinity_atLeastAtMost
  ** bc)
  apply (auto dest: has_integral_cmul [where k = (1 - k) *R j and c =
inverse (1 - k)])
  done
} note fj = this
show ?thesis
  using f k unfolding has_contour_integral_linepath
  by (simp add: linepath_def has_integral_combine [OF _ _ fi fj])
qed

```

lemma *continuous_on_closed_segment_transform*:

```

assumes f: continuous_on (closed_segment a b) f
  and k: 0 ≤ k k ≤ 1
  and c: c - a = k *R (b - a)
shows continuous_on (closed_segment a c) f

```

proof –

```

have c': c = (1 - k) *R a + k *R b
  using c by (simp add: algebra_simps)
have closed_segment a c ⊆ closed_segment a b
  by (metis c' ends_in_segment(1) in_segment(1) k subset_closed_segment)

```

```

then show continuous_on (closed_segment a c) f
  by (rule continuous_on_subset [OF f])
qed

```

lemma *contour_integral_split*:

```

assumes f: continuous_on (closed_segment a b) f
  and k:  $0 \leq k \leq 1$ 
  and c:  $c - a = k *_R (b - a)$ 
shows contour_integral(linepath a b) f = contour_integral(linepath a c) f +
contour_integral(linepath c b) f
proof -
  have c':  $c = (1 - k) *_R a + k *_R b$ 
  using c by (simp add: algebra_simps)
  have closed_segment a c  $\subseteq$  closed_segment a b
  by (metis c' ends_in_segment(1) in_segment(1) k subset_closed_segment)
  moreover have closed_segment c b  $\subseteq$  closed_segment a b
  by (metis c' ends_in_segment(2) in_segment(1) k subset_closed_segment)
  ultimately
  have *: continuous_on (closed_segment a c) f continuous_on (closed_segment
c b) f
  by (auto intro: continuous_on_subset [OF f])
  show ?thesis
  by (rule contour_integral_unique) (meson * c contour_integrable_continuous_linepath
has_contour_integral_integral has_contour_integral_split k)
qed

```

lemma *contour_integral_split_linepath*:

```

assumes f: continuous_on (closed_segment a b) f
  and c:  $c \in$  closed_segment a b
shows contour_integral(linepath a b) f = contour_integral(linepath a c) f +
contour_integral(linepath c b) f
using c by (auto simp: closed_segment_def algebra_simps intro!: contour_integral_split
[OF f])

```

1.12 Reversing the order in a double path integral

The condition is stronger than needed but it's often true in typical situations

```

lemma fst_im_cbox [simp]:  $\text{cbox } c \ d \neq \{\} \implies (\text{fst } \text{'cbox } (a,c) (b,d)) = \text{cbox } a \ b$ 
by (auto simp: cbox_Pair_eq)

```

```

lemma snd_im_cbox [simp]:  $\text{cbox } a \ b \neq \{\} \implies (\text{snd } \text{'cbox } (a,c) (b,d)) = \text{cbox } c \ d$ 
by (auto simp: cbox_Pair_eq)

```

proposition *contour_integral_swap*:

```

assumes fcon: continuous_on (path_image g  $\times$  path_image h)  $(\lambda(y1,y2). f \ y1 \ y2)$ 
and vp: valid_path g valid_path h
and gvcon: continuous_on  $\{0..1\}$   $(\lambda t. \text{vector\_derivative } g \ (at \ t))$ 

```

```

    and hvcon: continuous_on {0..1} (λt. vector_derivative h (at t))
  shows contour_integral g (λw. contour_integral h (f w)) =
    contour_integral h (λz. contour_integral g (λw. f w z))
proof -
  have gcon: continuous_on {0..1} g and hcon: continuous_on {0..1} h
  using assms by (auto simp: valid_path_def piecewise_C1_differentiable_on_def)
  have fgh1: λx. (λt. f (g x) (h t)) = (λ(y1,y2). f y1 y2) ∘ (λt. (g x, h t))
  by (rule ext) simp
  have fgh2: λx. (λt. f (g t) (h x)) = (λ(y1,y2). f y1 y2) ∘ (λt. (g t, h x))
  by (rule ext) simp
  have fcon_im1: λx. 0 ≤ x ⇒ x ≤ 1 ⇒ continuous_on ((λt. (g x, h t)) ‘
    {0..1}) (λ(x, y). f x y)
  by (rule continuous_on_subset [OF fcon]) (auto simp: path_image_def)
  have fcon_im2: λx. 0 ≤ x ⇒ x ≤ 1 ⇒ continuous_on ((λt. (g t, h x)) ‘
    {0..1}) (λ(x, y). f x y)
  by (rule continuous_on_subset [OF fcon]) (auto simp: path_image_def)
  have continuous_on (cbox (0, 0) (1, 1::real)) ((λx. vector_derivative g (at x))
    ∘ fst)
  continuous_on (cbox (0, 0) (1::real, 1)) ((λx. vector_derivative h (at x)) ∘
    snd)
  by (rule continuous_intros | simp add: gvcon hvcon)+
  then have gvcon': continuous_on (cbox (0, 0) (1, 1::real)) (λz. vector_derivative
    g (at (fst z)))
  and hvcon': continuous_on (cbox (0, 0) (1::real, 1)) (λx. vector_derivative
    h (at (snd x)))
  by auto
  have continuous_on (cbox (0, 0) (1, 1)) ((λ(y1, y2). f y1 y2) ∘ (λw. ((g ∘ fst)
    w, (h ∘ snd) w)))
  apply (intro gcon hcon continuous_intros | simp)+
  apply (auto simp: path_image_def intro: continuous_on_subset [OF fcon])
  done
  then have fgh: continuous_on (cbox (0, 0) (1, 1)) (λx. f (g (fst x)) (h (snd
    x)))
  by auto
  have integral {0..1} (λx. contour_integral h (f (g x)) * vector_derivative g (at
    x)) =
    integral {0..1} (λx. contour_integral h (λy. f (g x) y * vector_derivative g
    (at x)))
  proof (rule integral_cong [OF contour_integral_rmul [symmetric]])
    have λx. x ∈ {0..1} ⇒
      continuous_on {0..1} (λxa. f (g x) (h xa))
    by (subst fgh1) (rule fcon_im1 hcon continuous_intros | simp)+
    then show λx. x ∈ {0..1} ⇒ f (g x) contour_integrable_on h
    unfolding contour_integrable_on
    using continuous_on_mult hvcon integrable_continuous_real by blast
  qed
  also have ... = integral {0..1}
    (λy. contour_integral g (λx. f x (h y)) * vector_derivative h (at
    y)))

```

```

  unfolding contour_integral_integral
  apply (subst integral_swap_continuous [where 'a = real and 'b = real, of 0
0 1 1, simplified])
  subgoal
    by (rule fgh gvcon' hvcon' continuous_intros | simp add: split_def)+
  subgoal
    unfolding integral_mult_left [symmetric]
    by (simp only: mult_ac)
  done
  also have ... = contour_integral h (λz. contour_integral g (λw. f w z))
  unfolding contour_integral_integral integral_mult_left [symmetric]
  by (simp add: algebra_simps)
  finally show ?thesis
  by (simp add: contour_integral_integral)
qed

```

```

lemma valid_path_negatepath: valid_path γ ⇒ valid_path (uminus ∘ γ)
  unfolding o_def using piecewise_C1_differentiable_neg valid_path_def by
blast

```

```

lemma has_contour_integral_negatepath:
  assumes γ: valid_path γ and cint: ((λz. f (- z)) has_contour_integral - i) γ
  shows (f has_contour_integral i) (uminus ∘ γ)
proof -
  obtain S where cont: continuous_on {0..1} γ and finite S and diff: γ C1_differentiable_on
{0..1} - S
  using γ by (auto simp: valid_path_def piecewise_C1_differentiable_on_def)
  have ((λx. - (f (- γ x) * vector_derivative γ (at x within {0..1}))) has_integral
i) {0..1}
  using cint by (auto simp: has_contour_integral_def dest: has_integral_neg)
  then
  have ((λx. f (- γ x) * vector_derivative (uminus ∘ γ) (at x within {0..1}))
has_integral i) {0..1}
  proof (rule rev_iffD1 [OF _ has_integral_spike_eq])
    show negligible S
    by (simp add: ⟨finite S⟩ negligible_finite)
    show f (- γ x) * vector_derivative (uminus ∘ γ) (at x within {0..1}) =
- (f (- γ x) * vector_derivative γ (at x within {0..1}))
    if x ∈ {0..1} - S for x
  proof -
    have vector_derivative (uminus ∘ γ) (at x within cbox 0 1) = - vec-
tor_derivative γ (at x within cbox 0 1)
    proof (rule vector_derivative_within_cbox)
      show (uminus ∘ γ has_vector_derivative - vector_derivative γ (at x within
cbox 0 1)) (at x within cbox 0 1)
      using that unfolding o_def
      by (metis C1_differentiable_on_eq UNIV_I diff differentiable_subset
has_vector_derivative_minus subsetI that vector_derivative_works)
    qed (use that in auto)
  end
end

```

```

    then show ?thesis
      by simp
  qed
qed
then show ?thesis by (simp add: has_contour_integral_def)
qed

```

lemma *contour_integrable_negatepath*:

```

  assumes  $\gamma$ : valid_path  $\gamma$  and  $\pi$ :  $(\lambda z. f (-z))$  contour_integrable_on  $\gamma$ 
  shows  $f$  contour_integrable_on  $(\text{uminus} \circ \gamma)$ 
  by (metis  $\gamma$  add.inverse_inverse contour_integrable_on_def has_contour_integral_negatepath
   $\pi$ )

```

lemma *C1_differentiable_polynomial_function*:

```

  fixes  $p :: \text{real} \Rightarrow 'a::\text{euclidean\_space}$ 
  shows polynomial_function  $p \implies p$  C1_differentiable_on  $S$ 
  by (metis continuous_on_polynomial_function C1_differentiable_on_def has_vector_derivative_polynomial_function)

```

lemma *valid_path_polynomial_function*:

```

  fixes  $p :: \text{real} \Rightarrow 'a::\text{euclidean\_space}$ 
  shows polynomial_function  $p \implies \text{valid\_path } p$ 
  by (force simp: valid_path_def piecewise_C1_differentiable_on_def continuous_on_polynomial_function
  C1_differentiable_polynomial_function)

```

lemma *valid_path_subpath_trivial* [simp]:

```

  fixes  $g :: \text{real} \Rightarrow 'a::\text{euclidean\_space}$ 
  shows  $z \neq g\ x \implies \text{valid\_path } (\text{subpath } x\ x\ g)$ 
  by (simp add: subpath_def valid_path_polynomial_function)

```

1.13 Partial circle path

definition *part_circlepath* :: $[\text{complex}, \text{real}, \text{real}, \text{real}, \text{real}] \Rightarrow \text{complex}$

```

  where  $\text{part\_circlepath } z\ r\ s\ t \equiv \lambda x. z + \text{of\_real } r * \exp(i * \text{of\_real } (\text{linepath } s\ t\ x))$ 

```

lemma *pathstart_part_circlepath* [simp]:

```

  pathstart(part_circlepath  $z\ r\ s\ t$ ) =  $z + r * \exp(i * s)$ 
  by (metis part_circlepath_def pathstart_def pathstart_linepath)

```

lemma *pathfinish_part_circlepath* [simp]:

```

  pathfinish(part_circlepath  $z\ r\ s\ t$ ) =  $z + r * \exp(i * t)$ 
  by (metis part_circlepath_def pathfinish_def pathfinish_linepath)

```

lemma *reversepath_part_circlepath* [simp]:

```

  reversepath (part_circlepath  $z\ r\ s\ t$ ) = part_circlepath  $z\ r\ t\ s$ 
  unfolding part_circlepath_def reversepath_def linepath_def
  by (auto simp: algebra_simps)

```

lemma *has_vector_derivative_part_circlepath* [derivative_intros]:


```

    ((part_circlepath z r s t) has_vector_derivative
      (i * r * (of_real t - of_real s) * exp(i * linepath s t x)))
    (at x within X)
  unfolding part_circlepath_def linepath_def scaleR_conv_of_real
  by (rule has_vector_derivative_real_field derivative_eq_intros | simp)+

lemma differentiable_part_circlepath:
  part_circlepath c r a b differentiable at x within A
  using has_vector_derivative_part_circlepath[of c r a b x A] differentiableI_vector
  by blast

lemma vector_derivative_part_circlepath:
  vector_derivative (part_circlepath z r s t) (at x) =
    i * r * (of_real t - of_real s) * exp(i * linepath s t x)
  using has_vector_derivative_part_circlepath vector_derivative_at by blast

lemma vector_derivative_part_circlepath01:
  [[0 ≤ x; x ≤ 1]]
  ⇒ vector_derivative (part_circlepath z r s t) (at x within {0..1}) =
    i * r * (of_real t - of_real s) * exp(i * linepath s t x)
  using has_vector_derivative_part_circlepath
  by (auto simp: vector_derivative_at_within_ivl)

lemma valid_path_part_circlepath [simp]: valid_path (part_circlepath z r s t)
  unfolding valid_path_def
  by (auto simp: C1_differentiable_on_eq vector_derivative_works vector_derivative_part_circlepath
    has_vector_derivative_part_circlepath
    intro!: C1_differentiable_imp_piecewise continuous_intros)

lemma path_part_circlepath [simp]: path (part_circlepath z r s t)
  by (simp add: valid_path_imp_path)

proposition path_image_part_circlepath:
  assumes s ≤ t
  shows path_image (part_circlepath z r s t) = {z + r * exp(i * of_real x) | x.
s ≤ x ∧ x ≤ t}
proof -
  { fix z::real
    assume 0 ≤ z z ≤ 1
    with ‹s ≤ t› have ∃x. (exp (i * linepath s t z) = exp (i * of_real x)) ∧ s ≤ x
  ∧ x ≤ t
    apply (rule_tac x=(1 - z) * s + z * t in exI)
    apply (simp add: linepath_def scaleR_conv_of_real algebra_simps)
    by (metis (no_types) affine_ineq mult commute mult_left_mono)
  }
  moreover
  { fix z
    assume s ≤ z z ≤ t
    then have z + of_real r * exp (i * of_real z) ∈ (λx. z + of_real r * exp (i *

```

```

linepath s t x)) ' {0..1}
  apply (rule_tac x=(z - s)/(t - s) in image_eqI)
  apply (simp add: linepath_def scaleR_conv_of_real divide_simps exp_eq)
  apply (auto simp: field_split_simps)
  done
}
ultimately show ?thesis
  by (fastforce simp add: path_image_def part_circlepath_def)
qed

```

```

lemma path_image_part_circlepath':
  path_image (part_circlepath z r s t) = ( $\lambda x. z + r * cis x$ ) ' closed_segment s t
proof -
  have path_image (part_circlepath z r s t) =
    ( $\lambda x. z + r * exp(i * of_real x)$ ) ' linepath s t ' {0..1}
  by (simp add: image_image path_image_def part_circlepath_def)
  also have linepath s t ' {0..1} = closed_segment s t
  by (rule linepath_image_01)
  finally show ?thesis by (simp add: cis_conv_exp)
qed

```

```

lemma path_image_part_circlepath_subset:
  [ $s \leq t$ ;  $0 \leq r$ ]  $\implies$  path_image(part_circlepath z r s t)  $\subseteq$  sphere z r
by (auto simp: path_image_part_circlepath_sphere_def dist_norm algebra_simps
norm_mult)

```

```

lemma in_path_image_part_circlepath:
  assumes  $w \in$  path_image(part_circlepath z r s t)  $s \leq t$   $0 \leq r$ 
  shows  $norm(w - z) = r$ 
proof -
  have  $w \in \{c. dist z c = r\}$ 
  by (metis (no_types) path_image_part_circlepath_subset sphere_def subset_eq
assms)
  thus ?thesis
  by (simp add: dist_norm norm_minus_commute)
qed

```

```

lemma path_image_part_circlepath_subset':
  assumes  $r \geq 0$ 
  shows path_image (part_circlepath z r s t)  $\subseteq$  sphere z r
proof (cases  $s \leq t$ )
  case True
  thus ?thesis using path_image_part_circlepath_subset[of s t r z] assms by simp
next
  case False
  thus ?thesis using path_image_part_circlepath_subset[of t s r z] assms
  by (subst reversepath_part_circlepath [symmetric], subst path_image_reversepath)
simp_all
qed

```

lemma *part_circlepath_cnj*: $cnj (part_circlepath\ c\ r\ a\ b\ x) = part_circlepath\ (cnj\ c)\ r\ (-a)\ (-b)\ x$
by (*simp add: part_circlepath_def exp_cnj linepath_def algebra_simps*)

lemma *contour_integral_bound_part_circlepath*:

assumes $f\ contour_integrable_on\ part_circlepath\ c\ r\ a\ b$
assumes $B \geq 0\ r \geq 0 \wedge x. x \in path_image\ (part_circlepath\ c\ r\ a\ b) \implies norm\ (f\ x) \leq B$
shows $norm\ (contour_integral\ (part_circlepath\ c\ r\ a\ b)\ f) \leq B * r * |b - a|$
proof -
let $?I = integral\ \{0..1\}\ (\lambda x. f\ (part_circlepath\ c\ r\ a\ b\ x) * i * of_real\ (r * (b - a))) *$
 $exp\ (i * linepath\ a\ b\ x)$
have $norm\ ?I \leq integral\ \{0..1\}\ (\lambda x::real. B * 1 * (r * |b - a|) * 1)$
proof (*rule integral_norm_bound_integral, goal_cases*)
case 1
with *assms(1)* **show** ?case
by (*simp add: contour_integrable_on vector_derivative_part_circlepath mult_ac*)
next
case ($\exists x$)
with *assms(2-)* **show** ?case **unfolding** *norm_mult norm_of_real abs_mult*
by (*intro mult_mono*) (*auto simp: path_image_def*)
qed *auto*
also **have** $?I = contour_integral\ (part_circlepath\ c\ r\ a\ b)\ f$
by (*simp add: contour_integral_integral vector_derivative_part_circlepath mult_ac*)
finally **show** ?thesis **by** *simp*
qed

lemma *has_contour_integral_part_circlepath_iff*:

assumes $a < b$
shows $(f\ has_contour_integral\ I)\ (part_circlepath\ c\ r\ a\ b) \longleftrightarrow ((\lambda t. f\ (c + r * cis\ t) * r * i * cis\ t)\ has_integral\ I)\ \{a..b\}$
proof -
have $(f\ has_contour_integral\ I)\ (part_circlepath\ c\ r\ a\ b) \longleftrightarrow ((\lambda x. f\ (part_circlepath\ c\ r\ a\ b\ x) * vector_derivative\ (part_circlepath\ c\ r\ a\ b)\ (at\ x\ within\ \{0..1\}))\ has_integral\ I)\ \{0..1\}$
unfolding *has_contour_integral_def ..*
also **have** $\dots \longleftrightarrow ((\lambda x. f\ (part_circlepath\ c\ r\ a\ b\ x) * r * (b - a) * i * cis\ (linepath\ a\ b\ x))\ has_integral\ I)\ \{0..1\}$
by (*intro has_integral_cong, subst vector_derivative_part_circlepath01*) (*simp_all add: cis_conv_exp*)
also **have** $\dots \longleftrightarrow ((\lambda x. f\ (c + r * exp\ (i * linepath\ (of_real\ a)\ (of_real\ b)\ x)) * r * i * exp\ (i * linepath\ (of_real\ a)\ (of_real\ b)\ x) * vector_derivative\ (linepath\ (of_real\ a)\ (of_real\ b))\ (at\ x\ within\ \{0..1\}))\ has_integral\ I)\ \{0..1\}$

by (intro has_integral_cong, subst vector_derivative_linepath_within)
 (auto simp: part_circlepath_def cis_conv_exp of_real_linepath [symmetric])
 also have ... \longleftrightarrow (($\lambda z. f (c + r * \exp (i * z)) * r * i * \exp (i * z)$) has_contour_integral I)
 (linepath (of_real a) (of_real b))
 by (simp add: has_contour_integral_def)
 also have ... \longleftrightarrow (($\lambda t. f (c + r * \text{cis } t) * r * i * \text{cis } t$) has_integral I) {a..b}
 using assms
 by (subst has_contour_integral_linepath_Reals_iff) (simp_all add: cis_conv_exp)
 finally show ?thesis .
 qed

lemma contour_integrable_part_circlepath_iff:

assumes $a < b$
 shows f contour_integrable_on (part_circlepath c r a b) \longleftrightarrow
 ($\lambda t. f (c + r * \text{cis } t) * r * i * \text{cis } t$) integrable_on {a..b}
 using assms by (auto simp: contour_integrable_on_def integrable_on_def
 has_contour_integral_part_circlepath_iff)

lemma contour_integral_part_circlepath_eq:

assumes $a < b$
 shows contour_integral (part_circlepath c r a b) $f =$
 integral {a..b} ($\lambda t. f (c + r * \text{cis } t) * r * i * \text{cis } t$)
proof (cases f contour_integrable_on part_circlepath c r a b)
 case True
 hence ($\lambda t. f (c + r * \text{cis } t) * r * i * \text{cis } t$) integrable_on {a..b}
 using assms by (simp add: contour_integrable_part_circlepath_iff)
 with True show ?thesis
 using has_contour_integral_part_circlepath_iff[OF assms]
 contour_integral_unique has_integral_integrable_integral by blast
next
 case False
 hence \neg ($\lambda t. f (c + r * \text{cis } t) * r * i * \text{cis } t$) integrable_on {a..b}
 using assms by (simp add: contour_integrable_part_circlepath_iff)
 with False show ?thesis
 by (simp add: not_integrable_contour_integral not_integrable_integral)
 qed

lemma contour_integral_part_circlepath_reverse:

contour_integral (part_circlepath c r a b) $f = -$ contour_integral (part_circlepath
 c r b a) f
 by (metis contour_integral_reversepath reversepath_part_circlepath valid_path_part_circlepath)

lemma contour_integral_part_circlepath_reverse':

$b < a \implies$ contour_integral (part_circlepath c r a b) $f =$
 $-$ contour_integral (part_circlepath c r b a) f
 by (rule contour_integral_part_circlepath_reverse)

lemma finite_bounded_log: finite { $z::\text{complex}. \text{norm } z \leq b \wedge \exp z = w$ }

```

proof (cases w = 0)
  case True then show ?thesis by auto
next
  case False
  have *: finite {x. cmod ((2 * real_of_int x * pi) * i) ≤ b + cmod (Ln w)}
  proof (simp add: norm_mult finite_int_iff_bounded_le)
    show ∃k. abs ' {x. 2 * |of_int x| * pi ≤ b + cmod (Ln w)} ⊆ {..k}
    apply (rule_tac x=⌊(b + cmod (Ln w)) / (2*pi)⌋ in exI)
    apply (auto simp: field_split_simps le_floor_iff)
    done
  qed
  have [simp]: ⋀P f. {z. P z ∧ (∃n. z = f n)} = f ' {n. P (f n)}
    by blast
  have finite {z. cmod z ≤ b ∧ exp z = exp (Ln w)}
    using norm_add_leD by (fastforce intro: finite_subset [OF _ *] simp: exp_eq)
  then show ?thesis
    using False by auto
qed

lemma finite_bounded_log2:
  fixes a::complex
  assumes a ≠ 0
  shows finite {z. norm z ≤ b ∧ exp(a*z) = w}
proof -
  have *: finite ((λz. z / a) ' {z. cmod z ≤ b * cmod a ∧ exp z = w})
    by (rule finite_imageI [OF finite_bounded_log])
  show ?thesis
    by (rule finite_subset [OF _ *]) (force simp: assms norm_mult)
qed

lemma has_contour_integral_bound_part_circlepath_strong:
  assumes fi: (f has_contour_integral i) (part_circlepath z r s t)
  and finite k and le: 0 ≤ B 0 < r s ≤ t
  and B: ⋀x. x ∈ path_image(part_circlepath z r s t) - k ⇒ norm(f x) ≤ B
  shows cmod i ≤ B * r * (t - s)
proof -
  consider s = t | s < t using ⟨s ≤ t⟩ by linarith
  then show ?thesis
  proof cases
    case 1 with fi [unfolded has_contour_integral]
      have i = 0 by (simp add: vector_derivative_part_circlepath)
      with assms show ?thesis by simp
    next
      case 2
      have [simp]: |r| = r using ⟨r > 0⟩ by linarith
      have [simp]: cmod (complex_of_real t - complex_of_real s) = t - s
        by (metis 2 abs_of_pos diff_gt_0_iff_gt norm_of_real_of_real_diff)
      have finite (part_circlepath z r s t - ' {y} ∩ {0..1}) if y ∈ k for y
      proof -

```

```

    let ?w = (y - z)/of_real r / exp(i * of_real s)
    have fin: finite (of_real - '{z. cmod z ≤ 1 ∧ exp (i * complex_of_real (t -
s) * z) = ?w})
      using ‹s < t›
    by (intro finite_vimageI [OF finite_bounded_log2]) (auto simp: inj_of_real)
    show ?thesis
      unfolding part_circlepath_def linepath_def vimage_def
      using le
    by (intro finite_subset [OF _ fin]) (auto simp: algebra_simps scaleR_conv_of_real
exp_add exp_diff)
  qed
  then have fin01: finite ((part_circlepath z r s t) - 'k ∩ {0..1})
    by (rule finite_finite_vimage_IntI [OF ‹finite k›])
  have **: ((λx. if (part_circlepath z r s t x) ∈ k then 0
    else f(part_circlepath z r s t x) *
      vector_derivative (part_circlepath z r s t) (at x)) has_integral
i) {0..1}
    by (rule has_integral_spike [OF negligible_finite [OF fin01]]) (use fi has_contour_integral
in auto)
  have *: ‹λx. ‹0 ≤ x; x ≤ 1; part_circlepath z r s t x ∉ k› ⇒ cmod (f
(part_circlepath z r s t x)) ≤ B
    by (auto intro!: B [unfolded path_image_def image_def])
  show ?thesis
    apply (rule has_integral_bound [where 'a=real, simplified, OF _ **, sim-
plified])
    using assms le * 2 ‹r > 0› by (auto simp add: norm_mult vector_derivative_part_circlepath)
  qed
qed

corollary contour_integral_bound_part_circlepath_strong:
  assumes f_contour_integrable_on_part_circlepath z r s t
    and finite k and 0 ≤ B 0 < r s ≤ t
    and ‹λx. x ∈ path_image(part_circlepath z r s t) - k ⇒ norm(f x) ≤ B
  shows cmod (contour_integral (part_circlepath z r s t) f) ≤ B * r * (t - s)
  using assms has_contour_integral_bound_part_circlepath_strong has_contour_integral_integral
by blast

lemma has_contour_integral_bound_part_circlepath:
  ‹‹f has_contour_integral i› (part_circlepath z r s t);
  0 ≤ B; 0 < r; s ≤ t;
  ‹λx. x ∈ path_image(part_circlepath z r s t) ⇒ norm(f x) ≤ B›
  ⇒ norm i ≤ B*r*(t - s)
  by (auto intro: has_contour_integral_bound_part_circlepath_strong)

lemma contour_integrable_continuous_part_circlepath:
  continuous_on (path_image (part_circlepath z r s t)) f
  ⇒ f contour_integrable_on (part_circlepath z r s t)
  unfolding contour_integrable_on has_contour_integral_def vector_derivative_part_circlepath
path_image_def

```

by (best intro: integrable_continuous_real_path_part_circlepath [unfolded path_def]
continuous_intros

continuous_on_compose2 [where g=f, OF _ _ order_refl])

lemma simple_path_part_circlepath:

simple_path(part_circlepath z r s t) \longleftrightarrow ($r \neq 0 \wedge s \neq t \wedge |s - t| \leq 2 * \pi$)

proof (cases $r = 0 \vee s = t$)

case True

then show ?thesis

unfolding part_circlepath_def simple_path_def loop_free_def

by (rule disjE) (force intro: bexI [where $x = 1/4$] bexI [where $x = 1/3$])+

next

case False then have $r \neq 0 \wedge s \neq t$ by auto

have *: $\bigwedge x y z s t. i * ((1 - x) * s + x * t) = i * (((1 - y) * s + y * t)) + z \longleftrightarrow$
 $i * (x - y) * (t - s) = z$

by (simp add: algebra_simps)

have abs01: $\bigwedge x y :: \text{real}. 0 \leq x \wedge x \leq 1 \wedge 0 \leq y \wedge y \leq 1$

$\implies (x = y \vee x = 0 \wedge y = 1 \vee x = 1 \wedge y = 0 \longleftrightarrow |x - y| \in$

$\{0, 1\}$)

by auto

have **: $\bigwedge x y. (\exists n. (\text{complex_of_real } x - \text{of_real } y) * (\text{of_real } t - \text{of_real } s)$
 $= 2 * (\text{of_int } n * \text{of_real } \pi)) \longleftrightarrow$

$(\exists n. |x - y| * (t - s) = 2 * (\text{of_int } n * \pi))$

by (force simp: algebra_simps abs_if dest: arg_cong [where $f = \text{Re}$] arg_cong
[where $f = \text{complex_of_real}$]

intro: exI [where $x = -n$ for n])

have 1: $|s - t| \leq 2 * \pi$

if $\bigwedge x. 0 \leq x \wedge x \leq 1 \implies (\exists n. x * (t - s) = 2 * (\text{real_of_int } n * \pi)) \longrightarrow x$
 $= 0 \vee x = 1$

proof (rule ccontr)

assume $\neg |s - t| \leq 2 * \pi$

then have *: $\bigwedge n. t - s \neq \text{of_int } n * |s - t|$

using False that [of $2 * \pi / |t - s|$]

by (simp add: abs_minus_commute divide_simps)

show False

using * [of 1] * [of -1] by auto

qed

have 2: $|s - t| = |2 * (\text{real_of_int } n * \pi) / x|$ if $x \neq 0 \wedge x * (t - s) = 2 * (\text{real_of_int } n * \pi)$ for $x n$

proof -

have $t - s = 2 * (\text{real_of_int } n * \pi) / x$

using that by (simp add: field_simps)

then show ?thesis by (metis abs_minus_commute)

qed

have abs_away: $\bigwedge P. (\forall x \in \{0..1\}. \forall y \in \{0..1\}. P |x - y|) \longleftrightarrow (\forall x :: \text{real}. 0 \leq x$
 $\wedge x \leq 1 \longrightarrow P x)$

by force

show ?thesis using False

apply (simp add: simple_path_def loop_free_def)

```

apply (simp add: part_circlepath_def linepath_def exp_eq * ** abs01 del:
Set.insert_iff)
apply (subst abs_away)
apply (auto simp: 1)
apply (rule ccontr)
apply (auto simp: 2 field_split_simps abs_mult dest: of_int_leD)
done
qed

```

lemma arc_part_circlepath:

```

assumes  $r \neq 0$   $s \neq t$   $|s - t| < 2 * \pi$ 
shows arc (part_circlepath z r s t)
proof -
have *:  $x = y$  if eq:  $i * (\text{linepath } s \ t \ x) = i * (\text{linepath } s \ t \ y) + 2 * \text{of\_int } n * \text{complex\_of\_real } \pi * i$ 
and  $x \in \{0..1\}$  and  $y \in \{0..1\}$  for  $x \ y \ n$ 
proof (rule ccontr)
assume  $x \neq y$ 
have  $(\text{linepath } s \ t \ x) = (\text{linepath } s \ t \ y) + 2 * \text{of\_int } n * \text{complex\_of\_real } \pi$ 
by (metis add_divide_eq_iff complex_i_not_zero mult.commute nonzero_mult_div_cancel_left
eq)
then have  $s * y + t * x = s * x + (t * y + \text{of\_int } n * (\pi * 2))$ 
by (force simp: algebra_simps linepath_def dest: arg_cong [where f=Re])
with  $\langle x \neq y \rangle$  have st:  $s - t = (\text{of\_int } n * (\pi * 2)) / (y - x)$ 
by (force simp: field_simps)
have  $|\text{real\_of\_int } n| < |y - x|$ 
using assms  $\langle x \neq y \rangle$  by (simp add: st abs_mult field_simps)
then show False
using assms  $x \ y \ st$  by (auto dest: of_int_lessD)
qed
then have inj_on (part_circlepath z r s t)  $\{0..1\}$ 
using assms by (force simp add: part_circlepath_def inj_on_def exp_eq)
then show ?thesis
by (simp add: arc_def)
qed

```

1.14 Special case of one complete circle

definition circlepath :: $[\text{complex}, \text{real}, \text{real}] \Rightarrow \text{complex}$
where circlepath z r \equiv part_circlepath z r 0 (2 * π)

lemma circlepath: circlepath z r = $(\lambda x. z + r * \exp(2 * \text{of_real } \pi * i * \text{of_real } x))$
by (simp add: circlepath_def part_circlepath_def linepath_def algebra_simps)

lemma pathstart_circlepath [simp]: pathstart (circlepath z r) = $z + r$
by (simp add: circlepath_def)

lemma pathfinish_circlepath [simp]: pathfinish (circlepath z r) = $z + r$

by (simp add: circlepath_def) (metis exp_two_pi_i mult.commute)

lemma circlepath_minus: $\text{circlepath } z (-r) x = \text{circlepath } z r (x + 1/2)$

proof -

have $z + \text{of_real } r * \exp(2 * \text{pi} * i * (x + 1/2)) =$
 $z + \text{of_real } r * \exp(2 * \text{pi} * i * x + \text{pi} * i)$

by (simp add: divide_simps) (simp add: algebra_simps)

also have $\dots = z - r * \exp(2 * \text{pi} * i * x)$

by (simp add: exp_add)

finally show ?thesis

by (simp add: circlepath_path_image_def sphere_def dist_norm)

qed

lemma circlepath_add1: $\text{circlepath } z r (x+1) = \text{circlepath } z r x$

using circlepath_minus [of z r x+1/2] circlepath_minus [of z -r x]

by (simp add: add.commute)

lemma circlepath_add_half: $\text{circlepath } z r (x + 1/2) = \text{circlepath } z r (x - 1/2)$

using circlepath_add1 [of z r x-1/2]

by (simp add: add.commute)

lemma path_image_circlepath_minus_subset:

$\text{path_image } (\text{circlepath } z (-r)) \subseteq \text{path_image } (\text{circlepath } z r)$

proof -

have $\exists x \in \{0..1\}. \text{circlepath } z r (y + 1/2) = \text{circlepath } z r x$

if $0 \leq y \leq 1$ for y

proof (cases $y \leq 1/2$)

case False

with that show ?thesis

by (force simp: circlepath_add_half)

qed (use that in force)

then show ?thesis

by (auto simp add: path_image_def image_def circlepath_minus)

qed

lemma path_image_circlepath_minus: $\text{path_image } (\text{circlepath } z (-r)) = \text{path_image } (\text{circlepath } z r)$

using path_image_circlepath_minus_subset by fastforce

lemma has_vector_derivative_circlepath [derivative_intros]:

$((\text{circlepath } z r) \text{ has_vector_derivative } (2 * \text{pi} * i * r * \exp(2 * \text{of_real } \text{pi} * i * x)))$

(at x within X)

unfolding circlepath_def scaleR_conv_of_real

by (rule derivative_eq_intros) (simp add: algebra_simps)

lemma vector_derivative_circlepath:

$\text{vector_derivative } (\text{circlepath } z r) (\text{at } x) =$

$2 * \text{pi} * i * r * \exp(2 * \text{of_real } \text{pi} * i * x)$

using has_vector_derivative_circlepath vector_derivative_at by blast

lemma vector_derivative_circlepath01:

$\llbracket 0 \leq x; x \leq 1 \rrbracket$
 $\implies \text{vector_derivative } (\text{circlepath } z \ r) \ (\text{at } x \ \text{within } \{0..1\}) =$
 $2 * \pi * i * r * \exp(2 * \text{of_real } \pi * i * x)$

using has_vector_derivative_circlepath
 by (auto simp: vector_derivative_at_within_ivl)

lemma valid_path_circlepath [simp]: valid_path (circlepath z r)

by (simp add: circlepath_def)

lemma path_circlepath [simp]: path (circlepath z r)

by (simp add: valid_path_imp_path)

lemma path_image_circlepath_nonneg:

assumes $0 \leq r$ shows path_image (circlepath z r) = sphere z r

proof –

have *: $x \in (\lambda u. z + (\text{cmod } (x - z)) * \exp(i * (\text{of_real } u * (\text{of_real } \pi * 2))))$
 $\{0..1\}$ for x

proof (cases $x = z$)

case True then show ?thesis by force

next

case False

define w where $w = x - z$

then have $w \neq 0$ by (simp add: False)

have **: $\bigwedge t. \llbracket \text{Re } w = \cos t * \text{cmod } w; \text{Im } w = \sin t * \text{cmod } w \rrbracket \implies w = \text{of_real}$
 $(\text{cmod } w) * \exp(i * t)$

using cis_conv_exp_complex_eq_iff by auto

obtain t where $0 \leq t < 2 * \pi$ $\text{Re}(w / \text{norm } w) = \cos t$ $\text{Im}(w / \text{norm } w) = \sin t$

apply (rule sincos_total_2pi [of $\text{Re}(w / (\text{norm } w))$ $\text{Im}(w / (\text{norm } w))$])

by (auto simp add: divide_simps $\langle w \neq 0 \rangle$ cmod_power2 [symmetric])

then

show ?thesis

using False ** w_def $\langle w \neq 0 \rangle$

by (rule_tac $x=t / (2 * \pi)$ in image_eqI) (auto simp add: field_simps)

qed

show ?thesis

unfolding circlepath path_image_def sphere_def dist_norm

by (force simp: assms algebra_simps norm_mult norm_minus_commute intro:

*)

qed

lemma path_image_circlepath [simp]:

path_image (circlepath z r) = sphere z |r|

using path_image_circlepath_minus

by (force simp: path_image_circlepath_nonneg abs_if)

lemma has_contour_integral_bound_circlepath_strong:

```

  [(f has_contour_integral i) (circlepath z r);
   finite k; 0 ≤ B; 0 < r;
   ∧x. [norm(x - z) = r; x ∉ k] ⇒ norm(f x) ≤ B]
  ⇒ norm i ≤ B*(2*pi*r)
unfolding circlepath_def
by (auto simp: algebra_simps in_path_image_part_circlepath dest!: has_contour_integral_bound_part_circlepath)

```

```

lemma has_contour_integral_bound_circlepath:
  [(f has_contour_integral i) (circlepath z r);
   0 ≤ B; 0 < r; ∧x. norm(x - z) = r ⇒ norm(f x) ≤ B]
  ⇒ norm i ≤ B*(2*pi*r)
by (auto intro: has_contour_integral_bound_circlepath_strong)

```

```

lemma contour_integrable_continuous_circlepath:
  continuous_on (path_image (circlepath z r)) f
  ⇒ f contour_integrable_on (circlepath z r)
by (simp add: circlepath_def contour_integrable_continuous_part_circlepath)

```

```

lemma simple_path_circlepath: simple_path(circlepath z r) ↔ (r ≠ 0)
by (simp add: circlepath_def simple_path_part_circlepath)

```

```

lemma notin_path_image_circlepath [simp]: cmod (w - z) < r ⇒ w ∉ path_image
(circlepath z r)
by (simp add: sphere_def dist_norm norm_minus_commute)

```

```

lemma contour_integral_circlepath:
  assumes r > 0
  shows contour_integral (circlepath z r) (λw. 1 / (w - z)) = 2 * complex_of_real
pi * i
proof (rule contour_integral_unique)
  show ((λw. 1 / (w - z)) has_contour_integral 2 * complex_of_real pi * i)
(circlepath z r)
  unfolding has_contour_integral_def using assms has_integral_const_real [of
_ 0 1]
  apply (subst has_integral_cong)
  apply (simp add: vector_derivative_circlepath01)
  apply (force simp: circlepath)
  done
qed

```

1.15 Uniform convergence of path integral

Uniform convergence when the derivative of the path is bounded, and in particular for the special case of a circle.

```

proposition contour_integral_uniform_limit:
  assumes ev_fint: eventually (λn.: 'a. (f n) contour_integrable_on γ) F
  and ul_f: uniform_limit (path_image γ) f l F
  and noleB: ∧t. t ∈ {0..1} ⇒ norm (vector_derivative γ (at t)) ≤ B
  and γ: valid_path γ

```

```

    and [simp]:  $\neg$  trivial_limit F
    shows l contour_integrable_on  $\gamma$  (( $\lambda n$ . contour_integral  $\gamma$  (f n))  $\longrightarrow$  contour_integral  $\gamma$  l) F
  proof -
    have 0  $\leq$  B by (meson noleB [of 0] atLeastAtMost_iff norm_ge_zero order_refl order_trans zero_le_one)
    { fix e::real
      assume 0 < e
      then have 0 < e / (|B| + 1) by simp
      then have  $\forall_F n$  in F.  $\forall x \in \text{path\_image } \gamma$ . cmod (f n x - l x) < e / (|B| + 1)
        using ul_f [unfolded uniform_limit_iff dist_norm] by auto
      with ev_fint
      obtain a where fga:  $\bigwedge x$ .  $x \in \{0..1\} \implies$  cmod (f a ( $\gamma$  x) - l ( $\gamma$  x)) < e / (|B| + 1)
        and inta: ( $\lambda t$ . f a ( $\gamma$  t) * vector_derivative  $\gamma$  (at t)) integrable_on {0..1}
        using eventually_happens [OF eventually_conj]
        by (fastforce simp: contour_integrable_on_path_image_def)
      have Ble: B * e / (|B| + 1)  $\leq$  e
        using <0  $\leq$  B> <0 < e> by (simp add: field_split_simps)
      have  $\exists h$ . ( $\forall x \in \{0..1\}$ ). cmod (l ( $\gamma$  x) * vector_derivative  $\gamma$  (at x) - h x)  $\leq$  e)
         $\wedge$  h integrable_on {0..1}
      proof (intro exI conjI ballI)
        show cmod (l ( $\gamma$  x) * vector_derivative  $\gamma$  (at x) - f a ( $\gamma$  x) * vector_derivative  $\gamma$  (at x))  $\leq$  e
          if x  $\in$  {0..1} for x
          apply (rule order_trans [OF _ Ble])
          using noleB [OF that] fga [OF that] <0  $\leq$  B> <0 < e>
          apply (fastforce simp: mult_ac dest: mult_mono [OF less_imp_le] simp add: norm_mult left_diff_distrib [symmetric] norm_minus_commute divide_simps)
        done
      qed (rule inta)
    }
    then show lintg: l contour_integrable_on  $\gamma$ 
      unfolding contour_integrable_on by (metis (mono_tags, lifting) integrable_uniform_limit_real)
    { fix e::real
      define B' where B' = B + 1
      have B': B' > 0 B' > B using <0  $\leq$  B> by (auto simp: B'_def)
      assume 0 < e
      then have ev_no':  $\forall_F n$  in F.  $\forall x \in \text{path\_image } \gamma$ .  $2 * \text{cmod} (f n x - l x) < e / B'$ 
        using ul_f [unfolded uniform_limit_iff dist_norm, rule_format, of e / B'/2] B'
        by (simp add: field_simps)
      have ie: integral {0..1::real} ( $\lambda x$ . e/2) < e using <0 < e> by simp
      have *: cmod (f x ( $\gamma$  t) * vector_derivative  $\gamma$  (at t) - l ( $\gamma$  t) * vector_derivative  $\gamma$  (at t))  $\leq$  e/2
        if t: t  $\in$  {0..1} and leB':  $2 * \text{cmod} (f x (\gamma t) - l (\gamma t)) < e / B'$  for x t
      proof -

```

```

have 2 * cmod (f x (γ t) - l (γ t)) * cmod (vector_derivative γ (at t)) ≤ e
* (B/ B')
  using mult_mono [OF less_imp_le [OF leB] noleB] B' ‹0 < e› t by auto
also have ... < e
  by (simp add: B' ‹0 < e› mult_imp_div_pos_less)
finally have 2 * cmod (f x (γ t) - l (γ t)) * cmod (vector_derivative γ (at
t)) < e .
  then show ?thesis
  by (simp add: left_diff_distrib [symmetric] norm_mult)
qed
have le_e:  $\bigwedge x. [\forall u \in \{0..1\}. 2 * cmod (f x (\gamma u) - l (\gamma u)) < e / B'; f x$ 
contour_integrable_on  $\gamma]$ 
 $\implies cmod (integral \{0..1\}$ 
 $(\lambda u. f x (\gamma u) * vector\_derivative \gamma (at u) - l (\gamma u) * vector\_derivative$ 
 $\gamma (at u))) < e$ 
  apply (rule le_less_trans [OF integral_norm_bound_integral ie])
  apply (simp add: lintg_integrable_diff contour_integrable_on [symmetric])
  apply (blast intro: *)+
done
have  $\forall_F x$  in  $F. dist (contour\_integral \gamma (f x)) (contour\_integral \gamma l) < e$ 
  apply (rule eventually_mono [OF eventually_conj [OF ev_no' ev_fint]])
  apply (simp add: dist_norm contour_integrable_on_path_image_def con-
tour_integrable_integral)
  apply (simp add: lintg_integrable_diff [symmetric] contour_integrable_on
[symmetric] le_e)
done
}
then show  $((\lambda n. contour\_integral \gamma (f n)) \longrightarrow contour\_integral \gamma l) F$ 
by (rule tendstoI)
qed

```

corollary contour_integrable_uniform_limit_circlepath:

```

assumes  $\forall_F n::'a$  in  $F. (f n)$  contour_integrable_on (circlepath z r)
  and uniform_limit (sphere z r) f l F
  and  $\neg$  trivial_limit F 0 < r
shows l contour_integrable_on (circlepath z r)
   $((\lambda n. contour\_integral (circlepath z r) (f n)) \longrightarrow contour\_integral$ 
(circlepath z r) l) F
  using assms by (auto simp: vector_derivative_circlepath norm_mult intro!: con-
tour_integrable_uniform_limit)

```

end

2 Complex Path Integrals and Cauchy's Integral Theorem

By John Harrison et al. Ported from HOL Light by L C Paulson (2015)

theory Cauchy_Integral_Theorem

```

imports
  HOL-Analysis.Analysis
  Contour_Integration
begin

lemma leibniz_rule_holomorphic:
  fixes  $f::\text{complex} \Rightarrow 'b::\text{euclidean\_space} \Rightarrow \text{complex}$ 
  assumes  $\bigwedge x t. x \in U \implies t \in \text{cbox } a \ b \implies ((\lambda x. f \ x \ t) \text{ has\_field\_derivative } f \ x \ t)$  (at  $x$  within  $U$ )
  assumes  $\bigwedge x. x \in U \implies (f \ x) \text{ integrable\_on } \text{cbox } a \ b$ 
  assumes continuous_on  $(U \times (\text{cbox } a \ b)) (\lambda(x, t). f \ x \ t)$ 
  assumes convex  $U$ 
  shows  $(\lambda x. \text{integral } (\text{cbox } a \ b) (f \ x)) \text{ holomorphic\_on } U$ 
  using leibniz_rule_field_differentiable[OF assms(1-3) _ assms(4)]
  by (auto simp: holomorphic_on_def)

lemma Ln_measurable [measurable]:  $Ln \in \text{measurable borel borel}$ 
proof -
  have *:  $Ln \ (-\text{of\_real } x) = \text{of\_real } (\ln \ x) + i * \pi$  if  $x > 0$  for  $x$ 
    using that by (subst Ln_minus) (auto simp: Ln_of_real)
  have **:  $Ln \ (\text{of\_real } x) = \text{of\_real } (\ln \ (-x)) + i * \pi$  if  $x < 0$  for  $x$ 
    using *[of -x] that by simp
  have cont:  $(\lambda x. \text{indicat\_real } (-\ \mathbb{R}_{\leq 0}) \ x \ *_{\mathbb{R}} \ Ln \ x) \in \text{borel\_measurable borel}$ 
    by (intro borel_measurable_continuous_on_indicator continuous_intros) auto
  have  $(\lambda x. \text{if } x \in \mathbb{R}_{\leq 0} \ \text{then } \ln \ (-\text{Re } x) + i * \pi \ \text{else } \text{indicator } (-\mathbb{R}_{\leq 0}) \ x \ *_{\mathbb{R}} \ Ln \ x) \in \text{borel} \rightarrow_M \text{ borel}$ 
    (is  $?f \in \_$ ) by (rule measurable_If_set[OF _ cont]) auto
  hence  $(\lambda x. \text{if } x = 0 \ \text{then } Ln \ 0 \ \text{else } ?f \ x) \in \text{borel} \rightarrow_M \text{ borel}$  by measurable
  also have  $(\lambda x. \text{if } x = 0 \ \text{then } Ln \ 0 \ \text{else } ?f \ x) = Ln$ 
    by (auto simp: fun_eq_iff ** nonpos_Reals_def)
  finally show thesis .
qed

lemma powr_complex_measurable [measurable]:
  assumes [measurable]:  $f \in \text{measurable } M \ \text{borel } g \in \text{measurable } M \ \text{borel}$ 
  shows  $(\lambda x. f \ x \ \text{powr } g \ x :: \text{complex}) \in \text{measurable } M \ \text{borel}$ 
  using assms by (simp add: powr_def)

The special case of midpoints used in the main quadrisection

lemma has_contour_integral_midpoint:
  assumes  $(f \ \text{has\_contour\_integral } i) (\text{linepath } a \ (\text{midpoint } a \ b))$ 
     $(f \ \text{has\_contour\_integral } j) (\text{linepath } (\text{midpoint } a \ b) \ b)$ 
  shows  $(f \ \text{has\_contour\_integral } (i + j)) (\text{linepath } a \ b)$ 
proof (rule has_contour_integral_split)
  show  $\text{midpoint } a \ b - a = (1/2) *_{\mathbb{R}} (b - a)$ 
  using assms by (auto simp: midpoint_def scaleR_conv_of_real)
qed (use assms in auto)

lemma contour_integral_midpoint:

```

```

assumes continuous_on (closed_segment a b) f
shows contour_integral (linepath a b) f =
  contour_integral (linepath a (midpoint a b)) f + contour_integral (linepath
(midpoint a b) b) f
proof (rule contour_integral_split)
  show midpoint a b - a = (1/2) *R (b - a)
  using assms by (auto simp: midpoint_def scaleR_conv_of_real)
qed (use assms in auto)

```

A couple of special case lemmas that are useful below

lemma triangle_linear_has_chain_integral:

```

((λx. m*x + d) has_contour_integral 0) (linepath a b +++ linepath b c +++
linepath c a)

```

proof (rule Cauchy_theorem_primitive)

```

show ∧x. x ∈ UNIV ⇒ ((λx. m / 2 * x2 + d * x) has_field_derivative m *
x + d) (at x)

```

by (auto intro!: derivative_eq_intros)

qed auto

lemma has_chain_integral_chain_integral3:

```

assumes (f has_contour_integral i) (linepath a b +++ linepath b c +++ linepath
c d)

```

```

(is (f has_contour_integral i) ?g)

```

```

shows contour_integral (linepath a b) f + contour_integral (linepath b c) f +
contour_integral (linepath c d) f = i

```

```

(is ?lhs = _)

```

proof -

have f contour_integrable_on ?g

using assms contour_integrable_on_def **by** blast

then have ?lhs = contour_integral ?g f

by (simp add: valid_path_join has_contour_integral_integrable)

then show ?thesis

using assms contour_integral_unique **by** blast

qed

lemma has_chain_integral_chain_integral4:

```

assumes (f has_contour_integral i) (linepath a b +++ linepath b c +++ linepath
c d +++ linepath d e)

```

```

(is (f has_contour_integral i) ?g)

```

```

shows contour_integral (linepath a b) f + contour_integral (linepath b c) f +
contour_integral (linepath c d) f + contour_integral (linepath d e) f = i

```

```

(is ?lhs = _)

```

proof -

have f contour_integrable_on ?g

using assms contour_integrable_on_def **by** blast

then have ?lhs = contour_integral ?g f

by (simp add: valid_path_join has_contour_integral_integrable)

then show ?thesis

using assms contour_integral_unique **by** blast

qed

2.1 The key quadrisection step

lemma *norm_sum_half*:

assumes $\text{norm}(a + b) \geq e$

shows $\text{norm } a \geq e/2 \vee \text{norm } b \geq e/2$

proof –

have $e \leq \text{norm} (- a - b)$

by (*simp add: add commute assms norm_minus commute*)

thus *?thesis*

using *norm_triangle_ineq4 order_trans* **by** *fastforce*

qed

lemma *norm_sum_lemma*:

assumes $e \leq \text{norm} (a + b + c + d)$

shows $e / 4 \leq \text{norm } a \vee e / 4 \leq \text{norm } b \vee e / 4 \leq \text{norm } c \vee e / 4 \leq \text{norm } d$

proof –

have $e \leq \text{norm} ((a + b) + (c + d))$ **using** *assms*

by (*simp add: algebra_simps*)

then show *?thesis*

by (*auto dest!: norm_sum_half*)

qed

lemma *Cauchy_theorem_quadrisection*:

assumes *f: continuous_on (convex hull {a,b,c}) f*

and *dist: dist a b ≤ K dist b c ≤ K dist c a ≤ K*

and $e: e * K^2 \leq$

$\text{norm} (\text{contour_integral}(\text{linepath } a \ b) \ f + \text{contour_integral}(\text{linepath } b$
 $c) \ f + \text{contour_integral}(\text{linepath } c \ a) \ f)$

shows $\exists a' \ b' \ c'.$

$a' \in \text{convex hull } \{a,b,c\} \wedge b' \in \text{convex hull } \{a,b,c\} \wedge c' \in \text{convex hull}$
 $\{a,b,c\} \wedge$

$\text{dist } a' \ b' \leq K/2 \wedge \text{dist } b' \ c' \leq K/2 \wedge \text{dist } c' \ a' \leq K/2 \wedge$

$e * (K/2)^2 \leq \text{norm}(\text{contour_integral}(\text{linepath } a' \ b') \ f + \text{contour_integral}(\text{linepath}$
 $b' \ c') \ f + \text{contour_integral}(\text{linepath } c' \ a') \ f)$

(**is** $\exists x \ y \ z. \ ?\Phi \ x \ y \ z$)

proof –

note *divide_le_eq_numeral1 [simp del]*

define *a'* **where** $a' = \text{midpoint } b \ c$

define *b'* **where** $b' = \text{midpoint } c \ a$

define *c'* **where** $c' = \text{midpoint } a \ b$

have *fab*: *continuous_on (closed_segment a b) f* *continuous_on (closed_segment*
 $b \ c) \ f$ *continuous_on (closed_segment c a) f*

using *f continuous_on_subset segments_subset_convex_hull* **by** *metis+*

have *fcont'*: *continuous_on (closed_segment c' b') f*

continuous_on (closed_segment a' c') f

continuous_on (closed_segment b' a') f

unfolding *a'_def b'_def c'_def*


```

  by (rule continuous_on_subset [OF f],
      metis midpoints_in_convex_hull convex_hull_subset hull_subset in-
sert_subset segment_convex_hull)+
  define pathint where pathint x y  $\equiv$  contour_integral(linepath x y) f for x y
  have *: pathint a b + pathint b c + pathint c a =
    (pathint a c' + pathint c' b' + pathint b' a) +
    (pathint a' c' + pathint c' b + pathint b a') +
    (pathint a' c + pathint c b' + pathint b' a') +
    (pathint a' b' + pathint b' c' + pathint c' a')
  unfolding pathint_def
  by (simp add: fcont' contour_integral_reverse_linepath) (simp add: a'_def
b'_def c'_def contour_integral_midpoint fabc)
  have [simp]:  $\bigwedge x y. \text{cmod } (x * 2 - y * 2) = \text{cmod } (x - y) * 2$ 
  by (metis left_diff_distrib mult.commute norm_mult_numerals)
  have [simp]:  $\bigwedge x y. \text{cmod } (x - y) = \text{cmod } (y - x)$ 
  by (simp add: norm_minus_commute)
  consider  $e * K^2 / 4 \leq \text{cmod } (\text{pathint } a c' + \text{pathint } c' b' + \text{pathint } b' a) \mid$ 
     $e * K^2 / 4 \leq \text{cmod } (\text{pathint } a' c' + \text{pathint } c' b + \text{pathint } b a') \mid$ 
     $e * K^2 / 4 \leq \text{cmod } (\text{pathint } a' c + \text{pathint } c b' + \text{pathint } b' a') \mid$ 
     $e * K^2 / 4 \leq \text{cmod } (\text{pathint } a' b' + \text{pathint } b' c' + \text{pathint } c' a')$ 
  using assms by (metis * norm_sum_lemma pathint_def)
  then show ?thesis
  proof cases
    case 1 then have ? $\Phi$  a c' b'
      using assms unfolding pathint_def [symmetric]
      apply (clarsimp simp: c'_def b'_def midpoints_in_convex_hull hull_subset
[THEN subsetD])
      apply (auto simp: midpoint_def dist_norm scaleR_conv_of_real field_split_simps)
      done
    then show ?thesis by blast
  next
    case 2 then have ? $\Phi$  a' c' b
      using assms unfolding pathint_def [symmetric]
      apply (clarsimp simp: a'_def c'_def midpoints_in_convex_hull hull_subset
[THEN subsetD])
      apply (auto simp: midpoint_def dist_norm scaleR_conv_of_real field_split_simps)
      done
    then show ?thesis by blast
  next
    case 3 then have ? $\Phi$  a' c b'
      using assms unfolding pathint_def [symmetric]
      apply (clarsimp simp: a'_def b'_def midpoints_in_convex_hull hull_subset
[THEN subsetD])
      apply (auto simp: midpoint_def dist_norm scaleR_conv_of_real field_split_simps)
      done
    then show ?thesis by blast
  next
    case 4 then have ? $\Phi$  a' b' c'
      using assms unfolding pathint_def [symmetric]

```

```

apply (clarsimp simp: a'_def c'_def b'_def midpoints_in_convex_hull
hull_subset [THEN subsetD])
apply (auto simp: midpoint_def dist_norm scaleR_conv_of_real field_split_simps)
done
then show ?thesis by blast
qed
qed

```

2.2 Cauchy's theorem for triangles

lemma *triangle_points_closer*:

```

fixes a::complex
shows  $\llbracket x \in \text{convex hull } \{a,b,c\}; y \in \text{convex hull } \{a,b,c\} \rrbracket$ 
 $\implies \text{norm}(x - y) \leq \text{norm}(a - b) \vee$ 
 $\text{norm}(x - y) \leq \text{norm}(b - c) \vee$ 
 $\text{norm}(x - y) \leq \text{norm}(c - a)$ 
using simplex_extremal_le [of  $\{a,b,c\}$ ]
by (auto simp: norm_minus_commute)

```

lemma *holomorphic_point_small_triangle*:

```

assumes x:  $x \in S$ 
and f: continuous_on S f
and cd: f field_differentiable (at x within S)
and e:  $0 < e$ 
shows  $\exists k > 0. \forall a b c. \text{dist } a b \leq k \wedge \text{dist } b c \leq k \wedge \text{dist } c a \leq k \wedge$ 
 $x \in \text{convex hull } \{a,b,c\} \wedge \text{convex hull } \{a,b,c\} \subseteq S$ 
 $\longrightarrow \text{norm}(\text{contour\_integral}(\text{linepath } a b) f + \text{contour\_integral}(\text{linepath}$ 
 $b c) f +$ 
 $\text{contour\_integral}(\text{linepath } c a) f)$ 
 $\leq e * (\text{dist } a b + \text{dist } b c + \text{dist } c a)^2$ 
(is  $\exists k > 0. \forall a b c. \_ \longrightarrow ?\text{normle } a b c)$ 

```

proof –

```

have le_of_3:  $\bigwedge a x y z. \llbracket 0 \leq x*y; 0 \leq x*z; 0 \leq y*z; a \leq (e*(x + y + z))*x$ 
 $+ (e*(x + y + z))*y + (e*(x + y + z))*z \rrbracket$ 
 $\implies a \leq e*(x + y + z)^2$ 

```

by (simp add: algebra_simps power2_eq_square)

```

have disj_le:  $\llbracket x \leq a \vee x \leq b \vee x \leq c; 0 \leq a; 0 \leq b; 0 \leq c \rrbracket \implies x \leq a + b + c$ 
for x::real and a b c

```

by linarith

```

have fabc: f contour_integrable_on linepath a b f contour_integrable_on linepath
b c f contour_integrable_on linepath c a

```

if $\text{convex hull } \{a, b, c\} \subseteq S$ **for** a b c

using segments_subset_convex_hull that

by (metis continuous_on_subset f contour_integrable_continuous_linepath)+

note path_bound = has_contour_integral_bound_linepath [simplified norm_minus_commute,
OF has_contour_integral_integral]

{ **fix** f' a b c d

assume d: $0 < d$

```

    and f':  $\bigwedge y. \llbracket cmod (y - x) \leq d; y \in S \rrbracket \implies cmod (f y - f x - f' * (y - x))$ 
 $\leq e * cmod (y - x)$ 
    and le:  $cmod (a - b) \leq d \ cmod (b - c) \leq d \ cmod (c - a) \leq d$ 
    and xc:  $x \in convex \ hull \ \{a, b, c\}$ 
    and S:  $convex \ hull \ \{a, b, c\} \subseteq S$ 
    have pa:  $contour\_integral \ (linepath \ a \ b) \ f + contour\_integral \ (linepath \ b \ c) \ f$ 
+  $contour\_integral \ (linepath \ c \ a) \ f =$ 
       $contour\_integral \ (linepath \ a \ b) \ (\lambda y. f \ y - f \ x - f' * (y - x)) +$ 
       $contour\_integral \ (linepath \ b \ c) \ (\lambda y. f \ y - f \ x - f' * (y - x)) +$ 
       $contour\_integral \ (linepath \ c \ a) \ (\lambda y. f \ y - f \ x - f' * (y - x))$ 
    apply (simp add: contour_integral_diff contour_integral_lmul contour_integrable_lmul
contour_integrable_diff fabc [OF S])
    apply (simp add: field_simps)
    done
  { fix y
    assume yc:  $y \in convex \ hull \ \{a, b, c\}$ 
    have cmod (f y - f x - f' * (y - x))  $\leq e * norm (y - x)$ 
    proof (rule f')
      show  $cmod (y - x) \leq d$ 
      by (metis triangle_points_closer [OF xc yc] le norm_minus_commute
order_trans)
    qed (use S yc in blast)
    also have  $\dots \leq e * (cmod (a - b) + cmod (b - c) + cmod (c - a))$ 
      by (simp add: yc e xc disj_le [OF triangle_points_closer])
    finally have  $cmod (f y - f x - f' * (y - x)) \leq e * (cmod (a - b) + cmod$ 
(b - c) +  $cmod (c - a))$  .
  } note cm_le = this
  have ?normle a b c
  unfolding dist_norm pa
  using f' xc S e
  apply (intro le_of_3 norm_triangle_le add_mono path_bound)
  apply (simp_all add: contour_integral_diff contour_integral_lmul con-
tour_integrable_lmul contour_integrable_diff fabc)
  apply (blast intro: cm_le elim: dest: segments_subset_convex_hull [THEN
subsetD])+
  done
} note * = this
show ?thesis
using cd e
apply (simp add: field_differentiable_def has_field_derivative_def has_derivative_within_alt
approachable_lt_le2 Ball_def)
apply (clarify dest!: spec mp)
using * unfolding dist_norm
apply blast
done
qed

```

Hence the most basic theorem for a triangle.

locale Chain =

```

fixes  $x0$  At Follows
assumes  $At0$ :  $At\ x0\ 0$ 
and  $AtSuc$ :  $\bigwedge x\ n. At\ x\ n \implies \exists x'. At\ x'\ (Suc\ n) \wedge Follows\ x'\ x$ 
begin
primrec  $f$  where
   $f\ 0 = x0$ 
   $| f\ (Suc\ n) = (SOME\ x. At\ x\ (Suc\ n) \wedge Follows\ x\ (f\ n))$ 

lemma  $At$ :  $At\ (f\ n)\ n$ 
proof (induct n)
  case 0 show ?case
    by (simp add: At0)
next
  case ( $Suc\ n$ ) show ?case
    by (metis (no_types, lifting) AtSuc [OF Suc] f.simps(2) someI_ex)
qed

lemma  $Follows$ :  $Follows\ (f\ (Suc\ n))\ (f\ n)$ 
by (metis (no_types, lifting) AtSuc [OF At [of n]] f.simps(2) someI_ex)

declare  $f.simps(2)$  [simp del]
end

lemma  $Chain3$ :
assumes  $At0$ :  $At\ x0\ y0\ z0\ 0$ 
and  $AtSuc$ :  $\bigwedge x\ y\ z\ n. At\ x\ y\ z\ n \implies \exists x'\ y'\ z'. At\ x'\ y'\ z'\ (Suc\ n) \wedge Follows\ x'\ y'\ z'\ x\ y\ z$ 
obtains  $f\ g\ h$  where
   $f\ 0 = x0\ g\ 0 = y0\ h\ 0 = z0$ 
   $\bigwedge n. At\ (f\ n)\ (g\ n)\ (h\ n)\ n$ 
   $\bigwedge n. Follows\ (f\ (Suc\ n))\ (g\ (Suc\ n))\ (h\ (Suc\ n))\ (f\ n)\ (g\ n)\ (h\ n)$ 
proof –
interpret  $three$ :  $Chain\ (x0,y0,z0)\ \lambda(x,y,z). At\ x\ y\ z\ \lambda(x',y',z'). \lambda(x,y,z). Follows\ x'\ y'\ z'\ x\ y\ z$ 
proof qed (use At0 AtSuc in auto)
show ?thesis
proof
  show  $\bigwedge n. Follows\ (fst\ (three.f\ (Suc\ n)))\ (fst\ (snd\ (three.f\ (Suc\ n))))$ 
    ( $snd\ (snd\ (three.f\ (Suc\ n))))\ (fst\ (three.f\ n))$ )
    ( $fst\ (snd\ (three.f\ n))\ (snd\ (snd\ (three.f\ n)))$ )
   $\bigwedge n. At\ (fst\ (three.f\ n))\ (fst\ (snd\ (three.f\ n))\ (snd\ (snd\ (three.f\ n))))\ n$ 
using  $three.At\ three.Follows$ 
by (simp_all add: split_beta')
qed auto
qed

proposition  $Cauchy\_theorem\_triangle$ :
assumes  $f$  holomorphic_on (convex hull  $\{a,b,c\}$ )

```

```

  shows (f has_contour_integral 0) (linepath a b +++ linepath b c +++ linepath
c a)
proof -
  have contf: continuous_on (convex hull {a,b,c}) f
  by (metis assms holomorphic_on_imp_continuous_on)
  let ?pathint =  $\lambda x y. \text{contour\_integral}(\text{linepath } x y) f$ 
  { fix y::complex
    assume fy: (f has_contour_integral y) (linepath a b +++ linepath b c +++
linepath c a)
    and ynz:  $y \neq 0$ 
    define K where  $K = 1 + \max(\text{dist } a b) (\max(\text{dist } b c) (\text{dist } c a))$ 
    define e where  $e = \text{norm } y / K^2$ 
    have K1:  $K \geq 1$  by (simp add: K_def max.coboundedI1)
    then have K:  $K > 0$  by linarith
    have [iff]:  $\text{dist } a b \leq K \text{ dist } b c \leq K \text{ dist } c a \leq K$ 
    by (simp_all add: K_def)
    have e:  $e > 0$ 
    unfolding e_def using ynz K1 by simp
    define At where  $At \ x \ y \ z \ n \longleftrightarrow$ 
      convex hull {x,y,z}  $\subseteq$  convex hull {a,b,c}  $\wedge$ 
       $\text{dist } x y \leq K/2^n \wedge \text{dist } y z \leq K/2^n \wedge \text{dist } z x \leq K/2^n \wedge$ 
       $\text{norm}(\text{?pathint } x y + \text{?pathint } y z + \text{?pathint } z x) \geq e * (K/2^n)^2$ 
    for x y z n
    have At0:  $At \ a \ b \ c \ 0$ 
    using fy
    by (simp add: At_def e_def has_chain_integral_chain_integral3)
    { fix x y z n
      assume At:  $At \ x \ y \ z \ n$ 
      then have contf': continuous_on (convex hull {x,y,z}) f
      using contf At_def continuous_on_subset by metis
      have  $\exists x' y' z'. At \ x' y' z' (Suc \ n) \wedge \text{convex hull } \{x',y',z'\} \subseteq \text{convex hull}$ 
 $\{x,y,z\}$ 
      using At Cauchy_theorem_quadrisection [OF contf', of  $K/2^n$  e]
      apply (simp add: At_def algebra_simps)
      apply (meson convex_hull_subset empty_subsetI insert_subset subsetCE)
      done
    }
    note AtSuc = this
    obtain fa fb fc
    where f0 [simp]:  $fa \ 0 = a \ fb \ 0 = b \ fc \ 0 = c$ 
    and cosb:  $\bigwedge n. \text{convex hull } \{fa \ n, fb \ n, fc \ n\} \subseteq \text{convex hull } \{a,b,c\}$ 
    and dist:  $\bigwedge n. \text{dist } (fa \ n) (fb \ n) \leq K/2^n$ 
       $\bigwedge n. \text{dist } (fb \ n) (fc \ n) \leq K/2^n$ 
       $\bigwedge n. \text{dist } (fc \ n) (fa \ n) \leq K/2^n$ 
    and no:  $\bigwedge n. \text{norm}(\text{?pathint } (fa \ n) (fb \ n) +$ 
       $\text{?pathint } (fb \ n) (fc \ n) +$ 
       $\text{?pathint } (fc \ n) (fa \ n)) \geq e * (K/2^n)^2$ 
    and conv_le:  $\bigwedge n. \text{convex hull } \{fa(Suc \ n), fb(Suc \ n), fc(Suc \ n)\} \subseteq \text{convex}$ 
 $\text{hull } \{fa \ n, fb \ n, fc \ n\}$ 
    by (rule Chain3 [of At, OF At0 AtSuc]) (auto simp: At_def)
  }

```

```

obtain  $x$  where  $x: \bigwedge n. x \in \text{convex hull } \{fa\ n, fb\ n, fc\ n\}$ 
proof (rule bounded_closed_nest)
  show  $\bigwedge n. \text{closed } (\text{convex hull } \{fa\ n, fb\ n, fc\ n\})$ 
    by (simp add: compact_imp_closed finite_imp_compact convex_hull)
  show  $\bigwedge m\ n. m \leq n \implies \text{convex hull } \{fa\ n, fb\ n, fc\ n\} \subseteq \text{convex hull } \{fa\ m,$ 
fb  $m, fc\ m\}$ 
    by (erule transitive_stepwise_le) (auto simp: conv_le)
qed (fastforce intro: finite_imp_bounded_convex_hull)+
then have  $xin: x \in \text{convex hull } \{a,b,c\}$ 
  using assms  $f0$  by blast
then have  $fx: f \text{ field\_differentiable at } x \text{ within } (\text{convex hull } \{a,b,c\})$ 
  using assms holomorphic_on_def by blast
{ fix  $k\ n$ 
  assume  $k: 0 < k$ 
  and  $le:$ 
     $\bigwedge x'\ y'\ z'. \llbracket \text{dist } x'\ y' \leq k; \text{dist } y'\ z' \leq k; \text{dist } z'\ x' \leq k;$ 
 $x \in \text{convex hull } \{x',y',z'\};$ 
 $\text{convex hull } \{x',y',z'\} \subseteq \text{convex hull } \{a,b,c\} \rrbracket$ 
     $\implies$ 
 $\text{cmod } (?pathint\ x'\ y' + ?pathint\ y'\ z' + ?pathint\ z'\ x') * 10$ 
 $\leq e * (\text{dist } x'\ y' + \text{dist } y'\ z' + \text{dist } z'\ x')^2$ 
  and  $Kk: K / k < 2 \wedge n$ 
have  $K / 2 \wedge n < k$  using  $Kk\ k$ 
  by (auto simp: field_simps)
then have  $DD: \text{dist } (fa\ n)\ (fb\ n) \leq k \text{ dist } (fb\ n)\ (fc\ n) \leq k \text{ dist } (fc\ n)\ (fa\ n)$ 
 $\leq k$ 
    using dist [of  $n$ ]  $k$ 
    by linarith+
have  $dle: (\text{dist } (fa\ n)\ (fb\ n) + \text{dist } (fb\ n)\ (fc\ n) + \text{dist } (fc\ n)\ (fa\ n))^2$ 
 $\leq (3 * K / 2 \wedge n)^2$ 
    using dist [of  $n$ ]  $e\ K$ 
    by (simp add: abs_le_square_iff [symmetric])
have less10:  $\bigwedge x\ y::\text{real}. 0 < x \implies y \leq 9*x \implies y < x*10$ 
    by linarith
have  $e * (\text{dist } (fa\ n)\ (fb\ n) + \text{dist } (fb\ n)\ (fc\ n) + \text{dist } (fc\ n)\ (fa\ n))^2 \leq e * (3 * K / 2 \wedge n)^2$ 
    using ynz dle e mult_le_cancel_left_pos by blast
also have  $\dots <$ 
 $\text{cmod } (?pathint\ (fa\ n)\ (fb\ n) + ?pathint\ (fb\ n)\ (fc\ n) + ?pathint\ (fc\ n)\ (fa\ n)) * 10$ 
    using no [of  $n$ ]  $e\ K$ 
    by (simp add: e_def field_simps) (simp only: zero_less_norm_iff [symmetric])
finally have False
    using le [OF  $DD\ x\ cosb$ ] by auto
} then
have ?thesis
  using holomorphic_point_small_triangle [OF  $xin\ contf\ fx, of\ e/10$ ]  $e$ 
  apply clarsimp

```

```

    apply (rule_tac y1=K/k in exE [OF real_arch_pow[of 2]], force+)
  done
}
moreover have f contour_integrable_on (linepath a b +++ linepath b c +++
linepath c a)
  by simp (meson contf continuous_on_subset contour_integrable_continuous_linepath
segments_subset_convex_hull(1)
           segments_subset_convex_hull(3) segments_subset_convex_hull(5))
ultimately show ?thesis
  using has_contour_integral_integral by fastforce
qed

```

2.3 Version needing function holomorphic in interior only

lemma *Cauchy_theorem_flat_lemma:*

```

  assumes f: continuous_on (convex_hull {a,b,c}) f
    and c: c - a = k *R (b - a)
    and k: 0 ≤ k
  shows contour_integral (linepath a b) f + contour_integral (linepath b c) f +
        contour_integral (linepath c a) f = 0
proof -
  have fabc: continuous_on (closed_segment a b) f continuous_on (closed_segment
b c) f continuous_on (closed_segment c a) f
    using f continuous_on_subset segments_subset_convex_hull by metis+
  show ?thesis
  proof (cases k ≤ 1)
    case True show ?thesis
    by (simp add: contour_integral_split [OF fabc(1) k True c] contour_integral_reverse_linepath
fabc)
  next
    case False
    show ?thesis
    proof (subst contour_integral_split [symmetric])
      show b - a = (1/k) *R (c - a)
        using False c by force
      show contour_integral (linepath a c) f + contour_integral (linepath c a) f =
0
        by (simp add: contour_integral_reverse_linepath fabc(3))
      show continuous_on (closed_segment a c) f
        by (metis closed_segment_commute fabc(3))
    qed (use False in auto)
  qed
qed

```

lemma *Cauchy_theorem_flat:*

```

  assumes f: continuous_on (convex_hull {a,b,c}) f
    and c: c - a = k *R (b - a)
  shows contour_integral (linepath a b) f +
        contour_integral (linepath b c) f +

```

```

      contour_integral (linepath c a) f = 0
proof (cases 0 ≤ k)
  case True with assms show ?thesis
    by (blast intro: Cauchy_theorem_flat_lemma)
next
  case False
    have continuous_on (closed_segment a b) f continuous_on (closed_segment b
c) f continuous_on (closed_segment c a) f
      using f continuous_on_subset segments_subset_convex_hull by metis+
    moreover have contour_integral (linepath b a) f + contour_integral (linepath
a c) f +
      contour_integral (linepath c b) f = 0
    proof (rule Cauchy_theorem_flat_lemma [of b a c f 1-k])
      show continuous_on (convex_hull {b, a, c}) f
        by (simp add: f_insert_commute)
      show c - b = (1 - k) *R (a - b)
        using c by (auto simp: algebra_simps)
    qed (use False in auto)
    ultimately show ?thesis
      by (metis (no_types, lifting) contour_integral_reverse_linepath_eq_neg_iff_add_eq_0
minus_add_cancel)
qed

```

proposition *Cauchy_theorem_triangle_interior:*

```

  assumes contf: continuous_on (convex_hull {a,b,c}) f
    and holf: f holomorphic_on_interior (convex_hull {a,b,c})
    shows (f has_contour_integral 0) (linepath a b +++ linepath b c +++ linepath
c a)
proof -
  define pathint where pathint ≡ λx y. contour_integral(linepath x y) f
  have fabc: continuous_on (closed_segment a b) f continuous_on (closed_segment
b c) f continuous_on (closed_segment c a) f
    using contf continuous_on_subset segments_subset_convex_hull by metis+
  have bounded (f ‘ (convex_hull {a,b,c}))
    by (simp add: compact_continuous_image compact_convex_hull compact_imp_bounded
contf)
  then obtain B where 0 < B and Bnf: ∧x. x ∈ convex_hull {a,b,c} ⇒ norm
(f x) ≤ B
    by (auto simp: dest!: bounded_pos [THEN iffD1])
  have bounded (convex_hull {a,b,c})
    by (simp add: bounded_convex_hull)
  then obtain C where C: 0 < C and Cno: ∧y. y ∈ convex_hull {a,b,c} ⇒
norm y < C
    using bounded_pos_less by blast
  then have diff_2C: norm(x - y) ≤ 2*C
    if x: x ∈ convex_hull {a, b, c} and y: y ∈ convex_hull {a, b, c} for x y
proof -
  have cmod x ≤ C

```



```

    using x by (meson Cno not_le not_less_iff_gr_or_eq)
  hence cmod (x - y) ≤ C + C
  using y by (meson Cno add_mono_thms_linordered_field(4) less_eq_real_def
norm_triangle_ineq4 order_trans)
  thus cmod (x - y) ≤ 2 * C
  by (metis mult_2)
qed
have contf': continuous_on (convex hull {b,a,c}) f
  using contf by (simp add: insert_commute)
{ fix y::complex
  assume fy: (f has_contour_integral y) (linepath a b +++ linepath b c +++
linepath c a)
  and ynz: y ≠ 0
  have pi_eq_y: pathint a b + pathint b c + pathint c a = y
  unfolding pathint_def by (rule has_chain_integral_chain_integral3 [OF
fy])
  have ?thesis
  proof (cases c=a ∨ a=b ∨ b=c)
  case True then show ?thesis
    using Cauchy_theorem_flat [OF contf', of 0]
    using has_chain_integral_chain_integral3 [OF fy] ynz
    by (force simp: fabc contour_integral_reverse_linepath)
  next
  case False
  then have car3: card {a, b, c} = Suc (DIM(complex))
  by auto
  { assume interior(convex hull {a,b,c}) = {}
  then have collinear{a,b,c}
    using interior_convex_hull_eq_empty [OF car3]
    by (simp add: collinear_3_eq_affine_dependent)
  with False obtain d where c ≠ a a ≠ b b ≠ c c - b = d *R (a - b)
  by (auto simp: collinear_3 collinear_lemma)
  then have False
    using False Cauchy_theorem_flat [OF contf'] pi_eq_y ynz
    by (simp add: fabc add_eq_0_iff contour_integral_reverse_linepath
pathint_def)
  }
  then obtain d where d: d ∈ interior (convex hull {a, b, c})
  by blast
  { fix d1
  assume d1_pos: 0 < d1
  and d1: ∧x x'. [|x∈convex hull {a, b, c}; x'∈convex hull {a, b, c}; cmod
(x' - x) < d1|]
    ⇒ cmod (f x' - f x) < cmod y / (24 * C)
  define e where e = min 1 (min (d1/(4*C)) ((norm y / 24 / C) / B))
  define shrink where shrink x = x - e *R (x - d) for x
  have e: 0 < e ≤ 1 e ≤ d1 / (4 * C) e ≤ cmod y / 24 / C / B
  using d1_pos ‹C>0 ‹B>0 ynz by (simp_all add: e_def)
  have e_le_d1: e * (4 * C) ≤ d1

```

```

    using e <C>0 by (simp add: field_simps)
  have shrink a ∈ interior(convex hull {a,b,c})
    shrink b ∈ interior(convex hull {a,b,c})
    shrink c ∈ interior(convex hull {a,b,c})
  using d e by (auto simp: hull_inc mem_interior_convex_shrink shrink_def)
  then have fhp0: (f has_contour_integral 0)
    (linepath (shrink a) (shrink b) +++ linepath (shrink b) (shrink c)
  +++ linepath (shrink c) (shrink a))
    by (simp add: Cauchy_theorem_triangle holomorphic_on_subset [OF holf]
  hull_minimal)
  then have f_0_shrink: pathint (shrink a) (shrink b) + pathint (shrink b)
  (shrink c) + pathint (shrink c) (shrink a) = 0
    by (simp add: has_chain_integral_chain_integral3 pathint_def)
  have fpi_abc: f contour_integrable_on linepath (shrink a) (shrink b)
    f contour_integrable_on linepath (shrink b) (shrink c)
    f contour_integrable_on linepath (shrink c) (shrink a)
  using fhp0 by (auto simp: valid_path_join dest: has_contour_integral_integrable)
  have cmod_shr:  $\bigwedge x y. \text{cmod} (\text{shrink } y - \text{shrink } x - (y - x)) = e * \text{cmod}
  (x - y)$ 
    using e by (simp add: shrink_def real_vector.scale_right_diff_distrib
  [symmetric])
  have sh_eq:  $\bigwedge a b d::\text{complex}. (b - e *_R (b - d)) - (a - e *_R (a - d)) -
  (b - a) = e *_R (a - b)$ 
    by (simp add: algebra_simps)
  have cmod_y / (24 * C) ≤ cmod y / cmod (b - a) / 12
    using False <C>0 diff_2C [of b a] ynz
    by (auto simp: field_split_simps hull_inc)
  have less_C:  $x * \text{cmod } u < C$  if  $u \in \text{convex hull } \{a,b,c\}$   $0 \leq x \leq 1$  for
  x u
  proof (cases x=0)
  case False
  with that show ?thesis
    using Cno [of u] mult_left_le_one_le [of cmod u x] le_less_trans
  norm_ge_zero by blast
  qed (simp add: <0<C>)
  { fix u v
    assume uv:  $u \in \text{convex hull } \{a, b, c\}$   $v \in \text{convex hull } \{a, b, c\}$   $u \neq v$ 
    and fpi_uv: f contour_integrable_on linepath (shrink u) (shrink v)
    have shr_uv: shrink u ∈ interior(convex hull {a,b,c})
      shrink v ∈ interior(convex hull {a,b,c})
    using d e uv
    by (auto simp: hull_inc mem_interior_convex_shrink shrink_def)
    have cmod_fw:  $\bigwedge x. 0 \leq x \implies x \leq 1 \implies \text{cmod} (f (\text{linepath} (\text{shrink } u)
  (\text{shrink } v) x)) \leq B$ 
    using shr_uv by (blast intro: Bnflinepath_in_convex_hull_interior_subset
  [THEN subsetD])
    { fix x::real assume x:  $0 \leq x \leq 1$ 
      have |1 - x| * cmod u < C |x| * cmod v < C
        using uv x by (auto intro!: less_C)
    }
  }

```

```

moreover have  $|x| * cmod\ d < C$   $|1 - x| * cmod\ d < C$ 
using  $x\ d\ interior\_subset$  by (auto intro!: less_C)
ultimately
have  $cmod\_less\_4C$ :  $cmod\ ((1 - x) *_R\ u - (1 - x) *_R\ d) + cmod\ (x$ 
 $*_R\ v - x *_R\ d) < (C+C) + (C+C)$ 
by (metis add_strict_mono le_less_trans norm_scaleR norm_triangle_ineq4)
have  $ll$ :  $linepath\ (shrink\ u)\ (shrink\ v)\ x - linepath\ u\ v\ x = -e * ((1 -$ 
 $x) *_R\ (u - d) + x *_R\ (v - d))$ 
by (simp add: linepath_def shrink_def algebra_simps scaleR_conv_of_real)
have  $cmod\_less\_dt$ :  $cmod\ (linepath\ (shrink\ u)\ (shrink\ v)\ x - linepath\ u$ 
 $v\ x) < d1$ 
unfolding  $ll$   $norm\_mult\ scaleR\_diff\_right$ 
using  $\langle e > 0 \rangle$   $cmod\_less\_4C$  by (force intro: norm_triangle_lt
less_le_trans [OF  $_e\ le\ d1$ ])
have  $cmod\ (f\ (linepath\ (shrink\ u)\ (shrink\ v)\ x)) * cmod\ (shrink\ v -$ 
 $shrink\ u - (v - u)) +$ 
 $cmod\ (v - u) * cmod\ (f\ (linepath\ (shrink\ u)\ (shrink\ v)\ x) -$ 
 $f\ (linepath\ u\ v\ x))$ 
 $\leq B * (cmod\ y / 24 / C / B * 2 * C) + 2 * C * (cmod\ y /$ 
 $24 / C)$ 
proof (intro add_mono [OF mult_mono])
show  $cmod\ (f\ (linepath\ (shrink\ u)\ (shrink\ v)\ x)) \leq B$ 
using  $cmod\_fuv\ x$  by blast
have  $B * (12 * (e * cmod\ (u - v))) \leq 24 * e * C * B$ 
using  $e\ \langle B > 0 \rangle$   $diff\_2C$  [of  $u\ v$ ]  $uv$  by (auto simp: field_simps)
also have  $\dots \leq cmod\ y$ 
using  $\langle C > 0 \rangle$   $\langle B > 0 \rangle$   $e$  by (simp add: field_simps)
finally show  $cmod\ (shrink\ v - shrink\ u - (v - u)) \leq cmod\ y / 24 /$ 
 $C / B * 2 * C$ 
using  $\langle 0 < B \rangle$   $\langle 0 < C \rangle$  by (simp add: cmod_shr mult_ac divide_simps)
have  $cmod\ (f\ (linepath\ (shrink\ u)\ (shrink\ v)\ x) - f\ (linepath\ u\ v\ x))$ 
 $< cmod\ y / (24 * C)$ 
using  $x\ uv\ shr\_uv\ cmod\_less\_dt$ 
by (auto simp: hull_inc intro:  $d1\ interior\_subset$  [THEN subsetD]
linepath_in_convex_hull)
also have  $\dots \leq cmod\ y / cmod\ (v - u) / 12$ 
using  $False\ uv\ \langle C > 0 \rangle$   $diff\_2C$  [of  $v\ u$ ]  $ynz$ 
by (auto simp: field_split_simps hull_inc)
finally have  $cmod\ (f\ (linepath\ (shrink\ u)\ (shrink\ v)\ x) - f\ (linepath$ 
 $u\ v\ x)) \leq cmod\ y / cmod\ (v - u) / 12$ 
by simp
then show  $cmod\ (v - u) * cmod\ (f\ (linepath\ (shrink\ u)\ (shrink\ v)\ x)$ 
 $- f\ (linepath\ u\ v\ x))$ 
 $\leq 2 * C * (cmod\ y / 24 / C)$ 
using  $uv\ C$  by (simp add: field_simps)
qed (use  $\langle 0 < B \rangle$  in auto)
also have  $\dots \leq cmod\ y / 6$ 
by simp
finally have  $cmod\ (f\ (linepath\ (shrink\ u)\ (shrink\ v)\ x)) * cmod\ (shrink$ 

```

```

v - shrink u - (v - u) +
      cmod (v - u) * cmod (f (linepath (shrink u) (shrink v) x) -
f (linepath u v x))
      ≤ cmod y / 6 .
} note cmod_diff_le = this
have f_uv: continuous_on (closed_segment u v) f
by (blast intro: uv_continuous_on_subset [OF contf closed_segment_subset_convex_hull])
have **:  $\bigwedge f' x' f x::\text{complex}. f' * x' - f * x = f' * (x' - x) + x * (f' - f)$ 
  by (simp add: algebra_simps)
have norm (pathint (shrink u) (shrink v) - pathint u v)
  ≤ (B*(norm y / 24 / C / B)*2*C + (2*C)*(norm y / 24 / C)) * content
(cbox 0 (1::real))
  apply (rule has_integral_bound
    [of  $\lambda x. f(\text{linepath } (\text{shrink } u) (\text{shrink } v) x) * (\text{shrink } v - \text{shrink } u) - f(\text{linepath } u v x) * (v - u) - 0 \ 1$ ])
  using ynz <0 < B> <0 < C>
  apply (simp_all add: pathint_def has_integral_diff has_contour_integral_linepath
[symmetric] has_contour_integral_integral
  fpi_uv f_uv contour_integrable_continuous_linepath del: le_divide_eq_numeral1)
  apply (auto simp: ** norm_triangle_le norm_mult cmod_diff_le simp
del: le_divide_eq_numeral1)
  done
also have ... ≤ norm y / 6
  by simp
finally have norm (pathint (shrink u) (shrink v) - pathint u v) ≤ norm
y / 6 .
} note * = this
have norm (pathint (shrink a) (shrink b) - pathint a b) ≤ norm y / 6
  using False fpi_abc by (rule_tac *) (auto simp: hull_inc)
moreover
have norm (pathint (shrink b) (shrink c) - pathint b c) ≤ norm y / 6
  using False fpi_abc by (rule_tac *) (auto simp: hull_inc)
moreover
have norm (pathint (shrink c) (shrink a) - pathint c a) ≤ norm y / 6
  using False fpi_abc by (rule_tac *) (auto simp: hull_inc)
ultimately
have norm((pathint (shrink a) (shrink b) - pathint a b) +
  (pathint (shrink b) (shrink c) - pathint b c) + (pathint (shrink c)
  (shrink a) - pathint c a))
  ≤ norm y / 6 + norm y / 6 + norm y / 6
  by (metis norm_triangle_le add_mono)
also have ... = norm y / 2
  by simp
finally have norm((pathint (shrink a) (shrink b) + pathint (shrink b)
  (shrink c) + pathint (shrink c) (shrink a)) -
  (pathint a b + pathint b c + pathint c a))
  ≤ norm y / 2
  by (simp add: algebra_simps)

```

```

    then
      have norm(pathint a b + pathint b c + pathint c a) ≤ norm y / 2
        by (simp add: f_0_shrink) (metis (mono_tags) add commute minus_add_distrib norm_minus_cancel uminus_add_conv_diff)
      then have False
        using pi_eq_y ynz by auto
    }
    note * = this
    have uniformly_continuous_on (convex hull {a,b,c}) f
      by (simp add: contf compact_convex_hull compact_uniformly_continuous)
    moreover have norm y / (24 * C) > 0
      using ynz ⟨C > 0⟩ by auto
    ultimately obtain δ where δ > 0 and
      ∀ x ∈ convex hull {a, b, c}. ∀ x' ∈ convex hull {a, b, c}.
        dist x' x < δ → dist (f x') (f x) < cmod y / (24 * C)
      using ⟨C > 0⟩ ynz unfolding uniformly_continuous_on_def dist_norm
    by blast
      hence False using *[of δ] by (auto simp: dist_norm)
      then show ?thesis ..
    qed
  }
  moreover have f contour_integrable_on (linepath a b +++ linepath b c +++
linepath c a)
    using fabc contour_integrable_continuous_linepath by auto
  ultimately show ?thesis
    using has_contour_integral_integral by fastforce
  qed

```

2.4 Version allowing finite number of exceptional points

proposition *Cauchy_theorem_triangle_cofinite:*

assumes *continuous_on (convex hull {a,b,c}) f*

and *finite S*

and $(\bigwedge x. x \in \text{interior}(\text{convex hull } \{a,b,c\}) - S \implies f \text{ field_differentiable (at } x))$

shows $(f \text{ has_contour_integral } 0) (\text{linepath } a \text{ } b \text{ } +++ \text{ linepath } b \text{ } c \text{ } +++ \text{ linepath } c \text{ } a)$

using *assms*

proof (*induction card S arbitrary: a b c S rule: less_induct*)

case (*less S a b c*)

show *?case*

proof (*cases S={}*)

case *True with less show ?thesis*

by (*fastforce simp: holomorphic_on_def field_differentiable_at_within Cauchy_theorem_triangle_interior*)

next

case *False*

then obtain *d S' where d: S = insert d S' d ∉ S'*

by (*meson Set.set_insert all_not_in_conv*)

then show *?thesis*

```

proof (cases  $d \in \text{convex hull } \{a, b, c\}$ )
  case False
    show ( $f \text{ has\_contour\_integral } 0$ ) ( $\text{linepath } a \ b \ +++ \ \text{linepath } b \ c \ +++ \ \text{linepath } c \ a$ )
    proof (rule less.hyps)
      show  $\bigwedge x. x \in \text{interior } (\text{convex hull } \{a, b, c\}) - S' \implies f \text{ field\_differentiable at } x$ 
      using False d interior_subset by (auto intro!: less.prems)
    qed (use d less.prems in auto)
  next
    case True
    have *:  $\text{convex hull } \{a, b, d\} \subseteq \text{convex hull } \{a, b, c\}$ 
      by (meson True hull_subset insert_subset convex_hull_subset)
    have abd: ( $f \text{ has\_contour\_integral } 0$ ) ( $\text{linepath } a \ b \ +++ \ \text{linepath } b \ d \ +++ \ \text{linepath } d \ a$ )
    proof (rule less.hyps)
      show  $\bigwedge x. x \in \text{interior } (\text{convex hull } \{a, b, d\}) - S' \implies f \text{ field\_differentiable at } x$ 
      using d not_in_interior_convex_hull_3
      by (clarsimp intro!: less.prems) (metis * insert_absorb insert_subset interior_mono)
    qed (use d continuous_on_subset [OF _ *] less.prems in auto)
    have *:  $\text{convex hull } \{b, c, d\} \subseteq \text{convex hull } \{a, b, c\}$ 
      by (meson True hull_subset insert_subset convex_hull_subset)
    have bcd: ( $f \text{ has\_contour\_integral } 0$ ) ( $\text{linepath } b \ c \ +++ \ \text{linepath } c \ d \ +++ \ \text{linepath } d \ b$ )
    proof (rule less.hyps)
      show  $\bigwedge x. x \in \text{interior } (\text{convex hull } \{b, c, d\}) - S' \implies f \text{ field\_differentiable at } x$ 
      using d not_in_interior_convex_hull_3
      by (clarsimp intro!: less.prems) (metis * insert_absorb insert_subset interior_mono)
    qed (use d continuous_on_subset [OF _ *] less.prems in auto)
    have *:  $\text{convex hull } \{c, a, d\} \subseteq \text{convex hull } \{a, b, c\}$ 
      by (meson True hull_subset insert_subset convex_hull_subset)
    have cad: ( $f \text{ has\_contour\_integral } 0$ ) ( $\text{linepath } c \ a \ +++ \ \text{linepath } a \ d \ +++ \ \text{linepath } d \ c$ )
    proof (rule less.hyps)
      show  $\bigwedge x. x \in \text{interior } (\text{convex hull } \{c, a, d\}) - S' \implies f \text{ field\_differentiable at } x$ 
      using d not_in_interior_convex_hull_3
      by (clarsimp intro!: less.prems) (metis * insert_absorb insert_subset interior_mono)
    qed (use d continuous_on_subset [OF _ *] less.prems in auto)
    have  $f \text{ contour\_integrable\_on linepath } a \ b$ 
      using less.prems abd contour_integrable_joinD1 contour_integrable_on_def
by blast
    moreover have  $f \text{ contour\_integrable\_on linepath } b \ c$ 
      using less.prems bcd contour_integrable_joinD1 contour_integrable_on_def

```

```

by blast
  moreover have f contour_integrable_on linepath c a
    using less.prem3 cad contour_integrable_joinD1 contour_integrable_on_def
by blast
  ultimately have fpi: f contour_integrable_on (linepath a b +++ linepath b
c +++ linepath c a)
    by auto
  { fix y::complex
    assume fy: (f has_contour_integral y) (linepath a b +++ linepath b c +++
linepath c a)
      and ynz: y ≠ 0
      have cont_ad: continuous_on (closed_segment a d) f
      by (meson * continuous_on_subset less.prem3(1) segments_subset_convex_hull(3))
      have cont_bd: continuous_on (closed_segment b d) f
      by (meson True closed_segment_subset_convex_hull continuous_on_subset
hull_subset insert_subset less.prem3(1))
      have cont_cd: continuous_on (closed_segment c d) f
      by (meson * continuous_on_subset less.prem3(1) segments_subset_convex_hull(2))
      have contour_integral (linepath a b) f = - (contour_integral (linepath b
d) f + (contour_integral (linepath d a) f))
          contour_integral (linepath b c) f = - (contour_integral (linepath c d)
f + (contour_integral (linepath d b) f))
          contour_integral (linepath c a) f = - (contour_integral (linepath a d)
f + contour_integral (linepath d c) f)
      using has_chain_integral_chain_integral3 [OF abd]
          has_chain_integral_chain_integral3 [OF bcd]
          has_chain_integral_chain_integral3 [OF cad]
      by (simp_all add: algebra_simps add_eq_0_iff)
    then have ?thesis
      using cont_ad cont_bd cont_cd fy has_chain_integral_chain_integral3
contour_integral_reverse_linepath by fastforce
  }
  then show ?thesis
    using fpi contour_integrable_on_def by blast
qed
qed
qed

```

2.5 Cauchy's theorem for an open starlike set

lemma *starlike_convex_subset*:

assumes S : $a \in S$ *closed_segment* $b c \subseteq S$ **and** *subs*: $\bigwedge x. x \in S \implies \text{closed_segment } a x \subseteq S$

shows $\text{convex_hull } \{a, b, c\} \subseteq S$

proof –

have $\text{convex_hull } \{b, c\} \subseteq S$

using *assms*(2) *segment_convex_hull* **by** *auto*

then have $\bigwedge u v d. \llbracket 0 \leq u; 0 \leq v; u + v = 1; d \in \text{convex_hull } \{b, c\} \rrbracket \implies u *_{\mathbb{R}} a + v *_{\mathbb{R}} d \in S$

```

  by (meson subs convexD convex_closed_segment ends_in_segment subsetCE)
  then show ?thesis
  by (auto simp add: convex_hull_insert [of {b,c} a])
qed

```

lemma *triangle_contour_integrals_starlike_primitive:*

```

  assumes contf: continuous_on S f
    and S: a ∈ S open S
    and x: x ∈ S
    and subs:  $\bigwedge y. y \in S \implies \text{closed\_segment } a \ y \subseteq S$ 
    and zer:  $\bigwedge b \ c. \text{closed\_segment } b \ c \subseteq S$ 
       $\implies \text{contour\_integral } (\text{linepath } a \ b) \ f + \text{contour\_integral } (\text{linepath } b \ c) \ f +$ 
         $\text{contour\_integral } (\text{linepath } c \ a) \ f = 0$ 
  shows (( $\lambda x. \text{contour\_integral}(\text{linepath } a \ x) \ f$ ) has_field_derivative f x) (at x)
proof -
  let ?pathint =  $\lambda x \ y. \text{contour\_integral}(\text{linepath } x \ y) \ f$ 
  { fix e y
    assume e:  $0 < e$  and bxe:  $\text{ball } x \ e \subseteq S$  and close:  $\text{cmod } (y - x) < e$ 
    have y:  $y \in S$ 
      using bxe close by (force simp: dist_norm norm_minus_commute)
    have cont_ayf: continuous_on (closed_segment a y) f
      using contf continuous_on_subset subs y by blast
    have xys: closed_segment x y  $\subseteq S$ 
      by (metis bxe centre_in_ball close closed_segment_subset convex_ball dist_norm
        dual_order.trans e mem_ball norm_minus_commute)
    have ?pathint a y - ?pathint a x = ?pathint x y
      using zer [OF xys] contour_integral_reverse_linepath [OF cont_ayf] add_eq_0_iff
    by force
  } note [simp] = this
  { fix e::real
    assume e:  $0 < e$ 
    have cont_atx: continuous (at x) f
      using x S contf continuous_on_eq_continuous_at by blast
    then obtain d1 where d1:  $d1 > 0$  and d1_less:  $\bigwedge y. \text{cmod } (y - x) < d1 \implies$ 
       $\text{cmod } (f \ y - f \ x) < e/2$ 
      unfolding continuous_at Lim_at dist_norm using e
      by (drule_tac x=e/2 in spec) force
    obtain d2 where d2:  $d2 > 0$  ball x d2  $\subseteq S$  using  $\langle \text{open } S \rangle x$ 
      by (auto simp: open_contains_ball)
    have dpos:  $\min \ d1 \ d2 > 0$  using d1 d2 by simp
    { fix y
      assume yx:  $y \neq x$  and close:  $\text{cmod } (y - x) < \min \ d1 \ d2$ 
      have y:  $y \in S$ 
        using d2 close by (force simp: dist_norm norm_minus_commute)
      have closed_segment x y  $\subseteq S$ 
        using close d2 by (auto simp: dist_norm norm_minus_commute dest!:
          segment_bound(1))
      then have fry: f contour_integrable_on linepath x y

```



```

    by (metis contour_integrable_continuous_linepath continuous_on_subset
[OF contf])
  then obtain i where i: (f has_contour_integral i) (linepath x y)
    by (auto simp: contour_integrable_on_def)
  then have (( $\lambda w. f w - f x$ ) has_contour_integral (i - f x * (y - x)))
(linepath x y)
  by (rule has_contour_integral_diff [OF has_contour_integral_const_linepath])
  then have  $cmod (i - f x * (y - x)) \leq e / 2 * cmod (y - x)$ 
  proof (rule has_contour_integral_bound_linepath)
    show  $\bigwedge u. u \in \text{closed\_segment } x \ y \implies cmod (f u - f x) \leq e / 2$ 
      by (meson close_d1_less le_less_trans less_imp_le min.strict_boundedE
segment_bound1)
    qed (use e in simp)
  also have  $\dots < e * cmod (y - x)$ 
    by (simp add: e yx)
  finally have  $cmod (?pathint \ x \ y - f x * (y-x)) / cmod (y-x) < e$ 
    using i yx by (simp add: contour_integral_unique divide_less_eq)
  }
  then have  $\exists d > 0. \forall y. y \neq x \wedge cmod (y-x) < d \implies cmod (?pathint \ x \ y - f
x * (y-x)) / cmod (y-x) < e$ 
    using dpos by blast
  }
  then have  $(\lambda y. (?pathint \ x \ y - f x * (y - x)) /_R cmod (y - x)) -x \rightarrow 0$ 
    by (simp add: Lim_at_dist_norm_inverse_eq_divide)
  then have  $(\lambda y. (1 / cmod (y - x)) *_R (?pathint \ a \ y - (?pathint \ a \ x + f x * (y
- x)))) -x \rightarrow 0$ 
    using ‹open S› x
    by (force simp: dist_norm open_contains_ball inverse_eq_divide [symmetric]
eventually_at intro: Lim_transform [OF tendsto_eventually])
  then show ?thesis
    by (simp add: has_field_derivative_def has_derivative_at2 bounded_linear_mult_right)
qed

```

lemma *holomorphic_starlike_primitive:*

```

fixes f :: complex  $\Rightarrow$  complex
assumes contf: continuous_on S f
  and S: starlike S and os: open S
  and k: finite k
  and fcd:  $\bigwedge x. x \in S - k \implies f \text{ field\_differentiable at } x$ 
shows  $\exists g. \forall x \in S. (g \text{ has\_field\_derivative } f x) \text{ (at } x)$ 
proof -
  obtain a where a:  $a \in S$  and a_cs:  $\bigwedge x. x \in S \implies \text{closed\_segment } a \ x \subseteq S$ 
  using S by (auto simp: starlike_def)
  { fix x b c
    assume  $x \in S$   $\text{closed\_segment } b \ c \subseteq S$ 
    then have abcs:  $\text{convex\_hull } \{a, b, c\} \subseteq S$ 
      by (simp add: a a_cs starlike_convex_subset)
    then have continuous_on (convex_hull {a, b, c}) f

```

```

    by (simp add: continuous_on_subset [OF contf])
  then have (f has_contour_integral 0) (linepath a b +++ linepath b c +++
linepath c a)
    using abcs interior_subset by (force intro: fcd Cauchy_theorem_triangle_cofinite
[OF _ k])
  } note 0 = this
  show ?thesis
  proof (intro exI ballI)
    show  $\bigwedge x. x \in S \implies ((\lambda x. \text{contour\_integral (linepath a x) f}) \text{ has\_field\_derivative } f x) \text{ (at } x)$ 
      using 0 a a_cs contf has_chain_integral_chain_integral3 os triangle_contour_integrals_starlike_primitive
    by force
  qed
qed

```

lemma *Cauchy_theorem_starlike*:

```

[[open S; starlike S; finite k; continuous_on S f;
 $\bigwedge x. x \in S - k \implies f \text{ field\_differentiable at } x;$ 
valid_path g; path_image g  $\subseteq S$ ; pathfinish g = pathstart g]]
 $\implies (f \text{ has\_contour\_integral } 0) \ g$ 
by (metis holomorphic_starlike_primitive Cauchy_theorem_primitive at_within_open)

```

lemma *Cauchy_theorem_starlike_simple*:

```

[[open S; starlike S; f holomorphic_on S; valid_path g; path_image g  $\subseteq S$ ; pathfinish g = pathstart g]]
 $\implies (f \text{ has\_contour\_integral } 0) \ g$ 
using Cauchy_theorem_starlike [OF _ _ finite.emptyI]
by (simp add: holomorphic_on_imp_continuous_on holomorphic_on_imp_differentiable_at)

```

2.6 Cauchy's theorem for a convex set

For a convex set we can avoid assuming openness and boundary analyticity

lemma *triangle_contour_integrals_convex_primitive*:

```

assumes contf: continuous_on S f
  and S: a  $\in S$  convex S
  and x: x  $\in S$ 
  and zer:  $\bigwedge b \ c. [b \in S; c \in S]$ 
 $\implies \text{contour\_integral (linepath a b) } f + \text{contour\_integral (linepath } b \ c) \ f +$ 
 $\text{contour\_integral (linepath c a) } f = 0$ 
  shows  $((\lambda x. \text{contour\_integral (linepath a x) } f) \text{ has\_field\_derivative } f x) \text{ (at } x \text{ within } S)$ 

```

proof –

```

let ?pathint =  $\lambda x \ y. \text{contour\_integral (linepath x y) } f$ 
{ fix y
  assume y: y  $\in S$ 
  have cont_ayf: continuous_on (closed_segment a y) f
    using S y by (meson contf continuous_on_subset convex_contains_segment)
  have xys: closed_segment x y  $\subseteq S$ 

```

```

    using convex_contains_segment S x y by auto
  have ?pathint a y - ?pathint a x = ?pathint x y
  using zer [OF x y] contour_integral_reverse_linepath [OF cont_ayf] add_eq_0_iff
by force
} note [simp] = this
{ fix e::real
  assume e: 0 < e
  have cont_atx: continuous (at x within S) f
  using x S contf by (simp add: continuous_on_eq_continuous_within)
  then obtain d1 where d1: d1 > 0 and d1_less:  $\bigwedge y. \llbracket y \in S; cmod (y - x) < d1 \rrbracket \implies cmod (f y - f x) < e/2$ 
  unfolding continuous_within Lim_within dist_norm using e
  by (drule_tac x=e/2 in spec) force
  { fix y
    assume yx:  $y \neq x$  and close:  $cmod (y - x) < d1$  and y:  $y \in S$ 
    have fxy: f contour_integrable_on linepath x y
    using convex_contains_segment S x y
    by (blast intro!: contour_integrable_continuous_linepath continuous_on_subset [OF contf])
    then obtain i where i: (f has_contour_integral i) (linepath x y)
    by (auto simp: contour_integrable_on_def)
    then have (( $\lambda w. f w - f x$ ) has_contour_integral (i - f x * (y - x)))
    (linepath x y)
    by (rule has_contour_integral_diff [OF _ has_contour_integral_const_linepath])
    then have  $cmod (i - f x * (y - x)) \leq e / 2 * cmod (y - x)$ 
    proof (rule has_contour_integral_bound_linepath)
      show  $\bigwedge u. u \in \text{closed\_segment } x \ y \implies cmod (f u - f x) \leq e / 2$ 
      by (meson assms(3) close convex_contains_segment d1_less le_less_trans less_imp_le segment_bound1 subset_iff x y)
    qed (use e in simp)
    also have  $\dots < e * cmod (y - x)$ 
    by (simp add: e yx)
    finally have  $cmod (?pathint x y - f x * (y-x)) / cmod (y-x) < e$ 
    using i yx by (simp add: contour_integral_unique divide_less_eq)
  }
  then have  $\exists d > 0. \forall y \in S. y \neq x \wedge cmod (y-x) < d \longrightarrow cmod (?pathint x y - f x * (y-x)) / cmod (y-x) < e$ 
  using d1 by blast
}
then have (( $\lambda y. (?pathint x y - f x * (y - x)) /_R cmod (y - x)$ )  $\longrightarrow 0$ ) (at x within S)
by (simp add: Lim_within dist_norm inverse_eq_divide)
then have (( $\lambda y. (1 / cmod (y - x)) *_R (?pathint a y - (?pathint a x + f x * (y - x)))$ )  $\longrightarrow 0$ )
(at x within S)
using linordered_field_no_ub
by (force simp: inverse_eq_divide [symmetric] eventually_at intro: Lim_transform [OF _ tendsto_eventually])
then show ?thesis

```

by (simp add: has_field_derivative_def has_derivative_within bounded_linear_mult_right)
qed

lemma *contour_integral_convex_primitive*:

assumes *convex S continuous_on S f*
 $\bigwedge a\ b\ c. \llbracket a \in S; b \in S; c \in S \rrbracket \implies (f \text{ has_contour_integral } 0) (\text{linepath } a\ b\ \text{+++ linepath } b\ c\ \text{+++ linepath } c\ a)$
obtains *g where* $\bigwedge x. x \in S \implies (g \text{ has_field_derivative } f\ x) (\text{at } x \text{ within } S)$
proof (cases *S = {}*)
 case *False*
with *assms that show ?thesis*
 by (blast intro: *triangle_contour_integrals_convex_primitive has_chain_integral_chain_integral3*)
 qed *auto*

lemma *holomorphic_convex_primitive*:

fixes *f :: complex \Rightarrow complex*
 assumes *convex S finite K and contf: continuous_on S f*
 and *fd: $\bigwedge x. x \in \text{interior } S - K \implies f \text{ field_differentiable at } x$*
obtains *g where* $\bigwedge x. x \in S \implies (g \text{ has_field_derivative } f\ x) (\text{at } x \text{ within } S)$
proof (rule *contour_integral_convex_primitive [OF <convex S> contf Cauchy_theorem_triangle_cofinite]*)
 have **: convex hull {a, b, c} \subseteq S if $a \in S\ b \in S\ c \in S$ for $a\ b\ c$*
 by (simp add: <convex S> hull_minimal that)
 show *continuous_on (convex hull {a, b, c}) f if $a \in S\ b \in S\ c \in S$ for $a\ b\ c$*
 by (meson * contf continuous_on_subset that)
 show *f field_differentiable at x if $a \in S\ b \in S\ c \in S\ x \in \text{interior (convex hull } \{a, b, c\}) - K$ for $a\ b\ c\ x$*
 by (metis * DiffD1 DiffD2 DiffI fd interior_mono subsetCE that)
 qed (use *assms in <force+>*)

lemma *holomorphic_convex_primitive'*:

fixes *f :: complex \Rightarrow complex*
 assumes *convex S and open S and f holomorphic_on S*
obtains *g where* $\bigwedge x. x \in S \implies (g \text{ has_field_derivative } f\ x) (\text{at } x \text{ within } S)$
proof (rule *holomorphic_convex_primitive*)
 fix *x* assume *x \in interior S - {}*
 with *assms show f field_differentiable at x*
 by (auto intro!: *holomorphic_on_imp_differentiable_at simp: interior_open*)
 qed (use *assms in <auto intro: holomorphic_on_imp_continuous_on>*)

corollary *Cauchy_theorem_convex*:

$\llbracket \text{continuous_on } S\ f; \text{convex } S; \text{finite } K;$
 $\bigwedge x. x \in \text{interior } S - K \implies f \text{ field_differentiable at } x;$
 $\text{valid_path } g; \text{path_image } g \subseteq S; \text{pathfinish } g = \text{pathstart } g \rrbracket$
 $\implies (f \text{ has_contour_integral } 0) g$
 by (metis *holomorphic_convex_primitive Cauchy_theorem_primitive*)

corollary *Cauchy_theorem_convex_simple*:

assumes *holf: f holomorphic_on S*
 and *convex S valid_path g path_image g \subseteq S pathfinish g = pathstart g*

```

shows (f has_contour_integral 0) g
proof -
  have f holomorphic_on interior S
    by (meson holf holomorphic_on_subset interior_subset)
  with Cauchy_theorem_convex [where K = {}] show ?thesis
    using assms
    by (metis Diff_empty finite.emptyI holomorphic_on_imp_continuous_on holomorphic_on_imp_differentiable_at open_interior)
qed

```

In particular for a disc

```

corollary Cauchy_theorem_disc:
  [|finite K; continuous_on (cball a e) f;
   $\bigwedge x. x \in \text{ball } a \ e - K \implies f \text{ field\_differentiable at } x;$ 
  valid_path g; path_image g  $\subseteq$  cball a e;
  pathfinish g = pathstart g]  $\implies$  (f has_contour_integral 0) g
by (auto intro: Cauchy_theorem_convex)

```

```

corollary Cauchy_theorem_disc_simple:
  [|f holomorphic_on (ball a e); valid_path g; path_image g  $\subseteq$  ball a e;
  pathfinish g = pathstart g]  $\implies$  (f has_contour_integral 0) g
by (simp add: Cauchy_theorem_convex_simple)

```

2.7 Generalize integrability to local primitives

```

lemma contour_integral_local_primitive_lemma:
  fixes f :: complex  $\Rightarrow$  complex
  assumes gpd: g piecewise_differentiable_on {a..b}
    and dh:  $\bigwedge x. x \in S \implies (f \text{ has\_field\_derivative } f' \ x) \text{ (at } x \text{ within } S)$ 
    and gs:  $\bigwedge x. x \in \{a..b\} \implies g \ x \in S$ 
  shows
    ( $\lambda x. f' (g \ x) * \text{vector\_derivative } g \text{ (at } x \text{ within } \{a..b\})$ ) integrable_on {a..b}
proof (cases cbox a b = {})
  case False
  then show ?thesis
    unfolding integrable_on_def by (auto intro: assms contour_integral_primitive_lemma)
qed auto

```

```

lemma contour_integral_local_primitive_any:
  fixes f :: complex  $\Rightarrow$  complex
  assumes gpd: g piecewise_differentiable_on {a..b}
    and dh:  $\bigwedge x. x \in S$ 
       $\implies \exists d \ h. 0 < d \wedge$ 
      ( $\forall y. \text{norm}(y - x) < d \longrightarrow (h \text{ has\_field\_derivative } f \ y) \text{ (at } y$ 
      within S))
    and gs:  $\bigwedge x. x \in \{a..b\} \implies g \ x \in S$ 
  shows ( $\lambda x. f(g \ x) * \text{vector\_derivative } g \text{ (at } x)$ ) integrable_on {a..b}
proof -
  { fix x

```

```

assume  $x: a \leq x \leq b$ 
obtain  $d h$  where  $d: 0 < d$ 
      and  $h: (\bigwedge y. \text{norm}(y - g x) < d \implies (h \text{ has\_field\_derivative } f y))$  (at
       $y$  within  $S$ )
using  $x$   $gs$   $dh$  by (metis atLeastAtMost_iff)
have continuous_on  $\{a..b\}$   $g$  using gpd piecewise_differentiable_on_def by
blast
then obtain  $e$  where  $e: e > 0$  and lessd:  $\bigwedge x'. x' \in \{a..b\} \implies |x' - x| < e$ 
 $\implies$  cmod  $(g x' - g x) < d$ 
using  $x$   $d$  by (fastforce simp: dist_norm continuous_on_iff)
have  $\exists e > 0. \forall u v. u \leq x \wedge x \leq v \wedge \{u..v\} \subseteq \text{ball } x e \wedge (u \leq v \longrightarrow a \leq u \wedge$ 
 $v \leq b) \longrightarrow$ 
       $(\lambda x. f (g x) * \text{vector\_derivative } g \text{ (at } x)) \text{ integrable\_on } \{u..v\}$ 
proof -
  have  $(\lambda x. f (g x) * \text{vector\_derivative } g \text{ (at } x \text{ within } \{u..v\})) \text{ integrable\_on}$ 
 $\{u..v\}$ 
    if  $u \leq x \leq v$  and ball:  $\{u..v\} \subseteq \text{ball } x e$  and auvb:  $u \leq v \implies a \leq u \wedge v$ 
 $\leq b$ 
    for  $u v$ 
    proof (rule contour_integral_local_primitive_lemma)
      show  $g$  piecewise_differentiable_on  $\{u..v\}$ 
      by (metis atLeastAtMost_subset_iff gpd piecewise_differentiable_on_subset
      auvb)
      show  $\bigwedge x. x \in g^{-1} \{u..v\} \implies (h \text{ has\_field\_derivative } f x)$  (at  $x$  within  $g^{-1}$ 
 $\{u..v\}$ )
      using that by (force simp: ball_def dist_norm intro: lessd gs DERIV_subset
      [OF h])
    qed auto
    then show ?thesis
      using  $e$  integrable_on_localized_vector_derivative by blast
    qed
  } then
show ?thesis
  by (force simp: intro!: integrable_on_little_subintervals [of  $a b$ , simplified])
qed

```

lemma contour_integral_local_primitive:

```

fixes  $f :: \text{complex} \Rightarrow \text{complex}$ 
assumes  $g: \text{valid\_path } g \text{ path\_image } g \subseteq S$ 
      and  $dh: \bigwedge x. x \in S$ 
       $\implies \exists d h. 0 < d \wedge$ 
       $(\forall y. \text{norm}(y - x) < d \longrightarrow (h \text{ has\_field\_derivative } f y))$  (at  $y$ 
      within  $S$ )
shows  $f$  contour_integrable_on  $g$ 
proof -
have  $(\lambda x. f (g x) * \text{vector\_derivative } g \text{ (at } x)) \text{ integrable\_on } \{0..1\}$ 
using contour_integral_local_primitive_any [OF _ dh]  $g$ 
unfolding path_image_def valid_path_def
by (metis (no_types, lifting) image_subset_iff piecewise_C1_imp_differentiable)

```

```

then show ?thesis
  using contour_integrable_on by presburger
qed

```

In particular if a function is holomorphic

lemma *contour_integrable_holomorphic*:

```

assumes contf: continuous_on S f
  and os: open S
  and k: finite k
  and g: valid_path g path_image g  $\subseteq$  S
  and fed:  $\bigwedge x. x \in S - k \implies f$  field_differentiable at x
shows f contour_integrable_on g
proof -
  { fix z
    assume z: z  $\in$  S
    obtain d where d  $>$  0 and d: ball z d  $\subseteq$  S using  $\langle$ open S $\rangle$  z
      by (auto simp: open_contains_ball)
    then have contfb: continuous_on (ball z d) f
      using contf continuous_on_subset by blast
    obtain h where  $\forall y \in \text{ball } z \ d. (h \text{ has\_field\_derivative } f \ y) \text{ (at } y \text{ within ball } z \ d)$ 
      by (metis holomorphic_convex_primitive [OF convex_ball k contfb fed] d
interior_subset Diff_iff_subsetD)
    then have  $\forall y \in \text{ball } z \ d. (h \text{ has\_field\_derivative } f \ y) \text{ (at } y \text{ within } S)$ 
      by (metis open_ball_at_within_open d os subsetCE)
    then have  $\exists h. (\forall y. \text{cmod } (y - z) < d \implies (h \text{ has\_field\_derivative } f \ y) \text{ (at } y \text{ within } S))$ 
      by (force simp: dist_norm norm_minus_commute)
    then have  $\exists d \ h. 0 < d \wedge (\forall y. \text{cmod } (y - z) < d \implies (h \text{ has\_field\_derivative } f \ y) \text{ (at } y \text{ within } S))$ 
      using  $\langle$ 0  $<$  d $\rangle$  by blast
  }
then show ?thesis
  by (rule contour_integral_local_primitive [OF g])
qed

```

lemma *contour_integrable_holomorphic_simple*:

```

assumes fh: f holomorphic_on S
  and os: open S
  and g: valid_path g path_image g  $\subseteq$  S
shows f contour_integrable_on g
proof -
  have  $\bigwedge x. x \in S \implies f$  field_differentiable at x
    using fh holomorphic_on_imp_differentiable_at by blast
  moreover have continuous_on S f
    by (simp add: fh holomorphic_on_imp_continuous_on)
  ultimately show ?thesis
    by (metis Diff_empty contour_integrable_holomorphic finite.emptyI g os)
qed

```

lemma *continuous_on_inversediff*:

fixes $z:: 'a::real_normed_field$ **shows** $z \notin S \implies \text{continuous_on } S (\lambda w. 1 / (w - z))$
by (*rule continuous_intros | force*)⁺

lemma *contour_integrable_inversediff*:

assumes g : *valid_path* g
and *notin*: $z \notin \text{path_image } g$
shows $(\lambda w. 1 / (w - z))$ *contour_integrable_on* g
proof (*rule contour_integrable_holomorphic_simple*)
show $(\lambda w. 1 / (w - z))$ *holomorphic_on* $UNIV - \{z\}$
by (*auto simp: holomorphic_on_open open_delete intro!: derivative_eq_intros*)
qed (*use assms in auto*)

Key fact that path integral is the same for a "nearby" path. This is the main lemma for the homotopy form of Cauchy's theorem and is also useful if we want "without loss of generality" to assume some nice properties of a path (e.g. smoothness). It can also be used to define the integrals of analytic functions over arbitrary continuous paths. This is just done for winding numbers now.

A technical definition to avoid duplication of similar proofs, for paths joined at the ends versus looping paths

definition *linked_paths* :: *bool* \Rightarrow (*real* \Rightarrow 'a) \Rightarrow (*real* \Rightarrow 'a::*topological_space*) \Rightarrow *bool*

where *linked_paths attends* g h ==

(*if attends then pathstart* h = *pathstart* g \wedge *pathfinish* h = *pathfinish* g
else *pathfinish* g = *pathstart* g \wedge *pathfinish* h = *pathstart* h)

This formulation covers two cases: g and h share their start and end points; g and h both loop upon themselves.

lemma *contour_integral_nearby*:

assumes *os*: *open* S **and** p : *path* p *path_image* $p \subseteq S$
shows $\exists d. 0 < d \wedge$
 $(\forall g$ $h. \text{valid_path } g \wedge \text{valid_path } h \wedge$
 $(\forall t \in \{0..1\}. \text{norm}(g\ t - p\ t) < d \wedge \text{norm}(h\ t - p\ t) < d) \wedge$
linked_paths attends g h
 $\longrightarrow \text{path_image } g \subseteq S \wedge \text{path_image } h \subseteq S \wedge$
 $(\forall f. f \text{ holomorphic_on } S \longrightarrow \text{contour_integral } h\ f = \text{contour_integral } g\ f))$

proof –

have $\forall z. \exists e. z \in \text{path_image } p \longrightarrow 0 < e \wedge \text{ball } z\ e \subseteq S$
using *open_contains_ball os p(2)* **by** *blast*
then obtain *ee* **where** $ee: \bigwedge z. z \in \text{path_image } p \implies 0 < ee\ z \wedge \text{ball } z\ (ee\ z) \subseteq S$
by *metis*
define *cover* **where** $\text{cover} = (\lambda z. \text{ball } z\ (ee\ z/3)) \text{ ` } (\text{path_image } p)$
have *compact* (*path_image* p)


```

  by (metis p(1) compact_path_image)
moreover have path_image p  $\subseteq$  ( $\bigcup c \in \text{path\_image } p. \text{ball } c (ee \ c / 3)$ )
  using ee by auto
ultimately have  $\exists D \subseteq \text{cover}. \text{finite } D \wedge \text{path\_image } p \subseteq \bigcup D$ 
  by (simp add: compact_eq_Heine_Borel_cover_def)
then obtain D where D:  $D \subseteq \text{cover}$  finite D path_image p  $\subseteq \bigcup D$ 
  by blast
then obtain k where k:  $k \subseteq \{0..1\}$  finite k and D_eq:  $D = ((\lambda z. \text{ball } z (ee \ z / 3)) \circ p) \text{ ' } k$ 
  unfolding cover_def path_image_def image_comp
  by (meson finite_subset_image)
then have kne:  $k \neq \{\}$ 
  using D by auto
have pi:  $\bigwedge i. i \in k \implies p \ i \in \text{path\_image } p$ 
  using k by (auto simp: path_image_def)
then have eepi:  $\bigwedge i. i \in k \implies 0 < ee((p \ i))$ 
  by (metis ee)
define e where  $e = \text{Min}((ee \circ p) \text{ ' } k)$ 
have fin_eep: finite ((ee  $\circ$  p) ' k)
  using k by blast
have  $0 < e$ 
  using ee k by (simp add: kne e_def Min_gr_iff [OF fin_eep] eepi)
have uniformly_continuous_on {0..1} p
  using p by (simp add: path_def compact_uniformly_continuous)
then obtain d::real where d:  $d > 0$ 
  and de:  $\bigwedge x \ x'. |x' - x| < d \implies x \in \{0..1\} \implies x' \in \{0..1\} \implies cmod (p \ x' - p \ x) < e/3$ 
  unfolding uniformly_continuous_on_def dist_norm real_norm_def
  by (metis divide_pos_pos  $\langle 0 < e \rangle$  zero_less_numeral)
then obtain N::nat where N:  $N > 0$  inverse N  $< d$ 
  using real_arch_inverse [of d] by auto
show ?thesis
proof (intro exI conjI allI; clarify?)
  show  $e/3 > 0$ 
    using  $\langle 0 < e \rangle$  by simp
  fix g h
  assume g: valid_path g and ghp:  $\forall t \in \{0..1\}. cmod (g \ t - p \ t) < e / 3 \wedge cmod (h \ t - p \ t) < e / 3$ 
    and h: valid_path h
    and joins: linked_paths attends g h
  { fix t::real
    assume t:  $0 \leq t \ t \leq 1$ 
    then obtain u where u:  $u \in k$  and ptu:  $p \ t \in \text{ball}(p \ u) (ee(p \ u) / 3)$ 
      using  $\langle \text{path\_image } p \subseteq \bigcup D \rangle$  D_eq by (force simp: path_image_def)
    then have ele:  $e \leq ee (p \ u)$  using fin_eep
      by (simp add: e_def)
    have  $cmod (g \ t - p \ t) < e / 3$   $cmod (h \ t - p \ t) < e / 3$ 
      using ghp t by auto
    with ele have  $cmod (g \ t - p \ t) < ee (p \ u) / 3$ 

```

```

      cmod (h t - p t) < ee (p u) / 3
    by linarith+
  then have g t ∈ ball(p u) (ee(p u)) h t ∈ ball(p u) (ee(p u))
  using norm_diff_triangle_ineq [of g t p t p t p u]
       norm_diff_triangle_ineq [of h t p t p t p u] ptu eepi u
  by (force simp: dist_norm ball_def norm_minus_commute)+
  then have g t ∈ S h t ∈ S using ee u k
  by (auto simp: path_image_def ball_def)
}
then have ghs: path_image g ⊆ S path_image h ⊆ S
  by (auto simp: path_image_def)
moreover
{ fix f
  assume fhols: f holomorphic_on S
  then have fpa: f contour_integrable_on g f contour_integrable_on h
  using g ghs h holomorphic_on_imp_continuous_on os contour_integrable_holomorphic_simple
  by blast+
  have contf: continuous_on S f
  by (simp add: fhols holomorphic_on_imp_continuous_on)
  { fix z
    assume z: z ∈ path_image p
    have f holomorphic_on ball z (ee z)
    using fhols ee z holomorphic_on_subset by blast
    then have ∃ff. (∀ w ∈ ball z (ee z). (ff has_field_derivative f w) (at w))
    using holomorphic_convex_primitive [of ball z (ee z) {} f, simplified]
    by (metis open_ball_at_within_open holomorphic_on_def holomor-
    phic_on_imp_continuous_on mem_ball)
  }
  then obtain ff where ff:
    ⋀z w. [z ∈ path_image p; w ∈ ball z (ee z)] ⇒ (ff z has_field_derivative
  f w) (at w)
  by metis
  { fix n
    assume n: n ≤ N
    then have contour_integral(subpath 0 (n/N) h) f - contour_integral(subpath
  0 (n/N) g) f =
      contour_integral(linepath (g(n/N)) (h(n/N))) f - con-
  tour_integral(linepath (g 0) (h 0)) f
    proof (induct n)
      case 0 show ?case by simp
    next
      case (Suc n)
      obtain t where t: t ∈ k and p (n/N) ∈ ball(p t) (ee(p t) / 3)
      using ⟨path_image p ⊆ ⋃ D⟩ [THEN subsetD, where c=p (n/N)] D_eq
  N Suc.premis
      by (force simp: path_image_def)
      then have ptu: cmod (p t - p (n/N)) < ee (p t) / 3
      by (simp add: dist_norm)
      have e3le: e/3 ≤ ee (p t) / 3 using fin_eep t

```

```

    by (simp add: e_def)
  { fix x
    assume x:  $n/N \leq x \leq (1+n)/N$ 
    then have nN01:  $0 \leq n/N \wedge (1+n)/N \leq 1$ 
      using Suc.prem1 by auto
    then have x01:  $0 \leq x \leq 1$ 
      using x by linarith+
    have cmod (p t - p x) < ee (p t) / 3 + e/3
    proof (rule norm_diff_triangle_less [OF ptu de])
      show |real n / real N - x| < d
        using x N by (auto simp: field_simps)
    qed (use x01 Suc.prem1 in auto)
    then have ptx: cmod (p t - p x) < 2*ee (p t)/3
      using e3le eepi [OF t] by simp
    have cmod (p t - g x) < 2*ee (p t)/3 + e/3
      using ghp x01
    by (force simp add: norm_minus_commute intro!: norm_diff_triangle_less
      [OF ptx])
    also have ... ≤ ee (p t)
      using e3le eepi [OF t] by simp
    finally have gg: cmod (p t - g x) < ee (p t) .
    have cmod (p t - h x) < 2*ee (p t)/3 + e/3
      using ghp x01
    by (force simp add: norm_minus_commute intro!: norm_diff_triangle_less
      [OF ptx])
    also have ... ≤ ee (p t)
      using e3le eepi [OF t] by simp
    finally have cmod (p t - g x) < ee (p t) cmod (p t - h x) < ee (p t)
      using gg by auto
  } note ptgh_ee = this
  have closed_segment (g (n/N)) (h (n/N)) = path_image (linepath (h
(n/N)) (g (n/N)))
    by (simp add: closed_segment_commute)
  also have pi_hgn: ... ⊆ ball (p t) (ee (p t))
    using ptgh_ee [of n/N] Suc.prem1
  by (auto simp: field_simps dist_norm dest: segment_furthest_le [where
y=p t])
  finally have gh_ns: closed_segment (g (n/N)) (h (n/N)) ⊆ S
    using ee pi t by blast
  have pi_ghn': path_image (linepath (g ((1+n)/N)) (h ((1+n)/N)))
⊆ ball (p t) (ee (p t))
    using ptgh_ee [of (1+n)/N] Suc.prem1
  by (auto simp: field_simps dist_norm dest: segment_furthest_le [where
y=p t])
  then have gh_n's: closed_segment (g ((1+n)/N)) (h ((1+n)/N))
⊆ S
    using ⟨N>0⟩ Suc.prem1 ee pi t
  by (auto simp: Path_Connected.path_image_join field_simps)
  have pi_subset_ball:

```

```

      path_image (subpath (n/N) ((1+n) / N) g +++ linepath (g ((1+n)
/ N)) (h ((1+n) / N)) +++
      subpath ((1+n) / N) (n/N) h +++ linepath (h (n/N)) (g
(n/N)))
    ⊆ ball (p t) (ee (p t))
  proof (intro subset_path_image_join pi_hgn pi_ghn')
    show path_image (subpath (n/N) ((1+n) / N) g) ⊆ ball (p t) (ee (p t))
      path_image (subpath ((1+n) / N) (n/N) h) ⊆ ball (p t) (ee (p t))
    using ‹N>0› Suc.premis
    by (auto simp: path_image_subpath dist_norm_field_simps ptgh_ee)
  qed
  have pi0: (f has_contour_integral 0)
    (subpath (n / N) ((Suc n)/N) g +++ linepath(g ((Suc n) / N))
(h((Suc n) / N)) +++
    subpath ((Suc n) / N) (n/N) h +++ linepath(h (n/N)) (g
(n/N)))
  proof (rule Cauchy_theorem_primitive)
    show ∧x. x ∈ ball (p t) (ee (p t))
      ⇒ (ff (p t) has_field_derivative f x) (at x within ball (p t) (ee
(p t)))
    by (metis ff_open_ball at_within_open pi t)
  qed (use Suc.premis pi_subset_ball in ‹simp_all add: valid_path_subpath
g h›)
  have fpa1: f contour_integrable_on subpath (n/N) (real (Suc n) / real N)
g
    using Suc.premis by (simp add: contour_integrable_subpath g fpa)
  have fpa2: f contour_integrable_on linepath (g (real (Suc n) / real N)) (h
(real (Suc n) / real N))
    using gh_n's
    by (auto intro!: contour_integrable_continuous_linepath continu-
ous_on_subset [OF contf])
  have fpa3: f contour_integrable_on linepath (h (n/N)) (g (n/N))
    using gh_ns
    by (auto simp: closed_segment_commute intro!: contour_integrable_continuous_linepath
continuous_on_subset [OF contf])
  have eq0: contour_integral (subpath (n/N) ((Suc n) / real N) g) f +
    contour_integral (linepath (g ((Suc n) / N)) (h ((Suc n) / N))) f
+
    contour_integral (subpath ((Suc n) / N) (n/N) h) f +
    contour_integral (linepath (h (n/N)) (g (n/N))) f = 0
    using contour_integral_unique [OF pi0] Suc.premis
  by (simp add: g h fpa valid_path_subpath contour_integrable_subpath
fpa1 fpa2 fpa3 algebra_simps del: of_nat_Suc)
  have *: ∧hn he hn' gn gd gn' hgn ghn gh0 ghn'.
    [[hn - gn = ghn - gh0;
gd + ghn' + he + hgn = (0::complex);
hn - he = hn'; gn + gd = gn'; hgn = -ghn]] ⇒ hn' - gn' =
ghn' - gh0
    by (auto simp: algebra_simps)

```

```

      have contour_integral (subpath 0 (n/N) h) f - contour_integral (subpath
((Suc n) / N) (n/N) h) f =
        contour_integral (subpath 0 (n/N) h) f + contour_integral (subpath
(n/N) ((Suc n) / N) h) f
      unfolding reversepath_subpath [symmetric, of ((Suc n) / N)]
      using Suc.premis by (simp add: h fpa contour_integral_reversepath
valid_path_subpath contour_integrable_subpath)
      also have ... = contour_integral (subpath 0 ((Suc n) / N) h) f
      using Suc.premis by (simp add: contour_integral_subpath_combine h
fpa)
      finally have pi0_eq:
        contour_integral (subpath 0 (n/N) h) f - contour_integral (subpath
((Suc n) / N) (n/N) h) f =
          contour_integral (subpath 0 ((Suc n) / N) h) f .
      show ?case
      proof (rule * [OF Suc.hyps eq0 pi0_eq])
        show contour_integral (subpath 0 (n/N) g) f +
          contour_integral (subpath (n/N) ((Suc n) / N) g) f =
            contour_integral (subpath 0 ((Suc n) / N) g) f
          using Suc.premis contour_integral_subpath_combine fpa(1) g by auto
        show contour_integral (linepath (h (n/N)) (g (n/N))) f = - con-
tour_integral (linepath (g (n/N)) (h (n/N))) f
          by (metis contour_integral_unique fpa3 has_contour_integral_integral
has_contour_integral_reverse_linepath)
        qed (use Suc.premis in auto)
      qed
    } note ind = this
    have contour_integral h f = contour_integral g f
      using ind [OF order_refl] N joins
    by (simp add: linked_paths_def pathstart_def pathfinish_def split: if_split_asm)
  }
  ultimately
  show path_image g  $\subseteq$  S  $\wedge$  path_image h  $\subseteq$  S  $\wedge$  ( $\forall f. f$  holomorphic_on S  $\longrightarrow$ 
contour_integral h f = contour_integral g f)
    by metis
  qed
qed

```

lemma

```

assumes open S path p path_image p  $\subseteq$  S
shows contour_integral_nearby_ends:
   $\exists d. 0 < d \wedge$ 
    ( $\forall g h. valid\_path g \wedge valid\_path h \wedge$ 
      ( $\forall t \in \{0..1\}. norm(g t - p t) < d \wedge norm(h t - p t) < d) \wedge$ 
      pathstart h = pathstart g  $\wedge$  pathfinish h = pathfinish g
       $\longrightarrow$  path_image g  $\subseteq$  S  $\wedge$ 
        path_image h  $\subseteq$  S  $\wedge$ 
        ( $\forall f. f$  holomorphic_on S

```

```

       $\longrightarrow \text{contour\_integral } h f = \text{contour\_integral } g f))$ 
and contour_integral_nearby_loops:
   $\exists d. 0 < d \wedge$ 
     $(\forall g h. \text{valid\_path } g \wedge \text{valid\_path } h \wedge$ 
       $(\forall t \in \{0..1\}. \text{norm}(g t - p t) < d \wedge \text{norm}(h t - p t) < d) \wedge$ 
       $\text{pathfinish } g = \text{pathstart } g \wedge \text{pathfinish } h = \text{pathstart } h$ 
       $\longrightarrow \text{path\_image } g \subseteq S \wedge$ 
       $\text{path\_image } h \subseteq S \wedge$ 
       $(\forall f. f \text{ holomorphic\_on } S$ 
         $\longrightarrow \text{contour\_integral } h f = \text{contour\_integral } g f))$ 
using contour_integral_nearby [OF assms, where atends=True]
using contour_integral_nearby [OF assms, where atends=False]
unfolding linked_paths_def by simp_all

lemma contour_integral_bound_exists:
assumes S: open S
  and g: valid_path g
  and pag: path_image g  $\subseteq S$ 
shows  $\exists L. 0 < L \wedge$ 
   $(\forall f B. f \text{ holomorphic\_on } S \wedge (\forall z \in S. \text{norm}(f z) \leq B)$ 
     $\longrightarrow \text{norm}(\text{contour\_integral } g f) \leq L * B)$ 

proof –
  have path g using g
  by (simp add: valid_path_imp_path)
  then obtain d::real and p
  where  $0 < d$ 
    and p: polynomial_function p path_image p  $\subseteq S$ 
  and pi:  $\bigwedge f. f \text{ holomorphic\_on } S \implies \text{contour\_integral } g f = \text{contour\_integral } p f$ 

  using contour_integral_nearby_ends [OF S  $\langle \text{path } g \rangle$  pag]
  by (metis cancel_comm_monoid_add_class.diff_cancel g norm_zero path_approx_polynomial_function
valid_path_polynomial_function)
  then obtain p' where p': polynomial_function p'
     $\bigwedge x. (p \text{ has\_vector\_derivative } (p' x)) \text{ (at } x)$ 
  by (blast intro: has_vector_derivative_polynomial_function_that)
  then have bounded(p' ‘  $\{0..1\}$ )
    using continuous_on_polynomial_function
  by (force simp: intro!: compact_imp_bounded_compact_continuous_image)
  then obtain L where  $L: L > 0$  and nop':  $\bigwedge x. \llbracket 0 \leq x; x \leq 1 \rrbracket \implies \text{norm } (p' x) \leq L$ 
  by (force simp: bounded_pos)
  { fix f B
    assume f: f holomorphic_on S and B:  $\bigwedge z. z \in S \implies \text{cmod } (f z) \leq B$ 
    then have f contour_integrable_on p  $\wedge$  valid_path p
      using p S
    by (blast intro: valid_path_polynomial_function contour_integrable_holomorphic_simple
holomorphic_on_imp_continuous_on)
    moreover have cmod (vector_derivative p (at x)) * cmod (f (p x))  $\leq L * B$ 
    if  $0 \leq x \leq 1$  for x
  }

```

```

proof (rule mult_mono)
  show cmod (vector_derivative p (at x)) ≤ L
    by (metis nop' p'(2) that vector_derivative_at)
  show cmod (f (p x)) ≤ B
    by (metis B atLeastAtMost_iff imageI p(2) path_defs(4) subset_eq that)
qed (use ‹L>0› in auto)
ultimately
have cmod (integral {0..1} (λx. f (p x) * vector_derivative p (at x))) ≤ L * B
  by (intro order_trans [OF integral_norm_bound_integral])
    (auto simp: mult.commute norm_mult contour_integrable_on)
then have cmod (contour_integral g f) ≤ L * B
  using contour_integral_integral f pi by presburger
} then
show ?thesis using ‹L > 0›
  by (intro exI[of _ L]) auto
qed

```

2.8 Homotopy forms of Cauchy's theorem

lemma *Cauchy_theorem_homotopic*:

assumes *hom*: if attends then homotopic_paths S g h else homotopic_loops S g h

and *open S* **and** *f*: f holomorphic_on S

and *vpg*: valid_path g **and** *vph*: valid_path h

shows contour_integral g f = contour_integral h f

proof –

have *pathsf*: linked_paths attends g h

using *hom* **by** (auto simp: linked_paths_def homotopic_paths_imp_pathstart homotopic_paths_imp_pathfinish homotopic_loops_imp_loop)

obtain *k* :: real × real ⇒ complex

where *contk*: continuous_on ({0..1} × {0..1}) k

and *ks*: k ' ({0..1} × {0..1}) ⊆ S

and *k* [*simp*]: ∀ x. k (0, x) = g x ∀ x. k (1, x) = h x

and *ksf*: ∀ t ∈ {0..1}. linked_paths attends g (λx. k (t, x))

using *hom pathsf* **by** (auto simp: linked_paths_def homotopic_paths_def homotopic_loops_def homotopic_with_def split: if_split_asm)

have *ucontk*: uniformly_continuous_on ({0..1} × {0..1}) k

by (blast intro: compact_Times compact_uniformly_continuous [OF contk])

{ **fix** *t*::real **assume** *t*: t ∈ {0..1}

have *Pair t* ' {0..1} ⊆ {0..1} × {0..1}

using *t* **by** force

then have *pak*: path (k ∘ (λu. (t, u)))

unfolding *path_def*

by (intro continuous_intros continuous_on_subset [OF contk])+

have *pik*: path_image (k ∘ *Pair t*) ⊆ S

using *ks t* **by** (auto simp: path_image_def)

obtain *e* **where** *e*>0 **and** *e*:

∧ g h. [valid_path g; valid_path h;

∀ u ∈ {0..1}. cmod (g u - (k ∘ *Pair t*) u) < e ∧ cmod (h u - (k ∘

```

Pair t) u) < e;
  linked_paths attends g h]]
  ==> contour_integral h f = contour_integral g f
  using contour_integral_nearby [OF ‹open S› pak pik, of attends] f by metis
  obtain d where d>0 and d:
    ‹ $\bigwedge x x'. [x \in \{0..1\} \times \{0..1\}; x' \in \{0..1\} \times \{0..1\}; \text{norm } (x'-x) < d] \implies$ 
    norm (k x' - k x) < e/4›
    by (rule uniformly_continuous_onE [OF ucontk, of e/4]) (auto simp: dist_norm
    ‹e>0›)
    { fix t1 t2
      assume t1: 0 ≤ t1 t1 ≤ 1 and t2: 0 ≤ t2 t2 ≤ 1 and ltd: |t1 - t| < d |t2
      - t| < d
      have no2: norm(g1 - kt) < e if norm(g1 - k1) < e/4 norm(k1 - kt) <
      e/4 for g1 k1 kt :: complex
      proof (rule norm_triangle_half_l)
        show cmod (g1 - k1) < e/2 cmod (kt - k1) < e/2
          using ‹e > 0› that by (auto simp: norm_minus_commute intro: or-
          der_less_trans)
      qed
      have ‹ $\exists d>0. \forall g1 g2. \text{valid\_path } g1 \wedge \text{valid\_path } g2 \wedge$ 
      ‹ $(\forall u \in \{0..1\}. \text{cmod } (g1 u - k (t1, u)) < d \wedge \text{cmod } (g2 u - k$ 
      (t2, u)) < d) \wedge
      linked_paths attends g1 g2 ‹ $\implies$ 
      contour_integral g2 f = contour_integral g1 f›
      using t t1 t2 ltd ‹e > 0›
      by (rule_tac x=e/4 in exI) (auto intro!: e simp: d no2 simp del: less_divide_eq_numeral1)
    }
  then have ‹ $\exists e. 0 < e \wedge$ 
  ‹ $(\forall t1 t2. t1 \in \{0..1\} \wedge t2 \in \{0..1\} \wedge |t1 - t| < e \wedge |t2 - t| < e$ 
  ‹ $\implies (\exists d. 0 < d \wedge$ 
  ‹ $(\forall g1 g2. \text{valid\_path } g1 \wedge \text{valid\_path } g2 \wedge$ 
  ‹ $(\forall u \in \{0..1\}. \text{norm}(g1 u - k((t1,u))) < d \wedge \text{norm}(g2 u - k((t2,u))) < d) \wedge$ 
  linked_paths attends g1 g2
  ‹ $\implies \text{contour\_integral } g2 f = \text{contour\_integral } g1 f)$ ›››
  by (rule_tac x=d in exI) (simp add: ‹d > 0›)
}
  then obtain ee where ee:
    ‹ $\bigwedge t. t \in \{0..1\} \implies ee t > 0 \wedge$ 
    ‹ $(\forall t1 t2. t1 \in \{0..1\} \implies t2 \in \{0..1\} \implies |t1 - t| < ee t \implies |t2 - t| <$ 
    ee t
    ‹ $\implies (\exists d. 0 < d \wedge$ 
    ‹ $(\forall g1 g2. \text{valid\_path } g1 \wedge \text{valid\_path } g2 \wedge$ 
    ‹ $(\forall u \in \{0..1\}. \text{norm}(g1 u - k((t1,u))) < d \wedge \text{norm}(g2 u - k((t2,u))) < d) \wedge$ 
    linked_paths attends g1 g2
    ‹ $\implies \text{contour\_integral } g2 f = \text{contour\_integral } g1 f)$ ›››
    by metis
  note ee_rule = ee [THEN conjunct2, rule_format, of 0 0 0]

```



```

define  $C$  where  $C = (\lambda t. \text{ball } t \text{ (} ee \ t \ / \ 3)) \text{ ' } \{0..1\}$ 
obtain  $C'$  where  $C'$ :  $C' \subseteq C$  finite  $C'$  and  $C'01$ :  $\{0..1\} \subseteq \bigcup C'$ 
proof (rule compactE [OF compact_interval])
  show  $\{0..1\} \subseteq \bigcup C$ 
    using  $ee$  [THEN conjunct1] by (auto simp: C_def dist_norm)
qed (use C_def in auto)
define  $kk$  where  $kk = \{t \in \{0..1\}. \text{ball } t \text{ (} ee \ t \ / \ 3) \in C'\}$ 
have  $kk01$ :  $kk \subseteq \{0..1\}$  by (auto simp: kk_def)
define  $e$  where  $e = \text{Min } (ee \text{ ' } kk)$ 
have  $C'_eq$ :  $C' = (\lambda t. \text{ball } t \text{ (} ee \ t \ / \ 3)) \text{ ' } kk$ 
  using  $C'$  by (auto simp: kk_def C_def)
have  $ee\_pos$ [simp]:  $\bigwedge t. t \in \{0..1\} \implies ee \ t > 0$ 
  by (simp add: kk_def ee)
moreover have finite  $kk$ 
  using  $\langle \text{finite } C' \rangle$   $kk01$  by (force simp: C'_eq inj_on_def ball_eq_ball_iff dest: ee_pos finite_imageD)
moreover have  $kk \neq \{\}$  using  $\langle \{0..1\} \subseteq \bigcup C' \rangle$   $C'_eq$  by force
ultimately have  $e > 0$ 
  using finite_less Inf_iff [of ee ' kk 0]  $kk01$  by (force simp: e_def)
then obtain  $N::\text{nat}$  where  $N > 0$  and  $N$ :  $1/N < e/3$ 
  by (meson divide_pos_pos nat_approx_posE zero_less_Suc zero_less_numeral)
have  $e\_le\_ee$ :  $\bigwedge i. i \in kk \implies e \leq ee \ i$ 
  using  $\langle \text{finite } kk \rangle$  by (simp add: e_def Min_le_iff [of ee ' kk])
have  $plus$ :  $\exists t \in kk. x \in \text{ball } t \text{ (} ee \ t \ / \ 3)$  if  $x \in \{0..1\}$  for  $x$ 
  using  $C'$  subsetD [OF C'01 that] unfolding  $C'_eq$  by blast
have [OF order_refl]:
   $\exists d. 0 < d \wedge (\forall j. \text{valid\_path } j \wedge (\forall u \in \{0..1\}. \text{norm}(j \ u - k \ (n/N, \ u)) <$ 
d)  $\wedge \text{linked\_paths attends } g \ j$ 
   $\longrightarrow \text{contour\_integral } j \ f = \text{contour\_integral } g \ f)$ 
  if  $n \leq N$  for  $n$ 
using that
proof (induct n)
  case 0 show ?case
    using  $ee\_rule$ 
    by clarsimp (metis diff_self norm_eq_zero vpg)
next
  case (Suc n)
  then have  $N01$ :  $n/N \in \{0..1\}$   $(\text{Suc } n)/N \in \{0..1\}$  by auto
  then obtain  $t$  where  $t$ :  $t \in kk$   $n/N \in \text{ball } t \text{ (} ee \ t \ / \ 3)$ 
    using  $plus$  [of n/N] by blast
  then have  $nN\_less$ :  $|n/N - t| < ee \ t$ 
    by (simp add: dist_norm del: less_divide_eq_numeral1)
  have  $n'N\_less$ :  $|\text{real } (\text{Suc } n) / \text{real } N - t| < ee \ t$ 
    using  $t \ N \ \langle N > 0 \rangle$   $e\_le\_ee$  [of t]
  by (simp add: dist_norm add_divide_distrib abs_diff_less_iff del: less_divide_eq_numeral1)
  (simp add: field_simps)
  have  $t01$ :  $t \in \{0..1\}$  using  $\langle kk \subseteq \{0..1\} \rangle$   $\langle t \in kk \rangle$  by blast
  obtain  $d1$  where  $d1 > 0$  and  $d1$ :
     $\bigwedge g1 \ g2. \llbracket \text{valid\_path } g1; \text{valid\_path } g2;$ 

```

```

       $\forall u \in \{0..1\}. \text{cmod } (g1 \ u - k \ (n/N, \ u)) < d1 \wedge \text{cmod } (g2 \ u - k$ 
       $((\text{Suc } n) / N, \ u)) < d1;$ 
       $\llbracket \text{linked\_paths \textit{atends} } g1 \ g2 \rrbracket$ 
       $\implies \text{contour\_integral } g2 \ f = \text{contour\_integral } g1 \ f$ 
      using ee [THEN conjunct2, rule\_format, OF t01 N01 nN\_less n'N\_less] by
fastforce
      have  $n \leq N$  using Suc.prems by auto
      with Suc.hyps
      obtain d2 where  $d2 > 0$ 
      and  $d2: \bigwedge j. \llbracket \text{valid\_path } j; \forall u \in \{0..1\}. \text{cmod } (j \ u - k \ (n/N, \ u)) < d2;$ 
linked\_paths \textit{atends} } g \ j \rrbracket
       $\implies \text{contour\_integral } j \ f = \text{contour\_integral } g \ f$ 
      by auto
      have Pair  $(n / N) \text{ ' } \{0..1\} \subseteq \{0..1\} \times \{0..1\}$ 
      using N01 by auto
      then have continuous\_on  $\{0..1\} (k \circ (\lambda u. (n/N, \ u)))$ 
      by (intro continuous\_intros continuous\_on\_subset [OF contk])
      then have pkn: path  $(\lambda u. k \ (n/N, \ u))$ 
      by (simp add: path\_def)
      have min12:  $\min \ d1 \ d2 > 0$  by (simp add: <0 < d1> <0 < d2>)
      obtain p where polynomial\_function p
      and psf: pathstart p = pathstart  $(\lambda u. k \ (n/N, \ u))$ 
      pathfinish p = pathfinish  $(\lambda u. k \ (n/N, \ u))$ 
      and pk\_le:  $\bigwedge t. t \in \{0..1\} \implies \text{cmod } (p \ t - k \ (n/N, \ t)) < \min \ d1 \ d2$ 
      using path\_approx\_polynomial\_function [OF pkn min12] by blast
      then have vpp: valid\_path p using valid\_path\_polynomial\_function by blast
      have lpa: linked\_paths \textit{atends} } g \ p
      by (metis (mono\_tags, lifting) N01(1) ksf linked\_paths\_def pathfinish\_def
pathstart\_def psf)
      show ?case
      proof (intro exI; safe)
      fix j
      assume valid\_path j linked\_paths \textit{atends} } g \ j
      and  $\forall u \in \{0..1\}. \text{cmod } (j \ u - k \ (\text{real } (\text{Suc } n) / \text{real } N, \ u)) < \min \ d1 \ d2$ 
      then have contour\_integral j f = contour\_integral p f
      using pk\_le N01(1) ksf by (force intro!: vpp d1 simp add: linked\_paths\_def
psf)
      also have  $\dots = \text{contour\_integral } g \ f$ 
      using pk\_le by (force intro!: vpp d2 lpa)
      finally show contour\_integral j f = contour\_integral g f .
      qed (simp add: <0 < d1> <0 < d2>)
      qed
      then obtain d where  $0 < d$ 
       $\bigwedge j. \text{valid\_path } j \wedge (\forall u \in \{0..1\}. \text{norm}(j \ u - k \ (1, \ u)) < d) \wedge$ 
linked\_paths \textit{atends} } g \ j
       $\implies \text{contour\_integral } j \ f = \text{contour\_integral } g \ f$ 
      using  $\langle N > 0 \rangle$  by auto
      then have linked\_paths \textit{atends} } g \ h \implies \text{contour\_integral } h \ f = \text{contour\_integral } g \ f

```

```

  using <N>0> vph by fastforce
  then show ?thesis
  by (simp add: pathsf)
qed

```

proposition *Cauchy_theorem_homotopic_paths:*

```

  assumes hom: homotopic_paths S g h
  and open S and f: f holomorphic_on S
  and vpg: valid_path g and vph: valid_path h
  shows contour_integral g f = contour_integral h f
  using Cauchy_theorem_homotopic [of True S g h] assms by simp

```

proposition *Cauchy_theorem_homotopic_loops:*

```

  assumes hom: homotopic_loops S g h
  and open S and f: f holomorphic_on S
  and vpg: valid_path g and vph: valid_path h
  shows contour_integral g f = contour_integral h f
  using Cauchy_theorem_homotopic [of False S g h] assms by simp

```

lemma *has_contour_integral_newpath:*

```

  [(f has_contour_integral y) h; f contour_integrable_on g; contour_integral g f
= contour_integral h f]
  ⇒ (f has_contour_integral y) g
  using has_contour_integral_integral contour_integral_unique by auto

```

lemma *Cauchy_theorem_null_homotopic:*

```

  [(f holomorphic_on S; open S; valid_path g; homotopic_loops S g (linepath a
a))]
  ⇒ (f has_contour_integral 0) g
  by (metis Cauchy_theorem_homotopic_loops contour_integrable_holomorphic_simple
valid_path_linepath
contour_integral_trivial has_contour_integral_integral homotopic_loops_imp_subset)

```

end

3 Winding numbers

```

theory Winding_Numbers
  imports Cauchy_Integral_Theorem
begin

```

3.1 Definition

```

definition winding_number_prop :: [real ⇒ complex, complex, real, real ⇒ com-
plex, complex] ⇒ bool where
  winding_number_prop  $\gamma$  z e p n ≡
  valid_path p ∧ z ∉ path_image p ∧
  pathstart p = pathstart  $\gamma$  ∧
  pathfinish p = pathfinish  $\gamma$  ∧

```

$$(\forall t \in \{0..1\}. \text{norm}(\gamma t - p t) < e) \wedge$$

$$\text{contour_integral } p (\lambda w. 1/(w - z)) = 2 * \text{pi} * i * n$$

definition *winding_number*:: $[\text{real} \Rightarrow \text{complex}, \text{complex}] \Rightarrow \text{complex}$ **where**
winding_number $\gamma z \equiv \text{SOME } n. \forall e > 0. \exists p. \text{winding_number_prop } \gamma z e p n$

lemma *winding_number*:

assumes *path* $\gamma z \notin \text{path_image } \gamma 0 < e$

shows $\exists p. \text{winding_number_prop } \gamma z e p (\text{winding_number } \gamma z)$

proof –

have *path_image* $\gamma \subseteq \text{UNIV} - \{z\}$

using *assms* **by** *blast*

then obtain *d*

where *d*: $d > 0$

and *pi_eq*: $\bigwedge h1 h2. \text{valid_path } h1 \wedge \text{valid_path } h2 \wedge$

$$(\forall t \in \{0..1\}. \text{cmod } (h1 t - \gamma t) < d \wedge \text{cmod } (h2 t - \gamma t) < d) \wedge$$

$$\text{pathstart } h2 = \text{pathstart } h1 \wedge \text{pathfinish } h2 = \text{pathfinish } h1 \longrightarrow$$

$$\text{path_image } h1 \subseteq \text{UNIV} - \{z\} \wedge \text{path_image } h2 \subseteq \text{UNIV} -$$

$$\{z\} \wedge$$

$$(\forall f. f \text{ holomorphic_on } \text{UNIV} - \{z\} \longrightarrow \text{contour_integral } h2 f$$

$$= \text{contour_integral } h1 f)$$

using *contour_integral_nearby_ends* [of $\text{UNIV} - \{z\}$ γ] *assms* **by** (*auto simp: open_delete*)

then obtain *h* **where** *h*: *polynomial_function* *h* $\wedge \text{pathstart } h = \text{pathstart } \gamma \wedge$
pathfinish *h* = *pathfinish* $\gamma \wedge$

$$(\forall t \in \{0..1\}. \text{norm}(h t - \gamma t) < d/2)$$

using *path_approx_polynomial_function* [OF $\langle \text{path } \gamma \rangle$, of $d/2$] *d* **by** (*metis half_gt_zero_iff*)

define *nn* **where** $nn = 1/(2 * \text{pi} * i) * \text{contour_integral } h (\lambda w. 1/(w - z))$

have $\exists n. \forall e > 0. \exists p. \text{winding_number_prop } \gamma z e p n$

proof (*rule_tac* $x=nn$ **in** *exI*, *clarify*)

fix *e*::*real*

assume *e*: $e > 0$

obtain *p* **where** *p*: *polynomial_function* *p* \wedge

$$\text{pathstart } p = \text{pathstart } \gamma \wedge \text{pathfinish } p = \text{pathfinish } \gamma \wedge (\forall t \in \{0..1\}. \text{cmod } (p t - \gamma t) < \min e (d/2))$$

cmod ($p t - \gamma t$) $< \min e (d/2)$)

using *path_approx_polynomial_function* [OF $\langle \text{path } \gamma \rangle$, of $\min e (d/2)$] *d* $\langle 0 < e \rangle$

by (*metis min_less_iff_conj zero_less_divide_iff zero_less_numeral*)

have $(\lambda w. 1 / (w - z)) \text{ holomorphic_on } \text{UNIV} - \{z\}$

by (*auto simp: intro!: holomorphic_intros*)

then have *winding_number_prop* $\gamma z e p nn$

using *pi_eq* [of *h* *p*] *h* *p* *d*

by (*auto simp: valid_path_polynomial_function norm_minus_commute nn_def winding_number_prop_def*)

then show $\exists p. \text{winding_number_prop } \gamma z e p nn$

by *metis*

qed

```

then show ?thesis
  unfolding winding_number_def by (rule someI2_ex) (blast intro: ‹0 < e›)
qed

lemma winding_number_unique:
  assumes  $\gamma$ : path  $\gamma$   $z \notin \text{path\_image } \gamma$ 
  and  $\pi$ :  $\bigwedge e. e > 0 \implies \exists p. \text{winding\_number\_prop } \gamma z e p n$ 
  shows winding_number  $\gamma z = n$ 
proof -
  have path_image  $\gamma \subseteq \text{UNIV} - \{z\}$ 
  using assms by blast
  then obtain e
  where e:  $e > 0$ 
  and  $\pi$ _eq:  $\bigwedge h1 h2 f. [\text{valid\_path } h1; \text{valid\_path } h2;$ 
     $(\forall t \in \{0..1\}. \text{cmod } (h1 t - \gamma t) < e \wedge \text{cmod } (h2 t - \gamma t) < e);$ 
     $\text{pathstart } h2 = \text{pathstart } h1; \text{pathfinish } h2 = \text{pathfinish } h1; f$ 
     $\text{holomorphic\_on } \text{UNIV} - \{z\}] \implies$ 
     $\text{contour\_integral } h2 f = \text{contour\_integral } h1 f$ 
  using contour_integral_nearby_ends [of UNIV - {z}  $\gamma$ ] assms by (auto simp:
  open_delete)
  obtain p where p: winding_number_prop  $\gamma z e p n$ 
  using  $\pi$  [OF e] by blast
  obtain q where q: winding_number_prop  $\gamma z e q (\text{winding\_number } \gamma z)$ 
  using winding_number [OF  $\gamma$  e] by blast
  have  $2 * \text{complex\_of\_real } \pi * i * n = \text{contour\_integral } p (\lambda w. 1 / (w - z))$ 
  using p by (auto simp: winding_number_prop_def)
  also have ... = contour_integral q  $(\lambda w. 1 / (w - z))$ 
  proof (rule  $\pi$ _eq)
  show  $(\lambda w. 1 / (w - z)) \text{holomorphic\_on } \text{UNIV} - \{z\}$ 
  by (auto intro!: holomorphic_intros)
qed (use p q in ‹auto simp: winding_number_prop_def norm_minus_commute›)
  also have ... =  $2 * \text{complex\_of\_real } \pi * i * \text{winding\_number } \gamma z$ 
  using q by (auto simp: winding_number_prop_def)
  finally have  $2 * \text{complex\_of\_real } \pi * i * n = 2 * \text{complex\_of\_real } \pi * i * \text{winding\_number } \gamma z$  .
  then show ?thesis
  by simp
qed

```

```

lemma winding_number_unique_loop:
  assumes  $\gamma$ : path  $\gamma$   $z \notin \text{path\_image } \gamma$ 
  and loop: pathfinish  $\gamma = \text{pathstart } \gamma$ 
  and  $\pi$ :
     $\bigwedge e. e > 0 \implies \exists p. \text{valid\_path } p \wedge z \notin \text{path\_image } p \wedge$ 
     $\text{pathfinish } p = \text{pathstart } p \wedge$ 
     $(\forall t \in \{0..1\}. \text{norm } (\gamma t - p t) < e) \wedge$ 
     $\text{contour\_integral } p (\lambda w. 1 / (w - z)) = 2 * \pi * i * n$ 
  shows winding_number  $\gamma z = n$ 

```

proof –
have $\text{path_image } \gamma \subseteq \text{UNIV} - \{z\}$
using assms **by** blast
then obtain e
where $e: e > 0$
and $\text{pi_eq}: \bigwedge h1\ h2\ f. [\text{valid_path } h1; \text{valid_path } h2;$
 $(\forall t \in \{0..1\}. \text{cmod } (h1\ t - \gamma\ t) < e \wedge \text{cmod } (h2\ t - \gamma\ t) < e);$
 $\text{pathfinish } h1 = \text{pathstart } h1; \text{pathfinish } h2 = \text{pathstart } h2; f$
 $\text{holomorphic_on } \text{UNIV} - \{z\}] \implies$
 $\text{contour_integral } h2\ f = \text{contour_integral } h1\ f$
using $\text{contour_integral_nearby_loops}$ [of $\text{UNIV} - \{z\}$ γ] assms **by** (auto simp:
 open_delete)
obtain p **where** p :
 $\text{valid_path } p \wedge z \notin \text{path_image } p \wedge \text{pathfinish } p = \text{pathstart } p \wedge$
 $(\forall t \in \{0..1\}. \text{norm } (\gamma\ t - p\ t) < e) \wedge$
 $\text{contour_integral } p (\lambda w. 1/(w - z)) = 2 * \text{pi} * i * n$
using pi [OF e] **by** blast
obtain q **where** q : $\text{winding_number_prop } \gamma\ z\ e\ q$ ($\text{winding_number } \gamma\ z$)
using winding_number [OF $\gamma\ e$] **by** blast
have $2 * \text{complex_of_real } \text{pi} * i * n = \text{contour_integral } p (\lambda w. 1 / (w - z))$
using p **by** auto
also have $\dots = \text{contour_integral } q (\lambda w. 1 / (w - z))$
proof (rule pi_eq)
show $(\lambda w. 1 / (w - z)) \text{ holomorphic_on } \text{UNIV} - \{z\}$
by ($\text{auto intro!: holomorphic_intros}$)
qed ($\text{use } p\ q\ \text{loop in } \langle \text{auto simp: winding_number_prop_def norm_minus_commute} \rangle$)
also have $\dots = 2 * \text{complex_of_real } \text{pi} * i * \text{winding_number } \gamma\ z$
using q **by** ($\text{auto simp: winding_number_prop_def}$)
finally have $2 * \text{complex_of_real } \text{pi} * i * n = 2 * \text{complex_of_real } \text{pi} * i *$
 $\text{winding_number } \gamma\ z .$
then show $?thesis$
by simp
qed

proposition $\text{winding_number_valid_path}$:
assumes $\text{valid_path } \gamma\ z \notin \text{path_image } \gamma$
shows $\text{winding_number } \gamma\ z = 1 / (2 * \text{pi} * i) * \text{contour_integral } \gamma (\lambda w. 1 / (w - z))$
by ($\text{rule winding_number_unique}$)
 $(\text{use } \text{assms in } \langle \text{auto simp: valid_path_imp_path winding_number_prop_def} \rangle)$

proposition $\text{has_contour_integral_winding_number}$:
assumes $\gamma: \text{valid_path } \gamma\ z \notin \text{path_image } \gamma$
shows $((\lambda w. 1 / (w - z)) \text{ has_contour_integral } (2 * \text{pi} * i * \text{winding_number } \gamma\ z))$
 γ
by ($\text{simp add: winding_number_valid_path has_contour_integral_integral con-}$
 $\text{tour_integrable_inversediff } \text{assms}$)

lemma $\text{winding_number_trivial}$ [simp]: $z \neq a \implies \text{winding_number}(\text{linepath } a\ a)$
 $z = 0$

by (simp add: winding_number_valid_path)

lemma *winding_number_subpath_trivial* [simp]: $z \neq g \ x \implies \text{winding_number}(\text{subpath } x \ x \ g) \ z = 0$

by (simp add: path_image_subpath winding_number_valid_path)

lemma *winding_number_join*:

assumes $\gamma 1: \text{path } \gamma 1 \ z \notin \text{path_image } \gamma 1$

and $\gamma 2: \text{path } \gamma 2 \ z \notin \text{path_image } \gamma 2$

and $\text{pathfinish } \gamma 1 = \text{pathstart } \gamma 2$

shows $\text{winding_number}(\gamma 1 \ +++ \ \gamma 2) \ z = \text{winding_number } \gamma 1 \ z + \text{winding_number } \gamma 2 \ z$

proof (rule *winding_number_unique*)

show $\exists p. \text{winding_number_prop } (\gamma 1 \ +++ \ \gamma 2) \ z \ e \ p$

($\text{winding_number } \gamma 1 \ z + \text{winding_number } \gamma 2 \ z$) if $e > 0$ for e

proof –

obtain $p 1$ where $\text{winding_number_prop } \gamma 1 \ z \ e \ p 1$ ($\text{winding_number } \gamma 1 \ z$)

using $\langle 0 < e \rangle \ \gamma 1 \ \text{winding_number}$ by blast

moreover

obtain $p 2$ where $\text{winding_number_prop } \gamma 2 \ z \ e \ p 2$ ($\text{winding_number } \gamma 2 \ z$)

using $\langle 0 < e \rangle \ \gamma 2 \ \text{winding_number}$ by blast

ultimately

have $\text{winding_number_prop } (\gamma 1 \ +++ \ \gamma 2) \ z \ e \ (p 1 \ +++ \ p 2)$ ($\text{winding_number } \gamma 1 \ z + \text{winding_number } \gamma 2 \ z$)

using *assms*

apply (simp add: *winding_number_prop_def not_in_path_image_join contour_integrable_inversediff algebra_simps*)

apply (auto simp: *joinpaths_def*)

done

then show ?thesis

by blast

qed

qed (use *assms* in $\langle \text{auto simp: not_in_path_image_join} \rangle$)

lemma *winding_number_reversepath*:

assumes $\text{path } \gamma \ z \notin \text{path_image } \gamma$

shows $\text{winding_number}(\text{reversepath } \gamma) \ z = - (\text{winding_number } \gamma \ z)$

proof (rule *winding_number_unique*)

show $\exists p. \text{winding_number_prop } (\text{reversepath } \gamma) \ z \ e \ p$ ($-\text{winding_number } \gamma \ z$)

if $e > 0$ for e

proof –

obtain p where $\text{winding_number_prop } \gamma \ z \ e \ p$ ($\text{winding_number } \gamma \ z$)

using $\langle 0 < e \rangle \ \text{assms } \text{winding_number}$ by blast

then have $\text{winding_number_prop } (\text{reversepath } \gamma) \ z \ e \ (\text{reversepath } p)$ ($-\text{winding_number } \gamma \ z$)

using *assms unfolding winding_number_prop_def*

apply (simp add: *contour_integral_reversepath contour_integrable_inversediff valid_path_imp_reverse*)

apply (auto simp: *reversepath_def*)

```

done
then show ?thesis
  by blast
qed
qed (use assms in auto)

lemma winding_number_shiftpath:
  assumes  $\gamma: \text{path } \gamma \ z \notin \text{path\_image } \gamma$ 
    and  $\text{pathfinish } \gamma = \text{pathstart } \gamma \ a \in \{0..1\}$ 
  shows  $\text{winding\_number}(\text{shiftpath } a \ \gamma) \ z = \text{winding\_number } \gamma \ z$ 
proof (rule winding_number_unique_loop)
  show  $\exists p. \text{valid\_path } p \wedge z \notin \text{path\_image } p \wedge \text{pathfinish } p = \text{pathstart } p \wedge$ 
     $(\forall t \in \{0..1\}. \text{cmod}(\text{shiftpath } a \ \gamma \ t - p \ t) < e) \wedge$ 
     $\text{contour\_integral } p (\lambda w. 1 / (w - z)) =$ 
     $2 * \pi * i * \text{winding\_number } \gamma \ z$ 
  if  $e > 0$  for  $e$ 
proof -
  obtain  $p$  where  $\text{winding\_number\_prop } \gamma \ z \ e \ p$  ( $\text{winding\_number } \gamma \ z$ )
    using  $\langle 0 < e \rangle$  assms  $\text{winding\_number}$  by blast
  then show ?thesis
    apply (rule_tac  $x = \text{shiftpath } a \ p$  in  $\text{exI}$ )
    using assms that
    apply (auto simp:  $\text{winding\_number\_prop\_def}$   $\text{path\_image\_shiftpath}$   $\text{pathfinish\_shiftpath}$   $\text{pathstart\_shiftpath}$   $\text{contour\_integral\_shiftpath}$ )
    apply (simp add:  $\text{shiftpath\_def}$ )
    done
qed
qed (use assms in  $\langle \text{auto simp: path\_shiftpath path\_image\_shiftpath pathfinish\_shiftpath pathstart\_shiftpath} \rangle$ )

lemma winding_number_split_linepath:
  assumes  $c \in \text{closed\_segment } a \ b \ z \notin \text{closed\_segment } a \ b$ 
  shows  $\text{winding\_number}(\text{linepath } a \ b) \ z = \text{winding\_number}(\text{linepath } a \ c) \ z +$ 
 $\text{winding\_number}(\text{linepath } c \ b) \ z$ 
proof -
  have  $z \notin \text{closed\_segment } a \ c \ z \notin \text{closed\_segment } c \ b$ 
  using assms by ( $\text{meson convex\_contains\_segment convex\_segment\_ends\_in\_segment subsetCE}$ )+
  then show ?thesis
    using assms
    by (simp add:  $\text{winding\_number\_valid\_path}$   $\text{contour\_integral\_split\_linepath}$ 
     $[\text{symmetric}]$   $\text{continuous\_on\_inversediff}$   $\text{field\_simps}$ )
qed

lemma winding_number_cong:
   $(\bigwedge t. \llbracket 0 \leq t; t \leq 1 \rrbracket \implies p \ t = q \ t) \implies \text{winding\_number } p \ z = \text{winding\_number } q \ z$ 
  by (simp add:  $\text{winding\_number\_def}$   $\text{winding\_number\_prop\_def}$   $\text{pathstart\_def}$   $\text{pathfinish\_def}$ )

```


lemma *winding_number_constI*:

assumes $c \neq z$ **and** $g: \bigwedge t. [0 \leq t; t \leq 1] \implies g\ t = c$
shows $winding_number\ g\ z = 0$

proof –

have $winding_number\ g\ z = winding_number\ (linepath\ c\ c)\ z$

using $g\ winding_number_cong$ **by** *fastforce*

moreover have $winding_number\ (linepath\ c\ c)\ z = 0$

using $\langle c \neq z \rangle$ **by** *auto*

ultimately show *?thesis* **by** *auto*

qed

lemma *winding_number_offset*: $winding_number\ p\ z = winding_number\ (\lambda w. p\ w - z)\ 0$

unfolding *winding_number_def*

proof (*intro ext arg_cong [where f = Eps] arg_cong [where f = All] imp_cong refl, safe*)

fix $n\ e\ g$

assume $0 < e$ **and** $g: winding_number_prop\ p\ z\ e\ g\ n$

then show $\exists r. winding_number_prop\ (\lambda w. p\ w - z)\ 0\ e\ r\ n$

by (*rule_tac x = $\lambda t. g\ t - z$ in exI*)

(*force simp: winding_number_prop_def contour_integral_integral valid_path_def path_defs*)

(*vector_derivative_def has_vector_derivative_diff_const piecewise_C1_differentiable_diff C1_differentiable_imp_piecewise*)

next

fix $n\ e\ g$

assume $0 < e$ **and** $g: winding_number_prop\ (\lambda w. p\ w - z)\ 0\ e\ g\ n$

then have $winding_number_prop\ p\ z\ e\ (\lambda t. g\ t + z)\ n$

apply (*simp add: winding_number_prop_def contour_integral_integral valid_path_def path_defs*)

(*piecewise_C1_differentiable_add vector_derivative_def has_vector_derivative_add_const C1_differentiable_imp_piecewise*)

apply (*force simp: algebra_simps*)

done

then show $\exists r. winding_number_prop\ p\ z\ e\ r\ n$

by *metis*

qed

lemma *winding_number_negatepath*:

assumes $\gamma: valid_path\ \gamma$ **and** $0: 0 \notin path_image\ \gamma$

shows $winding_number(u\ minus \circ \gamma)\ 0 = winding_number\ \gamma\ 0$

proof –

have $(/) \ 1\ contour_integrable_on\ \gamma$

using $0\ \gamma\ contour_integrable_inversediff$ **by** *fastforce*

then have $((\lambda z. 1/z)\ has_contour_integral\ contour_integral\ \gamma\ ((/) \ 1))\ \gamma$

by (*rule has_contour_integral_integral*)

then have $((\lambda z. 1 / - z)\ has_contour_integral - contour_integral\ \gamma\ ((/) \ 1))$

γ

```

    using has_contour_integral_neg by auto
  then have contour_integral (uminus o  $\gamma$ ) ((/) 1) =
    contour_integral  $\gamma$  ((/) 1)
  using  $\gamma$  by (simp add: contour_integral_unique has_contour_integral_negatepath)
  then show ?thesis
  using assms by (simp add: winding_number_valid_path valid_path_negatepath
image_def path_defs)
qed

```

A combined theorem deducing several things piecewise.

lemma *winding_number_join_pos_combined*:

```

  [[valid_path  $\gamma_1$ ;  $z \notin \text{path\_image } \gamma_1$ ;  $0 < \text{Re}(\text{winding\_number } \gamma_1 z)$ ;
  valid_path  $\gamma_2$ ;  $z \notin \text{path\_image } \gamma_2$ ;  $0 < \text{Re}(\text{winding\_number } \gamma_2 z)$ ; pathfinish
 $\gamma_1 = \text{pathstart } \gamma_2$ ]
   $\implies \text{valid\_path}(\gamma_1 +++ \gamma_2) \wedge z \notin \text{path\_image}(\gamma_1 +++ \gamma_2) \wedge 0 <
\text{Re}(\text{winding\_number}(\gamma_1 +++ \gamma_2) z)$ 
  by (simp add: valid_path_join path_image_join winding_number_join valid_path_imp_path)

```

3.1.1 Useful sufficient conditions for the winding number to be positive

lemma *Re_winding_number*:

```

  [[valid_path  $\gamma$ ;  $z \notin \text{path\_image } \gamma$ ]
   $\implies \text{Re}(\text{winding\_number } \gamma z) = \text{Im}(\text{contour\_integral } \gamma (\lambda w. 1/(w - z))) /
(2*\pi)$ 
  by (simp add: winding_number_valid_path field_simps Re_divide_power2_eq_square)

```

lemma *winding_number_pos_le*:

```

  assumes  $\gamma$ : valid_path  $\gamma$   $z \notin \text{path\_image } \gamma$ 
  and  $ge$ :  $\bigwedge x. [0 < x; x < 1] \implies 0 \leq \text{Im}(\text{vector\_derivative } \gamma (\text{at } x) * \text{cnj}(\gamma
x - z))$ 
  shows  $0 \leq \text{Re}(\text{winding\_number } \gamma z)$ 

```

proof –

```

  have  $ge0$ :  $0 \leq \text{Im}(\text{vector\_derivative } \gamma (\text{at } x) / (\gamma x - z))$  if  $x$ :  $0 < x < 1$ 
for  $x$ 

```

```

  using  $ge$  by (simp add: Complex.Im_divide algebra_simps  $x$ )

```

```

  let  $?vd = \lambda x. 1 / (\gamma x - z) * \text{vector\_derivative } \gamma (\text{at } x)$ 

```

```

  let  $?int = \lambda z. \text{contour\_integral } \gamma (\lambda w. 1 / (w - z))$ 

```

```

  have  $0 \leq \text{Im} (?int z)$ 

```

```

proof (rule has_integral_component_nonneg [of  $i$ , simplified])

```

```

  show  $\bigwedge x. x \in \text{cbox } 0 1 \implies 0 \leq \text{Im}(\text{if } 0 < x \wedge x < 1 \text{ then } ?vd x \text{ else } 0)$ 

```

```

  by (force simp:  $ge0$ )

```

```

  have  $((\lambda a. 1 / (a - z)) \text{ has\_contour\_integral } \text{contour\_integral } \gamma (\lambda w. 1 / (w
- z))) \gamma$ 

```

```

  using  $\gamma$  by (simp flip: add: contour_integrable_inversediff has_contour_integral_integral)

```

```

  then have  $hi$ :  $(?vd \text{ has\_integral } ?int z) (\text{cbox } 0 1)$ 

```

```

  using has_contour_integral by auto

```

```

show  $((\lambda x. \text{if } 0 < x \wedge x < 1 \text{ then } ?vd x \text{ else } 0) \text{ has\_integral } ?int z) (\text{cbox } 0 1)$ 

```

```

  by (rule has_integral_spike_interior [OF  $hi$ ]) simp

```

```

qed
then show ?thesis
  by (simp add: Re_winding_number [OF  $\gamma$ ] field_simps)
qed

lemma winding_number_pos_lt_lemma:
  assumes  $\gamma$ : valid_path  $\gamma$   $z \notin \text{path\_image } \gamma$ 
    and  $e$ :  $0 < e$ 
    and  $ge$ :  $\bigwedge x. [0 < x; x < 1] \implies e \leq \text{Im } (\text{vector\_derivative } \gamma \text{ (at } x) / (\gamma x - z))$ 
  shows  $0 < \text{Re}(\text{winding\_number } \gamma z)$ 
proof -
  let ?vd =  $\lambda x. 1 / (\gamma x - z) * \text{vector\_derivative } \gamma \text{ (at } x)$ 
  let ?int =  $\lambda z. \text{contour\_integral } \gamma (\lambda w. 1 / (w - z))$ 
  have  $e \leq \text{Im } (\text{contour\_integral } \gamma (\lambda w. 1 / (w - z)))$ 
  proof (rule has_integral_component_le [of  $i \lambda x. i * e \ i * e \ \{0..1\}$ , simplified])
    have  $((\lambda a. 1 / (a - z)) \text{ has\_contour\_integral } \text{contour\_integral } \gamma (\lambda w. 1 / (w - z))) \ \gamma$ 
    thm has_integral_component_le [of  $i \lambda x. i * e \ i * e \ \{0..1\}$ , simplified]
    using  $\gamma$  by (simp add: contour_integrable_inversediff has_contour_integral_integral)
    then have  $hi$ :  $(?vd \text{ has\_integral } ?int z) \text{ (cbox } 0 \ 1)$ 
    using has_contour_integral by auto
    show  $((\lambda x. \text{if } 0 < x \wedge x < 1 \text{ then } ?vd \ x \ \text{else } i * e) \text{ has\_integral } ?int z) \ \{0..1\}$ 
    by (rule has_integral_spike_interior [OF  $hi$ , simplified box_real]) (use  $e$  in simp)
  show  $\bigwedge x. 0 \leq x \wedge x \leq 1 \implies e \leq \text{Im } (\text{if } 0 < x \wedge x < 1 \text{ then } ?vd \ x \ \text{else } i * e)$ 
    by (simp add:  $ge$ )
  qed
  qed (use has_integral_const_real [of  $\_ \ 0 \ 1$ ] in auto)
  with  $e$  show ?thesis
  by (simp add: Re_winding_number [OF  $\gamma$ ] field_simps)
qed

```

```

lemma winding_number_pos_lt:
  assumes  $\gamma$ : valid_path  $\gamma$   $z \notin \text{path\_image } \gamma$ 
    and  $e$ :  $0 < e$ 
    and  $ge$ :  $\bigwedge x. [0 < x; x < 1] \implies e \leq \text{Im } (\text{vector\_derivative } \gamma \text{ (at } x) * \text{cnj}(\gamma x - z))$ 
  shows  $0 < \text{Re } (\text{winding\_number } \gamma z)$ 
proof -
  have  $bm$ : bounded  $((\lambda w. w - z) \text{ ' (path\_image } \gamma))$ 
  using bounded_translation [of  $\_ -z$ ]  $\gamma$  by (simp add: bounded_valid_path_image)
  then obtain  $B$  where  $B$ :  $B > 0$  and  $Bno$ :  $\bigwedge x. x \in (\lambda w. w - z) \text{ ' (path\_image } \gamma) \implies \text{norm } x \leq B$ 
  using bounded_pos [THEN iffD1, OF  $bm$ ] by blast
  { fix  $x::\text{real}$  assume  $x$ :  $0 < x < 1$ 
    then have  $B2$ :  $\text{cmod } (\gamma x - z)^2 \leq B^2$  using  $Bno$  [of  $\gamma x - z$ ]
    by (simp add: path_image_def power2_eq_square mult_mono')
    with  $x$  have  $\gamma x \neq z$  using  $\gamma$ 
    using path_image_def by fastforce
  }

```

```

then have  $e / B^2 \leq e / (\text{cmod } (\gamma x - z))^2$ 
  using  $B B^2 e$  by (auto simp: divide_left_mono)
also have  $\dots \leq \text{Im } (\text{vector\_derivative } \gamma \text{ (at } x) * \text{cnj } (\gamma x - z)) / (\text{cmod } (\gamma x - z))^2$ 
  using  $ge [OF x]$  by (auto simp: divide_right_mono)
finally have  $e / B^2 \leq \text{Im } (\text{vector\_derivative } \gamma \text{ (at } x) * \text{cnj } (\gamma x - z)) / (\text{cmod } (\gamma x - z))^2$  .
  then have  $e / B^2 \leq \text{Im } (\text{vector\_derivative } \gamma \text{ (at } x) / (\gamma x - z))$ 
  by (simp add: complex_div_cnj [of _  $\gamma x - z$  for  $x$ ] del: complex_cnj_diff times_complex.sel)
} note  $*$  = this
show ?thesis
  using  $e B$  by (simp add: * winding_number_pos_lt_lemma [OF  $\gamma$ , of  $e/B^2$ ])
qed

```

3.2 The winding number is an integer

Proof from the book Complex Analysis by Lars V. Ahlfors, Chapter 4, section 2.1, Also on page 134 of Serge Lang's book with the name title, etc.

lemma *exp_fg*:

```

fixes  $z::\text{complex}$ 
assumes  $g: (g \text{ has\_vector\_derivative } g') \text{ (at } x \text{ within } s)$ 
  and  $f: (f \text{ has\_vector\_derivative } (g' / (g x - z))) \text{ (at } x \text{ within } s)$ 
  and  $z: g x \neq z$ 
shows  $((\lambda x. \text{exp}(-f x) * (g x - z)) \text{ has\_vector\_derivative } 0) \text{ (at } x \text{ within } s)$ 
proof -
  have  $*$ :  $(\text{exp} \circ (\lambda x. (- f x)) \text{ has\_vector\_derivative } - (g' / (g x - z)) * \text{exp}(- f x)) \text{ (at } x \text{ within } s)$ 
  using assms unfolding has_vector_derivative_def scaleR_conv_of_real
  by (auto intro!: derivative_eq_intros)
show ?thesis
  using  $z$  by (auto intro!: derivative_eq_intros * [unfolded o_def] g)
qed

```

lemma *winding_number_exp_integral*:

```

fixes  $z::\text{complex}$ 
assumes  $\gamma: \gamma \text{ piecewise\_C1\_differentiable\_on } \{a..b\}$ 
  and  $ab: a \leq b$ 
  and  $z: z \notin \gamma ' \{a..b\}$ 
shows  $(\lambda x. \text{vector\_derivative } \gamma \text{ (at } x) / (\gamma x - z)) \text{ integrable\_on } \{a..b\}$ 
  (is ?thesis1)
   $\text{exp}(- (\text{integral } \{a..b\} (\lambda x. \text{vector\_derivative } \gamma \text{ (at } x) / (\gamma x - z)))) * (\gamma b - z) = \gamma a - z$ 
  (is ?thesis2)
proof -
  let  $?D\gamma = \lambda x. \text{vector\_derivative } \gamma \text{ (at } x)$ 
  have [simp]:  $\bigwedge x. a \leq x \implies x \leq b \implies \gamma x \neq z$ 
  using  $z$  by force
  have con_g: continuous_on  $\{a..b\} \gamma$ 

```

```

    using  $\gamma$  by (simp add: piecewise_C1_differentiable_on_def)
  obtain  $k$  where fink: finite  $k$  and  $g\_C1\_diff: \gamma$  C1_differentiable_on  $\{a..b\} - k$ 
  using  $\gamma$  by (force simp: piecewise_C1_differentiable_on_def)
  have  $\circ: open \{a<..b\} - k$ 
    using  $\langle$ finite  $k\rangle$  by (simp add: finite_imp_closed open_Diff)
  moreover have  $\{a<..b\} - k \subseteq \{a..b\} - k$ 
    by force
  ultimately have  $g\_diff\_at: \bigwedge x. \llbracket x \notin k; x \in \{a<..b\} \rrbracket \implies \gamma$  differentiable at  $x$ 
    by (metis Diff_iff differentiable_on_subset C1_diff_imp_diff [OF  $g\_C1\_diff$ ] differentiable_on_def at_within_open)
  { fix  $w$ 
    assume  $w \neq z$ 
    have continuous_on (ball  $w$  (cmod  $(w - z)$ ))  $(\lambda w. 1 / (w - z))$ 
      by (auto simp: dist_norm intro!: continuous_intros)
    moreover have  $\bigwedge x. cmod (w - x) < cmod (w - z) \implies \exists f'. ((\lambda w. 1 / (w - z)))$  has_field_derivative  $f'$  (at  $x$ )
      by (auto simp: intro!: derivative_eq_intros)
    ultimately have  $\exists h. \forall y. norm(y - w) < norm(w - z) \longrightarrow (h$  has_field_derivative  $1/(y - z)$ ) (at  $y$ )
      using holomorphic_convex_primitive [of ball  $w$  (norm  $(w - z)$ )  $\{\}$   $\lambda w. 1/(w - z)$ ]
    by (force simp: field_differentiable_def Ball_def dist_norm at_within_open NO_MATCH norm_minus_commute)
  }
  then obtain  $h$  where  $h: \bigwedge w y. w \neq z \implies norm(y - w) < norm(w - z) \implies (h$  has_field_derivative  $1/(y - z)$ ) (at  $y$ )
    by meson
  have  $exy: \exists y. ((\lambda x. inverse (\gamma x - z) * ?D\gamma x)$  has_integral  $y$ )  $\{a..b\}$ 
    unfolding integrable_on_def [symmetric]
  proof (rule contour_integral_local_primitive_any [OF piecewise_C1_imp_differentiable [OF  $\gamma$ ]])
    show  $\exists d h. 0 < d \wedge (\forall y. cmod (y - w) < d \longrightarrow (h$  has_field_derivative  $inverse (y - z))$ ) (at  $y$  within  $-\{z\}$ )
      if  $w \in -\{z\}$  for  $w$ 
      using that inverse_eq_divide has_field_derivative_at_within  $h$ 
    by (metis Compl_insert DiffD2 insertCI right_minus_eq zero_less_norm_iff)
  qed simp
  have  $vg\_int: (\lambda x. ?D\gamma x / (\gamma x - z))$  integrable_on  $\{a..b\}$ 
    unfolding box_real [symmetric] divide_inverse_commute
  by (auto intro!: exy integrable_subinterval simp add: integrable_on_def ab)
  with  $ab$  show ?thesis1
    by (simp add: divide_inverse_commute integral_def integrable_on_def)
  { fix  $t$ 
    assume  $t: t \in \{a..b\}$ 
    have  $cball: continuous_on$  (ball  $(\gamma t)$  (dist  $(\gamma t) z$ ))  $(\lambda x. inverse (x - z))$ 
      using  $\gamma$  by (auto intro!: continuous_intros simp: dist_norm)

```

```

have icd:  $\bigwedge x. \text{cmod } (\gamma t - x) < \text{cmod } (\gamma t - z) \implies (\lambda w. \text{inverse } (w - z))$ 
field_differentiable at x
unfolding field_differentiable_def by (force simp: intro!: derivative_eq_intros)
obtain h where h:  $\bigwedge x. \text{cmod } (\gamma t - x) < \text{cmod } (\gamma t - z) \implies$ 
 $(h \text{ has\_field\_derivative } \text{inverse } (x - z))$  (at x within  $\{y. \text{cmod } (\gamma t - y) < \text{cmod } (\gamma t - z)\}$ )
using holomorphic_convex_primitive [where f =  $\lambda w. \text{inverse}(w - z)$ , OF
convex_ball finite.emptyI cball icd]
by simp (auto simp: ball_def dist_norm that)
have exp  $(- \text{integral } \{a..t\} (\lambda x. ?D\gamma x / (\gamma x - z))) * (\gamma t - z) = \gamma a - z$ 
proof (rule has_derivative_zero_unique_strong_interval [of  $\{a,b\} \cup k$  a b])
show continuous_on  $\{a..b\}$   $(\lambda b. \text{exp } (- \text{integral } \{a..b\} (\lambda x. ?D\gamma x / (\gamma x - z)))) * (\gamma b - z)$ 
by (auto intro!: continuous_intros con_g indefinite_integral_continuous_1
[OF vg_int])
show  $(\lambda b. \text{exp } (- \text{integral } \{a..b\} (\lambda x. ?D\gamma x / (\gamma x - z)))) * (\gamma b - z)$ 
has_derivative  $(\lambda h. 0)$ 
(at x within  $\{a..b\}$ )
if  $x \in \{a..b\} - (\{a, b\} \cup k)$  for x
proof -
have x:  $x \notin k$   $a < x < b$ 
using that by auto
then have x  $\in$  interior  $(\{a..b\} - k)$ 
using open_subset_interior [OF o] by fastforce
then have con: isCont  $?D\gamma x$ 
using g_C1_diff x by (auto simp: C1_differentiable_on_eq intro: continuous_on_interior)
then have con_vd: continuous (at x within  $\{a..b\}$ )  $(\lambda x. ?D\gamma x)$ 
by (rule continuous_at_imp_continuous_within)
have gdx:  $\gamma$  differentiable at x
using x by (simp add: g_diff_at)
then obtain d where d:  $(\gamma \text{ has\_derivative } (\lambda x. x *_R d))$  (at x)
by (auto simp add: differentiable_iff_scaleR)
show  $(\lambda c. \text{exp } (- \text{integral } \{a..c\} (\lambda x. ?D\gamma x / (\gamma x - z)))) * (\gamma c - z)$ 
has_derivative  $(\lambda h. 0)$ 
(at x within  $\{a..b\}$ )
proof (rule exp_fg [unfolded has_vector_derivative_def, simplified])
show  $(\gamma \text{ has\_derivative } (\lambda c. c *_R d))$  (at x within  $\{a..b\}$ )
using d by (blast intro: has_derivative_at_withinI)
have  $(\lambda x. \text{integral } \{a..x\} (\lambda x. ?D\gamma x / (\gamma x - z))) \text{ has\_vector\_derivative } d / (\gamma x - z)$ 
(at x within  $\{a..b\}$ )
proof (rule has_vector_derivative_eq_rhs [OF integral_has_vector_derivative_continuous_at
[where S =  $\{\}$ , simplified]])
show continuous (at x within  $\{a..b\}$ )  $(\lambda x. \text{vector\_derivative } \gamma$  (at x) /
 $(\gamma x - z))$ 
using continuous_at_imp_continuous_at_within differentiable_imp_continuous_within
gdx x
by (intro con_vd continuous_intros) (force+)

```

```

      show vector_derivative  $\gamma$  (at  $x$ ) / ( $\gamma x - z$ ) =  $d$  / ( $\gamma x - z$ )
      using d vector_derivative_at
      by (simp add: vector_derivative_at has_vector_derivative_def)
    qed (use x vg_int in auto)
  then show (( $\lambda x$ . integral {a..x} ( $\lambda x$ . ?D $\gamma x$  / ( $\gamma x - z$ ))) has_derivative
( $\lambda c$ .  $c *_R (d / (\gamma x - z))$ ))
    (at  $x$  within {a..b})
    by (auto simp: has_vector_derivative_def)
  qed (use x in auto)
qed
qed
}
with ab show ?thesis2
  by (simp add: divide_inverse_commute_integral_def)
qed

```

lemma winding_number_exp_2pi:

```

  [[path  $p$ ;  $z \notin \text{path\_image } p$ ]
   $\implies \text{pathfinish } p - z = \exp (2 * \pi * i * \text{winding\_number } p z) * (\text{pathstart } p - z)$ 
  using winding_number [of  $p z 1$ ] unfolding valid_path_def path_image_def pathstart_def pathfinish_def winding_number_prop_def
  by (force dest: winding_number_exp_integral(2) [of _ 0 1  $z$ ] simp: field_simps contour_integral_integral_exp_minus)

```

lemma integer_winding_number_eq:

```

  assumes  $\gamma$ : path  $\gamma$  and  $z$ :  $z \notin \text{path\_image } \gamma$ 
  shows winding_number  $\gamma z \in \mathbb{Z} \iff \text{pathfinish } \gamma = \text{pathstart } \gamma$ 
  proof -
    obtain  $p$  where  $p$ : valid_path  $p z \notin \text{path\_image } p$ 
      pathstart  $p = \text{pathstart } \gamma$  pathfinish  $p = \text{pathfinish } \gamma$ 
      and eq: contour_integral  $p (\lambda w. 1 / (w - z)) = \text{complex\_of\_real } (2 * \pi) * i * \text{winding\_number } \gamma z$ 
    using winding_number [OF assms, of 1] unfolding winding_number_prop_def
    by auto
    then have wneq: winding_number  $\gamma z = \text{winding\_number } p z$ 
      using eq winding_number_valid_path by force
    have iff: (winding_number  $\gamma z \in \mathbb{Z}$ )  $\iff (\exp (\text{contour\_integral } p (\lambda w. 1 / (w - z))) = 1)$ 
      using eq by (simp add: exp_eq_1 complex_is_Int_iff)
    have  $\gamma 0 \neq z$ 
      by (metis pathstart_def pathstart_in_path_image  $z$ )
    then have exp (contour_integral  $p (\lambda w. 1 / (w - z))) = (\gamma 1 - z) / (\gamma 0 - z)$ 
      using  $p$  winding_number_exp_integral(2) [of  $p 0 1 z$ ]
      by (simp add: valid_path_def path_defs contour_integral_integral_exp_minus field_split_simps)
    then have winding_number  $p z \in \mathbb{Z} \iff \text{pathfinish } p = \text{pathstart } p$ 
      using  $p$  wneq iff by (auto simp: path_defs)
    then show ?thesis using  $p$  eq
  end

```

by (auto simp: winding_number_valid_path)
qed

theorem integer_winding_number:

$\llbracket \text{path } \gamma; \text{pathfinish } \gamma = \text{pathstart } \gamma; z \notin \text{path_image } \gamma \rrbracket \implies \text{winding_number } \gamma$
 $z \in \mathbb{Z}$

by (metis integer_winding_number_eq)

If the winding number's magnitude is at least one, then the path must contain points in every direction.*) We can thus bound the winding number of a path that doesn't intersect a given ray.

lemma winding_number_pos_meets:

fixes $z::\text{complex}$

assumes $\gamma: \text{valid_path } \gamma$ and $z: z \notin \text{path_image } \gamma$ and $1: \text{Re } (\text{winding_number } \gamma z) \geq 1$

and $w: w \neq z$

shows $\exists a::\text{real}. 0 < a \wedge z + \text{of_real } a * (w - z) \in \text{path_image } \gamma$

proof –

have [simp]: $\bigwedge x. 0 \leq x \implies x \leq 1 \implies \gamma x \neq z$

using z by (auto simp: path_image_def)

have [simp]: $z \notin \gamma \text{ ' } \{0..1\}$

using path_image_def z by auto

have gpd: $\gamma \text{ piecewise_C1_differentiable_on } \{0..1\}$

using $\gamma \text{ valid_path_def}$ by blast

define r where $r = (w - z) / (\gamma 0 - z)$

have [simp]: $r \neq 0$

using $w z$ by (auto simp: r_def)

have cont: continuous_on $\{0..1\}$

$(\lambda x. \text{Im } (\text{integral } \{0..x\} (\lambda x. \text{vector_derivative } \gamma (at x) / (\gamma x - z))))$

by (intro continuous_intros indefinite_integral_continuous_1 winding_number_exp_integral [OF gpd]; simp)

have $\text{Arg}2\pi r \leq 2*\pi$

by (simp add: Arg2pi_less_eq_real_def)

also have $\dots \leq \text{Im } (\text{integral } \{0..1\} (\lambda x. \text{vector_derivative } \gamma (at x) / (\gamma x - z)))$

using 1

by (simp add: winding_number_valid_path [OF γz] contour_integral_integral Complex.Re_divide field_simps power2_eq_square)

finally have $\text{Arg}2\pi r \leq \text{Im } (\text{integral } \{0..1\} (\lambda x. \text{vector_derivative } \gamma (at x) / (\gamma x - z)))$.

then have $\exists t. t \in \{0..1\} \wedge \text{Im}(\text{integral } \{0..t\} (\lambda x. \text{vector_derivative } \gamma (at x) / (\gamma x - z))) = \text{Arg}2\pi r$

by (simp add: Arg2pi_ge_0 cont IVT')

then obtain t where $t: t \in \{0..1\}$

and eqArg: $\text{Im } (\text{integral } \{0..t\} (\lambda x. \text{vector_derivative } \gamma (at x) / (\gamma x - z))) = \text{Arg}2\pi r$

by blast

define i where $i = \text{integral } \{0..t\} (\lambda x. \text{vector_derivative } \gamma (at x) / (\gamma x - z))$

have gpd: $\gamma \text{ piecewise_C1_differentiable_on } \{0..t\}$


```

  by (metis atLeastAtMost_iff atLeastatMost_subset_iff order_refl piecewise_C1_differentiable_on_subset
    gpd t)
  have exp (- i) * (γ t - z) = γ 0 - z
    unfolding i_def
  proof (rule winding_number_exp_integral [OF gpd])
    show z ∉ γ ` {0..t}
      using t z unfolding path_image_def by force
  qed (use t in auto)
  then have *: γ t - z = exp i * (γ 0 - z)
    by (simp add: exp_minus_field_simps)
  then have (w - z) = r * (γ 0 - z)
    by (simp add: r_def)
  moreover have z + exp (Re i) * (exp (i * Im i) * (γ 0 - z)) = γ t
    using * by (simp add: exp_eq_polar_field_simps)
  moreover have Arg2pi r = Im i
    using eqArg by (simp add: i_def)
  ultimately have z + complex_of_real (exp (Re i)) * (w - z) / complex_of_real
    (cmod r) = γ t
    using Complex_Transcendental.Arg2pi_eq [of r] ⟨r ≠ 0⟩
    by (metis mult.left_commute nonzero_mult_div_cancel_left norm_eq_zero
      of_real_0_of_real_eq_iff times_divide_eq_left)
  with t show ?thesis
    by (rule_tac x=exp(Re i) / norm r in exI) (auto simp: path_image_def)
  qed

```

lemma winding_number_big_meets:

```

  fixes z::complex
  assumes γ: valid_path γ and z: z ∉ path_image γ and |Re (winding_number
    γ z)| ≥ 1
    and w: w ≠ z
  shows ∃ a::real. 0 < a ∧ z + of_real a * (w - z) ∈ path_image γ
  proof -
    { assume Re (winding_number γ z) ≤ - 1
      then have Re (winding_number (reversepath γ) z) ≥ 1
        by (simp add: γ valid_path_imp_path winding_number_reversepath z)
      moreover have valid_path (reversepath γ)
        using γ valid_path_imp_reverse by auto
      moreover have z ∉ path_image (reversepath γ)
        by (simp add: z)
      ultimately have ∃ a::real. 0 < a ∧ z + of_real a * (w - z) ∈ path_image
        (reversepath γ)
        using winding_number_pos_meets w by blast
      then have ?thesis
        by simp
    }
  then show ?thesis
    using assms
    by (simp add: abs_if_winding_number_pos_meets split: if_split_asm)
  qed

```

```

lemma winding_number_less_1:
  fixes z::complex
  shows
  [[valid_path  $\gamma$ ;  $z \notin \text{path\_image } \gamma$ ;  $w \neq z$ ;
    $\wedge a::\text{real. } 0 < a \implies z + \text{of\_real } a * (w - z) \notin \text{path\_image } \gamma$ ]
   $\implies \text{Re}(\text{winding\_number } \gamma z) < 1$ 
  by (auto simp: not_less dest: winding_number_big_meets)

```

One way of proving that $\text{WN}=1$ for a loop.

```

lemma winding_number_eq_1:
  fixes z::complex
  assumes  $\gamma$ : valid_path  $\gamma$  and  $z$ :  $z \notin \text{path\_image } \gamma$  and loop: pathfinish  $\gamma =$ 
  pathstart  $\gamma$ 
  and 0:  $0 < \text{Re}(\text{winding\_number } \gamma z)$  and 2:  $\text{Re}(\text{winding\_number } \gamma z) < 2$ 
  shows winding_number  $\gamma z = 1$ 
proof -
  have winding_number  $\gamma z \in \text{Ints}$ 
  by (simp add:  $\gamma$  integer_winding_number loop valid_path_imp_path z)
  then show ?thesis
  using 0 2 by (auto simp: Ints_def)
qed

```

3.3 Continuity of winding number and invariance on connected sets

```

theorem continuous_at_winding_number:
  fixes z::complex
  assumes  $\gamma$ : path  $\gamma$  and  $z$ :  $z \notin \text{path\_image } \gamma$ 
  shows continuous (at  $z$ ) (winding_number  $\gamma$ )
proof -
  obtain  $e$  where  $e > 0$  and  $cbg$ : cball  $z e \subseteq - \text{path\_image } \gamma$ 
  using open_contains_cball [of  $- \text{path\_image } \gamma$ ]  $z$ 
  by (force simp: closed_def [symmetric] closed_path_image [OF  $\gamma$ ])
  then have  $ppag$ : path_image  $\gamma \subseteq - \text{cball } z (e/2)$ 
  by (force simp: cball_def dist_norm)
  have  $oc$ : open  $(- \text{cball } z (e/2))$ 
  by (simp add: closed_def [symmetric])
  obtain  $d$  where  $d > 0$  and  $pi\_eq$ :
   $\wedge h1 h2. \llbracket \text{valid\_path } h1; \text{valid\_path } h2;
   (\forall t \in \{0..1\}. \text{cmod } (h1 t - \gamma t) < d \wedge \text{cmod } (h2 t - \gamma t) < d);
   \text{pathstart } h2 = \text{pathstart } h1; \text{pathfinish } h2 = \text{pathfinish } h1 \rrbracket$ 
   $\implies$ 
  path_image  $h1 \subseteq - \text{cball } z (e/2) \wedge$ 
  path_image  $h2 \subseteq - \text{cball } z (e/2) \wedge$ 
   $(\forall f. f \text{ holomorphic\_on } - \text{cball } z (e/2) \longrightarrow \text{contour\_integral } h2 f =$ 
  contour_integral  $h1 f)$ 
  using contour_integral_nearby_ends [OF  $oc \ \gamma \ ppag$ ] by metis
  obtain  $p$  where valid_path  $p \ z \notin \text{path\_image } p$ 

```

```

    and p: pathstart p = pathstart  $\gamma$  pathfinish p = pathfinish  $\gamma$ 
    and pg:  $\bigwedge t. t \in \{0..1\} \implies \text{cmod } (\gamma t - p t) < \min d e/2$ 
    and pi: contour_integral p ( $\lambda x. 1 / (x - z)$ ) = complex_of_real (2 *
pi) * i * winding_number  $\gamma z$ 
    using winding_number [OF  $\gamma z$ , of min d e/2] <d>0> <e>0> by (auto simp:
winding_number_prop_def)
  { fix w
    assume d2: cmod (w - z) < d/2 and e2: cmod (w - z) < e/2
    have wnotp: w  $\notin$  path_image p
    proof (clarsimp simp add: path_image_def)
      show False if w: w = p x and 0  $\leq$  x  $\leq$  1 for x
      proof -
        have cmod ( $\gamma x - p x$ ) < min d e/2
          using pg that by auto
        then have cmod (z -  $\gamma x$ ) < e
          by (metis e2 less_divide_eq_numeral1(1) min_less_iff_conj norm_minus_commute
norm_triangle_half_l w)
        then show ?thesis
          using cbg that by (auto simp add: path_image_def cball_def dist_norm
less_eq_real_def)
      qed
    qed
    have wnotg: w  $\notin$  path_image  $\gamma$ 
    using cbg e2 <e>0> by (force simp: dist_norm norm_minus_commute)
    { fix k::real
      assume k: k > 0
      then obtain q where q: valid_path q w  $\notin$  path_image q
        pathstart q = pathstart  $\gamma$   $\wedge$  pathfinish q = pathfinish  $\gamma$ 
        and qg:  $\bigwedge t. t \in \{0..1\} \implies \text{cmod } (\gamma t - q t) < \min k (\min d e)$ 
/ 2
        and qi: contour_integral q ( $\lambda u. 1 / (u - w)$ ) = complex_of_real
(2 * pi) * i * winding_number  $\gamma w$ 
        using winding_number [OF  $\gamma$  wnotg, of min k (min d e) / 2] <d>0> <e>0>
k
        by (force simp: min_divide_distrib_right winding_number_prop_def)
      moreover have  $\bigwedge t. t \in \{0..1\} \implies \text{cmod } (q t - \gamma t) < d \wedge \text{cmod } (p t - \gamma
t) < d$ 
        using pg qg <0 < d> by (fastforce simp add: norm_minus_commute)
      moreover have ( $\lambda u. 1 / (u - w)$ ) holomorphic_on - cball z (e/2)
        using e2 by (auto simp: dist_norm norm_minus_commute intro!: holo-
morphic_intros)
      ultimately have contour_integral p ( $\lambda u. 1 / (u - w)$ ) = contour_integral q
( $\lambda u. 1 / (u - w)$ )
        by (metis p <valid_path p> pi_eq)
      then have contour_integral p ( $\lambda x. 1 / (x - w)$ ) = complex_of_real (2 * pi)
* i * winding_number  $\gamma w$ 
        by (simp add: pi qi)
    } note pip = this
    have path p

```

```

    by (simp add: ‹valid_path p› valid_path_imp_path)
  moreover have  $\bigwedge e. e > 0 \implies \text{winding\_number\_prop } p \ w \ e \ p \ (\text{winding\_number } \gamma \ w)$ 
    by (simp add: ‹valid_path p› pip_winding_number_prop_def wnotp)
  ultimately have  $\text{winding\_number } p \ w = \text{winding\_number } \gamma \ w$ 
    using winding_number_unique wnotp by blast
} note wnw = this
obtain pe where  $pe > 0$  and cbp:  $\text{cball } z \ (3 / 4 * pe) \subseteq - \text{path\_image } p$ 
using ‹valid_path p› ‹ $z \notin \text{path\_image } p$ › open_contains_cball [of - path_image p]
by (force simp: closed_def [symmetric] closed_path_image [OF valid_path_imp_path])
obtain L
where  $L > 0$ 
and L:  $\bigwedge f \ B. \llbracket f \text{ holomorphic\_on } - \text{cball } z \ (3 / 4 * pe); \forall z \in - \text{cball } z \ (3 / 4 * pe). \text{cmod } (f \ z) \leq B \rrbracket \implies \text{cmod } (\text{contour\_integral } p \ f) \leq L * B$ 
using contour_integral_bound_exists [of - cball z (3/4*pe) p] cbp ‹valid_path p› by blast
{ fix e::real and w::complex
  assume  $e: 0 < e$  and  $w: \text{cmod } (w - z) < pe/4$   $\text{cmod } (w - z) < e * pe^2 / (8 * L)$ 
  then have [simp]:  $w \notin \text{path\_image } p$ 
    using cbp p(2) ‹ $0 < pe$ ›
  by (force simp: dist_norm norm_minus_commute path_image_def cball_def)
  have [simp]:  $\text{contour\_integral } p \ (\lambda x. 1/(x - w)) - \text{contour\_integral } p \ (\lambda x. 1/(x - z)) =$ 
 $\text{contour\_integral } p \ (\lambda x. 1/(x - w) - 1/(x - z))$ 
  by (simp add: ‹valid_path p› ‹ $z \notin \text{path\_image } p$ › contour_integrable_inversediff contour_integral_diff)
  { fix x
    assume  $pe: 3/4 * pe < \text{cmod } (z - x)$ 
    have  $\text{cmod } (w - x) < pe/4 + \text{cmod } (z - x)$ 
    by (meson add_less_cancel_right norm_diff_triangle_le order_refl order_trans_rules(21) w(1))
    then have  $wx: \text{cmod } (w - x) < 4/3 * \text{cmod } (z - x)$  using pe by simp
    have  $\text{cmod } (z - x) \leq \text{cmod } (z - w) + \text{cmod } (w - x)$ 
    using norm_diff_triangle_le by blast
    also have  $\dots < pe/4 + \text{cmod } (w - x)$ 
    using w by (simp add: norm_minus_commute)
    finally have  $pe/2 < \text{cmod } (w - x)$ 
    using pe by auto
    then have  $pe\_less: (pe/2)^2 < \text{cmod } (w - x)^2$ 
    by (simp add: ‹ $0 < pe$ › less_eq_real_def power_strict_mono)
    then have  $pe2: pe^2 < 4 * \text{cmod } (w - x)^2$ 
    by (simp add: power_divide)
    have  $8 * L * \text{cmod } (w - z) < e * pe^2$ 
    using w ‹ $L > 0$ › by (simp add: field_simps)
    also have  $\dots < e * 4 * \text{cmod } (w - x) * \text{cmod } (w - x)$ 
    using pe2 ‹ $e > 0$ › by (simp add: power2_eq_square)
  }
}

```

```

also have ... < e * 4 * cmod (w - x) * (4/3 * cmod (z - x))
  using <0 < pe> pe_less e_less_eq_real_def wx by fastforce
finally have L * cmod (w - z) < 2/3 * e * cmod (w - x) * cmod (z - x)
  by simp
also have ... ≤ e * cmod (w - x) * cmod (z - x)
  using e by simp
finally have Lwz: L * cmod (w - z) < e * cmod (w - x) * cmod (z - x) .
have L * cmod (1 / (x - w) - 1 / (x - z)) ≤ e
proof (cases x=z ∨ x=w)
  case True
  with pe <pe>0> w <L>0>
  show ?thesis
    by (force simp: norm_minus_commute)
  next
  case False
  with wx w(2) <L>0> pe pe2 Lwz
  show ?thesis
    by (auto simp: divide_simps mult_less_0_iff norm_minus_commute
norm_divide norm_mult power2_eq_square)
  qed
} note L_cmod_le = this
let ?f = (λx. 1 / (x - w) - 1 / (x - z))
have cmod (contour_integral p ?f) ≤ L * (e * pe2 / L / 4 * (inverse (pe /
2))2)
proof (rule L)
  show ?f holomorphic_on - cball z (3 / 4 * pe)
  using <pe>0> w
  by (force simp: dist_norm norm_minus_commute intro!: holomorphic_intros)
  show ∀ u ∈ - cball z (3 / 4 * pe). cmod (?f u) ≤ e * pe2 / L / 4 * (inverse
(pe / 2))2
  using <pe>0> w <L>0>
  by (auto simp: cball_def dist_norm field_simps L_cmod_le simp del:
less_divide_eq_numeral1 le_divide_eq_numeral1)
  qed
also have ... < 2*e
  using <L>0> e by (force simp: field_simps)
finally have cmod (winding_number p w - winding_number p z) < e
  using pi_ge_two e
  by (force simp: winding_number_valid_path <valid_path p> <z ∉ path_image
p> field_simps norm_divide norm_mult intro: less_le_trans)
} note cmod_wn_diff = this
have isCont (winding_number p) z
proof (clarsimp simp add: continuous_at_eps_delta)
  fix e::real assume e>0
  show ∃ d>0. ∀ x'. dist x' z < d → dist (winding_number p x') (winding_number
p z) < e
  using <pe>0> <L>0> <e>0>
  by (rule_tac x=min (pe/4) (e/2*pe2/L/4) in exI) (simp add: dist_norm
cmod_wn_diff)

```

```

qed
then show ?thesis
  apply (rule continuous_transform_within [where d = min d e/2])
  apply (auto simp: <d>0 <e>0 dist_norm wnw)
done
qed

```

corollary *continuous_on_winding_number*:

```

path  $\gamma \implies \text{continuous\_on } (- \text{path\_image } \gamma) (\lambda w. \text{winding\_number } \gamma w)
\text{by (simp add: continuous\_at\_imp\_continuous\_on continuous\_at\_winding\_number)}$ 
```

3.4 The winding number is constant on a connected region

lemma *winding_number_constant*:

```

assumes  $\gamma$ : path  $\gamma$  and loop: pathfinish  $\gamma = \text{pathstart } \gamma$  and cs: connected  $S$ 
and sg:  $S \cap \text{path\_image } \gamma = \{\}$ 

```

```

shows winding_number  $\gamma$  constant_on  $S$ 

```

proof –

```

have *:  $1 \leq \text{cmod } (\text{winding\_number } \gamma y - \text{winding\_number } \gamma z)$ 

```

```

  if ne: winding_number  $\gamma y \neq \text{winding\_number } \gamma z$  and  $y \in S z \in S$  for  $y z$ 

```

proof –

```

  have winding_number  $\gamma y \in \mathbb{Z}$  winding_number  $\gamma z \in \mathbb{Z}$ 

```

```

  using that integer_winding_number [OF  $\gamma$  loop] sg < $y \in S$ > by auto

```

```

  with ne show ?thesis

```

```

  by (auto simp: Ints_def simp flip: of_int_diff)

```

qed

```

have cont: continuous_on  $S$  ( $\lambda w. \text{winding\_number } \gamma w$ )

```

```

  using continuous_on_winding_number [OF  $\gamma$ ] sg

```

```

  by (meson continuous_on_subset disjoint_eq_subset_Compl)

```

show ?thesis

```

  using * zero_less_one

```

```

  by (blast intro: continuous_discrete_range_constant [OF cs cont])

```

qed

lemma *winding_number_eq*:

```

[[path  $\gamma$ ; pathfinish  $\gamma = \text{pathstart } \gamma$ ;  $w \in S$ ;  $z \in S$ ; connected  $S$ ;  $S \cap \text{path\_image } \gamma = \{\}$ ]

```

```

 $\implies \text{winding\_number } \gamma w = \text{winding\_number } \gamma z$ 

```

```

  using winding_number_constant by (metis constant_on_def)

```

lemma *open_winding_number_levelsets*:

```

assumes  $\gamma$ : path  $\gamma$  and loop: pathfinish  $\gamma = \text{pathstart } \gamma$ 

```

```

shows open  $\{z. z \notin \text{path\_image } \gamma \wedge \text{winding\_number } \gamma z = k\}$ 

```

proof (clarsimp simp: open_dist)

```

fix  $z$  assume  $z$ :  $z \notin \text{path\_image } \gamma$  and  $k$ :  $k = \text{winding\_number } \gamma z$ 

```

```

have open ( $- \text{path\_image } \gamma$ )

```

```

  by (simp add: closed_path_image  $\gamma$  open_Compl)

```

```

then obtain  $e$  where  $e > 0$  ball  $z e \subseteq - \text{path\_image } \gamma$ 

```

```

  using open_contains_ball [of  $- \text{path\_image } \gamma$ ]  $z$  by blast

```

```

then show  $\exists e > 0. \forall y. \text{dist } y \ z < e \longrightarrow y \notin \text{path\_image } \gamma \wedge \text{winding\_number}$ 
 $\gamma \ y = \text{winding\_number } \gamma \ z$ 
using  $\langle e > 0 \rangle$  by (force simp: norm_minus_commute dist_norm intro: wind-
ing_number_eq [OF assms, where  $S = \text{ball } z \ e$ ])
qed

```

3.5 Winding number is zero "outside" a curve

proposition *winding_number_zero_in_outside:*

assumes γ : path γ **and** loop: pathfinish $\gamma = \text{pathstart } \gamma$ **and** z : $z \in \text{outside}$
(*path_image* γ)

shows $\text{winding_number } \gamma \ z = 0$

proof –

obtain $B::\text{real}$ **where** $0 < B$ **and** B : $\text{path_image } \gamma \subseteq \text{ball } 0 \ B$

using *bounded_subset_ballD* [OF *bounded_path_image* [OF γ]] **by** *auto*

obtain $w::\text{complex}$ **where** $w \notin \text{ball } 0 \ (B + 1)$

by (*metis abs_of_nonneg le_less less_irrefl mem_ball_0 norm_of_real*)

have $-\text{ball } 0 \ (B + 1) \subseteq \text{outside} \ (\text{path_image } \gamma)$

using B *subset_ball* **by** (*intro outside_subset_convex*) *auto*

then have $w \in \text{outside} \ (\text{path_image } \gamma)$

using w **by** *blast*

moreover have $\text{winding_number } \gamma \ \text{constant_on } \text{outside} \ (\text{path_image } \gamma)$

using *winding_number_constant* [OF γ loop, of *outside(path_image γ)*] *con-*
nected_outside

by (*metis DIM_complex bounded_path_image dual_order.refl γ outside_no_overlap*)

ultimately have $\text{winding_number } \gamma \ z = \text{winding_number } \gamma \ w$

by (*metis (no_types, opaque_lifting) constant_on_def z*)

also have $\dots = 0$

proof –

have $w \notin \text{path_image } \gamma$ **using** w **by** (*simp add: outside_def*)

{ **fix** $e::\text{real}$ **assume** $0 < e$

obtain p **where** p : *polynomial_function* p $\text{pathstart } p = \text{pathstart } \gamma$ pathfinish
 $p = \text{pathfinish } \gamma$

and $pg1$: $(\bigwedge t. \llbracket 0 \leq t; t \leq 1 \rrbracket \Longrightarrow \text{cmod } (p \ t - \gamma \ t) < 1)$

and $pg2$: $(\bigwedge t. \llbracket 0 \leq t; t \leq 1 \rrbracket \Longrightarrow \text{cmod } (p \ t - \gamma \ t) < e)$

using *path_approx_polynomial_function* [OF γ , of *min 1 e*] $\langle e > 0 \rangle$

by (*metis atLeastAtMost_iff min_less_iff_conj zero_less_one*)

have $\exists p. \text{valid_path } p \wedge w \notin \text{path_image } p \wedge$

$\text{pathstart } p = \text{pathstart } \gamma \wedge \text{pathfinish } p = \text{pathfinish } \gamma \wedge$

$(\forall t \in \{0..1\}. \text{cmod } (\gamma \ t - p \ t) < e) \wedge \text{contour_integral } p \ (\lambda wa. 1$

$/ (wa - w)) = 0$

proof (*intro exI conjI*)

have $\bigwedge x. \llbracket 0 \leq x; x \leq 1 \rrbracket \Longrightarrow \text{cmod } (p \ x) < B + 1$

using B **unfolding** *image_subset_iff path_image_def*

by (*meson add_strict_mono atLeastAtMost_iff le_less_trans mem_ball_0*
norm_triangle_sub pg1)

then have $\text{pip: path_image } p \subseteq \text{ball } 0 \ (B + 1)$

by (*auto simp add: path_image_def dist_norm ball_def*)

then show $w \notin \text{path_image } p$ **using** w **by** *blast*

```

show vap: valid_path p
  by (simp add: p(1) valid_path_polynomial_function)
show  $\forall t \in \{0..1\}. \text{cmod } (\gamma t - p t) < e$ 
  by (metis atLeastAtMost_iff norm_minus_commute pge)
show contour_integral p ( $\lambda w a. 1 / (w a - w)$ ) = 0
proof (rule contour_integral_unique [OF Cauchy_theorem_convex_simple
[OF _convex_ball [of 0 B+1]]])
  have  $\bigwedge z. \text{cmod } z < B + 1 \implies z \neq w$ 
    using mem_ball_0 w by blast
  then show ( $\lambda z. 1 / (z - w)$ ) holomorphic_on ball 0 (B + 1)
    by (intro holomorphic_intros; simp add: dist_norm)
  qed (use p vap pip loop in auto)
qed (use p in auto)
}
then show ?thesis
by (auto intro: winding_number_unique [OF  $\gamma$ ] simp add: winding_number_prop_def
wnot)
qed
finally show ?thesis .
qed

```

```

corollary winding_number_zero_const: a  $\neq$  z  $\implies$  winding_number ( $\lambda t. a$ ) z =
0
by (rule winding_number_zero_in_outside)
(auto simp: pathfinish_def pathstart_def path_polynomial_function)

```

```

corollary winding_number_zero_outside:
 $\llbracket \text{path } \gamma; \text{convex } s; \text{pathfinish } \gamma = \text{pathstart } \gamma; z \notin s; \text{path\_image } \gamma \subseteq s \rrbracket \implies$ 
winding_number  $\gamma$  z = 0
by (meson convex_in_outside outside_mono subsetCE winding_number_zero_in_outside)

```

```

lemma winding_number_zero_at_infinity:
assumes  $\gamma: \text{path } \gamma$  and loop: pathfinish  $\gamma = \text{pathstart } \gamma$ 
shows  $\exists B. \forall z. B \leq \text{norm } z \longrightarrow \text{winding\_number } \gamma z = 0$ 
proof -
obtain B::real where  $0 < B$  and B: path_image  $\gamma \subseteq \text{ball } 0 B$ 
using bounded_subset_ballD [OF bounded_path_image [OF  $\gamma$ ]] by auto
have winding_number  $\gamma$  z = 0 if  $B + 1 \leq \text{cmod } z$  for z
proof (rule winding_number_zero_outside [OF  $\gamma$  convex_cball loop])
show  $z \notin \text{cball } 0 B$ 
using that by auto
show path_image  $\gamma \subseteq \text{cball } 0 B$ 
using B order.trans by blast
qed
then show ?thesis
by metis
qed

```

```

lemma bounded_winding_number_nz:

```



```

assumes path g pathfinish g = pathstart g
shows bounded {z. winding_number g z ≠ 0}
proof –
obtain B where  $\bigwedge x. \text{norm } x \geq B \implies \text{winding\_number } g \ x = 0$ 
using winding_number_zero_at_infinity[OF assms] by auto
thus ?thesis
unfolding bounded_iff by (intro exI[of _ B + 1]) force
qed

```

```

lemma winding_number_zero_point:
   $\llbracket \text{path } \gamma; \text{convex } S; \text{pathfinish } \gamma = \text{pathstart } \gamma; \text{open } S; \text{path\_image } \gamma \subseteq S \rrbracket$ 
   $\implies \exists z. z \in S \wedge \text{winding\_number } \gamma \ z = 0$ 
using outside_compact_in_open [of path_image  $\gamma$  S] path_image_nonempty
winding_number_zero_in_outside
by (fastforce simp add: compact_path_image)

```

If a path winds round a set, it winds rounds its inside.

```

lemma winding_number_around_inside:
assumes  $\gamma$ : path  $\gamma$  and loop: pathfinish  $\gamma = \text{pathstart } \gamma$ 
and cls: closed S and cos: connected S and S_disj:  $S \cap \text{path\_image } \gamma = \{\}$ 
and z:  $z \in S$  and wn_nz: winding_number  $\gamma \ z \neq 0$  and w:  $w \in S \cup \text{inside } S$ 
shows winding_number  $\gamma \ w = \text{winding\_number } \gamma \ z$ 
proof –
have ssb:  $S \subseteq \text{inside}(\text{path\_image } \gamma)$ 
proof
fix x :: complex
assume  $x \in S$ 
hence  $x \notin \text{path\_image } \gamma$ 
by (meson disjoint_iff_not_equal S_disj)
thus  $x \in \text{inside}(\text{path\_image } \gamma)$ 
by (metis Compl_iff S_disj UnE  $\gamma \ \langle x \in S \rangle$  cos inside_outside loop winding_number_eq winding_number_zero_in_outside wn_nz z)
qed
show ?thesis
proof (rule winding_number_eq [OF  $\gamma$  loop w])
show  $z \in S \cup \text{inside } S$ 
using z by blast
show connected (S  $\cup$  inside S)
by (simp add: cls connected_with_inside cos)
show  $(S \cup \text{inside } S) \cap \text{path\_image } \gamma = \{\}$ 
unfolding disjoint_iff Un_iff
by (metis ComplD UnI1  $\gamma$  cls compact_path_image connected_path_image inside_Un_outside inside_inside_compact_connected ssb subsetD)
qed
qed

```

Bounding a WN by 1/2 for a path and point in opposite halfspaces.

```

lemma winding_number_subpath_continuous:

```

```

assumes  $\gamma$ : valid_path  $\gamma$  and  $z$ :  $z \notin \text{path\_image } \gamma$ 
shows continuous_on  $\{0..1\}$   $(\lambda x. \text{winding\_number}(\text{subpath } 0 \ x \ \gamma) \ z)$ 
proof (rule continuous_on_eq)
let  $?f = \lambda x. \text{integral } \{0..x\} (\lambda t. \text{vector\_derivative } \gamma \ (\text{at } t) / (\gamma \ t - z))$ 
show continuous_on  $\{0..1\}$   $(\lambda x. 1 / (2 * \text{pi} * \text{i}) * ?f \ x)$ 
proof (intro indefinite_integral_continuous_1 winding_number_exp_integral
continuous_intros)
show  $\gamma$  piecewise_C1_differentiable_on  $\{0..1\}$ 
using  $\gamma$  valid_path_def by blast
qed (use path_image_def z in auto)
show  $1 / (2 * \text{pi} * \text{i}) * ?f \ x = \text{winding\_number} \ (\text{subpath } 0 \ x \ \gamma) \ z$ 
if  $x \in \{0..1\}$  for  $x$ 
proof -
have  $1 / (2 * \text{pi} * \text{i}) * ?f \ x = 1 / (2 * \text{pi} * \text{i}) * \text{contour\_integral} \ (\text{subpath } 0 \ x \ \gamma)$ 
 $(\lambda w. 1 / (w - z))$ 
using assms x
by (simp add: contour_integral_subcontour_integral [OF contour_integrable_inversediff])
also have  $\dots = \text{winding\_number} \ (\text{subpath } 0 \ x \ \gamma) \ z$ 
proof (subst winding_number_valid_path)
show  $z \notin \text{path\_image} \ (\text{subpath } 0 \ x \ \gamma)$ 
using assms x atLeastAtMost_iff path_image_subpath_subset by force
qed (use assms x valid_path_subpath in <force+>)
finally show ?thesis .
qed
qed

```

lemma *winding_number_ivt_pos*:

```

assumes  $\gamma$ : valid_path  $\gamma$  and  $z$ :  $z \notin \text{path\_image } \gamma$  and  $0 \leq w \leq \text{Re}(\text{winding\_number}$ 
 $\gamma \ z)$ 
shows  $\exists t \in \{0..1\}. \text{Re}(\text{winding\_number}(\text{subpath } 0 \ t \ \gamma) \ z) = w$ 
proof -
have continuous_on  $\{0..1\}$   $(\lambda x. \text{winding\_number} \ (\text{subpath } 0 \ x \ \gamma) \ z)$ 
using  $\gamma$  winding_number_subpath_continuous z by blast
moreover have  $\text{Re} \ (\text{winding\_number} \ (\text{subpath } 0 \ 0 \ \gamma) \ z) \leq w \leq \text{Re} \ (\text{winding\_number}$ 
 $(\text{subpath } 0 \ 1 \ \gamma) \ z)$ 
using assms by (auto simp: path_image_def image_def)
ultimately show ?thesis
using ivt_increasing_component_on_1 [of 0 1, where ?k = 1] by force
qed

```

lemma *winding_number_ivt_neg*:

```

assumes  $\gamma$ : valid_path  $\gamma$  and  $z$ :  $z \notin \text{path\_image } \gamma$  and  $\text{Re}(\text{winding\_number}$ 
 $\gamma \ z) \leq w \leq 0$ 
shows  $\exists t \in \{0..1\}. \text{Re}(\text{winding\_number}(\text{subpath } 0 \ t \ \gamma) \ z) = w$ 
proof -
have continuous_on  $\{0..1\}$   $(\lambda x. \text{winding\_number} \ (\text{subpath } 0 \ x \ \gamma) \ z)$ 
using  $\gamma$  winding_number_subpath_continuous z by blast
moreover have  $\text{Re} \ (\text{winding\_number} \ (\text{subpath } 0 \ 0 \ \gamma) \ z) \geq w \geq \text{Re} \ (\text{winding\_number}$ 
 $(\text{subpath } 0 \ 1 \ \gamma) \ z)$ 

```

```

    using assms by (auto simp: path_image_def image_def)
  ultimately show ?thesis
    using ivt_decreasing_component_on_1 [of 0 1, where ?k = 1] by force
qed

```

lemma *winding_number_ivt_abs*:

```

  assumes  $\gamma$ : valid_path  $\gamma$  and  $z$ :  $z \notin \text{path\_image } \gamma$  and  $0 \leq w \leq |\text{Re}(\text{winding\_number } \gamma z)|$ 
  shows  $\exists t \in \{0..1\}. |\text{Re}(\text{winding\_number } (\text{subpath } 0 t \gamma) z)| = w$ 
  using assms winding_number_ivt_pos [of  $\gamma z w$ ] winding_number_ivt_neg [of  $\gamma z -w$ ]
  by force

```

lemma *winding_number_lt_half_lemma*:

```

  assumes  $\gamma$ : valid_path  $\gamma$  and  $z$ :  $z \notin \text{path\_image } \gamma$  and  $az$ :  $a \cdot z \leq b$  and  $pag$ :
  path_image  $\gamma \subseteq \{w. a \cdot w > b\}$ 

```

```

  shows  $\text{Re}(\text{winding\_number } \gamma z) < 1/2$ 

```

proof –

```

  { assume  $\text{Re}(\text{winding\_number } \gamma z) \geq 1/2$ 
    then obtain  $t::\text{real}$  where  $t: 0 \leq t \leq 1$  and  $sub12$ :  $\text{Re}(\text{winding\_number } (\text{subpath } 0 t \gamma) z) = 1/2$ 

```

```

    using winding_number_ivt_pos [OF  $\gamma z$ , of  $1/2$ ] by auto

```

```

    have  $gt$ :  $\gamma t - z = -(\text{of\_real } (\exp(- (2 * \pi * \text{Im}(\text{winding\_number } (\text{subpath } 0 t \gamma) z)))) * (\gamma 0 - z))$ 

```

```

    using winding_number_exp_2pi [of subpath 0 t  $\gamma z$ ]

```

```

    apply (simp add: t  $\gamma$  valid_path_imp_path)

```

```

    using closed_segment_eq_real_ivl path_image_def t z by (fastforce simp:
  path_image_subpath Euler sub12)

```

```

    have  $b < a \cdot \gamma 0$ 

```

proof –

```

    have  $\gamma 0 \in \{c. b < a \cdot c\}$ 

```

```

    by (metis (no_types) pag atLeastAtMost_iff image_subset_iff order_refl
  path_image_def zero_le_one)

```

```

    thus ?thesis

```

```

    by blast

```

qed

```

  moreover have  $b < a \cdot \gamma t$ 

```

```

  by (metis atLeastAtMost_iff image_eqI mem_Collect_eq pag path_image_def
  subset_iff t)

```

```

  ultimately have  $0 < a \cdot (\gamma 0 - z) \wedge 0 < a \cdot (\gamma t - z)$  using az

```

```

  by (simp add: inner_diff_right)+

```

```

  then have False

```

```

  by (simp add: gt inner_mult_right mult_less_0_iff)
}

```

```

  then show ?thesis by force

```

qed

lemma *winding_number_lt_half*:

```

  assumes valid_path  $\gamma$   $a \cdot z \leq b$  path_image  $\gamma \subseteq \{w. a \cdot w > b\}$ 

```

```

  shows  $|Re (winding\_number \ \gamma \ z)| < 1/2$ 
proof -
  have  $z \notin path\_image \ \gamma$  using assms by auto
  with assms have  $Re (winding\_number \ \gamma \ z) < 0 \implies - Re (winding\_number \ \gamma \ z) < 1/2$ 
  by (metis complex_inner_1_right winding_number_lt_half_lemma [OF valid_path_imp_reverse, of  $\gamma \ z \ a \ b$ ])
  winding_number_reversepath valid_path_imp_path inner_minus_left path_image_reversepath)
  with assms show ?thesis
  using  $\langle z \notin path\_image \ \gamma \rangle$  winding_number_lt_half_lemma by fastforce
qed

```

lemma *winding_number_le_half*:

```

  assumes  $\gamma$ : valid_path  $\gamma$  and  $z$ :  $z \notin path\_image \ \gamma$ 
  and anz:  $a \neq 0$  and azb:  $a \cdot z \leq b$  and pag:  $path\_image \ \gamma \subseteq \{w. a \cdot w \geq b\}$ 
  shows  $|Re (winding\_number \ \gamma \ z)| \leq 1/2$ 
proof -
  { assume wnz_12:  $|Re (winding\_number \ \gamma \ z)| > 1/2$ 
  have isCont (winding_number  $\gamma$ )  $z$ 
  by (metis continuous_at_winding_number valid_path_imp_path  $\gamma \ z$ )
  then obtain  $d$  where  $d > 0$  and  $d$ :  $\bigwedge x'. dist \ x' \ z < d \implies dist (winding\_number \ \gamma \ x') (winding\_number \ \gamma \ z) < |Re(winding\_number \ \gamma \ z)| - 1/2$ 
  using continuous_at_eps_delta wnz_12 diff_gt_0_iff_gt by blast
  define  $z'$  where  $z' = z - (d / (2 * cmod \ a)) *_R \ a$ 
  have  $a \cdot z * 6 \leq d * cmod \ a + b * 6$ 
  by (metis  $\langle 0 < d \rangle$  add_increasing azb less_eq_real_def mult_nonneg_nonneg mult_right_mono norm_ge_zero norm_numeral)
  with anz have  $*$ :  $a \cdot z' \leq b - d / 3 * cmod \ a$ 
  unfolding  $z'$ _def inner_mult_right' divide_inverse
  by (simp add: field_split_simps algebra_simps dot_square_norm power2_eq_square)
  have cmod (winding_number  $\gamma \ z' - winding\_number \ \gamma \ z$ )  $< |Re (winding\_number \ \gamma \ z)| - 1/2$ 
  using  $d$  [of  $z'$ ] anz  $\langle d > 0 \rangle$  by (simp add: dist_norm  $z'$ _def)
  then have  $1/2 < |Re (winding\_number \ \gamma \ z)| - cmod (winding\_number \ \gamma \ z' - winding\_number \ \gamma \ z)$ 
  by simp
  then have  $1/2 < |Re (winding\_number \ \gamma \ z)| - |Re (winding\_number \ \gamma \ z') - Re (winding\_number \ \gamma \ z)|$ 
  using abs_Re_le_cmod [of winding_number  $\gamma \ z' - winding\_number \ \gamma \ z$ ]
  by simp
  then have wnz_12':  $|Re (winding\_number \ \gamma \ z')| > 1/2$ 
  by linarith
  moreover have  $|Re (winding\_number \ \gamma \ z')| < 1/2$ 
  proof (rule winding_number_lt_half [OF  $\gamma \ *$ ])
  show  $path\_image \ \gamma \subseteq \{w. b - d / 3 * cmod \ a < a \cdot w\}$ 
  using azb  $\langle d > 0 \rangle$  pag by (auto simp: add_strict_increasing anz field_split_simps dest!: subsetD)
  qed
  ultimately have False

```

```

    by simp
  }
  then show ?thesis by force
qed

```

lemma *winding_number_lt_half_linepath*:

assumes $z \notin \text{closed_segment } a \ b$ **shows** $|\text{Re}(\text{winding_number}(\text{linepath } a \ b) \ z)| < 1/2$

proof –

obtain $u \ v$ **where** $u \cdot z \leq v$ **and** $uv: \bigwedge x. x \in \text{closed_segment } a \ b \implies \text{inner } u \ x > v$

using *separating_hyperplane_closed_point* *assms* *closed_segment* *convex_closed_segment* *less_eq_real_def* **by** *metis*

moreover **have** $\text{path_image}(\text{linepath } a \ b) \subseteq \{w. v < u \cdot w\}$

using *in_segment(1)* *uv* **by** *auto*

ultimately **show** ?thesis

using *winding_number_lt_half* **by** *auto*

qed

Positivity of WN for a linepath.

lemma *winding_number_linepath_pos_lt*:

assumes $0 < \text{Im}((b - a) * \text{cnj}(b - z))$

shows $0 < \text{Re}(\text{winding_number}(\text{linepath } a \ b) \ z)$

proof –

have $z: z \notin \text{path_image}(\text{linepath } a \ b)$

using *assms*

by (*simp* *add: closed_segment_def*) (*force* *simp: algebra_simps*)

show ?thesis

by (*intro* *winding_number_pos_lt* [*OF* *valid_path_linepath* *z* *assms*]) (*simp* *add: linepath_def* *algebra_simps*)

qed

3.6 More winding number properties

including the fact that it's +1 inside a simple closed curve.

lemma *winding_number_homotopic_paths*:

assumes *homotopic_paths* $(-\{z\}) \ g \ h$

shows $\text{winding_number } g \ z = \text{winding_number } h \ z$

proof –

have *path* *g* *path* *h* **using** *homotopic_paths_imp_path* [*OF* *assms*] **by** *auto*

moreover **have** *pag*: $z \notin \text{path_image } g$ **and** *pah*: $z \notin \text{path_image } h$

using *homotopic_paths_imp_subset* [*OF* *assms*] **by** *auto*

ultimately **obtain** $d \ e$ **where** $d > 0 \ e > 0$

and $d: \bigwedge p. \llbracket \text{path } p; \text{pathstart } p = \text{pathstart } g; \text{pathfinish } p = \text{pathfinish } g; \forall t \in \{0..1\}. \text{norm}(p \ t - g \ t) < d \rrbracket$

$\implies \text{homotopic_paths}(-\{z\}) \ g \ p$

and $e: \bigwedge q. \llbracket \text{path } q; \text{pathstart } q = \text{pathstart } h; \text{pathfinish } q = \text{pathfinish } h; \forall t \in \{0..1\}. \text{norm}(q \ t - h \ t) < e \rrbracket$

$\implies \text{homotopic_paths}(-\{z\}) \ h \ q$

```

    using homotopic_nearby_paths [of g -{z}] homotopic_nearby_paths [of h
- {z}] by force
  obtain p where p:
    valid_path p z  $\notin$  path_image p
    pathstart p = pathstart g pathfinish p = pathfinish g
    and gp_less:  $\forall t \in \{0..1\}. \text{cmod } (g t - p t) < d$ 
    and pap:  $\text{contour\_integral } p (\lambda w. 1 / (w - z)) = \text{complex\_of\_real } (2 * \pi)$ 
  * i * winding_number g z
    using winding_number [OF  $\langle \text{path } g \rangle \text{ pap } \langle 0 < d \rangle$ ] unfolding winding_number_prop_def
  by blast
  obtain q where q:
    valid_path q z  $\notin$  path_image q
    pathstart q = pathstart h pathfinish q = pathfinish h
    and hq_less:  $\forall t \in \{0..1\}. \text{cmod } (h t - q t) < e$ 
    and paq:  $\text{contour\_integral } q (\lambda w. 1 / (w - z)) = \text{complex\_of\_real } (2 * \pi)$ 
  * i * winding_number h z
    using winding_number [OF  $\langle \text{path } h \rangle \text{ paq } \langle 0 < e \rangle$ ] unfolding winding_number_prop_def
  by blast
  have homotopic_paths (- {z}) g p
    by (simp add: d p valid_path_imp_path_norm_minus_commute gp_less)
  moreover have homotopic_paths (- {z}) h q
    by (simp add: e q valid_path_imp_path_norm_minus_commute hq_less)
  ultimately have homotopic_paths (- {z}) p q
    by (blast intro: homotopic_paths_trans homotopic_paths_sym assms)
  then have  $\text{contour\_integral } p (\lambda w. 1 / (w - z)) = \text{contour\_integral } q (\lambda w. 1 / (w - z))$ 
    by (rule Cauchy_theorem_homotopic_paths) (auto intro!: holomorphic_intros simp: p q)
  then show ?thesis
    by (simp add: pap paq)
qed

lemma winding_number_homotopic_loops:
  assumes homotopic_loops (-{z}) g h
  shows  $\text{winding\_number } g z = \text{winding\_number } h z$ 
proof -
  have path g path h using homotopic_loops_imp_path [OF assms] by auto
  moreover have pag:  $z \notin \text{path\_image } g$  and pah:  $z \notin \text{path\_image } h$ 
    using homotopic_loops_imp_subset [OF assms] by auto
  moreover have gloop:  $\text{pathfinish } g = \text{pathstart } g$  and hloop:  $\text{pathfinish } h = \text{pathstart } h$ 
    using homotopic_loops_imp_loop [OF assms] by auto
  ultimately obtain d e where  $d > 0 e > 0$ 
    and d:  $\bigwedge p. [\text{path } p; \text{pathfinish } p = \text{pathstart } p; \forall t \in \{0..1\}. \text{norm } (p t - g t) < d]$ 
     $\implies \text{homotopic\_loops } (-\{z\}) g p$ 
    and e:  $\bigwedge q. [\text{path } q; \text{pathfinish } q = \text{pathstart } q; \forall t \in \{0..1\}. \text{norm } (q t - h t) < e]$ 
     $\implies \text{homotopic\_loops } (-\{z\}) h q$ 

```

```

    using homotopic_nearby_loops [of g -{z}] homotopic_nearby_loops [of h
    -{z}] by force
  obtain p where p:
    valid_path p z  $\notin$  path_image p
    pathstart p = pathstart g pathfinish p = pathfinish g
    and gp_less:  $\forall t \in \{0..1\}. \text{cmod } (g t - p t) < d$ 
    and pap:  $\text{contour\_integral } p (\lambda w. 1 / (w - z)) = \text{complex\_of\_real } (2 * \pi)$ 
  * i * winding_number g z
    using winding_number [OF  $\langle$ path g $\rangle$  pag  $\langle$ 0 < d $\rangle$ ] unfolding winding_number_prop_def
  by blast
  obtain q where q:
    valid_path q z  $\notin$  path_image q
    pathstart q = pathstart h pathfinish q = pathfinish h
    and hq_less:  $\forall t \in \{0..1\}. \text{cmod } (h t - q t) < e$ 
    and paq:  $\text{contour\_integral } q (\lambda w. 1 / (w - z)) = \text{complex\_of\_real } (2 * \pi)$ 
  * i * winding_number h z
    using winding_number [OF  $\langle$ path h $\rangle$  pah  $\langle$ 0 < e $\rangle$ ] unfolding winding_number_prop_def
  by blast
  have gp: homotopic_loops (- {z}) g p
    by (simp add: gloop d gp_less norm_minus_commute p valid_path_imp_path)
  have hq: homotopic_loops (- {z}) h q
    by (simp add: e hloop hq_less norm_minus_commute q valid_path_imp_path)
  have contour_integral p ( $\lambda w. 1 / (w - z)$ ) = contour_integral q ( $\lambda w. 1 / (w - z)$ )
  proof (rule Cauchy_theorem_homotopic_loops)
    show homotopic_loops (- {z}) p q
      by (blast intro: homotopic_loops_trans homotopic_loops_sym gp hq assms)
  qed (auto intro!: holomorphic_intros simp: p q)
  then show ?thesis
    by (simp add: pap paq)
qed

```

lemma winding_number_paths_linear_eq:

```

[[path g; path h; pathstart h = pathstart g; pathfinish h = pathfinish g;
 $\wedge t. t \in \{0..1\} \implies z \notin \text{closed\_segment } (g t) (h t)$ ]]
 $\implies \text{winding\_number } h z = \text{winding\_number } g z$ 

```

by (blast intro: sym homotopic_paths_linear winding_number_homotopic_paths)

lemma winding_number_loops_linear_eq:

```

[[path g; path h; pathfinish g = pathstart g; pathfinish h = pathstart h;
 $\wedge t. t \in \{0..1\} \implies z \notin \text{closed\_segment } (g t) (h t)$ ]]
 $\implies \text{winding\_number } h z = \text{winding\_number } g z$ 

```

by (blast intro: sym homotopic_loops_linear winding_number_homotopic_loops)

lemma winding_number_nearby_paths_eq:

```

[[path g; path h; pathstart h = pathstart g; pathfinish h = pathfinish g;
 $\wedge t. t \in \{0..1\} \implies \text{norm}(h t - g t) < \text{norm}(g t - z)$ ]]
 $\implies \text{winding\_number } h z = \text{winding\_number } g z$ 

```

by (metis segment_bound(2) norm_minus_commute not_le winding_number_paths_linear_eq)

lemma *winding_number_nearby_loops_eq*:
 $\llbracket \text{path } g; \text{ path } h; \text{ pathfinish } g = \text{pathstart } g; \text{ pathfinish } h = \text{pathstart } h; \wedge t. t \in \{0..1\} \implies \text{norm}(h \ t - g \ t) < \text{norm}(g \ t - z) \rrbracket$
 $\implies \text{winding_number } h \ z = \text{winding_number } g \ z$
by (*metis segment_bound(2) norm_minus_commute not_le winding_number_loops_linear_eq*)

lemma *winding_number_subpath_combine*:
assumes *path g and notin*: $z \notin \text{path_image } g$ **and** $u \in \{0..1\} \ v \in \{0..1\} \ w \in \{0..1\}$
shows $\text{winding_number } (\text{subpath } u \ v \ g) \ z + \text{winding_number } (\text{subpath } v \ w \ g) \ z =$
 $\text{winding_number } (\text{subpath } u \ w \ g) \ z$ (**is** *?lhs = ?rhs*)

proof –
have *?lhs = winding_number (subpath u v g +++ subpath v w g) z*
using *assms*
by (*metis (no_types, lifting) path_image_subpath_subset path_subpath pathfinish_subpath pathstart_subpath_subsetD winding_number_join*)
also have $\dots = ?rhs$
by (*meson assms homotopic_join_subpaths subset_Compl_singleton winding_number_homotopic_paths*)
finally show *?thesis* .
qed

Winding numbers of circular contours

proposition *winding_number_part_circlepath_pos_less*:
assumes $s < t$ **and** *no*: $\text{norm}(w - z) < r$
shows $0 < \text{Re } (\text{winding_number}(\text{part_circlepath } z \ r \ s \ t) \ w)$
proof –
have $0 < r$ **by** (*meson no norm_not_less_zero not_le order.strict_trans2*)
note *valid_path_part_circlepath*
moreover have $w \notin \text{path_image } (\text{part_circlepath } z \ r \ s \ t)$
using *assms* **by** (*auto simp: path_image_def image_def part_circlepath_def norm_mult linpath_def*)
moreover have $0 < r * (t - s) * (r - \text{cmod } (w - z))$
using *assms* **by** (*metis <0 < r> diff_gt_0_iff_gt mult_pos_pos*)
ultimately show *?thesis*
apply (*rule winding_number_pos_lt [where e = r*(t - s)*(r - norm(w - z))]*)
apply (*simp add: vector_derivative_part_circlepath right_diff_distrib [symmetric] mult_ac mult_le_cancel_left_pos assms <0 < r>*)
using *Re_Im_le_cmod [of w-z linpath s t x for x]*
by (*simp add: exp_Euler cos_of_real sin_of_real part_circlepath_def algebra_simps cos_squared_eq [unfolded power2_eq_square]*)
qed

lemma *winding_number_circlepath_centre*: $0 < r \implies \text{winding_number } (\text{circlepath } z \ r) \ z = 1$
apply (*rule winding_number_unique_loop*)


```

apply (simp_all add: sphere_def valid_path_imp_path)
apply (rule_tac x=circlepath z r in exI)
apply (simp add: sphere_def contour_integral_circlepath)
done

```

proposition *winding_number_circlepath*:

assumes $\text{norm}(w - z) < r$ **shows** $\text{winding_number}(\text{circlepath } z \ r) \ w = 1$

proof (cases $w = z$)

case *True* **then show** *?thesis*

using *assms winding_number_circlepath_centre* **by** *auto*

next

case *False*

have [*simp*]: $r > 0$

using *assms le_less_trans norm_ge_zero* **by** *blast*

define r' **where** $r' = \text{norm}(w - z)$

have $r' < r$

by (*simp add: assms r'_def*)

have *disjo*: $\text{cball } z \ r' \cap \text{sphere } z \ r = \{\}$

using $\langle r' < r \rangle$ **by** (*force simp: cball_def sphere_def*)

have $\text{winding_number}(\text{circlepath } z \ r) \ w = \text{winding_number}(\text{circlepath } z \ r) \ z$

proof (*rule winding_number_around_inside* [**where** $S = \text{cball } z \ r'$])

show $\text{winding_number}(\text{circlepath } z \ r) \ z \neq 0$

by (*simp add: winding_number_circlepath_centre*)

show $\text{cball } z \ r' \cap \text{path_image}(\text{circlepath } z \ r) = \{\}$

by (*simp add: disjo less_eq_real_def*)

qed (*auto simp: r'_def dist_norm norm_minus_commute*)

also have $\dots = 1$

by (*simp add: winding_number_circlepath_centre*)

finally show *?thesis* .

qed

lemma *no_bounded_connected_component_imp_winding_number_zero*:

assumes $g: \text{path } g \ \text{path_image } g \subseteq S \ \text{pathfinish } g = \text{pathstart } g \ z \notin S$

and $\text{nb}: \bigwedge z. \text{bounded}(\text{connected_component_set}(-S) \ z) \implies z \in S$

shows $\text{winding_number } g \ z = 0$

proof –

have $z \in \text{outside}(\text{path_image } g)$

by (*metis nb [of z] $\langle \text{path_image } g \subseteq S \rangle \langle z \notin S \rangle \text{subsetD mem_Collect_eq}$*)

outside outside_mono)

then show *?thesis*

by (*simp add: g_winding_number_zero_in_outside*)

qed

lemma *no_bounded_path_component_imp_winding_number_zero*:

assumes $g: \text{path } g \ \text{path_image } g \subseteq S \ \text{pathfinish } g = \text{pathstart } g \ z \notin S$

and $\text{nb}: \bigwedge z. \text{bounded}(\text{path_component_set}(-S) \ z) \implies z \in S$

shows $\text{winding_number } g \ z = 0$

by (*simp add: bounded_subset nb path_component_subset_connected_component*)

no_bounded_connected_component_imp_winding_number_zero [OF *g*])

3.7 Winding number for a triangle

lemma *wn_triangle1*:

```

assumes  $0 \in \text{interior}(\text{convex hull } \{a,b,c\})$ 
shows  $\neg (Im(a/b) \leq 0 \wedge 0 \leq Im(b/c))$ 
proof -
  { assume  $0: Im(a/b) \leq 0 \ 0 \leq Im(b/c)$ 
    have  $0 \notin \text{interior}(\text{convex hull } \{a,b,c\})$ 
    proof (cases a=0  $\vee$  b=0  $\vee$  c=0)
      case True then show ?thesis
        by (auto simp: not_in_interior_convex_hull_3)
    next
      case False
      then have  $b \neq 0$  by blast
      { fix  $x y::\text{complex}$  and  $u::\text{real}$ 
        assume  $eq\_f': Im\ x * Re\ b \leq Im\ b * Re\ x \ Im\ y * Re\ b \leq Im\ b * Re\ y \ 0 \leq$ 
 $u \ u \leq 1$ 
        then have  $((1 - u) * Im\ x) * Re\ b \leq Im\ b * ((1 - u) * Re\ x)$ 
          by (simp add: mult_left_mono mult.assoc mult.left_commute [of Im b])
        then have  $((1 - u) * Im\ x + u * Im\ y) * Re\ b \leq Im\ b * ((1 - u) * Re\ x$ 
 $+ u * Re\ y)$ 
          using eq_f' ordered_comm_semiring_class.comm_mult_left_mono
          by (fastforce simp add: algebra_simps)
        }
      then have  $\text{convex } \{z. Im\ z * Re\ b \leq Im\ b * Re\ z\}$ 
        by (auto simp: algebra_simps convex_alt)
      with False 0 have  $\text{convex hull } \{a,b,c\} \leq \{z. Im\ z * Re\ b \leq Im\ b * Re\ z\}$ 
        by (simp add: subset_hull mult.commute Complex.Im_divide divide_simps
 $\text{complex_neq_0 [symmetric]}$ )
      moreover have  $0 \notin \text{interior}(\{z. Im\ z * Re\ b \leq Im\ b * Re\ z\})$ 
      proof
        assume  $0 \in \text{interior } \{z. Im\ z * Re\ b \leq Im\ b * Re\ z\}$ 
        then obtain  $e$  where  $e > 0$  and  $e: \text{ball } 0\ e \subseteq \{z. Im\ z * Re\ b \leq Im\ b * Re\ z\}$ 
          by (meson mem_interior)
        define  $z$  where  $z \equiv -\text{sgn } (Im\ b) * (e/3) + \text{sgn } (Re\ b) * (e/3) * i$ 
        have  $\text{cmod } z = \text{cmod } (e/3) * \text{cmod } (-\text{sgn } (Im\ b) + \text{sgn } (Re\ b) * i)$ 
          unfolding  $z\_def$  norm_mult [symmetric] by (simp add: algebra_simps)
        also have  $\dots < e$ 
          using  $\langle e > 0 \rangle$  by (auto simp: norm_mult intro: le_less_trans [OF
 $\text{norm_triangle_ineq4}]$ )
        finally have  $z \in \text{ball } 0\ e$ 
          using  $\langle e > 0 \rangle$  by simp
        then have  $Im\ z * Re\ b \leq Im\ b * Re\ z$ 
          using  $e$  by blast
        then have  $b: 0 < Re\ b \ 0 < Im\ b$  and  $\text{disj: } e * Re\ b < -(Im\ b * e) \vee e *$ 

```

```

Re b = - (Im b * e)
  using ‹e>0› ‹b ≠ 0›
  by (auto simp add: z_def dist_norm sgn_if less_eq_real_def mult_less_0_iff
complex.expand_split: if_split_asm)
  show False — or just one smt line
  using disj
  proof
    assume e * Re b < - (Im b * e)
    with b ‹0 < e› less_trans [of _ 0] show False
  by (metis (no_types) mult_pos_pos neg_less_0_iff_less not_less_iff_gr_or_eq)
  next
    assume e * Re b = - (Im b * e)
    with b ‹0 < e› show False
  by (metis mult_pos_pos neg_less_0_iff_less not_less_iff_gr_or_eq)
  qed
  qed
  ultimately show ?thesis
  using interior_mono by blast
  qed
} with assms show ?thesis by blast
qed

```

lemma *wn_triangle2_0*:

```

assumes 0 ∈ interior(convex hull {a,b,c})
shows
  0 < Im((b - a) * cnj (b)) ∧
  0 < Im((c - b) * cnj (c)) ∧
  0 < Im((a - c) * cnj (a))
  ∨
  Im((b - a) * cnj (b)) < 0 ∧
  0 < Im((b - c) * cnj (b)) ∧
  0 < Im((a - b) * cnj (a)) ∧
  0 < Im((c - a) * cnj (c))
proof —
  have [simp]: {b,c,a} = {a,b,c} {c,a,b} = {a,b,c} by auto
  show ?thesis
  using wn_triangle1 [OF assms] wn_triangle1 [of b c a] wn_triangle1 [of c a
b] assms
  by (auto simp: algebra_simps Im_complex_div_gt_0 Im_complex_div_lt_0
not_le not_less)
qed

```

lemma *wn_triangle2*:

```

assumes z ∈ interior(convex hull {a,b,c})
shows 0 < Im((b - a) * cnj (b - z)) ∧
  0 < Im((c - b) * cnj (c - z)) ∧
  0 < Im((a - c) * cnj (a - z))
  ∨
  Im((b - a) * cnj (b - z)) < 0 ∧

```

$$\begin{aligned}
& 0 < \operatorname{Im}((b - c) * \operatorname{cnj} (b - z)) \wedge \\
& 0 < \operatorname{Im}((a - b) * \operatorname{cnj} (a - z)) \wedge \\
& 0 < \operatorname{Im}((c - a) * \operatorname{cnj} (c - z))
\end{aligned}$$

proof –

have $0: 0 \in \operatorname{interior}(\operatorname{convex_hull} \{a-z, b-z, c-z\})$
using *assms convex_hull_translation* [of $-z \{a,b,c\}$]
interior_translation [of $-z$]
by (*simp cong: image_cong_simp*)
show *?thesis using wn_triangle2_0* [OF 0]
by *simp*

qed

lemma *wn_triangle3*:

assumes $z: z \in \operatorname{interior}(\operatorname{convex_hull} \{a,b,c\})$
and $0 < \operatorname{Im}((b-a) * \operatorname{cnj} (b-z))$
 $0 < \operatorname{Im}((c-b) * \operatorname{cnj} (c-z))$
 $0 < \operatorname{Im}((a-c) * \operatorname{cnj} (a-z))$
shows $\operatorname{winding_number} (\operatorname{linepath} a b \operatorname{+++} \operatorname{linepath} b c \operatorname{+++} \operatorname{linepath} c a) z$
 $= 1$

proof –

have *znot[simp]*: $z \notin \operatorname{closed_segment} a b \ z \notin \operatorname{closed_segment} b c \ z \notin \operatorname{closed_segment} c a$

using *z interior_of_triangle* [of $a b c$]
by (*auto simp: closed_segment_def*)

have *gt0*: $0 < \operatorname{Re} (\operatorname{winding_number} (\operatorname{linepath} a b \operatorname{+++} \operatorname{linepath} b c \operatorname{+++} \operatorname{linepath} c a) z)$

using *assms*

by (*simp add: winding_number_linepath_pos_lt_path_image_join winding_number_join_pos_combined*)

have *lt2*: $\operatorname{Re} (\operatorname{winding_number} (\operatorname{linepath} a b \operatorname{+++} \operatorname{linepath} b c \operatorname{+++} \operatorname{linepath} c a) z) < 2$

using *winding_number_lt_half_linepath* [of $a b$]

using *winding_number_lt_half_linepath* [of $b c$]

using *winding_number_lt_half_linepath* [of $c a$] *znot*

by (*fastforce simp add: winding_number_join_path_image_join*)

show *?thesis*

by (*rule winding_number_eq_1*) (*simp_all add: path_image_join gt0 lt2*)

qed

proposition *winding_number_triangle*:

assumes $z: z \in \operatorname{interior}(\operatorname{convex_hull} \{a,b,c\})$

shows $\operatorname{winding_number}(\operatorname{linepath} a b \operatorname{+++} \operatorname{linepath} b c \operatorname{+++} \operatorname{linepath} c a) z =$
(if $0 < \operatorname{Im}((b - a) * \operatorname{cnj} (b - z))$ *then* 1 *else* -1 *)*

proof –

have [*simp*]: $\{a,c,b\} = \{a,b,c\}$ **by** *auto*

have *znot[simp]*: $z \notin \operatorname{closed_segment} a b \ z \notin \operatorname{closed_segment} b c \ z \notin \operatorname{closed_segment} c a$

using *z interior_of_triangle* [of $a b c$]

by (*auto simp: closed_segment_def*)

then have [*simp*]: $z \notin \operatorname{closed_segment} b a \ z \notin \operatorname{closed_segment} c b \ z \notin \operatorname{closed_segment}$

```

a c
  using closed_segment_commute by blast+
  have *: winding_number (linepath a b +++ linepath b c +++ linepath c a) z =
    winding_number (reversepath (linepath a c +++ linepath c b +++
linepath b a)) z
  by (simp add: reversepath_joinpaths winding_number_join not_in_path_image_join)
  show ?thesis
  apply (rule disjE [OF wn_triangle2 [OF z]])
  subgoal
    by (simp add: wn_triangle3 z)
  subgoal
    by (simp add: path_image_join winding_number_reversepath * wn_triangle3
z)
  done
qed

```

3.8 Winding numbers for simple closed paths

lemma *winding_number_from_innerpath*:

```

  assumes simple_path c1 and c1: pathstart c1 = a pathfinish c1 = b
    and simple_path c2 and c2: pathstart c2 = a pathfinish c2 = b
    and simple_path c and c: pathstart c = a pathfinish c = b
    and c1c2: path_image c1 ∩ path_image c2 = {a,b}
    and c1c: path_image c1 ∩ path_image c = {a,b}
    and c2c: path_image c2 ∩ path_image c = {a,b}
    and ne_12: path_image c ∩ inside(path_image c1 ∪ path_image c2) ≠ {}
    and z: z ∈ inside(path_image c1 ∪ path_image c)
    and wn_d: winding_number (c1 +++ reversepath c) z = d
    and a ≠ b d ≠ 0
  obtains z ∈ inside(path_image c1 ∪ path_image c2) winding_number (c1 +++
reversepath c2) z = d
proof -
  obtain 0: inside(path_image c1 ∪ path_image c) ∩ inside(path_image c2 ∪
path_image c) = {}
    and 1: inside(path_image c1 ∪ path_image c) ∪ inside(path_image c2 ∪
path_image c) ∪
      (path_image c - {a,b}) = inside(path_image c1 ∪ path_image c2)
  by (rule split_inside_simple_closed_curve
[OF ⟨simple_path c1⟩ c1 ⟨simple_path c2⟩ c2 ⟨simple_path c⟩ c ⟨a ≠
b⟩ c1c2 c1c c2c ne_12])
  have znot: z ∉ path_image c z ∉ path_image c1 z ∉ path_image c2
  using union_with_outside z 1 by auto
  then have zout: z ∈ outside (path_image c ∪ path_image c2)
  by (metis 0 ComplI UnE disjoint_iff_not_equal sup.commute union_with_inside
z)
  have wn_cc2: winding_number (c +++ reversepath c2) z = 0
  by (rule winding_number_zero_in_outside; simp add: zout ⟨simple_path c2⟩
c c2 ⟨simple_path c⟩ simple_path_imp_path_image_join)
  show ?thesis

```

```

proof
  show  $z \in \text{inside } (\text{path\_image } c1 \cup \text{path\_image } c2)$ 
    using 1 z by blast
  have  $\text{winding\_number } c1 \ z - \text{winding\_number } c \ z = d$ 
    using assms znot
  by (metis wn_d winding_number_join_simple_path_imp_path_winding_number_reversepath
add.commute_path_image_reversepath_path_reversepath_pathstart_reversepath_minus_add_conv_diff)
  then show  $\text{winding\_number } (c1 +++ \text{reversepath } c2) \ z = d$ 
    using wn_cc2 by (simp add: winding_number_join_simple_path_imp_path
assms znot winding_number_reversepath)
qed
qed

```

lemma *simple_closed_path_wn1*:

```

fixes a::complex and e::real
assumes  $0 < e$ 
  and sp_pl: simple_path(p +++ linepath (a - e) (a + e)) (is simple_path
?pae)
  and psp: pathstart p = a + e
  and pfp: pathfinish p = a - e
  and disj:  $\text{ball } a \ e \cap \text{path\_image } p = \{\}$ 
obtains z where  $z \in \text{inside } (\text{path\_image } ?pae) \ \text{cmod } (\text{winding\_number } ?pae \ z) = 1$ 
proof -
  have arc p and arc_lp: arc (linepath (a - e) (a + e))
    and pap:  $\text{path\_image } p \cap \text{path\_image } (\text{linepath } (a - e) (a + e)) \subseteq \{\text{pathstart } p, a - e\}$ 
  using simple_path_join_loop_eq [of linepath (a - e) (a + e) p] assms by auto
  have mid_eq_a: midpoint (a - e) (a + e) = a
    by (simp add: midpoint_def)
  with assms have  $a \in \text{path\_image } ?pae$ 
    by (simp add: assms path_image_join) (metis midpoint_in_closed_segment)
  then have  $a \in \text{frontier}(\text{inside } (\text{path\_image } ?pae))$ 
    using assms by (simp add: Jordan_inside_outside)
  with  $\langle 0 < e \rangle$  obtain c where  $c \in \text{inside } (\text{path\_image } ?pae)$ 
    and dac:  $\text{dist } a \ c < e$ 
  by (auto simp: frontier_straddle)
  then have  $c \notin \text{path\_image } ?pae$ 
    using inside_no_overlap by blast
  then have  $c \notin \text{path\_image } p \ \wedge \ c \notin \text{closed\_segment } (a - e) (a + e)$ 
    by (simp_all add: assms path_image_join)
  with  $\langle 0 < e \rangle$  dac have  $c \notin \text{affine\_hull } \{a - e, a + e\}$ 
    by (simp add: segment_as_ball not_le)
  with  $\langle 0 < e \rangle$  have  $*$ :  $\neg \text{collinear } \{a - e, c, a + e\}$ 
    using collinear_3_affine_hull [of a-e a+e] by (auto simp: insert_commute)
  have 13:  $1/3 + 1/3 + 1/3 = (1::\text{real})$  by simp
  have  $(1/3) *_{\mathbb{R}} (a - \text{of\_real } e) + (1/3) *_{\mathbb{R}} c + (1/3) *_{\mathbb{R}} (a + \text{of\_real } e) \in \text{interior}(\text{convex\_hull } \{a - e, c, a + e\})$ 

```

```

    using interior_convex_hull_3_minimal [OF * DIM_complex]
    by clarsimp (metis 13 zero_less_divide_1_iff zero_less_numeral)
  then obtain z where z: z ∈ interior(convex hull {a - e, c, a + e}) by force
  have [simp]: z ∉ closed_segment (a - e) c
    by (metis DIM_complex Diff_iff IntD2 inf_sup_absorb interior_of_triangle
z)
  have [simp]: z ∉ closed_segment (a + e) (a - e)
    by (metis DIM_complex DiffD2 Un_iff interior_of_triangle z)
  have [simp]: z ∉ closed_segment c (a + e)
    by (metis (no_types, lifting) DIM_complex DiffD2 Un_insert_right inf_sup_aci(5)
insertCI interior_of_triangle mk_disjoint_insert z)
  show thesis
  proof
    have norm (winding_number (linepath (a - e) c +++ linepath c (a + e) +++
linepath (a + e) (a - e)) z) = 1
      using winding_number_triangle [OF z] by simp
    have zin: z ∈ inside (path_image (linepath (a + e) (a - e)) ∪ path_image p)
      and zeq: winding_number (linepath (a + e) (a - e) +++ reversepath p) z =
        winding_number (linepath (a - e) c +++ linepath c (a + e) +++
linepath (a + e) (a - e)) z
    proof (rule winding_number_from_innerpath
      [of linepath (a + e) (a - e) a+e a-e p
        linepath (a + e) c +++ linepath c (a - e) z
        winding_number (linepath (a - e) c +++ linepath c (a + e) +++
linepath (a + e) (a - e)) z])
      show sp_aec: simple_path (linepath (a + e) c +++ linepath c (a - e))
    proof (rule arc_imp_simple_path [OF arc_join])
      show arc (linepath (a + e) c)
        by (metis ‹c ∉ path_image p› arc_linepath pathstart_in_path_image psp)
      show arc (linepath c (a - e))
        by (metis ‹c ∉ path_image p› arc_linepath pathfinish_in_path_image pfp)
      show path_image (linepath (a + e) c) ∩ path_image (linepath c (a - e))
        ⊆ {pathstart (linepath c (a - e))}
      by clarsimp (metis * IntI Int_closed_segment closed_segment_commute
singleton_iff)
    qed auto
    show simple_path p
      using ‹arc p› arc_simple_path by blast
    show sp_ae2: simple_path (linepath (a + e) (a - e))
      using ‹arc p› arc_distinct_ends pfp psp by fastforce
    show pa: pathfinish (linepath (a + e) (a - e)) = a - e
      pathstart (linepath (a + e) c +++ linepath c (a - e)) = a + e
      pathfinish (linepath (a + e) c +++ linepath c (a - e)) = a - e
      pathstart p = a + e pathfinish p = a - e
      pathstart (linepath (a + e) (a - e)) = a + e
      by (simp_all add: assms)
    show 1: path_image (linepath (a + e) (a - e)) ∩ path_image p = {a + e,
a - e}
    proof

```

```

    show path_image (linepath (a + e) (a - e)) ∩ path_image p ⊆ {a + e, a
- e}
    using pap closed_segment_commute psp segment_convex_hull by fastforce
  show {a + e, a - e} ⊆ path_image (linepath (a + e) (a - e)) ∩ path_image
p
    using pap pathfinish_in_path_image pathstart_in_path_image pfp psp
by fastforce
  qed
  show 2: path_image (linepath (a + e) (a - e)) ∩ path_image (linepath (a +
e) c +++ linepath c (a - e)) =
    {a + e, a - e} (is ?lhs = ?rhs)
  proof
    have ¬ collinear {c, a + e, a - e}
      using * by (simp add: insert_commute)
    then have convex_hull {a + e, a - e} ∩ convex_hull {a + e, c} = {a + e}
      convex_hull {a + e, a - e} ∩ convex_hull {c, a - e} = {a - e}
    by (metis (full_types) Int_closed_segment insert_commute segment_convex_hull)+
    then show ?lhs ⊆ ?rhs
      by (metis Int_Un_distrib equalityD1 insert_is_Un path_image_join
path_image_linepath path_join_eq path_linepath segment_convex_hull simple_path_def
sp_aec)
    show ?rhs ⊆ ?lhs
      using segment_convex_hull by (simp add: path_image_join)
  qed
  have path_image p ∩ path_image (linepath (a + e) c) ⊆ {a + e}
  proof (clarsimp simp: path_image_join)
    fix x
    assume x ∈ path_image p and x_ac: x ∈ closed_segment (a + e) c
    then have dist x a ≥ e
      by (metis IntI all_not_in_conv disj_dist_commute mem_ball not_less)
    with x_ac dac ⟨e > 0⟩ show x = a + e
    by (auto simp: norm_minus_commute dist_norm closed_segment_eq_open
dest: open_segment_furthest_le [where y=a])
  qed
  moreover
  have path_image p ∩ path_image (linepath c (a - e)) ⊆ {a - e}
  proof (clarsimp simp: path_image_join)
    fix x
    assume x ∈ path_image p and x_ac: x ∈ closed_segment c (a - e)
    then have dist x a ≥ e
      by (metis IntI all_not_in_conv disj_dist_commute mem_ball not_less)
    with x_ac dac ⟨e > 0⟩ show x = a - e
    by (auto simp: norm_minus_commute dist_norm closed_segment_eq_open
dest: open_segment_furthest_le [where y=a])
  qed
  ultimately
  have path_image p ∩ path_image (linepath (a + e) c +++ linepath c (a -
e)) ⊆ {a + e, a - e}
    by (force simp: path_image_join)

```



```

then show 3: path_image p  $\cap$  path_image (linepath (a + e) c +++ linepath
c (a - e)) = {a + e, a - e}
  using 1 2 by blast
show 4: path_image (linepath (a + e) c +++ linepath c (a - e))  $\cap$ 
  inside (path_image (linepath (a + e) (a - e))  $\cup$  path_image p)  $\neq$  {}
apply (clar simp simp: path_image_join segment_convex_hull disjoint_iff_not_equal)
by (metis (no_types, opaque_lifting) UnI1 Un_commute c closed_segment_commute
ends_in_segment(1) path_image_join
  path_image_linepath pathstart_linepath pfp segment_convex_hull)
show zin_inside: z  $\in$  inside (path_image (linepath (a + e) (a - e))  $\cup$ 
  path_image (linepath (a + e) c +++ linepath c (a -
e)))
proof (simp add: path_image_join)
  show z  $\in$  inside (closed_segment (a + e) (a - e)  $\cup$  (closed_segment (a +
e) c  $\cup$  closed_segment c (a - e)))
  by (metis z inside_of_triangle DIM_complex Un_commute closed_segment_commute)
qed
show 5: winding_number
  (linepath (a + e) (a - e) +++ reversepath (linepath (a + e) c +++
linepath c (a - e))) z =
  winding_number (linepath (a - e) c +++ linepath c (a + e) +++
linepath (a + e) (a - e)) z
  by (simp add: reversepath_joinpaths path_image_join winding_number_join)
show 6: winding_number (linepath (a - e) c +++ linepath c (a + e) +++
linepath (a + e) (a - e)) z  $\neq$  0
  by (simp add: winding_number_triangle z)
show winding_number (linepath (a + e) (a - e) +++ reversepath p) z =
  winding_number (linepath (a - e) c +++ linepath c (a + e) +++
linepath (a + e) (a - e)) z
  by (metis 1 2 3 4 5 6 pa sp_aec sp_ae2 <arc p> <simple_path p> arc_distinct_ends
winding_number_from_innerpath zin_inside)
qed (use assms <e > 0> in auto)
show z_inside: z  $\in$  inside (path_image ?pae)
using zin by (simp add: assms path_image_join Un_commute closed_segment_commute)
have cmod (winding_number ?pae z) = cmod ((winding_number(reversepath
?pae) z))
proof (subst winding_number_reversepath)
  show path ?pae
  using simple_path_imp_path sp_pl by blast
  show z  $\notin$  path_image ?pae
  by (metis IntI emptyE inside_no_overlap z_inside)
qed (simp add: inside_def)
also have ... = cmod (winding_number(linepath (a + e) (a - e) +++ re-
versepath p) z)
  by (simp add: pfp reversepath_joinpaths)
also have ... = cmod (winding_number (linepath (a - e) c +++ linepath c (a
+ e) +++ linepath (a + e) (a - e)) z)
  by (simp add: zeq)
also have ... = 1

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    using z by (simp add: interior_of_triangle winding_number_triangle)
    finally show cmod (winding_number ?pae z) = 1 .
qed
qed

```

lemma *simple_closed_path_wn2*:

```

  fixes a::complex and d e::real
  assumes 0 < d 0 < e
    and sp_pl: simple_path(p +++ linepath (a - d) (a + e))
    and psp: pathstart p = a + e
    and pfp: pathfinish p = a - d
  obtains z where z ∈ inside (path_image (p +++ linepath (a - d) (a + e)))
    cmod (winding_number (p +++ linepath (a - d) (a + e)) z) = 1
  proof -
    have [simp]: a + of_real x ∈ closed_segment (a - α) (a - β) ↔ x ∈
      closed_segment (-α) (-β) for x α β::real
    using closed_segment_translation_eq [of a]
    by (metis (no_types, opaque_lifting) add_uminus_conv_diff of_real_minus
      of_real_closed_segment)
    have [simp]: a - of_real x ∈ closed_segment (a + α) (a + β) ↔ -x ∈
      closed_segment α β for x α β::real
    by (metis closed_segment_translation_eq diff_conv_add_uminus of_real_closed_segment
      of_real_minus)
    have arc p and arc_lp: arc (linepath (a - d) (a + e)) and path p
    and pap: path_image p ∩ closed_segment (a - d) (a + e) ⊆ {a+e, a-d}
    using simple_path_join_loop_eq [of linepath (a - d) (a + e) p] assms
    arc_imp_path by auto
    have 0 ∈ closed_segment (-d) e
    using ⟨0 < d⟩ ⟨0 < e⟩ closed_segment_eq_real_ivl by auto
    then have a ∈ path_image (linepath (a - d) (a + e))
    using of_real_closed_segment [THEN iffD2]
    by (force dest: closed_segment_translation_eq [of a, THEN iffD2] simp del:
      of_real_closed_segment)
    then have a ∉ path_image p
    using ⟨0 < d⟩ ⟨0 < e⟩ pap by auto
    then obtain k where 0 < k and k: ball a k ∩ (path_image p) = {}
    using ⟨0 < e⟩ ⟨path p⟩ not_on_path_ball by blast
    define kde where kde ≡ (min k (min d e)) / 2
    have 0 < kde kde < k kde < d kde < e
    using ⟨0 < k⟩ ⟨0 < d⟩ ⟨0 < e⟩ by (auto simp: kde_def)
    let ?q = linepath (a + kde) (a + e) +++ p +++ linepath (a - d) (a - kde)
    have -kde ∈ closed_segment (-d) e
    using ⟨0 < kde⟩ ⟨kde < d⟩ ⟨kde < e⟩ closed_segment_eq_real_ivl by auto
    then have a_diff_kde: a - kde ∈ closed_segment (a - d) (a + e)
    using of_real_closed_segment [THEN iffD2]
    by (force dest: closed_segment_translation_eq [of a, THEN iffD2] simp del:
      of_real_closed_segment)
    then have clsub2: closed_segment (a - d) (a - kde) ⊆ closed_segment (a -
      d) (a + e)

```

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  by (simp add: subset_closed_segment)
  then have path_image p  $\cap$  closed_segment (a - d) (a - kde)  $\subseteq$  {a + e, a -
d}
  using pap by force
  moreover
  have a + e  $\notin$  path_image p  $\cap$  closed_segment (a - d) (a - kde)
  using <0 < kde> <kde < d> <0 < e> by (auto simp: closed_segment_eq_real_ivl)
  ultimately have sub_a_diff_d: path_image p  $\cap$  closed_segment (a - d) (a -
kde)  $\subseteq$  {a - d}
  by blast
  have kde  $\in$  closed_segment (-d) e
  using <0 < kde> <kde < d> <kde < e> closed_segment_eq_real_ivl by auto
  then have a_diff_kde: a + kde  $\in$  closed_segment (a - d) (a + e)
  using of_real_closed_segment [THEN iffD2]
  by (force dest: closed_segment_translation_eq [of a, THEN iffD2] simp del:
of_real_closed_segment)
  then have clsub1: closed_segment (a + kde) (a + e)  $\subseteq$  closed_segment (a -
d) (a + e)
  by (simp add: subset_closed_segment)
  then have closed_segment (a + kde) (a + e)  $\cap$  path_image p  $\subseteq$  {a + e, a -
d}
  using pap by force
  moreover
  have closed_segment (a + kde) (a + e)  $\cap$  closed_segment (a - d) (a - kde) =
{}
  proof (clarsimp intro!: equals0I)
    fix y
    assume y1: y  $\in$  closed_segment (a + kde) (a + e)
    and y2: y  $\in$  closed_segment (a - d) (a - kde)
    obtain u::real where u: y = a + u and 0 < u
    proof -
      obtain  $\xi$  where  $\xi$ : y = (1 -  $\xi$ ) *R (a + kde) +  $\xi$  *R (a + e) and 0  $\leq$   $\xi$   $\xi$ 
 $\leq$  1
      using y1 by (auto simp: in_segment)
    show thesis
    proof
      show y = a + ((1 -  $\xi$ )*kde +  $\xi$ *e)
      using  $\xi$  by (auto simp: scaleR_conv_of_real algebra_simps)
      have (1 -  $\xi$ )*kde +  $\xi$ *e  $\geq$  0
      using <0 < kde> <0  $\leq$   $\xi\xi$   $\leq$  1> <0 < e> by auto
      moreover have (1 -  $\xi$ )*kde +  $\xi$ *e  $\neq$  0
      using <0 < kde> <0  $\leq$   $\xi\xi$   $\leq$  1> <0 < e> by (auto simp: add_nonneg_eq_0_iff)
      ultimately show (1 -  $\xi$ )*kde +  $\xi$ *e > 0 by simp
    qed
  qed
  moreover
  obtain v::real where v: y = a + v and v  $\leq$  0
  proof -
    obtain  $\xi$  where  $\xi$ : y = (1 -  $\xi$ ) *R (a - d) +  $\xi$  *R (a - kde) and 0  $\leq$   $\xi$   $\xi$ 

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≤ 1
  using y2 by (auto simp: in_segment)
show thesis
proof
  show  $y = a + (-((1 - \xi)*d + \xi*kde))$ 
    using  $\xi$  by (force simp: scaleR_conv_of_real algebra_simps)
  show  $-((1 - \xi)*d + \xi*kde) \leq 0$ 
    using  $\langle 0 < kde \rangle \langle 0 \leq \xi \rangle \langle \xi \leq 1 \rangle \langle 0 < d \rangle$ 
  by (metis add.left_neutral add_nonneg_nonneg le_diff_eq less_eq_real_def
neg_le_0_iff_le_split_mult_pos_le)

  qed
  qed
  ultimately show False
  by auto
  qed
  moreover have  $a - d \notin \text{closed\_segment } (a + kde) (a + e)$ 
    using  $\langle 0 < kde \rangle \langle kde < d \rangle \langle 0 < e \rangle$  by (auto simp: closed_segment_eq_real_ivl)
  ultimately have sub_a_plus_e:
     $\text{closed\_segment } (a + kde) (a + e) \cap (\text{path\_image } p \cup \text{closed\_segment } (a - d) (a - kde)) \subseteq \{a + e\}$ 
  by auto
  have  $kde \in \text{closed\_segment } (-kde) e$ 
    using  $\langle 0 < kde \rangle \langle kde < d \rangle \langle kde < e \rangle$  closed_segment_eq_real_ivl by auto
  then have a_add_kde:  $a + kde \in \text{closed\_segment } (a - kde) (a + e)$ 
    using of_real_closed_segment [THEN iffD2]
  by (force dest: closed_segment_translation_eq [of a, THEN iffD2] simp del:
of_real_closed_segment)
  have  $\text{closed\_segment } (a - kde) (a + kde) \cap \text{closed\_segment } (a + kde) (a + e) = \{a + kde\}$ 
  by (metis a_add_kde Int_closed_segment)
  moreover
  have  $\text{path\_image } p \cap \text{closed\_segment } (a - kde) (a + kde) = \{\}$ 
  proof (rule equals0I, clarify)
    fix y assume  $y \in \text{path\_image } p \ y \in \text{closed\_segment } (a - kde) (a + kde)$ 
    with equals0D [OF k, of y]  $\langle 0 < kde \rangle \langle kde < k \rangle$  show False
    by (auto simp: dist_norm dest: dist_decreases_closed_segment [where c=a])
  qed
  moreover
  have  $-kde \in \text{closed\_segment } (-d) kde$ 
    using  $\langle 0 < kde \rangle \langle kde < d \rangle \langle kde < e \rangle$  closed_segment_eq_real_ivl by auto
  then have a_diff_kde':  $a - kde \in \text{closed\_segment } (a - d) (a + kde)$ 
    using of_real_closed_segment [THEN iffD2]
  by (force dest: closed_segment_translation_eq [of a, THEN iffD2] simp del:
of_real_closed_segment)
  then have  $\text{closed\_segment } (a - d) (a - kde) \cap \text{closed\_segment } (a - kde) (a + kde) = \{a - kde\}$ 
  by (metis Int_closed_segment)
  ultimately

```

```

have pa_subset_pm_kde: path_image ?q  $\cap$  closed_segment (a - kde) (a + kde)
 $\subseteq$  {a - kde, a + kde}
by (auto simp: path_image_join assms)
have ge_kde1:  $\exists y. x = a + \text{of\_real } y \wedge y \geq \text{kde}$  if x: x  $\in$  closed_segment (a +
kde) (a + e) for x
proof -
  obtain u where  $0 \leq u \wedge u \leq 1$  and u: x = ( $1 - u$ ) *R (a + kde) + u *R (a
+ e)
  using x by (auto simp: in_segment)
  then have  $\text{kde} \leq (1 - u) * \text{kde} + u * e$ 
  using  $\langle \text{kde} < e \rangle$  segment_bound_lemma by auto
  moreover have x = a + ( $(1 - u) * \text{kde} + u * e$ )
  by (fastforce simp: u algebra_simps scaleR_conv_of_real)
  ultimately
  show ?thesis by blast
qed
have ge_kde2:  $\exists y. x = a + \text{of\_real } y \wedge y \leq -\text{kde}$  if x: x  $\in$  closed_segment (a
- d) (a - kde) for x
proof -
  obtain u where  $0 \leq u \wedge u \leq 1$  and u: x = ( $1 - u$ ) *R (a - d) + u *R (a -
kde)
  using x by (auto simp: in_segment)
  then have  $\text{kde} \leq ((1 - u) * d + u * \text{kde})$ 
  using  $\langle \text{kde} < d \rangle$  segment_bound_lemma by auto
  moreover have x = a - ( $(1 - u) * d + u * \text{kde}$ )
  by (fastforce simp: u algebra_simps scaleR_conv_of_real)
  ultimately show ?thesis
  by (metis add_uminus_conv_diff neg_le_iff_le of_real_minus)
qed
have notin_paq: x  $\notin$  path_image ?q if dist a x < kde for x
proof -
  have x  $\notin$  path_image p
  using k  $\langle \text{kde} < k \rangle$  that by force
  then show ?thesis
  using that assms by (auto simp: path_image_join dist_norm dest!: ge_kde1
ge_kde2)
qed
obtain z where zin: z  $\in$  inside (path_image (?q +++ linepath (a - kde) (a +
kde)))
  and z1: cmod (winding_number (?q +++ linepath (a - kde) (a + kde))
z) = 1
proof (rule simple_closed_path_wn1 [of kde ?q a])
  show simple_path (?q +++ linepath (a - kde) (a + kde))
  proof (intro simple_path_join_loop conjI)
  show arc ?q
  proof (rule arc_join)
  show arc (linepath (a + kde) (a + e))
  using  $\langle \text{kde} < e \rangle$   $\langle \text{arc } p \rangle$  by (force simp: pfp)
  show arc (p +++ linepath (a - d) (a - kde))

```

```

    using ‹kde < d› ‹kde < e› ‹arc p› sub_a_diff_d by (force simp: pfp intro:
arc_join)
    qed (auto simp: psp pfp path_image_join sub_a_plus_e)
    show arc (linepath (a - kde) (a + kde))
      using ‹0 < kde› by auto
    qed (use pa_subset_pm_kde in auto)
    qed (use ‹0 < kde› notin_paq in auto)
    have eq: path_image (?q +++ linepath (a - kde) (a + kde)) = path_image (p
+++ linepath (a - d) (a + e))
      (is ?lhs = ?rhs)
  proof
    show ?lhs ⊆ ?rhs
      using clsub1 clsub2 apply (auto simp: path_image_join assms)
      by (meson subsetCE subset_closed_segment)
    show ?rhs ⊆ ?lhs
      apply (simp add: path_image_join assms Un_ac)
      by (metis Un_closed_segment Un_assoc a_diff_kde a_diff_kde' le_supI2
subset_refl)
    qed
    show thesis
  proof
    show zzin: z ∈ inside (path_image (p +++ linepath (a - d) (a + e)))
      by (metis eq zin)
    then have znotin: z ∉ path_image p
      by (metis ComplD Un_iff inside_Un_outside path_image_join pathfin-
ish_linepath pathstart_reversepath pfp reversepath_linepath)
    have znotin_d_kde: z ∉ closed_segment (a - d) (a + kde)
      by (metis ComplD Un_iff Un_closed_segment a_diff_kde inside_Un_outside
path_image_join path_image_linepath pathstart_linepath pfp zzin)
    have znotin_d_e: z ∉ closed_segment (a - d) (a + e)
      by (metis IntI UnCI emptyE inside_no_overlap path_image_join path_image_linepath
pathstart_linepath pfp zzin)
    have znotin_kde_e: z ∉ closed_segment (a + kde) (a + e) and znotin_d_kde':
z ∉ closed_segment (a - d) (a - kde)
      using clsub1 clsub2 zzin
      by (metis (no_types, opaque_lifting) IntI Un_iff emptyE inside_no_overlap
path_image_join path_image_linepath pathstart_linepath pfp subsetD)+
    have winding_number (linepath (a - d) (a + e)) z =
      winding_number (linepath (a - d) (a + kde)) z + winding_number
(linepath (a + kde) (a + e)) z
    proof (rule winding_number_split_linepath)
      show a + complex_of_real kde ∈ closed_segment (a - d) (a + e)
        by (simp add: a_diff_kde)
      show z ∉ closed_segment (a - d) (a + e)
        by (metis ComplD Un_iff inside_Un_outside path_image_join path_image_linepath
pathstart_linepath pfp zzin)
    qed
    also have ... = winding_number (linepath (a + kde) (a + e)) z +
      (winding_number (linepath (a - d) (a - kde)) z + winding_number

```

```

(linepath (a - kde) (a + kde)) z)
  by (simp add: winding_number_split_linepath [of a-kde, symmetric] znotin_d_kde
a_diff_kde')
  finally have winding_number (p +++ linepath (a - d) (a + e)) z =
    winding_number p z + winding_number (linepath (a + kde) (a +
e)) z +
      (winding_number (linepath (a - d) (a - kde)) z +
winding_number (linepath (a - kde) (a + kde)) z)
  by (metis (no_types, lifting) ComplD Un_iff ‹arc p› add.assoc arc_imp_path
eq_path_image_join_path_join_path_ends_path_linepath_simple_path_imp_path
sp_pl union_with_outside winding_number_join zin)
  also have ... = winding_number (linepath (a + kde) (a + e)) z
    + winding_number (p +++ linepath (a - d) (a - kde)) z
    + winding_number (linepath (a - kde) (a + kde)) z
  using ‹path p› znotin assms
  by simp (metis Un_iff Un_closed_segment a_diff_kde' path_image_linepath
path_linepath pathstart_linepath winding_number_join znotin_d_kde)
  also have ... = winding_number ?q z + winding_number (linepath (a - kde)
(a + kde)) z
  using ‹path p› znotin assms by (simp add: path_image_join winding_number_join
znotin_kde_e znotin_d_kde')
  also have ... = winding_number (?q +++ linepath (a - kde) (a + kde)) z
  by (metis (mono_tags, lifting) ComplD UnCI ‹path p› eq_inside_outside
path_image_join_path_join_eq_path_linepath_pathfinish_join_pathfinish_linepath
pathstart_join_pathstart_linepath_pfp psp winding_number_join zzin)
  finally have winding_number (p +++ linepath (a - d) (a + e)) z =
    winding_number (?q +++ linepath (a - kde) (a + kde)) z .
  then show cmod (winding_number (p +++ linepath (a - d) (a + e)) z) = 1
  by (simp add: z1)
qed
qed

```

lemma *simple_closed_path_wn3*:

fixes $p :: \text{real} \Rightarrow \text{complex}$

assumes *simple_path p* **and** *loop: pathfinish p = pathstart p*

obtains z **where** $z \in \text{inside} (\text{path_image } p)$ $\text{cmod} (\text{winding_number } p \ z) = 1$

proof –

have $\text{ins}: \text{inside}(\text{path_image } p) \neq \{\}$ $\text{open}(\text{inside}(\text{path_image } p))$

$\text{connected}(\text{inside}(\text{path_image } p))$

and $\text{out}: \text{outside}(\text{path_image } p) \neq \{\}$ $\text{open}(\text{outside}(\text{path_image } p))$

$\text{connected}(\text{outside}(\text{path_image } p))$

and $\text{bo}: \text{bounded}(\text{inside}(\text{path_image } p)) \neg \text{bounded}(\text{outside}(\text{path_image } p))$

and $\text{ins_out}: \text{inside}(\text{path_image } p) \cap \text{outside}(\text{path_image } p) = \{\}$

$\text{inside}(\text{path_image } p) \cup \text{outside}(\text{path_image } p) = - \text{path_image } p$

and $\text{fro}: \text{frontier}(\text{inside}(\text{path_image } p)) = \text{path_image } p$

$\text{frontier}(\text{outside}(\text{path_image } p)) = \text{path_image } p$

using *Jordan_inside_outside* [*OF assms*] **by** *auto*

obtain a **where** $a \in \text{inside}(\text{path_image } p)$

using ‹ $\text{inside} (\text{path_image } p) \neq \{\}$ › **by** *blast*

```

obtain  $d::\text{real}$  where  $0 < d$  and  $d\_fro: a - d \in \text{frontier } (\text{inside } (\text{path\_image } p))$ 
and  $d\_int: \bigwedge \varepsilon. \llbracket 0 \leq \varepsilon; \varepsilon < d \rrbracket \implies (a - \varepsilon) \in \text{inside } (\text{path\_image } p)$ 
using  $\text{ray\_to\_frontier } [\text{of } \text{inside } (\text{path\_image } p) \ a \ -1] \ \text{bo } \text{ins } a \ \text{interior\_eq}$ 
by  $(\text{metis } \text{ab\_group\_add\_class.ab\_diff\_conv\_add\_uminus } \text{of\_real\_def } \text{scale\_minus\_right}$ 
 $\text{zero\_neq\_neq\_one})$ 
obtain  $e::\text{real}$  where  $0 < e$  and  $e\_fro: a + e \in \text{frontier } (\text{inside } (\text{path\_image } p))$ 
and  $e\_int: \bigwedge \varepsilon. \llbracket 0 \leq \varepsilon; \varepsilon < e \rrbracket \implies (a + \varepsilon) \in \text{inside } (\text{path\_image } p)$ 
using  $\text{ray\_to\_frontier } [\text{of } \text{inside } (\text{path\_image } p) \ a \ 1] \ \text{bo } \text{ins } a \ \text{interior\_eq}$ 
by  $(\text{metis } \text{of\_real\_def } \text{zero\_neq\_one})$ 
obtain  $t0$  where  $0 \leq t0$   $t0 \leq 1$  and  $pt: p \ t0 = a - d$ 
using  $a \ d \_fro \_fro$  by  $(\text{auto } \text{simp: } \text{path\_image\_def})$ 
obtain  $q$  where  $\text{simple\_path } q$  and  $q\_ends: \text{pathstart } q = a - d \ \text{pathfinish } q =$ 
 $a - d$ 
and  $q\_eq\_p: \text{path\_image } q = \text{path\_image } p$ 
and  $wn\_q\_eq\_wn\_p: \bigwedge z. z \in \text{inside}(\text{path\_image } p) \implies \text{winding\_number } q \ z$ 
 $= \text{winding\_number } p \ z$ 
proof
show  $\text{simple\_path } (\text{shiftpath } t0 \ p)$ 
by  $(\text{simp } \text{add: } \text{pathstart\_shiftpath } \text{pathfinish\_shiftpath}$ 
 $\ \text{simple\_path\_shiftpath } \text{path\_image\_shiftpath } \langle 0 \leq t0 \rangle \langle t0 \leq 1 \rangle \ \text{assms})$ 
show  $\text{pathstart } (\text{shiftpath } t0 \ p) = a - d$ 
using  $pt$  by  $(\text{simp } \text{add: } \langle t0 \leq 1 \rangle \ \text{pathstart\_shiftpath})$ 
show  $\text{pathfinish } (\text{shiftpath } t0 \ p) = a - d$ 
by  $(\text{simp } \text{add: } \langle 0 \leq t0 \rangle \ \text{loop } \text{pathfinish\_shiftpath } pt)$ 
show  $\text{path\_image } (\text{shiftpath } t0 \ p) = \text{path\_image } p$ 
by  $(\text{simp } \text{add: } \langle 0 \leq t0 \rangle \langle t0 \leq 1 \rangle \ \text{loop } \text{path\_image\_shiftpath})$ 
show  $\text{winding\_number } (\text{shiftpath } t0 \ p) \ z = \text{winding\_number } p \ z$ 
if  $z \in \text{inside } (\text{path\_image } p)$  for  $z$ 
by  $(\text{metis } \text{ComplD } \text{Un\_iff } \langle 0 \leq t0 \rangle \langle t0 \leq 1 \rangle \langle \text{simple\_path } p \rangle \ \text{atLeastAt}$ 
 $\ \text{Most\_iff } \text{inside\_Un\_outside}$ 
 $\ \text{loop } \text{simple\_path\_imp\_path } \text{that } \text{winding\_number\_shiftpath})$ 
qed
have  $ad\_not\_ae: a - d \neq a + e$ 
by  $(\text{metis } \langle 0 < d \rangle \langle 0 < e \rangle \ \text{add.left\_inverse } \text{add\_left\_cancel } \text{add\_uminus\_conv\_diff}$ 
 $\ \text{le\_add\_same\_cancel2 } \text{less\_eq\_real\_def } \text{not\_less\_of\_real\_add } \text{of\_real\_def}$ 
 $\ \text{of\_real\_eq\_0\_iff } pt)$ 
have  $ad\_ae\_q: \{a - d, a + e\} \subseteq \text{path\_image } q$ 
using  $\langle \text{path\_image } q = \text{path\_image } p \rangle \ d\_fro \ e\_fro \_fro(1)$  by  $\text{auto}$ 
have  $ada: \text{open\_segment } (a - d) \ a \subseteq \text{inside } (\text{path\_image } p)$ 
proof  $(\text{clarsimp } \text{simp: } \text{in\_segment})$ 
fix  $u::\text{real}$  assume  $0 < u$   $u < 1$ 
with  $d\_int$  have  $a - (1 - u) * d \in \text{inside } (\text{path\_image } p)$ 
by  $(\text{metis } \langle 0 < d \rangle \ \text{add.commute } \text{diff\_add\_cancel } \text{left\_diff\_distrib' } \text{less\_add\_same\_cancel2}$ 
 $\ \text{less\_eq\_real\_def } \text{mult.left\_neutral } \text{zero\_less\_mult\_iff})$ 
then show  $(1 - u) *_R (a - d) + u *_R a \in \text{inside } (\text{path\_image } p)$ 
by  $(\text{simp } \text{add: } \text{diff\_add\_eq } \text{of\_real\_def } \text{real\_vector.scale\_right\_diff\_distrib})$ 
qed

```



```

have aae: open_segment a (a + e)  $\subseteq$  inside (path_image p)
proof (clarsimp simp: in_segment)
  fix u::real assume 0 < u u < 1
  with e_int have a + u * e  $\in$  inside (path_image p)
    by (meson <0 < e> less_eq_real_def mult_less_cancel_right2 not_less
zero_less_mult_iff)
  then show (1 - u) *R a + u *R (a + e)  $\in$  inside (path_image p)
    by (metis (mono_tags, lifting) add.assoc of_real_mult pth_6 scaleR_collapse
scaleR_conv_of_real)
  qed
have complex_of_real (d * d + (e * e + d * (e + e)))  $\neq$  0
  using ad_not_ae
  by (metis <0 < d> <0 < e> add_strict_left_mono less_add_same_cancel1
not_sum_squares_lt_zero
of_real_eq_0_iff zero_less_double_add_iff_zero_less_single_add zero_less_mult_iff)
  moreover have  $\exists u > 0. u < 1 \wedge d = u * d + u * e$ 
  using <0 < d> <0 < e>
  by (rule_tac x=d / (d+e) in exI) (simp add: divide_simps scaleR_conv_of_real
flip: distrib_left)
  ultimately have a_in_de: a  $\in$  open_segment (a - d) (a + e)
  using ad_not_ae by (simp add: in_segment algebra_simps scaleR_conv_of_real
flip: of_real_add of_real_mult)
  then have open_segment (a - d) (a + e)  $\subseteq$  inside (path_image p)
  using ada a aae Un_open_segment [of a a-d a+e] by blast
  then have path_image q  $\cap$  open_segment (a - d) (a + e) = {}
  using inside_no_overlap by (fastforce simp: q_eq_p)
  with ad_ae_q have paq_Int_cs: path_image q  $\cap$  closed_segment (a - d) (a +
e) = {a - d, a + e}
  by (simp add: closed_segment_eq_open)
  obtain t where 0  $\leq$  t t  $\leq$  1 and qt: q t = a + e
  using a_e_fro ad_ae_q by (auto simp: path_defs)
  then have t  $\neq$  0
  by (metis ad_not_ae pathstart_def q_ends(1))
  then have t  $\neq$  1
  by (metis ad_not_ae pathfinish_def q_ends(2) qt)
  have q01: q 0 = a - d q 1 = a + e
  using q_ends by (auto simp: pathstart_def pathfinish_def)
  obtain z where zin: z  $\in$  inside (path_image (subpath t 0 q +++ linepath (a -
d) (a + e)))
  and z1: cmod (winding_number (subpath t 0 q +++ linepath (a - d)
(a + e)) z) = 1
  proof (rule simple_closed_path_wn2 [of d e subpath t 0 q a], simp_all add: q01)
  show simple_path (subpath t 0 q +++ linepath (a - d) (a + e))
  proof (rule simple_path_join_loop, simp_all add: qt q01)
  have inj_on q (closed_segment t 0)
  using <0  $\leq$  t> <simple_path q> <t  $\leq$  1> <t  $\neq$  0> <t  $\neq$  1>
  by (fastforce simp: simple_path_def loop_free_def inj_on_def closed_segment_eq_real_ivl)
  then show arc (subpath t 0 q)
  using <0  $\leq$  t> <simple_path q> <t  $\leq$  1> <t  $\neq$  0>

```

```

    by (simp add: arc_subpath_eq simple_path_imp_path)
  show arc (linepath (a - d) (a + e))
    by (simp add: ad_not_ae)
  show path_image (subpath t 0 q) ∩ closed_segment (a - d) (a + e) ⊆ {a +
e, a - d}
    using qt paq_Int_cs ⟨simple_path q⟩ ⟨0 ≤ t⟩ ⟨t ≤ 1⟩
    by (force simp: dest: rev_subsetD [OF _ path_image_subpath_subset] intro:
simple_path_imp_path)
  qed
  qed (auto simp: ⟨0 < d⟩ ⟨0 < e⟩ qt)
  have pa01_Un: path_image (subpath 0 t q) ∪ path_image (subpath 1 t q) =
path_image q
    unfolding path_image_subpath
    using ⟨0 ≤ t⟩ ⟨t ≤ 1⟩ by (force simp: path_image_def image_iff)
  with paq_Int_cs have pa_01q:
    (path_image (subpath 0 t q) ∪ path_image (subpath 1 t q)) ∩ closed_segment
(a - d) (a + e) = {a - d, a + e}
    by metis
  have z_notin_ed: z ∉ closed_segment (a + e) (a - d)
    using zin q01 by (simp add: path_image_join closed_segment_commute in-
side_def)
  have z_notin_0t: z ∉ path_image (subpath 0 t q)
    by (metis (no_types, opaque_lifting) IntI Un_upper1 subsetD empty_iff in-
side_no_overlap path_image_join
    path_image_subpath_commute pathfinish_subpath pathstart_def pathstart_linepath
q_ends(1) qt subpath_trivial zin)
  have [simp]: - winding_number (subpath t 0 q) z = winding_number (subpath
0 t q) z
    by (metis ⟨0 ≤ t⟩ ⟨simple_path q⟩ ⟨t ≤ 1⟩ atLeastAtMost_iff zero_le_one
    path_image_subpath_commute path_subpath_real_eq_0_iff_le_ge_0
    reversepath_subpath simple_path_imp_path winding_number_reversepath
z_notin_0t)
  obtain z_in_q: z ∈ inside(path_image q)
    and wn_q: winding_number (subpath 0 t q) +++ subpath t 1 q) z = -
winding_number (subpath t 0 q) +++ linepath (a - d) (a + e)) z
  proof (rule winding_number_from_innerpath
    [of subpath 0 t q a-d a+e subpath 1 t q linepath (a - d) (a + e)
    z - winding_number (subpath t 0 q) +++ linepath (a - d) (a + e)) z],
    simp_all add: q01 qt pa01_Un reversepath_subpath)
  show simple_path (subpath 0 t q) simple_path (subpath 1 t q)
    by (simp_all add: ⟨0 ≤ t⟩ ⟨simple_path q⟩ ⟨t ≤ 1⟩ ⟨t ≠ 0⟩ ⟨t ≠ 1⟩ sim-
ple_path_subpath)
  show simple_path (linepath (a - d) (a + e))
    using ad_not_ae by blast
  show path_image (subpath 0 t q) ∩ path_image (subpath 1 t q) = {a - d, a +
e} (is ?lhs = ?rhs)
  proof
    show ?lhs ⊆ ?rhs
      using ⟨0 ≤ t⟩ ⟨simple_path q⟩ ⟨t ≤ 1⟩ ⟨t ≠ 1⟩ q_ends qt q01

```

```

    by (force simp: pathfinish_def qt simple_path_def loop_free_def path_image_subpath)
    show ?rhs  $\subseteq$  ?lhs
      using  $\langle 0 \leq t \rangle \langle t \leq 1 \rangle$  q01 qt by (force simp: path_image_subpath)
  qed
  show path_image (subpath 0 t q)  $\cap$  closed_segment (a - d) (a + e) = {a -
d, a + e} (is ?lhs = ?rhs)
  proof
    show ?lhs  $\subseteq$  ?rhs using paq_Int_cs pa01_Un by fastforce
    show ?rhs  $\subseteq$  ?lhs using  $\langle 0 \leq t \rangle \langle t \leq 1 \rangle$  q01 qt by (force simp: path_image_subpath)
  qed
  show path_image (subpath 1 t q)  $\cap$  closed_segment (a - d) (a + e) = {a -
d, a + e} (is ?lhs = ?rhs)
  proof
    show ?lhs  $\subseteq$  ?rhs by (auto simp: pa_01q [symmetric])
    show ?rhs  $\subseteq$  ?lhs using  $\langle 0 \leq t \rangle \langle t \leq 1 \rangle$  q01 qt by (force simp: path_image_subpath)
  qed
  show closed_segment (a - d) (a + e)  $\cap$  inside (path_image q)  $\neq$  {}
    using a a_in_de open_closed_segment pa01_Un q_eq_p by fastforce
  show  $z \in$  inside (path_image (subpath 0 t q)  $\cup$  closed_segment (a - d) (a +
e))
    by (metis path_image_join path_image_linepath path_image_subpath_commute
pathfinish_subpath pathstart_linepath q01(1) zin)
  show winding_number (subpath 0 t q +++ linepath (a + e) (a - d)) z =
- winding_number (subpath t 0 q +++ linepath (a - d) (a + e)) z
    using z_notin_ed z_notin_0t  $\langle 0 \leq t \rangle \langle \text{simple\_path } q \rangle \langle t \leq 1 \rangle$ 
    by (simp add: simple_path_imp_path qt q01 path_image_subpath_commute
closed_segment_commute winding_number_join winding_number_reversepath [symmetric])
  show  $-d \neq e$ 
    using ad_not_ae by auto
  show winding_number (subpath t 0 q +++ linepath (a - d) (a + e)) z  $\neq$  0
    using z1 by auto
  qed
  show ?thesis
  proof
    show  $z \in$  inside (path_image p)
      using q_eq_p z_in_q by auto
    then have [simp]:  $z \notin$  path_image q
      by (metis disjoint_iff_not_equal inside_no_overlap q_eq_p)
    have [simp]:  $z \notin$  path_image (subpath 1 t q)
      using inside_def pa01_Un z_in_q by fastforce
    have winding_number (subpath 0 t q +++ subpath t 1 q) z = winding_number (subpath
0 1 q) z
      using z_notin_0t  $\langle 0 \leq t \rangle \langle \text{simple\_path } q \rangle \langle t \leq 1 \rangle$ 
      by (simp add: simple_path_imp_path qt path_image_subpath_commute wind-
ing_number_join winding_number_subpath_combine)
    with wn_q have winding_number (subpath t 0 q +++ linepath (a - d) (a +
e)) z = - winding_number q z
      by auto
    with z1 have cmod (winding_number q z) = 1

```

```

    by simp
  with z1 wn_q_eq_wn_p show cmod (winding_number p z) = 1
    using z1 wn_q_eq_wn_p by (simp add: ⟨z ∈ inside (path_image p)⟩)
  qed
qed

proposition simple_closed_path_winding_number_inside:
  assumes simple_path γ
  obtains  $\bigwedge z. z \in \text{inside}(\text{path\_image } \gamma) \implies \text{winding\_number } \gamma z = 1$ 
    |  $\bigwedge z. z \in \text{inside}(\text{path\_image } \gamma) \implies \text{winding\_number } \gamma z = -1$ 
proof (cases pathfinish γ = pathstart γ)
  case True
  have path γ
    by (simp add: assms simple_path_imp_path)
  then have const: winding_number γ constant_on inside(path_image γ)
  proof (rule winding_number_constant)
    show connected (inside(path_image γ))
      by (simp add: Jordan_inside_outside True assms)
  qed (use inside_no_overlap True in auto)
  obtain z where zin: z ∈ inside (path_image γ) and z1: cmod (winding_number
γ z) = 1
    using simple_closed_path_wn3 [of γ] True assms by blast
  have winding_number γ z ∈ ℤ
    using zin integer_winding_number [OF ⟨path γ⟩ True] inside_def by blast
  moreover have real_of_int x = - 1  $\longleftrightarrow$  x = -1 for x
    by linarith
  ultimately consider winding_number γ z = 1 | winding_number γ z = -1
    using z1 by (auto simp: Ints_def abs_if_split: if_split_asm)
  with that const zin show ?thesis
    unfolding constant_on_def by metis
next
  case False
  then show ?thesis
    using inside_simple_curve_imp_closed assms that(2) by blast
qed

lemma simple_closed_path_abs_winding_number_inside:
  assumes simple_path γ z ∈ inside(path_image γ)
  shows |Re (winding_number γ z)| = 1
  by (metis assms norm_minus_cancel norm_one one_complex.simps(1) real_norm_def
simple_closed_path_winding_number_inside uminus_complex.simps(1))

lemma simple_closed_path_norm_winding_number_inside:
  assumes simple_path γ z ∈ inside(path_image γ)
  shows norm (winding_number γ z) = 1
proof -
  have pathfinish γ = pathstart γ
    using assms inside_simple_curve_imp_closed by blast
  with assms integer_winding_number have winding_number γ z ∈ ℤ

```

```

  by (simp add: inside_def simple_path_def)
  then show ?thesis
  by (metis assms norm_minus_cancel norm_one simple_closed_path_winding_number_inside)
qed

```

lemma *simple_closed_path_winding_number_cases*:

```

  assumes simple_path  $\gamma$  pathfinish  $\gamma = \text{pathstart } \gamma$   $z \notin \text{path\_image } \gamma$ 
  shows winding_number  $\gamma$   $z \in \{-1, 0, 1\}$ 
  proof -
    consider  $z \in \text{inside } (\text{path\_image } \gamma) \mid z \in \text{outside } (\text{path\_image } \gamma)$ 
    by (metis ComplI UnE assms(3) inside_Un_outside)
    then show ?thesis
    proof cases
      case 1
      then show ?thesis
      using assms simple_closed_path_winding_number_inside by auto
    next
      case 2
      then show ?thesis
      using assms simple_path_def winding_number_zero_in_outside by blast
    qed
  qed

```

lemma *simple_closed_path_winding_number_pos*:

```

  [[simple_path  $\gamma$ ; pathfinish  $\gamma = \text{pathstart } \gamma$ ;  $z \notin \text{path\_image } \gamma$ ;  $0 < \text{Re}(\text{winding\_number } \gamma \ z)$ ]]
   $\implies \text{winding\_number } \gamma \ z = 1$ 
  using simple_closed_path_winding_number_cases
  by fastforce

```

3.9 Winding number for rectangular paths

proposition *winding_number_rectpath*:

```

  assumes  $z \in \text{box } a1 \ a3$ 
  shows winding_number (rectpath  $a1 \ a3$ )  $z = 1$ 
  proof -
    from assms have less:  $\text{Re } a1 < \text{Re } a3 \ \text{Im } a1 < \text{Im } a3$ 
    by (auto simp: in_box_complex_iff)
    define  $a2 \ a4$  where  $a2 = \text{Complex } (\text{Re } a3) (\text{Im } a1)$  and  $a4 = \text{Complex } (\text{Re } a1) (\text{Im } a3)$ 
    let  $?l1 = \text{linepath } a1 \ a2$  and  $?l2 = \text{linepath } a2 \ a3$ 
    and  $?l3 = \text{linepath } a3 \ a4$  and  $?l4 = \text{linepath } a4 \ a1$ 
    from assms and less have  $z \notin \text{path\_image } (\text{rectpath } a1 \ a3)$ 
    by (auto simp: path_image_rectpath_cbox_minus_box)
    also have path_image (rectpath  $a1 \ a3$ ) =
      path_image ?l1  $\cup$  path_image ?l2  $\cup$  path_image ?l3  $\cup$  path_image
    ?l4
    by (simp add: rectpath_def Let_def path_image_join Un_assoc a2_def a4_def)
    finally have  $z \notin \dots$ 

```

```

moreover have  $\forall l \in \{?l1, ?l2, ?l3, ?l4\}. \text{Re } (\text{winding\_number } l z) > 0$ 
unfolding ball_simps HOL.simp_thms a2_def a4_def
by (intro conjI; (rule winding_number_linepath_pos_lt;
  (insert assms, auto simp: a2_def a4_def in_box_complex_iff_mult_neg_neg))+)
ultimately have  $\text{Re } (\text{winding\_number } (\text{rectpath } a1 a3) z) > 0$ 
by (simp add: winding_number_join path_image_join rectpath_def Let_def
a2_def a4_def)
thus winding_number (rectpath a1 a3) z = 1 using assms less
by (intro simple_closed_path_winding_number_pos simple_path_rectpath)
  (auto simp: path_image_rectpath_cbox_minus_box)
qed

```

```

proposition winding_number_rectpath_outside:
assumes  $\text{Re } a1 \leq \text{Re } a3 \text{ Im } a1 \leq \text{Im } a3$ 
assumes  $z \notin \text{cbox } a1 a3$ 
shows winding_number (rectpath a1 a3) z = 0
using assms by (intro winding_number_zero_outside[OF ___ assms(3)]
  path_image_rectpath_subset_cbox) simp_all

```

A per-function version for continuous logs, a kind of monodromy

```

proposition winding_number_compose_exp:
assumes path p
shows winding_number (exp o p) 0 = (pathfinish p - pathstart p) / (2 * of_real
pi * i)
proof -
obtain e where  $0 < e$  and  $e: \bigwedge t. t \in \{0..1\} \implies e \leq \text{norm}(\text{exp}(p t))$ 
proof
have closed (path_image (exp o p))
by (simp add: assms closed_path_image holomorphic_on_exp holomor-
phic_on_imp_continuous_on path_continuous_image)
then show  $0 < \text{setdist } \{0\} (\text{path\_image } (\text{exp } \circ p))$ 
by (metis exp_not_eq_zero imageE image_comp infdist_eq_setdist infdist_pos_not_in_closed
path_defs(4) path_image_nonempty)
next
fix t::real
assume  $t \in \{0..1\}$ 
have  $\text{setdist } \{0\} (\text{path\_image } (\text{exp } \circ p)) \leq \text{dist } 0 (\text{exp } (p t))$ 
proof (rule setdist_le_dist)
show  $\text{exp } (p t) \in \text{path\_image } (\text{exp } \circ p)$ 
using  $\langle t \in \{0..1\} \rangle$  path_image_def by fastforce+
qed auto
then show  $\text{setdist } \{0\} (\text{path\_image } (\text{exp } \circ p)) \leq \text{cmod } (\text{exp } (p t))$ 
by simp
qed
have bounded (path_image p)
by (simp add: assms bounded_path_image)
then obtain B where  $0 < B$  and  $B: \text{path\_image } p \subseteq \text{cball } 0 B$ 
by (meson bounded_pos mem_cball_0 subsetI)
let ?B = cball (0::complex) (B+1)

```

```

have uniformly_continuous_on ?B exp
  using holomorphic_on_exp holomorphic_on_imp_continuous_on
  by (force intro: compact_uniformly_continuous)
then obtain d where d > 0
  and d:  $\bigwedge x x'. \llbracket x \in ?B; x' \in ?B; \text{dist } x' x < d \rrbracket \implies \text{norm } (\text{exp } x' - \text{exp } x) < e$ 
  using  $\langle e > 0 \rangle$  by (auto simp: uniformly_continuous_on_def dist_norm)
then have min 1 d > 0
  by force
then obtain g where pfg: polynomial_function g and g 0 = p 0 g 1 = p 1
  and gless:  $\bigwedge t. t \in \{0..1\} \implies \text{norm}(g t - p t) < \text{min } 1 d$ 
  using path_approx_polynomial_function [OF  $\langle \text{path } p \rangle$ ]  $\langle d > 0 \rangle$   $\langle 0 < e \rangle$ 
  unfolding pathfinish_def pathstart_def by blast
have winding_number (exp o p) 0 = winding_number (exp o g) 0
proof (rule winding_number_nearby_paths_eq [symmetric])
  show path (exp o p) path (exp o g)
  by (simp_all add: pfg assms holomorphic_on_exp holomorphic_on_imp_continuous_on
  path_continuous_image path_polynomial_function)
next
  fix t :: real
  assume t:  $t \in \{0..1\}$ 
  with gless have norm(g t - p t) < 1
    using min_less_iff_conj by blast
  moreover have ptB:  $\text{norm } (p t) \leq B$ 
    using B t by (force simp: path_image_def)
  ultimately have cmod (g t)  $\leq B + 1$ 
    by (meson add_mono thms linordered_field(4) le_less_trans less_imp_le
  norm_triangle_sub)
  with ptB gless t have cmod ((exp o g) t - (exp o p) t) < e
    by (auto simp: dist_norm d)
  with e t show cmod ((exp o g) t - (exp o p) t) < cmod ((exp o p) t - 0)
    by fastforce
qed (use  $\langle g 0 = p 0 \rangle$   $\langle g 1 = p 1 \rangle$  in  $\langle \text{auto simp: pathfinish_def pathstart_def} \rangle$ )
also have ... = 1 / (of_real (2 * pi) * i) * contour_integral (exp o g) ( $\lambda w. 1 / (w - 0)$ )
proof (rule winding_number_valid_path)
  have continuous_on (path_image g) (deriv exp)
    by (metis DERIV_exp DERIV_imp_deriv continuous_on_cong holomorphic_on_exp holomorphic_on_imp_continuous_on)
  then show valid_path (exp o g)
    by (simp add: field_differentiable_within_exp pfg valid_path_compose valid_path_polynomial_function)
  show  $0 \notin \text{path\_image } (\text{exp } \circ \text{g})$ 
    by (auto simp: path_image_def)
qed
also have ... = 1 / (of_real (2 * pi) * i) * integral {0..1} ( $\lambda x. \text{vector\_derivative } g \text{ (at } x)$ )
proof (simp add: contour_integral_integral, rule integral_cong)
  fix t :: real
  assume t:  $t \in \{0..1\}$ 
  show  $\text{vector\_derivative } (\text{exp } \circ \text{g}) \text{ (at } t) / \text{exp } (g t) = \text{vector\_derivative } g \text{ (at } t)$ 

```

```

proof –
  have (exp ∘ g has_vector_derivative vector_derivative (exp ∘ g) (at t)) (at t)
    by (meson DERIV_exp differentiable_def field_vector_diff_chain_at
has_vector_derivative_def
  has_vector_derivative_polynomial_function pfg vector_derivative_works)
  moreover have (exp ∘ g has_vector_derivative vector_derivative g (at t) *
exp (g t)) (at t)
    by (metis DERIV_exp field_vector_diff_chain_at has_vector_derivative_polynomial_function
pfg vector_derivative_at)
  ultimately show ?thesis
    by (simp add: divide_simps, rule vector_derivative_unique_at)
qed
qed
also have ... = (pathfinish p – pathstart p) / (2 * of_real pi * i)
proof –
  have ((λx. vector_derivative g (at x)) has_integral g 1 – g 0) {0..1}
    by (meson differentiable_at_polynomial_function fundamental_theorem_of_calculus
has_vector_derivative_at_within pfg vector_derivative_works
zero_le_one)
  then show ?thesis
    unfolding pathfinish_def pathstart_def
    using ⟨g 0 = p 0⟩ ⟨g 1 = p 1⟩ by auto
qed
finally show ?thesis .
qed
end

```

4 Cauchy's Integral Formula

```

theory Cauchy_Integral_Formula
  imports Winding_Numbers
begin

```

4.1 Proof

```

lemma Cauchy_integral_formula_weak:
  assumes S: convex S and finite k and conf: continuous_on S f
    and fcd: (∧x. x ∈ interior S – k ⇒ f field_differentiable at x)
    and z: z ∈ interior S – k and vpg: valid_path γ
    and pasz: path_image γ ⊆ S – {z} and loop: pathfinish γ = pathstart γ
  shows ((λw. f w / (w – z)) has_contour_integral (2*pi * i * winding_number
γ z * f z)) γ
proof –
  let ?fz = λw. (f w – f z)/(w – z)
  obtain f' where f': (f has_field_derivative f') (at z)
    using fcd [OF z] by (auto simp: field_differentiable_def)
  have pas: path_image γ ⊆ S and znotin: z ∉ path_image γ using pasz by

```



```

blast+
  have c: continuous (at x within S) ( $\lambda w. \text{if } w = z \text{ then } f' \text{ else } (f w - f z) / (w - z)$ ) if  $x \in S$  for  $x$ 
  proof (cases  $x = z$ )
    case True then show ?thesis
      using LIM_equal [of  $z$  ?fz  $\lambda w. \text{if } w = z \text{ then } f' \text{ else } ?fz w$ ] has_field_derivativeD [OF f']
      by (force simp add: continuous_within Lim_at_imp_Lim_at_within)
    next
      case False
      then have dxz:  $\text{dist } x z > 0$  by auto
      have cf: continuous (at x within S) f
        using conf continuous_on_eq_continuous_within that by blast
      have continuous (at x within S) ( $\lambda w. (f w - f z) / (w - z)$ )
        by (rule cf continuous_intros | simp add: False)+
      then show ?thesis
        using continuous_transform_within [OF dxz that] by (force simp: dist_commute)
  qed
  have fink': finite (insert  $z k$ ) using ⟨finite  $k$ ⟩ by blast
  have *: (( $\lambda w. \text{if } w = z \text{ then } f' \text{ else } ?fz w$ ) has_contour_integral 0)  $\gamma$ 
  proof (rule Cauchy_theorem_convex [OF S fink' vpg pas loop])
    show ( $\lambda w. \text{if } w = z \text{ then } f' \text{ else } ?fz w$ ) field_differentiable at  $w$ 
      if  $w \in \text{interior } S - \text{insert } z k$  for  $w$ 
    proof (rule field_differentiable_transform_within)
      show ( $\lambda w. ?fz w$ ) field_differentiable at  $w$ 
        using that by (intro derivative_intros fcd; simp)
    qed (use that in ⟨auto simp add: dist_pos_lt dist_commute⟩)
  qed (use c in ⟨force simp: continuous_on_eq_continuous_within⟩)
  show ?thesis
    apply (rule has_contour_integral_eq)
    using znotin has_contour_integral_add [OF has_contour_integral_lm mul [OF has_contour_integral_winding_number [OF vpg znotin], of f z] *]
    apply (auto simp: ac_simps divide_simps)
  done
qed

```

theorem Cauchy_integral_formula_convex_simple:

```

  assumes convex S and holf: f holomorphic_on S and  $z \in \text{interior } S$  valid_path
   $\gamma$  path_image  $\gamma \subseteq S - \{z\}$ 
  pathfinish  $\gamma = \text{pathstart } \gamma$ 
  shows (( $\lambda w. f w / (w - z)$ ) has_contour_integral ( $2 * \pi * i * \text{winding\_number } \gamma z * f z$ ))  $\gamma$ 
  proof -
    have  $\bigwedge x. x \in \text{interior } S \implies f$  field_differentiable at  $x$ 
      using holf at_within_interior holomorphic_onD interior_subset by fastforce
    then show ?thesis
      using assms
      by (intro Cauchy_integral_formula_weak [where  $k = \{z\}$ ]) (auto simp: holomorphic_on_imp_continuous_on)
  qed

```

qed

Hence the Cauchy formula for points inside a circle.

theorem *Cauchy_integral_circlepath:*

assumes *contf: continuous_on (cball z r) f* **and** *holf: f holomorphic_on (ball z r)* **and** *wz: norm(w - z) < r*
shows $((\lambda u. f u / (u - w)) \text{ has_contour_integral } (2 * \text{of_real } \pi * i * f w))$
(circlepath z r)

proof -

have $r > 0$

using *assms le_less_trans norm_ge_zero* **by** *blast*

have $((\lambda u. f u / (u - w)) \text{ has_contour_integral } (2 * \pi) * i * \text{winding_number}$
*(circlepath z r) w * f w)*
(circlepath z r)

proof (*rule Cauchy_integral_formula_weak [where S = cball z r and k = {}]*)

show $\bigwedge x. x \in \text{interior } (cball z r) - \{w\} \implies$

$f \text{ field_differentiable at } x$

using *holf holomorphic_on_imp_differentiable_at* **by** *auto*

have $w \notin \text{sphere } z r$

by *simp (metis dist_commute dist_norm_not_le order_refl wz)*

then show $\text{path_image } (circlepath z r) \subseteq cball z r - \{w\}$

using $\langle r > 0 \rangle$ **by** (*auto simp add: cball_def sphere_def*)

qed (*use wz in <simp_all add: dist_norm norm_minus_commute contf>*)

then show *?thesis*

by (*simp add: winding_number_circlepath assms*)

qed

corollary *Cauchy_integral_circlepath_simple:*

assumes *f holomorphic_on cball z r norm(w - z) < r*

shows $((\lambda u. f u / (u - w)) \text{ has_contour_integral } (2 * \text{of_real } \pi * i * f w))$
(circlepath z r)

using *assms* **by** (*force simp: holomorphic_on_imp_continuous_on holomorphic_on_subset Cauchy_integral_circlepath*)

4.2 General stepping result for derivative formulas

lemma *Cauchy_next_derivative:*

assumes *continuous_on (path_image γ) f'*

and *leB: $\bigwedge t. t \in \{0..1\} \implies \text{norm } (\text{vector_derivative } \gamma \text{ (at } t)) \leq B$*

and *int: $\bigwedge w. w \in S - \text{path_image } \gamma \implies ((\lambda u. f' u / (u - w)) \wedge k) \text{ has_contour_integral } f w) \gamma$*

and *k: $k \neq 0$*

and *open S*

and *γ : valid_path γ*

and *w: $w \in S - \text{path_image } \gamma$*

shows $(\lambda u. f' u / (u - w)) \wedge (Suc k) \text{ contour_integrable_on } \gamma$

and (*f has_field_derivative (k * contour_integral γ ($\lambda u. f' u / (u - w)) \wedge (Suc k))$*)

(at w) (is ?thes2)

```

proof -
  have open (S - path_image  $\gamma$ ) using ‹open S› closed_valid_path_image  $\gamma$  by
  blast
  then obtain d where d>0 and d: ball w d  $\subseteq$  S - path_image  $\gamma$  using w
  using open_contains_ball by blast
  have [simp]:  $\bigwedge n. \text{cmod } (1 + \text{of\_nat } n) = 1 + \text{of\_nat } n$ 
  by (metis norm_of_nat of_nat_Suc)
  have cint:  $\bigwedge x. \llbracket x \neq w; \text{cmod } (x - w) < d \rrbracket$ 
     $\implies (\lambda z. (f' z / (z - x) ^ k - f' z / (z - w) ^ k) / (x * k - w * k))$ 
  contour_integrable_on  $\gamma$ 
    using int w d
  apply (intro contour_integrable_div contour_integrable_diff has_contour_integral_integrable)
  by (force simp: dist_norm norm_minus_commute)
  have 1:  $\forall_F n$  in at w.  $(\lambda x. f' x * (\text{inverse } (x - n) ^ k - \text{inverse } (x - w) ^ k)$ 
  / (n - w) / of_nat k)
    contour_integrable_on  $\gamma$ 
  unfolding eventually_at
  apply (rule_tac x=d in exI)
  apply (simp add: ‹d > 0› dist_norm field_simps cint)
  done
  have bim_g: bounded (image f' (path_image  $\gamma$ ))
  by (simp add: compact_imp_bounded compact_continuous_image compact_valid_path_image
  assms)
  then obtain C where C > 0 and C:  $\bigwedge x. \llbracket 0 \leq x; x \leq 1 \rrbracket \implies \text{cmod } (f' (\gamma x))$ 
 $\leq C$ 
  by (force simp: bounded_pos path_image_def)
  have twom:  $\forall_F n$  in at w.
     $\forall x \in \text{path\_image } \gamma.$ 
     $\text{cmod } ((\text{inverse } (x - n) ^ k - \text{inverse } (x - w) ^ k) / (n - w) / k -$ 
  inverse (x - w) ^ Suc k) < e
  if 0 < e for e
  proof -
    have *:  $\text{cmod } ((\text{inverse } (x - u) ^ k - \text{inverse } (x - w) ^ k) / ((u - w) * k)$ 
  - inverse (x - w) ^ Suc k) < e
    if x: x  $\in$  path_image  $\gamma$  and u  $\neq$  w and uwd:  $\text{cmod } (u - w) < d/2$ 
    and uw_less:  $\text{cmod } (u - w) < e * (d/2) ^ (k+2) / (1 + \text{real } k)$ 
    for u x
  proof -
    define ff where [abs_def]:
      ff n w =
        (if n = 0 then inverse(x - w) ^ k
         else if n = 1 then k / (x - w) ^ (Suc k)
         else (k * of_real(Suc k)) / (x - w) ^ (k + 2)) for n :: nat and w
    have km1:  $\bigwedge z::\text{complex}. z \neq 0 \implies z ^ (k - \text{Suc } 0) = z ^ k / z$ 
    by (simp add: field_simps) (metis Suc_pred ‹k  $\neq$  0› neq0_conv power_Suc)
    have ff1: (ff i has_field_derivative ff (Suc i) z) (at z within ball w (d/2))
      if z  $\in$  ball w (d/2) i  $\leq$  1 for i z
  proof -
    have z  $\notin$  path_image  $\gamma$ 

```

```

    using ⟨x ∈ path_image γ⟩ d that ball_divide_subset_numeral by blast
  then have xz[simp]: x ≠ z using ⟨x ∈ path_image γ⟩ by blast
  then have neq: x * x + z * z ≠ x * (z * 2)
    by (blast intro: dest!: sum_sqs_eq)
  with xz have ∧v. v ≠ 0 ⇒ (x * x + z * z) * v ≠ (x * (z * 2) * v) by
auto
  then have neqq: ∧v. v ≠ 0 ⇒ x * (x * v) + z * (z * v) ≠ x * (z * (2 *
v))
    by (simp add: algebra_simps)
  show ?thesis using ⟨i ≤ 1⟩
    apply (simp add: ff_def dist_norm Nat.le_Suc_eq km1, safe)
    apply (rule derivative_eq_intros | simp add: km1 | simp add: field_simps
neq neqq)+
    done
  qed
  { fix a::real and b::real assume ab: a > 0 b > 0
    then have k * (1 + real k) * (1 / a) ≤ k * (1 + real k) * (4 / b) ⟷ b
≤ 4 * a
      by (subst mult_le_cancel_left_pos)
        (use ⟨k ≠ 0⟩ in ⟨auto simp: divide_simps⟩)
    with ab have real k * (1 + real k) / a ≤ (real k * 4 + real k * real k * 4)
/ b ⟷ b ≤ 4 * a
      by (simp add: field_simps)
    } note canc = this
  have ff2: cmod (ff (Suc 1) v) ≤ real (k * (k + 1)) / (d/2) ^ (k + 2)
    if v ∈ ball w (d/2) for v
  proof -
    have lessd: ∧z. cmod (γ z - v) < d/2 ⇒ cmod (w - γ z) < d
      by (metis that norm_minus_commute norm_triangle_half_r dist_norm
mem_ball)
    have d/2 ≤ cmod (x - v) using d x that
      using lessd d x
    by (auto simp add: dist_norm path_image_def ball_def not_less [symmetric]
del: divide_const_simps)
    then have d ≤ cmod (x - v) * 2
      by (simp add: field_split_simps)
    then have dpow_le: d ^ (k+2) ≤ (cmod (x - v) * 2) ^ (k+2)
      using ⟨0 < d⟩ order_less_imp_le power_mono by blast
    have x ≠ v using that
      using ⟨x ∈ path_image γ⟩ ball_divide_subset_numeral d by fastforce
    then show ?thesis
      using ⟨d > 0⟩ apply (simp add: ff_def norm_mult norm_divide norm_power
dist_norm canc)
        using dpow_le apply (simp add: field_split_simps)
      done
  qed
  have ub: u ∈ ball w (d/2)
    using uwd by (simp add: dist_commute dist_norm)
  have cmod (inverse (x - u)) ^ k - (inverse (x - w)) ^ k + of_nat k * (u -

```

```

w) / ((x - w) * (x - w) ^ k))
      ≤ (real k * 4 + real k * real k * 4) * (cmod (u - w) * cmod (u -
w) / (d * (d * (d/2) ^ k))
      using complex_Taylor [OF _ ff1 ff2 _ ub, of w, simplified]
      by (simp add: ff_def ‹0 < d›)
      then have cmod (inverse (x - u) ^ k - (inverse (x - w) ^ k + of_nat k *
(u - w) / ((x - w) * (x - w) ^ k)))
          ≤ (cmod (u - w) * real k) * (1 + real k) * cmod (u - w) / (d/2)
^ (k+2)
      by (simp add: field_simps)
      then have cmod (inverse (x - u) ^ k - (inverse (x - w) ^ k + of_nat k *
(u - w) / ((x - w) * (x - w) ^ k)))
          / (cmod (u - w) * real k)
          ≤ (1 + real k) * cmod (u - w) / (d/2) ^ (k+2)
      using ‹k ≠ 0› ‹u ≠ w› by (simp add: mult_ac zero_less_mult_iff
pos_divide_le_eq)
      also have ... < e
      using uw_less ‹0 < d› by (simp add: mult_ac divide_simps)
      finally have e: cmod (inverse (x-u) ^ k - (inverse (x-w) ^ k + of_nat k *
(u-w) / ((x-w) * (x-w) ^ k)))
          / cmod ((u - w) * real k) < e
      by (simp add: norm_mult)
      have x ≠ u
      using uwd ‹0 < d› x d by (force simp: dist_norm ball_def norm_minus_commute)
      show ?thesis
      apply (rule le_less_trans [OF _ e])
      using ‹k ≠ 0› ‹x ≠ u› ‹u ≠ w›
      apply (simp add: field_simps norm_divide [symmetric])
      done
qed
show ?thesis
  unfolding eventually_at
  apply (rule_tac x = min (d/2) ((e*(d/2)^(k + 2))/(Suc k)) in exI)
  apply (force simp: ‹d > 0› dist_norm that simp del: power_Suc intro: *)
  done
qed
have 2: uniform_limit (path_image γ) (λn x. f' x * (inverse (x - n) ^ k -
inverse (x - w) ^ k) / (n - w) / of_nat k) (λx. f' x / (x - w) ^ Suc k) (at w)
  unfolding uniform_limit_iff dist_norm
  proof clarify
    fix e::real
    assume 0 < e
    have *: cmod (f' (γ x) * (inverse (γ x - u) ^ k - inverse (γ x - w) ^ k) /
((u - w) * k) -
          f' (γ x) / ((γ x - w) * (γ x - w) ^ k)) < e
      if ec: cmod ((inverse (γ x - u) ^ k - inverse (γ x - w) ^ k) / ((u -
w) * k) -
          inverse (γ x - w) * inverse (γ x - w) ^ k) < e / C
      and x: 0 ≤ x x ≤ 1

```

```

      for u x
    proof (cases (f' (γ x)) = 0)
      case True then show ?thesis by (simp add: ‹0 < e›)
    next
      case False
      have cmod (f' (γ x) * (inverse (γ x - u) ^ k - inverse (γ x - w) ^ k) / ((u
- w) * k) -
          f' (γ x) / ((γ x - w) * (γ x - w) ^ k)) =
          cmod (f' (γ x) * ((inverse (γ x - u) ^ k - inverse (γ x - w) ^ k) / ((u
- w) * k) -
              inverse (γ x - w) * inverse (γ x - w) ^ k))
        by (simp add: field_simps)
      also have ... = cmod (f' (γ x)) *
          cmod ((inverse (γ x - u) ^ k - inverse (γ x - w) ^ k) / ((u -
w) * k) -
              inverse (γ x - w) * inverse (γ x - w) ^ k)
        by (simp add: norm_mult)
      also have ... < cmod (f' (γ x)) * (e/C)
        using False mult_strict_left_mono [OF ec] by force
      also have ... ≤ e using C
        by (metis False ‹0 < e› frac_le less_eq_real_def mult.commute pos_le_divide_eq
x zero_less_norm_iff)
      finally show ?thesis .
    qed
  show ∀F n in at w.
    ∀ x ∈ path_image γ.
      cmod (f' x * (inverse (x - n) ^ k - inverse (x - w) ^ k) / (n - w)
/ of_nat k - f' x / (x - w) ^ Suc k) < e
    using twom [OF divide_pos_pos [OF ‹0 < e› ‹C > 0›]] unfolding
path_image_def
    by (force intro: * elim: eventually_mono)
  qed
  show (λu. f' u / (u - w) ^ (Suc k)) contour_integrable_on γ
    by (rule contour_integral_uniform_limit [OF 1 2 leB γ]) auto
  have *: (λn. contour_integral γ (λx. f' x * (inverse (x - n) ^ k - inverse (x -
w) ^ k) / (n - w) / k))
    -w → contour_integral γ (λu. f' u / (u - w) ^ (Suc k))
    by (rule contour_integral_uniform_limit [OF 1 2 leB γ]) auto
  have **: contour_integral γ (λx. f' x * (inverse (x - u) ^ k - inverse (x - w)
^ k) / ((u - w) * k)) =
    (f u - f w) / (u - w) / k
    if dist u w < d for u
  proof -
    have u: u ∈ S - path_image γ
      by (metis subsetD d dist_commute mem_ball that)
    have §: ((λx. f' x * inverse (x - u) ^ k) has_contour_integral f u) γ
      ((λx. f' x * inverse (x - w) ^ k) has_contour_integral f w) γ
      using u w by (simp_all add: field_simps int)
    show ?thesis

```

```

    apply (rule contour_integral_unique)
    apply (simp add: diff_divide_distrib algebra_simps § has_contour_integral_diff
has_contour_integral_div)
  done
qed
show ?thes2
  apply (simp add: has_field_derivative_iff del: power_Suc)
  apply (rule Lim_transform_within [OF tendsto_mult_left [OF *] <0 < d> ])
  apply (simp add: <k ≠ 0> **)
  done
qed

```

lemma *Cauchy_next_derivative_circlepath:*

```

  assumes contf: continuous_on (path_image (circlepath z r)) f
    and int:  $\bigwedge w. w \in \text{ball } z \ r \implies ((\lambda u. f \ u / (u - w)^k) \text{ has\_contour\_integral } g \ w) \ (\text{circlepath } z \ r)$ 
    and k:  $k \neq 0$ 
    and w:  $w \in \text{ball } z \ r$ 
  shows  $(\lambda u. f \ u / (u - w) \wedge (\text{Suc } k)) \text{ contour\_integrable\_on } (\text{circlepath } z \ r)$ 
    (is ?thes1)
    and  $(g \text{ has\_field\_derivative } (k * \text{contour\_integral } (\text{circlepath } z \ r) (\lambda u. f \ u / (u - w) \wedge (\text{Suc } k)))) \ (\text{at } w)$ 
    (is ?thes2)

```

proof –

```

  have  $r > 0$  using w
    using ball_eq_empty by fastforce
  have wim:  $w \in \text{ball } z \ r - \text{path\_image } (\text{circlepath } z \ r)$ 
    using w by (auto simp: dist_norm)
  show ?thes1 ?thes2
    by (rule Cauchy_next_derivative [OF contf__int k open_ball valid_path_circlepath
wim, where  $B = 2 * \pi * |r|$ ];
      auto simp: vector_derivative_circlepath norm_mult)+
  qed

```

In particular, the first derivative formula.

lemma *Cauchy_derivative_integral_circlepath:*

```

  assumes contf: continuous_on (cball z r) f
    and holf: f holomorphic_on ball z r
    and w:  $w \in \text{ball } z \ r$ 
  shows  $(\lambda u. f \ u / (u - w)^2) \text{ contour\_integrable\_on } (\text{circlepath } z \ r)$ 
    (is ?thes1)
    and  $(f \text{ has\_field\_derivative } (1 / (2 * \text{of\_real } \pi * i) * \text{contour\_integral } (\text{circlepath } z \ r) (\lambda u. f \ u / (u - w)^2))) \ (\text{at } w)$ 
    (is ?thes2)

```

proof –

```

  have [simp]:  $r \geq 0$  using w
    using ball_eq_empty by fastforce
  have f: continuous_on (path_image (circlepath z r)) f
    by (rule continuous_on_subset [OF contf]) (force simp: cball_def sphere_def)

```

```

have int:  $\bigwedge w. \text{dist } z \ w < r \implies$ 
  (( $\lambda u. f \ u / (u - w)$ ) has_contour_integral ( $\lambda x. 2 * \text{of\_real } \pi * i * f \ x$ ) w) (circlepath z r)
  by (rule Cauchy_integral_circlepath [OF contf holf]) (simp add: dist_norm norm_minus_commute)
  show ?thes1
  apply (simp add: power2_eq_square)
  apply (rule Cauchy_next_derivative_circlepath [OF f __ w, where k=1, simplified])
  apply (blast intro: int)
  done
  have (( $\lambda x. 2 * \text{of\_real } \pi * i * f \ x$ ) has_field_derivative contour_integral (circlepath z r) ( $\lambda u. f \ u / (u - w)^2$ )) (at w)
  apply (simp add: power2_eq_square)
  apply (rule Cauchy_next_derivative_circlepath [OF f __ w, where k=1 and g = \lambda x. 2 * of_real pi * i * f x, simplified])
  apply (blast intro: int)
  done
  then have fder: (f has_field_derivative contour_integral (circlepath z r) ( $\lambda u. f \ u / (u - w)^2 / (2 * \text{of\_real } \pi * i)$ )) (at w)
  by (rule DERIV_cdivide [where  $f = \lambda x. 2 * \text{of\_real } \pi * i * f \ x$  and  $c = 2 * \text{of\_real } \pi * i$ , simplified])
  show ?thes2
  by simp (rule fder)
qed

```

4.3 Existence of all higher derivatives

proposition *derivative_is_holomorphic*:

assumes *open S*

and *fder*: $\bigwedge z. z \in S \implies (f \text{ has_field_derivative } f' \ z) \text{ (at } z)$

shows *f' holomorphic_on S*

proof –

have $*$: $\exists h. (f' \text{ has_field_derivative } h) \text{ (at } z) \text{ if } z \in S$ **for** *z*

proof –

obtain *r* **where** $r > 0$ **and** *r*: *cball* *z r* $\subseteq S$

using *open_contains_cball* $\langle z \in S \rangle \langle \text{open } S \rangle$ **by** *blast*

then have *holf_cball*: *f* *holomorphic_on* *cball* *z r*

unfolding *holomorphic_on_def*

using *field_differentiable_at_within* *field_differentiable_def* *fder* **by** *fastforce*

then have *continuous_on* (*path_image* (*circlepath* *z r*)) *f*

using $\langle r > 0 \rangle$ **by** (*force elim: holomorphic_on_subset* [*THEN holomorphic_on_imp_continuous_on*])

then have *contfpi*: *continuous_on* (*path_image* (*circlepath* *z r*)) ($\lambda x. 1 / (2 * \text{of_real } \pi * i) * f \ x$)

by (*auto intro: continuous_intros*)**+**

have *contf_cball*: *continuous_on* (*cball* *z r*) *f* **using** *holf_cball*

by (*simp add: holomorphic_on_imp_continuous_on holomorphic_on_subset*)

have *holf_ball*: *f* *holomorphic_on* *ball* *z r* **using** *holf_cball*


```

    using ball_subset_cball holomorphic_on_subset by blast
  { fix w assume w: w ∈ ball z r
    have intf: (λu. f u / (u - w)2) contour_integrable_on circlepath z r
      by (blast intro: w Cauchy_derivative_integral_circlepath [OF contf_cball
holf_ball])
    have fder': (f has_field_derivative 1 / (2 * of_real pi * i) * contour_integral
(circlepath z r) (λu. f u / (u - w)2))
      (at w)
      by (blast intro: w Cauchy_derivative_integral_circlepath [OF contf_cball
holf_ball])
    have f'_eq: f' w = contour_integral (circlepath z r) (λu. f u / (u - w)2) /
(2 * of_real pi * i)
      using fder' ball_subset_cball r w by (force intro: DERIV_unique [OF fder])
    have ((λu. f u / (u - w)2 / (2 * of_real pi * i)) has_contour_integral
contour_integral (circlepath z r) (λu. f u / (u - w)2) / (2 * of_real
pi * i))
      (circlepath z r)
      by (rule has_contour_integral_div [OF has_contour_integral_integral [OF
intf]])
    then have ((λu. f u / (2 * of_real pi * i * (u - w)2)) has_contour_integral
contour_integral (circlepath z r) (λu. f u / (u - w)2) / (2 * of_real
pi * i))
      (circlepath z r)
      by (simp add: algebra_simps)
    then have ((λu. f u / (2 * of_real pi * i * (u - w)2)) has_contour_integral
f' w) (circlepath z r)
      by (simp add: f'_eq)
  } note * = this
show ?thesis
  using Cauchy_next_derivative_circlepath [OF contfpi, of 2 f] ‹0 < r‗ *
  using centre_in_ball mem_ball by force
qed
show ?thesis
  by (simp add: holomorphic_on_open [OF ‹open S‗] *)
qed

```

lemma *holomorphic_deriv* [holomorphic_intros]:

[[f holomorphic_on S; open S]] ⇒ (deriv f) holomorphic_on S
by (metis DERIV_deriv_iff_field_differentiable_at_within_open derivative_is_holomorphic
holomorphic_on_def)

lemma *analytic_deriv* [analytic_intros]: f analytic_on S ⇒ (deriv f) analytic_on S

using analytic_on_holomorphic holomorphic_deriv by auto

lemma *holomorphic_higher_deriv* [holomorphic_intros]: [[f holomorphic_on S; open S]] ⇒ (deriv $\hat{\hat{}}$ n) f holomorphic_on S

by (induction n) (auto simp: holomorphic_deriv)

lemma *analytic_higher_deriv* [*analytic_intros*]: f *analytic_on* $S \implies (\text{deriv } \sim n)$
 f *analytic_on* S

unfolding *analytic_on_def* **using** *holomorphic_higher_deriv* **by** *blast*

lemma *has_field_derivative_higher_deriv*:

$\llbracket f$ *holomorphic_on* S ; *open* S ; $x \in S$ \rrbracket

$\implies ((\text{deriv } \sim n) f$ *has_field_derivative* $(\text{deriv } \sim (\text{Suc } n)) f x$) (*at* x)

by (*metis* (*no_types*, *opaque_lifting*) *DERIV_deriv_iff_field_differentiable_at_within_open*
comp_apply

funpow.simps(2) *holomorphic_higher_deriv* *holomorphic_on_def*)

lemma *higher_deriv_cmult*:

assumes f *holomorphic_on* A $x \in A$ *open* A

shows $(\text{deriv } \sim j) (\lambda x. c * f x) x = c * (\text{deriv } \sim j) f x$

using *assms*

proof (*induction* j *arbitrary*: $f x$)

case (*Suc* $j f x$)

have $\text{deriv } ((\text{deriv } \sim j) (\lambda x. c * f x)) x = \text{deriv } (\lambda x. c * (\text{deriv } \sim j) f x) x$

using *eventually_nhds_in_open*[*of* $A x$] *assms*(2,3) *Suc.prem*s

by (*intro* *deriv_cong_ev_refl*) (*auto* *elim!*: *eventually_mono* *simp*: *Suc.IH*)

also have $\dots = c * \text{deriv } ((\text{deriv } \sim j) f) x$ **using** *Suc.prem*s *assms*(2,3)

by (*intro* *deriv_cmult* *holomorphic_on_imp_differentiable_at* *holomorphic_higher_deriv*)
auto

finally show ?*case* **by** *simp*

qed *simp_all*

lemma *valid_path_compose_holomorphic*:

assumes *valid_path* g **and** *holo*: f *holomorphic_on* S **and** *open* S *path_image* g
 $\subseteq S$

shows *valid_path* $(f \circ g)$

by (*meson* *assms* *holomorphic_deriv* *holomorphic_on_imp_continuous_on* *holomorphic_on_imp_differentiable_at*

holomorphic_on_subset *subsetD* *valid_path_compose*)

4.4 Morera's theorem

lemma *Morera_local_triangle_ball*:

assumes $\bigwedge z. z \in S$

$\implies \exists e a. 0 < e \wedge z \in \text{ball } a e \wedge \text{continuous_on } (\text{ball } a e) f \wedge$

$(\forall b c. \text{closed_segment } b c \subseteq \text{ball } a e$

$\longrightarrow \text{contour_integral } (\text{linepath } a b) f +$

$\text{contour_integral } (\text{linepath } b c) f +$

$\text{contour_integral } (\text{linepath } c a) f = 0)$

shows f *analytic_on* S

proof –

{ **fix** z **assume** $z \in S$

with *assms* **obtain** $e a$ **where**

$0 < e$ **and** $z \in \text{ball } a e$ **and** *contf*: *continuous_on* $(\text{ball } a e) f$

and 0 : $\bigwedge b c. \text{closed_segment } b c \subseteq \text{ball } a e$

```

       $\implies \text{contour\_integral } (\text{linepath } a \ b) \ f +$ 
       $\text{contour\_integral } (\text{linepath } b \ c) \ f +$ 
       $\text{contour\_integral } (\text{linepath } c \ a) \ f = 0$ 
    by blast
  have az:  $\text{dist } a \ z < e$  using mem_ball z by blast
  have  $\exists e > 0. f \text{ holomorphic\_on ball } z \ e$ 
  proof (intro exI conjI)
    show  $f \text{ holomorphic\_on ball } z \ (e - \text{dist } a \ z)$ 
    proof (rule holomorphic_on_subset)
      show  $\text{ball } z \ (e - \text{dist } a \ z) \subseteq \text{ball } a \ e$ 
      by (simp add: dist_commute ball_subset_ball_iff)
      have sub_ball:  $\bigwedge y. \text{dist } a \ y < e \implies \text{closed\_segment } a \ y \subseteq \text{ball } a \ e$ 
      by (meson  $\langle 0 < e \rangle$  centre_in_ball convex_ball convex_contains_segment
    mem_ball)
      show  $f \text{ holomorphic\_on ball } a \ e$ 
      using triangle_contour_integrals_starlike_primitive [OF contf_open_ball,
    of a]
      derivative_is_holomorphic[OF open_ball]
      by (force simp add:  $0 < 0 < e$  sub_ball)
    qed
  qed (simp add: az)
}
then show ?thesis
  by (simp add: analytic_on_def)
qed

```

lemma Morera_local_triangle:

```

  assumes  $\bigwedge z. z \in S$ 
     $\implies \exists t. \text{open } t \wedge z \in t \wedge \text{continuous\_on } t \ f \wedge$ 
       $(\forall a \ b \ c. \text{convex\_hull } \{a, b, c\} \subseteq t$ 
       $\longrightarrow \text{contour\_integral } (\text{linepath } a \ b) \ f +$ 
       $\text{contour\_integral } (\text{linepath } b \ c) \ f +$ 
       $\text{contour\_integral } (\text{linepath } c \ a) \ f = 0)$ 
  shows  $f \text{ analytic\_on } S$ 
  proof -
    { fix z assume  $z \in S$ 
      with assms obtain t where
         $\text{open } t$  and  $z: z \in t$  and contf:  $\text{continuous\_on } t \ f$ 
        and 0:  $\bigwedge a \ b \ c. \text{convex\_hull } \{a, b, c\} \subseteq t$ 
           $\implies \text{contour\_integral } (\text{linepath } a \ b) \ f +$ 
           $\text{contour\_integral } (\text{linepath } b \ c) \ f +$ 
           $\text{contour\_integral } (\text{linepath } c \ a) \ f = 0$ 
        by force
      then obtain e where  $e > 0$  and e:  $\text{ball } z \ e \subseteq t$ 
        using open_contains_ball by blast
      have [simp]:  $\text{continuous\_on } (\text{ball } z \ e) \ f$  using contf
        using continuous_on_subset e by blast
      have eq0:  $\bigwedge b \ c. \text{closed\_segment } b \ c \subseteq \text{ball } z \ e \implies$ 
         $\text{contour\_integral } (\text{linepath } z \ b) \ f +$ 

```

```

      contour_integral (linepath b c) f +
      contour_integral (linepath c z) f = 0
    by (meson 0 z ⟨0 < e⟩ centre_in_ball closed_segment_subset convex_ball
dual_order.trans e starlike_convex_subset)
    have ∃ e a. 0 < e ∧ z ∈ ball a e ∧ continuous_on (ball a e) f ∧
      (∀ b c. closed_segment b c ⊆ ball a e →
        contour_integral (linepath a b) f + contour_integral (linepath b
c) f + contour_integral (linepath c a) f = 0)
      using ⟨e > 0⟩ eq0 by force
  }
  then show ?thesis
    by (simp add: Morera_local_triangle_ball)
qed

```

proposition *Morera_triangle*:

```

[[continuous_on S f; open S;
  ∧ a b c. convex_hull {a,b,c} ⊆ S
  → contour_integral (linepath a b) f +
    contour_integral (linepath b c) f +
    contour_integral (linepath c a) f = 0]]
⇒ f analytic_on S
using Morera_local_triangle by blast

```

4.5 Combining theorems for higher derivatives including Leibniz rule

lemma *higher_deriv_linear* [simp]:

```

(deriv ~ n) (λw. c*w) = (λz. if n = 0 then c*z else if n = 1 then c else 0)
by (induction n) auto

```

lemma *higher_deriv_const* [simp]: (deriv ~ n) (λw. c) = (λw. if n=0 then c else 0)

by (induction n) auto

lemma *higher_deriv_ident* [simp]:

```

(deriv ~ n) (λw. w) z = (if n = 0 then z else if n = 1 then 1 else 0)

```

proof (induction n)

case (Suc n)

then show ?case by (metis higher_deriv_linear lambda_one)

qed auto

lemma *higher_deriv_id* [simp]:

```

(deriv ~ n) id z = (if n = 0 then z else if n = 1 then 1 else 0)

```

by (simp add: id_def)

lemma *has_complex_derivative_funpow_1*:

```

[[f has_field_derivative 1] (at z); f z = z]] ⇒ (f ~ n has_field_derivative 1)
(at z)

```

proof (induction n)

```

    case 0
  then show ?case
    by (simp add: id_def)
next
  case (Suc n)
  then show ?case
    by (metis DERIV_chain funpow_Suc_right mult.right_neutral)
qed

lemma higher_deriv_uminus:
  assumes f holomorphic_on S open S and z: z ∈ S
  shows (deriv  $\hat{\hat{}}$  n) (λw. -(f w)) z = - ((deriv  $\hat{\hat{}}$  n) f z)
using z
proof (induction n arbitrary: z)
  case 0 then show ?case by simp
next
  case (Suc n z)
  have *: ((deriv  $\hat{\hat{}}$  n) f has_field_derivative deriv ((deriv  $\hat{\hat{}}$  n) f) z) (at z)
  using Suc.prem1 assms has_field_derivative_higher_deriv by auto
  have ∧x. x ∈ S ⇒ - (deriv  $\hat{\hat{}}$  n) f x = (deriv  $\hat{\hat{}}$  n) (λw. - f w) x
  by (auto simp add: Suc)
  then have ((deriv  $\hat{\hat{}}$  n) (λw. - f w) has_field_derivative - deriv ((deriv  $\hat{\hat{}}$  n)
f) z) (at z)
  using has_field_derivative_transform_within_open [of λw. -((deriv  $\hat{\hat{}}$  n) f
w)]
  using * DERIV_minus Suc.prem1 ‹open S› by blast
  then show ?case
  by (simp add: DERIV_imp_deriv)
qed

lemma higher_deriv_add:
  fixes z::complex
  assumes f holomorphic_on S g holomorphic_on S open S and z: z ∈ S
  shows (deriv  $\hat{\hat{}}$  n) (λw. f w + g w) z = (deriv  $\hat{\hat{}}$  n) f z + (deriv  $\hat{\hat{}}$  n) g z
using z
proof (induction n arbitrary: z)
  case 0 then show ?case by simp
next
  case (Suc n z)
  have *: ((deriv  $\hat{\hat{}}$  n) f has_field_derivative deriv ((deriv  $\hat{\hat{}}$  n) f) z) (at z)
    ((deriv  $\hat{\hat{}}$  n) g has_field_derivative deriv ((deriv  $\hat{\hat{}}$  n) g) z) (at z)
  using Suc.prem1 assms has_field_derivative_higher_deriv by auto
  have ∧x. x ∈ S ⇒ (deriv  $\hat{\hat{}}$  n) f x + (deriv  $\hat{\hat{}}$  n) g x = (deriv  $\hat{\hat{}}$  n) (λw. f
w + g w) x
  by (auto simp add: Suc)
  then have ((deriv  $\hat{\hat{}}$  n) (λw. f w + g w) has_field_derivative
deriv ((deriv  $\hat{\hat{}}$  n) f) z + deriv ((deriv  $\hat{\hat{}}$  n) g) z) (at z)
  using has_field_derivative_transform_within_open [of λw. (deriv  $\hat{\hat{}}$  n) f w
+ (deriv  $\hat{\hat{}}$  n) g w]

```

```

  using * Deriv.field_differentiable_add Suc.premis ‹open S› by blast
  then show ?case
  by (simp add: DERIV_imp_deriv)
qed

```

lemma *higher_deriv_diff*:

```

  fixes z::complex
  assumes f holomorphic_on S g holomorphic_on S open S z ∈ S
  shows (deriv ^^ n) (λw. f w - g w) z = (deriv ^^ n) f z - (deriv ^^ n) g z
  unfolding diff_conv_add_uminus_higher_deriv_add
  using assms higher_deriv_add higher_deriv_uminus holomorphic_on_minus
  by presburger

```

lemma *Suc_choose*: $Suc\ n\ choose\ k = (n\ choose\ k) + (if\ k = 0\ then\ 0\ else\ (n\ choose\ (k - 1)))$

```

  by (cases k) simp_all

```

lemma *higher_deriv_mult*:

```

  fixes z::complex
  assumes f holomorphic_on S g holomorphic_on S open S and z: z ∈ S
  shows (deriv ^^ n) (λw. f w * g w) z =
    (∑ i = 0..n. of_nat (n choose i) * (deriv ^^ i) f z * (deriv ^^ (n - i)) g
z)
  using z
  proof (induction n arbitrary: z)
  case 0 then show ?case by simp
  next
  case (Suc n z)
  have *: ∧ n. ((deriv ^^ n) f has_field_derivative deriv ((deriv ^^ n) f) z) (at z)
    ∧ n. ((deriv ^^ n) g has_field_derivative deriv ((deriv ^^ n) g) z) (at z)
  using Suc.premis assms has_field_derivative_higher_deriv by auto
  have sumeq: (∑ i = 0..n.
    of_nat (n choose i) * (deriv ((deriv ^^ i) f) z * (deriv ^^ (n - i)) g
z + deriv ((deriv ^^ (n - i)) g) z * (deriv ^^ i) f z)) =
    g z * deriv ((deriv ^^ n) f) z + (∑ i = 0..n. (deriv ^^ i) f z * (of_nat
(Suc n choose i) * (deriv ^^ (Suc n - i)) g z))
  apply (simp add: Suc_choose algebra_simps sum.distrib)
  apply (subst (4) sum_Suc_reindex)
  apply (auto simp: algebra_simps Suc_diff_le intro: sum.cong)
  done
  have ((deriv ^^ n) (λw. f w * g w) has_field_derivative
    (∑ i = 0..Suc n. (Suc n choose i) * (deriv ^^ i) f z * (deriv ^^ (Suc n -
i)) g z))
    (at z)
  apply (rule has_field_derivative_transform_within_open
    [of λw. (∑ i = 0..n. of_nat (n choose i) * (deriv ^^ i) f w * (deriv ^^ (n
- i)) g w) _ _ S])
  apply (simp add: algebra_simps)
  apply (rule derivative_eq_intros | simp)+

```

```

    apply (auto intro: DERIV_mult * ⟨open S⟩ Suc.premS Suc.IH [symmetric])
    by (metis (no_types, lifting) mult.commute sum.cong sumeq)
  then show ?case
    unfolding funpow.simps o_apply
    by (simp add: DERIV_imp_deriv)
qed

lemma higher_deriv_transform_within_open:
  fixes z::complex
  assumes f holomorphic_on S g holomorphic_on S open S and z: z ∈ S
    and fg:  $\bigwedge w. w \in S \implies f w = g w$ 
    shows (deriv  $\hat{\sim}$  i) f z = (deriv  $\hat{\sim}$  i) g z
using z
by (induction i arbitrary: z)
  (auto simp: fg intro: complex_derivative_transform_within_open holomorphic_higher_deriv
  assms)

lemma higher_deriv_compose_linear:
  fixes z::complex
  assumes f: f holomorphic_on T and S: open S and T: open T and z: z ∈ S
    and fg:  $\bigwedge w. w \in S \implies u * w \in T$ 
    shows (deriv  $\hat{\sim}$  n) ( $\lambda w. f (u * w)$ ) z =  $u \hat{\sim} n * (deriv \hat{\sim} n) f (u * z)$ 
using z
proof (induction n arbitrary: z)
  case 0 then show ?case by simp
next
  case (Suc n z)
  have holo0: f holomorphic_on (*) u ' S
    by (meson fg f holomorphic_on_subset image_subset_iff)
  have holo2: (deriv  $\hat{\sim}$  n) f holomorphic_on (*) u ' S
    by (meson fg f holomorphic_higher_deriv holomorphic_on_subset image_subset_iff
  T)
  have holo3: ( $\lambda z. u \hat{\sim} n * (deriv \hat{\sim} n) f (u * z)$ ) holomorphic_on S
    by (intro holo2 holomorphic_on_compose [where g=(deriv  $\hat{\sim}$  n) f, unfolded
  o_def] holomorphic_intros)
  have (*) u holomorphic_on S f holomorphic_on (*) u ' S
    by (rule holo0 holomorphic_intros)+
  then have holo1: ( $\lambda w. f (u * w)$ ) holomorphic_on S
    by (rule holomorphic_on_compose [where g=f, unfolded o_def])
  have deriv ((deriv  $\hat{\sim}$  n) ( $\lambda w. f (u * w)$ )) z = deriv ( $\lambda z. u \hat{\sim} n * (deriv \hat{\sim} n) f
  (u * z)$ ) z
  proof (rule complex_derivative_transform_within_open [OF_ holo3 S Suc.premS])
    show (deriv  $\hat{\sim}$  n) ( $\lambda w. f (u * w)$ ) holomorphic_on S
      by (rule holomorphic_higher_deriv [OF holo1 S])
  qed (simp add: Suc.IH)
  also have ... =  $u \hat{\sim} n * deriv (\lambda z. (deriv \hat{\sim} n) f (u * z)) z$ 
  proof -
    have (deriv  $\hat{\sim}$  n) f analytic_on T
      by (simp add: analytic_on_open f holomorphic_higher_deriv T)

```

```

then have ( $\lambda w. (\text{deriv } \overset{\sim}{\sim} n) f (u * w)$ ) analytic_on S
proof -
  have ( $\text{deriv } \overset{\sim}{\sim} n) f \circ (*) u$  holomorphic_on S
  by (simp add: holo2 holomorphic_on_compose)
  then show ?thesis
  by (simp add: S analytic_on_open o_def)
qed
then show ?thesis
by (intro deriv_cmult analytic_on_imp_differentiable_at [OF __ Suc.prem])
qed
also have ... =  $u * u^{\wedge} n * \text{deriv } ((\text{deriv } \overset{\sim}{\sim} n) f) (u * z)$ 
proof -
  have ( $\text{deriv } \overset{\sim}{\sim} n) f$  field_differentiable at (u * z)
  using Suc.prem T ffg holomorphic_higher_deriv holomorphic_on_imp_differentiable_at
by blast
  then show ?thesis
  by (simp add: deriv_compose_linear)
qed
finally show ?case
by simp
qed

```

lemma *higher_deriv_add_at*:

```

assumes f analytic_on {z} g analytic_on {z}
shows ( $\text{deriv } \overset{\sim}{\sim} n) (\lambda w. f w + g w) z = (\text{deriv } \overset{\sim}{\sim} n) f z + (\text{deriv } \overset{\sim}{\sim} n) g z$ 
using analytic_at_two assms higher_deriv_add by blast

```

lemma *higher_deriv_diff_at*:

```

assumes f analytic_on {z} g analytic_on {z}
shows ( $\text{deriv } \overset{\sim}{\sim} n) (\lambda w. f w - g w) z = (\text{deriv } \overset{\sim}{\sim} n) f z - (\text{deriv } \overset{\sim}{\sim} n) g z$ 
using analytic_at_two assms higher_deriv_diff by blast

```

lemma *higher_deriv_uminus_at*:

```

f analytic_on {z}  $\implies (\text{deriv } \overset{\sim}{\sim} n) (\lambda w. -(f w)) z = - ((\text{deriv } \overset{\sim}{\sim} n) f z)$ 
using higher_deriv_uminus by (auto simp: analytic_at)

```

lemma *higher_deriv_mult_at*:

```

assumes f analytic_on {z} g analytic_on {z}
shows ( $\text{deriv } \overset{\sim}{\sim} n) (\lambda w. f w * g w) z =$ 
  ( $\sum i = 0..n. \text{of\_nat } (n \text{ choose } i) * (\text{deriv } \overset{\sim}{\sim} i) f z * (\text{deriv } \overset{\sim}{\sim} (n - i)) g$ 
z)
using analytic_at_two assms higher_deriv_mult by blast

```

Nonexistence of isolated singularities and a stronger integral formula.

proposition *no_isolated_singularity*:

```

fixes z::complex
assumes f: continuous_on S f and holf: f holomorphic_on (S - K) and S:
open S and K: finite K
shows f holomorphic_on S

```



```

proof –
  { fix z
    assume  $z \in S$  and  $\text{cdf}: \bigwedge x. x \in S - K \implies f \text{ field\_differentiable at } x$ 
    have  $f \text{ field\_differentiable at } z$ 
    proof (cases  $z \in K$ )
      case False then show ?thesis by (blast intro: cdf  $\langle z \in S \rangle$ )
    next
      case True
      with finite_set_avoid [OF  $K$ , of  $z$ ]
      obtain  $d$  where  $d > 0$  and  $d: \bigwedge x. [x \in K; x \neq z] \implies d \leq \text{dist } z \ x$ 
        by blast
      obtain  $e$  where  $e > 0$  and  $e: \text{ball } z \ e \subseteq S$ 
        using  $S \langle z \in S \rangle$  by (force simp: open_contains_ball)
      have  $\text{fde}: \text{continuous\_on } (\text{ball } z \ (\text{min } d \ e)) \ f$ 
        by (metis Int_iff ball_min_Int continuous_on_subset e f subsetI)
      have  $\text{cont}: \{a, b, c\} \subseteq \text{ball } z \ (\text{min } d \ e) \implies \text{continuous\_on } (\text{convex hull } \{a, b, c\}) \ f$  for  $a \ b \ c$ 
        by (simp add: hull_minimal continuous_on_subset [OF fde])
      have  $\text{fd}: [\{a, b, c\} \subseteq \text{ball } z \ (\text{min } d \ e); x \in \text{interior } (\text{convex hull } \{a, b, c\}) - K]$ 
         $\implies f \text{ field\_differentiable at } x$  for  $a \ b \ c \ x$ 
        by (metis cdf Diff_iff Int_iff ball_min_Int subsetD convex_ball e interior_mono interior_subset subset_hull)
      obtain  $g$  where  $\bigwedge w. w \in \text{ball } z \ (\text{min } d \ e) \implies (g \text{ has\_field\_derivative } f \ w)$ 
        (at  $w$  within  $\text{ball } z \ (\text{min } d \ e)$ )
        apply (rule contour_integral_convex_primitive
          [OF  $\text{convex\_ball fde Cauchy\_theorem\_triangle\_cofinite [OF } - K]$ ])
        using  $\text{cont fd}$  by auto
      then have  $f \text{ holomorphic\_on } \text{ball } z \ (\text{min } d \ e)$ 
        by (metis open_ball at_within_open derivative_is_holomorphic)
      then show ?thesis
        unfolding holomorphic_on_def
        by (metis open_ball  $\langle 0 < d \rangle \langle 0 < e \rangle$  at_within_open centre_in_ball min_less_iff_conj)
      qed
    }
  with holf  $S \ K$  show ?thesis
    by (simp add: holomorphic_on_open open_Diff finite_imp_closed field_differentiable_def [symmetric])
  qed

lemma no_isolated_singularity':
  fixes  $z::\text{complex}$ 
  assumes  $f: \bigwedge z. z \in K \implies (f \longrightarrow f \ z)$  (at  $z$  within  $S$ )
    and holf:  $f \text{ holomorphic\_on } (S - K)$  and  $S: \text{open } S$  and  $K: \text{finite } K$ 
    shows  $f \text{ holomorphic\_on } S$ 
proof (rule no_isolated_singularity [OF assms(2-)])
  show  $\text{continuous\_on } S \ f$  unfolding continuous_on_def

```

```

proof
  fix  $z$  assume  $z: z \in S$ 
  have continuous_on ( $S - K$ )  $f$ 
    using holf holomorphic_on_imp_continuous_on by auto
  then show ( $f \longrightarrow f z$ ) (at z within S)
    by (metis Diff_iff K S at_within_interior continuous_on_def finite_imp_closed
interior_eq_open_Diff z)
  qed
qed

```

proposition *Cauchy_integral_formula_convex*:

```

assumes  $S$ : convex S and  $K$ : finite K and contf: continuous_on S f
and fcd: ( $\bigwedge x. x \in \text{interior } S - K \implies f \text{ field\_differentiable at } x$ )
and  $z: z \in \text{interior } S$  and vpg: valid_path  $\gamma$ 
and pasz: path_image  $\gamma \subseteq S - \{z\}$  and loop: pathfinish  $\gamma = \text{pathstart } \gamma$ 
shows ( $(\lambda w. f w / (w - z)) \text{ has\_contour\_integral } (2 * \pi * i * \text{winding\_number } \gamma z * f z)$ )  $\gamma$ 
proof -
  have  $*$ :  $\bigwedge x. x \in \text{interior } S \implies f \text{ field\_differentiable at } x$ 
    unfolding holomorphic_on_open [symmetric] field_differentiable_def
    using no_isolated_singularity [where S = interior S]
  by (meson K contf continuous_at_imp_continuous_on continuous_on_interior
fcd
    field_differentiable_at_within field_differentiable_def holomorphic_onI
holomorphic_on_imp_differentiable_at open_interior)
  show ?thesis
    by (rule Cauchy_integral_formula_weak [OF S finite.emptyI contf]) (use *
assms in auto)
qed

```

Formula for higher derivatives.

lemma *Cauchy_has_contour_integral_higher_derivative_circlepath*:

```

assumes contf: continuous_on (cball z r) f
and holf: f holomorphic_on ball z r
and  $w: w \in \text{ball } z r$ 
shows ( $(\lambda u. f u / (u - w) \wedge (\text{Suc } k)) \text{ has\_contour\_integral } ((2 * \pi * i) /$ 
(fact k) * (deriv  $\hat{\sim}$  k) f w)
  (circlepath z r)
using  $w$ 
proof (induction k arbitrary: w)
  case 0 then show ?case
    using assms by (auto simp: Cauchy_integral_circlepath dist_commute dist_norm)
  next
    case (Suc k)
    have [simp]:  $r > 0$  using  $w$ 
      using ball_eq_empty by fastforce
    have  $f$ : continuous_on (path_image (circlepath z r)) f
      by (rule continuous_on_subset [OF contf]) (force simp: cball_def sphere_def
less_imp_le)

```

```

obtain X where X: (( $\lambda u. f u / (u - w) \wedge \text{Suc } (\text{Suc } k)$ ) has_contour_integral
X) (circlepath z r)
  using Cauchy_next_derivative_circlepath(1) [OF f Suc.IH _ Suc.prems]
  by (auto simp: contour_integrable_on_def)
  then have con: contour_integral (circlepath z r) (( $\lambda u. f u / (u - w) \wedge \text{Suc } (\text{Suc } k)$ )) = X
    by (rule contour_integral_unique)
  have  $\bigwedge n. ((\text{deriv } \wedge n) f \text{ has\_field\_derivative } \text{deriv } ((\text{deriv } \wedge n) f) w) (\text{at } w)$ 
    using Suc.prems assms has_field_derivative_higher_deriv by auto
  then have dnf_diff:  $\bigwedge n. (\text{deriv } \wedge n) f \text{ field\_differentiable } (\text{at } w)$ 
    by (force simp: field_differentiable_def)
  have deriv ( $\lambda w. \text{complex\_of\_real } (2 * \text{pi}) * \text{i} / (\text{fact } k) * (\text{deriv } \wedge k) f w) w =$ 
of_nat (Suc k) * contour_integral (circlepath z r) ( $\lambda u. f u / (u - w) \wedge$ 
Suc (Suc k))
    by (force intro!: DERIV_imp_deriv Cauchy_next_derivative_circlepath [OF f
Suc.IH _ Suc.prems])
  also have  $\dots = \text{of\_nat } (\text{Suc } k) * X$ 
    by (simp only: con)
  finally have deriv ( $\lambda w. ((2 * \text{pi}) * \text{i} / (\text{fact } k)) * (\text{deriv } \wedge k) f w) w = \text{of\_nat}$ 
(Suc k) * X .
    then have  $((2 * \text{pi}) * \text{i} / (\text{fact } k)) * \text{deriv } (\lambda w. (\text{deriv } \wedge k) f w) w = \text{of\_nat}$ 
(Suc k) * X
    by (metis deriv_cmult dnf_diff)
  then have deriv ( $\lambda w. (\text{deriv } \wedge k) f w) w = \text{of\_nat } (\text{Suc } k) * X / ((2 * \text{pi}) * \text{i}$ 
/ (\text{fact } k))
    by (simp add: field_simps)
  then show ?case
  using of_nat_eq_0_iff X by fastforce
qed

```

lemma *Cauchy_higher_derivative_integral_circlepath:*

```

assumes contf: continuous_on (cball z r) f
  and holf: f holomorphic_on ball z r
  and w: w  $\in$  ball z r
shows  $(\lambda u. f u / (u - w) \wedge (\text{Suc } k)) \text{ contour\_integrable\_on } (\text{circlepath } z r)$ 
  (is ?thes1)
  and  $(\text{deriv } \wedge k) f w = (\text{fact } k) / (2 * \text{pi} * \text{i}) * \text{contour\_integral}(\text{circlepath } z$ 
r) ( $\lambda u. f u / (u - w) \wedge (\text{Suc } k)$ )
  (is ?thes2)
proof -
  have  $*$ :  $(\lambda u. f u / (u - w) \wedge \text{Suc } k) \text{ has\_contour\_integral } (2 * \text{pi}) * \text{i} / (\text{fact } k)$ 
 $* (\text{deriv } \wedge k) f w$ 
  (circlepath z r)
    using Cauchy_has_contour_integral_higher_derivative_circlepath [OF assms]
    by simp
  show ?thes1 using  $*$ 
    using contour_integrable_on_def by blast
  show ?thes2
    unfolding contour_integral_unique [OF *] by (simp add: field_split_simps)

```

qed

corollary *Cauchy_contour_integral_circlepath*:

assumes *continuous_on* (cball z r) f f *holomorphic_on* ball z r $w \in \text{ball } z \ r$
shows $\text{contour_integral}(\text{circlepath } z \ r) (\lambda u. f \ u / (u - w)^{\wedge} (\text{Suc } k)) = (2 * \pi * i) * (\text{deriv } \wedge^k) f \ w / (\text{fact } k)$
by (*simp add: Cauchy_higher_derivative_integral_circlepath [OF assms]*)

lemma *Cauchy_contour_integral_circlepath_2*:

assumes *continuous_on* (cball z r) f f *holomorphic_on* ball z r $w \in \text{ball } z \ r$
shows $\text{contour_integral}(\text{circlepath } z \ r) (\lambda u. f \ u / (u - w)^{\wedge} 2) = (2 * \pi * i) * \text{deriv } f \ w$
using *Cauchy_contour_integral_circlepath [OF assms, of 1]*
by (*simp add: power2_eq_square*)

4.6 A holomorphic function is analytic, i.e. has local power series

theorem *holomorphic_power_series*:

assumes *holfc*: *f* *holomorphic_on* ball z r
and $w \in \text{ball } z \ r$
shows $((\lambda n. (\text{deriv } \wedge^n) f \ z / (\text{fact } n) * (w - z)^{\wedge} n) \ \text{sums } f \ w)$

proof –

— Replacing r and the original (weak) premises with stronger ones

obtain r **where** $r > 0$ **and** *holfc*: *f* *holomorphic_on* cball z r **and** $w \in \text{ball } z \ r$

proof

have $\text{cball } z \ ((r + \text{dist } w \ z) / 2) \subseteq \text{ball } z \ r$

using w **by** (*simp add: dist_commute field_sum_of_halves subset_eq*)

then show *f* *holomorphic_on* cball z ((r + dist w z) / 2)

by (*rule holomorphic_on_subset [OF holfc]*)

have $r > 0$

using w **by** *clarsimp* (*metis dist_norm le_less_trans norm_ge_zero*)

then show $0 < (r + \text{dist } w \ z) / 2$

by *simp* (*use zero_le_dist [of w z] in linarith*)

qed (*use w in <auto simp: dist_commute>*)

then have *holfc*: *f* *holomorphic_on* ball z r

using *ball_subset_cball holomorphic_on_subset* **by** *blast*

have *contf*: *continuous_on* (cball z r) *f*

by (*simp add: holfc holomorphic_on_imp_continuous_on*)

have *cint*: $\bigwedge k. (\lambda u. f \ u / (u - z)^{\wedge} (\text{Suc } k)) \ \text{contour_integrable_on } \text{circlepath } z \ r$

by (*rule Cauchy_higher_derivative_integral_circlepath [OF contf holfc]*) (*simp add: <0 < r>*)

obtain B **where** $0 < B$ **and** $B: \bigwedge u. u \in \text{cball } z \ r \implies \text{norm}(f \ u) \leq B$

by (*metis* (*no_types*) *bounded_pos compact_cball compact_continuous_image compact_imp_bounded contf_image_eqI*)

obtain k **where** $0 < k \ k \leq r$ **and** *wz_eq*: $\text{norm}(w - z) = r - k$

and *kle*: $\bigwedge u. \text{norm}(u - z) = r \implies k \leq \text{norm}(u - w)$

proof

```

show  $\bigwedge u. \text{cmod } (u - z) = r \implies r - \text{dist } z \ w \leq \text{cmod } (u - w)$ 
  by (metis add_diff_eq diff_add_cancel dist_norm norm_diff_ineq)
qed (use w in <auto simp: dist_norm norm_minus_commute>)
have ul: uniform_limit (sphere z r) ( $\lambda n x. (\sum k < n. (w - z) ^ k * (f x / (x - z) ^ \text{Suc } k)$ )) ( $\lambda x. f x / (x - w)$ ) sequentially
  unfolding uniform_limit_iff dist_norm
proof clarify
  fix e::real
  assume  $0 < e$ 
  have rr:  $0 \leq (r - k) / r \ (r - k) / r < 1$  using k by auto
  obtain n where n:  $((r - k) / r) ^ n < e / B * k$ 
    using real_arch_pow_inv [of e/B*k (r - k)/r]  $\langle 0 < e \rangle \langle 0 < B \rangle$  k by force
  have norm  $((\sum k < N. (w - z) ^ k * f u / (u - z) ^ \text{Suc } k) - f u / (u - w)) < e$ 
    if  $n \leq N$  and r:  $r = \text{dist } z \ u$  for  $N \ u$ 
  proof -
    have N:  $((r - k) / r) ^ N < e / B * k$ 
      using le_less_trans [OF power_decreasing n]
      using  $\langle n \leq N \rangle$  k by auto
    have u [simp]:  $(u \neq z) \wedge (u \neq w)$ 
      using  $\langle 0 < r \rangle$  r by auto
    have wzu_not1:  $(w - z) / (u - z) \neq 1$ 
      by (metis (no_types) dist_norm divide_eq_1_iff less_irrefl mem_ball norm_minus_commute r w)
    have norm  $((\sum k < N. (w - z) ^ k * f u / (u - z) ^ \text{Suc } k) * (u - w) - f u) = \text{norm } ((\sum k < N. (((w - z) / (u - z)) ^ k)) * f u * (u - w) / (u - z) - f u)$ 
      unfolding sum_distrib_right sum_divide_distrib power_divide by (simp add: algebra_simps)
    also have ... =  $\text{norm } (((w - z) / (u - z)) ^ N - 1) * (u - w) / (((w - z) / (u - z) - 1) * (u - z) - 1) * \text{norm } (f u)$ 
      using  $\langle 0 < B \rangle$ 
      apply (auto simp: geometric_sum [OF wzu_not1])
      apply (simp add: field_simps norm_mult [symmetric])
      done
    also have ... =  $\text{norm } ((u - z) ^ N * (w - u) - ((w - z) ^ N - (u - z) ^ N) * (u - w)) / (r ^ N * \text{norm } (u - w)) * \text{norm } (f u)$ 
      using  $\langle 0 < r \rangle$  r by (simp add: divide_simps norm_mult norm_divide norm_power dist_norm norm_minus_commute)
    also have ... =  $\text{norm } ((w - z) ^ N * (w - u)) / (r ^ N * \text{norm } (u - w)) * \text{norm } (f u)$ 
      by (simp add: algebra_simps)
    also have ... =  $\text{norm } (w - z) ^ N * \text{norm } (f u) / r ^ N$ 
      by (simp add: norm_mult norm_power norm_minus_commute)
    also have ...  $\leq (((r - k) / r) ^ N) * B$ 
      using  $\langle 0 < r \rangle$  w k
      by (simp add: B divide_simps mult_mono r wz_eq)
    also have ...  $< e * k$ 
      using  $\langle 0 < B \rangle$  N by (simp add: divide_simps)
  qed

```

```

also have ...  $\leq e * \text{norm } (u - w)$ 
using r kle <0 < e> by (simp add: dist_commute dist_norm)
finally show ?thesis
by (simp add: field_split_simps norm_divide del: power_Suc)
qed
with <0 < r> show  $\forall_F n$  in sequentially.  $\forall x \in \text{sphere } z \ r.$ 
 $\text{norm } ((\sum k < n. (w - z) \wedge k * (f x / (x - z) \wedge \text{Suc } k)) - f x / (x - w)) < e$ 
by (auto simp: mult_ac less_imp_le eventually_sequentially Ball_def)
qed
have §:  $\bigwedge x \ k. k \in \{..<x\} \implies$ 
 $(\lambda u. (w - z) \wedge k * (f u / (u - z) \wedge \text{Suc } k)) \text{ contour\_integrable\_on } \text{circlepath } z \ r$ 
using contour_integrable_lmul [OF cint, of (w - z) ^ a for a] by (simp add: field_simps)
have eq:  $\forall_F x$  in sequentially.
 $\text{contour\_integral } (\text{circlepath } z \ r) (\lambda u. \sum k < x. (w - z) \wedge k * (f u / (u - z) \wedge \text{Suc } k)) =$ 
 $(\sum k < x. \text{contour\_integral } (\text{circlepath } z \ r) (\lambda u. f u / (u - z) \wedge \text{Suc } k) * (w - z) \wedge k)$ 
apply (rule eventuallyI)
apply (subst contour_integral_sum, simp)
apply (simp_all only: § contour_integrable_lmul cint algebra_simps)
done
have  $\bigwedge u \ k. k \in \{..<u\} \implies (\lambda x. f x / (x - z) \wedge \text{Suc } k) \text{ contour\_integrable\_on } \text{circlepath } z \ r$ 
using <0 < r> by (force intro!: Cauchy_higher_derivative_integral_circlepath [OF contf holf])
then have  $\bigwedge u. (\lambda y. \sum k < u. (w - z) \wedge k * (f y / (y - z) \wedge \text{Suc } k)) \text{ contour\_integrable\_on } \text{circlepath } z \ r$ 
by (intro contour_integrable_sum contour_integrable_lmul, simp)
then have  $(\lambda k. \text{contour\_integral } (\text{circlepath } z \ r) (\lambda u. f u / (u - z) \wedge (\text{Suc } k)) * (w - z) \wedge k)$ 
 $\text{sums contour\_integral } (\text{circlepath } z \ r) (\lambda u. f u / (u - w))$ 
unfolding sums_def using <0 < r>
by (intro Lim_transform_eventually [OF _ eq] contour_integral_uniform_limit_circlepath [OF eventuallyI ul]) auto
then have  $(\lambda k. \text{contour\_integral } (\text{circlepath } z \ r) (\lambda u. f u / (u - z) \wedge (\text{Suc } k)) * (w - z) \wedge k)$ 
 $\text{sums } (2 * \text{of\_real } \pi * i * f w)$ 
using w by (auto simp: dist_commute dist_norm contour_integral_unique [OF Cauchy_integral_circlepath_simple [OF holfc]])
then have  $(\lambda k. \text{contour\_integral } (\text{circlepath } z \ r) (\lambda u. f u / (u - z) \wedge \text{Suc } k) * (w - z) \wedge k / (i * (\text{of\_real } \pi * 2)))$ 
 $\text{sums } ((2 * \text{of\_real } \pi * i * f w) / (i * (\text{complex\_of\_real } \pi * 2)))$ 
by (rule sums_divide)
then have  $(\lambda n. (w - z) \wedge n * \text{contour\_integral } (\text{circlepath } z \ r) (\lambda u. f u / (u - z) \wedge \text{Suc } n) / (i * (\text{of\_real } \pi * 2)))$ 
 $\text{sums } f w$ 

```

```

  by (simp add: field_simps)
  then show ?thesis
  by (simp add: field_simps <0 < r> Cauchy_higher_derivative_integral_circlepath
[OF contf holf])
qed

```

4.7 The Liouville theorem and the Fundamental Theorem of Algebra

These weak Liouville versions don't even need the derivative formula.

```

lemma Liouville_weak_0:
  assumes holf: f holomorphic_on UNIV and inf: (f  $\longrightarrow$  0) at_infinity
  shows f z = 0
proof (rule ccontr)
  assume fz: f z  $\neq$  0
  with inf [unfolded Lim_at_infinity, rule_format, of norm(f z)/2]
  obtain B where B:  $\bigwedge x. B \leq cmod\ x \implies norm\ (f\ x) * 2 < cmod\ (f\ z)$ 
    by (auto simp: dist_norm)
  define R where R = 1 + |B| + norm z
  have R > 0
    unfolding R_def by (smt (verit) norm_ge_zero)
  have *: (( $\lambda u. f\ u / (u - z)$ ) has_contour_integral 2 * complex_of_real pi * i *
f z) (circlepath z R)
    using continuous_on_subset holf holomorphic_on_subset <0 < R>
    by (force intro: holomorphic_on_imp_continuous_on Cauchy_integral_circlepath)
  have cmod (x - z) = R  $\implies cmod\ (f\ x) * 2 < cmod\ (f\ z)$  for x
    unfolding R_def by (rule B) (use norm_triangle_ineq4 [of x z] in auto)
  with <R > 0> fz show False
    using has_contour_integral_bound_circlepath [OF *, of norm(f z)/2/R]
    by (auto simp: less_imp_le norm_mult norm_divide field_split_simps)
qed

```

```

proposition Liouville_weak:
  assumes f holomorphic_on UNIV and (f  $\longrightarrow$  l) at_infinity
  shows f z = l
  using Liouville_weak_0 [of  $\lambda z. f\ z - l$ ]
  by (simp add: assms holomorphic_on_diff LIM_zero)

```

```

proposition Liouville_weak_inverse:
  assumes f holomorphic_on UNIV and unbounded:  $\bigwedge B. eventually\ (\lambda x. norm\ (f\ x) \geq B)$  at_infinity
  obtains z where f z = 0
proof -
  { assume f:  $\bigwedge z. f\ z \neq 0$ 
    have 1: ( $\lambda x. 1 / f\ x$ ) holomorphic_on UNIV
      by (simp add: holomorphic_on_divide assms f)
    have 2: (( $\lambda x. 1 / f\ x$ )  $\longrightarrow$  0) at_infinity
    proof (rule tendstoI [OF eventually_mono])
      fix e::real

```

```

    assume  $e > 0$ 
    show eventually  $(\lambda x. 2/e \leq cmod (f x))$  at_infinity
      by (rule_tac  $B=2/e$  in unbounded)
    qed (simp add: dist_norm norm_divide field_split_simps)
    have False
      using Liouville_weak_0 [OF 1 2] f by simp
  }
  then show ?thesis
    using that by blast
qed

```

In particular we get the Fundamental Theorem of Algebra.

```

theorem fundamental_theorem_of_algebra:
  fixes  $a :: nat \Rightarrow complex$ 
  assumes  $a\ 0 = 0 \vee (\exists i \in \{1..n\}. a\ i \neq 0)$ 
  obtains  $z$  where  $(\sum_{i \leq n}. a\ i * z^i) = 0$ 
using assms
proof (elim disjE bexE)
  assume  $a\ 0 = 0$  then show ?thesis
    by (auto simp: that [of 0])
next
  fix  $i$ 
  assume  $i: i \in \{1..n\}$  and  $nz: a\ i \neq 0$ 
  have 1:  $(\lambda z. \sum_{i \leq n}. a\ i * z^i)$  holomorphic_on UNIV
    by (rule holomorphic_intros)+
  show thesis
  proof (rule Liouville_weak_inverse [OF 1])
    show  $\forall_F x$  in at_infinity.  $B \leq cmod (\sum_{i \leq n}. a\ i * x^i)$  for  $B$ 
      using  $i\ nz$  by (intro polyfun_extremal exI[of _ i]) auto
    qed (use that in auto)
  qed

```

4.8 Weierstrass convergence theorem

```

lemma holomorphic_uniform_limit:
  assumes  $cont: eventually (\lambda n. continuous\_on (cball\ z\ r) (f\ n) \wedge (f\ n)$  holomor-
  phic_on ball\ z\ r)  $F$ 
  and  $ulim: uniform\_limit (cball\ z\ r) f\ g\ F$ 
  and  $F: \neg trivial\_limit\ F$ 
  obtains  $continuous\_on (cball\ z\ r) g$   $g$  holomorphic_on ball\ z\ r
proof (cases  $r\ 0::real$  rule: linorder_cases)
  case less then show ?thesis by (force simp: ball_empty less_imp_le continu-
  ous_on_def holomorphic_on_def intro: that)
next
  case equal then show ?thesis
    by (force simp: holomorphic_on_def intro: that)
next
  case greater
  have  $contg: continuous\_on (cball\ z\ r) g$ 

```



```

using cont uniform_limit_theorem [OF eventually_mono ulim F] by blast
have path_image (circlepath z r)  $\subseteq$  cball z r
using ‹0 < r› by auto
then have 1: continuous_on (path_image (circlepath z r)) ( $\lambda x. 1 / (2 * \text{complex\_of\_real } \pi * i) * g x$ )
by (intro continuous_intros continuous_on_subset [OF contg])
have 2: (( $\lambda u. 1 / (2 * \text{of\_real } \pi * i) * g u / (u - w) ^ \wedge 1$ ) has_contour_integral
g w) (circlepath z r)
if w: w  $\in$  ball z r for w
proof -
define d where d = (r - norm(w - z))
have 0 < d d  $\leq$  r using w by (auto simp: norm_minus_commute d_def
dist_norm)
have dle:  $\bigwedge u. \text{cmod } (z - u) = r \implies d \leq \text{cmod } (u - w)$ 
unfolding d_def by (metis add_diff_eq diff_add_cancel norm_diff_ineq
norm_minus_commute)
have ev_int:  $\forall_F n \text{ in } F. (\lambda u. f n u / (u - w)) \text{ contour\_integrable\_on circlepath } z r$ 
using w
by (auto intro: eventually_mono [OF cont] Cauchy_higher_derivative_integral_circlepath
[where k=0, simplified])
have  $\bigwedge e. [0 < r; 0 < d; 0 < e]$ 
 $\implies \forall_F n \text{ in } F.$ 
 $\forall x \in \text{sphere } z r.$ 
 $x \neq w \longrightarrow$ 
 $\text{cmod } (f n x - g x) < e * \text{cmod } (x - w)$ 
apply (rule_tac e1=e * d in eventually_mono [OF uniform_limitD [OF
ulim]])
apply (force simp: dist_norm intro: dle mult_left_mono less_le_trans)+
done
then have ul_less: uniform_limit (sphere z r) ( $\lambda n x. f n x / (x - w)$ ) ( $\lambda x. g$ 
x / (x - w)) F
using greater ‹0 < d›
by (auto simp add: uniform_limit_iff dist_norm norm_divide diff_divide_distrib
[symmetric] divide_simps)
have g_cint: ( $\lambda u. g u / (u - w)$ ) contour_integrable_on circlepath z r
by (rule contour_integral_uniform_limit_circlepath [OF ev_int ul_less F ‹0
< r›])
have cif_tends_cig: (( $\lambda n. \text{contour\_integral}(\text{circlepath } z r) (\lambda u. f n u / (u -$ 
w)))  $\longrightarrow \text{contour\_integral}(\text{circlepath } z r) (\lambda u. g u / (u - w))$ ) F
by (rule contour_integral_uniform_limit_circlepath [OF ev_int ul_less F ‹0
< r›])
have f_tends_cig: (( $\lambda n. 2 * \text{of\_real } \pi * i * f n w$ )  $\longrightarrow \text{contour\_integral}$ 
(circlepath z r) ( $\lambda u. g u / (u - w)$ )) F
proof (rule Lim_transform_eventually)
show  $\forall_F x \text{ in } F. \text{contour\_integral } (\text{circlepath } z r) (\lambda u. f x u / (u - w))$ 
 $= 2 * \text{of\_real } \pi * i * f x w$ 
using w ‹0 < d› d_def
by (auto intro: eventually_mono [OF cont contour_integral_unique [OF

```

```

Cauchy_integral_circlepath]])
qed (auto simp: cif_tends_cig)
have  $\bigwedge e. 0 < e \implies \forall_F n \text{ in } F. \text{dist } (f \ n \ w) \ (g \ w) < e$ 
  by (rule eventually_mono [OF uniform_limitD [OF ulim]]) (use w in auto)
then have  $((\lambda n. 2 * \text{of\_real } \pi * i * f \ n \ w) \longrightarrow 2 * \text{of\_real } \pi * i * g \ w) \ F$ 
  by (rule tendsto_mult_left [OF tendstoI])
then have  $((\lambda u. g \ u / (u - w)) \text{ has\_contour\_integral } 2 * \text{of\_real } \pi * i * g \ w)$ 
  (circlepath z r)
  using has_contour_integral_integral [OF g_cint] tendsto_unique [OF F
f_tends_cig] w
  by fastforce
then have  $((\lambda u. g \ u / (2 * \text{of\_real } \pi * i * (u - w))) \text{ has\_contour\_integral } g \ w)$ 
  (circlepath z r)
  using has_contour_integral_div [where c = 2 * of_real pi * i]
  by (force simp: field_simps)
then show ?thesis
  by (simp add: dist_norm)
qed
show ?thesis
  using Cauchy_next_derivative_circlepath(2) [OF 1 2, simplified]
  by (fastforce simp add: holomorphic_on_open contg intro: that)
qed

```

Version showing that the limit is the limit of the derivatives.

```

proposition has_complex_derivative_uniform_limit:
  fixes z::complex
  assumes cont: eventually  $(\lambda n. \text{continuous\_on } (cball \ z \ r) \ (f \ n) \wedge$ 
     $(\forall w \in \text{ball } z \ r. ((f \ n) \text{ has\_field\_derivative } (f' \ n \ w)) \text{ (at } w))) \ F$ 
    and ulim: uniform_limit  $(cball \ z \ r) \ f \ g \ F$ 
    and F:  $\neg \text{trivial\_limit } F$  and  $0 < r$ 
  obtains g' where
    continuous_on  $(cball \ z \ r) \ g$ 
     $\bigwedge w. w \in \text{ball } z \ r \implies (g \text{ has\_field\_derivative } (g' \ w)) \text{ (at } w) \wedge ((\lambda n. f' \ n \ w) \longrightarrow g' \ w) \ F$ 
  proof -
  let ?conint = contour_integral (circlepath z r)
  have g: continuous_on  $(cball \ z \ r) \ g$  g holomorphic_on ball z r
    by (rule holomorphic_uniform_limit [OF eventually_mono [OF cont] ulim F];
      auto simp: holomorphic_on_open field_differentiable_def)
  then obtain g' where g':  $\bigwedge x. x \in \text{ball } z \ r \implies (g \text{ has\_field\_derivative } g' \ x) \text{ (at } x)$ 
    using DERIV_deriv_iff_has_field_derivative
    by (fastforce simp add: holomorphic_on_open)
  then have derg:  $\bigwedge x. x \in \text{ball } z \ r \implies \text{deriv } g \ x = g' \ x$ 
    by (simp add: DERIV_imp_deriv)
  have tends_f'n_g':  $((\lambda n. f' \ n \ w) \longrightarrow g' \ w) \ F$  if w:  $w \in \text{ball } z \ r$  for w
  proof -
  have eq_f': ?conint  $(\lambda x. f \ n \ x / (x - w)^2) - ?conint (\lambda x. g \ x / (x - w)^2) =$ 

```

```

(f' n w - g' w) * (2 * of_real pi * i)
  if cont_fn: continuous_on (cball z r) (f n)
  and fnd:  $\bigwedge w. w \in \text{ball } z \ r \implies (f \ n \ \text{has\_field\_derivative } f' \ n \ w) \ (\text{at } w)$ 
for n
  proof -
    have hol_fn: f n holomorphic_on ball z r
      using fnd by (force simp: holomorphic_on_open)
    have (f n has_field_derivative 1 / (2 * of_real pi * i) * ?conint ( $\lambda u. f \ n \ u / (u - w)^2$ )) (at w)
      by (rule Cauchy_derivative_integral_circlepath [OF cont_fn hol_fn w])
    then have f': f' n w = 1 / (2 * of_real pi * i) * ?conint ( $\lambda u. f \ n \ u / (u - w)^2$ )
      using DERIV_unique [OF fnd] w by blast
    show ?thesis
      by (simp add: f' Cauchy_contour_integral_circlepath_2 [OF g w] derg [OF w] field_split_simps)
    qed
    define d where d = (r - norm(w - z))^2
    have d > 0
      using w by (simp add: dist_commute dist_norm d_def)
    have dle: d ≤ cmod ((y - w)^2) if r = cmod (z - y) for y
    proof -
      have cmod (w - z) ≤ cmod (z - y)
        by (metis dist_commute dist_norm mem_ball order_less_imp_le that w)
      moreover have cmod (z - y) - cmod (w - z) ≤ cmod (y - w)
      by (metis diff_add_cancel diff_diff_eq2 norm_minus_commute norm_triangle_ineq2)
      ultimately show ?thesis
        using that by (simp add: d_def norm_power power_mono)
    qed
    have 1:  $\forall_F \ n \ \text{in } F. (\lambda x. f \ n \ x / (x - w)^2) \ \text{contour\_integrable\_on } \text{circlepath } z$ 
    r
      by (force simp: holomorphic_on_open intro: w Cauchy_derivative_integral_circlepath eventually_mono [OF cont])
    have 2: uniform_limit (sphere z r) ( $\lambda n \ x. f \ n \ x / (x - w)^2$ ) ( $\lambda x. g \ x / (x - w)^2$ ) F
      unfolding uniform_limit_iff
    proof clarify
      fix e::real
      assume e > 0
      with  $\langle r > 0 \rangle$ 
      have  $\forall_F \ n \ \text{in } F. \forall x. x \neq w \implies \text{cmod } (z - x) = r \implies \text{cmod } (f \ n \ x - g \ x) < e * \text{cmod } ((x - w)^2)$ 
      by (force simp:  $\langle 0 < d \rangle$  dist_norm dle intro: less_le_trans eventually_mono [OF uniform_limitD [OF ulim], of e*d])
      with  $\langle r > 0 \rangle \langle e > 0 \rangle$ 
      show  $\forall_F \ n \ \text{in } F. \forall x \in \text{sphere } z \ r. \text{dist } (f \ n \ x / (x - w)^2) (g \ x / (x - w)^2) < e$ 
      by (simp add: norm_divide field_split_simps sphere_def dist_norm)
    qed
    have (( $\lambda n. \text{contour\_integral } (\text{circlepath } z \ r) (\lambda x. f \ n \ x / (x - w)^2)$ )

```

```

       $\longrightarrow$  contour_integral (circlepath z r) (( $\lambda x. g x / (x - w)^2$ )) F
    by (rule contour_integral_uniform_limit_circlepath [OF 1 2 F <0 < r>])
    then have tendsto_0: (( $\lambda n. 1 / (2 * \text{of\_real } \pi * i) * (?conint (\lambda x. f n x / (x - w)^2) - ?conint (\lambda x. g x / (x - w)^2))$ )  $\longrightarrow$  0) F
    using Lim_null by (force intro!: tendsto_mult_right_zero)
    have (( $\lambda n. f' n w - g' w$ )  $\longrightarrow$  0) F
    apply (rule Lim_transform_eventually [OF tendsto_0])
    apply (force simp: divide_simps intro: eq_f' eventually_mono [OF cont])
    done
  then show ?thesis using Lim_null by blast
qed
obtain g' where  $\bigwedge w. w \in \text{ball } z \ r \implies (g \text{ has\_field\_derivative } (g' w)) \text{ (at } w) \wedge$ 
(( $\lambda n. f' n w$ )  $\longrightarrow$   $g' w$ ) F
  by (blast intro: tends_f'n_g' g')
then show ?thesis using g
  using that by blast
qed

```

4.9 Some more simple/convenient versions for applications

lemma holomorphic_uniform_sequence:

```

  assumes S: open S
    and hol_fn:  $\bigwedge n. (f n) \text{ holomorphic\_on } S$ 
    and ulim_g:  $\bigwedge x. x \in S \implies \exists d. 0 < d \wedge \text{cball } x \ d \subseteq S \wedge \text{uniform\_limit}$ 
(cball x d) f g sequentially
  shows g holomorphic_on S
proof -
  have  $\exists f'. (g \text{ has\_field\_derivative } f') \text{ (at } z) \text{ if } z \in S$  for z
proof -
  obtain r where  $0 < r$  and r:  $\text{cball } z \ r \subseteq S$ 
    and ul: uniform_limit (cball z r) f g sequentially
  using ulim_g [OF <z ∈ S>] by blast
  have *:  $\forall_F n \text{ in sequentially. continuous\_on } (\text{cball } z \ r) (f n) \wedge f n \text{ holomorphic\_on ball } z \ r$ 
  by (smt (verit, best) ball_subset_cball hol_fn holomorphic_on_imp_continuous_on

    holomorphic_on_subset not_eventuallyD r)
  show ?thesis
  using <0 < r> centre_in_ball ul
  by (auto simp: holomorphic_on_open intro: holomorphic_uniform_limit [OF
*])
qed
with S show ?thesis
  by (simp add: holomorphic_on_open)
qed

```

lemma has_complex_derivative_uniform_sequence:

```

  fixes S :: complex set
  assumes S: open S

```

```

    and hfd:  $\bigwedge n x. x \in S \implies ((f n) \text{ has\_field\_derivative } f' n x) \text{ (at } x)$ 
    and ulim_g:  $\bigwedge x. x \in S$ 
       $\implies \exists d. 0 < d \wedge \text{cball } x d \subseteq S \wedge \text{uniform\_limit } (\text{cball } x d) f g \text{ sequentially}$ 
    shows  $\exists g'. \forall x \in S. (g \text{ has\_field\_derivative } g' x) \text{ (at } x) \wedge ((\lambda n. f' n x) \longrightarrow g' x) \text{ sequentially}$ 
  proof -
    have y:  $\exists y. (g \text{ has\_field\_derivative } y) \text{ (at } z) \wedge (\lambda n. f' n z) \longrightarrow y$  if  $z \in S$ 
  for z
    proof -
      obtain r where  $0 < r$  and r:  $\text{cball } z r \subseteq S$ 
        and ul:  $\text{uniform\_limit } (\text{cball } z r) f g \text{ sequentially}$ 
      using ulim_g [OF  $\langle z \in S \rangle$ ] by blast
      have *:  $\forall_F n \text{ in sequentially. continuous\_on } (\text{cball } z r) (f n) \wedge$ 
         $(\forall w \in \text{ball } z r. ((f n) \text{ has\_field\_derivative } (f' n w)))$ 
      (at w)
      proof (intro eventuallyI conjI ballI)
        show  $\text{continuous\_on } (\text{cball } z r) (f x)$  for x
        by (meson S continuous_on_subset hfd holomorphic_on_imp_continuous_on
          holomorphic_on_open r)
        show  $w \in \text{ball } z r \implies (f x \text{ has\_field\_derivative } f' x w) \text{ (at } w)$  for w x
        using ball_subset_cball hfd r by blast
      qed
      show ?thesis
      by (rule has_complex_derivative_uniform_limit [OF *, of g]) (use  $\langle 0 < r \rangle$ 
        ul in  $\langle \text{force} \rangle$ )
    qed
    show ?thesis
    by (rule bchoice) (blast intro: y)
  qed

```

4.10 On analytic functions defined by a series

lemma series_and_derivative_comparison:

```

  fixes S :: complex set
  assumes S: open S
    and h: summable h
    and hfd:  $\bigwedge n x. x \in S \implies (f n \text{ has\_field\_derivative } f' n x) \text{ (at } x)$ 
    and to_g:  $\forall_F n \text{ in sequentially. } \forall x \in S. \text{norm } (f n x) \leq h n$ 
  obtains g g' where  $\forall x \in S. ((\lambda n. f n x) \text{ sums } g x) \wedge ((\lambda n. f' n x) \text{ sums } g' x)$ 
   $\wedge (g \text{ has\_field\_derivative } g' x) \text{ (at } x)$ 
  proof -
    obtain g where g:  $\text{uniform\_limit } S (\lambda n x. \sum i < n. f i x) g \text{ sequentially}$ 
      using Weierstrass_m_test_ev [OF to_g h] by force
    have *:  $\exists d > 0. \text{cball } x d \subseteq S \wedge \text{uniform\_limit } (\text{cball } x d) (\lambda n x. \sum i < n. f i x)$ 
      g sequentially
      if  $x \in S$  for x
      using open_contains_cball [of S]  $\langle x \in S \rangle$  S g uniform_limit_on_subset by
      blast
    have  $\bigwedge x. x \in S \implies (\lambda n. \sum i < n. f i x) \longrightarrow g x$ 

```

```

  by (metis tendsto_uniform_limitI [OF g])
  moreover have  $\exists g'. \forall x \in S. (g \text{ has\_field\_derivative } g' x) (at x) \wedge (\lambda n. \sum i < n. f' i x) \longrightarrow g' x$ 
  by (rule has_complex_derivative_uniform_sequence [OF S]) (auto intro: * hfd DERIV_sum)+
  ultimately show ?thesis
  by (metis sums_def that)
qed

```

A version where we only have local uniform/comparative convergence.

lemma series_and_derivative_comparison_local:

```

  fixes S :: complex set
  assumes S: open S
  and hfd:  $\bigwedge n x. x \in S \implies (f n \text{ has\_field\_derivative } f' n x) (at x)$ 
  and to_g:  $\bigwedge x. x \in S \implies \exists d h. 0 < d \wedge \text{summable } h \wedge (\forall_F n \text{ in sequentially. } \forall y \in \text{ball } x d \cap S. \text{norm } (f n y) \leq h n)$ 
  shows  $\exists g g'. \forall x \in S. ((\lambda n. f n x) \text{ sums } g x) \wedge ((\lambda n. f' n x) \text{ sums } g' x) \wedge (g \text{ has\_field\_derivative } g' x) (at x)$ 
  proof -
  have  $\exists y. (\lambda n. f n z) \text{ sums } (\sum n. f n z) \wedge (\lambda n. f' n z) \text{ sums } y \wedge ((\lambda x. \sum n. f n x) \text{ has\_field\_derivative } y) (at z)$ 
  if  $z \in S$  for  $z$ 
  proof -
  obtain  $d h$  where  $0 < d$  summable  $h$  and  $le\_h: \forall_F n \text{ in sequentially. } \forall y \in \text{ball } z d \cap S. \text{norm } (f n y) \leq h n$ 
  using to_g  $\langle z \in S \rangle$  by meson
  then obtain  $r$  where  $r > 0$  and  $r: \text{ball } z r \subseteq \text{ball } z d \cap S$  using  $\langle z \in S \rangle S$ 
  by (metis Int_iff open_ball centre_in_ball open_Int open_contains_ball_eq)
  have 1: open  $(\text{ball } z d \cap S)$ 
  by (simp add: open_Int S)
  have 2:  $\bigwedge n x. x \in \text{ball } z d \cap S \implies (f n \text{ has\_field\_derivative } f' n x) (at x)$ 
  by (auto simp: hfd)
  obtain  $g g'$  where  $gg': \forall x \in \text{ball } z d \cap S. ((\lambda n. f n x) \text{ sums } g x) \wedge ((\lambda n. f' n x) \text{ sums } g' x) \wedge (g \text{ has\_field\_derivative } g' x) (at x)$ 
  by (auto intro: le_h series_and_derivative_comparison [OF 1  $\langle \text{summable } h \rangle$  hfd])
  then have  $(\lambda n. f' n z) \text{ sums } g' z$ 
  by (meson  $\langle 0 < r \rangle$  centre_in_ball contra_subsetD r)
  moreover have  $(\lambda n. f n z) \text{ sums } (\sum n. f n z)$ 
  using summable_sums centre_in_ball  $\langle 0 < d \rangle \langle \text{summable } h \rangle le\_h$ 
  by (metis (full_types) Int_iff gg' summable_def that)
  moreover have  $((\lambda x. \sum n. f n x) \text{ has\_field\_derivative } g' z) (at z)$ 
  by (metis (no_types, lifting) 1 r  $\langle 0 < r \rangle$  gg' has_field_derivative_transform_within_open

  open_contains_ball_eq sums_unique)
  ultimately show ?thesis by auto
qed
then show ?thesis

```

by meson
qed

Sometimes convenient to compare with a complex series of positive reals.
(?)

lemma series_and_derivative_comparison_complex:

fixes $S :: \text{complex set}$
 assumes $S: \text{open } S$
 and $hfd: \bigwedge n x. x \in S \implies (f \ n \ \text{has_field_derivative } f' \ n \ x) \ (at \ x)$
 and $to_g: \bigwedge x. x \in S \implies \exists d h. 0 < d \wedge \text{summable } h \wedge \text{range } h \subseteq \mathbb{R}_{\geq 0} \wedge$
 $(\forall_F n \text{ in sequentially. } \forall y \in \text{ball } x \ d \cap S. \text{cmod}(f \ n \ y) \leq \text{cmod } (h \ n))$
 shows $\exists g \ g'. \forall x \in S. ((\lambda n. f \ n \ x) \ \text{sums } g \ x) \wedge ((\lambda n. f' \ n \ x) \ \text{sums } g' \ x) \wedge (g$
 $\text{has_field_derivative } g' \ x) \ (at \ x)$
 apply (rule series_and_derivative_comparison_local [OF S hfd], assumption)
 apply (rule ex_forward [OF to_g], assumption)
 apply (erule exE)
 apply (rule_tac $x = Re \circ h$ in exI)
 apply (force simp: summable_Re o_def nonneg_Reals_cmod_eq_Re image_subset_iff)
 done

Sometimes convenient to compare with a complex series of positive reals.
(?)

lemma series_differentiable_comparison_complex:

fixes $S :: \text{complex set}$
 assumes $S: \text{open } S$
 and $hfd: \bigwedge n x. x \in S \implies f \ n \ \text{field_differentiable } (at \ x)$
 and $to_g: \bigwedge x. x \in S \implies \exists d h. 0 < d \wedge \text{summable } h \wedge \text{range } h \subseteq \mathbb{R}_{\geq 0} \wedge (\forall_F$
 $n \text{ in sequentially. } \forall y \in \text{ball } x \ d \cap S. \text{cmod}(f \ n \ y) \leq \text{cmod } (h \ n))$
 obtains $g \ \text{where } \forall x \in S. ((\lambda n. f \ n \ x) \ \text{sums } g \ x) \wedge g \ \text{field_differentiable } (at \ x)$
 proof -
 have $hfd': \bigwedge n x. x \in S \implies (f \ n \ \text{has_field_derivative } \text{deriv } (f \ n) \ x) \ (at \ x)$
 using $hfd \ \text{field_differentiable_derivI}$ by blast
 show ?thesis
 by (metis field_differentiable_def that series_and_derivative_comparison_complex [OF S hfd' to_g])
 qed

In particular, a power series is analytic inside circle of convergence.

lemma power_series_and_derivative_0:

fixes $a :: \text{nat} \Rightarrow \text{complex}$ and $r :: \text{real}$
 assumes $\text{summable } (\lambda n. a \ n * r^n)$
 shows $\exists g \ g'. \forall z. \text{cmod } z < r \implies$
 $((\lambda n. a \ n * z^n) \ \text{sums } g \ z) \wedge ((\lambda n. \text{of_nat } n * a \ n * z^{n-1}) \ \text{sums}$
 $g' \ z) \wedge (g \ \text{has_field_derivative } g' \ z) \ (at \ z)$
 proof (cases $0 < r$)
 case True
 have $der: \bigwedge n z. ((\lambda x. a \ n * x^n) \ \text{has_field_derivative } \text{of_nat } n * a \ n * z^{n-1}) \ (at \ z)$

```

    by (rule derivative_eq_intros | simp)+
  have y_le: cmod y ≤ cmod (of_real r + of_real (cmod z)) / 2
    if cmod (z - y) * 2 < r - cmod z for z y
  by (smt (verit, best) field_sum_of_halves norm_minus_commute norm_of_real
norm_triangle_ineq2 of_real_add that)
  have summable (λn. a n * complex_of_real r ^ n)
    using assms ⟨r > 0⟩ by simp
  moreover have ∧z. cmod z < r ⇒ cmod ((of_real r + of_real (cmod z)) /
2) < cmod (of_real r)
    using ⟨r > 0⟩
    by (simp flip: of_real_add)
  ultimately have sum: ∧z. cmod z < r ⇒ summable (λn. of_real (cmod (a
n)) * ((of_real r + complex_of_real (cmod z)) / 2) ^ n)
    by (rule power_series_conv_imp_absconv_weak)
  have ∃g g'. ∀z ∈ ball 0 r. (λn. (a n) * z ^ n) sums g z ∧
    (λn. of_nat n * (a n) * z ^ (n - 1)) sums g' z ∧ (g has_field_derivative
g' z) (at z)
    apply (rule series_and_derivative_comparison_complex [OF open_ball der])
    apply (rule_tac x=(r - norm z)/2 in exI)
    apply (rule_tac x=λn. of_real(norm(a n)*((r + norm z)/2)^n) in exI)
    using ⟨r > 0⟩
    apply (auto simp: sum_eventually_sequentially norm_mult norm_power
dist_norm intro!: mult_left_mono power_mono y_le)
  done
  then show ?thesis
    by (simp add: ball_def)
next
case False then show ?thesis
  unfolding not_less using less_le_trans norm_not_less_zero by blast
qed

```

proposition *power_series_and_derivative*:

```

  fixes a :: nat ⇒ complex and r::real
  assumes summable (λn. a n * r ^ n)
  obtains g g' where ∀z ∈ ball w r.
    ((λn. a n * (z - w) ^ n) sums g z) ∧ ((λn. of_nat n * a n * (z - w) ^
(n - 1)) sums g' z) ∧
    (g has_field_derivative g' z) (at z)
  using power_series_and_derivative_0 [OF assms]
  apply clarify
  apply (rule_tac g=(λz. g(z - w)) in that)
  using DERIV_shift [where z=-w]
  apply (auto simp: norm_minus_commute Ball_def dist_norm)
  done

```

proposition *power_series_holomorphic*:

```

  assumes ∧w. w ∈ ball z r ⇒ ((λn. a n*(w - z) ^ n) sums f w)
  shows f holomorphic_on ball z r
proof -

```



```

have  $\exists f'. (f \text{ has\_field\_derivative } f') (at \ w)$  if  $w: \text{dist } z \ w < r$  for  $w$ 
proof -
  have  $wz: \text{cmod } (w - z) < r$  using  $w$ 
  by (auto simp: field_split_simps dist_norm norm_minus_commute)
  then have  $0 \leq r$ 
  by (meson less_eq_real_def norm_ge_zero order_trans)
  have  $inb: z + \text{complex\_of\_real } ((\text{dist } z \ w + r) / 2) \in \text{ball } z \ r$ 
  using  $w$  by (simp add: dist_norm ‹ $0 \leq r$ › flip: of_real_add)
  have  $sum: \text{summable } (\lambda n. a \ n * \text{of\_real } (((\text{cmod } (z - w) + r) / 2) ^ n))$ 
  using  $assms [OF inb]$  by (force simp: summable_def dist_norm)
  obtain  $g \ g'$  where  $gg': \bigwedge u. u \in \text{ball } z ((\text{cmod } (z - w) + r) / 2) \implies$ 
     $(\lambda n. a \ n * (u - z) ^ n) \text{ sums } g \ u \wedge$ 
     $(\lambda n. \text{of\_nat } n * a \ n * (u - z) ^ (n - 1)) \text{ sums } g' \ u \wedge (g$ 
has_field_derivative  $g' \ u) (at \ u)$ 
  by (rule power_series_and_derivative [OF sum, of z]) fastforce
  have [simp]:  $g \ u = f \ u$  if  $\text{cmod } (u - w) < (r - \text{cmod } (z - w)) / 2$  for  $u$ 
  proof -
    have  $less: \text{cmod } (z - u) * 2 < \text{cmod } (z - w) + r$ 
    using  $that \ \text{dist\_triangle2} [of \ z \ u \ w]$ 
    by (simp add: dist_norm [symmetric] algebra_simps)
    have  $(\lambda n. a \ n * (u - z) ^ n) \text{ sums } g \ u \ (\lambda n. a \ n * (u - z) ^ n) \text{ sums } f \ u$ 
    using  $gg' [of \ u] \ less \ w$  by (auto simp: assms dist_norm)
    then show ?thesis
    by (metis sums_unique2)
  qed
  have  $(f \text{ has\_field\_derivative } g' \ w) (at \ w)$ 
  by (rule has_field_derivative_transform_within [where  $d=(r - \text{norm}(z - w))/2$ ])
  (use  $w \ gg' [of \ w]$  in ‹(force simp: dist_norm)+›)
  then show ?thesis ..
  qed
  then show ?thesis by (simp add: holomorphic_on_open)
  qed

```

corollary *holomorphic_iff_power_series:*

```

f holomorphic_on ball z r  $\longleftrightarrow$ 
 $(\forall w \in \text{ball } z \ r. (\lambda n. (\text{deriv } ^ n) f \ z / (\text{fact } n) * (w - z) ^ n) \text{ sums } f \ w)$ 
apply (intro iffI ballI holomorphic_power_series, assumption+)
apply (force intro: power_series_holomorphic [where  $a = \lambda n. (\text{deriv } ^ n) f \ z / (\text{fact } n)$ ])
done

```

lemma *power_series_analytic:*

```

 $(\bigwedge w. w \in \text{ball } z \ r \implies (\lambda n. a \ n * (w - z) ^ n) \text{ sums } f \ w) \implies f \ \text{analytic\_on ball } z \ r$ 
by (force simp: analytic_on_open intro!: power_series_holomorphic)

```

lemma *analytic_iff_power_series:*

```

f analytic_on ball z r  $\longleftrightarrow$ 

```

$(\forall w \in \text{ball } z \ r. (\lambda n. (\text{deriv } \hat{\sim} n) f z / (\text{fact } n) * (w - z)^{\hat{n}}) \text{ sums } f w)$
by (*simp add: analytic_on_open holomorphic_iff_power_series*)

4.11 Equality between holomorphic functions, on open ball then connected set

lemma *holomorphic_fun_eq_on_ball*:

$\llbracket f \text{ holomorphic_on ball } z \ r; g \text{ holomorphic_on ball } z \ r;$
 $w \in \text{ball } z \ r;$
 $\bigwedge n. (\text{deriv } \hat{\sim} n) f z = (\text{deriv } \hat{\sim} n) g z \rrbracket$
 $\implies f w = g w$

by (*auto simp: holomorphic_iff_power_series sums_unique2 [of $\lambda n. (\text{deriv } \hat{\sim} n) f z / (\text{fact } n) * (w - z)^{\hat{n}}$]*)

lemma *holomorphic_fun_eq_0_on_ball*:

$\llbracket f \text{ holomorphic_on ball } z \ r; w \in \text{ball } z \ r;$
 $\bigwedge n. (\text{deriv } \hat{\sim} n) f z = 0 \rrbracket$
 $\implies f w = 0$

using *holomorphic_fun_eq_on_ball [where $g = \lambda z. 0$]* **by** *simp*

lemma *holomorphic_fun_eq_0_on_connected*:

assumes *holf: f holomorphic_on S and open S*
and *cons: connected S*
and *der: $\bigwedge n. (\text{deriv } \hat{\sim} n) f z = 0$*
and *$z \in S \ w \in S$*
shows *$f w = 0$*

proof –

have **: ball $x \ e \subseteq (\bigcap n. \{w \in S. (\text{deriv } \hat{\sim} n) f w = 0\})$*
if $\forall u. (\text{deriv } \hat{\sim} u) f x = 0$ *ball $x \ e \subseteq S$ for $x \ e$*

proof –

have *$(\text{deriv } \hat{\sim} m) ((\text{deriv } \hat{\sim} n) f) x = 0$ for $m \ n$*
by (*metis funpow_add o_apply that(1)*)
then have $\bigwedge x' \ n. \text{dist } x \ x' < e \implies (\text{deriv } \hat{\sim} n) f x' = 0$
using *$\langle \text{open } S \rangle$*

by (*meson holf holomorphic_fun_eq_0_on_ball holomorphic_higher_deriv holomorphic_on_subset mem_ball that(2)*)

with that show *?thesis by auto*

qed

obtain *e where $e > 0$ and $e: \text{ball } w \ e \subseteq S$ using* *openE [OF $\langle \text{open } S \rangle \ \langle w \in S \rangle$]*

then have *holfb: f holomorphic_on ball $w \ e$*

using *holf holomorphic_on_subset by blast*

have *open $(\bigcap n. \{w \in S. (\text{deriv } \hat{\sim} n) f w = 0\})$*

using *$\langle \text{open } S \rangle$*

apply (*simp add: open_contains_ball Ball_def image_iff*)

by (*metis (mono_tags) * mem_Collect_eq*)

then have *openin (top_of_set S) $(\bigcap n. \{w \in S. (\text{deriv } \hat{\sim} n) f w = 0\})$*

by (*force intro: open_subset*)

moreover have *closedin (top_of_set S) $(\bigcap n. \{w \in S. (\text{deriv } \hat{\sim} n) f w = 0\})$*

```

using assms
by (auto intro: continuous_closedin_preimage_constant holomorphic_on_imp_continuous_on
holomorphic_higher_deriv)
moreover have  $(\bigcap n. \{w \in S. (\text{deriv } \sim n) f w = 0\}) = S \implies f w = 0$ 
using  $\langle e > 0 \rangle$  e by (force intro: holomorphic_fun_eq_0_on_ball [OF holfb])
ultimately show ?thesis
using cons der  $\langle z \in S \rangle$ 
by (auto simp add: connected_clopen)
qed

```

```

lemma holomorphic_fun_eq_on_connected:
assumes f holomorphic_on S g holomorphic_on S and open S connected S
and  $\bigwedge n. (\text{deriv } \sim n) f z = (\text{deriv } \sim n) g z$ 
and  $z \in S \ w \in S$ 
shows  $f w = g w$ 
proof (rule holomorphic_fun_eq_0_on_connected [of  $\lambda x. f x - g x$  S z, simplified])
show  $(\lambda x. f x - g x)$  holomorphic_on S
by (intro assms holomorphic_intros)
show  $\bigwedge n. (\text{deriv } \sim n) (\lambda x. f x - g x) z = 0$ 
using assms higher_deriv_diff by auto
qed (use assms in auto)

```

```

lemma holomorphic_fun_eq_const_on_connected:
assumes holf: f holomorphic_on S and open S
and cons: connected S
and der:  $\bigwedge n. 0 < n \implies (\text{deriv } \sim n) f z = 0$ 
and  $z \in S \ w \in S$ 
shows  $f w = f z$ 
proof (rule holomorphic_fun_eq_0_on_connected [of  $\lambda w. f w - f z$  S z, simplified])
show  $(\lambda w. f w - f z)$  holomorphic_on S
by (intro assms holomorphic_intros)
show  $\bigwedge n. (\text{deriv } \sim n) (\lambda w. f w - f z) z = 0$ 
by (subst higher_deriv_diff) (use assms in  $\langle \text{auto } \text{intro}: \text{holomorphic\_intros} \rangle$ )
qed (use assms in auto)

```

4.12 Some basic lemmas about poles/singularities

```

lemma pole_lemma:
assumes holf: f holomorphic_on S and a:  $a \in \text{interior } S$ 
shows  $(\lambda z. \text{if } z = a \text{ then } \text{deriv } f a$ 
else } (f z - f a) / (z - a)) holomorphic_on S (is ?F holomorphic_on
S)
proof -
have *: ?F field_differentiable (at u within S) if  $u \in S \ u \neq a$  for u
proof -
have fcd: f field_differentiable at u within S
using holf holomorphic_on_def by (simp add:  $\langle u \in S \rangle$ )

```

```

have cd: ( $\lambda z. (f z - f a) / (z - a)$ ) field_differentiable at u within S
  by (rule fcd_derivative_intros | simp add: that) +
have 0 < dist a u using that dist_nz by blast
then show ?thesis
  by (rule field_differentiable_transform_within [OF _ _ _ cd]) (auto simp:
  ‹u ∈ S›)
qed
moreover
have ?F field_differentiable at a if 0 < e ball a e ⊆ S for e
proof -
  have holfb: f holomorphic_on ball a e
    by (rule holomorphic_on_subset [OF hol ‹ball a e ⊆ S›])
  have 2: ?F holomorphic_on ball a e - {a}
    using mem_ball that
    by (auto simp add: holomorphic_on_def simp flip: field_differentiable_def
intro: * field_differentiable_within_subset)
  have isCont ( $\lambda z. \text{if } z = a \text{ then deriv } f \ a \text{ else } (f z - f a) / (z - a)$ ) x
    if dist a x < e for x
proof (cases x=a)
  case True
    then have f field_differentiable at a
      using holfb ‹0 < e› holomorphic_on_imp_differentiable_at by auto
    with True show ?thesis
      by (smt (verit) DERIV_deriv_iff_field_differentiable LIM_equal continuous_at
has_field_derivativeD)
  next
    case False with 2 that show ?thesis
      by (simp add: field_differentiable_imp_continuous_at holomorphic_on_imp_differentiable_at
open_Diff)
  qed
then have 1: continuous_on (ball a e) ?F
  by (clarsimp simp: continuous_on_eq_continuous_at)
have ?F holomorphic_on ball a e
  by (auto intro: no_isolated_singularity [OF 1 2])
with that show ?thesis
  by (simp add: holomorphic_on_open field_differentiable_def [symmetric]
field_differentiable_at_within)
qed
ultimately show ?thesis
  by (metis (no_types, lifting) holomorphic_onI a field_differentiable_at_within
interior_subset openE open_interior subset_iff)
qed

lemma pole_theorem:
assumes holg: g holomorphic_on S and a: a ∈ interior S
and eq:  $\bigwedge z. z \in S - \{a\} \implies g z = (z - a) * f z$ 
shows ( $\lambda z. \text{if } z = a \text{ then deriv } g \ a$ 
  else  $f z - g a / (z - a)$ ) holomorphic_on S
using pole_lemma [OF holg a]

```

by (rule holomorphic_transform) (simp add: eq_field_split_simps)

lemma pole_lemma_open:

assumes f holomorphic_on S open S

shows $(\lambda z. \text{if } z = a \text{ then deriv } f \ a \text{ else } (f \ z - f \ a)/(z - a))$ holomorphic_on S

proof (cases $a \in S$)

case True with assms interior_eq_pole_lemma

show ?thesis by fastforce

next

case False with assms show ?thesis

apply (simp add: holomorphic_on_def field_differentiable_def [symmetric], clarify)

apply (rule field_differentiable_transform_within [where $f = \lambda z. (f \ z - f \ a)/(z - a)$ and $d = 1$])

apply (rule derivative_intros | force)+

done

qed

lemma pole_theorem_open:

assumes $holg: g$ holomorphic_on S and $S: \text{open } S$

and $eq: \bigwedge z. z \in S - \{a\} \implies g \ z = (z - a) * f \ z$

shows $(\lambda z. \text{if } z = a \text{ then deriv } g \ a$

else $f \ z - g \ a/(z - a))$ holomorphic_on S

using pole_lemma_open [OF $holg \ S$]

by (rule holomorphic_transform) (auto simp: eq_divide_simps)

lemma pole_theorem_0:

assumes $holg: g$ holomorphic_on S and $a: a \in \text{interior } S$

and $eq: \bigwedge z. z \in S - \{a\} \implies g \ z = (z - a) * f \ z$

and [simp]: $f \ a = \text{deriv } g \ a \ g \ a = 0$

shows f holomorphic_on S

using pole_theorem [OF $holg \ a \ eq$]

by (rule holomorphic_transform) (auto simp: eq_field_split_simps)

lemma pole_theorem_open_0:

assumes $holg: g$ holomorphic_on S and $S: \text{open } S$

and $eq: \bigwedge z. z \in S - \{a\} \implies g \ z = (z - a) * f \ z$

and [simp]: $f \ a = \text{deriv } g \ a \ g \ a = 0$

shows f holomorphic_on S

using pole_theorem_open [OF $holg \ S \ eq$]

by (rule holomorphic_transform) (auto simp: eq_field_split_simps)

lemma pole_theorem_analytic:

assumes $g: g$ analytic_on S

and $eq: \bigwedge z. z \in S$

$\implies \exists d. 0 < d \wedge (\forall w \in \text{ball } z \ d - \{a\}. g \ w = (w - a) * f \ w)$

shows $(\lambda z. \text{if } z = a \text{ then deriv } g \ a \text{ else } f \ z - g \ a/(z - a))$ analytic_on S (is ?F analytic_on S)

unfolding analytic_on_def

proof
fix x
assume $x \in S$
with g **obtain** e **where** $0 < e$ **and** e : g holomorphic_on ball x e
by (auto simp add: analytic_on_def)
obtain d **where** $0 < d$ **and** d : $\bigwedge w. w \in \text{ball } x \ d - \{a\} \implies g \ w = (w - a) * f \ w$
using $\langle x \in S \rangle$ eq **by** blast
have ?F holomorphic_on ball x (min d e)
using $d \ e \ \langle x \in S \rangle$ **by** (fastforce simp: holomorphic_on_subset subset_ball intro!: pole_theorem_open)
then show $\exists e > 0. ?F$ holomorphic_on ball x e
using $\langle 0 < d \rangle \ \langle 0 < e \rangle$ not_le **by** fastforce
qed

lemma pole_theorem_analytic_0:
assumes g : g analytic_on S
and eq: $\bigwedge z. z \in S \implies \exists d. 0 < d \wedge (\forall w \in \text{ball } z \ d - \{a\}. g \ w = (w - a) * f \ w)$
and [simp]: $f \ a = \text{deriv } g \ a \ g \ a = 0$
shows f analytic_on S
proof –
have [simp]: $(\lambda z. \text{if } z = a \ \text{then } \text{deriv } g \ a \ \text{else } f \ z - g \ a / (z - a)) = f$
by auto
show ?thesis
using pole_theorem_analytic [OF g eq] **by** simp
qed

lemma pole_theorem_analytic_open_superset:
assumes g : g analytic_on S **and** $S \subseteq T$ open T
and eq: $\bigwedge z. z \in T - \{a\} \implies g \ z = (z - a) * f \ z$
shows $(\lambda z. \text{if } z = a \ \text{then } \text{deriv } g \ a \ \text{else } f \ z - g \ a / (z - a))$ analytic_on S
proof (rule pole_theorem_analytic [OF g])
fix z
assume $z \in S$
then obtain e **where** $0 < e$ **and** e : ball z $e \subseteq T$
using assms openE **by** blast
then show $\exists d > 0. \forall w \in \text{ball } z \ d - \{a\}. g \ w = (w - a) * f \ w$
using eq **by** auto
qed

lemma pole_theorem_analytic_open_superset_0:
assumes g : g analytic_on S $S \subseteq T$ open T $\bigwedge z. z \in T - \{a\} \implies g \ z = (z - a) * f \ z$
and [simp]: $f \ a = \text{deriv } g \ a \ g \ a = 0$
shows f analytic_on S
proof –
have [simp]: $(\lambda z. \text{if } z = a \ \text{then } \text{deriv } g \ a \ \text{else } f \ z - g \ a / (z - a)) = f$

```

  by auto
  have (λz. if z = a then deriv g a else f z - g a/(z - a)) analytic_on S
    by (rule pole_theorem_analytic_open_superset [OF g])
  then show ?thesis by simp
qed

```

4.13 General, homology form of Cauchy's theorem

Proof is based on Dixon's, as presented in Lang's "Complex Analysis" book (page 147).

lemma *contour_integral_continuous_on_linepath_2D*:

assumes *open U and cont_dw*: $\bigwedge w. w \in U \implies F w$ *contour_integrable_on* (linepath a b)

and *cond_uu*: *continuous_on* (U × U) (λ(x,y). F x y)

and *abu*: *closed_segment* a b ⊆ U

shows *continuous_on* U (λw. *contour_integral* (linepath a b) (F w))

proof –

have *: $\exists d > 0. \forall x' \in U. \text{dist } x' w < d \implies$

$$\begin{aligned} & \text{dist } (\text{contour_integral } (\text{linepath } a \ b) \ (F \ x')) \\ & \quad (\text{contour_integral } (\text{linepath } a \ b) \ (F \ w)) \leq \varepsilon \end{aligned}$$

if $w \in U$ $0 < \varepsilon$ $a \neq b$ **for** $w \ \varepsilon$

proof –

obtain δ **where** $\delta > 0$ **and** δ : *cball* w $\delta \subseteq U$ **using** *open_contains_cball* $\langle \text{open } U \rangle$ $\langle w \in U \rangle$ **by** *force*

let ?TZ = *cball* w $\delta \times$ *closed_segment* a b

have *uniformly_continuous_on* ?TZ (λ(x,y). F x y)

proof (*rule compact_uniformly_continuous*)

show *continuous_on* ?TZ (λ(x,y). F x y)

by (*rule continuous_on_subset*[OF *cond_uu*]) (*use SigmaE* δ *abu* **in** *blast*)

show *compact* ?TZ

by (*simp add: compact_Times*)

qed

then obtain η **where** $\eta > 0$

and η : $\bigwedge x \ x'. \llbracket x \in ?TZ; x' \in ?TZ; \text{dist } x' \ x < \eta \rrbracket \implies$

$$\text{dist } ((\lambda(x,y). F \ x \ y) \ x') \ ((\lambda(x,y). F \ x \ y) \ x) < \varepsilon / \text{norm}(b - a)$$

using $\langle 0 < \varepsilon \rangle$ $\langle a \neq b \rangle$

by (*auto elim: uniformly_continuous_onE* [**where** $e = \varepsilon / \text{norm}(b - a)$])

have η : $\llbracket \text{norm}(w - x1) \leq \delta; x2 \in \text{closed_segment } a \ b;$

$\text{norm}(w - x1') \leq \delta; x2' \in \text{closed_segment } a \ b; \text{norm}((x1', x2') - (x1, x2)) < \eta \rrbracket$

$$\implies \text{norm}(F \ x1' \ x2' - F \ x1 \ x2) \leq \varepsilon / \text{cmod}(b - a)$$

for $x1 \ x2 \ x1' \ x2'$

using η [*of* (x1,x2) (x1',x2')] **by** (*force simp: dist_norm*)

have *le_ee*: $\text{cmod}(\text{contour_integral}(\text{linepath } a \ b) (\lambda x. F \ x' \ x - F \ w \ x)) \leq \varepsilon$

if $x' \in U$ $\text{cmod}(x' - w) < \delta$ $\text{cmod}(x' - w) < \eta$ **for** x'

proof –

have (λx. F x' x - F w x) *contour_integrable_on* linepath a b

by (*simp add:* $\langle w \in U \rangle$ *cont_dw* *contour_integrable_diff* *that*)

then have $\text{cmod}(\text{contour_integral}(\text{linepath } a \ b) (\lambda x. F \ x' \ x - F \ w \ x)) \leq$

```

 $\varepsilon / \text{norm}(b - a) * \text{norm}(b - a)$ 
  using has_contour_integral_bound_linepath [OF has_contour_integral_integral
  _  $\eta$ ]
    using  $\langle 0 < \varepsilon \rangle \langle 0 < \delta \rangle$  that by (force simp: norm_minus_commute)
    also have ... =  $\varepsilon$  using  $\langle a \neq b \rangle$  by simp
    finally show ?thesis .
  qed
  show ?thesis
    apply (rule_tac x=min  $\delta \ \eta$  in exI)
    using  $\langle 0 < \delta \rangle \langle 0 < \eta \rangle$ 
      by (auto simp: dist_norm contour_integral_diff [OF cont_dw cont_dw,
  symmetric]  $\langle w \in U \rangle$  intro: le_ee)
    qed
  show ?thesis
  proof (cases a=b)
    case False
      show ?thesis
        by (rule continuous_onI) (use False in  $\langle$ auto intro:  $\ast$  $\rangle$ )
    qed auto
  qed

```

This version has *polynomial_function* γ as an additional assumption.

lemma *Cauchy_integral_formula_global_weak*:

```

  assumes open U and holf: f holomorphic_on U
    and z: z  $\in$  U and  $\gamma$ : polynomial_function  $\gamma$ 
    and pasz: path_image  $\gamma \subseteq U - \{z\}$  and loop: pathfinish  $\gamma =$  pathstart  $\gamma$ 
    and zero:  $\bigwedge w. w \notin U \implies$  winding_number  $\gamma \ w = 0$ 
  shows (( $\lambda w. f \ w / (w - z)$ ) has_contour_integral ( $2 * \pi * i *$  winding_number
 $\gamma \ z * f \ z$ ))  $\gamma$ 
  proof -
    obtain  $\gamma'$  where pf $\gamma'$ : polynomial_function  $\gamma'$  and  $\gamma'$ :  $\bigwedge x. (\gamma$  has_vector_derivative
 $(\gamma' \ x))$  (at x)
      using has_vector_derivative_polynomial_function [OF  $\gamma$ ] by blast
    then have bounded(path_image  $\gamma')$ 
      by (simp add: path_image_def compact_imp_bounded compact_continuous_image
  continuous_on_polynomial_function)
    then obtain B where B $>0$  and B:  $\bigwedge x. x \in$  path_image  $\gamma' \implies$  norm x  $\leq$  B
      using bounded_pos by force
    define d where [abs_def]: d z w = (if w = z then deriv f z else (f w - f z)/(w
  - z)) for z w
    define v where v = {w. w  $\notin$  path_image  $\gamma \wedge$  winding_number  $\gamma \ w = 0$ }
    have path  $\gamma$  valid_path  $\gamma$  using  $\gamma$ 
      by (auto simp: path_polynomial_function valid_path_polynomial_function)
    then have ov: open v
      by (simp add: v_def open_winding_number_levelsets loop)
    have uv_Un: U  $\cup$  v = UNIV
      using pasz zero by (auto simp: v_def)
    have conf: continuous_on U f
      by (metis holf holomorphic_on_imp_continuous_on)

```



```

have hol_d: (d y) holomorphic_on U if y ∈ U for y
proof -
  have *: (λc. if c = y then deriv f y else (f c - f y) / (c - y)) holomorphic_on
  U
  by (simp add: holf pole_lemma_open ‹open U›)
  then have isCont (λx. if x = y then deriv f y else (f x - f y) / (x - y)) y
  using at_within_open field_differentiable_imp_continuous_at holomor-
  phic_on_def that ‹open U› by fastforce
  then have continuous_on U (d y)
  using * d_def holomorphic_on_imp_continuous_on by auto
  moreover have d y holomorphic_on U - {y}
  proof -
    have (λw. if w = y then deriv f y else (f w - f y) / (w - y)) field_differentiable
    at w
    if w ∈ U - {y} for w
    proof (rule field_differentiable_transform_within)
      show (λw. (f w - f y) / (w - y)) field_differentiable at w
      using that ‹open U› holf
      by (auto intro!: holomorphic_on_imp_differentiable_at derivative_intros)
      show dist w y > 0
      using that by auto
    qed (auto simp: dist_commute)
    then show ?thesis
    unfolding field_differentiable_def by (simp add: d_def holomorphic_on_open
    ‹open U› open_delete)
  qed
  ultimately show ?thesis
  by (rule no_isolated_singularity) (auto simp: ‹open U›)
qed
have cint_fxγ: (λx. (f x - f y) / (x - y)) contour_integrable_on γ if y ∉
path_image γ for y
proof (rule contour_integrable_holomorphic_simple [where S = U - {y}])
  show (λx. (f x - f y) / (x - y)) holomorphic_on U - {y}
  by (force intro: holomorphic_intros holomorphic_on_subset [OF holf])
  show path_image γ ⊆ U - {y}
  using pasz that by blast
qed (auto simp: ‹open U› open_delete ‹valid_path γ›)
define h where
  h z = (if z ∈ U then contour_integral γ (d z) else contour_integral γ (λw. f
  w/(w - z))) for z
have U: ((d z) has_contour_integral h z) γ if z ∈ U for z
proof -
  have d z holomorphic_on U
  by (simp add: hol_d that)
  with that show ?thesis
  by (metis Diff_subset ‹valid_path γ› ‹open U› contour_integrable_holomorphic_simple
  h_def has_contour_integral_integral pasz subset_trans)
qed
have V: ((λw. f w / (w - z)) has_contour_integral h z) γ if z: z ∈ v for z

```

```

proof -
  have 0: 0 = (f z) * 2 * of_real (2 * pi) * i * winding_number  $\gamma$  z
    using v_def z by auto
  then have (( $\lambda x$ . 1 / (x - z)) has_contour_integral 0)  $\gamma$ 
    using z v_def has_contour_integral_winding_number [OF ‹valid_path  $\gamma$ ›]
by fastforce
  then have (( $\lambda x$ . f z * (1 / (x - z))) has_contour_integral 0)  $\gamma$ 
    using has_contour_integral_lmul by fastforce
  then have (( $\lambda x$ . f z / (x - z)) has_contour_integral 0)  $\gamma$ 
    by (simp add: field_split_simps)
  moreover have (( $\lambda x$ . (f x - f z) / (x - z)) has_contour_integral con-
tour_integral  $\gamma$  (d z))  $\gamma$ 
    using z
    apply (simp add: v_def)
  apply (metis (no_types, lifting) contour_integrable_eq d_def has_contour_integral_eq
has_contour_integral_integral cint_fxy)
  done
  ultimately have *: (( $\lambda x$ . f z / (x - z) + (f x - f z) / (x - z)) has_contour_integral
(0 + contour_integral  $\gamma$  (d z)))  $\gamma$ 
    by (rule has_contour_integral_add)
  have (( $\lambda w$ . f w / (w - z)) has_contour_integral contour_integral  $\gamma$  (d z))  $\gamma$ 
    if z  $\in$  U
    using * by (auto simp: divide_simps has_contour_integral_eq)
  moreover have (( $\lambda w$ . f w / (w - z)) has_contour_integral contour_integral
 $\gamma$  ( $\lambda w$ . f w / (w - z)))  $\gamma$ 
    if z  $\notin$  U
  proof (rule has_contour_integral_integral [OF contour_integrable_holomorphic_simple
[where S=U]])
    show ( $\lambda w$ . f w / (w - z)) holomorphic_on U
      by (rule holomorphic_intros assms | use that in force)+
    qed (use ‹open U› pasz ‹valid_path  $\gamma$ › in auto)
    ultimately show ?thesis
      using z by (simp add: h_def)
  qed
  have znot: z  $\notin$  path_image  $\gamma$ 
    using pasz by blast
  obtain d0 where d0 > 0 and d0:  $\bigwedge x y$ . x  $\in$  path_image  $\gamma$   $\implies$  y  $\in$  - U  $\implies$  d0
 $\leq$  dist x y
    using separate_compact_closed [of path_image  $\gamma$  -U] pasz ‹open U› ‹path  $\gamma$ ›
compact_path_image
    by blast
  obtain dd where 0 < dd and dd: {y + k | y k. y  $\in$  path_image  $\gamma$   $\wedge$  k  $\in$  ball 0
dd}  $\subseteq$  U
  proof
    show 0 < d0 / 2 using ‹0 < d0› by auto
  qed (use ‹0 < d0› d0 in ‹force simp: dist_norm›)
  define T where T  $\equiv$  {y + k | y k. y  $\in$  path_image  $\gamma$   $\wedge$  k  $\in$  cball 0 (dd / 2)}
  have  $\bigwedge x x'$ .  $\llbracket x \in$  path_image  $\gamma$ ; dist x x' * 2 < dd  $\rrbracket \implies \exists y k$ . x' = y + k  $\wedge$  y
 $\in$  path_image  $\gamma$   $\wedge$  dist 0 k * 2  $\leq$  dd

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    apply (rule_tac x=x in exI)
    apply (rule_tac x=x'-x in exI)
    apply (force simp: dist_norm)
  done
then have subt: path_image  $\gamma \subseteq$  interior  $T$ 
  using <0 < dd>
  apply (clarsimp simp add: mem_interior T_def)
  apply (rule_tac x=dd/2 in exI, auto)
  done
have compact T
  unfolding T_def
  using <valid_path  $\gamma$ > compact_cball compact_sums compact_valid_path_image
by blast
have T:  $T \subseteq U$ 
  unfolding T_def using <0 < dd> dd by fastforce
obtain L where  $L > 0$ 
  and L:  $\bigwedge B. \llbracket f \text{ holomorphic\_on interior } T; \bigwedge z. z \in \text{interior } T \implies cmod$ 
 $(f z) \leq B \rrbracket \implies$ 
       $cmod (\text{contour\_integral } \gamma f) \leq L * B$ 
  using contour_integral_bound_exists [OF open_interior <valid_path  $\gamma$ > subt]
  by blast
have bounded(f ' T)
  by (meson <compact T> compact_continuous_image compact_imp_bounded
conf continuous_on_subset T)
then obtain D where  $D > 0$  and D:  $\bigwedge x. x \in T \implies \text{norm } (f x) \leq D$ 
  by (auto simp: bounded_pos)
obtain C where  $C > 0$  and C:  $\bigwedge x. x \in T \implies \text{norm } x \leq C$ 
  using <compact T> bounded_pos compact_imp_bounded by force
have dist (h y)  $0 \leq e$  if  $0 < e$  and le:  $D * L / e + C \leq cmod y$  for  $e y$ 
proof -
  have  $D * L / e > 0$  using <D>0> <L>0> <e>0> by simp
  with le have ybig:  $\text{norm } y > C$  by force
  with C have  $y \notin T$  by force
  then have ynot:  $y \notin \text{path\_image } \gamma$ 
    using subt interior_subset by blast
  have [simp]:  $\text{winding\_number } \gamma y = 0$ 
  proof (rule winding_number_zero_outside)
    show  $\text{path\_image } \gamma \subseteq \text{cball } 0 C$ 
    by (meson C interior_subset mem_cball_0 subset_eq subt)
  qed (use ybig loop <path  $\gamma$ > in auto)
  have [simp]:  $h y = \text{contour\_integral } \gamma (\lambda w. f w / (w - y))$ 
    by (rule contour_integral_unique [symmetric]) (simp add: v_def ynot V)
  have holint:  $(\lambda w. f w / (w - y)) \text{ holomorphic\_on interior } T$ 
  proof (intro holomorphic_intros)
    show  $f \text{ holomorphic\_on interior } T$ 
    using holf holomorphic_on_subset interior_subset T by blast
  qed (use <y  $\notin T$ > interior_subset in auto)
  have leD:  $cmod (f z / (z - y)) \leq D * (e / L / D)$  if  $z: z \in \text{interior } T$  for  $z$ 
  proof -

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have  $D * L / e + cmod\ z \leq cmod\ y$ 
  using  $le\ C\ [of\ z]\ z$  using interior_subset by force
then have  $DL2: D * L / e \leq cmod\ (z - y)$ 
  using norm_triangle_ineq2 [of  $y\ z$ ] by (simp add: norm_minus_commute)
have  $cmod\ (f\ z / (z - y)) = cmod\ (f\ z) * inverse\ (cmod\ (z - y))$ 
by (simp add: norm_mult norm_inverse Fields.field_class.field_divide_inverse)
also have  $\dots \leq D * (e / L / D)$ 
proof (rule mult_mono)
  show  $cmod\ (f\ z) \leq D$ 
    using  $D\ interior\_subset\ z$  by blast
  show  $inverse\ (cmod\ (z - y)) \leq e / L / D\ D \geq 0$ 
    using  $\langle L > 0 \rangle\ \langle e > 0 \rangle\ \langle D > 0 \rangle\ DL2$  by (auto simp: norm_divide field_split_simps)
qed auto
finally show ?thesis .
qed
have  $dist\ (h\ y)\ 0 = cmod\ (contour\_integral\ \gamma\ (\lambda w. f\ w / (w - y)))$ 
  by (simp add: dist_norm)
also have  $\dots \leq L * (D * (e / L / D))$ 
  by (rule  $L\ [OF\ holint\ leD]$ )
also have  $\dots = e$ 
  using  $\langle L > 0 \rangle\ \langle 0 < D \rangle$  by auto
finally show ?thesis .
qed
then have  $(h \longrightarrow 0)$  at_infinity
  by (meson  $Lim\_at\_infinityI$ )
moreover have  $h$  holomorphic_on UNIV
proof -
  have  $con\_ff: continuous\ (at\ (x,z))\ (\lambda(x,y). (f\ y - f\ x) / (y - x))$ 
    if  $x \in U\ z \in U\ x \neq z$  for  $x\ z$ 
  using that conf
  apply (simp add: split_def continuous_on_eq_continuous_at  $\langle open\ U \rangle$ )
  apply (simp | rule continuous_intros continuous_within_compose2 [where
 $g=f$ ])+
  done
  have  $con\_fstsnd: continuous\_on\ UNIV\ (\lambda x. (fst\ x - snd\ x) :: complex)$ 
    by (rule continuous_intros)+
  have  $open\_uu\_Id: open\ (U \times U - Id)$ 
proof (rule open_Diff)
  show  $open\ (U \times U)$ 
    by (simp add: open_Times  $\langle open\ U \rangle$ )
  show  $closed\ (Id :: complex\ rel)$ 
    using continuous_closed_preimage_constant [OF  $con\_fstsnd\ closed\_UNIV,$ 
of 0]
    by (auto simp: Id_fstsnd_eq algebra_simps)
qed
have  $con\_derf: continuous\ (at\ z)\ (deriv\ f)$  if  $z \in U$  for  $z$ 
  by (meson analytic_at analytic_at_imp_isCont assms(1) holf_holomor-
phic_deriv that)
have  $tendsto\_f': ((\lambda(x,y). if\ y = x\ then\ deriv\ f\ (x)$ 

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      else (f (y) - f (x)) / (y - x) ⟶ deriv f x
    (at (x, x) within U × U) if x ∈ U for x
  proof (rule Lim_withinI)
    fix e::real assume 0 < e
    obtain k1 where k1>0 and k1:  $\bigwedge x'. \text{norm } (x' - x) \leq k1 \implies \text{norm } (\text{deriv } f x' - \text{deriv } f x) < e$ 
      using ⟨0 < e⟩ continuous_within_E [OF con_derf [OF ⟨x ∈ U⟩]]
      by (metis UNIV_I dist_norm)
    obtain k2 where k2>0 and k2: ball x k2 ⊆ U
      by (blast intro: openE [OF ⟨open U⟩] ⟨x ∈ U⟩)
    have neg: norm ((f z' - f x') / (z' - x') - deriv f x) ≤ e
      if z' ≠ x' and less_k1: norm (x'-x, z'-x) < k1 and less_k2:
norm (x'-x, z'-x) < k2
      for x' z'
    proof -
      have cs_less: w ∈ closed_segment x' z' ⟹ cmod (w - x) ≤ norm (x'-x, z'-x) for w
        using segment_furthest_le [of w x' z' x]
        by (metis (no_types) dist_commute dist_norm normfst_le normsnd_le order_trans)
      have derf_le: w ∈ closed_segment x' z' ⟹ z' ≠ x' ⟹ cmod (deriv f w - deriv f x) ≤ e for w
        by (blast intro: cs_less less_k1 k1 [unfolded divide_const_simps dist_norm] less_imp_le le_less_trans)
      have f_has_der:  $\bigwedge x. x \in U \implies (f \text{ has\_field\_derivative } \text{deriv } f x) \text{ (at } x \text{ within } U)$ 
        by (metis DERIV_deriv_iff_field_differentiable at_within_open holf holomorphic_on_def ⟨open U⟩)
      have closed_segment x' z' ⊆ U
        by (rule order_trans [OF _ k2]) (simp add: cs_less le_less_trans [OF _ less_k2] dist_complex_def norm_minus_commute subset_iff)
      then have cint_derf: (deriv f has_contour_integral f z' - f x') (linepath x' z')
        using contour_integral_primitive [OF f_has_der valid_path_linepath]
        by simp
      then have *: (( $\lambda x. \text{deriv } f x / (z' - x')$ ) has_contour_integral (f z' - f x') / (z' - x')) (linepath x' z')
        by (rule has_contour_integral_div)
      have norm ((f z' - f x') / (z' - x') - deriv f x) ≤ e/norm(z' - x') * norm(z' - x')
        apply (rule has_contour_integral_bound_linepath [OF has_contour_integral_diff [OF *]])
        using has_contour_integral_div [where c = z' - x', OF has_contour_integral_const_linepath [of deriv f x z' x']]
        ⟨e > 0⟩ ⟨z' ≠ x'⟩
      apply (auto simp: norm_divide divide_simps derf_le)
      done
    also have ... ≤ e using ⟨0 < e⟩ by simp
    finally show ?thesis .
  
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qed
show  $\exists d > 0. \forall xa \in U \times U.$ 
       $0 < \text{dist } xa (x, x) \wedge \text{dist } xa (x, x) < d \longrightarrow$ 
       $\text{dist } (\text{case } xa \text{ of } (x, y) \Rightarrow \text{if } y = x \text{ then deriv } f x \text{ else } (f y - f x) / (y$ 
-  $x)) (\text{deriv } f x) \leq e$ 
      apply (rule_tac x=min k1 k2 in exI)
      using  $\langle k1 > 0 \rangle \langle k2 > 0 \rangle \langle e > 0 \rangle$ 
      by (force simp: dist_norm neq intro: dual_order.strict_trans2 k1 less_imp_le
norm_fst_le)
qed
have con_pa_f: continuous_on (path_image  $\gamma$ )  $f$ 
by (meson holf holomorphic_on_imp_continuous_on holomorphic_on_subset
interior_subset subt T)
have le_B:  $\bigwedge T. T \in \{0..1\} \Longrightarrow \text{cmod } (\text{vector\_derivative } \gamma (\text{at } T)) \leq B$ 
using  $\gamma' B$  by (simp add: path_image_def vector_derivative_at rev_image_eqI)
have f_has_cint:  $\bigwedge w. w \in v - \text{path\_image } \gamma \Longrightarrow ((\lambda u. f u / (u - w)) ^ 1)$ 
has_contour_integral h  $w$ )  $\gamma$ 
by (simp add: V)
have cond_uu: continuous_on  $(U \times U) (\lambda(x,y). d x y)$ 
apply (simp add: continuous_on_eq_continuous_within d_def continuous_
ous_within tendsto_f')
apply (simp add: tendsto_within_open_NO_MATCH open_Times  $\langle \text{open}$ 
 $U \rangle$ , clarify)
apply (rule Lim_transform_within_open [OF _ open_uu_Id, where  $f =$ 
 $(\lambda(x,y). (f y - f x) / (y - x))$ ])
using con_ff
apply (auto simp: continuous_within)
done
have holf_dw:  $(\lambda z. d z w)$  holomorphic_on  $U$  if  $w \in U$  for  $w$ 
proof -
have continuous_on  $U ((\lambda(x,y). d x y) \circ (\lambda z. (w,z)))$ 
by (rule continuous_on_compose continuous_intros continuous_on_subset
[OF cond_uu] | force intro: that)+
then have *: continuous_on  $U (\lambda z. \text{if } w = z \text{ then deriv } f z \text{ else } (f w - f z)$ 
 $/ (w - z))$ 
by (rule rev_iffD1 [OF _ continuous_on_cong [OF refl]]) (simp add: d_def
field_simps)
have **:  $(\lambda z. \text{if } w = z \text{ then deriv } f z \text{ else } (f w - f z) / (w - z))$  field_differentiable
at  $x$ 
if  $x \in U$   $x \neq w$  for  $x$ 
proof (rule_tac  $f = \lambda x. (f w - f x) / (w - x)$  and  $d = \text{dist } x w$  in
field_differentiable_transform_within)
show  $(\lambda x. (f w - f x) / (w - x))$  field_differentiable at  $x$ 
using that  $\langle \text{open } U \rangle$ 
by (intro derivative_intros holomorphic_on_imp_differentiable_at [OF
holf]; force)
qed (use that  $\langle \text{open } U \rangle$  in  $\langle \text{auto simp: dist\_commute} \rangle$ )
show ?thesis
unfolding d_def

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proof (rule no_isolated_singularity [OF * _ ‹open U›])
  show ( $\lambda z$ . if  $w = z$  then  $\text{deriv } f z$  else  $(f w - f z) / (w - z)$ ) holomorphic_on
  U - {w}
  by (auto simp: field_differentiable_def [symmetric] holomorphic_on_open
  open_Diff ‹open U› **)
qed auto
qed
{ fix a b
  assume abu: closed_segment a b  $\subseteq$  U
  have cont_cint_d: continuous_on U ( $\lambda w$ . contour_integral (linepath a b)
  ( $\lambda z$ . d z w))
  proof (rule contour_integral_continuous_on_linepath_2D [OF ‹open U› _
  _ abu])
    show  $\bigwedge w$ .  $w \in U \implies (\lambda z$ . d z w) contour_integrable_on (linepath a b)
    by (metis abu hol_dw continuous_on_subset contour_integrable_continuous_linepath
  holomorphic_on_imp_continuous_on)
    show continuous_on (U  $\times$  U) ( $\lambda(x, y)$ . d y x)
    by (auto intro: continuous_on_swap_args cond uu)
  qed
  have cont_cint_d $\gamma$ : continuous_on {0..1} (( $\lambda w$ . contour_integral (linepath
  a b) ( $\lambda z$ . d z w))  $\circ$   $\gamma$ )
  proof (rule continuous_on_compose)
    show continuous_on {0..1}  $\gamma$ 
    using ‹path  $\gamma$ › path_def by blast
    show continuous_on ( $\gamma$  ‘ {0..1}) ( $\lambda w$ . contour_integral (linepath a b) ( $\lambda z$ .
  d z w))
    using pasz unfolding path_image_def
    by (auto intro!: continuous_on_subset [OF cont_cint_d])
  qed
  have continuous_on {0..1} ( $\lambda x$ . vector_derivative  $\gamma$  (at x))
  using pf $\gamma$  by (simp add: continuous_on_polynomial_function vector_derivative_at
  [OF  $\gamma$ ])
  then have cint_cint: ( $\lambda w$ . contour_integral (linepath a b) ( $\lambda z$ . d z w))
  contour_integrable_on  $\gamma$ 
  apply (simp add: contour_integrable_on)
  apply (rule integrable_continuous_real)
  by (rule continuous_on_mult [OF cont_cint_d $\gamma$  [unfolded o_def]])
  have contour_integral (linepath a b) h = contour_integral (linepath a b) ( $\lambda z$ .
  contour_integral  $\gamma$  (d z))
  using abu by (force simp: h_def intro: contour_integral_eq)
  also have ... = contour_integral  $\gamma$  ( $\lambda w$ . contour_integral (linepath a b) ( $\lambda z$ .
  d z w))
  proof (rule contour_integral_swap)
    show continuous_on (path_image (linepath a b)  $\times$  path_image  $\gamma$ ) ( $\lambda(y1,
  y2)$ . d y1 y2)
    using abu pasz by (auto intro: continuous_on_subset [OF cond uu])
    show continuous_on {0..1} ( $\lambda t$ . vector_derivative (linepath a b) (at t))
    by (auto intro!: continuous_intros)
    show continuous_on {0..1} ( $\lambda t$ . vector_derivative  $\gamma$  (at t))

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    by (metis  $\gamma'$  continuous_on_eq path_def path_polynomial_function pf  $\gamma'$ 
vector_derivative_at)
  qed (use <valid_path  $\gamma$ > in auto)
  finally have cint_h_eq:
    contour_integral (linepath a b) h =
      contour_integral  $\gamma$  ( $\lambda w. contour\_integral (linepath a b) (\lambda z. d z$ 
w)) .
  note cint_cint cint_h_eq
} note cint_h = this
have conth_u: continuous_on U h
proof (simp add: continuous_on_sequentially, clarify)
  fix a x
  assume x:  $x \in U$  and au:  $\forall n. a n \in U$  and ax:  $a \longrightarrow x$ 
  then have A1:  $\forall_F n$  in sequentially.  $d (a n)$  contour_integrable_on  $\gamma$ 
    by (meson U contour_integrable_on_def eventuallyI)
  obtain dd where dd > 0 and dd:  $cball\ x\ dd \subseteq U$  using open_contains_cball
<open U> x by force
  have A2: uniform_limit (path_image  $\gamma$ ) ( $\lambda n. d (a n)$ ) ( $d x$ ) sequentially
    unfolding uniform_limit_iff dist_norm
  proof clarify
    fix ee::real
    assume 0 < ee
    show  $\forall_F n$  in sequentially.  $\forall \xi \in path\_image\ \gamma. cmod (d (a n) \xi - d x \xi) <$ 
ee
  proof -
    let ?ddpa =  $\{(w,z) \mid w z. w \in cball\ x\ dd \wedge z \in path\_image\ \gamma\}$ 
    have uniformly_continuous_on ?ddpa ( $\lambda(x,y). d x y$ )
  proof (rule compact_uniformly_continuous [OF continuous_on_subset [OF
cond_uu]])
    show compact  $\{(w, z) \mid w z. w \in cball\ x\ dd \wedge z \in path\_image\ \gamma\}$ 
      using <valid_path  $\gamma$ >
    by (auto simp: compact_Times compact_valid_path_image simp del:
mem_cball)
    qed (use dd pasz in auto)
    then obtain kk where kk > 0
      and kk:  $\bigwedge x x'. \llbracket x \in ?ddpa; x' \in ?ddpa; dist\ x'\ x < kk \rrbracket \implies$ 
 $dist ((\lambda(x,y). d x y) x') ((\lambda(x,y). d x y) x) < ee$ 
      by (rule uniformly_continuous_onE [where e = ee]) (use <0 < ee> in
auto)
    have kk:  $\llbracket norm (w - x) \leq dd; z \in path\_image\ \gamma; norm ((w, z) - (x, z))$ 
< kk  $\rrbracket \implies norm (d w z - d x z) < ee$ 
      for w z
      using <dd > 0> kk [of (x,z) (w,z)] by (force simp: norm_minus_commute
dist_norm)
    obtain no where  $\forall n \geq no. dist (a n) x < min\ dd\ kk$ 
      using ax unfolding lim_sequentially
      by (meson <0 < dd> <0 < kk> min_less_iff_conj)
    then show ?thesis
      using <dd > 0> <kk > 0> by (fastforce simp: eventually_sequentially kk

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dist_norm)
  qed
  qed
  have (λn. contour_integral γ (d (a n))) ⟶ contour_integral γ (d x)
    by (rule contour_integral_uniform_limit [OF A1 A2 le_B]) (auto simp:
⟨valid_path γ⟩)
  then have tendsto_hx: (λn. contour_integral γ (d (a n))) ⟶ h x
    by (simp add: h_def x)
  then show (h ∘ a) ⟶ h x
    by (simp add: h_def x au_o_def)
  qed
  show ?thesis
  proof (simp add: holomorphic_on_open field_differentiable_def [symmetric],
clarify)
    fix z0
    consider z0 ∈ v | z0 ∈ U using uv_Un by blast
    then show h field_differentiable at z0
    proof cases
      assume z0 ∈ v then show ?thesis
        using Cauchy_next_derivative [OF con_pa_f le_B f_has_cint _ ov] V
f_has_cint ⟨valid_path γ⟩
        by (auto simp: field_differentiable_def v_def)
      next
        assume z0 ∈ U then
          obtain e where e>0 and e: ball z0 e ⊆ U by (blast intro: openE [OF
⟨open U⟩])
          have *: contour_integral (linepath a b) h + contour_integral (linepath b c)
h + contour_integral (linepath c a) h = 0
            if abc_subset: convex_hull {a, b, c} ⊆ ball z0 e for a b c
          proof -
            have *: ∧x1 x2 z. z ∈ U ⟹ closed_segment x1 x2 ⊆ U ⟹ (λw. d w z)
contour_integrable_on linepath x1 x2
              using hol_dw holomorphic_on_imp_continuous_on ⟨open U⟩
              by (auto intro!: contour_integrable_holomorphic_simple)
            have abc: closed_segment a b ⊆ U closed_segment b c ⊆ U closed_segment
c a ⊆ U
              using that e segments_subset_convex_hull by fastforce+
            have eq0: ∧w. w ∈ U ⟹ contour_integral (linepath a b +++ linepath b
c +++ linepath c a) (λz. d z w) = 0
              proof (rule contour_integral_unique [OF Cauchy_theorem_triangle])
                show ∧w. w ∈ U ⟹ (λz. d z w) holomorphic_on convex_hull {a, b, c}
                  using e abc_subset by (auto intro: holomorphic_on_subset [OF
hol_dw])
              qed
            qed
          have contour_integral γ
            (λx. contour_integral (linepath a b) (λz. d z x) +
              contour_integral (linepath b c) (λz. d z x) +
              contour_integral (linepath c a) (λz. d z x)) = 0
            apply (rule contour_integral_eq_0)

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    using abc pasz U
    apply (subst contour_integral_join [symmetric], auto intro: eq0 *)+
    done
  then show ?thesis
    by (simp add: cint_h abc contour_integrable_add contour_integral_add
[symmetric] add_ac)
  qed
  show ?thesis
    using e <e > 0>
    by (auto intro!: holomorphic_on_imp_differentiable_at [OF __ open_ball]
analytic_imp_holomorphic
      Morera_triangle continuous_on_subset [OF conthu] *)
  qed
  qed
  qed
  ultimately have [simp]:  $h z = 0$  for  $z$ 
    by (meson Liouville_weak)
  have (( $\lambda w. 1 / (w - z)$ ) has_contour_integral complex_of_real (2 * pi) * i *
winding_number  $\gamma z$ )  $\gamma$ 
    by (rule has_contour_integral_winding_number [OF <valid_path  $\gamma$ > znot])
  then have (( $\lambda w. f z * (1 / (w - z))$ ) has_contour_integral complex_of_real (2
* pi) * i * winding_number  $\gamma z * f z$ )  $\gamma$ 
    by (metis mult.commute has_contour_integral_lmul)
  then have 1: (( $\lambda w. f z / (w - z)$ ) has_contour_integral complex_of_real (2 *
pi) * i * winding_number  $\gamma z * f z$ )  $\gamma$ 
    by (simp add: field_split_simps)
  moreover have 2: (( $\lambda w. (f w - f z) / (w - z)$ ) has_contour_integral 0)  $\gamma$ 
    using U [OF z] pasz d_def by (force elim: has_contour_integral_eq [where
 $g = \lambda w. (f w - f z)/(w - z)$ ])
  show ?thesis
    using has_contour_integral_add [OF 1 2] by (simp add: diff_divide_distrib)
  qed

```

theorem Cauchy_integral_formula_global:

```

  assumes S: open S and holf: f holomorphic_on S
  and z:  $z \in S$  and vpg: valid_path  $\gamma$ 
  and pasz: path_image  $\gamma \subseteq S - \{z\}$  and loop: pathfinish  $\gamma = \text{pathstart } \gamma$ 
  and zero:  $\bigwedge w. w \notin S \implies \text{winding\_number } \gamma w = 0$ 
  shows (( $\lambda w. f w / (w - z)$ ) has_contour_integral (2*pi * i * winding_number
 $\gamma z * f z$ ))  $\gamma$ 
  proof -
    have path  $\gamma$  using vpg by (blast intro: valid_path_imp_path)
    have hols: ( $\lambda w. f w / (w - z)$ ) holomorphic_on  $S - \{z\}$  ( $\lambda w. 1 / (w - z)$ )
holomorphic_on  $S - \{z\}$ 
      by (rule holomorphic_intros holomorphic_on_subset [OF holf] | force)+
    then have cint_fw: ( $\lambda w. f w / (w - z)$ ) contour_integrable_on  $\gamma$ 
      by (meson contour_integrable_holomorphic_simple holomorphic_on_imp_continuous_on
open_delete S vpg pasz)
    obtain d where  $d > 0$ 

```

```

and  $d$ :  $\bigwedge g h. \llbracket \text{valid\_path } g; \text{ valid\_path } h; \forall t \in \{0..1\}. \text{cmod } (g \ t - \gamma \ t) < d$ 
 $\wedge \text{cmod } (h \ t - \gamma \ t) < d;$ 
       $\text{pathstart } h = \text{pathstart } g \wedge \text{pathfinish } h = \text{pathfinish } g \rrbracket$ 
       $\implies \text{path\_image } h \subseteq S - \{z\} \wedge (\forall f. f \text{ holomorphic\_on } S - \{z\}$ 
 $\longrightarrow \text{contour\_integral } h \ f = \text{contour\_integral } g \ f)$ 
      using contour_integral_nearby_ends [OF  $\langle \text{path } \gamma \rangle \text{ pasz}$ ]  $S$  by (simp add:
open_Diff) metis
      obtain  $p$  where polyp: polynomial_function  $p$ 
      and  $ps$ :  $\text{pathstart } p = \text{pathstart } \gamma$  and  $pf$ :  $\text{pathfinish } p = \text{pathfinish } \gamma$ 
and  $led$ :  $\forall t \in \{0..1\}. \text{cmod } (p \ t - \gamma \ t) < d$ 
      using path_approx_polynomial_function [OF  $\langle \text{path } \gamma \rangle \langle d > 0 \rangle$ ] by metis
      then have  $ploop$ :  $\text{pathfinish } p = \text{pathstart } p$  using loop by auto
      have  $vpp$ :  $\text{valid\_path } p$  using polyp valid_path_polynomial_function by blast
      have [simp]:  $z \notin \text{path\_image } \gamma$  using pasz by blast
      have  $paps$ :  $\text{path\_image } p \subseteq S - \{z\}$  and  $cint\_eq$ :  $(\bigwedge f. f \text{ holomorphic\_on } S -$ 
 $\{z\} \implies \text{contour\_integral } p \ f = \text{contour\_integral } \gamma \ f)$ 
      using pf ps led d [OF  $\langle vpg \ vpp \rangle \langle d > 0 \rangle$ ] by auto
      have  $wn\_eq$ :  $\text{winding\_number } p \ z = \text{winding\_number } \gamma \ z$ 
      using vpp paps
      by (simp add: subset_Diff_insert vpg valid_path_polynomial_function wind-
ing_number_valid_path_cint_eq hols)
      have  $\text{winding\_number } p \ w = \text{winding\_number } \gamma \ w$  if  $w \notin S$  for  $w$ 
      proof  $-$ 
      have  $hol$ :  $(\lambda v. 1 / (v - w)) \text{ holomorphic\_on } S - \{z\}$ 
      using that by (force intro: holomorphic_intros holomorphic_on_subset [OF
holf])
      have  $w \notin \text{path\_image } p \ w \notin \text{path\_image } \gamma$  using paps pasz that by auto
      then show ?thesis
      using vpp vpg by (simp add: subset_Diff_insert valid_path_polynomial_function
winding_number_valid_path_cint_eq [OF hol])
      qed
      then have  $wn0$ :  $\bigwedge w. w \notin S \implies \text{winding\_number } p \ w = 0$ 
      by (simp add: zero)
      show ?thesis
      using Cauchy_integral_formula_global_weak [OF  $S$  holf  $z$  polyp paps ploop
wn0] hols
      by (metis wn_eq cint_eq has_contour_integral_eqpath cint_fw cint_eq)
      qed

theorem Cauchy_theorem_global:
  assumes  $S$ : open  $S$  and  $holf$ :  $f \text{ holomorphic\_on } S$ 
    and  $vpg$ :  $\text{valid\_path } \gamma$  and  $loop$ :  $\text{pathfinish } \gamma = \text{pathstart } \gamma$ 
    and  $pas$ :  $\text{path\_image } \gamma \subseteq S$ 
    and  $zero$ :  $\bigwedge w. w \notin S \implies \text{winding\_number } \gamma \ w = 0$ 
  shows  $(f \text{ has\_contour\_integral } 0) \ \gamma$ 
proof  $-$ 
  obtain  $z$  where  $z \in S$  and  $znot$ :  $z \notin \text{path\_image } \gamma$ 
  proof  $-$ 
  have  $\text{path\_image } \gamma \neq S$ 

```

by (*metis compact_valid_path_image vpg compact_open path_image_nonempty*
S)
with pas show *?thesis* **by** (*blast intro: that*)
qed
then have *pasz*: $\text{path_image } \gamma \subseteq S - \{z\}$ **using** *pas* **by** *blast*
have *hol*: $(\lambda w. (w - z) * f w)$ *holomorphic_on S*
by (*rule holomorphic_intros holf*)
show *?thesis*
using *Cauchy_integral_formula_global* [*OF S hol* $\langle z \in S \rangle$ *vpg pasz loop zero*]
by (*auto simp: znot elim!: has_contour_integral_eq*)
qed

corollary *Cauchy_theorem_global_outside*:

assumes *open S f holomorphic_on S valid_path* γ *pathfinish* $\gamma = \text{pathstart } \gamma$
path_image $\gamma \subseteq S$
 $\bigwedge w. w \notin S \implies w \in \text{outside}(\text{path_image } \gamma)$
shows (*f has_contour_integral 0*) γ
by (*metis Cauchy_theorem_global* *assms winding_number_zero_in_outside valid_path_imp_path*)

lemma *simply_connected_imp_winding_number_zero*:

assumes *simply_connected S path g*
 $\text{path_image } g \subseteq S$ *pathfinish* $g = \text{pathstart } g$ $z \notin S$
shows *winding_number g z = 0*
proof –
have *hom*: *homotopic_loops S g* (*linepath* (*pathstart g*) (*pathstart g*))
by (*meson* *assms homotopic_paths_imp_homotopic_loops pathfinish_linepath*
simply_connected_eq_contractible_path)
then have *homotopic_paths* ($-\{z\}$) *g* (*linepath* (*pathstart g*) (*pathstart g*))
by (*meson* $\langle z \notin S \rangle$ *homotopic_loops_imp_homotopic_paths_null* *homotopic_paths_subset*
subset_Comp1_singleton)
then have *winding_number g z = winding_number*(*linepath* (*pathstart g*) (*pathstart*
g)) *z*
by (*rule winding_number_homotopic_paths*)
also have $\dots = 0$
using *assms* **by** (*force intro: winding_number_trivial*)
finally show *?thesis* .
qed

lemma *Cauchy_theorem_simply_connected*:

assumes *open S simply_connected S f holomorphic_on S valid_path g*
 $\text{path_image } g \subseteq S$ *pathfinish* $g = \text{pathstart } g$
shows (*f has_contour_integral 0*) g
by (*meson* *assms Cauchy_theorem_global simply_connected_imp_winding_number_zero*
valid_path_imp_path)

proposition *holomorphic_logarithm_exists*:

assumes *A: convex A open A*
and *f: f holomorphic_on A* $\bigwedge x. x \in A \implies f x \neq 0$
and *z0: z0* $\in A$

```

obtains  $g$  where  $g$  holomorphic_on  $A$  and  $\bigwedge x. x \in A \implies \exp (g x) = f x$ 
proof –
  note  $f' = \text{holomorphic\_derivI}$  [OF  $f(1)$   $A(2)$ ]
  obtain  $g$  where  $g: \bigwedge x. x \in A \implies (g \text{ has\_field\_derivative } \text{deriv } f x / f x) \text{ (at } x)$ 
  proof (rule holomorphic_convex_primitive' [OF  $A$ ])
    show  $(\lambda x. \text{deriv } f x / f x)$  holomorphic_on  $A$ 
      by (intro holomorphic_intros  $f A$ )
    qed (auto simp: A at_within_open[of  $A$ ])
    define  $h$  where  $h = (\lambda x. -g z0 + \ln (f z0) + g x)$ 
    from  $g$  and  $A$  have  $g \text{ holo: } g \text{ holomorphic\_on } A$ 
      by (auto simp: holomorphic_on_def at_within_open[of  $A$ ] field_differentiable_def)
    hence  $h \text{ holo: } h \text{ holomorphic\_on } A$ 
      by (auto simp: h_def intro!: holomorphic_intros)
    note [simp] = at_within_open[OF  $\langle \text{open } A \rangle$ ]
    have  $\exists c. \forall x \in A. f x / \exp (h x) - 1 = c$ 
      using  $\langle \text{convex } A \rangle z0 f$ 
      by (force simp: h_def exp_diff field_simps intro!: has_field_derivative_zero_constant
derivative_eq_intros  $g f'$ )
    then obtain  $c$  where  $c: \bigwedge x. x \in A \implies f x / \exp (h x) - 1 = c$ 
      by blast
    from  $c$  [OF  $z0$ ] and  $z0$  and  $f$  have  $c = 0$ 
      by (simp add: h_def)
    with  $c$  have  $\bigwedge x. x \in A \implies \exp (h x) = f x$  by simp
    from that [OF  $h \text{ holo this}$ ] show ?thesis .
qed

```

4.14 Cauchy's inequality and more versions of Liouville

lemma Cauchy_higher_deriv_bound:

```

assumes  $holf: f \text{ holomorphic\_on } (\text{ball } z r)$ 
and  $contf: \text{continuous\_on } (\text{cball } z r) f$ 
and  $fin : \bigwedge w. w \in \text{ball } z r \implies f w \in \text{ball } y B0$ 
and  $0 < r$  and  $0 < n$ 
shows norm  $((\text{deriv } \sim n) f z) \leq (\text{fact } n) * B0 / r^n$ 
proof –
  have  $0 < B0$  using  $\langle 0 < r \rangle fin$  [of  $z$ ]
    by (metis ball_eq_empty ex_in_conv fin not_less)
  have  $le\_B0: \text{cmod } (f w - y) \leq B0$  if  $\text{cmod } (w - z) \leq r$  for  $w$ 
    proof (rule continuous_on_closure_norm_le [of  $\text{ball } z r \lambda w. f w - y$ ], use  $\langle 0 < r \rangle$ 
    in simp_all)
      show continuous_on  $(\text{cball } z r) (\lambda w. f w - y)$ 
        by (intro continuous_intros contf)
      show  $\text{dist } z w \leq r$ 
        by (simp add: dist_commute dist_norm that)
    qed (use fin in  $\langle \text{auto simp: dist_norm less_eq_real_def norm_minus_commute} \rangle$ )
  have  $(\text{deriv } \sim n) f z = (\text{deriv } \sim n) (\lambda w. f w) z - (\text{deriv } \sim n) (\lambda w. y) z$ 
    using  $\langle 0 < n \rangle$  by simp
  also have  $\dots = (\text{deriv } \sim n) (\lambda w. f w - y) z$ 
    by (rule higher_deriv_diff [OF  $holf, symmetric$ ]) (auto simp:  $\langle 0 < r \rangle$ )

```

```

finally have (deriv  $\hat{\sim}$  n) f z = (deriv  $\hat{\sim}$  n) ( $\lambda$ w. f w - y) z .
have contf': continuous_on (cball z r) ( $\lambda$ u. f u - y)
  by (rule contf continuous_intros)+
have holf': ( $\lambda$ u. (f u - y)) holomorphic_on (ball z r)
  by (simp add: holf holomorphic_on_diff)
define a where a = (2 * pi)/(fact n)
have 0 < a by (simp add: a_def)
have B0/r^(Suc n)*2 * pi * r = a*((fact n)*B0/r^n)
  using <0 < r> by (simp add: a_def field_split_simps)
have der_dif: (deriv  $\hat{\sim}$  n) ( $\lambda$ w. f w - y) z = (deriv  $\hat{\sim}$  n) f z
  using <0 < r> <0 < n>
  by (auto simp: higher_deriv_diff [OF holf holomorphic_on_const])
have norm ((2 * of_real pi * i)/(fact n) * (deriv  $\hat{\sim}$  n) ( $\lambda$ w. f w - y) z)
  ≤ (B0/r^(Suc n)) * (2 * pi * r)
  apply (rule has_contour_integral_bound_circlepath [of ( $\lambda$ u. (f u - y)/(u -
z)^(Suc n)) _ z])
  using Cauchy_has_contour_integral_higher_derivative_circlepath [OF contf'
holf]
  using <0 < B0> <0 < r>
  apply (auto simp: norm_divide norm_mult norm_power divide_simps le_B0)
  done
then show ?thesis
  using <0 < r>
  by (auto simp: norm_divide norm_mult norm_power field_simps der_dif
le_B0)
qed

```

lemma Cauchy_inequality:

```

assumes holf: f holomorphic_on (ball  $\xi$  r)
and contf: continuous_on (cball  $\xi$  r) f
and 0 < r
and nof:  $\bigwedge$ x. norm( $\xi$ -x) = r  $\implies$  norm(f x) ≤ B
shows norm ((deriv  $\hat{\sim}$  n) f  $\xi$ ) ≤ (fact n) * B / r^n
proof -
obtain x where norm ( $\xi$ -x) = r
  by (metis <0 < r> dist_norm order_less_imp_le vector_choose_dist)
then have 0 ≤ B
  by (metis nof norm_not_less_zero not_le order_trans)
have  $\xi \in$  ball  $\xi$  r
  using <0 < r> by simp
then have (( $\lambda$ u. f u / (u- $\xi$ ) ^ Suc n) has_contour_integral (2 * pi) * i / fact
n * (deriv  $\hat{\sim}$  n) f  $\xi$ )
  (circlepath  $\xi$  r)
  by (rule Cauchy_has_contour_integral_higher_derivative_circlepath [OF contf
holf])
have norm ((2 * pi * i)/(fact n) * (deriv  $\hat{\sim}$  n) f  $\xi$ ) ≤ (B / r^(Suc n)) * (2 *
pi * r)
proof (rule has_contour_integral_bound_circlepath)
  have  $\xi \in$  ball  $\xi$  r

```

```

    using ‹0 < r› by simp
  then show ((λu. f u / (u-ξ) ^ Suc n) has_contour_integral (2 * pi) * i /
    fact n * (deriv ^^ n) f ξ)
    (circlepath ξ r)
  by (rule Cauchy_has_contour_integral_higher_derivative_circlepath [OF
    contf holf])
  show ∧x. cmod (x-ξ) = r ⇒ cmod (f x / (x-ξ) ^ Suc n) ≤ B / r ^ Suc n
  using ‹0 ≤ B› ‹0 < r›
  by (simp add: norm_divide norm_power nof_frac_le norm_minus_commute
    del: power_Suc)
  qed (use ‹0 ≤ B› ‹0 < r› in auto)
  then show ?thesis using ‹0 < r›
  by (simp add: norm_divide norm_mult field_simps)
qed

```

lemma *Liouville_polynomial*:

```

  assumes holf: f holomorphic_on UNIV
  and nof: ∧z. A ≤ norm z ⇒ norm(f z) ≤ B * norm z ^ n
  shows f ξ = (∑ k≤n. (deriv ^^ k) f 0 / fact k * ξ ^ k)
proof (cases rule: le_less_linear [THEN disjE])
  assume B ≤ 0
  then have ∧z. A ≤ norm z ⇒ norm(f z) = 0
  by (metis nof less_le_trans zero_less_mult_iff neqE norm_not_less_zero
    norm_power not_le)
  then have f0: (f ⟶ 0) at_infinity
  using Lim_at_infinity by force
  then have [simp]: f = (λw. 0)
  using Liouville_weak [OF holf, of 0]
  by (simp add: eventually_at_infinity f0) meson
  show ?thesis by simp
next
  assume 0 < B
  have ((λk. (deriv ^^ k) f 0 / (fact k) * (ξ - 0) ^ k) sums f ξ)
  proof (rule holomorphic_power_series [where r = norm ξ + 1])
    show f holomorphic_on ball 0 (cmod ξ + 1) ξ ∈ ball 0 (cmod ξ + 1)
    using holf holomorphic_on_subset by auto
  qed
  then have sumsf: ((λk. (deriv ^^ k) f 0 / (fact k) * ξ ^ k) sums f ξ) by simp
  have (deriv ^^ k) f 0 / fact k * ξ ^ k = 0 if k>n for k
  proof (cases (deriv ^^ k) f 0 = 0)
    case True then show ?thesis by simp
  next
    case False
    define w where w = complex_of_real (fact k * B / cmod ((deriv ^^ k) f 0)
    + (|A| + 1))
    have 1 ≤ abs (fact k * B / cmod ((deriv ^^ k) f 0) + (|A| + 1))
    using ‹0 < B› by simp
    then have wge1: 1 ≤ norm w
    by (metis norm_of_real w_def)

```

```

then have  $w \neq 0$  by auto
have  $kB: 0 < \text{fact } k * B$ 
  using  $\langle 0 < B \rangle$  by simp
then have  $0 \leq \text{fact } k * B / \text{cmod } ((\text{deriv } \sim k) f 0)$ 
  by simp
then have  $wgeA: A \leq \text{cmod } w$ 
  by (simp only:  $w\_def$   $\text{norm\_of\_real}$ )
have  $\text{fact } k * B / \text{cmod } ((\text{deriv } \sim k) f 0) < \text{abs } (\text{fact } k * B / \text{cmod } ((\text{deriv } \sim k) f 0) + (|A| + 1))$ 
  using  $\langle 0 < B \rangle$  by simp
then have  $wge: \text{fact } k * B / \text{cmod } ((\text{deriv } \sim k) f 0) < \text{norm } w$ 
  by (metis  $\text{norm\_of\_real } w\_def$ )
then have  $\text{fact } k * B / \text{norm } w < \text{cmod } ((\text{deriv } \sim k) f 0)$ 
  using False by (simp add:  $\text{field\_split\_simps}$   $\text{mult.commute}$   $\text{split: if\_split\_asm}$ )
also have  $\dots \leq \text{fact } k * (B * \text{norm } w ^ n) / \text{norm } w ^ k$ 
proof (rule  $\text{Cauchy\_inequality}$ )
  show  $f$  holomorphic_on ball 0 (cmod w)
    using holf holomorphic_on_subset by force
  show continuous_on (cball 0 (cmod w)) f
    using holf holomorphic_on_imp_continuous_on holomorphic_on_subset
by blast
  show  $\bigwedge x. \text{cmod } (0 - x) = \text{cmod } w \implies \text{cmod } (f x) \leq B * \text{cmod } w ^ n$ 
    by (metis  $\text{nof } wgeA$   $\text{dist\_0\_norm}$   $\text{dist\_norm}$ )
qed (use  $\langle w \neq 0 \rangle$  in auto)
also have  $\dots = \text{fact } k * B / \text{cmod } w ^ (k-n)$ 
  using  $\langle k > n \rangle$  by (simp add:  $\text{divide\_simps}$   $\text{flip: power\_add}$ )
finally have  $\text{fact } k * B / \text{cmod } w < \text{fact } k * B / \text{cmod } w ^ (k - n)$  .
then have  $1 / \text{cmod } w < 1 / \text{cmod } w ^ (k - n)$ 
  by (metis  $kB$   $\text{divide\_inverse}$   $\text{inverse\_eq\_divide}$   $\text{mult\_less\_cancel\_left\_pos}$ )
then have  $\text{cmod } w ^ (k - n) < \text{cmod } w$ 
  by (smt (verit, best)  $\langle w \neq 0 \rangle$   $\text{frac\_le\_zero\_less\_norm\_iff}$ )
with  $\text{self\_le\_power}$  [OF  $wge1$ ] show ?thesis
  by (meson  $\text{diff\_is\_0\_eq}$   $\text{not\_gr0}$   $\text{not\_le}$  that)
qed
then have  $(\text{deriv } \sim (k + \text{Suc } n)) f 0 / \text{fact } (k + \text{Suc } n) * \xi ^ (k + \text{Suc } n) = 0$  for  $k$ 
  using  $\text{not\_less\_eq}$  by blast
then have  $(\lambda i. (\text{deriv } \sim (i + \text{Suc } n)) f 0 / \text{fact } (i + \text{Suc } n) * \xi ^ (i + \text{Suc } n))$ 
  sums 0
  by (rule  $\text{sums\_0}$ )
with  $\text{sums\_split\_initial\_segment}$  [OF  $\text{sumsf}$ , where  $n = \text{Suc } n$ ]
show ?thesis
  using  $\text{atLeast0AtMost lessThan\_Suc\_atMost}$   $\text{sums\_unique2}$  by fastforce
qed

```

Every bounded entire function is a constant function.

theorem *Liouville_theorem*:

assumes $\text{holf: } f$ holomorphic_on UNIV
 and $\text{bf: } f$ bounded (range f)

shows f constant on UNIV
 using Liouville_polynomial [OF holf, of 0 _ 0, simplified]
 by (metis bf bounded_iff constant_on_def rangeI)

A holomorphic function f has only isolated zeros unless f is 0.

lemma *power_0_nonzero*:
 fixes $a :: \text{nat} \Rightarrow 'a :: \{\text{real_normed_field}, \text{banach}\}$
 assumes $r: 0 < r$
 and $sm: \bigwedge x. \text{norm } (x-\xi) < r \implies (\lambda n. a\ n * (x-\xi) ^ n) \text{ sums } (f\ x)$
 and [simp]: $f\ \xi = 0$
 and $m0: a\ m \neq 0$ and $m > 0$
 obtains s where $0 < s$ and $\bigwedge z. z \in \text{cball } \xi\ s - \{\xi\} \implies f\ z \neq 0$
proof -
 have $r \leq \text{conv_radius } a$
 using sm sums_summable by (auto simp: le_conv_radius_iff [where $\xi=\xi$])
 obtain m where $am: a\ m \neq 0$ and az [simp]: $(\bigwedge n. n < m \implies a\ n = 0)$
proof
 show a (LEAST $n. a\ n \neq 0$) $\neq 0$
 by (metis (mono_tags, lifting) m0 LeastI)
qed (fastforce dest!: not_less_Least)
 define b where $b\ i = a\ (i+m) / a\ m$ for i
 define g where $g\ x = \text{suminf } (\lambda i. b\ i * (x-\xi) ^ i)$ for x
 have [simp]: $b\ 0 = 1$
 by (simp add: am b_def)
 { fix $x::'a$
 assume $\text{norm } (x-\xi) < r$
 then have $(\lambda n. (a\ m * (x-\xi) ^ m) * (b\ n * (x-\xi) ^ n)) \text{ sums } (f\ x)$
 using $am\ az\ sm$ sums_zero_iff_shift [of m $(\lambda n. a\ n * (x-\xi) ^ n)$ $f\ x$]
 by (simp add: b_def monoid_mult_class.power_add algebra_simps)
 then have $x \neq \xi \implies (\lambda n. b\ n * (x-\xi) ^ n) \text{ sums } (f\ x / (a\ m * (x-\xi) ^ m))$
 using am by (simp add: sums_mult_D)
 } note $b\text{sums} = \text{this}$
 then have $\text{norm } (x-\xi) < r \implies \text{summable } (\lambda n. b\ n * (x-\xi) ^ n)$ for x
 using sums_summable by (cases $x=\xi$) auto
 then have $r \leq \text{conv_radius } b$
 by (simp add: le_conv_radius_iff [where $\xi=\xi$])
 then have $r/2 < \text{conv_radius } b$
 using not_le order_trans r by fastforce
 then have continuous_on $(\text{cball } \xi\ (r/2))\ g$
 using power_continuous_suminf [of $r/2\ b\ \xi$] by (simp add: g_def)
 then obtain s where $s > 0$ $\bigwedge x. [\text{norm } (x-\xi) \leq s; \text{norm } (x-\xi) \leq r/2] \implies \text{dist}$
 $(g\ x)\ (g\ \xi) < 1/2$
proof (rule continuous_onE)
 show $\xi \in \text{cball } \xi\ (r / 2)\ 1/2 > (0::\text{real})$
 using r by auto
qed (auto simp: dist_commute dist_norm)
 moreover have $g\ \xi = 1$
 by (simp add: g_def)
 ultimately have $gnz: \bigwedge x. [\text{norm } (x-\xi) \leq s; \text{norm } (x-\xi) \leq r/2] \implies (g\ x) \neq 0$

```

  by fastforce
  have  $f x \neq 0$  if  $x \neq \xi$  norm  $(x-\xi) \leq s$  norm  $(x-\xi) \leq r/2$  for  $x$ 
  using bsums [of  $x$ ] that gncz [of  $x$ ] r sums_iff unfolding g_def by fastforce
  then show ?thesis
  apply (rule_tac s=min s (r/2) in that)
  using  $\langle 0 < r \rangle \langle 0 < s \rangle$  by (auto simp: dist_commute dist_norm)
qed

```

4.15 Complex functions and power series

The following defines the power series expansion of a complex function at a given point (assuming that it is analytic at that point).

definition $\text{fps_expansion} :: (\text{complex} \Rightarrow \text{complex}) \Rightarrow \text{complex} \Rightarrow \text{complex fps}$
where

$$\text{fps_expansion } f \ z0 = \text{Abs_fps } (\lambda n. (\text{deriv } \hat{\ } n) f \ z0 / \text{fact } n)$$

lemma $\text{fps_expansion_cong}$:

assumes $\forall_F w$ in nhds $x. f w = g w$

shows $\text{fps_expansion } f \ x = \text{fps_expansion } g \ x$

unfolding fps_expansion_def **using** $\text{assms_higher_deriv_cong_ev}$ **by** fastforce

lemma

fixes $r :: \text{ereal}$

assumes f holomorphic_on eball $z0 \ r$

shows $\text{conv_radius_fps_expansion}: \text{fps_conv_radius } (\text{fps_expansion } f \ z0) \geq r$

and $\text{eval_fps_expansion}: \bigwedge z. z \in \text{eball } z0 \ r \implies \text{eval_fps } (\text{fps_expansion } f \ z0) (z - z0) = f \ z$

and $\text{eval_fps_expansion}': \bigwedge z. \text{norm } z < r \implies \text{eval_fps } (\text{fps_expansion } f \ z0) z = f (z0 + z)$

proof -

have $(\lambda n. \text{fps_nth } (\text{fps_expansion } f \ z0) \ n * (z - z0) \hat{\ } n) \text{ sums } f \ z$

if $z \in \text{ball } z0 \ r'$ $\text{ereal } r' < r$ **for** $z \ r'$

proof -

have f holomorphic_on ball $z0 \ r'$

using holomorphic_on_subset[OF __ ball_eball_mono] assms that **by** force

then show ?thesis

using fps_expansion_def holomorphic_power_series that **by** auto

qed

hence *: $(\lambda n. \text{fps_nth } (\text{fps_expansion } f \ z0) \ n * (z - z0) \hat{\ } n) \text{ sums } f \ z$

if $z \in \text{eball } z0 \ r$ **for** z

using that **by** (subst (asm) eball_conv_UNION_balls) blast

show $\text{fps_conv_radius } (\text{fps_expansion } f \ z0) \geq r$ **unfolding** $\text{fps_conv_radius_def}$

proof (rule conv_radius_geI_ex)

fix $r' :: \text{real}$ **assume** $r': r' > 0$ $\text{ereal } r' < r$

thus $\exists z. \text{norm } z = r' \wedge \text{summable } (\lambda n. \text{fps_nth } (\text{fps_expansion } f \ z0) \ n * z \hat{\ } n)$

using *[of $z0 + \text{of_real } r'$]

by (intro exI[of __ of_real r']) (auto simp: summable_def dist_norm)

qed

```

show eval_fps (fps_expansion f z0) (z - z0) = f z if z ∈ eball z0 r for z
using *[OF that] by (simp add: eval_fps_def sums_iff)
show eval_fps (fps_expansion f z0) z = f (z0 + z) if ereal (norm z) < r for z
using *[of z0 + z] and that by (simp add: eval_fps_def sums_iff dist_norm)
qed

```

We can now show several more facts about power series expansions (at least in the complex case) with relative ease that would have been trickier without complex analysis.

lemma

```

fixes f :: complex fps and r :: ereal
assumes  $\bigwedge z. \text{ereal} (\text{norm } z) < r \implies \text{eval\_fps } f z \neq 0$ 
shows fps_conv_radius_inverse: fps_conv_radius (inverse f)  $\geq \min r$  (fps_conv_radius f)
and eval_fps_inverse:  $\bigwedge z. \text{ereal} (\text{norm } z) < \text{fps\_conv\_radius } f \implies \text{ereal} (\text{norm } z) < r \implies$ 
 $\text{eval\_fps } (\text{inverse } f) z = \text{inverse } (\text{eval\_fps } f z)$ 

```

proof –

```

define R where R = min (fps_conv_radius f) r
have *: fps_conv_radius (inverse f)  $\geq \min r$  (fps_conv_radius f)  $\wedge$ 
 $(\forall z \in \text{eball } 0 (\min (\text{fps\_conv\_radius } f) r). \text{eval\_fps } (\text{inverse } f) z = \text{inverse}$ 
 $(\text{eval\_fps } f z))$ 
proof (cases min r (fps_conv_radius f) > 0)
case True
define f' where f' = fps_expansion ( $\lambda z. \text{inverse } (\text{eval\_fps } f z)$ ) 0
have holo: ( $\lambda z. \text{inverse } (\text{eval\_fps } f z)$ ) holomorphic_on eball 0 (min r (fps_conv_radius f))
using assms by (intro holomorphic_intros) auto
from holo have radius: fps_conv_radius f'  $\geq \min r$  (fps_conv_radius f)
unfolding f'_def by (rule conv_radius_fps_expansion)
have eval_f': eval_fps f' z = inverse (eval_fps f z)
if norm z < fps_conv_radius f norm z < r for z
using that unfolding f'_def by (subst eval_fps_expansion'[OF holo]) auto

```

have f * f' = 1

proof (rule eval_fps_eqD)

```

from radius and True have 0 < min (fps_conv_radius f) (fps_conv_radius f')
f')

```

by (auto simp: min_def split: if_splits)

also have ... $\leq \text{fps_conv_radius } (f * f')$ **by** (rule fps_conv_radius_mult)

finally show ... > 0 .

next

from True **have** R > 0 **by** (auto simp: R_def)

hence eventually ($\lambda z. z \in \text{eball } 0 R$) (nhds 0)

by (intro eventually_nhds_in_open) (auto simp: zero_ereal_def)

thus eventually ($\lambda z. \text{eval_fps } (f * f') z = \text{eval_fps } 1 z$) (nhds 0)

proof eventually_elim

case (elim z)

hence eval_fps (f * f') z = eval_fps f z * eval_fps f' z

```

    using radius by (intro eval_fps_mult)
                      (auto simp: R_def min_def split: if_splits intro: less_trans)
  also have eval_fps f' z = inverse (eval_fps f z)
    using elim by (intro eval_f') (auto simp: R_def)
  also from elim have eval_fps f z ≠ 0
    by (intro assms) (auto simp: R_def)
  hence eval_fps f z * inverse (eval_fps f z) = eval_fps 1 z
    by simp
  finally show eval_fps (f * f') z = eval_fps 1 z .
qed
qed simp_all
hence f' = inverse f
  by (intro fps_inverse_unique [symmetric]) (simp_all add: mult_ac)
with eval_f' and radius show ?thesis by simp
next
case False
hence *: eball 0 R = {}
  by (intro eball_empty) (auto simp: R_def min_def split: if_splits)
show ?thesis
proof safe
  from False have min r (fps_conv_radius f) ≤ 0
    by (simp add: min_def)
  also have 0 ≤ fps_conv_radius (inverse f)
    by (simp add: fps_conv_radius_def conv_radius_nonneg)
  finally show min r (fps_conv_radius f) ≤ ... .
qed (unfold * [unfolded R_def], auto)
qed

from * show fps_conv_radius (inverse f) ≥ min r (fps_conv_radius f) by blast
from * show eval_fps (inverse f) z = inverse (eval_fps f z)
  if ereal (norm z) < fps_conv_radius f ereal (norm z) < r for z
  using that by auto
qed

lemma
  fixes f g :: complex fps and r :: ereal
  defines R ≡ Min {r, fps_conv_radius f, fps_conv_radius g}
  assumes fps_conv_radius f > 0 fps_conv_radius g > 0 r > 0
  assumes nz:  $\bigwedge z. z \in \text{eball } 0 r \implies \text{eval\_fps } g z \neq 0$ 
  shows fps_conv_radius_divide':  $\text{fps\_conv\_radius } (f / g) \geq R$ 
    and eval_fps_divide':
       $\text{ereal } (\text{norm } z) < R \implies \text{eval\_fps } (f / g) z = \text{eval\_fps } f z / \text{eval\_fps } g z$ 
proof -
  from nz[of 0] and  $\langle r > 0 \rangle$  have nz':  $\text{fps\_nth } g 0 \neq 0$ 
    by (auto simp: eval_fps_at_0 zero_ereal_def)
  have  $R \leq \text{min } r (\text{fps\_conv\_radius } g)$ 
    by (auto simp: R_def intro: min.coboundedI2)
  also have  $\text{min } r (\text{fps\_conv\_radius } g) \leq \text{fps\_conv\_radius } (\text{inverse } g)$ 
    by (intro fps_conv_radius_inverse assms) (auto simp: zero_ereal_def)

```

```

finally have radius: fps_conv_radius (inverse g) ≥ R .
have R ≤ min (fps_conv_radius f) (fps_conv_radius (inverse g))
  by (intro radius min.boundedI) (auto simp: R_def intro: min.coboundedI1
min.coboundedI2)
also have ... ≤ fps_conv_radius (f * inverse g)
  by (rule fps_conv_radius_mult)
also have f * inverse g = f / g
  by (intro fps_divide_unit [symmetric] nz')
finally show fps_conv_radius (f / g) ≥ R .

assume z: ereal (norm z) < R
have eval_fps (f * inverse g) z = eval_fps f z * eval_fps (inverse g) z
  using radius by (intro eval_fps_mult less_le_trans[OF z])
  (auto simp: R_def intro: min.coboundedI1 min.coboundedI2)
also have eval_fps (inverse g) z = inverse (eval_fps g z) using ‹r > 0›
  by (intro eval_fps_inverse[where r = r] less_le_trans[OF z] nz)
  (auto simp: R_def intro: min.coboundedI1 min.coboundedI2)
also have f * inverse g = f / g by fact
finally show eval_fps (f / g) z = eval_fps f z / eval_fps g z
  by (simp add: field_split_simps)
qed

```

lemma

```

fixes f g :: complex fps and r :: ereal
defines R ≡ Min {r, fps_conv_radius f, fps_conv_radius g}
assumes subdegree g ≤ subdegree f
assumes fps_conv_radius f > 0 fps_conv_radius g > 0 r > 0
assumes ∧z. z ∈ eball 0 r ⇒ z ≠ 0 ⇒ eval_fps g z ≠ 0
shows fps_conv_radius_divide: fps_conv_radius (f / g) ≥ R
  and eval_fps_divide:
    ereal (norm z) < R ⇒ c = fps_nth f (subdegree g) / fps_nth g (subdegree
g) ⇒
    eval_fps (f / g) z = (if z = 0 then c else eval_fps f z / eval_fps g z)
proof –
  define f' g' where f' = fps_shift (subdegree g) f and g' = fps_shift (subdegree
g) g
  have f_eq: f = f' * fps_X ^ subdegree g and g_eq: g = g' * fps_X ^ subdegree
g
  unfolding f'_def g'_def by (rule subdegree_decompose' le_refl | fact)+
  have subdegree: subdegree f' = subdegree f – subdegree g subdegree g' = 0
  using assms(2) by (simp_all add: f'_def g'_def)
  have [simp]: fps_conv_radius f' = fps_conv_radius f fps_conv_radius g' =
fps_conv_radius g
  by (simp_all add: f'_def g'_def)
  have [simp]: fps_nth f' 0 = fps_nth f (subdegree g)
    fps_nth g' 0 = fps_nth g (subdegree g) by (simp_all add: f'_def
g'_def)
  have g_nz: g ≠ 0
proof –

```

```

define  $z :: \text{complex}$  where  $z = (\text{if } r = \infty \text{ then } 1 \text{ else of\_real (real\_of\_ereal } r / 2))$ 
from  $\langle r > 0 \rangle$  have  $z \in \text{eball } 0 \ r$ 
  by (cases  $r$ ) (auto simp:  $z\_def$   $\text{eball\_def}$ )
moreover have  $z \neq 0$  using  $\langle r > 0 \rangle$ 
  by (cases  $r$ ) (auto simp:  $z\_def$ )
ultimately have  $\text{eval\_fps } g \ z \neq 0$  by (rule  $\text{assms}(6)$ )
thus  $g \neq 0$  by auto
qed
have  $fg: f / g = f' * \text{inverse } g'$ 
  by (subst  $f\_eq$ , subst (2)  $g\_eq$ ) (insert  $g\_nz$ , simp add:  $\text{fps\_divide\_unit}$ )

have  $g'_nz: \text{eval\_fps } g' \ z \neq 0$  if  $z: \text{norm } z < \min r (\text{fps\_conv\_radius } g)$  for  $z$ 
proof (cases  $z = 0$ )
  case False
    with  $\text{assms}$  and  $z$  have  $\text{eval\_fps } g \ z \neq 0$  by auto
    also from  $z$  have  $\text{eval\_fps } g \ z = \text{eval\_fps } g' \ z * z^{\wedge} \text{subdegree } g$ 
      by (subst  $g\_eq$ ) (auto simp:  $\text{eval\_fps\_mult}$ )
    finally show ?thesis by auto
qed (use  $\langle g \neq 0 \rangle$  in  $\langle \text{auto simp: } g'_def \ \text{eval\_fps\_at\_0} \rangle$ )

have  $R \leq \min (\min r (\text{fps\_conv\_radius } g)) (\text{fps\_conv\_radius } g')$ 
  by (auto simp:  $R\_def$   $\text{min.coboundedI1}$   $\text{min.coboundedI2}$ )
also have  $\dots \leq \text{fps\_conv\_radius } (\text{inverse } g')$ 
  using  $g'_nz$  by (rule  $\text{fps\_conv\_radius\_inverse}$ )
finally have  $\text{conv\_radius\_inv}: R \leq \text{fps\_conv\_radius } (\text{inverse } g')$  .
hence  $R \leq \text{fps\_conv\_radius } (f' * \text{inverse } g')$ 
  by (intro  $\text{order.trans}[OF \ \text{fps\_conv\_radius\_mult}]$ )
  (auto simp:  $R\_def$   $\text{intro: min.coboundedI1 min.coboundedI2}$ )
thus  $\text{fps\_conv\_radius } (f / g) \geq R$  by (simp add:  $fg$ )

fix  $z \ c :: \text{complex}$  assume  $z: \text{ereal } (\text{norm } z) < R$ 
assume  $c: c = \text{fps\_nth } f (\text{subdegree } g) / \text{fps\_nth } g (\text{subdegree } g)$ 
show  $\text{eval\_fps } (f / g) \ z = (\text{if } z = 0 \text{ then } c \text{ else } \text{eval\_fps } f \ z / \text{eval\_fps } g \ z)$ 
proof (cases  $z = 0$ )
  case False
    from  $z$  and  $\text{conv\_radius\_inv}$  have  $\text{ereal } (\text{norm } z) < \text{fps\_conv\_radius } (\text{inverse } g')$ 
      by  $\text{simp}$ 
    with  $z$  have  $\text{eval\_fps } (f / g) \ z = \text{eval\_fps } f' \ z * \text{eval\_fps } (\text{inverse } g') \ z$ 
      unfolding  $fg$  by (subst  $\text{eval\_fps\_mult}$ ) (auto simp:  $R\_def$ )
    also have  $\text{eval\_fps } (\text{inverse } g') \ z = \text{inverse } (\text{eval\_fps } g' \ z)$ 
      using  $z$  by (intro  $\text{eval\_fps\_inverse}[of \ \min r (\text{fps\_conv\_radius } g')] \ g'_nz$ )
      (auto simp:  $R\_def$ )
    also have  $\text{eval\_fps } f' \ z * \dots = \text{eval\_fps } f \ z / \text{eval\_fps } g \ z$ 
      using  $z$  False  $\text{assms}(2)$  by (simp add:  $f'_def \ g'_def \ \text{eval\_fps\_shift } R\_def$ )
    finally show ?thesis using False by  $\text{simp}$ 
qed (simp_all add:  $\text{eval\_fps\_at\_0 } fg \ \text{field\_simps } c$ )
qed

```

```

lemma has_fps_expansion_fps_expansion [intro]:
  assumes open A 0 ∈ A f holomorphic_on A
  shows f has_fps_expansion fps_expansion f 0
proof –
  from assms obtain r where r > 0 and r: ball 0 r ⊆ A
    by (auto simp: open_contains_ball)
  with assms have holo: f holomorphic_on eball 0 (ereal r)
    by auto
  have r ≤ fps_conv_radius (fps_expansion f 0)
    using holo by (intro conv_radius_fps_expansion) auto
  then have ... > 0
    by (simp add: ereal_le_less ‹r > 0› zero_ereal_def)
  moreover have eventually (λz. z ∈ ball 0 r) (nhds 0)
    using ‹r > 0› by (intro eventually_nhds_in_open) auto
  hence eventually (λz. eval_fps (fps_expansion f 0) z = f z) (nhds 0)
    by eventually_elim (subst eval_fps_expansion'[OF holo], auto)
  ultimately show ?thesis using ‹r > 0› by (auto simp: has_fps_expansion_def)
qed

```

```

lemma fps_conv_radius_tan:
  fixes c :: complex
  assumes c ≠ 0
  shows fps_conv_radius (fps_tan c) ≥ pi / (2 * norm c)
proof –
  have fps_conv_radius (fps_tan c) ≥
    Min {pi / (2 * norm c), fps_conv_radius (fps_sin c), fps_conv_radius
(fps_cos c)}
    unfolding fps_tan_def
  proof (rule fps_conv_radius_divide)
    fix z :: complex assume z ∈ eball 0 (pi / (2 * norm c))
    with cos_eq_zero_imp_norm_ge[of c*z] assms
      show eval_fps (fps_cos c) z ≠ 0 by (auto simp: norm_mult_field_simps)
  qed (insert assms, auto)
  thus ?thesis by (simp add: min_def)
qed

```

```

lemma eval_fps_tan:
  fixes c :: complex
  assumes norm z < pi / (2 * norm c)
  shows eval_fps (fps_tan c) z = tan (c * z)
proof (cases c = 0)
  case False
  show ?thesis unfolding fps_tan_def
  proof (subst eval_fps_divide'[where r = pi / (2 * norm c)])
    fix z :: complex assume z ∈ eball 0 (pi / (2 * norm c))
    with cos_eq_zero_imp_norm_ge[of c*z] assms
      show eval_fps (fps_cos c) z ≠ 0 using False by (auto simp: norm_mult
field_simps)

```

```
qed (use False assms in ⟨auto simp: field_simps tan_def⟩)
qed simp_all
```

```
end
```

5 Conformal Mappings and Consequences of Cauchy's Integral Theorem

By John Harrison et al. Ported from HOL Light by L C Paulson (2016)

Also Cauchy's residue theorem by Wenda Li (2016)

```
theory Conformal_Mappings
imports Cauchy_Integral_Formula
```

```
begin
```

5.1 Analytic continuation

proposition *isolated_zeros*:

```
assumes holf: f holomorphic_on S
and open S connected S ξ ∈ S f ξ = 0 β ∈ S f β ≠ 0
obtains r where 0 < r and ball ξ r ⊆ S and
  ∧z. z ∈ ball ξ r - {ξ} ⇒ f z ≠ 0
```

proof –

```
obtain r where 0 < r and r: ball ξ r ⊆ S
using ⟨open S⟩ ⟨ξ ∈ S⟩ open_contains_ball_eq by blast
have powf: ((λn. (deriv ~ n) f ξ / (fact n) * (z - ξ) ^ n) sums f z) if z ∈ ball
ξ r for z
by (intro holomorphic_power_series [OF _ that] holomorphic_on_subset [OF
holf r])
obtain m where m: (deriv ~ m) f ξ / (fact m) ≠ 0
using holomorphic_fun_eq_0_on_connected [OF holf ⟨open S⟩ ⟨connected S⟩
_ ⟨ξ ∈ S⟩ ⟨β ∈ S⟩] ⟨f β ≠ 0⟩
by auto
then have m ≠ 0 using assms(5) funpow_0 by fastforce
obtain s where 0 < s and s: ∧z. z ∈ cball ξ s - {ξ} ⇒ f z ≠ 0
using powser_0_nonzero [OF ⟨0 < r⟩ powf ⟨f ξ = 0⟩ m]
by (metis ⟨m ≠ 0⟩ dist_norm mem_ball norm_minus_commute not_gr_zero)
have 0 < min r s by (simp add: ⟨0 < r⟩ ⟨0 < s⟩)
then show thesis
apply (rule that)
using r s by auto
qed
```

proposition *analytic_continuation*:

```
assumes holf: f holomorphic_on S
and open S and connected S
and U ⊆ S and ξ ∈ S
```



```

    and  $\xi$  islimpt  $U$ 
    and  $fU0$  [simp]:  $\bigwedge z. z \in U \implies f z = 0$ 
    and  $w \in S$ 
    shows  $f w = 0$ 
  proof -
    obtain  $e$  where  $0 < e$  and  $e$ :  $cball \ \xi \ e \subseteq S$ 
    using  $\langle open \ S \rangle \langle \xi \in S \rangle open\_contains\_cball\_eq$  by blast
    define  $T$  where  $T = cball \ \xi \ e \cap U$ 
    have  $contf$ :  $continuous\_on$  (closure  $T$ )  $f$ 
    by (metis  $T\_def$  closed_cball closure_minimal  $e$  holf holomorphic_on_imp_continuous_on
        holomorphic_on_subset inf.cobounded1)
    have  $fT0$  [simp]:  $\bigwedge x. x \in T \implies f x = 0$ 
    by (simp add:  $T\_def$ )
    have  $\bigwedge r. [\forall e > 0. \exists x' \in U. x' \neq \xi \wedge dist \ x' \ \xi < e; 0 < r] \implies \exists x' \in cball \ \xi \ e \cap$ 
 $U. x' \neq \xi \wedge dist \ x' \ \xi < r$ 
    by (metis  $\langle 0 < e \rangle$  IntI dist_commute less_eq_real_def mem_cball min_less_iff_conj)
    then have  $\xi$  islimpt  $T$  using  $\langle \xi$  islimpt  $U \rangle$ 
    by (auto simp:  $T\_def$  islimpt_approachable)
    then have  $\xi \in closure \ T$ 
    by (simp add: closure_def)
    then have  $f \ \xi = 0$ 
    by (auto simp: continuous_constant_on_closure [OF  $contf$ ])
    moreover have  $\bigwedge r. [0 < r; \bigwedge z. z \in ball \ \xi \ r - \{\xi\} \implies f z \neq 0] \implies False$ 
    by (metis open_ball  $\langle \xi$  islimpt  $T \rangle$  centre_in_ball  $fT0$  insertE insert_Diff
        islimptE)
    ultimately show ?thesis
    by (metis  $\langle open \ S \rangle \langle connected \ S \rangle \langle \xi \in S \rangle \langle w \in S \rangle$  holf isolated_zeros)
  qed

```

```

corollary analytic_continuation_open:
  assumes  $open \ s$  and  $open \ s'$  and  $s \neq \{\}$  and  $connected \ s'$ 
  and  $s \subseteq s'$ 
  assumes  $f$  holomorphic_on  $s'$  and  $g$  holomorphic_on  $s'$ 
  and  $\bigwedge z. z \in s \implies f z = g z$ 
  assumes  $z \in s'$ 
  shows  $f z = g z$ 
  proof -
    from  $\langle s \neq \{\} \rangle$  obtain  $\xi$  where  $\xi \in s$  by auto
    with  $\langle open \ s \rangle$  have  $\xi$ :  $\xi$  islimpt  $s$ 
    by (intro interior_limit_point) (auto simp: interior_open)
    have  $f z - g z = 0$ 
    by (rule analytic_continuation[of  $\lambda z. f z - g z \ s' \ s \ \xi$ ])
    (insert assms  $\langle \xi \in s \rangle \ \xi$ , auto intro: holomorphic_intros)
    thus ?thesis by simp
  qed

```

```

corollary analytic_continuation':
  assumes  $f$  holomorphic_on  $S$  open  $S$  connected  $S$ 
  and  $U \subseteq S \ \xi \in S \ \xi$  islimpt  $U$ 

```

```

    and  $f$  constant_on  $U$ 
  shows  $f$  constant_on  $S$ 
proof -
  obtain  $c$  where  $c: \bigwedge x. x \in U \implies f x - c = 0$ 
    by (metis  $\langle f$  constant_on  $U \rangle$  constant_on_def diff_self)
  have  $(\lambda z. f z - c)$  holomorphic_on  $S$ 
    using assms by (intro holomorphic_intros)
  with  $c$  analytic_continuation assms have  $\bigwedge x. x \in S \implies f x - c = 0$ 
    by blast
  then show ?thesis
    unfolding constant_on_def by force
qed

```

```

lemma holomorphic_compact_finite_zeros:
  assumes  $S: f$  holomorphic_on  $S$  open  $S$  connected  $S$ 
    and compact  $K$   $K \subseteq S$ 
    and  $\neg f$  constant_on  $S$ 
  shows finite  $\{z \in K. f z = 0\}$ 
proof (rule ccontr)
  assume infinite  $\{z \in K. f z = 0\}$ 
  then obtain  $z$  where  $z \in K$  and  $z: z$  islimpt  $\{z \in K. f z = 0\}$ 
    using  $\langle$ compact  $K \rangle$  by (auto simp: compact_eq_Bolzano_Weierstrass)
  moreover have  $\{z \in K. f z = 0\} \subseteq S$ 
    using  $\langle K \subseteq S \rangle$  by blast
  ultimately show False
    using assms analytic_continuation [OF  $S$ ] unfolding constant_on_def
    by blast
qed

```

```

lemma holomorphic_countable_zeros:
  assumes  $S: f$  holomorphic_on  $S$  open  $S$  connected  $S$  and  $f$ sigma  $S$ 
    and  $\neg f$  constant_on  $S$ 
  shows countable  $\{z \in S. f z = 0\}$ 
proof -
  obtain  $F::\text{nat} \Rightarrow$  complex set
    where  $F: \text{range } F \subseteq \text{Collect compact}$  and  $\text{Seq}: S = (\bigcup i. F i)$ 
    using  $\langle f$ sigma  $S \rangle$  by (meson  $f$ sigma_Union_compact)
  have  $\text{fin}: \text{finite } \{z \in F i. f z = 0\}$  for  $i$ 
    using holomorphic_compact_finite_zeros assms  $F$  Seq Union_iff by blast
  have  $\{z \in S. f z = 0\} = (\bigcup i. \{z \in F i. f z = 0\})$ 
    using Seq by auto
  with  $\text{fin}$  show ?thesis
    by (simp add: countable_finite)
qed

```

```

lemma holomorphic_countable_equal:
  assumes  $f$  holomorphic_on  $S$   $g$  holomorphic_on  $S$  open  $S$  connected  $S$  and  $f$ sigma  $S$ 
    and  $\text{eq}: \text{uncountable } \{z \in S. f z = g z\}$ 

```

```

shows  $S \subseteq \{z \in S. f z = g z\}$ 
proof -
  obtain  $z$  where  $z: z \in S \ f z = g z$ 
    using eq not_finite_existsD uncountable_infinite by blast
  have  $(\lambda x. f x - g x)$  holomorphic_on  $S$ 
    by (simp add: assms holomorphic_on_diff)
  then have  $(\lambda x. f x - g x)$  constant_on  $S$ 
    using holomorphic_countable_zeros assms by force
  with  $z$  have  $\bigwedge x. x \in S \implies f x - g x = 0$ 
    unfolding constant_on_def by force
  then show ?thesis
    by auto
qed

```

```

lemma holomorphic_countable_equal_UNIV:
  assumes  $fg: f$  holomorphic_on UNIV  $g$  holomorphic_on UNIV
    and  $eq: \text{uncountable } \{z. f z = g z\}$ 
  shows  $f=g$ 
  using holomorphic_countable_equal [OF fg] eq by fastforce

```

5.2 Open mapping theorem

```

lemma holomorphic_contract_to_zero:
  assumes  $contf: \text{continuous\_on } (cball \ \xi \ r)$   $f$ 
    and  $holf: f$  holomorphic_on ball  $\xi \ r$ 
    and  $0 < r$ 
    and  $norm\_less: \bigwedge z. norm(\xi - z) = r \implies norm(f \ \xi) < norm(f \ z)$ 
  obtains  $z$  where  $z \in ball \ \xi \ r \ f z = 0$ 
proof -
  { assume  $fnz: \bigwedge w. w \in ball \ \xi \ r \implies f w \neq 0$ 
    then have  $0 < norm(f \ \xi)$ 
      by (simp add: <0 < r>)
    have  $fnz': \bigwedge w. w \in cball \ \xi \ r \implies f w \neq 0$ 
      by (metis norm_less dist_norm fnz less_eq_real_def mem_ball mem_cball
norm_not_less_zero norm_zero)
    have  $frontier(cball \ \xi \ r) \neq \{\}$ 
      using <0 < r> by simp
    define  $g$  where  $[abs\_def]: g \ z = \text{inverse } (f \ z)$  for  $z$ 
    have  $contg: \text{continuous\_on } (cball \ \xi \ r)$   $g$ 
      unfolding  $g\_def$  using  $contf$  continuous_on_inverse  $fnz'$  by blast
    have  $holg: g$  holomorphic_on ball  $\xi \ r$ 
      unfolding  $g\_def$  using  $fnz$  holf holomorphic_on_inverse by blast
    have  $frontier(cball \ \xi \ r) \subseteq cball \ \xi \ r$ 
      by (simp add: subset_iff)
    then have  $contf': \text{continuous\_on } (frontier(cball \ \xi \ r))$   $f$ 
      and  $contg': \text{continuous\_on } (frontier(cball \ \xi \ r))$   $g$ 
      by (blast intro:  $contf$   $contg$  continuous_on_subset)+
    have  $froc: frontier(cball \ \xi \ r) \neq \{\}$ 
      using <0 < r> by simp
  }

```

```

moreover have continuous_on (frontier (cball  $\xi$   $r$ )) (norm o  $f$ )
  using contf' continuous_on_compose continuous_on_norm_id by blast
ultimately obtain  $w$  where  $w: w \in \text{frontier}(cball \xi r)$ 
  and now:  $\bigwedge x. x \in \text{frontier}(cball \xi r) \implies \text{norm}(f w) \leq \text{norm}$ 
( $f x$ )
  using continuous_attains_inf [OF compact_frontier [OF compact_cball]]
  by (metis comp_apply)
then have  $fw: 0 < \text{norm}(f w)$ 
  by (simp add: fnz')
have continuous_on (frontier (cball  $\xi$   $r$ )) (norm o  $g$ )
  using contg' continuous_on_compose continuous_on_norm_id by blast
then obtain  $v$  where  $v: v \in \text{frontier}(cball \xi r)$ 
  and nov:  $\bigwedge x. x \in \text{frontier}(cball \xi r) \implies \text{norm}(g v) \geq \text{norm}(g x)$ 
  using continuous_attains_sup [OF compact_frontier [OF compact_cball]]
froc] by force
then have  $fv: 0 < \text{norm}(f v)$ 
  by (simp add: fnz')
have  $\text{norm}((\text{deriv } \hat{\sim} 0) g \xi) \leq \text{fact } 0 * \text{norm}(g v) / r \wedge 0$ 
  by (rule Cauchy_inequality [OF holg contg <0 < r]) (simp add: dist_norm
nov)
then have  $cmod(g \xi) \leq cmod(g v)$ 
  by simp
moreover have  $cmod(\xi - w) = r$ 
  by (metis (no_types) dist_norm frontier_cball mem_sphere  $w$ )
ultimately obtain  $wr: \text{norm}(\xi - w) = r$  and  $nfw: \text{norm}(f w) \leq \text{norm}(f \xi)$ 
  unfolding  $g\_def$ 
  by (metis (no_types) <0 < cmod( $f \xi$ )> less_imp_inverse_less norm_inverse
not_le now order_trans  $v$ )
  with  $fw$  have False
  using norm_less by force
}
with that show ?thesis by blast
qed

```

theorem open_mapping_thm:

```

assumes holgf:  $f$  holomorphic_on  $S$ 
  and  $S$ : open  $S$  and connected  $S$ 
  and open  $U$  and  $U \subseteq S$ 
  and  $fne: \neg f$  constant_on  $S$ 
shows open ( $f ' U$ )

```

proof –

```

have *: open ( $f ' U$ )

```

```

  if  $U \neq \{\}$  and  $U$ : open  $U$  connected  $U$  and  $f$  holomorphic_on  $U$  and
fneU:  $\bigwedge x. \exists y \in U. f y \neq x$ 

```

```

  for  $U$ 

```

```

proof (clarsimp simp: open_contains_ball)

```

```

  fix  $\xi$  assume  $\xi: \xi \in U$ 

```

```

  show  $\exists e > 0. \text{ball}(f \xi) e \subseteq f ' U$ 

```

```

proof –

```

```

have hol: ( $\lambda z. f z - f \xi$ ) holomorphic_on U
  by (rule holomorphic_intros that)+
obtain s where  $0 < s$  and sbU: ball  $\xi$  s  $\subseteq U$ 
  and sne:  $\bigwedge z. z \in \text{ball } \xi \ s - \{\xi\} \implies (\lambda z. f z - f \xi) z \neq 0$ 
  using isolated_zeros [OF hol U  $\xi$ ] by (metis fneU right_minus_eq)
obtain r where  $0 < r$  and r: cball  $\xi$  r  $\subseteq \text{ball } \xi$  s
  using  $\langle 0 < s \rangle$  by (rule_tac r=s/2 in that) auto
have cball  $\xi$  r  $\subseteq U$ 
  using sbU r by blast
then have frsbU: frontier (cball  $\xi$  r)  $\subseteq U$ 
  using Diff_subset frontier_def order_trans by fastforce
then have cof: compact (frontier (cball  $\xi$  r))
  by blast
have frne: frontier (cball  $\xi$  r)  $\neq \{\}$ 
  using  $\langle 0 < r \rangle$  by auto
have contfr: continuous_on (frontier (cball  $\xi$  r)) ( $\lambda z. \text{norm } (f z - f \xi)$ )
  by (metis continuous_on_norm continuous_on_subset frsbU hol holomorphic_on_imp_continuous_on)
obtain w where norm ( $\xi - w$ ) = r
  and w: ( $\bigwedge z. \text{norm } (\xi - z) = r \implies \text{norm } (f w - f \xi) \leq \text{norm } (f z - f \xi)$ )
using continuous_attains_inf [OF cof frne contfr] by (auto simp: dist_norm)
moreover define  $\varepsilon$  where  $\varepsilon \equiv \text{norm } (f w - f \xi) / 3$ 
ultimately have  $0 < \varepsilon$ 
  using  $\langle 0 < r \rangle$  dist_complex_def r sne by auto
have ball (f  $\xi$ )  $\varepsilon \subseteq f^{-1} U$ 
proof
  fix  $\gamma$ 
  assume  $\gamma: \gamma \in \text{ball } (f \xi) \ \varepsilon$ 
  have *: cmod ( $\gamma - f \xi$ ) < cmod ( $\gamma - f z$ ) if cmod ( $\xi - z$ ) = r for z
  proof -
    have lt: cmod (f w - f  $\xi$ ) / 3 < cmod ( $\gamma - f z$ )
      using w [OF that]  $\gamma$ 
      using dist_triangle2 [of f  $\xi$   $\gamma$  f z] dist_triangle2 [of f  $\xi$  f z  $\gamma$ ]
      by (simp add:  $\varepsilon$ _def dist_norm norm_minus_commute)
    show ?thesis
      by (metis  $\varepsilon$ _def dist_commute dist_norm less_trans lt mem_ball  $\gamma$ )
  qed
have continuous_on (cball  $\xi$  r) ( $\lambda z. \gamma - f z$ )
  using  $\langle \text{cball } \xi \ r \subseteq U \rangle$   $\langle f \text{ holomorphic\_on } U \rangle$ 
by (force intro: continuous_intros continuous_on_subset holomorphic_on_imp_continuous_on)
moreover have ( $\lambda z. \gamma - f z$ ) holomorphic_on ball  $\xi$  r
  using  $\langle \text{cball } \xi \ r \subseteq U \rangle$  ball_subset_cball holomorphic_on_subset that(4)
  by (intro holomorphic_intros) blast
ultimately obtain z where  $z \in \text{ball } \xi \ r$   $\gamma - f z = 0$ 
  using *  $\langle 0 < r \rangle$  holomorphic_contract_to_zero by blast
then show  $\gamma \in f^{-1} U$ 
  using  $\langle \text{cball } \xi \ r \subseteq U \rangle$  by fastforce
qed

```

```

    then show ?thesis using ‹0 < ε› by blast
  qed
qed
have open (f ' X) if X ∈ components U for X
proof -
  have holfU: f holomorphic_on U
    using ‹U ⊆ S› holf holomorphic_on_subset by blast
  have X ≠ {}
    using that by (simp add: in_components_nonempty)
  moreover have open X
    using that ‹open U› open_components by auto
  moreover have connected X
    using that in_components_maximal by blast
  moreover have f holomorphic_on X
    by (meson that holfU holomorphic_on_subset in_components_maximal)
  moreover have ∃ y ∈ X. f y ≠ x for x
  proof (rule ccontr)
    assume not: ¬ (∃ y ∈ X. f y ≠ x)
    have X ⊆ S
      using ‹U ⊆ S› in_components_subset that by blast
    obtain w where w: w ∈ X using ‹X ≠ {}› by blast
    have wis: w islimpt X
      using w ‹open X› interior_eq by auto
    have hol: (λz. f z - x) holomorphic_on S
      by (simp add: holf holomorphic_on_diff)
    with fne [unfolded constant_on_def]
      analytic_continuation[OF hol S ‹connected S› ‹X ⊆ S› _ wis] not ‹X ⊆
S› w
    show False by auto
  qed
ultimately show ?thesis
  by (rule *)
qed
then have open (f ' ⋃ (components U))
  by (metis (no_types, lifting) imageE image_Union open_Union)
then show ?thesis
  by force
qed

```

No need for S to be connected. But the nonconstant condition is stronger.

corollary *open_mapping_thm2*:

```

assumes holf: f holomorphic_on S
  and S: open S
  and open U U ⊆ S
  and fnc: ⋀ X. [[open X; X ⊆ S; X ≠ {}]] ⇒ ¬ f constant_on X
shows open (f ' U)
proof -
  have S = ⋃ (components S) by simp
  with ‹U ⊆ S› have U = (⋃ C ∈ components S. C ∩ U) by auto

```

```

then have  $f' U = (\bigcup C \in \text{components } S. f' (C \cap U))$ 
  using image_UN by fastforce
moreover
{ fix C assume  $C \in \text{components } S$ 
  with  $S \langle C \in \text{components } S \rangle \text{ open\_components in\_components\_connected}$ 
  have  $C: \text{open } C \text{ connected } C$  by auto
  have  $C \subseteq S$ 
    by (metis  $\langle C \in \text{components } S \rangle \text{ in\_components\_maximal}$ )
  have  $\text{nf}: \neg f \text{ constant\_on } C$ 
    using  $\langle \text{open } C \rangle \langle C \in \text{components } S \rangle \langle C \subseteq S \rangle \text{ fnc in\_components\_nonempty}$ 
  by blast
  have  $f \text{ holomorphic\_on } C$ 
    by (metis holf holomorphic_on_subset  $\langle C \subseteq S \rangle$ )
  then have  $\text{open } (f' (C \cap U))$ 
    by (meson  $C \langle \text{open } U \rangle \text{ inf\_le1 nf open\_Int open\_mapping\_thm}$ )
} ultimately show ?thesis
  by force
qed

```

```

corollary open_mapping_thm3:
  assumes holf:  $f \text{ holomorphic\_on } S$ 
    and open S and injf:  $\text{inj\_on } f S$ 
  shows  $\text{open } (f' S)$ 
proof (rule open_mapping_thm2 [OF holf])
  show  $\bigwedge X. [\text{open } X; X \subseteq S; X \neq \{\}] \implies \neg f \text{ constant\_on } X$ 
    using inj_on_subset injective_not_constant injf by blast
qed (use assms in auto)

```

5.3 Maximum modulus principle

If f is holomorphic, then its norm (modulus) cannot exhibit a true local maximum that is properly within the domain of f .

proposition *maximum_modulus_principle*:

```

assumes holf:  $f \text{ holomorphic\_on } S$ 
  and  $S: \text{open } S \text{ and connected } S$ 
  and  $\text{open } U \text{ and } U \subseteq S \text{ and } \xi \in U$ 
  and  $\text{no}: \bigwedge z. z \in U \implies \text{norm}(f z) \leq \text{norm}(f \xi)$ 
shows  $f \text{ constant\_on } S$ 

```

proof (rule ccontr)

```

assume  $\neg f \text{ constant\_on } S$ 

```

```

then have  $\text{open } (f' U)$ 

```

```

  using open_mapping_thm assms by blast

```

```

moreover have  $\neg \text{open } (f' U)$ 

```

proof –

```

have  $\exists t. \text{cmod } (f \xi - t) < e \wedge t \notin f' U$  if  $0 < e$  for  $e$ 

```

```

  using that

```

```

  apply (rule_tac  $x = \text{if } 0 < \text{Re}(f \xi) \text{ then } f \xi + (e/2) \text{ else } f \xi - (e/2)$  in exI)

```

```

  apply (simp add: dist_norm)

```

```

  apply (fastforce simp: cmod_Re_le_iff dest!: no dest: sym)

```

```

done
then show ?thesis
  unfolding open_contains_ball by (metis ⟨ξ ∈ U⟩ contra_subsetD dist_norm
imageI mem_ball)
qed
ultimately show False
  by blast
qed

```

proposition *maximum_modulus_frontier*:

```

assumes holf: f holomorphic_on (interior S)
  and conf: continuous_on (closure S) f
  and bos: bounded S
  and leB:  $\bigwedge z. z \in \text{frontier } S \implies \text{norm}(f z) \leq B$ 
  and ξ ∈ S
shows norm(f ξ) ≤ B
proof -
  have compact (closure S) using bos
    by (simp add: bounded_closure compact_eq_bounded_closed)
  moreover have continuous_on (closure S) (cmod ∘ f)
    using conf continuous_on_compose continuous_on_norm_id by blast
  ultimately obtain z where z ∈ closure S and z:  $\bigwedge y. y \in \text{closure } S \implies (cmod$ 
  ∘ f) y ≤ (cmod ∘ f) z
    using continuous_attains_sup [of closure S norm ∘ f] ⟨ξ ∈ S⟩ by auto
  then consider z ∈ frontier S | z ∈ interior S using frontier_def by auto
  then have norm(f z) ≤ B
  proof cases
    case 1 then show ?thesis using leB by blast
  next
    case 2
    have f_constant_on (connected_component_set (interior S) z)
    proof (rule maximum_modulus_principle)
      show f holomorphic_on connected_component_set (interior S) z
        by (metis connected_component_subset holf holomorphic_on_subset)
      show zin: z ∈ connected_component_set (interior S) z
        by (simp add: 2)
      show  $\bigwedge W. W \in \text{connected\_component\_set } (interior\ S)\ z \implies cmod (f W)$ 
      ≤ cmod (f z)
        using closure_def connected_component_subset z by fastforce
    qed (auto simp: open_connected_component)
    then obtain c where c:  $\bigwedge w. w \in \text{connected\_component\_set } (interior\ S)\ z$ 
     $\implies f w = c$ 
      by (auto simp: constant_on_def)
    have f 'closure(connected_component_set (interior S) z) ⊆ {c}
    proof (rule image_closure_subset)
      show continuous_on (closure (connected_component_set (interior S) z)) f
        by (meson closure_mono connected_component_subset conf continuous_on_subset interior_subset)
    qed (use c in auto)

```



```

then have cc:  $\bigwedge w. w \in \text{closure}(\text{connected\_component\_set}(\text{interior } S) z) \implies$ 
 $f w = c$  by blast
have connected_component (interior S) z z
  by (simp add: 2)
moreover have connected_component_set (interior S) z  $\neq$  UNIV
by (metis bos bounded_interior connected_component_eq_UNIV not_bounded_UNIV)
ultimately have frontier(connected_component_set (interior S) z)  $\neq$  {}
  by (meson 2 connected_component_eq_empty frontier_not_empty)
then obtain w where w:  $w \in \text{frontier}(\text{connected\_component\_set}(\text{interior } S) z)$ 
  by auto
then have norm (f z) = norm (f w) by (simp add: 2 c cc frontier_def)
also have ...  $\leq B$ 
  using w frontier_interior_subset frontier_of_connected_component_subset
  by (blast intro: leB)
finally show ?thesis .
qed
then show ?thesis
  using z  $\langle \xi \in S \rangle$  closure_subset by fastforce
qed

```

corollary maximum_real_frontier:

```

assumes holf: f holomorphic_on (interior S)
  and conf: continuous_on (closure S) f
  and bos: bounded S
  and leB:  $\bigwedge z. z \in \text{frontier } S \implies \text{Re}(f z) \leq B$ 
  and  $\xi \in S$ 
shows  $\text{Re}(f \xi) \leq B$ 
using maximum_modulus_frontier [of exp o f S exp B]
  Transcendental.continuous_on_exp holomorphic_on_compose holomorphic_on_exp
  assms
by auto

```

5.4 Factoring out a zero according to its order

lemma holomorphic_factor_order_of_zero:

```

assumes holf: f holomorphic_on S
  and os: open S
  and  $\xi \in S$   $0 < n$ 
  and dnz: (deriv  $\hat{\sim}$  n) f  $\xi \neq 0$ 
  and dfz:  $\bigwedge i. \llbracket 0 < i; i < n \rrbracket \implies (\text{deriv } \hat{\sim} i) f \xi = 0$ 
obtains g r where  $0 < r$ 
  g holomorphic_on ball  $\xi$  r
   $\bigwedge w. w \in \text{ball } \xi r \implies f w - f \xi = (w - \xi)^{\hat{\sim} n} * g w$ 
   $\bigwedge w. w \in \text{ball } \xi r \implies g w \neq 0$ 

```

proof –

```

obtain r where  $r > 0$  and r: ball  $\xi$  r  $\subseteq S$  using assms by (blast elim!: openE)
then have holfb: f holomorphic_on ball  $\xi$  r
  using holf holomorphic_on_subset by blast

```

```

define g where  $g\ w = \text{suminf } (\lambda i. (\text{deriv } \widehat{\widehat{(i+n)}}) f\ \xi / (\text{fact}(i+n)) * (w - \xi)^{\widehat{i}})$  for w
have sumsg:  $(\lambda i. (\text{deriv } \widehat{\widehat{(i+n)}}) f\ \xi / (\text{fact}(i+n)) * (w - \xi)^{\widehat{i}})$  sums g w
and feq:  $f\ w - f\ \xi = (w - \xi)^{\widehat{n}} * g\ w$ 
if w:  $w \in \text{ball } \xi\ r$  for w
proof -
  define powf where  $\text{powf} = (\lambda i. (\text{deriv } \widehat{\widehat{i}}) f\ \xi / (\text{fact } i) * (w - \xi)^{\widehat{i}})$ 
  have [simp]:  $\text{powf } 0 = f\ \xi$ 
  by (simp add: powf_def)
  have sing:  $\{.. < n\} - \{i. \text{powf } i = 0\} = (\text{if } f\ \xi = 0 \text{ then } \{\} \text{ else } \{0\})$ 
  unfolding powf_def using  $\langle 0 < n \rangle$  dfz by (auto simp: dfz; metis funpow_0 not_gr0)
  have powf sums f w
  unfolding powf_def by (rule holomorphic_power_series [OF holfb w])
  moreover have  $(\sum i < n. \text{powf } i) = f\ \xi$ 
  by (subst sum.setdiff_irrelevant [symmetric]; simp add: dfz sing)
  ultimately have fsums:  $(\lambda i. \text{powf } (i+n))$  sums  $(f\ w - f\ \xi)$ 
  using w sums_iff_shift' by metis
  then have *:  $\text{summable } (\lambda i. (w - \xi)^{\widehat{n}} * ((\text{deriv } \widehat{\widehat{(i+n)}}) f\ \xi * (w - \xi)^{\widehat{i}} / \text{fact } (i+n)))$ 
  unfolding powf_def using sums_summable
  by (auto simp: power_add mult_ac)
  have summable  $(\lambda i. (\text{deriv } \widehat{\widehat{(i+n)}}) f\ \xi * (w - \xi)^{\widehat{i}} / \text{fact } (i+n))$ 
  proof (cases w=ξ)
    case False then show ?thesis
    using summable_mult [OF *, of 1 / (w - ξ)^n] by simp
  next
    case True then show ?thesis
    by (auto simp: Power.semiring_1_class.power_0_left intro!: summable_finite [of {0}])
    split: if_split_asm)
  qed
  then show sumsg:  $(\lambda i. (\text{deriv } \widehat{\widehat{(i+n)}}) f\ \xi / (\text{fact}(i+n)) * (w - \xi)^{\widehat{i}})$ 
sums g w
  by (simp add: summable_sums_iff g_def)
  show  $f\ w - f\ \xi = (w - \xi)^{\widehat{n}} * g\ w$ 
  using sums_mult [OF sumsg, of (w - ξ)^n]
  by (intro sums_unique2 [OF fsums]) (auto simp: power_add mult_ac powf_def)
qed
then have holg: g holomorphic_on ball  $\xi\ r$ 
by (meson sumsg power_series_holomorphic)
then have contg: continuous_on (ball  $\xi\ r$ ) g
by (blast intro: holomorphic_on_imp_continuous_on)
have  $g\ \xi \neq 0$ 
using dnz unfolding g_def
by (subst suminf_finite [of {0}]) auto
obtain d where  $0 < d$  and d:  $\bigwedge w. w \in \text{ball } \xi\ d \implies g\ w \neq 0$ 
using  $\langle 0 < r \rangle$  continuous_on_avoid [OF contg  $\langle g\ \xi \neq 0 \rangle$ ]
by (metis centre_in_ball le_cases mem_ball mem_ball_leI)

```

```

show ?thesis
proof
  show  $g$  holomorphic_on ball  $\xi$  (min  $r$   $d$ )
    using  $holg$  by (auto simp: feq holomorphic_on_subset subset_ball  $d$ )
  qed (use <0 <  $r$ > <0 <  $d$ > in <auto simp: feq  $d$ >)
qed

lemma holomorphic_factor_order_of_zero_strong:
  assumes  $holf$ :  $f$  holomorphic_on  $S$  open  $S$   $\xi \in S$   $0 < n$ 
    and (deriv  $\hat{\sim} n$ )  $f \xi \neq 0$ 
    and  $\bigwedge i. [0 < i; i < n] \implies (deriv \hat{\sim} i) f \xi = 0$ 
  obtains  $g$   $r$  where  $0 < r$ 
     $g$  holomorphic_on ball  $\xi$   $r$ 
     $\bigwedge w. w \in \text{ball } \xi \ r \implies f w - f \xi = ((w - \xi) * g w) \hat{\sim} n$ 
     $\bigwedge w. w \in \text{ball } \xi \ r \implies g w \neq 0$ 

proof -
  obtain  $g$   $r$  where  $0 < r$ 
    and  $holg$ :  $g$  holomorphic_on ball  $\xi$   $r$ 
    and  $feq$ :  $\bigwedge w. w \in \text{ball } \xi \ r \implies f w - f \xi = (w - \xi) \hat{\sim} n * g w$ 
    and  $gne$ :  $\bigwedge w. w \in \text{ball } \xi \ r \implies g w \neq 0$ 
  by (auto intro: holomorphic_factor_order_of_zero [OF  $assms$ ])
  have  $con$ : continuous_on (ball  $\xi$   $r$ ) ( $\lambda z. \text{deriv } g \ z / g \ z$ )
  by (rule continuous_intros) (auto simp:  $gne$   $holg$  holomorphic_deriv holomorphic_on_imp_continuous_on)
  have  $cd$ : ( $\lambda z. \text{deriv } g \ z / g \ z$ ) field_differentiable at  $x$  if  $\text{dist } \xi \ x < r$  for  $x$ 
  proof (intro derivative_intros)
    show  $\text{deriv } g$  field_differentiable at  $x$ 
      using that  $holg$  mem_ball by (blast intro: holomorphic_deriv holomorphic_on_imp_differentiable_at)
    show  $g$  field_differentiable at  $x$ 
      by (metis that open_ball at_within_open  $holg$  holomorphic_on_def mem_ball)
    qed (simp add:  $gne$  that)
  obtain  $h$  where  $h$ :  $\bigwedge x. x \in \text{ball } \xi \ r \implies (h \text{ has\_field\_derivative } \text{deriv } g \ x / g \ x)$  (at  $x$ )
    using holomorphic_convex_primitive [of ball  $\xi$   $r$   $\{ \} \lambda z. \text{deriv } g \ z / g \ z$ ]
    by (metis (no_types, lifting) Diff_empty_at_within_interior  $cd$  con convex_ball infinite_imp_nonempty interior_ball mem_ball)
  then have continuous_on (ball  $\xi$   $r$ )  $h$ 
    by (metis open_ball holomorphic_on_imp_continuous_on holomorphic_on_open)
  then have  $con$ : continuous_on (ball  $\xi$   $r$ ) ( $\lambda x. \text{exp } (h \ x) / g \ x$ )
    by (auto intro!: continuous_intros simp add:  $holg$  holomorphic_on_imp_continuous_on  $gne$ )
  have  $0$ :  $\text{dist } \xi \ x < r \implies ((\lambda x. \text{exp } (h \ x) / g \ x) \text{ has\_field\_derivative } 0)$  (at  $x$ )
  for  $x$ 
    apply (rule  $h$  derivative_eq_intros DERIV_deriv_iff_field_differentiable [THEN iffD2] | simp)+
    using  $holg$  by (auto simp: holomorphic_on_imp_differentiable_at  $gne$   $h$ )
  obtain  $c$  where  $c$ :  $\bigwedge x. x \in \text{ball } \xi \ r \implies \text{exp } (h \ x) / g \ x = c$ 

```

```

  by (rule DERIV_zero_connected_constant [of ball  $\xi$   $r$  {}]  $\lambda x. \exp(h x) / g x$ )
(auto simp: con 0)
  have hol: ( $\lambda z. \exp ((Ln (\text{inverse } c) + h z) / \text{of\_nat } n)$ ) holomorphic_on ball  $\xi$   $r$ 
  proof (intro holomorphic_intros holomorphic_on_compose [unfolded o_def,
where  $g = \exp$ ])
    show  $h$  holomorphic_on ball  $\xi$   $r$ 
      using  $h$  holomorphic_on_open by blast
  qed (use  $\langle 0 < n \rangle$  in auto)
  show ?thesis
  proof
    show  $\bigwedge w. w \in \text{ball } \xi \ r \implies f w - f \xi = ((w - \xi) * \exp ((Ln (\text{inverse } c) + h w) / \text{of\_nat } n)) \wedge^n$ 
      using  $\langle 0 < n \rangle$ 
      by (auto simp: feq power_mult_distrib exp_divide_power_eq exp_add gne
simp flip: c)
    qed (use hol  $\langle 0 < r \rangle$  in auto)
  qed

```

lemma

```

  fixes  $k :: 'a::wellorder$ 
  assumes  $a\_def: a == \text{LEAST } x. P x$  and  $P: P k$ 
  shows  $\text{def\_LeastI}: P a$  and  $\text{def\_Least\_le}: a \leq k$ 
  unfolding  $a\_def$ 
  by (rule LeastI Least_le; rule P)+

```

lemma holomorphic_factor_zero_nonconstant:

```

  assumes  $\text{hol}f: f$  holomorphic_on  $S$  and  $S: \text{open } S$  connected  $S$ 
  and  $\xi \in S$   $f \xi = 0$ 
  and  $\text{nonconst}: \neg f \text{ constant\_on } S$ 
  obtains  $g$   $r$   $n$ 
  where  $0 < n$   $0 < r$   $\text{ball } \xi \ r \subseteq S$ 
  and  $g$  holomorphic_on ball  $\xi$   $r$ 
  and  $\bigwedge w. w \in \text{ball } \xi \ r \implies f w = (w - \xi) \wedge^n * g w$ 
  and  $\bigwedge w. w \in \text{ball } \xi \ r \implies g w \neq 0$ 
  proof (cases  $\forall n > 0. (\text{deriv } \wedge^n) f \xi = 0$ )
  case True then show ?thesis
    using holomorphic_fun_eq_const_on_connected [OF  $\text{hol}f$   $S$   $\langle \xi \in S \rangle$ ] non-const by (simp add: constant_on_def)
  next
  case False
    then obtain  $n0$  where  $n0 > 0$  and  $n0: (\text{deriv } \wedge^{n0}) f \xi \neq 0$  by blast
    obtain  $r0$  where  $r0 > 0$   $\text{ball } \xi \ r0 \subseteq S$  using  $S$  openE  $\langle \xi \in S \rangle$  by auto
    define  $n$  where  $n \equiv \text{LEAST } n. (\text{deriv } \wedge^n) f \xi \neq 0$ 
    have  $n\_ne: (\text{deriv } \wedge^n) f \xi \neq 0$ 
      by (rule  $\text{def\_LeastI}$  [OF  $n\_def$ ]) (rule  $n0$ )
    then have  $0 < n$  using  $\langle f \xi = 0 \rangle$ 
      using funpow_0 by fastforce
    have  $n\_min: \bigwedge k. k < n \implies (\text{deriv } \wedge^k) f \xi = 0$ 

```

```

  using def_Least_le [OF n_def] not_le by blast
then obtain g r1
  where g: 0 < r1 g holomorphic_on ball ξ r1
        and geq:  $\bigwedge w. w \in \text{ball } \xi \ r1 \implies f w = (w - \xi) ^ n * g w$ 
        and g0:  $\bigwedge w. w \in \text{ball } \xi \ r1 \implies g w \neq 0$ 
  by (auto intro: holomorphic_factor_order_of_zero [OF holf ⟨open S⟩ ⟨ξ ∈ S⟩
⟨n > 0⟩ n_ne] simp: ⟨f ξ = 0⟩)
  show ?thesis
proof
  show g holomorphic_on ball ξ (min r0 r1)
    using g by auto
  show  $\bigwedge w. w \in \text{ball } \xi \ (\text{min } r0 \ r1) \implies f w = (w - \xi) ^ n * g w$ 
    by (simp add: geq)
qed (use ⟨0 < n⟩ ⟨0 < r0⟩ ⟨0 < r1⟩ ⟨ball ξ r0 ⊆ S⟩ g0 in auto)
qed

```

lemma holomorphic_lower_bound_difference:

```

  assumes holf: f holomorphic_on S and S: open S connected S
        and ξ ∈ S and φ ∈ S
        and fne: f φ ≠ f ξ
  obtains k n r
    where 0 < k 0 < r
          ball ξ r ⊆ S
           $\bigwedge w. w \in \text{ball } \xi \ r \implies k * \text{norm}(w - \xi) ^ n \leq \text{norm}(f w - f \xi)$ 
proof -
  define n where n = (LEAST n. 0 < n ∧ (deriv  $\hat{\hat{}}$  n) f ξ ≠ 0)
  obtain n0 where 0 < n0 and n0: (deriv  $\hat{\hat{}}$  n0) f ξ ≠ 0
  using fne holomorphic_fun_eq_const_on_connected [OF holf S] ⟨ξ ∈ S⟩ ⟨φ ∈ S⟩ by blast
  then have 0 < n and n_ne: (deriv  $\hat{\hat{}}$  n) f ξ ≠ 0
  unfolding n_def by (metis (mono_tags, lifting) LeastI)+
  have n_min:  $\bigwedge k. \llbracket 0 < k; k < n \rrbracket \implies (\text{deriv } \hat{\hat{}} k) f \xi = 0$ 
  unfolding n_def by (blast dest: not_less_Least)
  then obtain g r
    where 0 < r and holf: g holomorphic_on ball ξ r
          and fne:  $\bigwedge w. w \in \text{ball } \xi \ r \implies f w - f \xi = (w - \xi) ^ n * g w$ 
          and gnz:  $\bigwedge w. w \in \text{ball } \xi \ r \implies g w \neq 0$ 
    by (auto intro: holomorphic_factor_order_of_zero [OF holf ⟨open S⟩ ⟨ξ ∈ S⟩
⟨n > 0⟩ n_ne])
  obtain e where e>0 and e: ball ξ e ⊆ S using assms by (blast elim!: openE)
  then have holfb: f holomorphic_on ball ξ e
    using holf holomorphic_on_subset by blast
  define d where d = (min e r) / 2
  have 0 < d using ⟨0 < r⟩ ⟨0 < e⟩ by (simp add: d_def)
  have d < r
    using ⟨0 < r⟩ by (auto simp: d_def)
  then have cbb: cball ξ d ⊆ ball ξ r
    by (auto simp: cball_subset_ball_iff)

```

```

then have g holomorphic_on cball ξ d
  by (rule holomorphic_on_subset [OF holg])
then have closed (g ` cball ξ d)
  by (simp add: compact_imp_closed compact_continuous_image holomorphic_on_imp_continuous_on)
moreover have g ` cball ξ d ≠ {}
  using ⟨0 < d⟩ by auto
ultimately obtain x where x: x ∈ g ` cball ξ d and ∧y. y ∈ g ` cball ξ d ⇒
dist 0 x ≤ dist 0 y
  by (rule distance_attains_inf) blast
then have leg: ∧w. w ∈ cball ξ d ⇒ norm x ≤ norm (g w)
  by auto
have ball ξ d ⊆ cball ξ d by auto
also have ... ⊆ ball ξ e using ⟨0 < d⟩ d_def by auto
also have ... ⊆ S by (rule e)
finally have dS: ball ξ d ⊆ S .
have x ≠ 0 using gnz x ⟨d < r⟩ by auto
show thesis
proof
  show ∧w. w ∈ ball ξ d ⇒ cmod x * cmod (w - ξ) ^ n ≤ cmod (f w - f ξ)
    using ⟨d < r⟩ leg by (auto simp: fne norm_mult norm_power algebra_simps
mult_right_mono)
qed (use dS ⟨x ≠ 0⟩ ⟨d > 0⟩ in auto)
qed

```

lemma

```

assumes holg: f holomorphic_on (S - {ξ}) and ξ: ξ ∈ interior S
shows holomorphic_on_extend_lim:
  (∃g. g holomorphic_on S ∧ (∀z ∈ S - {ξ}. g z = f z)) ↔
  ((λz. (z - ξ) * f z) → 0) (at ξ)
  (is ?P = ?Q)
and holomorphic_on_extend_bounded:
  (∃g. g holomorphic_on S ∧ (∀z ∈ S - {ξ}. g z = f z)) ↔
  (∃B. eventually (λz. norm(f z) ≤ B) (at ξ))
  (is ?P = ?R)

```

proof -

```

obtain δ where 0 < δ and δ: ball ξ δ ⊆ S

```

```

  using ξ mem_interior by blast

```

```

have ?R if holg: g holomorphic_on S and gf: ∧z. z ∈ S - {ξ} ⇒ g z = f z

```

for g

proof -

```

  have §: cmod (f x) ≤ cmod (g ξ) + 1 if x ≠ ξ dist x ξ < δ dist (g x) (g ξ) <
1 for x

```

proof -

```

  have x ∈ S

```

```

    by (metis δ dist_commute mem_ball subsetD that(2))

```

```

  with that gf [of x] show ?thesis

```

```

    using norm_triangle_ineq2 [of f x g ξ] dist_complex_def by auto

```

qed

```

then have *: ∀F z in at ξ. dist (g z) (g ξ) < 1 → cmod (f z) ≤ cmod (g ξ)

```

```

+ 1
  using ‹0 < δ› eventually_at by blast
  have continuous_on (interior S) g
  by (meson continuous_on_subset holg holomorphic_on_imp_continuous_on
interior_subset)
  then have  $\bigwedge x. x \in \text{interior } S \implies (g \longrightarrow g x) \text{ (at } x)$ 
  using continuous_on_interior continuous_within holg holomorphic_on_imp_continuous_on
by blast
  then have  $(g \longrightarrow g \xi) \text{ (at } \xi)$ 
  by (simp add:  $\xi$ )
  then show ?thesis
  apply (rule_tac x=norm(g  $\xi$ ) + 1 in exI)
  apply (rule eventually_mp [OF * tendstoD [where e=1]], auto)
  done
qed
moreover have ?Q if  $\forall_F z \text{ in at } \xi. \text{ cmod } (f z) \leq B \text{ for } B$ 
  by (rule lim_null_mult_right_bounded [OF _ that]) (simp add: LIM_zero)
moreover have ?P if  $(\lambda z. (z - \xi) * f z) -\xi \rightarrow 0$ 
proof -
  define h where [abs_def]:  $h z = (z - \xi)^{\wedge 2} * f z$  for z
  have h0: (h has_field_derivative 0) (at  $\xi$ )
  apply (simp add: h_def has_field_derivative_iff)
  apply (auto simp: field_split_simps power2_eq_square Lim_transform_within
[OF that, of 1])
  done
  have holh: h holomorphic_on S
  proof (simp add: holomorphic_on_def, clarify)
    fix z assume z ∈ S
    show h field_differentiable at z within S
    proof (cases z =  $\xi$ )
      case True then show ?thesis
      using field_differentiable_at_within field_differentiable_def h0 by blast
    next
      case False
      then have f field_differentiable at z within S
      using holomorphic_onD [OF holf, of z] ‹z ∈ S›
      unfolding field_differentiable_def has_field_derivative_iff
      by (force intro: exI [where x=dist  $\xi$ ] elim: Lim_transform_within_set
[unfolded eventually_at])
      then show ?thesis
      by (simp add: h_def power2_eq_square derivative_intros)
    qed
  qed
  define g where [abs_def]:  $g z = (\text{if } z = \xi \text{ then deriv } h \ \xi \text{ else } (h z - h \xi) / (z - \xi))$  for z
  have holg: g holomorphic_on S
  unfolding g_def by (rule pole_lemma [OF holh  $\xi$ ])
  have  $\S: \forall z \in S - \{\xi\}. (g z - g \xi) / (z - \xi) = f z$ 
  using h0 by (auto simp: g_def power2_eq_square divide_simps DERIV_imp_deriv

```

```

h_def)
  show ?thesis
    apply (intro exI conjI)
    apply (rule pole_lemma [OF holg ξ])
    apply (simp add: §)
    done
qed
ultimately show ?P = ?Q and ?P = ?R
  by meson+
qed

lemma pole_at_infinity:
  assumes holg: f holomorphic_on UNIV and lim: ((inverse o f) → l) at_infinity
  obtains a n where  $\bigwedge z. f z = (\sum_{i \leq n} a_i * z^i)$ 
proof (cases l = 0)
  case False
  show thesis
  proof
    show  $f z = (\sum_{i \leq 0} \text{inverse } l * z^i)$  for z
      using False tendsto_inverse [OF lim] by (simp add: Liouville_weak [OF
holg])
    qed
  next
  case True
  then have [simp]: l = 0 .
  show ?thesis
  proof (cases  $\exists r. 0 < r \wedge (\forall z \in \text{ball } 0 \ r - \{0\}. f(\text{inverse } z) \neq 0)$ )
    case True
    then obtain r where  $0 < r$  and r:  $\bigwedge z. z \in \text{ball } 0 \ r - \{0\} \implies f(\text{inverse } z) \neq 0$ 
      by auto
    have 1:  $\text{inverse} \circ f \circ \text{inverse}$  holomorphic_on ball 0 r - {0}
    by (rule holomorphic_on_compose holomorphic_intros holomorphic_on_subset
[OF holg] | force simp: r)+
    have 2:  $0 \in \text{interior} (\text{ball } 0 \ r)$ 
      using  $\langle 0 < r \rangle$  by simp
    have  $\exists B. 0 < B \wedge \text{eventually } (\lambda z. \text{cmod } ((\text{inverse} \circ f \circ \text{inverse}) z) \leq B)$  (at
0)
      apply (rule exI [where x=1])
      using tendstoD [OF lim [unfolded lim_at_infinity_0] zero_less_one]
      by (simp add: eventually_mono)
    with holomorphic_on_extend_bounded [OF 1 2]
    obtain g where holg: g holomorphic_on ball 0 r
      and geq:  $\bigwedge z. z \in \text{ball } 0 \ r - \{0\} \implies g z = (\text{inverse} \circ f \circ \text{inverse}) z$ 
      by meson
    have ifi0:  $(\text{inverse} \circ f \circ \text{inverse}) - 0 \rightarrow 0$ 
      using  $\langle l = 0 \rangle$  lim lim_at_infinity_0 by blast
    have g2g0:  $g - 0 \rightarrow g \ 0$ 
      using  $\langle 0 < r \rangle$  centre_in_ball continuous_at continuous_on_eq continuous_at

```



```

holg
  by (blast intro: holomorphic_on_imp_continuous_on)
have g2g1:  $g - 0 \rightarrow 0$ 
  apply (rule Lim_transform_within_open [OF if0 open_ball [of 0 r]])
  using  $\langle 0 < r \rangle$  by (auto simp: geq)
have [simp]:  $g 0 = 0$ 
  by (rule tendsto_unique [OF _ g2g0 g2g1]) simp
have ball 0 r -  $\{0::\text{complex}\} \neq \{\}$ 
using  $\langle 0 < r \rangle$  by (metis 2 Diff_iff insert_Diff interior_ball interior_singleton)
then obtain  $w::\text{complex}$  where  $w \neq 0$  and  $w: \text{norm } w < r$  by force
then have  $g w \neq 0$  by (simp add: geq r)
obtain B n e where  $0 < B 0 < e e \leq r$ 
  and leg:  $\bigwedge w. \text{norm } w < e \implies B * \text{cmod } w ^ n \leq \text{cmod } (g w)$ 
proof (rule holomorphic_lower_bound_difference [OF holg open_ball con-
nected_ball])
  show  $g w \neq g 0$ 
    by (simp add:  $\langle g w \neq 0 \rangle$ )
  show  $w \in \text{ball } 0 r$ 
    using mem_ball_0 w by blast
qed (use  $\langle 0 < r \rangle$  in  $\langle \text{auto simp: ball_subset_ball_iff} \rangle$ )
have  $\S: \text{cmod } (f z) \leq \text{cmod } z ^ n / B$  if  $2/e \leq \text{cmod } z$  for z
proof -
  have ize:  $inverse z \in \text{ball } 0 e - \{0\}$  using that  $\langle 0 < e \rangle$ 
    by (auto simp: norm_divide field_split_simps algebra_simps)
  then have [simp]:  $z \neq 0$  and izr:  $inverse z \in \text{ball } 0 r - \{0\}$  using  $\langle e \leq$ 
r)
    by auto
  then have [simp]:  $f z \neq 0$ 
    using r [of inverse z] by simp
  have [simp]:  $f z = inverse (g (inverse z))$ 
    using izr geq [of inverse z] by simp
  show ?thesis using ize leg [of inverse z]  $\langle 0 < B \rangle$   $\langle 0 < e \rangle$ 
    by (simp add: field_split_simps norm_divide algebra_simps)
qed
show thesis
proof
  show  $f z = (\sum_{i \leq n}. (\text{deriv } ^ i) f 0 / \text{fact } i * z ^ i)$  for z
    using  $\S$  by (rule_tac A = 2/e and B = 1/B in Liouville_polynomial
[OF holf], simp)
  qed
next
case False
then have fi0:  $\bigwedge r. r > 0 \implies \exists z \in \text{ball } 0 r - \{0\}. f (inverse z) = 0$ 
  by simp
  have fz0:  $f z = 0$  if  $0 < r$  and lt1:  $\bigwedge x. x \neq 0 \implies \text{cmod } x < r \implies inverse$ 
(cmod (f (inverse x))) < 1
    for z r
  proof -
  have f0: (f  $\longrightarrow$  0) at_infinity

```

```

proof –
  have  $DIM\_complex[intro]: 2 \leq DIM(complex)$  — should not be necessary!
  by simp
  have  $f(inverse\ x) \neq 0 \implies x \neq 0 \implies cmod\ x < r \implies 1 < cmod(f(inverse\ x))$ 
for  $x$ 
    using lt1[of x] by (auto simp: field_simps)
    then have  $** : cmod(f(inverse\ x)) \leq 1 \implies x \neq 0 \implies cmod\ x < r \implies f(inverse\ x) = 0$ 
for  $x$ 
    by force
    then have  $*$ :  $(f \circ inverse)^{-1}(ball\ 0\ r - \{0\}) \subseteq \{0\} \cup -ball\ 0\ 1$ 
    by force
    have continuous_on  $(inverse^{-1}(ball\ 0\ r - \{0\}))\ f$ 
    using continuous_on_subset holf holomorphic_on_imp_continuous_on
by blast
    then have connected  $((f \circ inverse)^{-1}(ball\ 0\ r - \{0\}))$ 
    using connected_punctured_ball
    by (intro connected_continuous_image continuous_intros; force)
    then have  $\{0\} \cap (f \circ inverse)^{-1}(ball\ 0\ r - \{0\}) = \{0\} \vee -ball\ 0\ 1 \cap (f \circ inverse)^{-1}(ball\ 0\ r - \{0\}) = \{0\}$ 
    by (rule connected_closedD) (use * in auto)
    then have  $f(inverse\ w) = 0$  if  $w \neq 0$   $cmod\ w < r$  for  $w$ 
    using  $**[of\ w]\ fi0\ \langle 0 < r \rangle$  that by force
    then show ?thesis
    unfolding lim_at_infinity_0
    using eventually_at  $\langle r > 0 \rangle$  by (force simp add: intro: tendsto_eventually)
qed
obtain  $w$  where  $w \in ball\ 0\ r - \{0\}$  and  $f(inverse\ w) = 0$ 
using False  $\langle 0 < r \rangle$  by blast
then show ?thesis
by (auto simp: f0 Liouville_weak [OF holf, of 0])
qed
show thesis
proof
  show  $\bigwedge z. f\ z = (\sum_{i \leq 0}. 0 * z^i)$ 
  using lim
  apply (simp add: lim_at_infinity_0 Lim_at_dist_norm norm_inverse)
  using fz0 zero_less_one by blast
qed
qed
qed

```

5.5 Entire proper functions are precisely the non-trivial polynomials

lemma *proper_map_polyfun*:

fixes $c :: nat \Rightarrow 'a :: \{real_normed_div_algebra, heine_borel\}$

assumes *closed S and compact K and* $c\ i \neq 0\ 1 \leq i \leq n$

shows *compact* $(S \cap \{z. (\sum_{i \leq n}. c\ i * z^i) \in K\})$

proof –

```

obtain  $B$  where  $B > 0$  and  $B: \bigwedge x. x \in K \implies \text{norm } x \leq B$ 
  by (metis compact_imp_bounded ⟨compact  $K$ ⟩ bounded_pos)
have *:  $\text{norm } x \leq b$ 
  if  $\bigwedge x. b \leq \text{norm } x \implies B + 1 \leq \text{norm } (\sum_{i \leq n}. c_i * x^i)$ 
     $(\sum_{i \leq n}. c_i * x^i) \in K$  for  $b$   $x$ 
proof -
  have  $\text{norm } (\sum_{i \leq n}. c_i * x^i) \leq B$ 
    using  $B$  that by blast
  moreover have  $\neg B + 1 \leq B$ 
    by simp
  ultimately show  $\text{norm } x \leq b$ 
    using that by (metis (no_types) less_eq_real_def not_less order_trans)
qed
have bounded  $\{z. (\sum_{i \leq n}. c_i * z^i) \in K\}$ 
  using Limits.polyfun_extremal [where  $c=c$  and  $B=B+1$ ,  $OF$   $c$ ]
  by (auto simp: bounded_pos eventually_at_infinity_pos *)
moreover have closed  $(\lambda z. (\sum_{i \leq n}. c_i * z^i))^{-1} K$ 
  using ⟨compact  $K$ ⟩ compact_eq_bounded_closed
  by (intro allI continuous_closed_vimage continuous_intros; force)
ultimately show ?thesis
  using closed_Int_compact [ $OF$  ⟨closed  $S$ ⟩] compact_eq_bounded_closed
  by (auto simp add: vimage_def)
qed

lemma proper_map_polyfun_univ:
  fixes  $c :: \text{nat} \Rightarrow 'a::\{\text{real\_normed\_div\_algebra, heine\_borel}\}$ 
  assumes compact  $K$   $c_i \neq 0$   $1 \leq i \leq n$ 
  shows compact  $\{z. (\sum_{i \leq n}. c_i * z^i) \in K\}$ 
  using proper_map_polyfun [of  $UNIV$   $K$   $c$   $i$   $n$ ] assms by simp

lemma proper_map_polyfun_eq:
  assumes  $f$  holomorphic_on  $UNIV$ 
  shows  $(\forall k. \text{compact } k \longrightarrow \text{compact } \{z. f z \in k\}) \longleftrightarrow$ 
     $(\exists c\ n. 0 < n \wedge (c\ n \neq 0) \wedge f = (\lambda z. \sum_{i \leq n}. c_i * z^i))$ 
    (is ?lhs = ?rhs)

proof
  assume compf [rule_format]: ?lhs
  have 2:  $\exists k. 0 < k \wedge a\ k \neq 0 \wedge f = (\lambda z. \sum_{i \leq k}. a_i * z^i)$ 
    if  $\bigwedge z. f z = (\sum_{i \leq n}. a_i * z^i)$  for  $a$   $n$ 
  proof (cases  $\forall i \leq n. 0 < i \longrightarrow a_i = 0$ )
    case True
      then have [simp]:  $\bigwedge z. f z = a\ 0$ 
        by (simp add: that sum.atMost_shift)
      have False using compf [of  $\{a\ 0\}$ ] by simp
      then show ?thesis ..
    next
      case False
      then obtain  $k$  where  $0 < k$   $k \leq n$   $a\ k \neq 0$  by force
      define  $m$  where  $m = (\text{GREATEST } k. k \leq n \wedge a\ k \neq 0)$ 

```

```

have m:  $m \leq n \wedge a \ m \neq 0$ 
  unfolding m_def
  using GreatestI_nat [where b = n] k by (metis (mono_tags, lifting))
have [simp]:  $a \ i = 0$  if  $m < i \ i \leq n$  for i
  using Greatest_le_nat [where b = n and  $P = \lambda k. k \leq n \wedge a \ k \neq 0$ ]
  using m_def not_le that by auto
have  $k \leq m$ 
  unfolding m_def
  using Greatest_le_nat [where b = n] k by (metis (mono_tags, lifting))
with k m show ?thesis
by (rule_tac x=m in exI) (auto simp: that comm_monoid_add_class.sum_mono_neutral_right)
qed
have §:  $((inverse \circ f) \longrightarrow 0)$  at_infinity
proof (rule Lim_at_infinityI)
fix e::real assume  $0 < e$ 
with compf [of cball 0 (inverse e)]
show  $\exists B. \forall x. B \leq cmod \ x \longrightarrow dist \ ((inverse \circ f) \ x) \ 0 \leq e$ 
  apply simp
  apply (clarsimp simp add: compact_eq_bounded_closed bounded_pos norm_inverse)
  by (metis (no_types, opaque_lifting) inverse_inverse_eq le_less_trans less_eq_real_def
less_imp_inverse_less linordered_field_no_ub not_less)
qed
then obtain a n where  $\bigwedge z. f \ z = (\sum_{i \leq n}. a \ i * z^i)$ 
  using assms pole_at_infinity by blast
with § 2 show ?rhs by blast
next
assume ?rhs
then obtain c n where  $0 < n \ c \ n \neq 0 \ f = (\lambda z. \sum_{i \leq n}. c \ i * z^i)$  by blast
then have compact {z. f z ∈ k} if compact k for k
  by (auto intro: proper_map_polyfun_univ [OF that])
then show ?lhs by blast
qed

```

5.6 Relating invertibility and nonvanishing of derivative

lemma has_complex_derivative_locally_injective:

```

assumes holf: f holomorphic_on S
  and S:  $\xi \in S$  open S
  and dnz: deriv f  $\xi \neq 0$ 
obtains r where  $r > 0$  ball  $\xi \ r \subseteq S$  inj_on f (ball  $\xi \ r$ )
proof -
have *:  $\exists d > 0. \forall x. dist \ \xi \ x < d \longrightarrow onorm \ (\lambda v. deriv \ f \ x * v - deriv \ f \ \xi * v) < e$  if  $e > 0$  for e
proof -
have contdf: continuous_on S (deriv f)
  by (simp add: holf holomorphic_deriv holomorphic_on_imp_continuous_on (open S))
obtain  $\delta > 0$  and  $\delta$ :  $\bigwedge x. \llbracket x \in S; dist \ x \ \xi \leq \delta \rrbracket \implies cmod \ (deriv \ f \ x - deriv \ f \ \xi) \leq e/2$ 

```

```

    using continuous_onE [OF contdf ⟨ξ ∈ S⟩, of e/2] ⟨0 < e⟩
    by (metis dist_complex_def half_gt_zero less_imp_le)
    have §:  $\bigwedge \zeta z. [\zeta \in S; \text{dist } \xi \zeta < \delta] \implies \text{cmod } (\text{deriv } f \zeta - \text{deriv } f \xi) * \text{cmod } z \leq e/2 * \text{cmod } z$ 
    by (intro mult_right_mono [OF δ]) (auto simp: dist_commute)
    obtain ε where ε > 0 ball ξ ε ⊆ S
    by (metis openE [OF ⟨open S⟩ ⟨ξ ∈ S⟩])
    with ⟨δ > 0⟩ have  $\exists \delta > 0. \forall x. \text{dist } \xi x < \delta \longrightarrow \text{onorm } (\lambda v. \text{deriv } f x * v - \text{deriv } f \xi * v) \leq e/2$ 
    using §
    apply (rule_tac x=min δ ε in exI)
    apply (intro conjI allI impI Operator_Norm.onorm_le)
    apply (force simp: mult_right_mono norm_mult [symmetric] left_diff_distrib δ)+
    done
    with ⟨e > 0⟩ show ?thesis by force
  qed
  have inj ((* ) (deriv f ξ))
  using dnz by simp
  then obtain g' where g': linear g' g' o ((* ) (deriv f ξ)) = id
  using linear_injective_left_inverse [of ((* ) (deriv f ξ))] by auto
  show ?thesis
  apply (rule has_derivative_locally_injective [OF S, where f=f and f' = λz
h. deriv f z * h and g' = g'])
  using g' *
  apply (simp_all add: linear_conv_bounded_linear that)
  using ⟨open S⟩ has_field_derivative_imp_has_derivative holf_holomorphic_derivI
  by blast
  qed

lemma has_complex_derivative_locally_invertible:
  assumes holf: f holomorphic_on S
  and S: ξ ∈ S open S
  and dnz: deriv f ξ ≠ 0
  obtains r where r > 0 ball ξ r ⊆ S open (f ' (ball ξ r)) inj_on f (ball ξ r)
  proof -
    obtain r where r > 0 ball ξ r ⊆ S inj_on f (ball ξ r)
    by (blast intro: that has_complex_derivative_locally_injective [OF assms])
    then have ξ: ξ ∈ ball ξ r by simp
    then have nc: ¬ f constant_on ball ξ r
    using ⟨inj_on f (ball ξ r)⟩ injective_not_constant by fastforce
    have holf': f holomorphic_on ball ξ r
    using ⟨ball ξ r ⊆ S⟩ holf_holomorphic_on_subset by blast
    have open (f ' ball ξ r)
    by (simp add: ⟨inj_on f (ball ξ r)⟩ holf' open_mapping_thm3)
    then show ?thesis
    using ⟨0 < r⟩ ⟨ball ξ r ⊆ S⟩ ⟨inj_on f (ball ξ r)⟩ that by blast
  qed

```

```

lemma holomorphic_injective_imp_regular:
  assumes holf: f holomorphic_on S
    and open S and injf: inj_on f S
    and  $\xi \in S$ 
  shows deriv f  $\xi \neq 0$ 
proof -
  obtain r where r>0 and r: ball  $\xi$  r  $\subseteq S$  using assms by (blast elim!: openE)
  have holf': f holomorphic_on ball  $\xi$  r
    using  $\langle$ ball  $\xi$  r  $\subseteq S\rangle$  holf holomorphic_on_subset by blast
  show ?thesis
  proof (cases  $\forall n>0. (\text{deriv } \sim n) f \xi = 0$ )
    case True
      have fcon: f w = f  $\xi$  if w  $\in$  ball  $\xi$  r for w
        by (meson open_ball True  $\langle 0 < r\rangle$  centre_in_ball connected_ball holf'
          holomorphic_fun_eq_const_on_connected that)
      have False
        using fcon [of  $\xi + r/2$ ]  $\langle 0 < r\rangle$  r injf unfolding inj_on_def
        by (metis  $\langle \xi \in S\rangle$  contra_subsetD dist_commute fcon mem_ball perfect_choose_dist)
      then show ?thesis ..
    next
      case False
      then obtain n0 where n0: n0 > 0  $\wedge$  (deriv  $\sim$  n0) f  $\xi \neq 0$  by blast
      define n where [abs_def]: n = (LEAST n. n > 0  $\wedge$  (deriv  $\sim$  n) f  $\xi \neq 0$ )
      have n_ne: n > 0 (deriv  $\sim$  n) f  $\xi \neq 0$ 
        using def_LeastI [OF n_def n0] by auto
      have n_min:  $\bigwedge k. 0 < k \implies k < n \implies (\text{deriv } \sim k) f \xi = 0$ 
        using def_Least_le [OF n_def] not_le by auto
      obtain g  $\delta$  where 0 <  $\delta$ 
        and holg: g holomorphic_on ball  $\xi$   $\delta$ 
        and fd:  $\bigwedge w. w \in \text{ball } \xi \delta \implies f w - f \xi = ((w - \xi) * g w) \wedge n$ 
        and gnz:  $\bigwedge w. w \in \text{ball } \xi \delta \implies g w \neq 0$ 
        by (blast intro: n_min holomorphic_factor_order_of_zero_strong [OF holf
           $\langle$ open S $\rangle$   $\langle \xi \in S\rangle$  n_ne])
      show ?thesis
      proof (cases n=1)
        case True
          with n_ne show ?thesis by auto
        next
          case False
          have g holomorphic_on ball  $\xi$  (min r  $\delta$ )
            using holg by (simp add: holomorphic_on_subset subset_ball)
          then have holgw:  $(\lambda w. (w - \xi) * g w)$  holomorphic_on ball  $\xi$  (min r  $\delta$ )
            by (intro holomorphic_intros)
          have gd:  $\bigwedge w. \text{dist } \xi w < \delta \implies (g \text{ has\_field\_derivative } \text{deriv } g w)$  (at w)
            using holg
            by (simp add: DERIV_deriv_iff_field_differentiable holomorphic_on_def
              at_within_open_NO_MATCH)
          have *:  $\bigwedge w. w \in \text{ball } \xi$  (min r  $\delta$ )
             $\implies ((\lambda w. (w - \xi) * g w) \text{ has\_field\_derivative } ((w - \xi) * \text{deriv } g w + g$ 

```

```

w))
  (at w)
  by (rule gd_derivative_eq_intros | simp)+
  have [simp]: deriv ( $\lambda w. (w - \xi) * g w$ )  $\xi \neq 0$ 
  using * [of  $\xi$ ]  $\langle 0 < \delta \rangle \langle 0 < r \rangle$  by (simp add: DERIV_imp_deriv gnz)
  obtain T where  $\xi \in T$  open T and Tsb:  $T \subseteq \text{ball } \xi (\min r \delta)$  and oimT:
  open (( $\lambda w. (w - \xi) * g w$ ) ' T)
  using  $\langle 0 < r \rangle \langle 0 < \delta \rangle$  has_complex_derivative_locally_invertible [OF
  holgw, of  $\xi$ ]
  by force
  define U where  $U = (\lambda w. (w - \xi) * g w)$  ' T
  have open U by (metis oimT U_def)
  moreover have  $0 \in U$ 
  using  $\langle \xi \in T \rangle$  by (auto simp: U_def intro: image_eqI [where  $x = \xi$ ])
  ultimately obtain  $\varepsilon$  where  $\varepsilon > 0$  and  $\varepsilon$ :  $\text{cball } 0 \varepsilon \subseteq U$ 
  using  $\langle \text{open } U \rangle$  open_contains_cball by blast
  then have  $\varepsilon * \exp(2 * \text{of\_real } \pi * i * (0/n)) \in \text{cball } 0 \varepsilon$ 
   $\varepsilon * \exp(2 * \text{of\_real } \pi * i * (1/n)) \in \text{cball } 0 \varepsilon$ 
  by (auto simp: norm_mult)
  with  $\varepsilon$  have  $\varepsilon * \exp(2 * \text{of\_real } \pi * i * (0/n)) \in U$ 
   $\varepsilon * \exp(2 * \text{of\_real } \pi * i * (1/n)) \in U$  by blast+
  then obtain  $y0 y1$  where  $y0 \in T$  and  $y0$ :  $(y0 - \xi) * g y0 = \varepsilon * \exp(2 * \text{of\_real } \pi * i * (0/n))$ 
  and  $y1 \in T$  and  $y1$ :  $(y1 - \xi) * g y1 = \varepsilon * \exp(2 * \text{of\_real } \pi * i * (1/n))$ 
  by (auto simp: U_def)
  then have  $y0 \in \text{ball } \xi \delta$   $y1 \in \text{ball } \xi \delta$  using Tsb by auto
  then have  $f y0 - f \xi = ((y0 - \xi) * g y0) ^ n$   $f y1 - f \xi = ((y1 - \xi) * g$ 
 $y1) ^ n$ 
  using fd by blast+
  moreover have  $y0 \neq y1$ 
  using  $y0 y1 \langle \varepsilon > 0 \rangle$  complex_root_unity_eq_1 [of n 1]  $\langle n > 0 \rangle$  False by
  auto
  moreover have  $T \subseteq S$ 
  by (meson Tsb min.cobounded1 order_trans r subset_ball)
  ultimately have False
  using inj_onD [OF injf, of  $y0 y1$ ]  $\langle y0 \in T \rangle \langle y1 \in T \rangle$ 
  using complex_root_unity [of n 1]
  apply (simp add:  $y0 y1$  power_mult_distrib)
  apply (force simp: algebra_simps)
  done
  then show ?thesis ..
qed
qed
qed

```

5.6.1 Hence a nice clean inverse function theorem

lemma has_field_derivative_inverse_strong:

```

fixes f :: 'a::{euclidean_space,real_normed_field}  $\Rightarrow$  'a
shows  $\llbracket$ DERIV f x :> f'; f'  $\neq$  0; open S; x  $\in$  S; continuous_on S f;
 $\bigwedge$ z. z  $\in$  S  $\implies$  g (f z) = z $\rrbracket$ 
 $\implies$  DERIV g (f x) :> inverse (f')
unfolding has_field_derivative_def
by (rule has_derivative_inverse_strong [of S x f g]) auto

lemma has_field_derivative_inverse_strong_x:
fixes f :: 'a::{euclidean_space,real_normed_field}  $\Rightarrow$  'a
shows  $\llbracket$ DERIV f (g y) :> f'; f'  $\neq$  0; open S; continuous_on S f; g y  $\in$  S; f(g
y) = y;
 $\bigwedge$ z. z  $\in$  S  $\implies$  g (f z) = z $\rrbracket$ 
 $\implies$  DERIV g y :> inverse (f')
unfolding has_field_derivative_def
by (rule has_derivative_inverse_strong_x [of S g y f]) auto

proposition holomorphic_has_inverse:
assumes holf: f holomorphic_on S
and open S and injf: inj_on f S
obtains g where g holomorphic_on (f ' S)
 $\bigwedge$ z. z  $\in$  S  $\implies$  deriv f z * deriv g (f z) = 1
 $\bigwedge$ z. z  $\in$  S  $\implies$  g(f z) = z

proof –
have ofs: open (f ' S)
by (rule open_mapping_thm3 [OF assms])
have contf: continuous_on S f
by (simp add: holf holomorphic_on_imp_continuous_on)
have *: (the_inv_into S f has_field_derivative inverse (deriv f z)) (at (f z)) if
z  $\in$  S for z
proof –
have 1: (f has_field_derivative deriv f z) (at z)
using DERIV_deriv_iff_field_differentiable  $\langle$ z  $\in$  S $\rangle$   $\langle$ open S $\rangle$  holf holomor-
phic_on_imp_differentiable_at
by blast
have 2: deriv f z  $\neq$  0
using  $\langle$ z  $\in$  S $\rangle$   $\langle$ open S $\rangle$  holf holomorphic_injective_imp_regular injf by blast
show ?thesis
proof (rule has_field_derivative_inverse_strong [OF 1 2  $\langle$ open S $\rangle$   $\langle$ z  $\in$  S $\rangle$ ])
show continuous_on S f
by (simp add: holf holomorphic_on_imp_continuous_on)
show  $\bigwedge$ z. z  $\in$  S  $\implies$  the_inv_into S f (f z) = z
by (simp add: injf the_inv_into_f_f)
qed
qed
show ?thesis
proof
show the_inv_into S f holomorphic_on f ' S
by (simp add: holomorphic_on_open ofs) (blast intro: *)
next

```



```

fix z assume z ∈ S
have deriv f z ≠ 0
  using ⟨z ∈ S⟩ ⟨open S⟩ holf holomorphic_injective_imp_regular injf by
blast
then show deriv f z * deriv (the_inv_into S f) (f z) = 1
  using * [OF ⟨z ∈ S⟩] by (simp add: DERIV_imp_deriv)
next
fix z assume z ∈ S
show the_inv_into S f (f z) = z
  by (simp add: ⟨z ∈ S⟩ injf the_inv_into_f_f)
qed
qed

```

5.6.2 Holomorphism of covering maps and lifts.

lemma *covering_space_lift_is_holomorphic*:

```

assumes cov: covering_space C p S
  and C: open C p holomorphic_on C
  and holf: f holomorphic_on U and fim: f ∈ U → S and gim: g ∈ U → C
  and contg: continuous_on U g and pg_f:  $\bigwedge x. x \in U \implies p(g x) = f x$ 
  shows g holomorphic_on U
unfolding holomorphic_on_def
proof (intro strip)
fix z
assume z ∈ U
with fim have f z ∈ S by blast
then obtain T  $\mathcal{V}$  where f z ∈ T and opeT: openin (top_of_set S) T
  and UV:  $\bigcup \mathcal{V} = C \cap p^{-1} T$ 
  and  $\bigwedge W. W \in \mathcal{V} \implies \text{openin } (top\_of\_set C) W$ 
  and disV: pairwise disjnt  $\mathcal{V}$  and homeV:  $\bigwedge W. W \in \mathcal{V} \implies \exists q. \text{homeomor-}$ 
phism W T p q
  using cov fim unfolding covering_space_def by meson
then have  $\bigwedge W. W \in \mathcal{V} \implies \text{open } W \wedge W \subseteq C$ 
  by (metis ⟨open C⟩ inf_le1 open_Int openin_open)
then obtain V where V ∈  $\mathcal{V}$  g z ∈ V open V V ⊆ C
  by (metis IntI UnionE image_subset_iff vimageI UV ⟨f z ∈ T⟩ ⟨z ∈ U⟩ gim
pg_f image_subset_iff_funcset)
have holf: p holomorphic_on V
  using ⟨V ⊆ C⟩ ⟨p holomorphic_on C⟩ holomorphic_on_subset by blast
moreover have injp: inj_on p V
  by (metis ⟨V ∈  $\mathcal{V}$ ⟩ homeV homeomorphism_def inj_on_inverseI)
ultimately
obtain p' where holf': p' holomorphic_on (p^{-1} V) and pp':  $\bigwedge z. z \in V \implies p'(p$ 
z) = z
  using ⟨open V⟩ holomorphic_has_inverse by metis
have z ∈ U ∩ g^{-1} V
  using ⟨g z ∈ V⟩ ⟨z ∈ U⟩ by blast
moreover have openin (top_of_set U) (U ∩ g^{-1} V)
  using continuous_openin_preimage [OF contg gim]

```

```

  by (meson ‹open V› contg continuous_openin_preimage_eq)
  ultimately obtain  $\varepsilon$  where  $\varepsilon > 0$  and  $e: \text{ball } z \ \varepsilon \cap U \subseteq g^{-1} V$ 
  by (force simp add: openin_contains_ball)
  show  $g$  field_differentiable at  $z$  within  $U$ 
  proof (rule field_differentiable_transform_within)
    show  $(0::\text{real}) < \varepsilon$ 
    by (simp add: ‹ $0 < \varepsilon$ ›)
    show  $z \in U$ 
    by (simp add: ‹ $z \in U$ ›)
    show  $(p' \circ f) x' = g x'$  if  $x' \in U$  dist  $x' z < \varepsilon$  for  $x'$ 
    using that
    by (metis Int_iff comp_apply dist_commute e mem_ball pg_f pp' subsetD
vimage_eq)
  have open  $(p^{-1} V)$ 
  using ‹open V› holp injp open_mapping_thm3 by force
  moreover have  $f z \in p^{-1} V$ 
  by (metis ‹ $z \in U$ › image_iff pg_f ‹ $g z \in V$ ›)
  ultimately have  $p'$  field_differentiable at  $(f z)$ 
  using holomorphic_on_imp_differentiable_at holp' by blast
  moreover have  $f$  field_differentiable at  $z$  within  $U$ 
  by (metis (no_types) ‹ $z \in U$ › holf holomorphic_on_def)
  ultimately show  $(p' \circ f)$  field_differentiable at  $z$  within  $U$ 
  by (metis (no_types) field_differentiable_at_within field_differentiable_compose_within)
qed
qed

```

lemma *covering_space_lift_holomorphic*:

```

  assumes cov: covering_space C p S
  and C: open C p holomorphic_on C
  and f: f holomorphic_on U f ∈ U → S
  and U: simply_connected U locally_path_connected U
  obtains g where g holomorphic_on U g ∈ U → C  $\wedge$   $\forall y. y \in U \implies p(g y) = f y$ 
  proof –
    obtain g where continuous_on U g g ∈ U → C  $\wedge$   $\forall y. y \in U \implies p(g y) = f y$ 
    using covering_space_lift [OF cov U] f holomorphic_on_imp_continuous_on
  by blast
  then show ?thesis
  by (metis C cov covering_space_lift_is_holomorphic f image_subset_iff funcset
that)
qed

```

5.7 The Schwarz Lemma

lemma *Schwarz1*:

```

  assumes holf: f holomorphic_on S
  and contf: continuous_on (closure S) f
  and S: open S connected S
  and boS: bounded S
  and S ≠ {}

```

```

obtains  $w$  where  $w \in \text{frontier } S$ 
   $\bigwedge z. z \in \text{closure } S \implies \text{norm } (f z) \leq \text{norm } (f w)$ 
proof -
  have  $\text{conff}: \text{continuous\_on } (\text{closure } S) (\text{norm } \circ f)$ 
    using  $\text{contf continuous\_on\_compose continuous\_on\_norm\_id}$  by  $\text{blast}$ 
  have  $\text{coc}: \text{compact } (\text{closure } S)$ 
    by  $(\text{simp add: } \langle \text{bounded } S \rangle \text{ bounded\_closure compact\_eq\_bounded\_closed})$ 
  then obtain  $x$  where  $x: x \in \text{closure } S$  and  $x\text{max}: \bigwedge z. z \in \text{closure } S \implies \text{norm}(f z) \leq \text{norm}(f x)$ 
    using  $\text{continuous\_attains\_sup } [OF \_ \_ \text{conff}] \langle S \neq \{\} \rangle$  by  $\text{auto}$ 
  then show  $?thesis$ 
  proof  $(\text{cases } x \in \text{frontier } S)$ 
    case  $\text{True}$ 
      then show  $?thesis$  using  $\text{that } x\text{max}$  by  $\text{blast}$ 
    next
      case  $\text{False}$ 
      then have  $x \in S$ 
        using  $\langle \text{open } S \rangle \text{ frontier\_def interior\_eq } x$  by  $\text{auto}$ 
      then have  $f$   $\text{constant\_on } S$ 
        proof  $(\text{rule maximum\_modulus\_principle } [OF \text{ holf } S \langle \text{open } S \rangle \text{ order\_refl}])$ 
          show  $\bigwedge z. z \in S \implies \text{cmod } (f z) \leq \text{cmod } (f x)$ 
            using  $\text{closure\_subset}$  by  $(\text{blast intro: } x\text{max})$ 
          qed
        then have  $f$   $\text{constant\_on } (\text{closure } S)$ 
          by  $(\text{rule constant\_on\_closureI } [OF \_ \_ \text{contf}])$ 
        then obtain  $c$  where  $c: \bigwedge x. x \in \text{closure } S \implies f x = c$ 
          by  $(\text{meson constant\_on\_def})$ 
        obtain  $w$  where  $w \in \text{frontier } S$ 
          by  $(\text{metis coc all\_not\_in\_conv assms}(6) \text{ closure\_UNIV frontier\_eq\_empty not\_compact\_UNIV})$ 
        then show  $?thesis$ 
          by  $(\text{simp add: } c \text{ frontier\_def that})$ 
        qed
      qed

```

lemma *Schwarz2*:

```

 $\llbracket f \text{ holomorphic\_on ball } 0 r;$ 
 $0 < s; \text{ball } w s \subseteq \text{ball } 0 r;$ 
 $\bigwedge z. \text{norm } (w - z) < s \implies \text{norm}(f z) \leq \text{norm}(f w) \rrbracket$ 
 $\implies f \text{ constant\_on ball } 0 r$ 
by  $(\text{rule maximum\_modulus\_principle } [\text{where } U = \text{ball } w s \text{ and } \xi = w]) (\text{simp\_all add: dist\_norm})$ 

```

lemma *Schwarz3*:

```

assumes  $\text{holf}: f \text{ holomorphic\_on } (\text{ball } 0 r)$  and  $[\text{simp}]: f 0 = 0$ 
obtains  $h$  where  $h \text{ holomorphic\_on } (\text{ball } 0 r)$  and  $\bigwedge z. \text{norm } z < r \implies f z = z * (h z)$  and  $\text{deriv } f 0 = h 0$ 
proof -
  define  $h$  where  $h z = (\text{if } z = 0 \text{ then deriv } f 0 \text{ else } f z / z)$  for  $z$ 

```

```

have d0: deriv f 0 = h 0
  by (simp add: h_def)
moreover have h holomorphic_on (ball 0 r)
  by (rule pole_theorem_open_0 [OF holf, of 0]) (auto simp: h_def)
moreover have norm z < r  $\implies$  f z = z * h z for z
  by (simp add: h_def)
ultimately show ?thesis
  using that by blast
qed

```

proposition *Schwarz Lemma:*

```

assumes holf: f holomorphic_on (ball 0 1) and [simp]: f 0 = 0
  and no:  $\bigwedge z. \text{norm } z < 1 \implies \text{norm } (f z) < 1$ 
  and xi: norm xi < 1
shows norm (f xi)  $\leq$  norm xi and norm(deriv f 0)  $\leq$  1
  and (( $\exists z. \text{norm } z < 1 \wedge z \neq 0 \wedge \text{norm}(f z) = \text{norm } z$ )
     $\vee$  norm(deriv f 0) = 1)
 $\implies \exists \alpha. (\forall z. \text{norm } z < 1 \longrightarrow f z = \alpha * z) \wedge \text{norm } \alpha = 1$ 
(is ?P  $\implies$  ?Q)
proof -
  obtain h where holf: h holomorphic_on (ball 0 1)
    and fz_eq:  $\bigwedge z. \text{norm } z < 1 \implies f z = z * (h z)$  and df0: deriv f 0 = h
  0
  by (rule Schwarz3 [OF holf]) auto
  have noh_le: norm (h z)  $\leq$  1 if z: norm z < 1 for z
  proof -
    have norm (h z) < a if a: 1 < a for a
    proof -
      have max (inverse a) (norm z) < 1
        using z a by (simp_all add: inverse_less_1_iff)
      then obtain r where r: max (inverse a) (norm z) < r and r < 1
        using Rats_dense_in_real by blast
      then have nzs: norm z < r and ira: inverse r < a
        using z a less_imp_inverse_less by force+
      then have 0 < r
        by (meson norm_not_less_zero not_le order.strict_trans2)
      have holf': h holomorphic_on ball 0 r
        by (meson holf  $\langle r < 1 \rangle$  holomorphic_on_subset less_eq_real_def subset_ball)
      have conth': continuous_on (cball 0 r) h
      by (meson  $\langle r < 1 \rangle$  dual_order.trans holf holomorphic_on_imp_continuous_on
        holomorphic_on_subset mem_ball_0 mem_cball_0 not_less subsetI)
      obtain w where w: norm w = r and lenw:  $\bigwedge z. \text{norm } z < r \implies \text{norm}(h z) \leq \text{norm}(h w)$ 
      apply (rule Schwarz1 [OF holf']) using conth'  $\langle 0 < r \rangle$  by auto
      have h w = f w / w using fz_eq  $\langle r < 1 \rangle$  nzs w by auto
      then have cmod (h z) < inverse r
        by (metis  $\langle 0 < r \rangle \langle r < 1 \rangle$  divide_strict_right_mono inverse_eq_divide
          le_less_trans lenw no norm_divide nzs w)

```

```

    then show ?thesis using ira by linarith
  qed
  then show norm (h z) ≤ 1
    using not_le by blast
  qed
  show cmod (f ξ) ≤ cmod ξ
  proof (cases ξ = 0)
    case True then show ?thesis by auto
  next
    case False
    then show ?thesis
      by (simp add: noh_le fz_eq ξ mult_left_le norm_mult)
  qed
  show no_df0: norm (deriv f 0) ≤ 1
    by (simp add: ‹ $\bigwedge z. \text{cmod } z < 1 \implies \text{cmod } (h z) \leq 1$ › df0)
  show ?Q if ?P
    using that
  proof
    assume  $\exists z. \text{cmod } z < 1 \wedge z \neq 0 \wedge \text{cmod } (f z) = \text{cmod } z$ 
    then obtain  $\gamma$  where  $\gamma: \text{cmod } \gamma < 1 \wedge \gamma \neq 0 \wedge \text{cmod } (f \gamma) = \text{cmod } \gamma$  by blast
    then have [simp]: norm (h  $\gamma$ ) = 1
      by (simp add: fz_eq norm_mult)
    have  $\S: \text{ball } \gamma (1 - \text{cmod } \gamma) \subseteq \text{ball } 0 1$ 
      by (simp add: ball_subset_ball_iff)
    moreover have  $\bigwedge z. \text{cmod } (\gamma - z) < 1 - \text{cmod } \gamma \implies \text{cmod } (h z) \leq \text{cmod } (h$ 
 $\gamma)$ 
      by (metis ‹ $\text{cmod } (h \gamma) = 1$ ›  $\S$  dist_0_norm dist_complex_def in_mono mem_ball noh_le)
    ultimately obtain c where  $c: \bigwedge z. \text{norm } z < 1 \implies h z = c$ 
      using Schwarz2 [OF holh, of 1 - norm  $\gamma$   $\gamma$ , unfolded constant_on_def]  $\gamma$ 
    by auto
    then have norm c = 1
      using  $\gamma$  by force
    with c show ?thesis
      using fz_eq by auto
  next
    assume [simp]: cmod (deriv f 0) = 1
    then obtain c where  $c: \bigwedge z. \text{norm } z < 1 \implies h z = c$ 
      using Schwarz2 [OF holh zero_less_one, of 0, unfolded constant_on_def]
    df0 noh_le
      by auto
    moreover have norm c = 1 using df0 c by auto
    ultimately show ?thesis
      using fz_eq by auto
  qed
  qed

```

corollary Schwarz_Lemma':

```

  assumes holf: f holomorphic_on (ball 0 1) and [simp]: f 0 = 0

```

and no: $\bigwedge z. \text{norm } z < 1 \implies \text{norm } (f z) < 1$
shows $((\forall \xi. \text{norm } \xi < 1 \longrightarrow \text{norm } (f \xi) \leq \text{norm } \xi)$
 $\wedge \text{norm}(\text{deriv } f 0) \leq 1)$
 $\wedge (((\exists z. \text{norm } z < 1 \wedge z \neq 0 \wedge \text{norm}(f z) = \text{norm } z)$
 $\vee \text{norm}(\text{deriv } f 0) = 1)$
 $\longrightarrow (\exists \alpha. (\forall z. \text{norm } z < 1 \longrightarrow f z = \alpha * z) \wedge \text{norm } \alpha = 1))$
using *Schwarz_Lemma* [OF *assms*]
by (*metis* (*no_types*) *norm_eq_zero zero_less_one*)

5.8 The Schwarz reflection principle

lemma *hol_pal_lem0*:

assumes $d \cdot a \leq k \leq d \cdot b$

obtains *c* **where**

$c \in \text{closed_segment } a b \ d \cdot c = k$

$\bigwedge z. z \in \text{closed_segment } a c \implies d \cdot z \leq k$

$\bigwedge z. z \in \text{closed_segment } c b \implies k \leq d \cdot z$

proof –

obtain *c* **where** *cin*: $c \in \text{closed_segment } a b$ **and** *kek*: $k = d \cdot c$

using *connected_ivt_hyperplane* [of *closed_segment a b a b d k*]

by (*auto simp: assms*)

have $\text{closed_segment } a c \subseteq \{z. d \cdot z \leq k\}$ $\text{closed_segment } c b \subseteq \{z. k \leq d \cdot z\}$

unfolding *segment_convex_hull* **using** *assms kek*

by (*auto simp: convex_halfspace_le convex_halfspace_ge hull_minimal*)

then show *?thesis* **using** *cin* **that** **by** *fastforce*

qed

lemma *hol_pal_lem1*:

assumes *convex S* *open S*

and *abc*: $a \in S \ b \in S \ c \in S$

$d \neq 0$ **and** *lek*: $d \cdot a \leq k \leq d \cdot b \leq k \leq d \cdot c \leq k$

and *hol1*: f *holomorphic_on* $\{z. z \in S \wedge d \cdot z < k\}$

and *contf*: *continuous_on S f*

shows $\text{contour_integral } (\text{linepath } a b) f +$

$\text{contour_integral } (\text{linepath } b c) f +$

$\text{contour_integral } (\text{linepath } c a) f = 0$

proof –

have $\text{interior } (\text{convex hull } \{a, b, c\}) \subseteq \text{interior}(S \cap \{x. d \cdot x \leq k\})$

proof (*intro interior_mono hull_minimal*)

show $\{a, b, c\} \subseteq S \cap \{x. d \cdot x \leq k\}$

by (*simp add: abc lek*)

show $\text{convex } (S \cap \{x. d \cdot x \leq k\})$

by (*rule convex_Int* [OF $\langle \text{convex } S \rangle$ *convex_halfspace_le*])

qed

also have $\dots \subseteq \{z \in S. d \cdot z < k\}$

by (*force simp: interior_open* [OF $\langle \text{open } S \rangle$] $\langle d \neq 0 \rangle$)

finally have $*$: $\text{interior } (\text{convex hull } \{a, b, c\}) \subseteq \{z \in S. d \cdot z < k\}$.

have *continuous_on* $(\text{convex hull } \{a, b, c\}) f$

using $\langle \text{convex } S \rangle$ *contf abc continuous_on_subset subset_hull*

```

  by fastforce
  moreover have f holomorphic_on interior (convex hull {a,b,c})
    by (rule holomorphic_on_subset [OF holf1 *])
  ultimately show ?thesis
    using Cauchy_theorem_triangle_interior_has_chain_integral_chain_integral3
    by blast
qed

```

lemma hol_pal_lem2:

```

  assumes S: convex S open S
    and abc: a ∈ S b ∈ S c ∈ S
    and d ≠ 0 and lek: d · a ≤ k d · b ≤ k
    and holf1: f holomorphic_on {z. z ∈ S ∧ d · z < k}
    and holf2: f holomorphic_on {z. z ∈ S ∧ k < d · z}
    and contf: continuous_on S f
  shows contour_integral (linepath a b) f +
    contour_integral (linepath b c) f +
    contour_integral (linepath c a) f = 0
proof (cases d · c ≤ k)
  case True show ?thesis
    by (rule hol_pal_lem1 [OF S abc ⟨d ≠ 0⟩ lek True holf1 contf])
  next
  case False
  then have d · c > k by force
  obtain a' where a': a' ∈ closed_segment b c and d · a' = k
    and ba': ∧z. z ∈ closed_segment b a' ⇒ d · z ≤ k
    and a'c: ∧z. z ∈ closed_segment a' c ⇒ k ≤ d · z
    using False hol_pal_lem0 [of d b k c, OF ⟨d · b ≤ k⟩] by auto
  obtain b' where b': b' ∈ closed_segment a c and d · b' = k
    and ab': ∧z. z ∈ closed_segment a b' ⇒ d · z ≤ k
    and b'c: ∧z. z ∈ closed_segment b' c ⇒ k ≤ d · z
    using False hol_pal_lem0 [of d a k c, OF ⟨d · a ≤ k⟩] by auto
  have a'b': a' ∈ S ∧ b' ∈ S
    using a' abc b' convex_contains_segment ⟨convex S⟩ by auto
  have continuous_on (closed_segment c a) f
    by (meson abc contf continuous_on_subset convex_contains_segment ⟨convex
S⟩)
  then have 1: contour_integral (linepath c a) f =
    contour_integral (linepath c b') f + contour_integral (linepath b' a) f
    using b' closed_segment_commute contour_integral_split_linepath by blast
  have continuous_on (closed_segment b c) f
    by (meson abc contf continuous_on_subset convex_contains_segment ⟨convex
S⟩)
  then have 2: contour_integral (linepath b c) f =
    contour_integral (linepath b a') f + contour_integral (linepath a' c) f
    by (rule contour_integral_split_linepath [OF _ a'])
  have 3: contour_integral (reversepath (linepath b' a')) f =
    - contour_integral (linepath b' a') f
    by (rule contour_integral_reversepath [OF valid_path_linepath])

```

```

have fcd_le: f field_differentiable at x
  if x ∈ interior S ∧ x ∈ interior {x. d · x ≤ k} for x
proof -
  have f holomorphic_on S ∩ {c. d · c < k}
    by (metis (no_types) Collect_conj_eq Collect_mem_eq holf1)
  then have ∃ C D. x ∈ interior C ∩ interior D ∧ f holomorphic_on interior C
    ∩ interior D
    using that
    by (metis Collect_mem_eq Int_Collect ⟨d ≠ 0⟩ interior_halfspace_le interior_open ⟨open S⟩)
  then show f field_differentiable at x
    by (metis at_within_interior holomorphic_on_def interior_Int interior_interior)
qed
have ab_le: ∧x. x ∈ closed_segment a b ⇒ d · x ≤ k
proof -
  fix x :: complex
  assume x ∈ closed_segment a b
  then have ∧C. x ∈ C ∨ b ∉ C ∨ a ∉ C ∨ ¬ convex C
    by (meson contra_subsetD convex_contains_segment)
  then show d · x ≤ k
    by (metis lek_convex_halfspace_le mem_Collect_eq)
qed
have cs: closed_segment a' b' ⊆ {x. d · x ≤ k} ∧ closed_segment b' a' ⊆ {x. d · x ≤ k}
  by (simp add: ⟨d · a' = k⟩ ⟨d · b' = k⟩ closed_segment_subset convex_halfspace_le lek(1))
have continuous_on (S ∩ {x. d · x ≤ k}) f using contf
  by (simp add: continuous_on_subset)
then have (f has_contour_integral 0)
  (linepath a b +++ linepath b a' +++ linepath a' b' +++ linepath b' a)
  apply (rule Cauchy_theorem_convex [where K = {}])
  by (simp_all add: path_image_join convex_Int convex_halfspace_le ⟨convex S⟩ fcd_le ab_le
    closed_segment_subset abc a'b' ba' cs)
then have 4: contour_integral (linepath a b) f +
  contour_integral (linepath b a') f +
  contour_integral (linepath a' b') f +
  contour_integral (linepath b' a) f = 0
  by (rule has_chain_integral_chain_integral4)
have fcd_ge: f field_differentiable at x
  if x ∈ interior S ∧ x ∈ interior {x. k ≤ d · x} for x
proof -
  have f2: f holomorphic_on S ∩ {c. k < d · c}
    by (metis (full_types) Collect_conj_eq Collect_mem_eq holf2)
  have f3: interior S = S
    by (simp add: interior_open ⟨open S⟩)
  then have x ∈ S ∩ interior {c. k ≤ d · c}
    using that by simp
  then show f field_differentiable at x

```



```

    using f3 f2 unfolding holomorphic_on_def
    by (metis (no_types) ‹d ≠ 0› at_within_interior interior_Int interior_halfspace_ge
interior_interior)
  qed
  have cs: closed_segment c b' ⊆ {x. k ≤ d · x} ∧ closed_segment b' a' ⊆ {x. k
≤ d · x}
    by (simp add: ‹d · a' = k› b'c closed_segment_subset convex_halfspace_ge)
  have continuous_on (S ∩ {x. k ≤ d · x}) f using contf
    by (simp add: continuous_on_subset)
  then have (f has_contour_integral 0) (linepath a' c +++ linepath c b' +++
linepath b' a')
    apply (rule Cauchy_theorem_convex [where K = {}])
    by (simp_all add: path_image_join convex_Int convex_halfspace_ge ‹convex
S›
        fcd_ge closed_segment_subset abc a'b' a'c cs)
  then have 5: contour_integral (linepath a' c) f + contour_integral (linepath c
b') f + contour_integral (linepath b' a') f = 0
    by (rule has_chain_integral_chain_integral3)
  show ?thesis
    using 1 2 3 4 5 by (metis add.assoc eq_neg_iff_add_eq_0 reversepath_linepath)
  qed

```

lemma hol_pal_lem3:

```

  assumes S: convex S open S
    and abc: a ∈ S b ∈ S c ∈ S
    and d ≠ 0 and lek: d · a ≤ k
    and holf1: f holomorphic_on {z. z ∈ S ∧ d · z < k}
    and holf2: f holomorphic_on {z. z ∈ S ∧ k < d · z}
    and contf: continuous_on S f
  shows contour_integral (linepath a b) f +
        contour_integral (linepath b c) f +
        contour_integral (linepath c a) f = 0
proof (cases d · b ≤ k)
  case True show ?thesis
    by (rule hol_pal_lem2 [OF S abc ‹d ≠ 0› lek True holf1 holf2 contf])
next
  case False
  show ?thesis
  proof (cases d · c ≤ k)
    case True
    have contour_integral (linepath c a) f +
          contour_integral (linepath a b) f +
          contour_integral (linepath b c) f = 0
      by (rule hol_pal_lem2 [OF S ‹c ∈ S› ‹a ∈ S› ‹b ∈ S› ‹d ≠ 0› ‹d · c ≤ k›
lek holf1 holf2 contf])
    then show ?thesis
      by (simp add: algebra_simps)
  next
    case False

```

```

have contour_integral (linepath b c) f +
      contour_integral (linepath c a) f +
      contour_integral (linepath a b) f = 0
using hol_pal_lem2 [OF S ‹b ∈ S› ‹c ∈ S› ‹a ∈ S›, of -d -k]
using ‹d ≠ 0› ‹¬ d · b ≤ k› False by (simp_all add: holf1 holf2 contf)
then show ?thesis
by (simp add: algebra_simps)
qed
qed

```

lemma *hol_pal_lem4*:

```

assumes S: convex S open S
and abc: a ∈ S b ∈ S c ∈ S and d ≠ 0
and holf1: f holomorphic_on {z. z ∈ S ∧ d · z < k}
and holf2: f holomorphic_on {z. z ∈ S ∧ k < d · z}
and contf: continuous_on S f
shows contour_integral (linepath a b) f +
      contour_integral (linepath b c) f +
      contour_integral (linepath c a) f = 0
proof (cases d · a ≤ k)
case True show ?thesis
by (rule hol_pal_lem3 [OF S abc ‹d ≠ 0› True holf1 holf2 contf])
next
case False
show ?thesis
using ‹d ≠ 0› hol_pal_lem3 [OF S abc, of -d -k] False
by (simp_all add: holf1 holf2 contf)
qed

```

lemma *holomorphic_on_paste_across_line*:

```

assumes S: open S and d ≠ 0
and holf1: f holomorphic_on (S ∩ {z. d · z < k})
and holf2: f holomorphic_on (S ∩ {z. k < d · z})
and contf: continuous_on S f
shows f holomorphic_on S
proof -
have *: ∃ t. open t ∧ p ∈ t ∧ continuous_on t f ∧
      (∀ a b c. convex hull {a, b, c} ⊆ t →
       contour_integral (linepath a b) f +
       contour_integral (linepath b c) f +
       contour_integral (linepath c a) f = 0)
if p ∈ S for p
proof -
obtain e where e > 0 and e: ball p e ⊆ S
using ‹p ∈ S› openE S by blast
then have continuous_on (ball p e) f
using contf continuous_on_subset by blast
moreover have f holomorphic_on {z. dist p z < e ∧ d · z < k}
apply (rule holomorphic_on_subset [OF holf1])

```

```

    using e by auto
  moreover have f holomorphic_on {z. dist p z < e ∧ k < d · z}
    apply (rule holomorphic_on_subset [OF holf2])
    using e by auto
  ultimately show ?thesis
    apply (rule_tac x=ball p e in exI)
    using ‹e > 0› e ‹d ≠ 0› hol_pal_lem4 [of ball p e _ _ d _ k]
    by (force simp add: subset_hull)
qed
show ?thesis
  by (blast intro: * Morera_local_triangle_analytic_imp_holomorphic)
qed

```

proposition Schwarz_reflection:

```

  assumes open S and cnjs: cnj ' S ⊆ S
    and holf: f holomorphic_on (S ∩ {z. 0 < Im z})
    and contf: continuous_on (S ∩ {z. 0 ≤ Im z}) f
    and f: ∧z. [z ∈ S; z ∈ ℝ] ⇒ (f z) ∈ ℝ
  shows (λz. if 0 ≤ Im z then f z else cnj(f (cnj z))) holomorphic_on S
proof -
  have 1: (λz. if 0 ≤ Im z then f z else cnj (f (cnj z))) holomorphic_on (S ∩ {z.
0 < Im z})
  by (force intro: iffD1 [OF holomorphic_cong [OF refl] holf])
  have cont_cfc: continuous_on (S ∩ {z. Im z ≤ 0}) (cnj o f o cnj)
  using cnjs
  by (intro continuous_intros continuous_on_compose continuous_on_subset
[OF contf]) auto
  have cnj o f o cnj field_differentiable at x within S ∩ {z. Im z < 0}
    if x ∈ S Im x < 0 f field_differentiable at (cnj x) within S ∩ {z. 0 < Im z}
for x
  using that
  apply (clarsimp simp add: field_differentiable_def has_field_derivative_iff
Lim_within dist_norm)
  apply (rule_tac x=cnj f' in exI)
  apply (elim_all_forward ex_forward conj_forward imp_forward asm_rl, clarify)
  apply (drule_tac x=cnj xa in bspec)
  using cnjs apply force
  apply (metis complex_cnj_cnj complex_cnj_diff complex_cnj_divide complex_mod_cnj)
  done
  then have hol_cfc: (cnj o f o cnj) holomorphic_on (S ∩ {z. Im z < 0})
  using holf cnjs
  by (force simp: holomorphic_on_def)
  have 2: (λz. if 0 ≤ Im z then f z else cnj (f (cnj z))) holomorphic_on (S ∩ {z.
Im z < 0})
  apply (rule iffD1 [OF holomorphic_cong [OF refl]])
  using hol_cfc by auto
  have [simp]: (S ∩ {z. 0 ≤ Im z}) ∪ (S ∩ {z. Im z ≤ 0}) = S

```

```

  by force
  have eq:  $\bigwedge z. \llbracket z \in S; \text{Im } z \leq 0; 0 \leq \text{Im } z \rrbracket \implies f z = \text{cnj } (f (\text{cnj } z))$ 
    using f Reals_cnj_iff complex_is_Real_iff by auto
  have continuous_on ((S  $\cap$  {z. 0  $\leq$  Im z})  $\cup$  (S  $\cap$  {z. Im z  $\leq$  0}))
    ( $\lambda z. \text{if } 0 \leq \text{Im } z \text{ then } f z \text{ else } \text{cnj } (f (\text{cnj } z))$ )
    apply (rule continuous_on_cases_local)
    using cont_cfc contf
  by (simp_all add: closedin_closed_Int closed_halfspace_Im_le closed_halfspace_Im_ge
  eq)
  then have  $\exists$ : continuous_on S ( $\lambda z. \text{if } 0 \leq \text{Im } z \text{ then } f z \text{ else } \text{cnj } (f (\text{cnj } z))$ )
    by force
  show ?thesis
    apply (rule holomorphic_on_paste_across_line [OF ⟨open S⟩, of - i 0])
    using 1 2 3 by auto
qed

```

5.9 Bloch's theorem

```

lemma Bloch_lemma_0:
  assumes holf: f holomorphic_on cball 0 r and 0 < r
    and [simp]: f 0 = 0
    and le:  $\bigwedge z. \text{norm } z < r \implies \text{norm}(\text{deriv } f z) \leq 2 * \text{norm}(\text{deriv } f 0)$ 
    shows ball 0 ((3 - 2 * sqrt 2) * r * norm(deriv f 0))  $\subseteq$  f ` ball 0 r
proof -
  have sqrt 2 < 3/2
    by (rule real_less_sqrt) (auto simp: power2_eq_square)
  then have sq3: 0 < 3 - 2 * sqrt 2 by simp
  show ?thesis
proof (cases deriv f 0 = 0)
  case True then show ?thesis by simp
next
  case False
  define C where C = 2 * norm(deriv f 0)
  have 0 < C using False by (simp add: C_def)
  have holf': f holomorphic_on ball 0 r using holf
    using ball_subset_cball holomorphic_on_subset by blast
  then have holdf': deriv f holomorphic_on ball 0 r
    by (rule holomorphic_deriv [OF _ open_ball])
  have Le1: norm(deriv f z - deriv f 0)  $\leq$  norm z / (r - norm z) * C
    if norm z < r for z
proof -
  have T1: norm(deriv f z - deriv f 0)  $\leq$  norm z / (R - norm z) * C
    if R: norm z < R R < r for R
proof -
  have 0 < R using R
    by (metis less_trans norm_zero zero_less_norm_iff)
  have df_le:  $\bigwedge x. \text{norm } x < r \implies \text{norm}(\text{deriv } f x) \leq C$ 
    using le by (simp add: C_def)
  have hol_df: deriv f holomorphic_on cball 0 R

```

```

    using R holdf' holomorphic_on_subset by auto
    have *: (( $\lambda w. \text{deriv } f w / (w - z)$ ) has_contour_integral  $2 * \pi * i * \text{deriv } f z$ ) (circlepath 0 R)
      if norm z < R for z
      using <0 < R> that Cauchy_integral_formula_convex_simple [OF convex_cball hol_df, of_circlepath 0 R]
      by (force simp: winding_number_circlepath)
    have **: (( $\lambda x. \text{deriv } f x / (x - z) - \text{deriv } f x / x$ ) has_contour_integral_of_real  $(2 * \pi) * i * (\text{deriv } f z - \text{deriv } f 0)$ ) (circlepath 0 R)
      using has_contour_integral_diff [OF * [of z] * [of 0]] <0 < R> that
      by (simp add: algebra_simps)
    have [simp]:  $\bigwedge x. \text{norm } x = R \implies x \neq z$  using that(1) by blast
    have norm ( $\text{deriv } f x / (x - z) - \text{deriv } f x / x$ )
       $\leq C * \text{norm } z / (R * (R - \text{norm } z))$ 
      if norm x = R for x
    proof -
      have [simp]:  $\text{norm } (\text{deriv } f x * x - \text{deriv } f x * (x - z)) = \text{norm } (\text{deriv } f x) * \text{norm } z$ 
        by (simp add: norm_mult_right_diff_distrib')
      show ?thesis
        using <0 < R> <0 < C> R that
        by (auto simp add: norm_mult norm_divide divide_simps df_le_mult_mono norm_triangle_ineq2)
      qed
    then show ?thesis
      using has_contour_integral_bound_circlepath [OF **, of  $C * \text{norm } z / (R * (R - \text{norm } z))$ ] <0 < R> <0 < C> R
      apply (simp add: norm_mult norm_divide)
      apply (simp add: divide_simps mult.commute)
      done
    qed
  obtain r' where r':  $\text{norm } z < r' r' < r$ 
    using Rats_dense_in_real [of norm z r] <norm z < r> by blast
  then have [simp]:  $\text{closure } \{r' <..< r\} = \{r' .. r\}$  by simp
  show ?thesis
    apply (rule continuous_ge_on_closure
      [where f =  $\lambda r. \text{norm } z / (r - \text{norm } z) * C$  and s =  $\{r' <..< r\}$ , OF _ T1])
    using that r'
    by (auto simp: not_le intro!: continuous_intros)
  qed
  have *:  $(\text{norm } z - \text{norm } z^2 / (r - \text{norm } z)) * \text{norm } (\text{deriv } f 0) \leq \text{norm } (f z)$ 
    if r:  $\text{norm } z < r$  for z
  proof -
    have 1:  $\bigwedge x. x \in \text{ball } 0 r \implies ((\lambda z. f z - \text{deriv } f 0 * z) \text{ has\_field\_derivative } \text{deriv } f x - \text{deriv } f 0)$ 
      (at x within ball 0 r)

```

```

    by (rule derivative_eq_intros holomorphic_derivI holf' | simp)+
  have 2: closed_segment 0 z  $\subseteq$  ball 0 r
    by (metis <0 < r> convex_ball convex_contains_segment dist_self mem_ball
mem_ball_0 that)
  have 4: norm (deriv f (x *R z) - deriv f 0) * norm z  $\leq$  norm z * norm z *
x * C / (r - norm z)
    if x: 0  $\leq$  x x  $\leq$  1 for x
  proof -
    have [simp]: x * norm z < r
    using r x by (meson le_less_trans mult_le_cancel_right2 norm_not_less_zero)
    have norm (deriv f (x *R z) - deriv f 0)  $\leq$  norm (x *R z) / (r - norm (x
*_R z)) * C
      apply (rule Le1) using r x <0 < r> by simp
    also have ...  $\leq$  norm (x *R z) / (r - norm z) * C
      using r x <0 < r>
    apply (simp add: field_split_simps)
    by (simp add: <0 < C> mult.assoc mult_left_le_one_le ordered_comm_semiring_class.comm_mult_le)
    finally have norm (deriv f (x *R z) - deriv f 0) * norm z  $\leq$  norm (x *R
z) / (r - norm z) * C * norm z
      by (rule mult_right_mono) simp
    with x show ?thesis by (simp add: algebra_simps)
  qed
  have le_norm: abc  $\leq$  norm d - e  $\implies$  norm(f - d)  $\leq$  e  $\implies$  abc  $\leq$  norm f
for abc d e and f::complex
  by (metis add_diff_cancel_left' add_diff_eq diff_left_mono norm_diff_ineq
order_trans)
  have norm (integral {0..1} ( $\lambda$ x. (deriv f (x *R z) - deriv f 0) * z))
 $\leq$  integral {0..1} ( $\lambda$ t. (norm z)2 * t / (r - norm z) * C)
  proof (rule integral_norm_bound_integral)
    show ( $\lambda$ x. (deriv f (x *R z) - deriv f 0) * z) integrable_on {0..1}
      using contour_integral_primitive [OF 1, of linepath 0 z] 2
    by (simp add: has_contour_integral_linepath_has_integral_integrable_integral)
    have (*) ((cmod z)2) integrable_on {0..1}
    by (metis ident_integrable_on integrable_0 integrable_eq integrable_on_cmult_iff
lambda_zero)
    then show ( $\lambda$ t. (norm z)2 * t / (r - norm z) * C) integrable_on {0..1}
      using integrable_on_cmult_right[where 'b=real, simplified] integrable_on_divide
[where 'b=real, simplified]
      by blast
  qed (simp add: norm_mult_power2_eq_square 4)
  then have int_le: norm (f z - deriv f 0 * z)  $\leq$  (norm z)2 * norm(deriv f 0)
/ ((r - norm z))
    using contour_integral_primitive [OF 1, of linepath 0 z] 2
  by (simp add: has_contour_integral_linepath_has_integral_integrable_integral
C_def)
  have norm z * (norm (deriv f 0) * (r - norm z - norm z))  $\leq$  norm z *
(norm (deriv f 0) * (r - norm z) - norm (deriv f 0) * norm z)
    by (simp add: algebra_simps)
  then have  $\S$ : (norm z * (r - norm z) - norm z * norm z) * norm (deriv f

```

```

0) ≤ norm (deriv f 0) * norm z * (r - norm z) - norm z * norm z * norm (deriv
f 0)
  by (simp add: algebra_simps)
show ?thesis
  apply (rule le_norm [OF _ int_le])
  using ‹norm z < r›
  by (simp add: power2_eq_square divide_simps C_def norm_mult §)
qed
have sq201 [simp]: 0 < (1 - sqrt 2 / 2) (1 - sqrt 2 / 2) < 1
  by (auto simp: sqrt2_less_2)
have 1: continuous_on (closure (ball 0 ((1 - sqrt 2 / 2) * r))) f
proof (rule continuous_on_subset [OF holomorphic_on_imp_continuous_on
[OF holf]])
  show closure (ball 0 ((1 - sqrt 2 / 2) * r)) ⊆ cball 0 r
  proof -
    have (1 - sqrt 2 / 2) * r ≤ r
      by (simp add: ‹0 < r›)
    then show ?thesis
      by (meson ball_subset_cball closed_cball closure_minimal dual_order.trans
subset_ball)
  qed
qed
have 2: open (f ` interior (ball 0 ((1 - sqrt 2 / 2) * r)))
proof (rule open_mapping_thm [OF holf' open_ball connected_ball])
  show interior (ball 0 ((1 - sqrt 2 / 2) * r)) ⊆ ball (0::complex) r
    using ‹0 < r› mult_pos_pos sq201 by (simp add: ball_subset_ball_iff)
  show ¬ f constant_on ball 0 r
    using False ‹0 < r› centre_in_ball holf' holomorphic_nonconstant by blast
qed auto
have ball 0 ((3 - 2 * sqrt 2) * r * norm (deriv f 0)) =
  ball (f 0) ((3 - 2 * sqrt 2) * r * norm (deriv f 0))
  by simp
also have ... ⊆ f ` ball 0 ((1 - sqrt 2 / 2) * r)
proof -
  have 3: (3 - 2 * sqrt 2) * r * norm (deriv f 0) ≤ norm (f z)
    if norm z = (1 - sqrt 2 / 2) * r for z
    apply (rule order_trans [OF _ *])
    using ‹0 < r›
    apply (simp_all add: field_simps power2_eq_square that)
    apply (simp add: mult.assoc [symmetric])
    done
  show ?thesis
    apply (rule ball_subset_open_map_image [OF 1 2 _ bounded_ball])
    using ‹0 < r› sq201 3 C_def ‹0 < C› sq3 by auto
qed
also have ... ⊆ f ` ball 0 r
proof -
  have ∧x. (1 - sqrt 2 / 2) * r ≤ r
    using ‹0 < r› by (auto simp: field_simps)

```

```

    then show ?thesis
      by auto
    qed
    finally show ?thesis .
  qed
qed

lemma Bloch_lemma:
  assumes holf: f holomorphic_on cball a r and 0 < r
    and le:  $\bigwedge z. z \in \text{ball } a \ r \implies \text{norm}(\text{deriv } f \ z) \leq 2 * \text{norm}(\text{deriv } f \ a)$ 
    shows ball (f a) ((3 - 2 * sqrt 2) * r * norm(deriv f a))  $\subseteq$  f ' ball a r (is ?lhs
 $\subseteq$  ?rhs)
  proof -
    have fz:  $(\lambda z. f (a + z)) = f \circ (\lambda z. (a + z))$ 
      by (simp add: o_def)
    have hol0:  $(\lambda z. f (a + z))$  holomorphic_on cball 0 r
      unfolding fz by (intro holomorphic_intros holf holomorphic_on_compose |
simp)+
    then have [simp]:  $\bigwedge x. \text{norm } x < r \implies (\lambda z. f (a + z))$  field_differentiable at x
      by (metis open_ball at_within_open ball_subset_cball diff_0 dist_norm holomorphic_on_def holomorphic_on_subset mem_ball norm_minus_cancel)
    have [simp]:  $\bigwedge z. \text{norm } z < r \implies f$  field_differentiable at (a + z)
      by (metis holf open_ball add_diff_cancel_left' dist_complex_def holomorphic_on_imp_differentiable_at holomorphic_on_subset interior_cball interior_subset mem_ball norm_minus_commute)
    then have [simp]: f field_differentiable at a
      by (metis add.comm_neutral <0 < r> norm_eq_zero)
    have hol1:  $(\lambda z. f (a + z) - f a)$  holomorphic_on cball 0 r
      by (intro holomorphic_intros hol0)
    then have  $\S$ : ball 0 ((3 - 2 * sqrt 2) * r * norm (deriv  $(\lambda z. f (a + z) - f a)$ 
0))
       $\subseteq (\lambda z. f (a + z) - f a)$  ' ball 0 r
    apply (rule Bloch_lemma_0)
    using <0 < r>
      apply (simp_all add: <0 < r>)
    apply (simp add: fz deriv_chain dist_norm le)
    done
  show ?thesis
  proof clarify
    fix x
    assume x  $\in$  ?lhs
    with subsetD [OF  $\S$ , of x - f a] show x  $\in$  ?rhs
      by (force simp: fz <0 < r> dist_norm deriv_chain field_differentiable_compose)
    qed
  qed
qed

```

proposition Bloch_unit:

assumes holf: f holomorphic_on ball a 1 and [simp]: deriv f a = 1
 obtains b r where $1/12 < r$ and ball b r \subseteq f ' (ball a 1)


```

proof –
  define  $r :: \text{real}$  where  $r = 249/256$ 
  have  $0 < r \wedge r < 1$  by (auto simp: r_def)
  define  $g$  where  $g\ z = \text{deriv } f\ z * \text{of\_real}(r - \text{norm}(z - a))$  for  $z$ 
  have  $\text{deriv } f$  holomorphic_on ball  $a\ 1$ 
    by (rule holomorphic_deriv [OF holf open_ball])
  then have continuous_on (ball  $a\ 1$ ) (deriv  $f$ )
    using holomorphic_on_imp_continuous_on by blast
  then have continuous_on (cball  $a\ r$ ) (deriv  $f$ )
    by (rule continuous_on_subset) (simp add: cball_subset_ball_iff < $r < 1$ >)
  then have continuous_on (cball  $a\ r$ )  $g$ 
    by (simp add: g_def continuous_intros)
  then have 1: compact (g ‘ cball  $a\ r$ )
    by (rule compact_continuous_image [OF _ compact_cball])
  have 2:  $g\ ‘\ \text{cball } a\ r \neq \{\}$ 
    using < $r > 0$ > by auto
  obtain  $p$  where  $pr: p \in \text{cball } a\ r$ 
    and  $pge: \bigwedge y. y \in \text{cball } a\ r \implies \text{norm } (g\ y) \leq \text{norm } (g\ p)$ 
    using distance_attains_sup [OF 1 2, of 0] by force
  define  $t$  where  $t = (r - \text{norm}(p - a)) / 2$ 
  have  $\text{norm } (p - a) \neq r$ 
    using  $pge$  [of  $a$ ] < $r > 0$ > by (auto simp: g_def norm_mult)
  then have  $\text{norm } (p - a) < r$  using  $pr$ 
    by (simp add: norm_minus_commute dist_norm)
  then have  $0 < t$ 
    by (simp add: t_def)
  have  $cpt: \text{cball } p\ t \subseteq \text{ball } a\ r$ 
    using < $0 < t$ > by (simp add: cball_subset_ball_iff dist_norm t_def field_simps)
  have  $\text{gen\_le\_dfp}: \text{norm } (\text{deriv } f\ y) * (r - \text{norm } (y - a)) / (r - \text{norm } (p - a))$ 
 $\leq \text{norm } (\text{deriv } f\ p)$ 
    if  $y \in \text{cball } a\ r$  for  $y$ 
  proof –
    have [simp]:  $\text{norm } (y - a) \leq r$ 
      using that by (simp add: dist_norm norm_minus_commute)
    have  $\text{norm } (g\ y) \leq \text{norm } (g\ p)$ 
      using  $pge$  [OF that] by simp
    then have  $\text{norm } (\text{deriv } f\ y) * \text{abs } (r - \text{norm } (y - a)) \leq \text{norm } (\text{deriv } f\ p) * \text{abs } (r - \text{norm } (p - a))$ 
      by (simp only: dist_norm g_def norm_mult norm_of_real)
    with that < $\text{norm } (p - a) < r$ > show ?thesis
      by (simp add: dist_norm field_split_simps)
  qed
  have  $\text{le\_norm\_dfp}: r / (r - \text{norm } (p - a)) \leq \text{norm } (\text{deriv } f\ p)$ 
    using  $\text{gen\_le\_dfp}$  [of  $a$ ] < $r > 0$ > by auto
  have 1:  $f$  holomorphic_on cball  $p\ t$ 
    using  $cpt$  < $r < 1$ > order_subst1 subset_ball
    by (force simp add: intro!: holomorphic_on_subset [OF holf])
  have 2:  $\text{norm } (\text{deriv } f\ z) \leq 2 * \text{norm } (\text{deriv } f\ p)$  if  $z \in \text{ball } p\ t$  for  $z$ 
  proof –

```

```

have  $z: z \in \text{cball } a \ r$ 
  by (meson ball_subset_cball subsetD cpt that)
then have  $\text{norm}(z - a) < r$ 
  by (metis ball_subset_cball contra_subsetD cpt dist_norm mem_ball norm_minus_commute
that)
have  $\text{norm}(\text{deriv } f \ z) * (r - \text{norm}(z - a)) / (r - \text{norm}(p - a)) \leq \text{norm}$ 
( $\text{deriv } f \ p$ )
  using gen_le_dfp [OF  $z$ ] by simp
with  $\langle \text{norm}(z - a) < r \rangle \langle \text{norm}(p - a) < r \rangle$ 
have  $\text{norm}(\text{deriv } f \ z) \leq (r - \text{norm}(p - a)) / (r - \text{norm}(z - a)) * \text{norm}$ 
( $\text{deriv } f \ p$ )
  by (simp add: field_simps)
also have  $\dots \leq 2 * \text{norm}(\text{deriv } f \ p)$ 
proof (rule mult_right_mono)
  show  $(r - \text{cmod}(p - a)) / (r - \text{cmod}(z - a)) \leq 2$ 
    using that  $\langle \text{norm}(p - a) < r \rangle \langle \text{norm}(z - a) < r \rangle$  dist_triangle3 [of  $z \ a \ p$ ]
    by (simp add: field_simps t_def dist_norm [symmetric])
qed auto
finally show ?thesis .
qed
have sqrt2:  $\text{sqrt } 2 < 2113/1494$ 
  by (rule real_less_sqrt) (auto simp: power2_eq_square)
then have sq3:  $0 < 3 - 2 * \text{sqrt } 2$  by simp
have  $1 / 12 / ((3 - 2 * \text{sqrt } 2) / 2) < r$ 
  using sq3 sqrt2 by (auto simp: field_simps r_def)
also have  $\dots \leq \text{cmod}(\text{deriv } f \ p) * (r - \text{cmod}(p - a))$ 
  using  $\langle \text{norm}(p - a) < r \rangle$  le_norm_dfp by (simp add: pos_divide_le_eq)
finally have  $1 / 12 < \text{cmod}(\text{deriv } f \ p) * (r - \text{cmod}(p - a)) * ((3 - 2 * \text{sqrt}$ 
2) / 2)
  using pos_divide_less_eq half_gt_zero_iff sq3 by blast
then have **:  $1 / 12 < (3 - 2 * \text{sqrt } 2) * t * \text{norm}(\text{deriv } f \ p)$ 
  using sq3 by (simp add: mult.commute t_def)
have ball (f p)  $((3 - 2 * \text{sqrt } 2) * t * \text{norm}(\text{deriv } f \ p)) \subseteq f \text{' ball } p \ t$ 
  by (rule Bloch_lemma [OF 1  $\langle 0 < t \rangle$  2])
also have  $\dots \subseteq f \text{' ball } a \ 1$ 
proof -
  have ball a r  $\subseteq$  ball a 1
    using  $\langle 0 < t \rangle \langle r < 1 \rangle$  by (simp add: ball_subset_ball_iff dist_norm)
  then show ?thesis
    using ball_subset_cball cpt by blast
qed
finally have ball (f p)  $((3 - 2 * \text{sqrt } 2) * t * \text{norm}(\text{deriv } f \ p)) \subseteq f \text{' ball } a \ 1$  .
with ** show ?thesis
  by (rule that)
qed

```

theorem Bloch:

```

assumes holf:  $f$  holomorphic_on ball a r and  $0 < r$ 
and r':  $r' \leq r * \text{norm}(\text{deriv } f \ a) / 12$ 

```

```

obtains b where  $\text{ball } b \ r' \subseteq f \text{ ` } (\text{ball } a \ r)$ 
proof (cases  $\text{deriv } f \ a = 0$ )
  case True with r' show ?thesis
    using ball_eq_empty that by fastforce
next
  case False
    define C where  $C = \text{deriv } f \ a$ 
    have  $0 < \text{norm } C$  using False by (simp add: C_def)
    have dfa:  $f \text{ field\_differentiable at } a$ 
      using  $\langle 0 < r \rangle$  holomorphic_on_imp_differentiable_at [OF holf] by auto
    have fo:  $(\lambda z. f (a + \text{of\_real } r * z)) = f \circ (\lambda z. (a + \text{of\_real } r * z))$ 
      by (simp add: o_def)
    have holf':  $f \text{ holomorphic\_on } (\lambda z. a + \text{complex\_of\_real } r * z) \text{ ` ball } 0 \ 1$ 
      using  $\langle 0 < r \rangle$  holomorphic_on_subset [OF holf] by (force simp: dist_norm
norm_mult)
    have 1:  $(\lambda z. f (a + r * z) / (C * r)) \text{ holomorphic\_on ball } 0 \ 1$ 
      using  $\langle 0 < r \rangle \ \langle 0 < \text{norm } C \rangle$ 
      by (intro holomorphic_intros holomorphic_on_compose holf'; simp add: fo)
    have  $((\lambda z. f (a + \text{of\_real } r * z) / (C * \text{of\_real } r)) \text{ has\_field\_derivative}$ 
       $(\text{deriv } f (a + \text{of\_real } r * z) / C)) \text{ (at } z)$ 
      if  $\text{norm } z < 1$  for z
    proof -
      have fd:  $f \text{ field\_differentiable at } (a + \text{complex\_of\_real } r * z)$ 
        using  $\langle 0 < r \rangle$  by (simp_all add: dist_norm norm_mult holomorphic_on_imp_differentiable_at
[OF holf] that)
      have *:  $((\lambda x. f (a + \text{of\_real } r * x)) \text{ has\_field\_derivative}$ 
         $(\text{deriv } f (a + \text{of\_real } r * z) * \text{of\_real } r)) \text{ (at } z)$ 
        by (rule fd DERIV_chain [OF field_differentiable_derivI] derivative_eq_intros
| simp add: fo)
      show ?thesis
        apply (rule derivative_eq_intros * | simp)+
        using  $\langle 0 < r \rangle$  by (auto simp: C_def False)
    qed
    have  $\text{deriv } (\lambda z. f (a + \text{of\_real } r * z) / (C * \text{of\_real } r)) \ 0 = \text{deriv } (\lambda z. f (a +$ 
complex_of_real  $r * z)) \ 0 /$ 
       $(C * \text{complex\_of\_real } r)$ 
      apply (rule deriv_cdivide_right)
      by (metis (no_types) DERIV_chain2 add.right_neutral dfa field_differentiable_add_const
field_differentiable_def field_differentiable_linear fo mult_zero_right)
    also have ... = 1
      using  $\langle 0 < r \rangle$  by (simp add: C_def False fo derivative_intros dfa deriv_chain)
    finally have 2:  $\text{deriv } (\lambda z. f (a + \text{of\_real } r * z) / (C * \text{of\_real } r)) \ 0 = 1$  .
    have sb1:  $(*) \ (C * r) \text{ ` } (\lambda z. f (a + \text{of\_real } r * z) / (C * r)) \text{ ` ball } 0 \ 1$ 
       $\subseteq f \text{ ` ball } a \ r$ 
      using  $\langle 0 < r \rangle$  by (auto simp: dist_norm norm_mult C_def False)
    have sb2:  $\text{ball } (C * r * b) \ r' \subseteq (*) \ (C * r) \text{ ` ball } b \ t$ 
      if  $1 / 12 < t$  for b t
    proof -
      have *:  $r * \text{cmod } (\text{deriv } f \ a) / 12 \leq r * (t * \text{cmod } (\text{deriv } f \ a))$ 

```

```

    using that <0 < r> less_eq_real_def mult.commute mult.right_neutral mult_left_mono
norm_ge_zero times_divide_eq_right
    by auto
    show ?thesis
    apply clarify
    apply (rule_tac x=x / (C * r) in image_eqI)
    using <0 < r> apply (simp_all add: dist_norm norm_mult norm_divide
C_def False field_simps)
    using * r' by linarith
    qed
    show ?thesis
    apply (rule Bloch_unit [OF 1 2])
    using image_mono sb1 sb2 that by fastforce
    qed

```

corollary *Bloch_general*:

```

    assumes holf: f holomorphic_on S and a ∈ S
    and tle:  $\bigwedge z. z \in \text{frontier } S \implies t \leq \text{dist } a z$ 
    and rle:  $r \leq t * \text{norm}(\text{deriv } f a) / 12$ 
    obtains b where ball b r  $\subseteq f^{-1} S$ 
proof -
    consider  $r \leq 0 \mid 0 < t * \text{norm}(\text{deriv } f a) / 12$  using rle by force
    then show ?thesis
    proof cases
    case 1 then show ?thesis
    by (simp add: ball_empty that)
    next
    case 2
    show ?thesis
    proof (cases deriv f a = 0)
    case True then show ?thesis
    using rle by (simp add: ball_empty that)
    next
    case False
    then have  $t > 0$ 
    using 2 by (force simp: zero_less_mult_iff)
    have  $\neg \text{ball } a t \subseteq S \implies \text{ball } a t \cap \text{frontier } S \neq \{\}$ 
    by (metis Diff_eq_empty_iff <0 < t> <a ∈ S> closure_Int_ball_not_empty
closure_subset connected_Int_frontier connected_ball inf.commute)
    with tle have *: ball a t  $\subseteq S$  by fastforce
    then have 1: f holomorphic_on ball a t
    using holf using holomorphic_on_subset by blast
    show ?thesis
    using Bloch [OF 1 <t > 0> rle] * by (metis image_mono order_trans that)
    qed
    qed
    qed
end

```

```

theory Complex_Singularities
  imports Conformal_Mappings
begin

```

5.10 Non-essential singular points

```

definition is_pole ::
  ('a::topological_space  $\Rightarrow$  'b::real_normed_vector)  $\Rightarrow$  'a  $\Rightarrow$  bool where
  is_pole f a = (LIM x (at a). f x  $:\>$  at_infinity)

```

```

lemma is_pole_cong:
  assumes eventually ( $\lambda x. f x = g x$ ) (at a) a=b
  shows is_pole f a  $\longleftrightarrow$  is_pole g b
  unfolding is_pole_def using assms by (intro filterlim_cong, auto)

```

```

lemma is_pole_transform:
  assumes is_pole f a eventually ( $\lambda x. f x = g x$ ) (at a) a=b
  shows is_pole g b
  using is_pole_cong assms by auto

```

```

lemma is_pole_shift_iff:
  fixes f :: ('a::real_normed_vector  $\Rightarrow$  'b::real_normed_vector)
  shows is_pole (f  $\circ$  (+) d) a = is_pole f (a + d)
  by (metis add_diff_cancel_right' filterlim_shift_iff is_pole_def)

```

```

lemma is_pole_tendsto:
  fixes f::('a::topological_space  $\Rightarrow$  'b::real_normed_div_algebra)
  shows is_pole f x  $\implies$  ((inverse o f)  $\longrightarrow$  0) (at x)
  unfolding is_pole_def
  by (auto simp add:filterlim_inverse_at_iff[symmetric] comp_def filterlim_at)

```

```

lemma is_pole_shift_0:
  fixes f :: ('a::real_normed_vector  $\Rightarrow$  'b::real_normed_vector)
  shows is_pole f z  $\longleftrightarrow$  is_pole ( $\lambda x. f (z + x)$ ) 0
  unfolding is_pole_def by (subst at_to_0) (auto simp: filterlim_filtermap add_ac)

```

```

lemma is_pole_shift_0':
  fixes f :: ('a::real_normed_vector  $\Rightarrow$  'b::real_normed_vector)
  shows NO_MATCH 0 z  $\implies$  is_pole f z  $\longleftrightarrow$  is_pole ( $\lambda x. f (z + x)$ ) 0
  by (metis is_pole_shift_0)

```

```

lemma is_pole_compose_iff:
  assumes filtermap g (at x) = (at y)
  shows is_pole (f  $\circ$  g) x  $\longleftrightarrow$  is_pole f y
  unfolding is_pole_def filterlim_def filtermap_compose assms ..

```

```

lemma is_pole_inverse_holomorphic:
  assumes open s
  and f_holo:f holomorphic_on (s- $\{z\}$ )

```

```

  and pole:is_pole f z
  and non_z:∀ x∈s-{z}. f x≠0
shows (λx. if x=z then 0 else inverse (f x)) holomorphic_on s
proof -
  define g where g ≡ λx. if x=z then 0 else inverse (f x)
  have isCont g z unfolding isCont_def using is_pole_tendsto[OF pole]
    by (simp add: g_def cong: LIM_cong)
  moreover have continuous_on (s-{z}) f using f_holo holomorphic_on_imp_continuous_on
  by auto
  hence continuous_on (s-{z}) (inverse o f) unfolding comp_def
    by (auto elim!:continuous_on_inverse simp add:non_z)
  hence continuous_on (s-{z}) g unfolding g_def
    using continuous_on_eq by fastforce
  ultimately have continuous_on s g using open_delete[OF ‹open s›] ‹open s›
    by (auto simp add:continuous_on_eq_continuous_at)
  moreover have (inverse o f) holomorphic_on (s-{z})
    unfolding comp_def using f_holo
    by (auto elim!:holomorphic_on_inverse simp add:non_z)
  hence g holomorphic_on (s-{z})
    using g_def holomorphic_transform by force
  ultimately show ?thesis unfolding g_def using ‹open s›
    by (auto elim!: no_isolated_singularity)
qed

```

lemma not_is_pole_holomorphic:

```

  assumes open A x ∈ A f holomorphic_on A
  shows ¬is_pole f x
proof -
  have continuous_on A f by (intro holomorphic_on_imp_continuous_on) fact
  with assms have isCont f x by (simp add: continuous_on_eq_continuous_at)
  hence f -x→ f x by (simp add: isCont_def)
  thus ¬is_pole f x unfolding is_pole_def
    using not_tendsto_and_filterlim_at_infinity[of at x f f x] by auto
qed

```

lemma is_pole_inverse_power: $n > 0 \implies is_pole (\lambda z::complex. 1 / (z - a) ^ n) a$

```

  unfolding is_pole_def inverse_eq_divide [symmetric]
  by (intro filterlim_compose[OF filterlim_inverse_at_infinity] tendsto_intros)
    (auto simp: filterlim_at_eventually_at intro!: exI[of _ 1] tendsto_eq_intros)

```

lemma is_pole_cmult_iff [simp]:

$c \neq 0 \implies is_pole (\lambda z. c * f z :: 'a :: real_normed_field) z \iff is_pole f z$

proof

```

  assume *: c ≠ 0 is_pole (λz. c * f z) z
  have is_pole (λz. inverse c * (c * f z)) z unfolding is_pole_def
    by (rule tendsto_mult_filterlim_at_infinity tendsto_const)+ (use * in ‹auto
  simp: is_pole_def›)
  thus is_pole f z

```

```

    using *(1) by (simp add: field_simps)
next
  assume *: c ≠ 0 is_pole f z
  show is_pole (λz. c * f z) z unfolding is_pole_def
    by (rule tendsto_mult_filterlim_at_infinity tendsto_const)+ (use * in ⟨auto
simp: is_pole_def⟩)
qed

```

```

lemma is_pole_uminus_iff [simp]: is_pole (λz. -f z :: 'a :: real_normed_field)
z ⟷ is_pole f z
  using is_pole_cmult_iff[of -1 f] by (simp del: is_pole_cmult_iff)

```

```

lemma is_pole_inverse: is_pole (λz::complex. 1 / (z - a)) a
  using is_pole_inverse_power[of 1 a] by simp

```

```

lemma is_pole_divide:
  fixes f :: 'a :: t2_space ⇒ 'b :: real_normed_field
  assumes isCont f z filterlim g (at 0) (at z) f z ≠ 0
  shows is_pole (λz. f z / g z) z
proof -
  have filterlim (λz. f z * inverse (g z)) at_infinity (at z)
    using assms filterlim_compose filterlim_inverse_at_infinity isCont_def
    tendsto_mult_filterlim_at_infinity by blast
  thus ?thesis by (simp add: field_split_simps is_pole_def)
qed

```

```

lemma is_pole_basic:
  assumes f holomorphic_on A open A z ∈ A f z ≠ 0 n > 0
  shows is_pole (λw. f w / (w - z) ^ n) z
proof (rule is_pole_divide)
  have continuous_on A f by (rule holomorphic_on_imp_continuous_on) fact
  with assms show isCont f z by (auto simp: continuous_on_eq_continuous_at)
  have filterlim (λw. (w - z) ^ n) (nhds 0) (at z)
    using assms by (auto intro!: tendsto_eq_intros)
  thus filterlim (λw. (w - z) ^ n) (at 0) (at z)
    by (intro filterlim_atI tendsto_eq_intros)
    (insert assms, auto simp: eventually_at_filter)
qed fact+

```

```

lemma is_pole_basic':
  assumes f holomorphic_on A open A 0 ∈ A f 0 ≠ 0 n > 0
  shows is_pole (λw. f w / w ^ n) 0
  using is_pole_basic[of f A 0] assms by simp

```

```

lemma is_pole_compose:
  assumes is_pole f w g -z → w eventually (λz. g z ≠ w) (at z)
  shows is_pole (λx. f (g x)) z
  using assms(1) unfolding is_pole_def
  by (rule filterlim_compose) (use assms in ⟨auto simp: filterlim_at⟩)

```

```

lemma is_pole_plus_const_iff:
  is_pole f z  $\longleftrightarrow$  is_pole ( $\lambda x. f\ x + c$ ) z
proof
  assume is_pole f z
  then have filterlim f at_infinity (at z) unfolding is_pole_def .
  moreover have ( $\lambda \_. c$ )  $\longrightarrow$  c (at z) by auto
  ultimately have LIM x (at z). f x + c  $:$  at_infinity
    using tendsto_add_filterlim_at_infinity'[of f at z] by auto
  then show is_pole ( $\lambda x. f\ x + c$ ) z unfolding is_pole_def .
next
  assume is_pole ( $\lambda x. f\ x + c$ ) z
  then have filterlim ( $\lambda x. f\ x + c$ ) at_infinity (at z)
    unfolding is_pole_def .
  moreover have ( $\lambda \_. -c$ )  $\longrightarrow$   $-c$  (at z) by auto
  ultimately have LIM x (at z). f x  $:$  at_infinity
    using tendsto_add_filterlim_at_infinity'[of ( $\lambda x. f\ x + c$ )
      at z ( $\lambda \_. -c$ )  $-c$ ]
    by auto
  then show is_pole f z unfolding is_pole_def .
qed

lemma is_pole_minus_const_iff:
  is_pole ( $\lambda x. f\ x - c$ ) z  $\longleftrightarrow$  is_pole f z
  using is_pole_plus_const_iff [of f z  $-c$ ] by simp

lemma is_pole_alt:
  is_pole f x = ( $\forall B > 0. \exists U. \text{open } U \wedge x \in U \wedge (\forall y \in U. y \neq x \longrightarrow \text{norm } (f\ y) \geq B)$ )
  unfolding is_pole_def
  unfolding filterlim_at_infinity[of 0, simplified] eventually_at_topological
  by auto

lemma is_pole_mult_analytic_nonzero1:
  assumes is_pole g x f analytic_on {x} f x  $\neq 0$ 
  shows is_pole ( $\lambda x. f\ x * g\ x$ ) x
  unfolding is_pole_def
proof (rule tendsto_mult_filterlim_at_infinity)
  show f  $-x \rightarrow f\ x$ 
  using assms by (simp add: analytic_at_imp_isCont isContD)
qed (use assms in  $\langle$ auto simp: is_pole_def $\rangle$ )

lemma is_pole_mult_analytic_nonzero2:
  assumes is_pole f x g analytic_on {x} g x  $\neq 0$ 
  shows is_pole ( $\lambda x. f\ x * g\ x$ ) x
  by (subst mult.commute, rule is_pole_mult_analytic_nonzero1) (use assms in
auto)

lemma is_pole_mult_analytic_nonzero1_iff:
  assumes f analytic_on {x} f x  $\neq 0$ 

```



```

shows is_pole ( $\lambda x. f x * g x$ )  $x \longleftrightarrow$  is_pole  $g x$ 
proof
  assume is_pole  $g x$ 
  thus is_pole ( $\lambda x. f x * g x$ )  $x$ 
    by (intro is_pole_mult_analytic_nonzero1 assms)
next
  assume is_pole ( $\lambda x. f x * g x$ )  $x$ 
  hence is_pole ( $\lambda x. \text{inverse } (f x) * (f x * g x)$ )  $x$ 
    by (rule is_pole_mult_analytic_nonzero1)
    (use assms in <auto intro!: analytic_intros>)
  also have ?this  $\longleftrightarrow$  is_pole  $g x$ 
  proof (rule is_pole_cong)
    have eventually ( $\lambda x. f x \neq 0$ ) (at x)
      using assms by (simp add: analytic_at_neq_imp_eventually_neq)
    thus eventually ( $\lambda x. \text{inverse } (f x) * (f x * g x) = g x$ ) (at x)
      by eventually_elim auto
    qed auto
  finally show is_pole  $g x$  .
qed

```

```

lemma is_pole_mult_analytic_nonzero2_iff:
  assumes  $g$  analytic_on  $\{x\}$   $g x \neq 0$ 
  shows is_pole ( $\lambda x. f x * g x$ )  $x \longleftrightarrow$  is_pole  $f x$ 
  by (subst mult.commute, rule is_pole_mult_analytic_nonzero1_iff) (fact assms)+

```

```

lemma frequently_const_imp_not_is_pole:
  fixes  $z :: 'a::\text{first\_countable\_topology}$ 
  assumes frequently ( $\lambda w. f w = c$ ) (at z)
  shows  $\neg$  is_pole  $f z$ 
proof
  assume is_pole  $f z$ 
  from assms have  $z$  islimpt  $\{w. f w = c\}$ 
    by (simp add: islimpt_conv_frequently_at)
  then obtain  $g$  where  $g: \bigwedge n. g n \in \{w. f w = c\} - \{z\}$   $g \longrightarrow z$ 
    unfolding islimpt_sequential by blast
  then have  $(f \circ g) \longrightarrow c$ 
    by (simp add: tendsto_eventually)
  moreover have  $*$ : filterlim  $g$  (at z) sequentially
    using  $g$  by (auto simp: filterlim_at)
  have filterlim  $(f \circ g)$  at_infinity sequentially
    unfolding o_def by (rule filterlim_compose [OF _ *])
    (use <is_pole f z> in <simp add: is_pole_def>)
  ultimately show False
    using not_tendsto_and_filterlim_at_infinity_trivial_limit_sequentially by
blast
qed

```

The proposition $\exists x. f -z \rightarrow x \vee \text{is_pole } f z$ can be interpreted as the complex function f has a non-essential singularity at z (i.e. the singularity is

either removable or a pole).

definition *not_essential*::[*complex* \Rightarrow *complex*, *complex*] \Rightarrow *bool* **where**
not_essential *f* *z* = ($\exists x. f - z \rightarrow x \vee \text{is_pole } f z$)

definition *isolated_singularity_at*::[*complex* \Rightarrow *complex*, *complex*] \Rightarrow *bool* **where**
isolated_singularity_at *f* *z* = ($\exists r > 0. f \text{ analytic_on ball } z r - \{z\}$)

lemma *not_essential_cong*:

assumes *eventually* ($\lambda x. f x = g x$) (*at* *z*) *z* = *z'*

shows *not_essential* *f* *z* \longleftrightarrow *not_essential* *g* *z'*

unfolding *not_essential_def* **using** *assms filterlim_cong is_pole_cong* **by** *fast-force*

lemma *not_essential_compose_iff*:

assumes *filtermap* *g* (*at* *z*) = *at* *z'*

shows *not_essential* (*f* \circ *g*) *z* = *not_essential* *f* *z'*

unfolding *not_essential_def filterlim_def filtermap_compose assms is_pole_compose_iff* [*OF assms*]

by *blast*

lemma *isolated_singularity_at_cong*:

assumes *eventually* ($\lambda x. f x = g x$) (*at* *z*) *z* = *z'*

shows *isolated_singularity_at* *f* *z* \longleftrightarrow *isolated_singularity_at* *g* *z'*

proof –

have *isolated_singularity_at* *g* *z*

if *isolated_singularity_at* *f* *z* *eventually* ($\lambda x. f x = g x$) (*at* *z*) **for** *f* *g*

proof –

from *that(1)* **obtain** *r* **where** *r*: *r* > 0 *f* *analytic_on* *ball* *z* *r* – {*z*}

by (*auto simp: isolated_singularity_at_def*)

from *that(2)* **obtain** *r'* **where** *r'*: *r'* > 0 $\forall x \in \text{ball } z r' - \{z\}. f x = g x$

unfolding *eventually_at_filter eventually_nhds_metric* **by** (*auto simp: dist_commute*)

have *f* *holomorphic_on* *ball* *z* *r* – {*z*}

using *r(2)* **by** (*subst (asm) analytic_on_open*) *auto*

hence *f* *holomorphic_on* *ball* *z* (*min* *r* *r'*) – {*z*}

by (*rule holomorphic_on_subset*) *auto*

also have *?this* \longleftrightarrow *g* *holomorphic_on* *ball* *z* (*min* *r* *r'*) – {*z*}

using *r'* **by** (*intro holomorphic_cong*) *auto*

also have ... \longleftrightarrow *g* *analytic_on* *ball* *z* (*min* *r* *r'*) – {*z*}

by (*subst analytic_on_open*) *auto*

finally show *?thesis*

unfolding *isolated_singularity_at_def*

by (*intro exI*[*of* _ *min* *r* *r'*]) (*use* $\langle r > 0 \rangle \langle r' > 0 \rangle$) **in** *auto*)

qed

from *this*[*of* *f* *g*] *this*[*of* *g* *f*] *assms* **show** *?thesis*

by (*auto simp: eq_commute*)

qed

```

lemma removable_singularity:
  assumes  $f$  holomorphic_on  $A - \{x\}$  open  $A$ 
  assumes  $f -x \rightarrow c$ 
  shows  $(\lambda y. \text{if } y = x \text{ then } c \text{ else } f y)$  holomorphic_on  $A$  (is  $?g$  holomorphic_on
  _)
proof -
  have continuous_on  $A$   $?g$ 
    unfolding continuous_on_def
  proof
    fix  $y$  assume  $y \in A$ 
    show  $(?g \longrightarrow ?g y)$  (at  $y$  within  $A$ )
    proof (cases  $y = x$ )
      case False
        have continuous_on  $(A - \{x\})$   $f$ 
          using assms(1) by (meson holomorphic_on_imp_continuous_on)
        hence  $(f \longrightarrow ?g y)$  (at  $y$  within  $A - \{x\}$ )
          using False by (auto simp: continuous_on_def)
        also have  $?this \longleftrightarrow (?g \longrightarrow ?g y)$  (at  $y$  within  $A - \{x\}$ )
          by (intro filterlim_cong refl) (auto simp: eventually_at_filter)
        also have at  $y$  within  $A - \{x\} = \text{at } y \text{ within } A$ 
          using  $y$  assms False
          by (intro at_within_nhd[of _ A - \{x\}]) auto
        finally show  $?thesis$  .
      next
        case [simp]: True
          have  $f -x \rightarrow c$ 
            by fact
          also have  $?this \longleftrightarrow (?g \longrightarrow ?g y)$  (at  $y$ )
            by (intro filterlim_cong) (auto simp: eventually_at_filter)
          finally show  $?thesis$ 
            by (meson Lim_at_imp_Lim_at_within)
    qed
  qed
  moreover {
    have  $?g$  holomorphic_on  $A - \{x\}$ 
      using assms(1) holomorphic_transform by fastforce
    }
  ultimately show  $?thesis$ 
    by (rule no_isolated_singularity) (use assms in auto)
qed

lemma removable_singularity':
  assumes isolated_singularity_at  $f$   $z$ 
  assumes  $f -z \rightarrow c$ 
  shows  $(\lambda y. \text{if } y = z \text{ then } c \text{ else } f y)$  analytic_on  $\{z\}$ 
proof -
  from assms obtain  $r$  where  $r: r > 0$  analytic_on ball  $z$   $r - \{z\}$ 
    by (auto simp: isolated_singularity_at_def)
  have  $(\lambda y. \text{if } y = z \text{ then } c \text{ else } f y)$  holomorphic_on ball  $z$   $r$ 

```

```

proof (rule removable_singularity)
  show  $f$  holomorphic_on ball  $z$   $r - \{z\}$ 
    using  $r(2)$  by (subst (asm) analytic_on_open) auto
qed (use assms in auto)
thus ?thesis
  using  $r(1)$  unfolding analytic_at by (intro exI[of _ ball  $z$   $r$ ]) auto
qed

```

named_theorems singularity_intros introduction rules for singularities

lemma holomorphic_factor_unique:

```

fixes  $f::\text{complex} \Rightarrow \text{complex}$  and  $z::\text{complex}$  and  $r::\text{real}$  and  $m\ n::\text{int}$ 
assumes  $r>0$   $g\ z\neq 0$   $h\ z\neq 0$ 
  and  $asm:\forall w\in\text{ball } z\ r-\{z\}. f\ w = g\ w * (w-z)^{\text{powi } n} \wedge g\ w\neq 0 \wedge f\ w = h\ w$ 
  *  $(w-z)^{\text{powi } m} \wedge h\ w\neq 0$ 
  and  $g\_holo:g$  holomorphic_on ball  $z$   $r$  and  $h\_holo:h$  holomorphic_on ball  $z$   $r$ 
shows  $n=m$ 
proof -
  have [simp]:at  $z$  within ball  $z$   $r \neq \text{bot}$  using  $\langle r>0 \rangle$ 
    by (auto simp add:at_within_ball_bot_iff)
  have False when  $n>m$ 
proof -
  have  $(h \longrightarrow 0)$  (at  $z$  within ball  $z$   $r$ )
proof (rule Lim_transform_within[OF _  $\langle r>0 \rangle$ , where  $f=\lambda w. (w-z)^{\text{powi } (n-m)} * g\ w$ 
  ( $n-m$ ) *  $g\ w$ ])
  have  $\forall w\in\text{ball } z\ r-\{z\}. h\ w = (w-z)^{\text{powi } (n-m)} * g\ w$ 
    using  $\langle n>m \rangle$   $asm\ \langle r>0 \rangle$  by (simp add: field_simps power_int_diff) force
  then show  $\llbracket x' \in \text{ball } z\ r; 0 < \text{dist } x'\ z; \text{dist } x'\ z < r \rrbracket$ 
     $\implies (x' - z)^{\text{powi } (n-m)} * g\ x' = h\ x'$  for  $x'$  by auto
next
define  $F$  where  $F \equiv \text{at } z \text{ within ball } z\ r$ 
define  $f'$  where  $f' \equiv \lambda x. (x - z)^{\text{powi } (n-m)}$ 
have  $f'\ z=0$  using  $\langle n>m \rangle$  unfolding  $f'_\text{def}$  by auto
moreover have continuous  $F\ f'$  unfolding  $f'_\text{def}\ F_\text{def}$  continuous_def
  using  $\langle n>m \rangle$ 
  by (auto simp add: Lim_ident_at intro!:tendsto_powr_complex_0 tendsto_eq_intros)
ultimately have  $(f' \longrightarrow 0)\ F$  unfolding  $F_\text{def}$ 
  by (simp add: continuous_within)
moreover have  $(g \longrightarrow g\ z)\ F$ 
  unfolding  $F_\text{def}$ 
using  $\langle r>0 \rangle$  centre_in_ball continuous_on_def  $g\_holo$  holomorphic_on_imp_continuous_on
by blast
ultimately show  $((\lambda w. f'\ w * g\ w) \longrightarrow 0)\ F$  using tendsto_mult by
fastforce
qed
moreover have  $(h \longrightarrow h\ z)$  (at  $z$  within ball  $z$   $r$ )
  using holomorphic_on_imp_continuous_on[OF  $h\_holo$ ]
  by (auto simp add:continuous_on_def  $\langle r>0 \rangle$ )

```

```

ultimately have  $h z=0$  by (auto intro!: tendsto_unique)
thus False using  $\langle h z \neq 0 \rangle$  by auto
qed
moreover have False when  $m > n$ 
proof -
  have  $(g \longrightarrow 0)$  (at  $z$  within ball  $z r$ )
  proof (rule Lim_transform_within[OF  $\langle r > 0 \rangle$ ], where  $f = \lambda w. (w - z) \text{ powi } (m - n) * h w$ )
    have  $\forall w \in \text{ball } z r - \{z\}. g w = (w - z) \text{ powi } (m - n) * h w$  using  $\langle m > n \rangle$  asm
    by (simp add: field_simps power_int_diff) force
    then show  $\llbracket x' \in \text{ball } z r; 0 < \text{dist } x' z; \text{dist } x' z < r \rrbracket$ 
       $\implies (x' - z) \text{ powi } (m - n) * h x' = g x'$  for  $x'$  by auto
  next
  define  $F$  where  $F \equiv \text{at } z \text{ within ball } z r$ 
  define  $f'$  where  $f' \equiv \lambda x. (x - z) \text{ powi } (m - n)$ 
  have  $f' z = 0$  using  $\langle m > n \rangle$  unfolding  $f'_\text{def}$  by auto
  moreover have continuous  $F f'$  unfolding  $f'_\text{def} F_\text{def}$  continuous_def
    using  $\langle m > n \rangle$ 
  by (auto simp: Lim_ident_at intro!: tendsto_powr_complex_0 tendsto_eq_intros)
  ultimately have  $(f' \longrightarrow 0) F$  unfolding  $F_\text{def}$ 
    by (simp add: continuous_within)
  moreover have  $(h \longrightarrow h z) F$ 
    using holomorphic_on_imp_continuous_on[OF  $h_\text{holo}$ , unfolded continuous_on_on_def]  $\langle r > 0 \rangle$ 
    unfolding  $F_\text{def}$  by auto
  ultimately show  $((\lambda w. f' w * h w) \longrightarrow 0) F$  using tendsto_mult by
fastforce
qed
moreover have  $(g \longrightarrow g z)$  (at  $z$  within ball  $z r$ )
  using holomorphic_on_imp_continuous_on[OF  $g_\text{holo}$ ]
  by (auto simp add: continuous_on_def  $\langle r > 0 \rangle$ )
ultimately have  $g z = 0$  by (auto intro!: tendsto_unique)
thus False using  $\langle g z \neq 0 \rangle$  by auto
qed
ultimately show  $n = m$  by fastforce
qed

```

lemma holomorphic_factor_puncture:

```

assumes  $f_\text{iso:isolated_singularity\_at } f z$ 
  and  $\text{not\_essential } f z$  —  $f$  has either a removable singularity or a pole at  $z$ 
  and  $\text{non\_zero: } \exists F w \text{ in } (\text{at } z). f w \neq 0$  —  $f$  will not be constantly zero in a
neighbour of  $z$ 
shows  $\exists ! n :: \text{int}. \exists g r. 0 < r \wedge g \text{ holomorphic\_on } \text{cball } z r \wedge g z \neq 0$ 
   $\wedge (\forall w \in \text{cball } z r - \{z\}. f w = g w * (w - z) \text{ powi } n \wedge g w \neq 0)$ 
proof -
  define  $P$  where  $P = (\lambda f n g r. 0 < r \wedge g \text{ holomorphic\_on } \text{cball } z r \wedge g z \neq 0$ 
   $\wedge (\forall w \in \text{cball } z r - \{z\}. f w = g w * (w - z) \text{ powi } n \wedge g w \neq 0))$ 
  have  $\text{imp\_unique: } \exists ! n :: \text{int}. \exists g r. P f n g r$  when  $\exists n g r. P f n g r$ 
  proof (rule ex_ex1I[OF that])

```

```

fix n1 n2 :: int
assume g1_asm:  $\exists g1\ r1. P\ f\ n1\ g1\ r1$  and g2_asm:  $\exists g2\ r2. P\ f\ n2\ g2\ r2$ 
define fac where fac  $\equiv \lambda n\ g\ r. \forall w \in \text{cball}\ z\ r - \{z\}. f\ w = g\ w * (w - z)\ \text{powi}\ n \wedge g\ w \neq 0$ 
obtain g1 r1 where  $0 < r1$  and g1_holo: g1 holomorphic_on cball z r1 and
g1 z $\neq 0$ 
and fac n1 g1 r1 using g1_asm unfolding P_def fac_def by auto
obtain g2 r2 where  $0 < r2$  and g2_holo: g2 holomorphic_on cball z r2 and
g2 z $\neq 0$ 
and fac n2 g2 r2 using g2_asm unfolding P_def fac_def by auto
define r where r  $\equiv \min\ r1\ r2$ 
have r>0 using <r1>0> <r2>0> unfolding r_def by auto
moreover have  $\forall w \in \text{ball}\ z\ r - \{z\}. f\ w = g1\ w * (w - z)\ \text{powi}\ n1 \wedge g1\ w \neq 0$ 
 $\wedge f\ w = g2\ w * (w - z)\ \text{powi}\ n2 \wedge g2\ w \neq 0$ 
using <fac n1 g1 r1> <fac n2 g2 r2> unfolding fac_def r_def
by fastforce
ultimately show n1=n2 using g1_holo g2_holo <g1 z $\neq 0$ > <g2 z $\neq 0$ >
apply (elim holomorphic_factor_unique)
by (auto simp add:r_def)
qed

have P_exist:  $\exists n\ g\ r. P\ h\ n\ g\ r$  when
 $\exists z'. (h \longrightarrow z')\ (at\ z)\ \text{isolated\_singularity\_at}\ h\ z\ \exists_F w\ \text{in}\ (at\ z). h\ w \neq 0$ 
for h
proof -
from that(2) obtain r where r>0 and r: h analytic_on ball z r - {z}
unfolding isolated_singularity_at_def by auto
obtain z' where (h  $\longrightarrow$  z') (at z) using < $\exists z'. (h \longrightarrow z')\ (at\ z)$ > by auto
define h' where h' = ( $\lambda x. \text{if}\ x = z\ \text{then}\ z'\ \text{else}\ h\ x$ )
have h' holomorphic_on ball z r
proof (rule no_isolated_singularity'[of {z}])
show  $\bigwedge w. w \in \{z\} \implies (h' \longrightarrow h' w)\ (at\ w\ \text{within}\ \text{ball}\ z\ r)$ 
by (simp add: LIM_cong Lim_at_imp_Lim_at_within <h - z  $\rightarrow$  z'> h'_def)
show h' holomorphic_on ball z r - {z}
using r analytic_imp_holomorphic h'_def holomorphic_transform by
fastforce
qed auto
have ?thesis when z'=0
proof -
have h' z=0 using that unfolding h'_def by auto
moreover have  $\neg h'$  constant_on ball z r
using < $\exists_F w\ \text{in}\ (at\ z). h\ w \neq 0$ > unfolding constant_on_def frequently_def
eventually_at h'_def
by (metis <0 < r> centre_in_ball dist_commute mem_ball that)
moreover note <h' holomorphic_on ball z r>
ultimately obtain g r1 n where  $0 < n\ 0 < r1\ \text{ball}\ z\ r1 \subseteq \text{ball}\ z\ r$  and
g: g holomorphic_on ball z r1
 $\bigwedge w. w \in \text{ball}\ z\ r1 \implies h' w = (w - z) \wedge^n * g w$ 
 $\bigwedge w. w \in \text{ball}\ z\ r1 \implies g w \neq 0$ 

```

```

    using holomorphic_factor_zero_nonconstant[of _ ball z r z thesis,simplified,
      OF ‹h' holomorphic_on ball z r› ‹r>0› ‹h' z=0› ‹¬ h' constant_on
ball z r›]
    by (auto simp add:dist_commute)
  define rr where rr=r1/2
  have P h' n g rr
    unfolding P_def rr_def
    using ‹n>0› ‹r1>0› g by (auto simp add:powr_nat)
  then have P h n g rr
    unfolding h'_def P_def by auto
  then show ?thesis unfolding P_def by blast
qed
moreover have ?thesis when z'≠0
proof -
  have h' z≠0 using that unfolding h'_def by auto
  obtain r1 where r1>0 cball z r1 ⊆ ball z r ∀ x∈cball z r1. h' x≠0
  proof -
    have isCont h' z h' z≠0
      by (auto simp add: Lim_cong_within ‹h -z→ z'› ‹z'≠0› continuous_at
h'_def)
    then obtain r2 where r2:r2>0 ∀ x∈ball z r2. h' x≠0
      using continuous_at_avoid[of z h' 0 ] unfolding ball_def by auto
    define r1 where r1=min r2 r / 2
    have 0 < r1 cball z r1 ⊆ ball z r
      using ‹r2>0› ‹r>0› unfolding r1_def by auto
    moreover have ∀ x∈cball z r1. h' x ≠ 0
      using r2 unfolding r1_def by simp
    ultimately show ?thesis using that by auto
  qed
  then have P h' 0 h' r1 using ‹h' holomorphic_on ball z r› unfolding P_def
by auto
  then have P h 0 h' r1 unfolding P_def h'_def by auto
  then show ?thesis unfolding P_def by blast
qed
ultimately show ?thesis by auto
qed

have ?thesis when ∃ x. (f ⟶ x) (at z)
  apply (rule_tac imp_unique[unfolded P_def])
  using P_exist[OF that(1) f_iso non_zero] unfolding P_def .
moreover have ?thesis when is_pole f z
proof (rule imp_unique[unfolded P_def])
  obtain e where [simp]:e>0 and e_holo:f holomorphic_on ball z e - {z} and
e_nz: ∀ x∈ball z e-{z}. f x≠0
  proof -
    have ∀_F z in at z. f z ≠ 0
      using ‹is_pole f z› filterlim_at_infinity_imp_eventually_ne unfolding
is_pole_def
      by auto

```

```

then obtain e1 where e1:e1>0  $\forall x \in \text{ball } z \ e1 - \{z\}. f \ x \neq 0$ 
  using that eventually_at[of  $\lambda x. f \ x \neq 0 \ z \ \text{UNIV}, \text{simplified}$ ] by (auto simp
add:dist_commute)
obtain e2 where e2:e2>0  $f \ \text{holomorphic\_on } \text{ball } z \ e2 - \{z\}$ 
using f_iso analytic_imp_holomorphic unfolding isolated_singularity_at_def
by auto
show ?thesis
  using e1 e2 by (force intro: that[of min e1 e2])
qed

define h where h  $\equiv \lambda x. \text{inverse } (f \ x)$ 
have  $\exists n \ g \ r. P \ h \ n \ g \ r$ 
proof -
  have  $(\lambda x. \text{inverse } (f \ x)) \ \text{analytic\_on } \text{ball } z \ e - \{z\}$ 
  by (metis e_holo e_nz open_ball analytic_on_open holomorphic_on_inverse
open_delete)
  moreover have  $h \ -z \rightarrow 0$ 
  using Lim_transform_within_open assms(2) h_def is_pole_tendsto that
by fastforce
  moreover have  $\exists_F w \ \text{in } (\text{at } z). h \ w \neq 0$ 
  using non_zero by (simp add: h_def)
  ultimately show ?thesis
  using P_exist[of h]  $\langle e > 0 \rangle$ 
  unfolding isolated_singularity_at_def h_def
  by blast
qed
then obtain n g r
  where  $0 < r$  and
    g_holo:g holomorphic_on cball z r and  $g \ z \neq 0$  and
    g_fac: $(\forall w \in \text{cball } z \ r - \{z\}. h \ w = g \ w * (w - z) \ \text{powi } n \ \wedge \ g \ w \neq 0)$ 
  unfolding P_def by auto
have  $P \ f \ (-n) \ (\text{inverse } o \ g) \ r$ 
proof -
  have  $f \ w = \text{inverse } (g \ w) * (w - z) \ \text{powi } (-n)$  when  $w \in \text{cball } z \ r - \{z\}$  for
w
  by (metis g_fac h_def inverse_inverse_eq inverse_mult_distrib power_int_minus
that)
  then show ?thesis
  unfolding P_def comp_def
  using  $\langle r > 0 \rangle$  g_holo g_fac  $\langle g \ z \neq 0 \rangle$  by (auto intro:holomorphic_intros)
qed
then show  $\exists x \ g \ r. 0 < r \ \wedge \ g \ \text{holomorphic\_on } \text{cball } z \ r \ \wedge \ g \ z \neq 0$ 
 $\wedge \ (\forall w \in \text{cball } z \ r - \{z\}. f \ w = g \ w * (w - z) \ \text{powi } x \ \wedge \ g \ w \neq 0)$ 
  unfolding P_def by blast
qed
ultimately show ?thesis using  $\langle \text{not\_essential } f \ z \rangle$  unfolding not_essential_def
by presburger
qed

```


lemma *not_essential_transform*:

assumes *not_essential* $g\ z$
assumes $\forall_F w$ *in* $(at\ z)$. $g\ w = f\ w$
shows *not_essential* $f\ z$
using *assms* **unfolding** *not_essential_def*
by (*simp* *add*: *filterlim_cong* *is_pole_cong*)

lemma *isolated_singularity_at_transform*:

assumes *isolated_singularity_at* $g\ z$
assumes $\forall_F w$ *in* $(at\ z)$. $g\ w = f\ w$
shows *isolated_singularity_at* $f\ z$

proof –

obtain $r1$ **where** $r1 > 0$ **and** $r1: g$ *analytic_on* *ball* $z\ r1 - \{z\}$
using *assms*(1) **unfolding** *isolated_singularity_at_def* **by** *auto*
obtain $r2$ **where** $r2 > 0$ **and** $r2: \forall x. x \neq z \wedge dist\ x\ z < r2 \longrightarrow g\ x = f\ x$
using *assms*(2) **unfolding** *eventually_at* **by** *auto*
define $r3$ **where** $r3 = \min\ r1\ r2$
have $r3 > 0$ **unfolding** $r3_def$ **using** $\langle r1 > 0 \rangle\ \langle r2 > 0 \rangle$ **by** *auto*
moreover **have** f *analytic_on* *ball* $z\ r3 - \{z\}$

proof –

have g *holomorphic_on* *ball* $z\ r3 - \{z\}$
using $r1$ **unfolding** $r3_def$ **by** (*subst* (*asm*) *analytic_on_open*, *auto*)
then **have** f *holomorphic_on* *ball* $z\ r3 - \{z\}$
using $r2$ **unfolding** $r3_def$
by (*auto* *simp* *add*: *dist_commute* *elim*!: *holomorphic_transform*)
then **show** *?thesis* **by** (*subst* *analytic_on_open*, *auto*)

qed

ultimately **show** *?thesis* **unfolding** *isolated_singularity_at_def* **by** *auto*

qed

lemma *not_essential_pour*[*singularity_intros*]:

assumes *LIM* w $(at\ z)$. $f\ w :> (at\ x)$
shows *not_essential* $(\lambda w. (f\ w)\ powi\ n)\ z$

proof –

define fp **where** $fp = (\lambda w. (f\ w)\ powi\ n)$
have *?thesis* **when** $n > 0$

proof –

have $(\lambda w. (f\ w)\ ^{(nat\ n)}) -z \rightarrow x\ ^{nat\ n}$
using *that* *assms* **unfolding** *filterlim_at* **by** (*auto* *intro*!: *tendsto_eq_intros*)
then **have** $fp -z \rightarrow x\ ^{nat\ n}$ **unfolding** fp_def
by (*smt* (*verit*) *LIM_cong* *power_int_def* *that*)
then **show** *?thesis* **unfolding** *not_essential_def* fp_def **by** *auto*

qed

moreover **have** *?thesis* **when** $n = 0$

proof –

have $\forall_F x$ *in* $at\ z$. $fp\ x = 1$
using *that* *filterlim_at_within_not_equal*[*OF* *assms*] **by** (*auto* *simp*: fp_def)
then **have** $fp -z \rightarrow 1$
by (*simp* *add*: *tendsto_eventually*)

```

  then show ?thesis unfolding fp_def not_essential_def by auto
qed
moreover have ?thesis when n<0
proof (cases x=0)
  case True
  have  $(\lambda x. f x \wedge \text{nat } (-n)) -z \rightarrow 0$ 
  using assms True that unfolding filterlim_at by (auto intro!: tendsto_eq_intros)
  moreover have  $\forall_F x \text{ in } \text{at } z. f x \wedge \text{nat } (-n) \neq 0$ 
  by (smt (verit) True assms eventually_at_topological filterlim_at power_eq_0_iff)
  ultimately have LIM w (at z). inverse ((f w)  $\wedge$  (nat (-n)))  $\rightarrow$  at_infinity
  by (metis filterlim_atI filterlim_compose filterlim_inverse_at_infinity)
  then have LIM w (at z). fp w  $\rightarrow$  at_infinity
  proof (elim filterlim_mono_eventually)
    show  $\forall_F x \text{ in } \text{at } z. \text{inverse } (f x \wedge \text{nat } (-n)) = \text{fp } x$ 
    using filterlim_at_within_not_equal[OF assms, of 0] unfolding fp_def
    by (smt (verit) eventuallyI power_int_def power_inverse that)
  qed auto
  then show ?thesis unfolding fp_def not_essential_def is_pole_def by auto
next
  case False
  let ?xx= inverse (x  $\wedge$  (nat (-n)))
  have  $(\lambda w. \text{inverse } ((f w) \wedge \text{nat } (-n))) -z \rightarrow ?xx$ 
  using assms False unfolding filterlim_at by (auto intro!: tendsto_eq_intros)
  then have fp  $-z \rightarrow ?xx$ 
  by (smt (verit, best) LIM_cong fp_def power_int_def power_inverse that)
  then show ?thesis unfolding fp_def not_essential_def by auto
qed
ultimately show ?thesis by linarith
qed

lemma isolated_singularity_at_powr[singularity_intros]:
  assumes isolated_singularity_at f z  $\forall_F w \text{ in } \text{at } z. f w \neq 0$ 
  shows isolated_singularity_at  $(\lambda w. (f w) \text{ powi } n) z$ 
proof -
  obtain r1 where r1>0 f analytic_on ball z r1 - {z}
  using assms(1) unfolding isolated_singularity_at_def by auto
  then have r1:f holomorphic_on ball z r1 - {z}
  using analytic_on_open[of ball z r1 - {z} f] by blast
  obtain r2 where r2>0 and r2: $\forall w. w \neq z \wedge \text{dist } w z < r2 \rightarrow f w \neq 0$ 
  using assms(2) unfolding eventually_at by auto
  define r3 where r3=min r1 r2
  have  $(\lambda w. (f w) \text{ powi } n) \text{ holomorphic\_on ball } z r3 - \{z\}$ 
  by (intro holomorphic_on_power_int) (use r1 r2 in <auto simp: dist_commute r3_def>)
  moreover have r3>0 unfolding r3_def using <0 < r1> <0 < r2> by linarith
  ultimately show ?thesis
  by (meson open_ball analytic_on_open isolated_singularity_at_def open_delete)
qed

```

```

lemma non_zero_neighbour:
  assumes f_iso:isolated_singularity_at f z
    and f_ness:not_essential f z
    and f_nconst: $\exists_F w$  in (at z). f w  $\neq 0$ 
  shows  $\forall_F w$  in (at z). f w  $\neq 0$ 
proof -
  obtain fn fp fr where [simp]:fp z  $\neq 0$  and fr > 0
    and fr: fp holomorphic_on cball z fr
       $\forall w \in \text{cball } z \text{ fr} - \{z\}. f w = fp w * (w - z)^{powi fn} \wedge fp w \neq 0$ 
  using holomorphic_factor_puncture[OF f_iso f_ness f_nconst, THEN ex1_implies_ex]
by auto
  have f w  $\neq 0$  when w  $\neq z$  dist w z < fr for w
  proof -
    have f w = fp w * (w - z)^{powi fn} fp w  $\neq 0$ 
      using fr(2)[rule_format, of w] using that by (auto simp add:dist_commute)
    moreover have (w - z)^{powi fn}  $\neq 0$ 
      unfolding powr_eq_0_iff using  $\langle w \neq z \rangle$  by auto
    ultimately show ?thesis by auto
  qed
  then show ?thesis using  $\langle fr > 0 \rangle$  unfolding eventually_at by auto
qed

lemma non_zero_neighbour_pole:
  assumes is_pole f z
  shows  $\forall_F w$  in (at z). f w  $\neq 0$ 
  using assms filterlim_at_infinity_imp_eventually_ne[of f at z 0]
  unfolding is_pole_def by auto

lemma non_zero_neighbour_alt:
  assumes holo: f holomorphic_on S
    and open S connected S z  $\in S$   $\beta \in S$  f  $\beta \neq 0$ 
  shows  $\forall_F w$  in (at z). f w  $\neq 0$   $\wedge w \in S$ 
proof (cases f z = 0)
  case True
    from isolated_zeros[OF holo  $\langle$ open S $\rangle$   $\langle$ connected S $\rangle$   $\langle$ z  $\in S$  $\rangle$  True  $\langle$  $\beta \in S$  $\rangle$   $\langle$ f  $\beta \neq 0$  $\rangle$ ]
    obtain r where 0 < r ball z r  $\subseteq S$   $\forall w \in \text{ball } z \text{ r} - \{z\}. f w \neq 0$  by metis
    then show ?thesis
      by (smt (verit) open_ball centre_in_ball eventually_at_topological insertE insert_Diff subsetD)
  next
  case False
    obtain r1 where r1:r1>0  $\forall y. \text{dist } z \ y < r1 \longrightarrow f y \neq 0$ 
      using continuous_at_avoid[of z f, OF _ False] assms(2,4) continuous_on_eq_continuous_at
        holo holomorphic_on_imp_continuous_on by blast
    obtain r2 where r2:r2>0 ball z r2  $\subseteq S$ 
      using assms openE by blast
    show ?thesis unfolding eventually_at
      by (metis (no_types) dist_commute dual_order.strict_trans linorder_less_linear)

```

mem_ball r1 r2 subsetD)
qed

lemma *not_essential_times[singularity_intros]:*

assumes *f_ness: not_essential f z* **and** *g_ness: not_essential g z*

assumes *f_iso: isolated_singularity_at f z* **and** *g_iso: isolated_singularity_at g z*

shows *not_essential* ($\lambda w. f w * g w$) *z*

proof –

define *fg* **where** *fg* = ($\lambda w. f w * g w$)

have *?thesis* **when** $\neg ((\exists_{Fw} \text{in } (at\ z). f\ w \neq 0) \wedge (\exists_{Fw} \text{in } (at\ z). g\ w \neq 0))$

proof –

have $\forall_{Fw} \text{in } (at\ z). fg\ w = 0$

using *that[unfolded_frequently_def, simplified]* **unfolding** *fg_def*

by (*auto elim: eventually_rev_mp*)

from *tendsto_cong[OF this]* **have** *fg* $-z \rightarrow 0$ **by** *auto*

then show *?thesis* **unfolding** *not_essential_def fg_def* **by** *auto*

qed

moreover have *?thesis* **when** *f_nconst: $\exists_{Fw} \text{in } (at\ z). f\ w \neq 0$* **and** *g_nconst: $\exists_{Fw} \text{in } (at\ z). g\ w \neq 0$*

proof –

obtain *fn fp fr* **where** *[simp]: fp z $\neq 0$* **and** *fr > 0*

and *fr: fp holomorphic_on cball z fr*

$\forall w \in cball\ z\ fr - \{z\}. f\ w = fp\ w * (w - z)^{powi\ fn} \wedge fp\ w \neq 0$

using *holomorphic_factor_puncture[OF f_iso f_ness f_nconst, THEN ex1_implies_ex]*

by *auto*

obtain *gn gp gr* **where** *[simp]: gp z $\neq 0$* **and** *gr > 0*

and *gr: gp holomorphic_on cball z gr*

$\forall w \in cball\ z\ gr - \{z\}. g\ w = gp\ w * (w - z)^{powi\ gn} \wedge gp\ w \neq 0$

using *holomorphic_factor_puncture[OF g_iso g_ness g_nconst, THEN ex1_implies_ex]*

by *auto*

define *r1* **where** *r1 = (min fr gr)*

have *r1 > 0* **unfolding** *r1_def* **using** $\langle fr > 0 \rangle \langle gr > 0 \rangle$ **by** *auto*

have *fg_times: fg w = (fp w * gp w) * (w - z)^{powi (fn+gn)}* **and** *fgp_nz: fp w * gp w $\neq 0$*

when $w \in ball\ z\ r1 - \{z\}$ **for** *w*

proof –

have $f\ w = fp\ w * (w - z)^{powi\ fn}$ *fp w $\neq 0$*

using *fr(2)[rule_format, of w]* **that** **unfolding** *r1_def* **by** *auto*

moreover have $g\ w = gp\ w * (w - z)^{powi\ gn}$ *gp w $\neq 0$*

using *gr(2)[rule_format, of w]* **that** **unfolding** *r1_def* **by** *auto*

ultimately show $fg\ w = (fp\ w * gp\ w) * (w - z)^{powi (fn+gn)}$ *fp w * gp w $\neq 0$*

using *that* **unfolding** *fg_def* **by** (*auto simp add: power_int_add*)

qed

have [*intro*]: $fp -z \rightarrow fp\ z$ $gp -z \rightarrow gp\ z$

using *fr(1) $\langle fr > 0 \rangle$ gr(1) $\langle gr > 0 \rangle$*

by (*meson open_ball ball_subset_cball centre_in_ball*)

```

    continuous_on_eq_continuous_at continuous_within holomorphic_on_imp_continuous_on
    holomorphic_on_subset)+
  have ?thesis when  $fn+gn>0$ 
  proof -
    have  $(\lambda w. (fp\ w * gp\ w) * (w - z) ^{\text{nat}(fn+gn)}) -z \rightarrow 0$ 
      using that by (auto intro!: tendsto_eq_intros)
    then have  $fg -z \rightarrow 0$ 
      apply (elim Lim_transform_within[OF _ <r1>0])
    by (smt (verit, best) Diff_iff dist_commute fg_times mem_ball power_int_def
    singletonD that zero_less_dist_iff)
    then show ?thesis unfolding not_essential_def fg_def by auto
  qed
  moreover have ?thesis when  $fn+gn=0$ 
  proof -
    have  $(\lambda w. fp\ w * gp\ w) -z \rightarrow fp\ z * gp\ z$ 
      using that by (auto intro!: tendsto_eq_intros)
    then have  $fg -z \rightarrow fp\ z * gp\ z$ 
      apply (elim Lim_transform_within[OF _ <r1>0])
      apply (subst fg_times)
      by (auto simp add: dist_commute that)
    then show ?thesis unfolding not_essential_def fg_def by auto
  qed
  moreover have ?thesis when  $fn+gn<0$ 
  proof -
    have  $LIM\ x\ at\ z. (x - z) ^{\text{nat}(- (fn + gn))} :> at\ 0$ 
      using eventually_at_topological that
      by (force intro!: tendsto_eq_intros filterlim_atI)
    moreover have  $\exists c. (\lambda c. fp\ c * gp\ c) -z \rightarrow c \wedge 0 \neq c$ 
      using <fp -z -> fp z> <gp -z -> gp z> tendsto_mult by fastforce
    ultimately have  $LIM\ w\ (at\ z). fp\ w * gp\ w / (w-z)^{\text{nat}(- (fn+gn))} :>$ 
    at_infinity
      using filterlim_divide_at_infinity by blast
    then have is_pole fg z unfolding is_pole_def
      apply (elim filterlim_transform_within[OF _ _ <r1>0])
      using that
      by (simp_all add: dist_commute fg_times of_int_of_nat divide_simps
    power_int_def del: minus_add_distrib)
    then show ?thesis unfolding not_essential_def fg_def by auto
  qed
  ultimately show ?thesis unfolding not_essential_def fg_def by fastforce
  qed
  ultimately show ?thesis by auto
  qed

lemma not_essential_inverse[singularity_intros]:
  assumes f_ess: not_essential f z
  assumes f_iso: isolated_singularity_at f z
  shows not_essential  $(\lambda w. inverse\ (f\ w))\ z$ 
  proof -

```

```

define vf where vf = ( $\lambda w.$  inverse (f w))
have ?thesis when  $\neg(\exists_{Fw}$  in (at z). f w  $\neq 0$ )
proof -
  have  $\forall_{Fw}$  in (at z). f w = 0
  using that[unfolded frequently_def, simplified] by (auto elim: eventually_rev_mp)
  then have  $\forall_{Fw}$  in (at z). vf w = 0
  unfolding vf_def by auto
  from tendsto_cong[OF this] have vf  $-z \rightarrow 0$  unfolding vf_def by auto
  then show ?thesis unfolding not_essential_def vf_def by auto
qed
moreover have ?thesis when is_pole f z
proof -
  have vf  $-z \rightarrow 0$ 
  using that filterlim_at filterlim_inverse_at_iff unfolding is_pole_def vf_def
by blast
  then show ?thesis unfolding not_essential_def vf_def by auto
qed
moreover have ?thesis when  $\exists x.$  f  $-z \rightarrow x$  and f_nconst: $\exists_{Fw}$  in (at z). f w  $\neq 0$ 
proof -
  from that obtain fz where fz:f  $-z \rightarrow fz$  by auto
  have ?thesis when fz = 0
  proof -
    have ( $\lambda w.$  inverse (vf w))  $-z \rightarrow 0$ 
    using fz that unfolding vf_def by auto
    moreover have  $\forall_{Fw}$  in at z. inverse (vf w)  $\neq 0$ 
    using non_zero_neighbour[OF f_iso f_ness f_nconst]
    unfolding vf_def by auto
    ultimately have is_pole vf z
    using filterlim_inverse_at_iff[of vf at z] unfolding filterlim_at is_pole_def
by auto
    then show ?thesis unfolding not_essential_def vf_def by auto
  qed
  moreover have ?thesis when fz  $\neq 0$ 
  proof -
    have vf  $-z \rightarrow$  inverse fz
    using fz that unfolding vf_def by (auto intro:tendsto_eq_intros)
    then show ?thesis unfolding not_essential_def vf_def by auto
  qed
  ultimately show ?thesis by auto
qed
ultimately show ?thesis using f_ness unfolding not_essential_def by auto
qed

lemma isolated_singularity_at_inverse[singularity_intros]:
  assumes f_iso:isolated_singularity_at f z
  and f_ness:not_essential f z
  shows isolated_singularity_at ( $\lambda w.$  inverse (f w)) z
proof -
  define vf where vf = ( $\lambda w.$  inverse (f w))

```

```

have ?thesis when  $\neg(\exists_{Fw} \text{ in } (at\ z). f\ w \neq 0)$ 
proof -
  have  $\forall_{Fw} \text{ in } (at\ z). f\ w = 0$ 
  using that[unfolded frequently_def, simplified] by (auto elim: eventually_rev_mp)
  then have  $\forall_{Fw} \text{ in } (at\ z). \text{vf}\ w = 0$ 
  unfolding vf_def by auto
  then obtain d1 where  $d1 > 0$  and  $d1: \forall x. x \neq z \wedge \text{dist}\ x\ z < d1 \longrightarrow \text{vf}\ x = 0$ 
  unfolding eventually_at by auto
  then have vf_holomorphic_on_ball_z_d1_minus_z
  apply (rule_tac holomorphic_transform[of  $\lambda_. 0$ ])
  by (auto simp add: dist_commute)
  then have vf_analytic_on_ball_z_d1_minus_z
  by (simp add: analytic_on_open open_delete)
  then show ?thesis using  $\langle d1 > 0 \rangle$  unfolding isolated_singularity_at_def
  vf_def by auto
qed
moreover have ?thesis when  $f\_nconst: \exists_{Fw} \text{ in } (at\ z). f\ w \neq 0$ 
proof -
  have  $\forall_{Fw} \text{ in } at\ z. f\ w \neq 0$  using non_zero_neighbour[OF f_iso f_ness
  f_nconst].
  then obtain d1 where  $d1: d1 > 0 \forall x. x \neq z \wedge \text{dist}\ x\ z < d1 \longrightarrow f\ x \neq 0$ 
  unfolding eventually_at by auto
  obtain d2 where  $d2 > 0$  and  $d2: f\ \text{analytic\_on\_ball}\ z\ d2 - \{z\}$ 
  using f_iso unfolding isolated_singularity_at_def by auto
  define d3 where  $d3 = \min\ d1\ d2$ 
  have  $d3 > 0$  unfolding d3_def using  $\langle d1 > 0 \rangle \langle d2 > 0 \rangle$  by auto
  moreover
  have f_analytic_on_ball_z_d3_minus_z
  by (smt (verit, best) Diff_iff analytic_on_analytic_at d2 d3_def mem_ball)
  then have vf_analytic_on_ball_z_d3_minus_z
  unfolding vf_def
  by (intro analytic_on_inverse; simp add: d1(2) d3_def dist_commute)
  ultimately show ?thesis unfolding isolated_singularity_at_def vf_def by
  auto
qed
ultimately show ?thesis by auto
qed

lemma not_essential_divide[singularity_intros]:
  assumes f_ness: not_essential f z and g_ness: not_essential g z
  assumes f_iso: isolated_singularity_at f z and g_iso: isolated_singularity_at g
  z
  shows not_essential  $(\lambda w. f\ w / g\ w)$  z
proof -
  have not_essential  $(\lambda w. f\ w * \text{inverse}\ (g\ w))$  z
  by (simp add: f_iso f_ness g_iso g_ness isolated_singularity_at_inverse
  not_essential_inverse not_essential_times)
  then show ?thesis by (simp add: field_simps)
qed

```

lemma

assumes $f_iso:isolated_singularity_at\ f\ z$
and $g_iso:isolated_singularity_at\ g\ z$
shows $isolated_singularity_at_times[singularity_intros]:$
 $isolated_singularity_at\ (\lambda w. f\ w * g\ w)\ z$ **and**
 $isolated_singularity_at_add[singularity_intros]:$
 $isolated_singularity_at\ (\lambda w. f\ w + g\ w)\ z$

proof –

obtain $d1\ d2$ **where** $d1>0\ d2>0$
and $d1:f\ analytic_on\ ball\ z\ d1 - \{z\}$ **and** $d2:g\ analytic_on\ ball\ z\ d2 - \{z\}$
using $f_iso\ g_iso$ **unfolding** $isolated_singularity_at_def$ **by** $auto$
define $d3$ **where** $d3=min\ d1\ d2$
have $d3>0$ **unfolding** $d3_def$ **using** $\langle d1>0\rangle\ \langle d2>0\rangle$ **by** $auto$

have $fan: f\ analytic_on\ ball\ z\ d3 - \{z\}$
by $(smt\ (verit,\ best)\ Diff_iff\ analytic_on_analytic_at\ d1\ d3_def\ mem_ball)$
have $gan: g\ analytic_on\ ball\ z\ d3 - \{z\}$
by $(smt\ (verit,\ best)\ Diff_iff\ analytic_on_analytic_at\ d2\ d3_def\ mem_ball)$
have $(\lambda w. f\ w * g\ w)\ analytic_on\ ball\ z\ d3 - \{z\}$
using $analytic_on_mult\ fan\ gan$ **by** $blast$
then show $isolated_singularity_at\ (\lambda w. f\ w * g\ w)\ z$
using $\langle d3>0\rangle$ **unfolding** $isolated_singularity_at_def$ **by** $auto$
have $(\lambda w. f\ w + g\ w)\ analytic_on\ ball\ z\ d3 - \{z\}$
using $analytic_on_add\ fan\ gan$ **by** $blast$
then show $isolated_singularity_at\ (\lambda w. f\ w + g\ w)\ z$
using $\langle d3>0\rangle$ **unfolding** $isolated_singularity_at_def$ **by** $auto$

qed

lemma $isolated_singularity_at_uminus[singularity_intros]:$

assumes $f_iso:isolated_singularity_at\ f\ z$
shows $isolated_singularity_at\ (\lambda w. - f\ w)\ z$
using $assms$ **unfolding** $isolated_singularity_at_def$ **using** $analytic_on_neg$ **by**
 $blast$

lemma $isolated_singularity_at_id[singularity_intros]:$

$isolated_singularity_at\ (\lambda w. w)\ z$
unfolding $isolated_singularity_at_def$ **by** $(simp\ add: gt_ex)$

lemma $isolated_singularity_at_minus[singularity_intros]:$

assumes $f_iso:isolated_singularity_at\ f\ z$
and $g_iso:isolated_singularity_at\ g\ z$
shows $isolated_singularity_at\ (\lambda w. f\ w - g\ w)\ z$
using $isolated_singularity_at_uminus[THEN\ isolated_singularity_at_add[OF\ f_iso,of\ \lambda w. - g\ w]$
 $,OF\ g_iso]$ **by** $simp$

lemma $isolated_singularity_at_divide[singularity_intros]:$

assumes $f_iso:isolated_singularity_at\ f\ z$


```

    and g_iso:isolated_singularity_at g z
    and g_ness:not_essential g z
    shows isolated_singularity_at ( $\lambda w. f w / g w$ ) z
    using isolated_singularity_at_inverse[THEN isolated_singularity_at_times[OF
f_iso,
    of  $\lambda w. inverse (g w)$ ],OF g_iso g_ness] by (simp add:field_simps)

```

```

lemma isolated_singularity_at_const[singularity_intros]:
  isolated_singularity_at ( $\lambda w. c$ ) z
  unfolding isolated_singularity_at_def by (simp add: gt_ex)

```

```

lemma isolated_singularity_at_holomorphic:
  assumes f_holomorphic_on s - {z} open s z ∈ s
  shows isolated_singularity_at f z
  using assms unfolding isolated_singularity_at_def
  by (metis analytic_on_holomorphic centre_in_ball insert_Diff openE open_delete
subset_insert_iff)

```

```

lemma isolated_singularity_at_altdef:
  isolated_singularity_at f z  $\longleftrightarrow$  eventually ( $\lambda z. f analytic\_on \{z\}$ ) (at z)

```

proof

```

  assume isolated_singularity_at f z
  then obtain r where r: r > 0 f analytic_on ball z r - {z}
    unfolding isolated_singularity_at_def by blast
  have eventually ( $\lambda w. w \in ball z r - \{z\}$ ) (at z)
    using r(1) by (intro eventually_at_in_open) auto
  thus eventually ( $\lambda z. f analytic\_on \{z\}$ ) (at z)
    by eventually_elim (use r analytic_on_subset in auto)
next
  assume eventually ( $\lambda z. f analytic\_on \{z\}$ ) (at z)
  then obtain A where A: open A z ∈ A  $\wedge w. w \in A - \{z\} \implies f analytic\_on \{w\}$ 
    unfolding eventually_at_topological by blast
  then show isolated_singularity_at f z
    by (meson analytic_imp_holomorphic analytic_on_analytic_at isolated_singularity_at_holomorphic)
qed

```

```

lemma isolated_singularity_at_shift:

```

```

  assumes isolated_singularity_at ( $\lambda x. f (x + w)$ ) z
  shows isolated_singularity_at f (z + w)

```

proof –

```

  from assms obtain r where r: r > 0 and ana: ( $\lambda x. f (x + w)$ ) analytic_on
ball z r - {z}
  unfolding isolated_singularity_at_def by blast
  have (( $\lambda x. f (x + w)$ )  $\circ$  ( $\lambda x. x - w$ )) analytic_on (ball (z + w) r - {z + w})
    by (rule analytic_on_compose_gen[OF ana])
    (auto simp: dist_norm algebra_simps intro!: analytic_intros)
  hence f analytic_on (ball (z + w) r - {z + w})
    by (simp add: o_def)

```

```

thus ?thesis using r
  unfolding isolated_singularity_at_def by blast
qed

```

```

lemma isolated_singularity_at_shift_iff:
  isolated_singularity_at f (z + w)  $\longleftrightarrow$  isolated_singularity_at ( $\lambda x. f (x + w)$ ) z
  using isolated_singularity_at_shift[of f w z]
    isolated_singularity_at_shift[of  $\lambda x. f (x + w) -w w + z$ ]
  by (auto simp: algebra_simps)

```

```

lemma isolated_singularity_at_shift_0:
  NO_MATCH 0 z  $\implies$  isolated_singularity_at f z  $\longleftrightarrow$  isolated_singularity_at
  ( $\lambda x. f (z + x)$ ) 0
  using isolated_singularity_at_shift_iff[of f 0 z] by (simp add: add_ac)

```

```

lemma not_essential_shift:
  assumes not_essential ( $\lambda x. f (x + w)$ ) z
  shows not_essential f (z + w)
proof -
  from assms consider c where ( $\lambda x. f (x + w)$ )  $-z \rightarrow c$  | is_pole ( $\lambda x. f (x + w)$ ) z
  unfolding not_essential_def by blast
  thus ?thesis
proof cases
  case (1 c)
  hence f  $-z + w \rightarrow c$ 
  by (smt (verit, ccfv_SIG) LIM_cong add.assoc filterlim_at_to_0)
  thus ?thesis
  by (auto simp: not_essential_def)
next
  case 2
  hence is_pole f (z + w)
  by (subst is_pole_shift_iff [symmetric]) (auto simp: o_def add_ac)
  thus ?thesis
  by (auto simp: not_essential_def)
qed
qed

```

```

lemma not_essential_shift_iff: not_essential f (z + w)  $\longleftrightarrow$  not_essential ( $\lambda x. f (x + w)$ ) z
  using not_essential_shift[of f w z]
    not_essential_shift[of  $\lambda x. f (x + w) -w w + z$ ]
  by (auto simp: algebra_simps)

```

```

lemma not_essential_shift_0:
  NO_MATCH 0 z  $\implies$  not_essential f z  $\longleftrightarrow$  not_essential ( $\lambda x. f (z + x)$ ) 0
  using not_essential_shift_iff[of f 0 z] by (simp add: add_ac)

```

```

lemma not_essential_holomorphic:

```

```

assumes  $f$  holomorphic_on  $A$   $x \in A$  open  $A$ 
shows not_essential  $f$   $x$ 
proof -
  have continuous_on  $A$   $f$ 
    using assms holomorphic_on_imp_continuous_on by blast
  hence  $f -x \rightarrow f x$ 
    using assms continuous_on_eq_continuous_at isContD by blast
  thus ?thesis
    by (auto simp: not_essential_def)
qed

```

```

lemma not_essential_analytic:
  assumes  $f$  analytic_on  $\{z\}$ 
  shows not_essential  $f$   $z$ 
  using analytic_at assms not_essential_holomorphic by blast

```

```

lemma not_essential_id [singularity_intros]: not_essential  $(\lambda w. w)$   $z$ 
  by (simp add: not_essential_analytic)

```

```

lemma is_pole_imp_not_essential [intro]: is_pole  $f$   $z \implies$  not_essential  $f$   $z$ 
  by (auto simp: not_essential_def)

```

```

lemma tendsto_imp_not_essential [intro]:  $f -z \rightarrow c \implies$  not_essential  $f$   $z$ 
  by (auto simp: not_essential_def)

```

```

lemma eventually_not_pole:
  assumes isolated_singularity_at  $f$   $z$ 
  shows eventually  $(\lambda w. \neg \text{is\_pole } f w)$  (at  $z$ )
proof -
  from assms obtain  $r$  where  $r > 0$  and  $r$ :  $f$  analytic_on ball  $z$   $r - \{z\}$ 
    by (auto simp: isolated_singularity_at_def)
  then have eventually  $(\lambda w. w \in \text{ball } z$   $r - \{z\})$  (at  $z$ )
    by (intro eventually_at_in_open) auto
  thus eventually  $(\lambda w. \neg \text{is\_pole } f w)$  (at  $z$ )
  proof eventually_elim
    case (elim  $w$ )
    with  $r$  show ?case
    using analytic_imp_holomorphic not_is_pole_holomorphic open_delete by
blast
  qed
qed

```

```

lemma not_islimpt_poles:
  assumes isolated_singularity_at  $f$   $z$ 
  shows  $\neg z$  islimpt  $\{w. \text{is\_pole } f w\}$ 
  using eventually_not_pole [OF assms]
  by (auto simp: islimpt_conv_frequently_at frequently_def)

```

```

lemma analytic_at_imp_no_pole:  $f$  analytic_on  $\{z\} \implies \neg \text{is\_pole } f$   $z$ 

```

using *analytic_at not_is_pole_holomorphic* by *blast*

lemma *not_essential_const* [*singularity_intros*]: *not_essential* ($\lambda_. c$) *z*
by *blast*

lemma *not_essential_uminus* [*singularity_intros*]:

assumes *f_ness*: *not_essential* *f z*

assumes *f_iso*: *isolated_singularity_at* *f z*

shows *not_essential* ($\lambda w. -f w$) *z*

proof –

have *not_essential* ($\lambda w. -1 * f w$) *z*

by (*intro assms singularity_intros*)

thus ?*thesis* by *simp*

qed

lemma *isolated_singularity_at_analytic*:

assumes *f_analytic_on* {*z*}

shows *isolated_singularity_at* *f z*

proof –

from *assms* obtain *r* where *r*: *r* > 0 *f* *holomorphic_on* *ball z r*

by (*auto simp: analytic_on_def*)

show ?*thesis*

by (*rule isolated_singularity_at_holomorphic*[*of f ball z r*])

(*use* <*r* > 0> **in** <*auto intro!*: *holomorphic_on_subset*[*OF r*(2)]>)

qed

5.11 The order of non-essential singularities (i.e. removable singularities or poles)

definition *zorder* :: (*complex* \Rightarrow *complex*) \Rightarrow *complex* \Rightarrow *int* **where**

zorder f z = (*THE* *n*. ($\exists h r. r > 0 \wedge h$ *holomorphic_on* *cball z r* $\wedge h z \neq 0$
 $\wedge (\forall w \in \text{cball } z r - \{z\}. f w = h w * (w - z)^{\text{powi } n}$
 $\wedge h w \neq 0)$)))

definition *zor_poly*

:: [*complex* \Rightarrow *complex*, *complex*] \Rightarrow *complex* \Rightarrow *complex* **where**

zor_poly f z = (*SOME* *h*. $\exists r. r > 0 \wedge h$ *holomorphic_on* *cball z r* $\wedge h z \neq 0$
 $\wedge (\forall w \in \text{cball } z r - \{z\}. f w = h w * (w - z)^{\text{powi } (zorder f z)}$
 $\wedge h w \neq 0)$))

lemma *zorder_exist*:

fixes *f*::*complex* \Rightarrow *complex* **and** *z*::*complex*

defines *n* \equiv *zorder f z* **and** *g* \equiv *zor_poly f z*

assumes *f_iso*: *isolated_singularity_at* *f z*

and *f_ness*: *not_essential* *f z*

and *f_nconst*: $\exists_F w$ *in* (*at z*). *f w* $\neq 0$

shows *g z* $\neq 0 \wedge (\exists r. r > 0 \wedge g$ *holomorphic_on* *cball z r*

$\wedge (\forall w \in \text{cball } z r - \{z\}. f w = g w * (w - z)^{\text{powi } n} \wedge g w \neq 0)$)

proof –

```

define P where P = ( $\lambda n g r. 0 < r \wedge g \text{ holomorphic\_on cball } z r \wedge g z \neq 0$ 
   $\wedge (\forall w \in \text{cball } z r - \{z\}. f w = g w * (w-z) \text{ powi } n \wedge g w \neq 0)$ )
have  $\exists! n. \exists g r. P n g r$ 
  using holomorphic_factor_puncture[OF assms(3-)] unfolding P_def by auto
then have  $\exists g r. P n g r$ 
  unfolding n_def P_def zorder_def
  by (drule_tac theI', argo)
then have  $\exists r. P n g r$ 
  unfolding P_def zor_poly_def g_def n_def
  by (drule_tac someI_ex, argo)
then obtain r1 where P n g r1 by auto
then show ?thesis unfolding P_def by auto
qed

```

lemma zorder_shift:

```

shows zorder f z = zorder ( $\lambda u. f (u + z)$ ) 0
unfolding zorder_def
apply (rule arg_cong [of concl: The])
apply (auto simp: fun_eq_iff Ball_def dist_norm)
subgoal for x h r
  apply (rule_tac x=h o (+)z in exI)
  apply (rule_tac x=r in exI)
  apply (intro conjI holomorphic_on_compose holomorphic_intros)
  apply (simp_all flip: cball_translation)
  apply (simp add: add commute)
  done
subgoal for x h r
  apply (rule_tac x=h o ( $\lambda u. u-z$ ) in exI)
  apply (rule_tac x=r in exI)
  apply (intro conjI holomorphic_on_compose holomorphic_intros)
  apply (simp_all add: flip: cball_translation_subtract)
  by (metis diff_add_cancel eq_iff_diff_eq_0 norm_minus_commute)
done

```

lemma zorder_shift': NO_MATCH $0 z \implies \text{zorder } f z = \text{zorder } (\lambda u. f (u + z)) 0$
by (rule zorder_shift)

lemma

```

fixes f::complex  $\Rightarrow$  complex and z::complex
assumes f_iso:isolated_singularity_at f z
  and f_ness:not_essential f z
  and f_nconst: $\exists_F w \text{ in } (at z). f w \neq 0$ 
shows zorder_inverse: zorder ( $\lambda w. \text{inverse } (f w)$ ) z = - zorder f z
  and zor_poly_inverse:  $\forall_F w \text{ in } (at z). \text{zor\_poly } (\lambda w. \text{inverse } (f w)) z w$ 
  = inverse (zor_poly f z w)

```

proof -

```

define vf where vf = ( $\lambda w. \text{inverse } (f w)$ )
define fn vfn where
  fn = zorder f z and vfn = zorder vf z

```

```

define fp vfp where
  fp = zor_poly f z and vfp = zor_poly vf z

obtain fr where [simp]:fp z ≠ 0 and fr > 0
  and fr: fp holomorphic_on cball z fr
    ∀ w∈cball z fr - {z}. f w = fp w * (w - z) powi fn ∧ fp w ≠ 0
  using zorder_exist[OF f_iso f_nconst,folded fn_def fp_def]
  by auto
have fr_inverse: vf w = (inverse (fp w)) * (w - z) powi (-fn)
  and fr_nz: inverse (fp w) ≠ 0
  when w∈ball z fr - {z} for w
proof -
  have f w = fp w * (w - z) powi fn fp w ≠ 0
  using fr(2)[rule_format,of w] that by auto
  then show vf w = (inverse (fp w)) * (w - z) powi (-fn) inverse (fp w) ≠ 0
  by (simp_all add: power_int_minus vf_def)
qed
obtain vfr where [simp]:vfp z ≠ 0 and vfr>0 and vfr:vfp holomorphic_on cball
z vfr
  (∀ w∈cball z vfr - {z}. vf w = vfp w * (w - z) powi vfn ∧ vfp w ≠ 0)
proof -
  have isolated_singularity_at vf z
  using isolated_singularity_at_inverse[OF f_iso f_nconst] unfolding vf_def .
  moreover have not_essential vf z
  using not_essential_inverse[OF f_nconst f_iso] unfolding vf_def .
  moreover have ∃F w in at z. vf w ≠ 0
  using f_nconst unfolding vf_def by (auto elim:frequently_elim1)
  ultimately show ?thesis using zorder_exist[of vf z, folded vfn_def vfp_def]
that by auto
qed

define r1 where r1 = min fr vfr
have r1>0 using ⟨fr>0 ⟨vfr>0 unfolding r1_def by simp
show vfn = - fn
proof (rule holomorphic_factor_unique)
  have §: ∧w. [fp w = 0; dist z w < fr] ⇒ False
  using fr_nz by force
  then show ∀ w∈ball z r1 - {z}.
    vf w = vfp w * (w - z) powi vfn ∧
    vfp w ≠ 0 ∧ vf w = inverse (fp w) * (w - z) powi (-fn) ∧
    inverse (fp w) ≠ 0
  using fr_inverse r1_def vfr(2)
  by (smt (verit) Diff_iff inverse_nonzero_iff_nonzero mem_ball mem_cball)
show vfp holomorphic_on ball z r1
  using r1_def vfr(1) by auto
show (λw. inverse (fp w)) holomorphic_on ball z r1
  by (metis § ball_subset_cball fr(1) holomorphic_on_inverse holomorphic_on_subset
mem_ball min.cobounded2 min commute r1_def subset_ball)

```

```

qed (use <r1>0> in auto)
have vfp w = inverse (fp w) when w ∈ ball z r1 - {z} for w
proof -
  have w ∈ ball z fr - {z} w ∈ cball z vfr - {z} w ≠ z using that unfolding
  r1_def by auto
  from fr_inverse[OF this(1)] fr_nz[OF this(1)] vfr(2)[rule_format, OF this(2)]
  <vfn = - fn> <w ≠ z>
  show ?thesis by auto
qed
then show  $\forall_F w$  in (at z). vfp w = inverse (fp w)
  unfolding eventually_at using <r1>0>
  by (metis DiffI dist_commute mem_ball singletonD)
qed

```

```

lemma zor_poly_shift:
  assumes iso1: isolated_singularity_at f z
  and ness1: not_essential f z
  and nzero1:  $\exists_F w$  in at z. f w ≠ 0
  shows  $\forall_F w$  in nhds z. zor_poly f z w = zor_poly ( $\lambda u. f (z + u)$ ) 0 (w - z)
proof -
  obtain r1 where r1 > 0 zor_poly f z z ≠ 0 and
    holo1: zor_poly f z holomorphic_on cball z r1 and
    rball1:  $\forall w \in \text{cball } z \ r1 - \{z\}.$ 
      f w = zor_poly f z w * (w - z) powi (zorder f z)  $\wedge$ 
      zor_poly f z w ≠ 0
  using zorder_exist[OF iso1 ness1 nzero1] by blast

```

```

define ff where ff = ( $\lambda u. f (z + u)$ )
have isolated_singularity_at ff 0
  unfolding ff_def
  using iso1 isolated_singularity_at_shift_iff[of f 0 z]
  by (simp add: algebra_simps)
moreover have not_essential ff 0
  unfolding ff_def
  using ness1 not_essential_shift_iff[of f 0 z]
  by (simp add: algebra_simps)
moreover have  $\exists_F w$  in at 0. ff w ≠ 0
  unfolding ff_def using nzero1
  by (smt (verit, ccfv_SIG) add_commute eventually_at_to_0
    eventually_mono not_frequently)
ultimately obtain r2 where r2 > 0 zor_poly ff 0 0 ≠ 0 and
  holo2: zor_poly ff 0 holomorphic_on cball 0 r2 and
  rball2:  $\forall w \in \text{cball } 0 \ r2 - \{0\}.$ 
    ff w = zor_poly ff 0 w * w powi (zorder ff 0)  $\wedge$ 
    zor_poly ff 0 w ≠ 0
  using zorder_exist[of ff 0] by auto

```

```

define r where r = min r1 r2
have r > 0 using <r1>0> <r2>0> unfolding r_def by auto

```

```

have zor_poly f z w = zor_poly ff 0 (w - z)
  if w ∈ ball z r - {z} for w
proof -
  define n where n ≡ zorder f z

  have f w = zor_poly f z w * (w - z) powi n
  proof -
    have w ∈ cball z r1 - {z}
    using r_def that by auto
    from rball1[rule_format, OF this]
    show ?thesis unfolding n_def by auto
  qed

  moreover have f w = zor_poly ff 0 (w - z) * (w - z) powi n
  proof -
    have w - z ∈ cball 0 r2 - {0}
    using r_def that by (auto simp:dist_complex_def)
    from rball2[rule_format, OF this]
    have ff (w - z) = zor_poly ff 0 (w - z) * (w - z)
      powi (zorder ff 0)

    by simp
    moreover have of_int (zorder ff 0) = n
    unfolding n_def ff_def by (simp add:zorder_shift' add.commute)
    ultimately show ?thesis unfolding ff_def by auto
  qed

  ultimately have zor_poly f z w * (w - z) powi n
    = zor_poly ff 0 (w - z) * (w - z) powi n
    by auto
  moreover have (w - z) powi n ≠ 0
    using that by auto
  ultimately show ?thesis
    using mult_cancel_right by blast
  qed

then have ∀F w in at z. zor_poly f z w
  = zor_poly ff 0 (w - z)
  unfolding eventually_at
  apply (rule_tac x=r in exI)
  using ⟨r>0⟩ by (auto simp:dist_commute)
moreover have isCont (zor_poly f z) z
  using holo1[THEN holomorphic_on_imp_continuous_on]
  apply (elim continuous_on_interior)
  using ⟨r1>0⟩ by auto
moreover have isCont (λw. zor_poly ff 0 (w - z)) z
proof -
  have isCont (zor_poly ff 0) 0
  using holo2[THEN holomorphic_on_imp_continuous_on]
  apply (elim continuous_on_interior)
  using ⟨r2>0⟩ by auto

```



```

    then show ?thesis
      unfolding isCont_iff by simp
    qed
  ultimately show  $\forall_F w \text{ in nhds } z. \text{zor\_poly } f z w$ 
    =  $\text{zor\_poly } ff 0 (w - z)$ 
    by (elim at_within_isCont_imp_nhds; auto)
  qed

lemma
  fixes  $f g :: \text{complex} \Rightarrow \text{complex}$  and  $z :: \text{complex}$ 
  assumes  $f\_iso: \text{isolated\_singularity\_at } f z$  and  $g\_iso: \text{isolated\_singularity\_at } g z$ 
    and  $f\_ness: \text{not\_essential } f z$  and  $g\_ness: \text{not\_essential } g z$ 
    and  $fg\_nconst: \exists_F w \text{ in } (at z). f w * g w \neq 0$ 
  shows  $\text{zorder\_times: zorder } (\lambda w. f w * g w) z = \text{zorder } f z + \text{zorder } g z$  and
     $\text{zor\_poly\_times: } \forall_F w \text{ in } (at z). \text{zor\_poly } (\lambda w. f w * g w) z w$ 
    =  $\text{zor\_poly } f z w * \text{zor\_poly } g z w$ 

proof -
  define fg where  $fg = (\lambda w. f w * g w)$ 
  define fn gn fgn where
     $fn = \text{zorder } f z$  and  $gn = \text{zorder } g z$  and  $fgn = \text{zorder } fg z$ 
  define fp gp fgp where
     $fp = \text{zor\_poly } f z$  and  $gp = \text{zor\_poly } g z$  and  $fgp = \text{zor\_poly } fg z$ 
  have  $f\_nconst: \exists_F w \text{ in } (at z). f w \neq 0$  and  $g\_nconst: \exists_F w \text{ in } (at z). g w \neq 0$ 
    using  $fg\_nconst$  by (auto elim!: frequently_elim1)
  obtain fr where  $[simp]: fp z \neq 0$  and  $fr > 0$ 
    and  $fr: fp \text{ holomorphic\_on } cball z fr$ 
     $\forall w \in cball z fr - \{z\}. f w = fp w * (w - z) \text{ powi } fn \wedge fp w \neq 0$ 
    using  $\text{zorder\_exist}[OF f\_iso f\_ness f\_nconst, folded fp\_def fn\_def]$  by auto
  obtain gr where  $[simp]: gp z \neq 0$  and  $gr > 0$ 
    and  $gr: gp \text{ holomorphic\_on } cball z gr$ 
     $\forall w \in cball z gr - \{z\}. g w = gp w * (w - z) \text{ powi } gn \wedge gp w \neq 0$ 
    using  $\text{zorder\_exist}[OF g\_iso g\_ness g\_nconst, folded gn\_def gp\_def]$  by auto
  define r1 where  $r1 = \min fr gr$ 
  have  $r1 > 0$  unfolding  $r1\_def$  using  $\langle fr > 0 \rangle \langle gr > 0 \rangle$  by auto
  have  $fg\_times: fg w = (fp w * gp w) * (w - z) \text{ powi } (fn + gn)$  and  $fgp\_nz: fp w * gp w \neq 0$ 
    when  $w \in ball z r1 - \{z\}$  for  $w$ 
  proof -
    have  $f w = fp w * (w - z) \text{ powi } fn$  and  $fp w \neq 0$ 
      using  $fr(2)[rule\_format, of w]$  that unfolding  $r1\_def$  by auto
    moreover have  $g w = gp w * (w - z) \text{ powi } gn$  and  $gp w \neq 0$ 
      using  $gr(2)[rule\_format, of w]$  that unfolding  $r1\_def$  by auto
    ultimately show  $fg w = (fp w * gp w) * (w - z) \text{ powi } (fn + gn)$  and  $fgp w \neq 0$ 
      using that
      unfolding  $fg\_def$  by (auto simp add: power_int_add)
    qed
  obtain fgr where  $[simp]: fgp z \neq 0$  and  $fgr > 0$ 

```

```

    and fgr: fgp holomorphic_on cball z fgr
       $\forall w \in \text{cball } z \text{ fgr} - \{z\}. fg \ w = fgp \ w * (w - z) \text{ powi } fgn \wedge fgp \ w \neq 0$ 
  proof -
    have fgp z  $\neq 0 \wedge (\exists r > 0. fgp \text{ holomorphic\_on } \text{cball } z \ r$ 
       $\wedge (\forall w \in \text{cball } z \ r - \{z\}. fg \ w = fgp \ w * (w - z) \text{ powi } fgn \wedge fgp \ w \neq 0))$ 
    apply (rule zorder_exist[of fg z, folded fgn_def fgp_def])
    subgoal unfolding fg_def using isolated_singularity_at_times[OF f_iso
g_iso] .
    subgoal unfolding fg_def using not_essential_times[OF f_ess g_ess
f_iso g_iso] .
    subgoal unfolding fg_def using fg_nconst .
    done
    then show ?thesis using that by blast
  qed
  define r2 where r2 = min fgr r1
  have r2 > 0 using <r1 > 0 <fgr > 0 unfolding r2_def by simp
  show fgn = fn + gn
    apply (rule holomorphic_factor_unique[of r2 fgp z  $\lambda w. fp \ w * gp \ w \ fg$ ])
    subgoal using <r2 > 0 by simp
    subgoal by simp
    subgoal by simp
    subgoal
    proof (rule ballI)
      fix w assume w  $\in \text{ball } z \ r2 - \{z\}$ 
      then have w  $\in \text{ball } z \ r1 - \{z\} \ w \in \text{cball } z \ \text{fgr} - \{z\}$  unfolding r2_def by
auto
      from fg_times[OF this(1)] fgp_nz[OF this(1)] fgr(2)[rule_format, OF this(2)]
      show fg w = fgp w * (w - z) powi fgn  $\wedge fgp \ w \neq 0$ 
         $\wedge fg \ w = fp \ w * gp \ w * (w - z) \text{ powi } (fn + gn) \wedge fp \ w * gp \ w \neq 0$  by
auto
    qed
    subgoal using fgr(1) unfolding r2_def r1_def by (auto intro!: holomorphic_intros)
    subgoal using fr(1) gr(1) unfolding r2_def r1_def by (auto intro!: holomorphic_intros)
    done

  have fgp w = fp w * gp w when w  $\in \text{ball } z \ r2 - \{z\}$  for w
  proof -
    have w  $\in \text{ball } z \ r1 - \{z\} \ w \in \text{cball } z \ \text{fgr} - \{z\} \ w \neq z$  using that unfolding
r2_def by auto
    from fg_times[OF this(1)] fgp_nz[OF this(1)] fgr(2)[rule_format, OF this(2)]
    <fgn = fn + gn> <w  $\neq z$ >
    show ?thesis by auto
  qed
  then show  $\forall_F w \text{ in } (at \ z). fgp \ w = fp \ w * gp \ w$ 
    using <r2 > 0 unfolding eventually_at by (auto simp add: dist_commute)
  qed

lemma
  fixes f g :: complex  $\Rightarrow$  complex and z :: complex

```

assumes $f_iso:isolated_singularity_at\ f\ z$ **and** $g_iso:isolated_singularity_at\ g\ z$
and $f_ness:not_essential\ f\ z$ **and** $g_ness:not_essential\ g\ z$
and $fg_nconst:\exists_{F}w\ in\ (at\ z). f\ w * g\ w \neq 0$
shows $zorder_divide:zorder\ (\lambda w. f\ w / g\ w)\ z = zorder\ f\ z - zorder\ g\ z$ **and**
 $zor_poly_divide:\forall_{F}w\ in\ (at\ z). zor_poly\ (\lambda w. f\ w / g\ w)\ z\ w$
 $= zor_poly\ f\ z\ w / zor_poly\ g\ z\ w$

proof –

have $f_nconst:\exists_{F}w\ in\ (at\ z). f\ w \neq 0$ **and** $g_nconst:\exists_{F}w\ in\ (at\ z). g\ w \neq 0$
using fg_nconst **by** (*auto elim!:frequently_elim1*)
define vg **where** $vg=(\lambda w. inverse\ (g\ w))$
have $zorder\ (\lambda w. f\ w * vg\ w)\ z = zorder\ f\ z + zorder\ vg\ z$
apply (*rule zorder_times[OF f_iso _ f_ness,of vg]*)
subgoal unfolding vg_def **using** $isolated_singularity_at_inverse[OF\ g_iso\ g_ness]$.
subgoal unfolding vg_def **using** $not_essential_inverse[OF\ g_ness\ g_iso]$.
subgoal unfolding vg_def **using** fg_nconst **by** (*auto elim!:frequently_elim1*)
done
then show $zorder\ (\lambda w. f\ w / g\ w)\ z = zorder\ f\ z - zorder\ g\ z$
using $zorder_inverse[OF\ g_iso\ g_ness\ g_nconst,folded\ vg_def]$ **unfolding**
 vg_def
by (*auto simp add:field_simps*)

have $\forall_{F} w\ in\ at\ z. zor_poly\ (\lambda w. f\ w * vg\ w)\ z\ w = zor_poly\ f\ z\ w * zor_poly\ vg\ z\ w$
apply (*rule zor_poly_times[OF f_iso _ f_ness,of vg]*)
subgoal unfolding vg_def **using** $isolated_singularity_at_inverse[OF\ g_iso\ g_ness]$.
subgoal unfolding vg_def **using** $not_essential_inverse[OF\ g_ness\ g_iso]$.
subgoal unfolding vg_def **using** fg_nconst **by** (*auto elim!:frequently_elim1*)
done
then show $\forall_{F}w\ in\ (at\ z). zor_poly\ (\lambda w. f\ w / g\ w)\ z\ w = zor_poly\ f\ z\ w /$
 $zor_poly\ g\ z\ w$
using $zor_poly_inverse[OF\ g_iso\ g_ness\ g_nconst,folded\ vg_def]$ **unfolding**
 vg_def
by *eventually_elim* (*auto simp add:field_simps*)

qed

lemma $zorder_exist_zero$:

fixes $f::complex \Rightarrow complex$ **and** $z::complex$
defines $n \equiv zorder\ f\ z$ **and** $g \equiv zor_poly\ f\ z$
assumes $holo: f\ holomorphic_on\ s$ **and**
 $open\ s\ connected\ s\ z \in s$
and $non_const: \exists w \in s. f\ w \neq 0$
shows (*if* $f\ z = 0$ *then* $n > 0$ *else* $n = 0$) $\wedge (\exists r. r > 0 \wedge cball\ z\ r \subseteq s \wedge g\ holomorphic_on\ cball\ z\ r$
 $\wedge (\forall w \in cball\ z\ r. f\ w = g\ w * (w - z) ^ nat\ n \wedge g\ w \neq 0))$

proof –

obtain r **where** $g\ z \neq 0$ **and** $r: r > 0\ cball\ z\ r \subseteq s\ g\ holomorphic_on\ cball\ z\ r$

```

      ( $\forall w \in \text{cball } z \ r - \{z\}. f \ w = g \ w * (w - z) \text{ powi } n \wedge g \ w \neq 0$ )
proof -
  have  $g \ z \neq 0 \wedge (\exists r > 0. g \text{ holomorphic\_on } \text{cball } z \ r$ 
     $\wedge (\forall w \in \text{cball } z \ r - \{z\}. f \ w = g \ w * (w - z) \text{ powi } n \wedge g \ w \neq 0))$ 
  proof (rule zorder_exist[of f z, folded g_def n_def])
  show isolated_singularity_at f z unfolding isolated_singularity_at_def
    using holo_assms(4,6)
    by (meson Diff_subset open_ball analytic_on_holomorphic holomor-
      phic_on_subset openE)
  show not_essential f z unfolding not_essential_def
  using assms(4,6) at_within_open continuous_on_holo holomorphic_on_imp_continuous_on
    by fastforce
  have  $\forall_F w \text{ in } \text{at } z. f \ w \neq 0 \wedge w \in s$ 
  proof -
    obtain w where  $w \in s \ f \ w \neq 0$  using non_const by auto
    then show ?thesis
      by (rule non_zero_neighbour_alt[OF holo ⟨open s⟩ ⟨connected s⟩ ⟨z ∈ s⟩])
  qed
  then show  $\exists_F w \text{ in } \text{at } z. f \ w \neq 0$ 
    apply (elim eventually_frequentlyE)
    by auto
  qed
then obtain r1 where  $g \ z \neq 0 \ r1 > 0$  and  $r1 : g \text{ holomorphic\_on } \text{cball } z \ r1$ 
  ( $\forall w \in \text{cball } z \ r1 - \{z\}. f \ w = g \ w * (w - z) \text{ powi } n \wedge g \ w \neq 0$ )
  by auto
obtain r2 where  $r2 : r2 > 0 \ \text{cball } z \ r2 \subseteq s$ 
  using assms(4,6) open_contains_cball_eq by blast
define r3 where  $r3 = \min \ r1 \ r2$ 
have  $r3 > 0 \ \text{cball } z \ r3 \subseteq s$  using ⟨r1 > 0⟩ r2 unfolding r3_def by auto
moreover have  $g \text{ holomorphic\_on } \text{cball } z \ r3$ 
  using r1(1) unfolding r3_def by auto
moreover have ( $\forall w \in \text{cball } z \ r3 - \{z\}. f \ w = g \ w * (w - z) \text{ powi } n \wedge g \ w \neq$ 
0)
  using r1(2) unfolding r3_def by auto
  ultimately show ?thesis using that[of r3] ⟨g z ≠ 0⟩ by auto
qed

have fz_lim:  $f - z \rightarrow f \ z$ 
by (metis assms(4,6) at_within_open continuous_on_holo holomorphic_on_imp_continuous_on)
have gz_lim:  $g - z \rightarrow g \ z$ 
  using r
by (meson Elementary_Metric_Spaces.open_ball analytic_at analytic_at_imp_isCont

  ball_subset_cball centre_in_ball holomorphic_on_subset isContD)
have if_0: if  $f \ z = 0$  then  $n > 0$  else  $n = 0$ 
proof -
  have ( $\lambda w. g \ w * (w - z) \text{ powi } n$ )  $-z \rightarrow f \ z$ 
    using fz_lim Lim_transform_within_open[where  $s = \text{ball } z \ r$ ] r by fastforce
  then have ( $\lambda w. (g \ w * (w - z) \text{ powi } n) / g \ w$ )  $-z \rightarrow f \ z / g \ z$ 

```

```

    using gz_lim ⟨g z ≠ 0⟩ tendsto_divide by blast
  then have powi_tendsto:(λw. (w - z) powi n) -z→ f z/g z
    using Lim_transform_within_open[where s=ball z r] r by fastforce

  have ?thesis when n≥0 f z=0
  proof -
    have (λw. (w - z) ^ nat n) -z→ f z/g z
      using Lim_transform_within[OF powi_tendsto, where d=r]
      by (meson power_int_def r(1) that(1))
    then have *:(λw. (w - z) ^ nat n) -z→ 0 using ⟨f z=0⟩ by simp
    moreover have False when n=0
    proof -
      have (λw. (w - z) ^ nat n) -z→ 1
        using ⟨n=0⟩ by auto
      then show False using * using LIM_unique zero_neq_one by blast
    qed
    ultimately show ?thesis using that by fastforce
  qed
  moreove have ?thesis when n≥0 f z≠0
  proof -
    have False when n>0
    proof -
      have (λw. (w - z) ^ nat n) -z→ f z/g z
        using Lim_transform_within[OF powi_tendsto, where d=r]
        by (meson ⟨0 ≤ n⟩ power_int_def r(1))
      moreover have (λw. (w - z) ^ nat n) -z→ 0
        using ⟨n>0⟩ by (auto intro!:tendsto_eq_intros)
      ultimately show False using ⟨f z≠0⟩ ⟨g z≠0⟩ using LIM_unique di-
vide_eq_0_iff by blast
    qed
    then show ?thesis using that by force
  qed
  moreover have False when n<0
  proof -
    have (λw. inverse ((w - z) ^ nat (- n))) -z→ f z/g z
      by (smt (verit) LIM_cong power_int_def power_inverse powi_tendsto that)
    moreover
    have (λw.((w - z) ^ nat (- n))) -z→ 0
      using that by (auto intro!:tendsto_eq_intros)
    ultimately
    have (λx. inverse ((x - z) ^ nat (- n)) * (x - z) ^ nat (- n)) -z→ 0
      using tendsto_mult by fastforce
    then have (λx. 1::complex) -z→ 0
      by (elim Lim_transform_within_open[where s=UNIV],auto)
    then show False using LIM_const_eq by fastforce
  qed
  ultimately show ?thesis by fastforce
  qed
  moreover have f w = g w * (w - z) ^ nat n ∧ g w ≠ 0 when w∈cball z r for w

```

```

proof (cases w=z)
  case True
  then have  $f -z \rightarrow f w$ 
    using fz_lim by blast
  then have  $(\lambda w. g w * (w-z) ^{\wedge} \text{nat } n) -z \rightarrow f w$ 
  proof (elim Lim_transform_within[OF_ <r>0])
    fix x assume  $0 < \text{dist } x z \text{ dist } x z < r$ 
    then have  $x \in \text{cball } z r - \{z\} \ x \neq z$ 
      unfolding cball_def by (auto simp add: dist_commute)
    then have  $f x = g x * (x - z) \text{ powi } n$ 
      using r(4)[rule_format,of x] by simp
    also have  $\dots = g x * (x - z) ^{\wedge} \text{nat } n$ 
      by (smt (verit, best) if_0_int_nat_eq_power_int_of_nat)
    finally show  $f x = g x * (x - z) ^{\wedge} \text{nat } n$  .
  qed
  moreover have  $(\lambda w. g w * (w-z) ^{\wedge} \text{nat } n) -z \rightarrow g w * (w-z) ^{\wedge} \text{nat } n$ 
    using True by (auto intro!:tendsto_eq_intros gz_lim)
  ultimately have  $f w = g w * (w-z) ^{\wedge} \text{nat } n$  using LIM_unique by blast
  then show ?thesis using <g z≠0> True by auto
next
  case False
  then have  $f w = g w * (w - z) \text{ powi } n \wedge g w \neq 0$ 
    using r(4) that by auto
  then show ?thesis
    by (smt (verit, best) False if_0_int_nat_eq_power_int_of_nat)
  qed
ultimately show ?thesis using r by auto
qed

lemma zorder_exist_pole:
  fixes f::complex  $\Rightarrow$  complex and z::complex
  defines n $\equiv$ zorder f z and g $\equiv$ zor_poly f z
  assumes holo: f holomorphic_on s- $\{z\}$  and
    open s z $\in$ s
    and is_pole f z
  shows  $n < 0 \wedge g z \neq 0 \wedge (\exists r. r > 0 \wedge \text{cball } z r \subseteq s \wedge g \text{ holomorphic\_on } \text{cball } z r$ 
     $\wedge (\forall w \in \text{cball } z r - \{z\}. f w = g w / (w-z) ^{\wedge} \text{nat } (-n) \wedge g w \neq 0))$ 
proof -
  obtain r where  $g z \neq 0$  and  $r: r > 0 \text{ cball } z r \subseteq s \wedge g \text{ holomorphic\_on } \text{cball } z r$ 
     $(\forall w \in \text{cball } z r - \{z\}. f w = g w * (w - z) \text{ powi } n \wedge g w \neq 0)$ 
  proof -
  have  $g z \neq 0 \wedge (\exists r > 0. g \text{ holomorphic\_on } \text{cball } z r$ 
     $\wedge (\forall w \in \text{cball } z r - \{z\}. f w = g w * (w - z) \text{ powi } n \wedge g w \neq 0))$ 
  proof (rule zorder_exist[of f z,folded g_def n_def])
    show isolated_singularity_at f z unfolding isolated_singularity_at_def
      using holo assms(4,5)
    by (metis analytic_on_holomorphic_centre_in_ball insert_Diff openE
open_delete_subset_insert_iff)
    show not_essential f z unfolding not_essential_def

```

```

    using assms(4,6) at_within_open continuous_on holo holomorphic_on_imp_continuous_on
      by fastforce
    from non_zero_neighbour_pole[OF ‹is_pole f z›] show  $\exists_F w$  in at z.  $f w \neq$ 
0
      apply (elim eventually_frequentlyE)
      by auto
    qed
    then obtain r1 where  $g z \neq 0$   $r1 > 0$  and  $r1: g$  holomorphic_on cball z r1
      ( $\forall w \in$  cball z r1 - {z}.  $f w = g w * (w - z)^{\text{powi } n} \wedge g w \neq 0$ )
      by auto
    obtain r2 where  $r2: r2 > 0$  cball z r2  $\subseteq s$ 
      using assms(4,5) open_contains_cball_eq by metis
    define r3 where  $r3 = \min r1 r2$ 
    have  $r3 > 0$  cball z r3  $\subseteq s$  using ‹ $r1 > 0$ › r2 unfolding r3_def by auto
    moreover have  $g$  holomorphic_on cball z r3
      using r1(1) unfolding r3_def by auto
    moreover have ( $\forall w \in$  cball z r3 - {z}.  $f w = g w * (w - z)^{\text{powi } n} \wedge g w \neq$ 
0)
      using r1(2) unfolding r3_def by auto
    ultimately show ?thesis using that[of r3] ‹ $g z \neq 0$ › by auto
  qed

  have  $n < 0$ 
  proof (rule ccontr)
    assume  $\neg n < 0$ 
    define c where  $c = (\text{if } n = 0 \text{ then } g z \text{ else } 0)$ 
    have [simp]:  $g -z \rightarrow g z$ 
      by (metis open_ball at_within_open ball_subset_cball centre_in_ball
        continuous_on holomorphic_on_imp_continuous_on holomorphic_on_subset
        r(1) r(3))
    have  $\forall_F x$  in at z.  $f x = g x * (x - z)^{\wedge \text{nat } n}$ 
      unfolding eventually_at_topological
      apply (rule_tac exI[where  $x = \text{ball } z r$ ])
      by (simp add: ‹ $\neg n < 0$ › linorder_not_le power_int_def r(1) r(4))
    moreover have ( $\lambda x. g x * (x - z)^{\wedge \text{nat } n} -z \rightarrow c$ )
    proof (cases  $n = 0$ )
      case True
      then show ?thesis unfolding c_def by simp
    next
      case False
      then have ( $\lambda x. (x - z)^{\wedge \text{nat } n} -z \rightarrow 0$ ) using ‹ $\neg n < 0$ ›
        by (auto intro!: tendsto_eq_intros)
      from tendsto_mult[OF _ this, of  $g z$ , simplified]
      show ?thesis unfolding c_def using False by simp
    qed
    ultimately have  $f -z \rightarrow c$  using tendsto_cong by fast
  then show False using ‹is_pole f z› at_neq_bot not_tendsto_and_filterlim_at_infinity
    unfolding is_pole_def by blast
  qed

```

moreover have $\forall w \in \text{cball } z \ r - \{z\}. f \ w = g \ w / (w - z) \wedge \text{nat } (-n) \wedge g \ w \neq 0$
using $r(4) \langle n < 0 \rangle$
by (*smt* (*verit*) *inverse_eq_divide* *mult.right_neutral* *power_int_def* *power_inverse* *times_divide_eq_right*)
ultimately show *?thesis* **using** $r(1-3) \langle g \ z \neq 0 \rangle$ **by** *auto*
qed

lemma *zorder_eqI*:

assumes *open* $s \ z \in s \ g \ \text{holomorphic_on } s \ g \ z \neq 0$
assumes $fg_eq: \bigwedge w. \llbracket w \in s; w \neq z \rrbracket \implies f \ w = g \ w * (w - z) \ \text{powi } n$
shows $zorder \ f \ z = n$

proof –

have *continuous_on* $s \ g$ **by** (*rule* *holomorphic_on_imp_continuous_on*) *fact*
moreover have *open* $(-\{0::\text{complex}\})$ **by** *auto*
ultimately have *open* $((g - '(-\{0\})) \cap s)$
unfolding *continuous_on_open_vimage*[*OF* $\langle \text{open } s \rangle$] **by** *blast*
moreover from *assms* **have** $z \in (g - '(-\{0\})) \cap s$ **by** *auto*
ultimately obtain r **where** $r: r > 0 \ \text{cball } z \ r \subseteq s \cap (g - '(-\{0\}))$
unfolding *open_contains_cball* **by** *blast*

let *?gg* = $(\lambda w. g \ w * (w - z) \ \text{powi } n)$

define P **where** $P = (\lambda n \ g \ r. 0 < r \wedge g \ \text{holomorphic_on } \text{cball } z \ r \wedge g \ z \neq 0$
 $\wedge (\forall w \in \text{cball } z \ r - \{z\}. f \ w = g \ w * (w - z) \ \text{powi } n \wedge g \ w \neq 0))$

have $P \ n \ g \ r$

unfolding *P_def* **using** $r \ \text{assms}(3,4,5)$ **by** *auto*

then have $\exists g \ r. P \ n \ g \ r$ **by** *auto*

moreover have *unique*: $\exists! n. \exists g \ r. P \ n \ g \ r$ **unfolding** *P_def*

proof (*rule* *holomorphic_factor_puncture*)

have $\text{ball } z \ r - \{z\} \subseteq s$ **using** r **using** *ball_subset_cball* **by** *blast*

then have *?gg* *holomorphic_on* $\text{ball } z \ r - \{z\}$

using $\langle g \ \text{holomorphic_on } s \rangle \ r$ **by** (*auto* *intro!*: *holomorphic_intros*)

then have $f \ \text{holomorphic_on } \text{ball } z \ r - \{z\}$

by (*smt* (*verit*, *best*) *DiffD2* $\langle \text{ball } z \ r - \{z\} \subseteq s \rangle \ fg_eq \ \text{holomorphic_cong}$ *singleton_iff_subset_iff*)

then show *isolated_singularity_at* $f \ z$ **unfolding** *isolated_singularity_at_def*

using *analytic_on_open* *open_delete* $r(1)$ **by** *blast*

next

have *not_essential* *?gg* z

proof (*intro* *singularity_intros*)

show *not_essential* $g \ z$

by (*meson* $\langle \text{continuous_on } s \rangle \ \text{assms } \text{continuous_on_eq_continuous_at}$ *isCont_def* *not_essential_def*)

show $\forall_F w \ \text{in_at } z. w - z \neq 0$ **by** (*simp* *add*: *eventually_at_filter*)

then show *LIM* $w \ \text{at } z. w - z :> \text{at } 0$

unfolding *filterlim_at* **by** (*auto* *intro*: *tendsto_eq_intros*)

show *isolated_singularity_at* $g \ z$

by (*meson* *Diff_subset_open_ball* *analytic_on_holomorphic* *assms* *holomorphic_on_subset* *isolated_singularity_at_def* *openE*)

qed


```

moreover
have  $\forall_F w \text{ in } at\ z. g\ w * (w - z)^{pow\ n} = f\ w$ 
  unfolding eventually_at_topological using assms fg_eq by force
ultimately show not_essential f z
  using not_essential_transform by blast
show  $\exists_F w \text{ in } at\ z. f\ w \neq 0$  unfolding frequently_at
proof (intro strip)
  fix d::real assume  $0 < d$ 
  define z' where  $z' \equiv z + \min\ d\ r / 2$ 
  have  $z' \neq z \wedge \text{dist}\ z'\ z < d$ 
    unfolding z'_def using  $\langle d > 0 \rangle \langle r > 0 \rangle$  by (auto simp add: dist_norm)
  moreover have  $f\ z' \neq 0$ 
  proof (subst fg_eq[OF _  $\langle z' \neq z \rangle$ ])
    have  $z' \in \text{cball}\ z\ r$ 
      unfolding z'_def using  $\langle r > 0 \rangle \langle d > 0 \rangle$  by (auto simp add: dist_norm)
    then show  $z' \in s$  using r(2) by blast
    show  $g\ z' * (z' - z)^{pow\ n} \neq 0$ 
      using P_def  $\langle P\ n\ g\ r \rangle \langle z' \in \text{cball}\ z\ r \rangle \langle z' \neq z \rangle$  by auto
  qed
  ultimately show  $\exists x \in UNIV. x \neq z \wedge \text{dist}\ x\ z < d \wedge f\ x \neq 0$  by auto
qed
qed
ultimately have (THE  $n. \exists g\ r. P\ n\ g\ r$ ) = n
  by (rule_tac the1_equality)
then show ?thesis unfolding zorder_def P_def by blast
qed

```

lemma *simple_zeroI*:

```

assumes open s z  $z \in s$  g holomorphic_on s  $g\ z \neq 0$ 
assumes  $\bigwedge w. w \in s \implies f\ w = g\ w * (w - z)$ 
shows  $zorder\ f\ z = 1$ 
using assms zorder_eqI by force

```

lemma *higher_deriv_power*:

```

shows  $(\text{deriv } \sim j) (\lambda w. (w - z)^{\wedge n}) w =$ 
   $\text{pochhammer } (\text{of\_nat } (\text{Suc } n - j))\ j * (w - z)^{\wedge (n - j)}$ 
proof (induction j arbitrary: w)
  case 0
  thus ?case by auto
next
  case (Suc j w)
  have  $(\text{deriv } \sim \text{Suc } j) (\lambda w. (w - z)^{\wedge n}) w = \text{deriv } ((\text{deriv } \sim j) (\lambda w. (w - z)^{\wedge n})) w$ 
    by simp
  also have  $(\text{deriv } \sim j) (\lambda w. (w - z)^{\wedge n}) =$ 
     $(\lambda w. \text{pochhammer } (\text{of\_nat } (\text{Suc } n - j))\ j * (w - z)^{\wedge (n - j)})$ 
    using Suc by (intro Suc.IH ext)
  also {
    have ( $\dots$  has_field_derivative of_nat (n - j) *

```

```

      pochhammer (of_nat (Suc n - j)) j * (w - z) ^ (n - Suc j)) (at w)
    using Suc.premis by (auto intro!: derivative_eq_intros)
  also have of_nat (n - j) * pochhammer (of_nat (Suc n - j)) j =
      pochhammer (of_nat (Suc n - Suc j)) (Suc j)
    by (cases Suc j ≤ n, subst pochhammer_rec)
      (insert Suc.premis, simp_all add: algebra_simps Suc_diff_le pochhammer_0_left)
  finally have deriv (λw. pochhammer (of_nat (Suc n - j)) j * (w - z) ^ (n - j)) w =
      ... * (w - z) ^ (n - Suc j)
    by (rule DERIV_imp_deriv)
  }
  finally show ?case .
qed

```

lemma zorder_zero_eqI:

```

  assumes f_holo: f holomorphic_on s and open s z ∈ s
  assumes zero: ∧i. i < nat n ⇒ (deriv ^^ i) f z = 0
  assumes nz: (deriv ^^ nat n) f z ≠ 0 and n ≥ 0
  shows zorder f z = n

```

proof -

```

  obtain r where [simp]: r > 0 and ball z r ⊆ s

```

```

  using ⟨open s⟩ ⟨z ∈ s⟩ openE by blast

```

```

  have nz': ∃ w ∈ ball z r. f w ≠ 0

```

proof (rule ccontr)

```

  assume ¬ (∃ w ∈ ball z r. f w ≠ 0)

```

```

  then have eventually (λu. f u = 0) (nhds z)

```

```

    using open_ball ⟨0 < r⟩ centre_in_ball eventually_nhds by blast

```

```

  then have (deriv ^^ nat n) f z = (deriv ^^ nat n) (λ_. 0) z

```

```

    by (intro higher_deriv_cong_ev) auto

```

```

  also have (deriv ^^ nat n) (λ_. 0) z = 0

```

```

    by (induction n) simp_all

```

```

  finally show False using nz by contradiction

```

qed

```

  define zn g where zn = zorder f z and g = zor_poly f z

```

```

  obtain e where e_if: if f z = 0 then 0 < zn else zn = 0 and

```

```

    [simp]: e > 0 and cball z e ⊆ ball z r and

```

```

    g_holo: g holomorphic_on cball z e and

```

```

    e_fac: (∀ w ∈ cball z e. f w = g w * (w - z) ^ nat zn ∧ g w ≠ 0)

```

proof -

```

  have f_holo_on_ball z r

```

```

    using f_holo ⟨ball z r ⊆ s⟩ by auto

```

```

  from that zorder_exist_zero[of f ball z r z, simplified, OF this nz', folded zn_def g_def]

```

```

  show ?thesis by blast

```

qed

```

  then obtain zn ≥ 0 g z ≠ 0

```

```

    by (metis centre_in_cball less_le_not_le order_refl)

```

```

define A where A  $\equiv$  ( $\lambda i.$  of_nat (i choose (nat zn)) * fact (nat zn) * (deriv  $\sim$ 
(i - nat zn)) g z)
have deriv_A:(deriv  $\sim$  i) f z = (if zn  $\leq$  int i then A i else 0) for i
proof -
  have eventually ( $\lambda w.$  w  $\in$  ball z e) (nhds z)
    using <ball z e  $\subseteq$  ball z r> <e>0> by (intro eventually_nhds_in_open) auto
  hence eventually ( $\lambda w.$  f w = (w - z)  $\wedge$  (nat zn) * g w) (nhds z)
    using e_fac eventually_mono by fastforce
  hence (deriv  $\sim$  i) f z = (deriv  $\sim$  i) ( $\lambda w.$  (w - z)  $\wedge$  nat zn * g w) z
    by (intro higher_deriv_cong_ev) auto
  also have ... = ( $\sum_{j=0..i}$  of_nat (i choose j) *
    (deriv  $\sim$  j) ( $\lambda w.$  (w - z)  $\wedge$  nat zn) z * (deriv  $\sim$  (i - j)) g z)
    using g_holo <e>0>
    by (intro higher_deriv_mult[of _ ball z e]) (auto intro!: holomorphic_intros)
  also have ... = ( $\sum_{j=0..i}$  if j = nat zn then
    of_nat (i choose nat zn) * fact (nat zn) * (deriv  $\sim$  (i - nat zn))
g z else 0)
proof (intro sum.cong_refl, goal_cases)
  case (1 j)
  have (deriv  $\sim$  j) ( $\lambda w.$  (w - z)  $\wedge$  nat zn) z =
    pochhammer (of_nat (Suc (nat zn) - j)) j * 0  $\wedge$  (nat zn - j)
    by (subst higher_deriv_power) auto
  also have ... = (if j = nat zn then fact j else 0)
    by (auto simp: not_less pochhammer_0_left pochhammer_fact)
  also have of_nat (i choose j) * ... * (deriv  $\sim$  (i - j)) g z =
    (if j = nat zn then of_nat (i choose (nat zn)) * fact (nat zn)
    * (deriv  $\sim$  (i - nat zn)) g z else 0)
    by simp
  finally show ?case .
qed
also have ... = (if i  $\geq$  zn then A i else 0)
  by (auto simp: A_def)
finally show (deriv  $\sim$  i) f z = ... .
qed

have False when n < zn
proof -
  have (deriv  $\sim$  nat n) f z = 0
    using deriv_A[of nat n] that <n  $\geq$  0> by auto
  with nz show False by auto
qed
moreover have n  $\leq$  zn
proof -
  have g z  $\neq$  0 using e_fac[rule_format, of z] <e>0> by simp
  then have (deriv  $\sim$  nat zn) f z  $\neq$  0
    using deriv_A[of nat zn] by (auto simp add: A_def)
  then have nat zn  $\geq$  nat n using zero[of nat zn] by linarith
  moreover have zn  $\geq$  0 using e_if by (auto split: if_splits)

```

```

ultimately show ?thesis using nat_le_eq_zle by blast
qed
ultimately show ?thesis unfolding zn_def by fastforce
qed

lemma
  assumes eventually ( $\lambda z. f z = g z$ ) (at z)  $z = z'$ 
  shows zorder_cong: zorder  $f z = zorder g z'$  and zor_poly_cong: zor_poly  $f z =$ 
  zor_poly  $g z'$ 
proof -
  define P where  $P = (\lambda ff n h r. 0 < r \wedge h \text{ holomorphic\_on } cball z r \wedge h z \neq 0$ 
     $\wedge (\forall w \in cball z r - \{z\}. \text{ff } w = h w * (w - z) \text{ powi } n \wedge h w \neq 0))$ 
  have  $(\exists r. P f n h r) = (\exists r. P g n h r)$  for  $n h$ 
  proof -
    have *:  $\exists r. P g n h r$  if  $\exists r. P f n h r$  and eventually  $(\lambda x. f x = g x)$  (at z)
  for  $f g$ 
  proof -
    from that(1) obtain r1 where  $r1\_P: P f n h r1$  by auto
    from that(2) obtain r2 where  $r2 > 0$  and  $r2\_dist: \forall x. x \neq z \wedge dist x z \leq$ 
 $r2 \longrightarrow f x = g x$ 
    unfolding eventually_at_le by auto
    define r where  $r = \min r1 r2$ 
    have  $r > 0$   $h z \neq 0$  using r1_P <math>r2 > 0</math> unfolding r_def P_def by auto
    moreover have  $h \text{ holomorphic\_on } cball z r$ 
    using r1_P unfolding P_def r_def by auto
    moreover have  $g w = h w * (w - z) \text{ powi } n \wedge h w \neq 0$  when  $w \in cball z r$ 
  -  $\{z\}$  for  $w$ 
  proof -
    have  $f w = h w * (w - z) \text{ powi } n \wedge h w \neq 0$ 
    using r1_P that unfolding P_def r_def by auto
    moreover have  $f w = g w$  using r2_dist[rule_format, of w] that unfolding
  r_def
    by (simp add: dist_commute)
    ultimately show ?thesis by simp
  qed
  ultimately show ?thesis unfolding P_def by auto
qed
from assms have eq': eventually  $(\lambda z. g z = f z)$  (at z)
by (simp add: eq_commute)
show ?thesis
by (rule iffI[OF *[OF _ assms(1)] *[OF _ eq']])
qed
then show zorder  $f z = zorder g z'$  zor_poly  $f z = zor\_poly g z'$ 
using <math>z = z'</math> unfolding P_def zorder_def zor_poly_def by auto
qed

lemma zorder_times_analytic':
  assumes isolated_singularity_at  $f z$  not_essential  $f z$ 
  assumes  $g$  analytic_on  $\{z\}$  frequently  $(\lambda z. f z * g z \neq 0)$  (at z)

```

shows $\text{zorder } (\lambda x. f x * g x) z = \text{zorder } f z + \text{zorder } g z$
proof (rule *zorder_times*)
show *isolated_singularity_at* $g z$ *not_essential* $g z$
by (intro *isolated_singularity_at_analytic* *not_essential_analytic* *assms*)
qed (use *assms* in *auto*)

lemma *zorder_cmult*:

assumes $c \neq 0$
shows $\text{zorder } (\lambda z. c * f z) z = \text{zorder } f z$
proof –
define P **where**
 $P = (\lambda f n h r. 0 < r \wedge h \text{ holomorphic_on } \text{cball } z r \wedge$
 $h z \neq 0 \wedge (\forall w \in \text{cball } z r - \{z\}. f w = h w * (w - z) \text{ powi } n \wedge h w \neq$
 $0))$
have $*$: $P (\lambda x. c * f x) n (\lambda x. c * h x) r$ **if** $P f n h r c \neq 0$ **for** $f n h r c$
using *that* **unfolding** P_def **by** (*auto* *intro!*: *holomorphic_intros*)
have $(\exists h r. P (\lambda x. c * f x) n h r) \longleftrightarrow (\exists h r. P f n h r)$ **for** n
using $*$ [*of* $f n _ _ c$] $*$ [*of* $\lambda x. c * f x n _ _ \text{inverse } c$] $\langle c \neq 0 \rangle$
by (*fastforce* *simp*: *field_simps*)
hence (*THE* $n. \exists h r. P (\lambda x. c * f x) n h r$) = (*THE* $n. \exists h r. P f n h r$)
by *simp*
thus *?thesis*
by (*simp* *add*: *zorder_def* P_def)
qed

lemma *zorder_nonzero_div_power*:

assumes sz : *open* $s z \in s$ *f holomorphic_on* $s f z \neq 0$ **and** $n > 0$
shows $\text{zorder } (\lambda w. f w / (w - z) ^ n) z = - n$
by (intro *zorder_eqI* [*OF* sz]) (*simp* *add*: *inverse_eq_divide* *power_int_minus*)

lemma *zor_poly_eq*:

assumes *isolated_singularity_at* $f z$ *not_essential* $f z \exists_F w$ *in* *at* $z. f w \neq 0$
shows *eventually* $(\lambda w. \text{zor_poly } f z w = f w * (w - z) \text{ powi } - \text{zorder } f z)$ (*at* z)
proof –
obtain r **where** $r:r>0$
 $(\forall w \in \text{cball } z r - \{z\}. f w = \text{zor_poly } f z w * (w - z) \text{ powi } (\text{zorder } f z))$
using *zorder_exist* [*OF* *assms*] **by** *blast*
then **have** $*$: $\forall w \in \text{ball } z r - \{z\}. \text{zor_poly } f z w = f w * (w - z) \text{ powi } - \text{zorder } f z$
by (*auto* *simp*: *field_simps* *power_int_minus*)
have *eventually* $(\lambda w. w \in \text{ball } z r - \{z\})$ (*at* z)
using *r* *eventually_at_ball'* [*of* $r z UNIV$] **by** *auto*
thus *?thesis* **by** *eventually_elim* (*insert* $*$, *auto*)
qed

lemma *zor_poly_zero_eq*:

assumes *f holomorphic_on* s *open* s *connected* $s z \in s \exists w \in s. f w \neq 0$
shows *eventually* $(\lambda w. \text{zor_poly } f z w = f w / (w - z) ^ \text{nat } (\text{zorder } f z))$ (*at* z)
proof –

obtain r **where** $r:r>0$
 $(\forall w \in \text{cball } z \ r - \{z\}. f \ w = \text{zor_poly } f \ z \ w * (w - z) ^{\wedge} \text{nat } (\text{zorder } f \ z))$
using $\text{zorder_exist_zero}[OF \ \text{assms}]$ **by** auto
then have $*$: $\forall w \in \text{ball } z \ r - \{z\}. \text{zor_poly } f \ z \ w = f \ w / (w - z) ^{\wedge} \text{nat } (\text{zorder } f \ z)$
by $(\text{auto } \text{simp: field_simps } \text{powr_minus})$
have $\text{eventually } (\lambda w. w \in \text{ball } z \ r - \{z\})$ $(\text{at } z)$
using $r \ \text{eventually_at_ball}'[\text{of } r \ z \ \text{UNIV}]$ **by** auto
thus $?thesis$ **by** $\text{eventually_elim } (\text{insert } *, \text{auto})$
qed

lemma zor_poly_pole_eq :
assumes $f_iso:\text{isolated_singularity_at } f \ z \ \text{is_pole } f \ z$
shows $\text{eventually } (\lambda w. \text{zor_poly } f \ z \ w = f \ w * (w - z) ^{\wedge} \text{nat } (- \text{zorder } f \ z))$ $(\text{at } z)$
proof $-$
obtain e **where** $[simp]:e>0$ **and** $f_holo:f \ \text{holomorphic_on } \text{ball } z \ e - \{z\}$
using $f_iso \ \text{analytic_imp_holomorphic}$ **unfolding** $\text{isolated_singularity_at_def}$
by blast
obtain r **where** $r:r>0$
 $(\forall w \in \text{cball } z \ r - \{z\}. f \ w = \text{zor_poly } f \ z \ w / (w - z) ^{\wedge} \text{nat } (- \text{zorder } f \ z))$
using $\text{zorder_exist_pole}[OF \ f_holo, \text{simplified}, OF \ \langle \text{is_pole } f \ z \rangle]$ **by** auto
then have $*$: $\forall w \in \text{ball } z \ r - \{z\}. \text{zor_poly } f \ z \ w = f \ w * (w - z) ^{\wedge} \text{nat } (- \text{zorder } f \ z)$
by $(\text{auto } \text{simp: field_simps})$
have $\text{eventually } (\lambda w. w \in \text{ball } z \ r - \{z\})$ $(\text{at } z)$
using $r \ \text{eventually_at_ball}'[\text{of } r \ z \ \text{UNIV}]$ **by** auto
thus $?thesis$ **by** $\text{eventually_elim } (\text{insert } *, \text{auto})$
qed

lemma zor_poly_eqI :
fixes $f :: \text{complex} \Rightarrow \text{complex}$ **and** $z0 :: \text{complex}$
defines $n \equiv \text{zorder } f \ z0$
assumes $\text{isolated_singularity_at } f \ z0 \ \text{not_essential } f \ z0 \ \exists_F \ w \ \text{in } \text{at } z0. f \ w \neq 0$
assumes $\text{lim: } ((\lambda x. f \ (g \ x) * (g \ x - z0) \ \text{powi } - n) \longrightarrow c) \ F$
assumes $g: \text{filterlim } g \ (\text{at } z0) \ F$ **and** $F \neq \text{bot}$
shows $\text{zor_poly } f \ z0 \ z0 = c$
proof $-$
from $\text{zorder_exist}[OF \ \text{assms}(2-4)]$ **obtain** r **where**
 $r: r > 0 \ \text{zor_poly } f \ z0 \ \text{holomorphic_on } \text{cball } z0 \ r$
 $\bigwedge w. w \in \text{cball } z0 \ r - \{z0\} \Longrightarrow f \ w = \text{zor_poly } f \ z0 \ w * (w - z0) \ \text{powi } n$
unfolding n_def **by** blast
from $r(1)$ **have** $\text{eventually } (\lambda w. w \in \text{ball } z0 \ r \wedge w \neq z0)$ $(\text{at } z0)$
using $\text{eventually_at_ball}'[\text{of } r \ z0 \ \text{UNIV}]$ **by** auto
hence $\text{eventually } (\lambda w. \text{zor_poly } f \ z0 \ w = f \ w * (w - z0) \ \text{powi } - n)$ $(\text{at } z0)$
by $\text{eventually_elim } (\text{insert } r, \text{auto } \text{simp: field_simps } \text{power_int_minus})$
moreover have $\text{continuous_on } (\text{ball } z0 \ r)$ $(\text{zor_poly } f \ z0)$
using r **by** $(\text{intro } \text{holomorphic_on_imp_continuous_on}) \ \text{auto}$
with $r(1,2)$ **have** $\text{isCont } (\text{zor_poly } f \ z0) \ z0$

by (auto simp: continuous_on_eq_continuous_at)
 hence (zor_poly f z0 \longrightarrow zor_poly f z0 z0) (at z0)
 unfolding isCont_def .
 ultimately have (($\lambda w. f w * (w - z0)$ powi - n) \longrightarrow zor_poly f z0 z0) (at z0)
 by (blast intro: Lim_transform_eventually)
 hence (($\lambda x. f (g x) * (g x - z0)$ powi - n) \longrightarrow zor_poly f z0 z0) F
 by (rule filterlim_compose[OF _ g])
 from tendsto_unique[OF $\langle F \neq \text{bot} \rangle$ this lim] show ?thesis .
 qed

lemma zor_poly_zero_eqI:

fixes f :: complex \Rightarrow complex and z0 :: complex
 defines n \equiv zorder f z0
 assumes f holomorphic_on A open A connected A z0 \in A $\exists z \in A. f z \neq 0$
 assumes lim: (($\lambda x. f (g x) / (g x - z0)$ ^ nat n) \longrightarrow c) F
 assumes g: filterlim g (at z0) F and F \neq bot
 shows zor_poly f z0 z0 = c

proof -

from zorder_exist_zero[OF assms(2-6)] obtain r where

r: r > 0 cball z0 r \subseteq A zor_poly f z0 holomorphic_on cball z0 r

$\bigwedge w. w \in \text{cball } z0 \ r \implies f w = \text{zor_poly } f \ z0 \ w * (w - z0)^{\wedge} \text{ nat } n$

unfolding n_def by blast

from r(1) have eventually ($\lambda w. w \in \text{ball } z0 \ r \wedge w \neq z0$) (at z0)

using eventually_at_ball'[of r z0 UNIV] by auto

hence eventually ($\lambda w. \text{zor_poly } f \ z0 \ w = f w / (w - z0)^{\wedge} \text{ nat } n$) (at z0)

by eventually_elim (insert r, auto simp: field_simps)

moreover have continuous_on (ball z0 r) (zor_poly f z0)

using r by (intro holomorphic_on_imp_continuous_on) auto

with r(1,2) have isCont (zor_poly f z0) z0

by (auto simp: continuous_on_eq_continuous_at)

hence (zor_poly f z0 \longrightarrow zor_poly f z0 z0) (at z0)

unfolding isCont_def .

ultimately have (($\lambda w. f w / (w - z0)^{\wedge} \text{ nat } n$) \longrightarrow zor_poly f z0 z0) (at z0)

by (blast intro: Lim_transform_eventually)

hence (($\lambda x. f (g x) / (g x - z0)^{\wedge} \text{ nat } n$) \longrightarrow zor_poly f z0 z0) F

by (rule filterlim_compose[OF _ g])

from tendsto_unique[OF $\langle F \neq \text{bot} \rangle$ this lim] show ?thesis .

qed

lemma zor_poly_pole_eqI:

fixes f :: complex \Rightarrow complex and z0 :: complex

defines n \equiv zorder f z0

assumes f_iso: isolated_singularity_at f z0 and is_pole f z0

assumes lim: (($\lambda x. f (g x) * (g x - z0)^{\wedge} \text{ nat } (-n)$) \longrightarrow c) F

assumes g: filterlim g (at z0) F and F \neq bot

shows zor_poly f z0 z0 = c

proof -

obtain r where r: r > 0 zor_poly f z0 holomorphic_on cball z0 r

proof –
have $\exists_F w \text{ in at } z0. f w \neq 0$
using *non_zero_neighbour_pole*[*OF* $\langle is_pole f z0 \rangle$] **by** (*auto elim:eventually_frequentlyE*)
moreover have *not_essential_f z0 unfolding not_essential_def using* $\langle is_pole f z0 \rangle$ **by** *simp*
ultimately show *?thesis using that zorder_exist*[*OF* *f_iso,folded n_def*] **by** *auto*
qed
from *r(1)* **have** *eventually* $(\lambda w. w \in \text{ball } z0 \ r \wedge w \neq z0)$ (at *z0*)
using *eventually_at_ball'*[*of* *r z0 UNIV*] **by** *auto*
have *eventually* $(\lambda w. \text{zor_poly } f \ z0 \ w = f w * (w - z0) \wedge \text{nat } (-n))$ (at *z0*)
using *zor_poly_pole_eq*[*OF* *f_iso* $\langle is_pole f z0 \rangle$] **unfolding** *n_def* .
moreover have *continuous_on* $(\text{ball } z0 \ r)$ $(\text{zor_poly } f \ z0)$
using *r* **by** (*intro holomorphic_on_imp_continuous_on*) *auto*
with *r(1,2)* **have** *isCont* $(\text{zor_poly } f \ z0)$ *z0*
by (*auto simp: continuous_on_eq_continuous_at*)
hence $(\text{zor_poly } f \ z0 \longrightarrow \text{zor_poly } f \ z0 \ z0)$ (at *z0*)
unfolding *isCont_def* .
ultimately have $((\lambda w. f w * (w - z0) \wedge \text{nat } (-n)) \longrightarrow \text{zor_poly } f \ z0 \ z0)$ (at *z0*)
by (*blast intro: Lim_transform_eventually*)
hence $((\lambda x. f (g x) * (g x - z0) \wedge \text{nat } (-n)) \longrightarrow \text{zor_poly } f \ z0 \ z0)$ *F*
by (*rule filterlim_compose*[*OF* *_ g*])
from *tendsto_unique*[*OF* $\langle F \neq \text{bot} \rangle$ *this lim*] **show** *?thesis* .
qed

lemma

assumes *is_pole* *f* $(x :: \text{complex})$ *open* *A* $x \in A$
assumes $\bigwedge y. y \in A - \{x\} \implies (f \text{ has_field_derivative } f' y)$ (at *y*)
shows *is_pole_deriv'*: *is_pole* $f' x$
and *zorder_deriv'*: *zorder* $f' x = \text{zorder } f x - 1$
proof –
have *holo*: *f* *holomorphic_on* $A - \{x\}$
using *assms* **by** (*subst holomorphic_on_open*) *auto*
obtain *r* **where** *r*: $r > 0$ $\text{ball } x \ r \subseteq A$
using *assms*(2,3) *openE* **by** *blast*
moreover have *open* $(\text{ball } x \ r - \{x\})$
by *auto*
ultimately have *isolated_singularity_at* *f* *x*
by (*auto simp: isolated_singularity_at_def analytic_on_open*
intro!: exI[*of* *_ r*] *holomorphic_on_subset*[*OF* *holo*])
hence *ev*: $\forall_F w \text{ in at } x. \text{zor_poly } f \ x \ w = f w * (w - x) \wedge \text{nat } (-\text{zorder } f \ x)$
using $\langle is_pole f x \rangle$ *zor_poly_pole_eq* **by** *blast*

define *P* **where** $P = \text{zor_poly } f \ x$

define *n* **where** $n = \text{nat } (-\text{zorder } f \ x)$

obtain *r* **where** *r*: $r > 0$ $\text{cball } x \ r \subseteq A$ *P* *holomorphic_on* $\text{cball } x \ r$ *zorder* $f \ x < 0$ $P \ x \neq 0$


```

   $\forall w \in \text{cball } x \ r - \{x\}. f \ w = P \ w / (w - x)^{\wedge} n \wedge P \ w \neq 0$ 
  unfolding P_def n_def using zorder_exist_pole[OF holo_assms(2,3,1)] by
blast
  have n:  $n > 0$ 
  using r(4) by (auto simp: n_def)

  have [derivative_intros]: (P has_field_derivative deriv P w) (at w)
  if  $w \in \text{ball } x \ r$  for w
  using that by (intro holomorphic_derivI[OF holomorphic_on_subset[OF r(3),
of ball x r]]) auto

  define D where  $D = (\lambda w. (\text{deriv } P \ w * (w - x) - \text{of\_nat } n * P \ w) / (w - x)^{\wedge} (n + 1))$ 
  define n' where  $n' = n - 1$ 
  have n':  $n = \text{Suc } n'$ 
  using n by (simp add: n'_def)

  have eventually ( $\lambda w. w \in \text{ball } x \ r$ ) (nhds x)
  using  $\langle r > 0 \rangle$  by (intro eventually_nhds_in_open) auto
  hence ev'': eventually ( $\lambda w. w \in \text{ball } x \ r - \{x\}$ ) (at x)
  by (auto simp: eventually_at_filter elim: eventually_mono)

  {
  fix w assume w:  $w \in \text{ball } x \ r - \{x\}$ 
  have ev': eventually ( $\lambda w. w \in \text{ball } x \ r - \{x\}$ ) (nhds w)
  using w by (intro eventually_nhds_in_open) auto

  have (( $\lambda w. P \ w / (w - x)^{\wedge} n$ ) has_field_derivative D w) (at w)
  apply (rule derivative_eq_intros refl | use w in force)+
  using w
  apply (simp add: divide_simps D_def)
  apply (simp add: n' algebra_simps)
  done
  also have ?this  $\longleftrightarrow$  (f has_field_derivative D w) (at w)
  using r by (intro has_field_derivative_cong_ev refl eventually_mono[OF
ev'']) auto
  finally have (f has_field_derivative D w) (at w) .
  moreover have (f has_field_derivative f' w) (at w)
  using w r by (intro assms) auto
  ultimately have  $D \ w = f' \ w$ 
  using DERIV_unique by blast
  } note D_eq = this

  have is_pole D x
  unfolding D_def using n  $\langle r > 0 \rangle$   $\langle P \ x \neq 0 \rangle$ 
  by (intro is_pole_basic[where A = ball x r] holomorphic_intros holomor-
phic_on_subset[OF r(3)]) auto
  also have ?this  $\longleftrightarrow$  is_pole f' x
  by (intro is_pole_cong eventually_mono[OF ev''] D_eq) auto

```

```

finally show is_pole f' x .

have zorder f' x = -int (Suc n)
proof (rule zorder_eqI)
  show open (ball x r) x ∈ ball x r
    using  $\langle r > 0 \rangle$  by auto
  show  $f' w = (\text{deriv } P w * (w - x) - \text{of\_nat } n * P w) * (w - x)^{\text{powi } (- \text{int } (\text{Suc } n))}$ 
    if  $w \in \text{ball } x r$   $w \neq x$  for w
    using that D_eq[of w] n by (auto simp: D_def power_int_diff power_int_minus powr_nat' divide_simps)
  qed (use r n in  $\langle \text{auto intro!} \rangle$  holomorphic_intros)
  thus zorder f' x = zorder f x - 1
    using n by (simp add: n_def)
qed

```

lemma

```

assumes is_pole f (x :: complex) isolated_singularity_at f x
shows is_pole_deriv: is_pole (deriv f) x
  and zorder_deriv: zorder (deriv f) x = zorder f x - 1
proof -
  from assms(2) obtain r where  $r > 0$  f analytic_on ball x r - {x}
    by (auto simp: isolated_singularity_at_def)
  hence holo: f holomorphic_on ball x r - {x}
    by (subst (asm) analytic_on_open) auto
  have  $*$ :  $x \in \text{ball } x r$  open (ball x r) open (ball x r - {x})
    using  $\langle r > 0 \rangle$  by auto
  show is_pole (deriv f) x zorder (deriv f) x = zorder f x - 1
    by (rule is_pole_deriv' zorder_deriv', (rule assms * holomorphic_derivI holo
| assumption)+)+
qed

```

lemma *removable_singularity_deriv'*:

```

assumes  $f -x \rightarrow c$   $x \in A$  open (A :: complex set)
assumes  $\bigwedge y. y \in A - \{x\} \implies (f \text{ has\_field\_derivative } f' y) (at y)$ 
shows  $\exists c. f' -x \rightarrow c$ 

```

proof -

```

have holo: f holomorphic_on A - {x}
  using assms by (subst holomorphic_on_open) auto

```

```

define g where  $g = (\lambda y. \text{if } y = x \text{ then } c \text{ else } f y)$ 

```

```

have deriv_g_eq: deriv g y = f' y if  $y \in A - \{x\}$  for y

```

proof -

```

have ev: eventually  $(\lambda y. y \in A - \{x\})$  (nhds y)
  using that assms by (intro eventually_nhds_in_open) auto

```

```

have (f has_field_derivative f' y) (at y)

```

```

  using assms that by auto

```

```

also have  $?this \longleftrightarrow (g \text{ has\_field\_derivative } f' y) (at y)$ 

```

```

  by (intro has_field_derivative_cong_ev refl eventually_mono[OF ev]) (auto)

```

```

simp: g_def)
  finally show ?thesis
    by (intro DERIV_imp_deriv assms)
qed

have g holomorphic_on A
  unfolding g_def using assms assms(1) holo by (intro removable_singularity)
auto
hence deriv g holomorphic_on A
  by (intro holomorphic_deriv assms)
hence continuous_on A (deriv g)
  by (meson holomorphic_on_imp_continuous_on)
hence (deriv g  $\longrightarrow$  deriv g x) (at x within A)
  using assms by (auto simp: continuous_on_def)
also have ?this  $\longleftrightarrow$  (f'  $\longrightarrow$  deriv g x) (at x within A)
  by (intro filterlim_cong refl) (auto simp: eventually_at_filter deriv_g_eq)
finally have f'  $-x \rightarrow$  deriv g x
  using <open A> <x  $\in$  A> by (meson tendsto_within_open)
thus ?thesis
  by blast
qed

lemma removable_singularity_deriv:
  assumes f  $-x \rightarrow$  c isolated_singularity_at f x
  shows  $\exists c. \text{deriv } f -x \rightarrow c$ 
proof -
  from assms(2) obtain r where r: r > 0 f analytic_on ball x r - {x}
    by (auto simp: isolated_singularity_at_def)
  hence holo: f holomorphic_on ball x r - {x}
    using analytic_imp_holomorphic by blast
  show ?thesis
    using assms(1)
  proof (rule removable_singularity_deriv')
    show x  $\in$  ball x r open (ball x r)
      using <r > 0> by auto
    qed (auto intro!: holomorphic_derivI[OF holo])
  qed

lemma not_essential_deriv':
  assumes not_essential f x x  $\in$  A open A
  assumes  $\bigwedge y. y \in A - \{x\} \implies (f \text{ has\_field\_derivative } f' y) (at y)$ 
  shows not_essential f' x
proof -
  have holo: f holomorphic_on A - {x}
    using assms by (subst holomorphic_on_open) auto
  from assms consider is_pole f x | c where f  $-x \rightarrow$  c
    by (auto simp: not_essential_def)
  thus ?thesis
  proof cases

```

```

case 1
hence is_pole  $f' x$ 
  using is_pole_deriv' assms by blast
thus ?thesis by (auto simp: not_essential_def)
next
case (2 c)
from 2 have  $\exists c. f' -x \rightarrow c$ 
  by (rule removable_singularity_deriv'[OF _ assms(2-4)])
thus ?thesis
  by (auto simp: not_essential_def)
qed
qed

```

```

lemma not_essential_deriv[singularity_intros]:
  assumes not_essential  $f x$  isolated_singularity_at  $f x$ 
  shows not_essential (deriv  $f$ )  $x$ 
proof -
  from assms(2) obtain  $r$  where  $r > 0$  f analytic_on ball  $x r - \{x\}$ 
    by (auto simp: isolated_singularity_at_def)
  hence holo: f holomorphic_on ball  $x r - \{x\}$ 
    by (subst (asm) analytic_on_open) auto
  show ?thesis
    using assms(1)
  proof (rule not_essential_deriv')
    show  $x \in \text{ball } x r \text{ open } (\text{ball } x r)$ 
      using  $\langle r > 0 \rangle$  by auto
    qed (auto intro!: holomorphic_derivI[OF holo])
qed

```

```

lemma not_essential_frequently_0_imp_tendsto_0:
  fixes  $f :: \text{complex} \Rightarrow \text{complex}$ 
  assumes sing: isolated_singularity_at  $f z$  not_essential  $f z$ 
  assumes freq: frequently  $(\lambda z. f z = 0)$  (at  $z$ )
  shows  $f -z \rightarrow 0$ 
proof -
  from freq obtain  $g :: \text{nat} \Rightarrow \text{complex}$  where filterlim  $g$  (at  $z$ ) at_top  $\wedge n. f$ 
    ( $g n$ ) = 0
    using frequently_atE by blast
  have eventually  $(\lambda x. f (g x) = 0)$  sequentially
    using  $g$  by auto
  hence  $fg: (\lambda x. f (g x)) \longrightarrow 0$ 
    by (simp add: tendsto_eventually)

  from assms(2) consider  $c$  where  $f -z \rightarrow c$  | is_pole  $f z$ 
    unfolding not_essential_def by blast
  thus ?thesis
proof cases
  case (1 c)
    have  $(\lambda x. f (g x)) \longrightarrow c$ 

```

```

    by (rule filterlim_compose[OF 1 g(1)])
  with fg have c = 0
    using LIMSEQ_unique by blast
  with 1 show ?thesis by simp
next
case 2
have filterlim ( $\lambda x. f (g x)$ ) at_infinity sequentially
  by (rule filterlim_compose[OF _ g(1)]) (use 2 in ‹auto simp: is_pole_def›)
with fg have False
  by (meson not_tendsto_and_filterlim_at_infinity_sequentially_bot)
thus ?thesis ..
qed
qed

```

lemma not_essential_frequently_0_imp_eventually_0:

```

  fixes f :: complex  $\Rightarrow$  complex
  assumes sing: isolated_singularity_at f z not_essential f z
  assumes freq: frequently ( $\lambda z. f z = 0$ ) (at z)
  shows eventually ( $\lambda z. f z = 0$ ) (at z)
proof -
  from sing obtain r where r:  $r > 0$  and f_analytic_on_ball z r - {z}
    by (auto simp: isolated_singularity_at_def)
  hence holo: f_holomorphic_on_ball z r - {z}
    by (subst (asm) analytic_on_open) auto
  have eventually ( $\lambda w. w \in \text{ball } z \ r - \{z\}$ ) (at z)
    using r by (intro eventually_at_in_open) auto
  from freq and this have frequently ( $\lambda w. f w = 0 \wedge w \in \text{ball } z \ r - \{z\}$ ) (at z)
    using frequently_eventually_frequently by blast
  hence frequently ( $\lambda w. w \in \{w \in \text{ball } z \ r - \{z\}. f w = 0\}$ ) (at z)
    by (simp add: conj_commute)
  hence limpt: z_islimpt {w ∈ ball z r - {z}. f w = 0}
    using islimpt_conv_frequently_at by blast

  define g where g = ( $\lambda w. \text{if } w = z \text{ then } 0 \text{ else } f w$ )
  have f_neg_z → 0
    by (intro not_essential_frequently_0_imp_tendsto_0 assms)
  hence g_holo: g_holomorphic_on_ball z r
    unfolding g_def by (intro removable_singularity_holo) auto

  have g_eq_0: g w = 0 if w ∈ ball z r for w
  proof (rule analytic_continuation[where f = g])
    show open (ball z r) connected (ball z r)
      using r by auto
    show z_islimpt {w ∈ ball z r - {z}. f w = 0}
      by fact
    show g w = 0 if w ∈ {w ∈ ball z r - {z}. f w = 0} for w
      using that by (auto simp: g_def)
  qed (use r that g_holo in auto)

```

```

have eventually ( $\lambda w. w \in \text{ball } z \ r - \{z\}$ ) (at z)
  using r by (intro eventually_at_in_open) auto
thus eventually ( $\lambda w. f \ w = 0$ ) (at z)
proof eventually_elim
  case (elim w)
  thus ?case using g_eq_0[of w]
    by (auto simp: g_def)
qed
qed

```

lemma pole_imp_not_constant:

```

fixes f :: 'a :: {perfect_space}  $\Rightarrow$  _
assumes is_pole f x open A x  $\in$  A A  $\subseteq$  insert x B
shows  $\neg$ f constant_on B
proof
assume *: f constant_on B
then obtain c where c:  $\forall x \in B. f \ x = c$ 
  by (auto simp: constant_on_def)
have eventually ( $\lambda y. y \in A - \{x\}$ ) (at x)
  using assms by (intro eventually_at_in_open) auto
hence eventually ( $\lambda y. f \ y = c$ ) (at x)
  by eventually_elim (use c assms in auto)
hence **: f  $-x \rightarrow c$ 
  by (simp add: tendsto_eventually)
show False
  using not_tendsto_and_filterlim_at_infinity[OF_ ** assms(1)[unfolded is_pole_def]]
by simp
qed

```

lemma neg_zorder_imp_is_pole:

```

assumes iso:isolated_singularity_at f z and f_ness:not_essential f z
  and zorder f z < 0 and fre_nz: $\exists_F w$  in at z. f w  $\neq$  0
shows is_pole f z
proof -
define P where P = zorder_poly f z
define n where n = zorder f z
have n < 0 unfolding n_def by (simp add: assms(3))
define nn where nn = nat (-n)

obtain r where P z  $\neq$  0 r > 0 and r_holo:P holomorphic_on cball z r and
  w_Pn:( $\forall w \in \text{cball } z \ r - \{z\}. f \ w = P \ w * (w - z)^{\text{powi } n} \wedge P \ w \neq 0$ )
  using zorder_exist[OF iso f_ness fre_nz,folded P_def n_def] by auto

have is_pole ( $\lambda w. P \ w * (w - z)^{\text{powi } n}$ ) z
  unfolding is_pole_def
proof (rule tendsto_mult_filterlim_at_infinity)
show P  $-z \rightarrow P \ z$ 
  by (meson open_ball ‹0 < r› ball_subset_cball centre_in_ball

```

```

    continuous_on_eq_continuous_at continuous_on_subset
    holomorphic_on_imp_continuous_on isContD r_holo)
show  $P z \neq 0$  by (simp add:  $\langle P z \neq 0 \rangle$ )

have LIM  $x$  at  $z$ . inverse  $((x - z) \wedge^{\text{nat}} (-n)) :> \text{at\_infinity}$ 
  apply (subst filterlim_inverse_at_iff[symmetric])
  using  $\langle n < 0 \rangle$ 
  by (auto intro!: tendsto_eq_intros filterlim_atI
      simp add: eventually_at_filter)
then show LIM  $x$  at  $z$ .  $(x - z) \text{ powi } n :> \text{at\_infinity}$ 
proof (elim filterlim_mono_eventually)
  have inverse  $((x - z) \wedge^{\text{nat}} (-n)) = (x - z) \text{ powi } n$ 
    if  $x \neq z$  for  $x$ 
  by (metis  $\langle n < 0 \rangle$  linorder_not_le power_int_def power_inverse)
  then show  $\forall_F x$  in at  $z$ . inverse  $((x - z) \wedge^{\text{nat}} (-n))$ 
    =  $(x - z) \text{ powi } n$ 
    by (simp add: eventually_at_filter)
qed auto
qed
moreover have  $\forall_F w$  in at  $z$ .  $f w = P w * (w - z) \text{ powi } n$ 
  unfolding eventually_at_le
  apply (rule exI[where  $x=r$ ])
  using  $w \_Pn \langle r > 0 \rangle$  by (simp add: dist_commute)
ultimately show ?thesis using is_pole_cong by fast
qed

lemma is_pole_divide_zorder:
  fixes  $f g :: \text{complex} \Rightarrow \text{complex}$  and  $z :: \text{complex}$ 
  assumes  $f\_iso: \text{isolated\_singularity\_at } f z$  and  $g\_iso: \text{isolated\_singularity\_at } g z$ 
  and  $f\_ness: \text{not\_essential } f z$  and  $g\_ness: \text{not\_essential } g z$ 
  and  $fg\_nconst: \exists_F w$  in  $(\text{at } z)$ .  $f w * g w \neq 0$ 
  and  $z\_less: \text{zorder } f z < \text{zorder } g z$ 
  shows is_pole  $(\lambda z. f z / g z) z$ 
proof -
  define  $fn gn fg$  where  $fn = \text{zorder } f z$  and  $gn = \text{zorder } g z$ 
  and  $fg = (\lambda w. f w / g w)$ 

  have isolated_singularity_at  $fg z$ 
  unfolding  $fg\_def$  using  $f\_iso g\_iso g\_ness$ 
  by (auto intro: singularity_intros)
  moreover have not_essential  $fg z$ 
  unfolding  $fg\_def$  using  $f\_iso g\_iso g\_ness f\_ness$ 
  by (auto intro: singularity_intros)
  moreover have  $\text{zorder } fg z < 0$ 
proof -
  have  $\text{zorder } fg z = fn - gn$ 
  using  $\text{zorder\_divide}[OF f\_iso g\_iso f\_ness g\_ness fg\_nconst, folded fn\_def gn\_def fg\_def]$  .

```

```

    then show ?thesis
      using z_less by (simp add: fn_def gn_def)
    qed
  moreover have  $\exists_F w$  in at z.  $fg\ w \neq 0$ 
    using fg_nconst unfolding fg_def by force
  ultimately show is_pole fg z
    using neg_zorder_imp_is_pole by auto
  qed

lemma isolated_pole_imp_nzero_times:
  assumes f_iso:isolated_singularity_at f z
    and is_pole f z
  shows  $\exists_F w$  in (at z).  $deriv\ f\ w * f\ w \neq 0$ 
proof (rule ccontr)
  assume  $\neg (\exists_F w$  in at z.  $deriv\ f\ w * f\ w \neq 0)$ 
  then have  $\forall_F x$  in at z.  $deriv\ f\ x * f\ x = 0$ 
    unfolding not_frequently by simp
  moreover have  $\forall_F w$  in at z.  $f\ w \neq 0$ 
    using non_zero_neighbour_pole[OF  $\langle is\_pole\ f\ z \rangle$ ].
  moreover have  $\forall_F w$  in at z.  $deriv\ f\ w \neq 0$ 
    using is_pole_deriv[OF  $\langle is\_pole\ f\ z \rangle$  f_iso, THEN non_zero_neighbour_pole]
  .
  ultimately have  $\forall_F w$  in at z. False
    apply eventually_elim
    by auto
  then show False by auto
  qed

lemma isolated_pole_imp_neg_zorder:
  assumes f_iso:isolated_singularity_at f z
    and is_pole f z
  shows zorder f z < 0
proof -
  obtain e where [simp]:  $e > 0$  and f_holo: f holomorphic_on ball z e - {z}
    using f_iso analytic_imp_holomorphic unfolding isolated_singularity_at_def
  by blast
  show ?thesis
    using zorder_exist_pole[OF f_holo, simplified, OF  $\langle is\_pole\ f\ z \rangle$ ]
    by auto
  qed

lemma isolated_singularity_at_deriv[singularity_intros]:
  assumes isolated_singularity_at f x
  shows isolated_singularity_at (deriv f) x
proof -
  obtain r where  $r > 0$  f_analytic_on ball x r - {x}
    using assms unfolding isolated_singularity_at_def by auto
  from analytic_deriv[OF this(2)]
  have deriv_f_analytic_on ball x r - {x} .

```



```

with  $\langle r > 0 \rangle$  show ?thesis
  unfolding isolated_singularity_at_def by auto
qed

lemma zorder_deriv_minus_1:
  fixes  $f g :: \text{complex} \Rightarrow \text{complex}$  and  $z :: \text{complex}$ 
  assumes  $f\_iso: \text{isolated\_singularity\_at } f z$ 
    and  $f\_ness: \text{not\_essential } f z$ 
    and  $f\_nconst: \exists_F w \text{ in at } z. f w \neq 0$ 
    and  $f\_ord: \text{zorder } f z \neq 0$ 
  shows  $\text{zorder } (\text{deriv } f) z = \text{zorder } f z - 1$ 
proof -
  define  $P$  where  $P = \text{zor\_poly } f z$ 
  define  $n$  where  $n = \text{zorder } f z$ 
  have  $n \neq 0$  unfolding  $n\_def$  using  $f\_ord$  by auto

  obtain  $r$  where  $P z \neq 0$   $r > 0$  and  $P\_holo: P \text{ holomorphic\_on } \text{cball } z r$ 
    and  $(\forall w \in \text{cball } z r - \{z\}. f w = P w * (w - z)^{\text{powi } n} \wedge P w \neq 0)$ 
  using  $\text{zorder\_exist}[OF f\_iso f\_ness f\_nconst, folded P\_def n\_def]$  by auto
  from  $\text{this}(4)$ 
  have  $f\_eq: (\forall w \in \text{cball } z r - \{z\}. f w = P w * (w - z)^{\text{powi } n} \wedge P w \neq 0)$ 
    using  $\text{complex\_powr\_of\_int } f\_ord n\_def$  by presburger

  define  $D$  where  $D = (\lambda w. (\text{deriv } P w * (w - z) + \text{of\_int } n * P w) * (w - z)^{\text{powi } (n - 1)})$ 

  have  $\text{deriv } f\_eq: \text{deriv } f w = D w$  if  $w \in \text{ball } z r - \{z\}$  for  $w$ 
  proof -
  have  $ev'$ : eventually  $(\lambda w. w \in \text{ball } z r - \{z\})$   $(\text{nhds } w)$ 
    using that by  $(\text{intro eventually\_nhds\_in\_open})$  auto

  define  $wz$  where  $wz = w - z$ 

  have  $wz \neq 0$  unfolding  $wz\_def$  using that by auto
  moreover have  $(P \text{ has\_field\_derivative } \text{deriv } P w)$   $(\text{at } w)$ 
    by  $(\text{meson } \text{DiffD1 } \text{Elementary\_Metric\_Spaces.open\_ball } P\_holo$ 
       $\text{ball\_subset\_cball } \text{holomorphic\_derivI } \text{holomorphic\_on\_subset } \text{that})$ 
  ultimately have  $((\lambda w. P w * (w - z)^{\text{powi } n}) \text{ has\_field\_derivative } D w)$   $(\text{at } w)$ 
    unfolding  $D\_def$  using that
    apply  $(\text{auto intro!: derivative\_eq\_intros})$ 
    apply  $(\text{fold } wz\_def)$ 
    by  $(\text{auto simp: algebra\_simps simp flip: power\_int\_add\_1'})$ 
  also have  $?this \longleftrightarrow (f \text{ has\_field\_derivative } D w)$   $(\text{at } w)$ 
    using  $f\_eq$ 
    by  $(\text{intro has\_field\_derivative\_cong\_ev refl eventually\_mono}[OF ev'])$  auto
  ultimately have  $(f \text{ has\_field\_derivative } D w)$   $(\text{at } w)$  by simp

```

```

moreover have (f has_field_derivative deriv f w) (at w)
  by (metis DERIV_imp_deriv calculation)
ultimately show ?thesis using DERIV_imp_deriv by blast
qed

show zorder (deriv f) z = n - 1
proof (rule zorder_eqI)
  show open (ball z r) z ∈ ball z r
    using ⟨r > 0⟩ by auto
  define g where g=(λw. (deriv P w * (w - z) + of_int n * P w))
  show g holomorphic_on ball z r
    unfolding g_def using P_holo
    by (auto intro!:holomorphic_intros)
  show g z ≠ 0
    unfolding g_def using ⟨P z ≠ 0⟩ ⟨n≠0⟩ by auto
  show deriv f w =
    (deriv P w * (w - z) + of_int n * P w) * (w - z) powi (n - 1)
    if w ∈ ball z r w ≠ z for w
    using D_def deriv_f_eq that by blast
qed
qed

lemma deriv_divide_is_pole: — Generalises  $\llbracket is\_pole\ ?f\ ?x; isolated\_singularity\_at\ ?f\ ?x \rrbracket \implies zorder\ (deriv\ ?f)\ ?x = zorder\ ?f\ ?x - 1$ 
  fixes f g::complex ⇒ complex and z::complex
  assumes f_iso:isolated_singularity_at f z
    and f_ness:not_essential f z
    and fg_nconst: ∃Fw in (at z). deriv f w * f w ≠ 0
    and f_ord:zorder f z ≠ 0
  shows is_pole (λz. deriv f z / f z) z
proof (rule neg_zorder_imp_is_pole)
  define ff where ff=(λw. deriv f w / f w)
  show isolated_singularity_at ff z
    using f_iso f_ness unfolding ff_def
    by (auto intro:singularity_intros)
  show not_essential ff z
    unfolding ff_def using f_ness f_iso
    by (auto intro:singularity_intros)

  have zorder ff z = zorder (deriv f) z - zorder f z
    unfolding ff_def using f_iso f_ness fg_nconst
    apply (rule_tac zorder_divide)
    by (auto intro:singularity_intros)
  moreover have zorder (deriv f) z = zorder f z - 1
  proof (rule zorder_deriv_minus_1)
    show ∃F w in at z. f w ≠ 0
      using fg_nconst frequently_elim1 by fastforce
  qed (use f_iso f_ness f_ord in auto)

```

ultimately show $\text{zorder } ff \ z < 0$ **by auto**

show $\exists_F w \text{ in at } z. ff \ w \neq 0$
unfolding ff_def **using** fg_nconst **by auto**
qed

lemma $is_pole_deriv_divide_is_pole$:
fixes $f \ g :: complex \Rightarrow complex$ **and** $z :: complex$
assumes $f_iso: isolated_singularity_at \ f \ z$
and $is_pole \ f \ z$
shows $is_pole \ (\lambda z. deriv \ f \ z / f \ z) \ z$
proof ($rule \ deriv_divide_is_pole[OF \ f_iso]$)
show $not_essential \ f \ z$
using $\langle is_pole \ f \ z \rangle$ **unfolding** $not_essential_def$ **by auto**
show $\exists_F w \text{ in at } z. deriv \ f \ w * f \ w \neq 0$
apply ($rule \ isolated_pole_imp_nzero_times$)
using $assms$ **by auto**
show $\text{zorder } f \ z \neq 0$
using $isolated_pole_imp_neg_zorder \ assms$ **by fastforce**
qed

5.12 Isolated zeroes

definition $isolated_zero :: (complex \Rightarrow complex) \Rightarrow complex \Rightarrow bool$ **where**
 $isolated_zero \ f \ z \longleftrightarrow f \ z = 0 \wedge eventually \ (\lambda z. f \ z \neq 0) \ (at \ z)$

lemma $isolated_zero_altdef$: $isolated_zero \ f \ z \longleftrightarrow f \ z = 0 \wedge \neg z \ islimpt \ \{z. f \ z = 0\}$
unfolding $isolated_zero_def$ $eventually_at_filter$ $eventually_nhds$ $islimpt_def$
by blast

lemma $isolated_zero_mult1$:
assumes $isolated_zero \ f \ x$ $isolated_zero \ g \ x$
shows $isolated_zero \ (\lambda x. f \ x * g \ x) \ x$
proof –
have $eventually \ (\lambda x. f \ x \neq 0) \ (at \ x)$ $eventually \ (\lambda x. g \ x \neq 0) \ (at \ x)$
using $assms$ **unfolding** $isolated_zero_def$ **by auto**
hence $eventually \ (\lambda x. f \ x * g \ x \neq 0) \ (at \ x)$
by $eventually_elim \ auto$
with $assms$ **show** $?thesis$
by ($auto \ simp: isolated_zero_def$)
qed

lemma $isolated_zero_mult2$:
assumes $isolated_zero \ f \ x$ $g \ x \neq 0$ $g \ analytic_on \ \{x\}$
shows $isolated_zero \ (\lambda x. f \ x * g \ x) \ x$
proof –
have $eventually \ (\lambda x. f \ x \neq 0) \ (at \ x)$
using $assms$ **unfolding** $isolated_zero_def$ **by auto**

```

moreover have eventually ( $\lambda x. g\ x \neq 0$ ) (at x)
  using analytic_at_neq_imp_eventually_neq[of g x 0] assms by auto
ultimately have eventually ( $\lambda x. f\ x * g\ x \neq 0$ ) (at x)
  by eventually_elim auto
thus ?thesis
  using assms(1) by (auto simp: isolated_zero_def)
qed

```

```

lemma isolated_zero_mult3:
  assumes isolated_zero f x g x  $\neq 0$  g analytic_on {x}
  shows isolated_zero ( $\lambda x. g\ x * f\ x$ ) x
  using isolated_zero_mult2[OF assms] by (simp add: mult_ac)

```

```

lemma isolated_zero_prod:
  assumes  $\bigwedge x. x \in I \implies$  isolated_zero (f x) z  $I \neq \{\}$  finite I
  shows isolated_zero ( $\lambda y. \prod_{x \in I}. f\ x\ y$ ) z
  using assms(3,2,1) by (induction rule: finite_ne_induct) (auto intro: isolated_zero_mult1)

```

```

lemma non_isolated_zero':
  assumes isolated_singularity_at f z not_essential f z f z = 0  $\neg$ isolated_zero f z
  shows eventually ( $\lambda z. f\ z = 0$ ) (at z)
proof (rule not_essential_frequently_0_imp_eventually_0)
  from assms show frequently ( $\lambda z. f\ z = 0$ ) (at z)
  by (auto simp: frequently_def isolated_zero_def)
qed fact+

```

```

lemma non_isolated_zero:
  assumes  $\neg$ isolated_zero f z f analytic_on {z} f z = 0
  shows eventually ( $\lambda z. f\ z = 0$ ) (nhds z)
proof -
  have eventually ( $\lambda z. f\ z = 0$ ) (at z)
    by (rule non_isolated_zero')
    (use assms in  $\langle$ auto intro: not_essential_analytic_isolated_singularity_at_analytic $\rangle$ )
  with  $\langle f\ z = 0 \rangle$  show ?thesis
  unfolding eventually_at_filter by (auto elim!: eventually_mono)
qed

```

```

lemma not_essential_compose:
  assumes not_essential f (g z) g analytic_on {z}
  shows not_essential ( $\lambda x. f\ (g\ x)$ ) z
proof (cases isolated_zero ( $\lambda w. g\ w - g\ z$ ) z)
  case False
  hence eventually ( $\lambda w. g\ w - g\ z = 0$ ) (nhds z)
    by (rule non_isolated_zero) (use assms in  $\langle$ auto intro!: analytic_intros $\rangle$ )
  hence not_essential ( $\lambda x. f\ (g\ x)$ ) z  $\longleftrightarrow$  not_essential ( $\lambda \_. f\ (g\ z)$ ) z
    by (intro not_essential_cong_refl)
    (auto elim!: eventually_mono simp: eventually_at_filter)
  thus ?thesis
    by (simp add: not_essential_const)

```

```

next
  case True
  hence ev: eventually ( $\lambda w. g w \neq g z$ ) (at z)
    by (auto simp: isolated_zero_def)
  from assms consider c where  $f - g z \rightarrow c \mid is\_pole f (g z)$ 
    by (auto simp: not_essential_def)
  have isCont g z
    by (rule analytic_at_imp_isCont) fact
  hence lim:  $g - z \rightarrow g z$ 
    using isContD by blast

  from assms(1) consider c where  $f - g z \rightarrow c \mid is\_pole f (g z)$ 
    unfolding not_essential_def by blast
  thus ?thesis
  proof cases
    fix c assume  $f - g z \rightarrow c$ 
    hence ( $\lambda x. f (g x)$ )  $-z \rightarrow c$ 
      by (rule filterlim_compose) (use lim ev in <auto simp: filterlim_at>)
    thus ?thesis
      by (auto simp: not_essential_def)
  next
    assume is_pole f (g z)
    hence is_pole ( $\lambda x. f (g x)$ ) z
      by (rule is_pole_compose) fact+
    thus ?thesis
      by (auto simp: not_essential_def)
  qed
qed

```

5.13 Isolated points

definition *isolated_points_of* :: complex set \Rightarrow complex set **where**
 $isolated_points_of A = \{z \in A. eventually (\lambda w. w \notin A) (at z)\}$

lemma *isolated_points_of_altdef*: $isolated_points_of A = \{z \in A. \neg z \text{ islimpt } A\}$
unfolding *isolated_points_of_def islimpt_def eventually_at_filter eventually_nhds*
by *blast*

lemma *isolated_points_of_empty* [simp]: $isolated_points_of \{\} = \{\}$
and *isolated_points_of_UNIV* [simp]: $isolated_points_of UNIV = \{\}$
by (auto simp: isolated_points_of_def)

lemma *isolated_points_of_open_is_empty* [simp]: $open A \implies isolated_points_of A = \{\}$
unfolding *isolated_points_of_altdef*
by (simp add: interior_limit_point interior_open)

lemma *isolated_points_of_subset*: $isolated_points_of A \subseteq A$
by (auto simp: isolated_points_of_def)

```

lemma isolated_points_of_discrete:
  assumes discrete A
  shows isolated_points_of A = A
  using assms by (auto simp: isolated_points_of_def discrete_altdef)

lemmas uniform_discreteI1 = uniformI1
lemmas uniform_discreteI2 = uniformI2

lemma isolated_singularity_at_compose:
  assumes isolated_singularity_at f (g z) g analytic_on {z}
  shows isolated_singularity_at ( $\lambda x. f (g x)$ ) z
proof (cases isolated_zero ( $\lambda w. g w - g z$ ) z)
  case False
  hence eventually ( $\lambda w. g w - g z = 0$ ) (nhds z)
    by (rule non_isolated_zero) (use assms in <auto intro!: analytic_intros>)
  hence isolated_singularity_at ( $\lambda x. f (g x)$ ) z  $\longleftrightarrow$  isolated_singularity_at ( $\lambda \_.$ 
  f (g z)) z
    by (intro isolated_singularity_at_cong refl)
      (auto elim!: eventually_mono simp: eventually_at_filter)
  thus ?thesis
    by (simp add: isolated_singularity_at_const)
next
  case True
  from assms(1) obtain r where r: r > 0 f analytic_on ball (g z) r - {g z}
    by (auto simp: isolated_singularity_at_def)
  hence holo_f: f holomorphic_on ball (g z) r - {g z}
    by (subst (asm) analytic_on_open) auto
  from assms(2) obtain r' where r': r' > 0 g holomorphic_on ball z r'
    by (auto simp: analytic_on_def)

  have continuous_on (ball z r') g
    using holomorphic_on_imp_continuous_on r' by blast
  hence isCont g z
    using r' by (subst (asm) continuous_on_eq_continuous_at) auto
  hence g - z  $\rightarrow$  g z
    using isContD by blast
  hence eventually ( $\lambda w. g w \in \text{ball } (g z) r$ ) (at z)
    using <r > 0 unfolding tendsto_def by force
  moreover have eventually ( $\lambda w. g w \neq g z$ ) (at z) using True
    by (auto simp: isolated_zero_def elim!: eventually_mono)
  ultimately have eventually ( $\lambda w. g w \in \text{ball } (g z) r - \{g z\}$ ) (at z)
    by eventually_elim auto
  then obtain r'' where r'': r'' > 0  $\forall w \in \text{ball } z r'' - \{z\}. g w \in \text{ball } (g z) r - \{g$ 
  z}
    unfolding eventually_at_filter eventually_nhds_metric ball_def
    by (auto simp: dist_commute)
  have f  $\circ$  g holomorphic_on ball z (min r' r'') - {z}
  proof (rule holomorphic_on_compose_gen)

```

```

show  $g$  holomorphic_on ball  $z$  ( $\min r' r''$ ) -  $\{z\}$ 
  by (rule holomorphic_on_subset[OF  $r'(2)$ ]) auto
show  $f$  holomorphic_on ball ( $g z$ )  $r - \{g z\}$ 
  by fact
show  $g^{-1}(\text{ball } z (\min r' r'') - \{z\}) \subseteq \text{ball } (g z) r - \{g z\}$ 
  using  $r''$  by force
qed
hence  $f \circ g$  analytic_on ball  $z$  ( $\min r' r''$ ) -  $\{z\}$ 
  by (subst analytic_on_open) auto
thus ?thesis using  $\langle r' > 0 \rangle \langle r'' > 0 \rangle$ 
  by (auto simp: isolated_singularity_at_def o_def intro!: exI[of _  $\min r' r''$ ])
qed

```

```

lemma is_pole_power_int_0:
  assumes  $f$  analytic_on  $\{x\}$  isolated_zero  $f x n < 0$ 
  shows is_pole  $(\lambda x. f x \text{ powi } n) x$ 
proof -
  have  $f -x \rightarrow f x$ 
    using  $\text{assms}(1)$  by (simp add: analytic_at_imp_isCont isContD)
  with  $\text{assms}$  show ?thesis
    unfolding is_pole_def
    by (intro filterlim_power_int_neg_at_infinity) (auto simp: isolated_zero_def)
qed

```

```

lemma isolated_zero_imp_not_constant_on:
  assumes isolated_zero  $f x x \in A$  open  $A$ 
  shows  $\neg f$  constant_on  $A$ 
proof
  assume  $f$  constant_on  $A$ 
  then obtain  $c$  where  $c: \bigwedge x. x \in A \implies f x = c$ 
    by (auto simp: constant_on_def)
  from  $\text{assms}$  and  $c[\text{of } x]$  have  $[simp]: c = 0$ 
    by (auto simp: isolated_zero_def)
  have eventually  $(\lambda x. f x \neq 0)$  (at  $x$ )
    using  $\text{assms}$  by (auto simp: isolated_zero_def)
  moreover have eventually  $(\lambda x. x \in A)$  (at  $x$ )
    using  $\text{assms}$  by (intro eventually_at_in_open') auto
  ultimately have eventually  $(\lambda x. \text{False})$  (at  $x$ )
    by eventually_elim (use  $c$  in auto)
  thus False
    by simp
qed

```

```

end
theory Complex_Residues
  imports Complex_Singularities
begin

```

5.14 Definition of residues

Wenda Li and LC Paulson (2016). A Formal Proof of Cauchy's Residue Theorem. Interactive Theorem Proving

definition *residue* :: (complex \Rightarrow complex) \Rightarrow complex \Rightarrow complex **where**
residue $f z = (\text{SOME } \text{int}. \exists e > 0. \forall \varepsilon > 0. \varepsilon < e \longrightarrow (f \text{ has_contour_integral } 2 * \text{pi} * i * \text{int}) (\text{circlepath } z \ \varepsilon))$

lemma *residue_cong*:

assumes *eq*: *eventually* ($\lambda z. f z = g z$) (*at* z) **and** $z = z'$
shows *residue* $f z = \text{residue } g z'$

proof –

from *assms* **have** *eq'*: *eventually* ($\lambda z. g z = f z$) (*at* z)

by (*simp add: eq_commute*)

let $?P = \lambda f c e. (\forall \varepsilon > 0. \varepsilon < e \longrightarrow$

$(f \text{ has_contour_integral_of_real } (2 * \text{pi}) * i * c) (\text{circlepath } z \ \varepsilon))$

have *residue* $f z = \text{residue } g z$ **unfolding** *residue_def*

proof (*rule Eps_cong*)

fix $c :: \text{complex}$

have $\exists e > 0. ?P g c e$

if $\exists e > 0. ?P f c e$ **and** *eventually* ($\lambda z. f z = g z$) (*at* z) **for** $f g$

proof –

from *that*(1) **obtain** e **where** $e > 0$ $?P f c e$

by *blast*

from *that*(2) **obtain** e' **where** $e' > 0 \wedge z'. z' \neq z \implies \text{dist } z' z < e' \implies$

$f z' = g z'$

unfolding *eventually_at* **by** *blast*

have $?P g c (\text{min } e e')$

proof (*intro allI exI impI, goal_cases*)

case (1 ε)

hence $(f \text{ has_contour_integral_of_real } (2 * \text{pi}) * i * c) (\text{circlepath } z \ \varepsilon)$

using $e(2)$ **by** *auto*

thus *?case*

proof (*rule has_contour_integral_eq*)

fix z' **assume** $z' \in \text{path_image } (\text{circlepath } z \ \varepsilon)$

hence $\text{dist } z' z < e'$ **and** $z' \neq z$

using 1 **by** (*auto simp: dist_commute*)

with $e'(2)[\text{of } z']$ **show** $f z' = g z'$ **by** *simp*

qed

qed

moreover **from** e **and** e' **have** $\text{min } e e' > 0$ **by** *auto*

ultimately **show** *?thesis* **by** *blast*

qed

from *this*[*OF _ eq*] **and** *this*[*OF _ eq*]

show $(\exists e > 0. ?P f c e) \longleftrightarrow (\exists e > 0. ?P g c e)$

by *blast*

qed

with *assms* **show** *?thesis* **by** *simp*

qed


```

lemma contour_integral_circlepath_eq:
  assumes open s and f_holo:f holomorphic_on (s - {z}) and  $0 < e1 \leq e2$ 
  and e2_cball:cball z e2  $\subseteq$  s
  shows
    f contour_integrable_on circlepath z e1
    f contour_integrable_on circlepath z e2
    contour_integral (circlepath z e2) f = contour_integral (circlepath z e1) f
proof -
  define l where  $l \equiv \text{linepath } (z + e2) (z + e1)$ 
  have [simp]:valid_path l pathstart l = z + e2 pathfinish l = z + e1 unfolding l_def
by auto
  have  $e2 > 0$  using  $\langle e1 > 0 \rangle \langle e1 \leq e2 \rangle$  by auto
  have z_l_img: $z \notin \text{path\_image } l$ 
  proof
    assume  $z \in \text{path\_image } l$ 
    then have  $e2 \leq \text{cmod } (e2 - e1)$ 
    using segment_furthest_le[of z z + e2 z + e1 z + e2, simplified]  $\langle e1 > 0 \rangle \langle e2 > 0 \rangle$ 
unfolding l_def
    by (auto simp add: closed_segment_commute)
    thus False using  $\langle e2 > 0 \rangle \langle e1 > 0 \rangle \langle e1 \leq e2 \rangle$ 
    apply (subst (asm) norm_of_real)
    by auto
  qed
  define g where  $g \equiv \text{circlepath } z \ e2 \ +++ \ l \ +++ \ \text{reversepath } (\text{circlepath } z \ e1)$ 
   $+++ \ \text{reversepath } l$ 
  show [simp]: f contour_integrable_on circlepath z e2 f contour_integrable_on
  (circlepath z e1)
  proof -
    show f contour_integrable_on circlepath z e2
    apply (intro contour_integrable_continuous_circlepath[OF
    continuous_on_subset[OF holomorphic_on_imp_continuous_on[OF
f_holo]]])
    using  $\langle e2 > 0 \rangle \ e2\_cball$  by auto
    show f contour_integrable_on (circlepath z e1)
    apply (intro contour_integrable_continuous_circlepath[OF
    continuous_on_subset[OF holomorphic_on_imp_continuous_on[OF
f_holo]]])
    using  $\langle e1 > 0 \rangle \langle e1 \leq e2 \rangle \ e2\_cball$  by auto
  qed
  have [simp]:f contour_integrable_on l
  proof -
    have closed_segment (z + e2) (z + e1)  $\subseteq$  cball z e2 using  $\langle e2 > 0 \rangle \langle e1 > 0 \rangle$ 
 $\langle e1 \leq e2 \rangle$ 
    by (intro closed_segment_subset, auto simp add: dist_norm)
    hence closed_segment (z + e2) (z + e1)  $\subseteq$  s - {z} using z_l_img e2_cball
unfolding l_def
    by auto
    then show f contour_integrable_on l unfolding l_def

```

```

apply (intro contour_integrable_continuous_linepath[OF
        continuous_on_subset[OF holomorphic_on_imp_continuous_on[OF
f_holo]]])
by auto
qed
let ?ig= $\lambda g.$  contour_integral g f
have (f has_contour_integral 0) g
proof (rule Cauchy_theorem_global[OF _ f_holo])
show open (s - {z}) using ‹open s› by auto
show valid_path g unfolding g_def l_def by auto
show pathfinish g = pathstart g unfolding g_def l_def by auto
next
have path_img:path_image g  $\subseteq$  cball z e2
proof -
have closed_segment (z + e2) (z + e1)  $\subseteq$  cball z e2 using ‹e2>0› ‹e1>0›
‹e1 $\leq$ e2›
by (intro closed_segment_subset,auto simp add:dist_norm)
moreover have sphere z |e1|  $\subseteq$  cball z e2 using ‹e2>0› ‹e1 $\leq$ e2› ‹e1>0›
by auto
ultimately show ?thesis unfolding g_def l_def using ‹e2>0›
by (simp add: path_image_join closed_segment_commute)
qed
show path_image g  $\subseteq$  s - {z}
proof -
have z $\notin$ path_image g using zl_img
unfolding g_def l_def by (auto simp add: path_image_join closed_segment_commute)
moreover note ‹cball z e2  $\subseteq$  s› and path_img
ultimately show ?thesis by auto
qed
show winding_number g w = 0 when w  $\notin$  s - {z} for w
proof -
have winding_number g w = 0 when w $\notin$ s using that e2_cball
apply (intro winding_number_zero_outside[OF _ _ _ _ path_img])
by (auto simp add:g_def l_def)
moreover have winding_number g z=0
proof -
let ?Wz= $\lambda g.$  winding_number g z
have ?Wz g = ?Wz (circlepath z e2) + ?Wz l + ?Wz (reversepath
(circlepath z e1))
+ ?Wz (reversepath l)
using ‹e2>0› ‹e1>0› zl_img unfolding g_def l_def
by (subst winding_number_join,auto simp add:path_image_join
closed_segment_commute)+
also have ... = ?Wz (circlepath z e2) + ?Wz (reversepath (circlepath
z e1))
using zl_img
apply (subst (2) winding_number_reversepath)
by (auto simp add:l_def closed_segment_commute)
also have ... = 0

```

```

proof –
  have ?Wz (circlepath z e2) = 1 using ⟨e2>0⟩
    by (auto intro: winding_number_circlepath_centre)
  moreover have ?Wz (reversepath (circlepath z e1)) = -1 using
⟨e1>0⟩
    apply (subst winding_number_reversepath)
    by (auto intro: winding_number_circlepath_centre)
  ultimately show ?thesis by auto
qed
finally show ?thesis .
qed
ultimately show ?thesis using that by auto
qed
qed
then have 0 = ?ig g using contour_integral_unique by simp
also have ... = ?ig (circlepath z e2) + ?ig l + ?ig (reversepath (circlepath z e1))
  + ?ig (reversepath l)
  unfolding g_def
  by (auto simp add: contour_integrable_reversepath_eq)
also have ... = ?ig (circlepath z e2) - ?ig (circlepath z e1)
  by (auto simp add: contour_integral_reversepath)
finally show contour_integral (circlepath z e2) f = contour_integral (circlepath
z e1) f
  by simp
qed

lemma base_residue:
  assumes open s z∈s r>0 and f_holo:f holomorphic_on (s - {z})
  and r_cball:cball z r ⊆ s
  shows (f has_contour_integral 2 * pi * i * (residue f z)) (circlepath z r)
proof –
  obtain e where e>0 and e_cball:cball z e ⊆ s
  using open_contains_cball[of s] ⟨open s⟩ ⟨z∈s⟩ by auto
  define c where c ≡ 2 * pi * i
  define i where i ≡ contour_integral (circlepath z e) f / c
  have (f has_contour_integral c*i) (circlepath z e) when ε>0 ε<e for ε
  proof –
    have contour_integral (circlepath z e) f = contour_integral (circlepath z ε) f
      f contour_integrable_on circlepath z ε
      f contour_integrable_on circlepath z e
    using ⟨ε<e⟩
    by (intro contour_integral_circlepath_eq[OF ⟨open s⟩ f_holo ⟨ε>0⟩ _
e_cball],auto)+
    then show ?thesis unfolding i_def c_def
    by (auto intro:has_contour_integral_integral)
  qed
  then have ∃ e>0. ∀ ε>0. ε<e → (f has_contour_integral c * (residue f z))
(circlepath z ε)
  unfolding residue_def c_def

```

```

  apply (rule_tac someI[of _ i],intro exI[where x=e])
  by (auto simp add:⟨e>0⟩ c_def)
then obtain e' where e'>0
  and e'_def:∀ε>0. ε<e' ⟶ (f has_contour_integral c * (residue f z))
(circlepath z ε)
  by auto
let ?int=λe. contour_integral (circlepath z e) f
define ε where ε ≡ Min {r,e'} / 2
have ε>0 ε≤r ε<e' using ⟨r>0⟩ ⟨e'>0⟩ unfolding ε_def by auto
have (f has_contour_integral c * (residue f z)) (circlepath z ε)
  using e'_def[rule_format,OF ⟨ε>0⟩ ⟨ε<e'⟩] .
then show ?thesis unfolding c_def
  using contour_integral_circlepath_eq[OF ⟨open s⟩ f_holo ⟨ε>0⟩ ⟨ε≤r⟩ r_cball]
  by (auto elim: has_contour_integral_eqpath[of _ _ circlepath z ε circlepath z
r])
qed

```

lemma residue_holo:

```

  assumes open s z ∈ s and f_holo: f holomorphic_on s
  shows residue f z = 0
proof -
  define c where c ≡ 2 * pi * i
  obtain e where e>0 and e_cball:cball z e ⊆ s using ⟨open s⟩ ⟨z∈s⟩
  using open_contains_cball_eq by blast
  have (f has_contour_integral c*residue f z) (circlepath z e)
  using f_holo
  by (auto intro: base_residue[OF ⟨open s⟩ ⟨z∈s⟩ ⟨e>0⟩ _ e_cball,folded c_def])
  moreover have (f has_contour_integral 0) (circlepath z e)
  using f_holo e_cball ⟨e>0⟩
  by (auto intro: Cauchy_theorem_convex_simple[of _ cball z e])
  ultimately have c*residue f z = 0
  using has_contour_integral_unique by blast
  thus ?thesis unfolding c_def by auto
qed

```

lemma residue_const:residue (λ_. c) z = 0

```

  by (intro residue_holo[of UNIV::complex set],auto intro:holomorphic_intros)

```

lemma residue_add:

```

  assumes open s z ∈ s and f_holo: f holomorphic_on s - {z}
  and g_holo:g holomorphic_on s - {z}
  shows residue (λz. f z + g z) z = residue f z + residue g z
proof -
  define c where c ≡ 2 * pi * i
  define fg where fg ≡ (λz. f z + g z)
  obtain e where e>0 and e_cball:cball z e ⊆ s using ⟨open s⟩ ⟨z∈s⟩
  using open_contains_cball_eq by blast
  have (fg has_contour_integral c * residue fg z) (circlepath z e)
  unfolding fg_def using f_holo g_holo

```

```

apply (intro base_residue[OF ‹open s› ‹z∈s› ‹e>0› _ e_cball,folded c_def])
by (auto intro:holomorphic_intros)
moreover have (fg has_contour_integral c*residue f z + c* residue g z) (circlepath
z e)
  unfolding fg_def using f_holo g_holo
  by (auto intro: has_contour_integral_add base_residue[OF ‹open s› ‹z∈s›
‹e>0› _ e_cball,folded c_def])
ultimately have c*(residue f z + residue g z) = c * residue fg z
  using has_contour_integral_unique by (auto simp add:distrib_left)
thus ?thesis unfolding fg_def
  by (auto simp add:c_def)
qed

```

lemma residue_lmul:

```

assumes open s z ∈ s and f_holo: f holomorphic_on s - {z}
shows residue (λz. c * (f z)) z = c * residue f z
proof (cases c=0)
  case True
    thus ?thesis using residue_const by auto
  next
    case False
      define c' where c' ≡ 2 * pi * i
      define f' where f' ≡ (λz. c * (f z))
      obtain e where e>0 and e_cball: cball z e ⊆ s using ‹open s› ‹z∈s›
        using open_contains_cball_eq by blast
      have (f' has_contour_integral c' * residue f' z) (circlepath z e)
        unfolding f'_def using f_holo
        apply (intro base_residue[OF ‹open s› ‹z∈s› ‹e>0› _ e_cball,folded c'_def])
        by (auto intro:holomorphic_intros)
      moreover have (f' has_contour_integral c * (c' * residue f z)) (circlepath z e)
        unfolding f'_def using f_holo
        by (auto intro: has_contour_integral_lmul
          base_residue[OF ‹open s› ‹z∈s› ‹e>0› _ e_cball,folded c'_def])
      ultimately have c' * residue f' z = c * (c' * residue f z)
        using has_contour_integral_unique by auto
      thus ?thesis unfolding f'_def c'_def using False
        by (auto simp add:field_simps)
    qed

```

lemma residue_rmul:

```

assumes open s z ∈ s and f_holo: f holomorphic_on s - {z}
shows residue (λz. (f z) * c) z = residue f z * c
using residue_lmul[OF assms,of c] by (auto simp add:algebra_simps)

```

lemma residue_div:

```

assumes open s z ∈ s and f_holo: f holomorphic_on s - {z}
shows residue (λz. (f z) / c) z = residue f z / c
using residue_lmul[OF assms,of 1/c] by (auto simp add:algebra_simps)

```

lemma *residue_neg*:

assumes *open s z ∈ s* **and** *f_holo: f holomorphic_on s - {z}*
shows *residue (λz. - (f z)) z = - residue f z*
using *residue_lmMul[OF assms, of -1]* **by** *auto*

lemma *residue_diff*:

assumes *open s z ∈ s* **and** *f_holo: f holomorphic_on s - {z}*
and *g_holo: g holomorphic_on s - {z}*
shows *residue (λz. f z - g z) z = residue f z - residue g z*
using *residue_add[OF assms(1,2,3), of λz. - g z]* *residue_neg[OF assms(1,2,4)]*
by (*auto intro: holomorphic_intros g_holo*)

lemma *residue_simple*:

assumes *open s z ∈ s* **and** *f_holo: f holomorphic_on s*
shows *residue (λw. f w / (w - z)) z = f z*
proof -
define *c* **where** *c ≡ 2 * pi * i*
define *f'* **where** *f' ≡ λw. f w / (w - z)*
obtain *e* **where** *e > 0* **and** *e_cball: cball z e ⊆ s* **using** *⟨open s⟩ ⟨z ∈ s⟩*
using *open_contains_cball_eq* **by** *blast*
have (*f' has_contour_integral c * f z*) (*circlepath z e*)
unfolding *f'_def c_def* **using** *⟨e > 0⟩ f_holo e_cball*
by (*auto intro!: Cauchy_integral_circlepath_simple holomorphic_intros*)
moreover **have** (*f' has_contour_integral c * residue f' z*) (*circlepath z e*)
unfolding *f'_def* **using** *f_holo*
apply (*intro base_residue[OF ⟨open s⟩ ⟨z ∈ s⟩ ⟨e > 0⟩ _ e_cball, folded c_def]*)
by (*auto intro!: holomorphic_intros*)
ultimately **have** *c * f z = c * residue f' z*
using *has_contour_integral_unique* **by** *blast*
thus *?thesis* **unfolding** *c_def f'_def* **by** *auto*
qed

lemma *residue_simple'*:

assumes *s: open s z ∈ s* **and** *holo: f holomorphic_on (s - {z})*
and *lim: ((λw. f w * (w - z)) -> c) (at z)*
shows *residue f z = c*
proof -
define *g* **where** *g = (λw. if w = z then c else f w * (w - z))*
from *holo* **have** (*λw. f w * (w - z)*) *holomorphic_on (s - {z})* (**is** *?P*)
by (*force intro: holomorphic_intros*)
also **have** *?P ⟷ g holomorphic_on (s - {z})*
by (*intro holomorphic_cong refl*) (*simp_all add: g_def*)
finally **have** **: g holomorphic_on (s - {z})* .

note *lim*

also **have** (*λw. f w * (w - z)*) *-z → c ⟷ g -z → g z*
by (*intro filterlim_cong refl*) (*simp_all add: g_def [abs_def] eventually_at_filter*)
finally **have** ***:* *g -z → g z* .

```

have g_holo: g holomorphic_on s
  by (rule no_isolated_singularity'[where K = {z}])
    (insert assms * **, simp_all add: at_within_open_NO_MATCH)
from s and this have residue ( $\lambda w. g w / (w - z)$ ) z = g z
  by (rule residue_simple)
also have  $\forall_F za$  in at z. g za / (za - z) = f za
  unfolding eventually_at by (auto intro!: exI[of _ 1] simp: field_simps g_def)
hence residue ( $\lambda w. g w / (w - z)$ ) z = residue f z
  by (intro residue_cong refl)
finally show ?thesis
  by (simp add: g_def)
qed

```

```

lemma residue_holomorphic_over_power:
  assumes open A z0  $\in$  A f holomorphic_on A
  shows residue ( $\lambda z. f z / (z - z0) ^ Suc n$ ) z0 = (deriv  $\hat{\hat{}}$  n) f z0 / fact n
proof -
  let ?f =  $\lambda z. f z / (z - z0) ^ Suc n$ 
  from assms(1,2) obtain r where r: r > 0 cball z0 r  $\subseteq$  A
  by (auto simp: open_contains_cball)
  have (?f has_contour_integral 2 * pi * i * residue ?f z0) (circlepath z0 r)
  using r assms by (intro base_residue[of A]) (auto intro!: holomorphic_intros)
  moreover have (?f has_contour_integral 2 * pi * i / fact n * (deriv  $\hat{\hat{}}$  n) f
z0) (circlepath z0 r)
  using assms r
  by (intro Cauchy_has_contour_integral_higher_derivative_circlepath)
    (auto intro!: holomorphic_on_subset[OF assms(3)] holomorphic_on_imp_continuous_on)
  ultimately have 2 * pi * i * residue ?f z0 = 2 * pi * i / fact n * (deriv  $\hat{\hat{}}$  n)
f z0
  by (rule has_contour_integral_unique)
  thus ?thesis by (simp add: field_simps)
qed

```

```

lemma residue_holomorphic_over_power':
  assumes open A 0  $\in$  A f holomorphic_on A
  shows residue ( $\lambda z. f z / z ^ Suc n$ ) 0 = (deriv  $\hat{\hat{}}$  n) f 0 / fact n
  using residue_holomorphic_over_power[OF assms] by simp

```

```

theorem residue_fps_expansion_over_power_at_0:
  assumes f has_fps_expansion F
  shows residue ( $\lambda z. f z / z ^ Suc n$ ) 0 = fps_nth F n
proof -
  from has_fps_expansion_imp_holomorphic[OF assms] obtain s
  where open s 0  $\in$  s f holomorphic_on s  $\wedge$  z. z  $\in$  s  $\implies$  f z = eval_fps F z
  by auto
  with assms have residue ( $\lambda z. f z / (z - 0) ^ Suc n$ ) 0 = (deriv  $\hat{\hat{}}$  n) f 0 / fact
n
  unfolding has_fps_expansion_def
  by (intro residue_holomorphic_over_power[of s]) (auto simp: zero_ereal_def)

```

also from *assms* have ... = *fps_nth F n*
 by (*subst fps_nth_fps_expansion*) *auto*
 finally show ?*thesis* by *simp*
 qed

lemma *residue_pole_order*:

fixes *f::complex* \Rightarrow *complex* **and** *z::complex*
defines *n* \equiv *nat* ($-$ *zorder f z*) **and** *h* \equiv *zor_poly f z*
assumes *f_iso:isolated_singularity_at f z*
and *pole:is_pole f z*
shows *residue f z* = $((\text{deriv } \hat{\sim} (n - 1)) h z / \text{fact } (n-1))$
proof $-$
define *g* **where** *g* \equiv $\lambda x.$ if $x=z$ then 0 else *inverse (f x)*
obtain *e* **where** [*simp*]: $e>0$ **and** *f_holo:f holomorphic_on ball z e - {z}*
using *f_iso analytic_imp_holomorphic* **unfolding** *isolated_singularity_at_def*
by *blast*
obtain *r* **where** $0 < n$ $0 < r$ **and** *r_cball:cball z r \subseteq ball z e* **and** *h_holo: h holomorphic_on cball z r*
and *h_divide:($\forall w \in \text{cball } z r. (w \neq z \longrightarrow f w = h w / (w - z) ^ n) \wedge h w \neq 0$)*
proof $-$
obtain *r* **where** *r:zorder f z < 0* *h z $\neq 0$* $r>0$ *cball z r \subseteq ball z e* *h holomorphic_on cball z r*
 $(\forall w \in \text{cball } z r - \{z\}. f w = h w / (w - z) ^ n \wedge h w \neq 0)$
using *zorder_exist_pole[OF f_holo,simplified,OF <is_pole f z>,folded n_def h_def]* **by** *auto*
have $n>0$ **using** $\langle \text{zorder } f z < 0 \rangle$ **unfolding** *n_def* **by** *simp*
moreover **have** $(\forall w \in \text{cball } z r. (w \neq z \longrightarrow f w = h w / (w - z) ^ n) \wedge h w \neq 0)$
using $\langle h z \neq 0 \rangle$ *r(6)* **by** *blast*
ultimately show ?*thesis* **using** *r(3,4,5)* **that** **by** *blast*
 qed
have *r_nonzero: $\bigwedge w. w \in \text{ball } z r - \{z\} \implies f w \neq 0$*
using *h_divide* **by** *simp*
define *c* **where** *c* $\equiv 2 * \text{pi} * i$
define *der_f* **where** *der_f* $\equiv ((\text{deriv } \hat{\sim} (n - 1)) h z / \text{fact } (n-1))$
define *h'* **where** *h'* $\equiv \lambda u. h u / (u - z) ^ n$
have $(h' \text{ has_contour_integral } c / \text{fact } (n - 1) * (\text{deriv } \hat{\sim} (n - 1)) h z)$
 $(\text{circlepath } z r)$
unfolding *h'_def*
proof (*rule Cauchy_has_contour_integral_higher_derivative_circlepath*[*of z r h z n-1,*
folded c_def Suc_pred'[OF <n>0]])
show *continuous_on (cball z r) h* **using** *holomorphic_on_imp_continuous_on h_holo* **by** *simp*
show *h holomorphic_on ball z r* **using** *h_holo* **by** *auto*
show $z \in \text{ball } z r$ **using** $\langle r>0 \rangle$ **by** *auto*
 qed
then **have** $(h' \text{ has_contour_integral } c * \text{der}_f)$ $(\text{circlepath } z r)$ **unfolding**
der_f_def **by** *auto*


```

then have (f has_contour_integral c * der_f) (circlepath z r)
  proof (elim has_contour_integral_eq)
    fix x assume x ∈ path_image (circlepath z r)
    hence x ∈ cball z r - {z} using ‹r>0› by auto
    then show h' x = f x using h_divide unfolding h'_def by auto
  qed
moreover have (f has_contour_integral c * residue f z) (circlepath z r)
  using base_residue[of ‹ball z e› z,simplified,OF ‹r>0› f_holo r_cball,folded
c_def]
  unfolding c_def by simp
ultimately have c * der_f = c * residue f z using has_contour_integral_unique
by blast
hence der_f = residue f z unfolding c_def by auto
thus ?thesis unfolding der_f_def by auto
qed

```

```

lemma residue_simple_pole:
  assumes isolated_singularity_at f z0
  assumes is_pole f z0 zorder f z0 = - 1
  shows residue f z0 = zor_poly f z0 z0
  using assms by (subst residue_pole_order) simp_all

```

```

lemma residue_simple_pole_limit:
  assumes isolated_singularity_at f z0
  assumes is_pole f z0 zorder f z0 = - 1
  assumes ((λx. f (g x) * (g x - z0)) → c) F
  assumes filterlim g (at z0) F F ≠ bot
  shows residue f z0 = c
proof -
  have residue f z0 = zor_poly f z0 z0
    by (rule residue_simple_pole assms)+
  also have ... = c
    apply (rule zor_poly_pole_eqI)
    using assms by auto
  finally show ?thesis .
qed

```

```

lemma
  assumes f_holo:f holomorphic_on s and g_holo:g holomorphic_on s
  and open s connected s z ∈ s
  assumes g_deriv:(g has_field_derivative g') (at z)
  assumes f z ≠ 0 g z = 0 g' ≠ 0
  shows porder_simple_pole_deriv: zorder (λw. f w / g w) z = - 1
  and residue_simple_pole_deriv: residue (λw. f w / g w) z = f z / g'
proof -
  have [simp]:isolated_singularity_at f z isolated_singularity_at g z
  using isolated_singularity_at_holomorphic[OF ‹open s› ‹z ∈ s›] f_holo g_holo
  by (meson Diff_subset holomorphic_on_subset)+
  have [simp]:not_essential f z not_essential g z

```

```

  unfolding not_essential_def using f_holo g_holo assms(3,5)
  by (meson continuous_on_eq_continuous_at continuous_within holomorphic_on_imp_continuous_on)+
  have g_nconst:  $\exists_F w \text{ in } at z. g w \neq 0$ 
  proof (rule ccontr)
    assume  $\neg (\exists_F w \text{ in } at z. g w \neq 0)$ 
    then have  $\forall_F w \text{ in } nhds z. g w = 0$ 
      unfolding eventually_at eventually_nhds frequently_at using  $\langle g z = 0 \rangle$ 
      by (metis open_ball UNIV_I centre_in_ball dist_commute mem_ball)
    then have  $deriv\ g\ z = deriv\ (\lambda_. 0)\ z$ 
      by (intro deriv_cong_ev) auto
    then have  $deriv\ g\ z = 0$  by auto
    then have  $g' = 0$  using g_deriv DERIV_imp_deriv by blast
    then show False using  $\langle g' \neq 0 \rangle$  by auto
  qed

  have zorder  $(\lambda w. f w / g w)\ z = zorder\ f\ z - zorder\ g\ z$ 
  proof -
    have  $\forall_F w \text{ in } at z. f w \neq 0 \wedge w \in s$ 
      apply (rule non_zero_neighbour_alt)
      using assms by auto
    with g_nconst have  $\exists_F w \text{ in } at z. f w * g w \neq 0$ 
      by (elim frequently_rev_mp eventually_rev_mp, auto)
    then show ?thesis using zorder_divide[of f z g] by auto
  qed

  moreover have  $zorder\ f\ z = 0$ 
    apply (rule zorder_zero_eqI[OF f_holo  $\langle open\ s \rangle$   $\langle z \in s \rangle$ ])
    using  $\langle f z \neq 0 \rangle$  by auto
  moreover have  $zorder\ g\ z = 1$ 
    apply (rule zorder_zero_eqI[OF g_holo  $\langle open\ s \rangle$   $\langle z \in s \rangle$ ])
    subgoal using assms(8) by auto
    subgoal using DERIV_imp_deriv assms(9) g_deriv by auto
    subgoal by simp
    done
  ultimately show  $zorder\ (\lambda w. f w / g w)\ z = - 1$  by auto

  show residue  $(\lambda w. f w / g w)\ z = f z / g'$ 
  proof (rule residue_simple_pole_limit[where  $g=id$  and  $F=at\ z, simplified$ ])
    show  $zorder\ (\lambda w. f w / g w)\ z = - 1$  by fact
    show isolated_singularity_at  $(\lambda w. f w / g w)\ z$ 
      by (auto intro: singularity_intros)
    show is_pole  $(\lambda w. f w / g w)\ z$ 
    proof (rule is_pole_divide)
      have  $\forall_F x \text{ in } at z. g x \neq 0$ 
        apply (rule non_zero_neighbour)
        using g_nconst by auto
      moreover have  $g -z \rightarrow 0$ 
        using DERIV_isCont assms(8) continuous_at g_deriv by force
      ultimately show filterlim  $g$  (at 0) (at z) unfolding filterlim_at by simp
    show isCont f z
  
```

```

      using assms(3,5) continuous_on_eq_continuous_at f_holo holomor-
phic_on_imp_continuous_on
    by auto
    show  $f z \neq 0$  by fact
  qed
  show filterlim id (at z) (at z) by (simp add: filterlim_iff)
  have  $((\lambda w. (f w * (w - z)) / g w) \longrightarrow f z / g')$  (at z)
  proof (rule lhopital_complex_simple)
    show  $((\lambda w. f w * (w - z)) \text{ has\_field\_derivative } f z)$  (at z)
      using assms by (auto intro!: derivative_eq_intros holomorphic_derivI[OF
f_holo])
    show  $(g \text{ has\_field\_derivative } g')$  (at z) by fact
  qed (insert assms, auto)
  then show  $((\lambda w. (f w / g w) * (w - z)) \longrightarrow f z / g')$  (at z)
    by (simp add: field_split_simps)
  qed
qed

```

5.15 Poles and residues of some well-known functions

```

lemma is_pole_Gamma: is_pole Gamma (-of_nat n)
  unfolding is_pole_def using Gamma_poles .

```

```

lemma Gamma_residue:

```

```

  residue Gamma (-of_nat n) =  $(-1)^n / \text{fact } n$ 
proof (rule residue_simple')
  show open  $(-\mathbb{Z}_{<0} - \{-of\_nat\ n\}) :: \text{complex set}$ 
    by (intro open_Compl closed_subset_Ints) auto
  show Gamma holomorphic_on  $(-\mathbb{Z}_{<0} - \{-of\_nat\ n\}) - \{-of\_nat\ n\}$ 
    by (rule holomorphic_Gamma) auto
  show  $(\lambda w. \text{Gamma } w * (w - (-of\_nat\ n))) \rightarrow (-of\_nat\ n) \rightarrow (-1)^n / \text{fact } n$ 
    using Gamma_residues[of n] by simp
qed auto

```

```

end

```

6 The Residue Theorem, the Argument Principle and Rouché's Theorem

```

theory Residue_Theorem

```

```

  imports Complex_Residues HOL-Library.Landau_Symbols
begin

```

Several theorems that could be moved up, IF there were a previous theory importing both Landau Symbols and Elementary Metric Spaces

```

lemma continuous_bounded_at_infinity_imp_bounded:

```

```

  fixes  $f :: \text{real} \Rightarrow 'a :: \text{real\_normed\_field}$ 
  assumes  $f \in O[at\_bot](\lambda_. 1)$ 

```

```

assumes  $f \in O[at\_top](\lambda_. 1)$ 
assumes continuous_on UNIV f
shows bounded (range f)
proof -
from assms(1) obtain  $c1$  where eventually  $(\lambda x. norm (f x) \leq c1)$  at_bot
by (auto elim!: landau_o.bigE)
then obtain  $x1$  where  $x1: \bigwedge x. x \leq x1 \implies norm (f x) \leq c1$ 
by (auto simp: eventually_at_bot_linorder)
from assms(2) obtain  $c2$  where eventually  $(\lambda x. norm (f x) \leq c2)$  at_top
by (auto elim!: landau_o.bigE)
then obtain  $x2$  where  $x2: \bigwedge x. x \geq x2 \implies norm (f x) \leq c2$ 
by (auto simp: eventually_at_top_linorder)
have compact  $(f \text{ ` } \{x1..x2\})$ 
by (intro compact_continuous_image continuous_on_subset[OF assms(3)])
auto
hence bounded  $(f \text{ ` } \{x1..x2\})$ 
by (rule compact_imp_bounded)
then obtain  $c3$  where  $c3: \bigwedge x. x \in \{x1..x2\} \implies norm (f x) \leq c3$ 
unfolding bounded_iff by fast
have  $norm (f x) \leq \text{Max } \{c1, c2, c3\}$  for  $x$ 
by (cases  $x \leq x1$ ; cases  $x \geq x2$ ) (use  $x1\ x2\ c3$  in auto simp: le_max_iff_disj)
thus ?thesis
unfolding bounded_iff by blast
qed

```

lemma *holomorphic_on_extend*:

```

assumes f holomorphic_on S - {ξ} ξ ∈ interior S f ∈ O[at ξ](λ_. 1)
shows  $(\exists g. g \text{ holomorphic\_on } S \wedge (\forall z \in S - \{\xi\}. g z = f z))$ 
by (subst holomorphic_on_extend_bounded) (insert assms, auto elim!: landau_o.bigE)

```

lemma *removable_singularities*:

```

assumes finite X X ⊆ interior S f holomorphic_on (S - X)
assumes  $\bigwedge z. z \in X \implies f \in O[at z](\lambda_. 1)$ 
shows  $\exists g. g \text{ holomorphic\_on } S \wedge (\forall z \in S - X. g z = f z)$ 
using assms
proof (induction arbitrary: f rule: finite_induct)
case empty
thus ?case by auto
next
case (insert z0 X f)
from insert.prem1 and insert.hyps have  $z0: z0 \in \text{interior } (S - X)$ 
by (auto simp: interior_diff finite_imp_closed)
hence  $\exists g. g \text{ holomorphic\_on } (S - X) \wedge (\forall z \in S - X - \{z0\}. g z = f z)$ 
using insert.prem1 insert.hyps by (intro holomorphic_on_extend) auto
then obtain  $g$  where  $g: g \text{ holomorphic\_on } (S - X) \wedge (\forall z \in S - X - \{z0\}. g z =$ 
f z by blast
have  $\exists h. h \text{ holomorphic\_on } S \wedge (\forall z \in S - X. h z = g z)$ 
proof (rule insert.IH)
fix  $z0'$  assume  $z0' \in X$ 

```

hence *eventually* ($\lambda z. z \in \text{interior } S - (X - \{z0'\}) - \{z0\}$) (*nhds* $z0'$)
using *insert.prem*s *insert.hyps*
by (*intro eventually_nhds_in_open open_Diff finite_imp_closed*) *auto*
hence *ev*: *eventually* ($\lambda z. z \in S - X - \{z0\}$) (*at* $z0'$)
unfolding *eventually_at_filter*
by *eventually_elim* (*insert* $z0'$ *insert.hyps interior_subset*[*of* S], *auto*)
have $g \in \Theta[\text{at } z0'](f)$
by (*intro bigthetaI_cong eventually_mono*[*OF* *ev*]) (*insert* g , *auto*)
also have $f \in O[\text{at } z0'](\lambda_. 1)$
using $z0'$ **by** (*intro insert.prem*s) *auto*
finally show $g \in \dots$
qed (*insert insert.prem*s g , *auto*)
then obtain h **where** h *holomorphic_on* $S \forall z \in S - X. h z = g z$ **by** *blast*
with g **have** h *holomorphic_on* $S \forall z \in S - \text{insert } z0 X. h z = f z$ **by** *auto*
thus *?case* **by** *blast*
qed

lemma *continuous_imp_bigo_1*:
assumes *continuous* (*at* x *within* A) f
shows $f \in O[\text{at } x \text{ within } A](\lambda_. 1)$
proof (*rule bigoI_tendsto*)
from *assms* **show** ($\lambda x. f x / 1 \longrightarrow f x$) (*at* x *within* A)
by (*auto simp: continuous_within*)
qed *auto*

lemma *taylor_bigo_linear*:
assumes f *field_differentiable* *at* $x0$ *within* A
shows ($\lambda x. f x - f x0 \in O[\text{at } x0 \text{ within } A](\lambda x. x - x0)$)
proof –
from *assms* **obtain** f' **where** (f *has_field_derivative* f') (*at* $x0$ *within* A)
by (*auto simp: field_differentiable_def*)
hence ($\lambda x. (f x - f x0) / (x - x0) \longrightarrow f'$) (*at* $x0$ *within* A)
by (*auto simp: has_field_derivative_iff*)
thus *?thesis* **by** (*intro bigoI_tendsto*[**where** $c = f'$]) (*auto simp: eventually_at_filter*)
qed

6.1 Cauchy's residue theorem

lemma *get_integrable_path*:
assumes *open* s *connected* (s -*pts*) *finite* *pts* f *holomorphic_on* (s -*pts*) $a \in s$ -*pts*
 $b \in s$ -*pts*
obtains g **where** *valid_path* g *pathstart* $g = a$ *pathfinish* $g = b$
path_image $g \subseteq s$ -*pts* f *contour_integrable_on* g **using** *assms*
proof (*induct arbitrary:s thesis a rule:finite_induct*[*OF* \langle *finite pts* \rangle])
case 1
obtain g **where** *valid_path* g *path_image* $g \subseteq s$ *pathstart* $g = a$ *pathfinish* $g = b$
using *connected_open_polynomial_connected*[*OF* \langle *open s* \rangle ,*of* a b \rangle \langle *connected*
 $(s - \{\}$ \rangle
valid_path_polynomial_function 1.*prems*(6) 1.*prems*(7) **by** *auto*

```

moreover have  $f$  contour_integrable_on  $g$ 
  using contour_integrable_holomorphic_simple[ $OF$   $\langle open\ s \rangle$   $\langle valid\_path\ g \rangle$ 
 $\langle path\_image\ g \subseteq s \rangle$ ,  $of\ f$ ]
   $\langle f$  holomorphic_on  $s - \{ \}$   $\rangle$ 
  by auto
ultimately show ?case using  $1(1)[of\ g]$  by auto
next
case idt:( $2\ p\ pts$ )
obtain  $e$  where  $e > 0$  and  $e: \forall w \in ball\ a\ e. w \in s \wedge (w \neq a \longrightarrow w \notin insert\ p\ pts)$ 
  using finite_ball_avoid[ $OF$   $\langle open\ s \rangle$   $\langle finite\ (insert\ p\ pts) \rangle$ ,  $of\ a$ ]
   $\langle a \in s - insert\ p\ pts \rangle$ 
  by auto
define  $a'$  where  $a' \equiv a + e/2$ 
have  $a' \in s - \{p\} - pts$  using  $e[rule\_format, of\ a + e/2]$   $\langle e > 0 \rangle$ 
  by (auto simp add: dist_complex_def a'_def)
then obtain  $g'$  where  $g'[simp]: valid\_path\ g'$   $pathstart\ g' = a'$   $pathfinish\ g' = b$ 
   $path\_image\ g' \subseteq s - \{p\} - pts$   $f$  contour_integrable_on  $g'$ 
  using idt.hyps( $3$ )[ $of\ a' s - \{p\}$ ] idt.prems idt.hyps( $1$ )
  by (metis Diff_insert2 open_delete)
define  $g$  where  $g \equiv linepath\ a\ a' +++ g'$ 
have valid_path  $g$  unfolding  $g\_def$  by (auto intro: valid_path_join)
moreover have  $pathstart\ g = a$  and  $pathfinish\ g = b$  unfolding  $g\_def$  by auto
moreover have  $path\_image\ g \subseteq s - insert\ p\ pts$  unfolding  $g\_def$ 
proof (rule subset_path_image_join)
  have closed_segment  $a\ a' \subseteq ball\ a\ e$  using  $\langle e > 0 \rangle$ 
  by (auto dest!: segment_bound1 simp: a'_def dist_complex_def norm_minus_commute)
  then show  $path\_image\ (linepath\ a\ a') \subseteq s - insert\ p\ pts$  using  $e\ idt(9)$ 
  by auto
next
  show  $path\_image\ g' \subseteq s - insert\ p\ pts$  using  $g'(4)$  by blast
qed
moreover have  $f$  contour_integrable_on  $g$ 
proof -
  have closed_segment  $a\ a' \subseteq ball\ a\ e$  using  $\langle e > 0 \rangle$ 
  by (auto dest!: segment_bound1 simp: a'_def dist_complex_def norm_minus_commute)
  then have continuous_on (closed_segment  $a\ a'$ )  $f$ 
  using  $e\ idt.prem$ s( $6$ ) holomorphic_on_imp_continuous_on[ $OF\ idt.prem$ s( $5$ )]
  apply (elim continuous_on_subset)
  by auto
  then have  $f$  contour_integrable_on linepath  $a\ a'$ 
  using contour_integrable_continuous_linepath by auto
  then show ?thesis unfolding  $g\_def$ 
  apply (rule contour_integrable_joinI)
  by (auto simp add:  $\langle e > 0 \rangle$ )
qed
ultimately show ?case using  $idt.prem$ s( $1$ )[ $of\ g$ ] by auto
qed

```

lemma *Cauchy_theorem_aux*:

```

assumes open s connected (s-pts) finite pts pts  $\subseteq$  s f holomorphic_on s-pts
  valid_path g pathfinish g = pathstart g path_image g  $\subseteq$  s-pts
   $\forall z. (z \notin s) \longrightarrow \text{winding\_number } g \ z = 0$ 
   $\forall p \in s. h \ p > 0 \wedge (\forall w \in \text{cball } p (h \ p). w \in s \wedge (w \neq p \longrightarrow w \notin \text{pts}))$ 
shows contour_integral g f =  $(\sum p \in \text{pts}. \text{winding\_number } g \ p * \text{contour\_integral } (\text{circlepath } p (h \ p)) \ f)$ 
using assms
proof (induct arbitrary:s g rule:finite_induct[OF <finite pts>])
  case 1
  then show ?case by (simp add: Cauchy_theorem_global contour_integral_unique)
next
  case (2 p pts)
  note fin[simp] = <finite (insert p pts)>
  and connected = <connected (s - insert p pts)>
  and valid[simp] = <valid_path g>
  and g_loop[simp] = <pathfinish g = pathstart g>
  and holo[simp] = <f holomorphic_on s - insert p pts>
  and path_img = <path_image g  $\subseteq$  s - insert p pts>
  and winding = < $\forall z. z \notin s \longrightarrow \text{winding\_number } g \ z = 0$ >
  and h = < $\forall p a \in s. 0 < h \ p a \wedge (\forall w \in \text{cball } p a (h \ p a). w \in s \wedge (w \neq p a \longrightarrow w \notin \text{insert } p \ \text{pts}))$ >
  have h p > 0 and p  $\in$  s
  and h_p:  $\forall w \in \text{cball } p (h \ p). w \in s \wedge (w \neq p \longrightarrow w \notin \text{insert } p \ \text{pts})$ 
  using h <insert p pts  $\subseteq$  s> by auto
  obtain pg where pg[simp]: valid_path pg pathstart pg = pathstart g pathfinish
pg=p+h p
  path_image pg  $\subseteq$  s-insert p pts f contour_integrable_on pg
  proof -
  have p + h p  $\in$  cball p (h p) using h[rule_format,of p]
  by (simp add: <p  $\in$  s> dist_norm)
  then have p + h p  $\in$  s - insert p pts using h[rule_format,of p] <insert p
pts  $\subseteq$  s>
  by fastforce
  moreover have pathstart g  $\in$  s - insert p pts using path_img by auto
  ultimately show ?thesis
  using get_integrable_path[OF <open s> connected fin holo,of pathstart g p+h
p] that
  by blast
  qed
obtain n::int where n=winding_number g p
  using integer_winding_number[OF _ g_loop,of p] valid_path_img
  by (metis DiffD2 Ints_cases insertI1 subset_eq valid_path_imp_path)
define p_circ where p_circ  $\equiv$  circlepath p (h p)
define p_circ_pt where p_circ_pt  $\equiv$  linepath (p+h p) (p+h p)
define n_circ where n_circ  $\equiv$   $\lambda n. ((+++)$  p_circ  $\widehat{\sim}$  n) p_circ_pt
define cp where cp  $\equiv$  if n  $\geq$  0 then reversepath (n_circ (nat n)) else n_circ (nat
(- n))
  have n_circ:valid_path (n_circ k)
  winding_number (n_circ k) p = k

```

```

    pathstart (n_circ k) = p + h p pathfinish (n_circ k) = p + h p
    path_image (n_circ k) = (if k=0 then {p + h p} else sphere p (h p))
    p ∉ path_image (n_circ k)
    ∧ p'. p' ∉ s - pts ⇒ winding_number (n_circ k) p'=0 ∧ p' ∉ path_image
(n_circ k)
  f contour_integrable_on (n_circ k)
  contour_integral (n_circ k) f = k * contour_integral p_circ f
  for k
proof (induct k)
  case 0
  show valid_path (n_circ 0)
  and path_image (n_circ 0) = (if 0=0 then {p + h p} else sphere p (h p))
  and winding_number (n_circ 0) p = of_nat 0
  and pathstart (n_circ 0) = p + h p
  and pathfinish (n_circ 0) = p + h p
  and p ∉ path_image (n_circ 0)
  unfolding n_circ_def p_circ_pt_def using ⟨h p > 0⟩
  by (auto simp add: dist_norm)
  show winding_number (n_circ 0) p'=0 ∧ p' ∉ path_image (n_circ 0) when
p' ∉ s - pts for p'
  unfolding n_circ_def p_circ_pt_def
  apply (auto intro!: winding_number_trivial)
  by (metis Diff_iff pathfinish_in_path_image pg(3) pg(4) subsetCE sub-
set_insertI that)+
  show f contour_integrable_on (n_circ 0)
  unfolding n_circ_def p_circ_pt_def
  by (auto intro!: contour_integrable_continuous_linepath simp add: continuous_on_sing)
  show contour_integral (n_circ 0) f = of_nat 0 * contour_integral p_circ f
  unfolding n_circ_def p_circ_pt_def by auto
next
  case (Suc k)
  have n_Suc:n_circ (Suc k) = p_circ +++ n_circ k unfolding n_circ_def
by auto
  have pcirc:p ∉ path_image p_circ valid_path p_circ pathfinish p_circ =
pathstart (n_circ k)
  using Suc(3) unfolding p_circ_def using ⟨h p > 0⟩ by (auto simp add:
p_circ_def)
  have pcirc_image:path_image p_circ ⊆ s - insert p pts
  proof -
  have path_image p_circ ⊆ cball p (h p) using ⟨0 < h p⟩ p_circ_def by
auto
  then show ?thesis using h_p pcirc(1) by auto
  qed
  have pcirc_integrable:f contour_integrable_on p_circ
  by (auto simp add:p_circ_def intro!: pcirc_image[unfolded p_circ_def]
contour_integrable_continuous_circlepath holomorphic_on_imp_continuous_on
holomorphic_on_subset[OF holo])
  show valid_path (n_circ (Suc k))
  using valid_path_join[OF pcirc(2) Suc(1) pcirc(3)] unfolding n_circ_def

```



```

by auto
  show path_image (n_circ (Suc k))
    = (if Suc k = 0 then {p + complex_of_real (h p)} else sphere p (h p))
  proof -
    have path_image p_circ = sphere p (h p)
      unfolding p_circ_def using ‹0 < h p› by auto
    then show ?thesis unfolding n_Suc using Suc.hyps(5) ‹h p > 0›
      by (auto simp add: path_image_join[OF pcirc(3)] dist_norm)
    qed
  then show p ∉ path_image (n_circ (Suc k)) using ‹h p > 0› by auto
  show winding_number (n_circ (Suc k)) p = of_nat (Suc k)
  proof -
    have winding_number p_circ p = 1
      by (simp add: ‹h p > 0› p_circ_def winding_number_circlepath_centre)
    moreover have p ∉ path_image (n_circ k) using Suc(5) ‹h p > 0› by
auto
    then have winding_number (p_circ +++ n_circ k) p
      = winding_number p_circ p + winding_number (n_circ k) p
      using valid_path_imp_path Suc.hyps(1) Suc.hyps(2) pcirc
      apply (intro winding_number_join)
      by auto
    ultimately show ?thesis using Suc(2) unfolding n_circ_def
      by auto
    qed
  show pathstart (n_circ (Suc k)) = p + h p
    by (simp add: n_circ_def p_circ_def)
  show pathfinish (n_circ (Suc k)) = p + h p
    using Suc(4) unfolding n_circ_def by auto
  show winding_number (n_circ (Suc k)) p' = 0 ∧ p' ∉ path_image (n_circ
(Suc k)) when p' ∉ s-pts for p'
  proof -
    have p' ∉ path_image p_circ using ‹p ∈ s› h p_circ_def that using
pcirc_image by blast
    moreover have p' ∉ path_image (n_circ k)
      using Suc.hyps(7) that by blast
    moreover have winding_number p_circ p' = 0
    proof -
      have path_image p_circ ⊆ cball p (h p)
        using h unfolding p_circ_def using ‹p ∈ s› by fastforce
      moreover have p' ∉ cball p (h p) using ‹p ∈ s› h that 2.hyps(2) by
fastforce
    ultimately show ?thesis unfolding p_circ_def
      apply (intro winding_number_zero_outside)
      by auto
    qed
  qed
  ultimately show ?thesis
    unfolding n_Suc
    apply (subst winding_number_join)
  by (auto simp: valid_path_imp_path pcirc Suc that not_in_path_image_join

```

```

Suc.hyps(7)[OF that]
  qed
  show f contour_integrable_on (n_circ (Suc k))
    unfolding n_Suc
    by (rule contour_integrable_joinI[OF pcirc_integrable Suc(8) pcirc(2)
Suc(1)])
  show contour_integral (n_circ (Suc k)) f = (Suc k) * contour_integral
p_circ f
    unfolding n_Suc
    by (auto simp add: contour_integral_join[OF pcirc_integrable Suc(8) pcirc(2)
Suc(1)])
  Suc(9) algebra_simps
  qed
have cp[simp]: pathstart cp = p + h p pathfinish cp = p + h p
  valid_path cp path_image cp  $\subseteq$  s - insert p pts
  winding_number cp p = - n
   $\wedge p'. p' \notin s - pts \implies$  winding_number cp p'=0  $\wedge p' \notin$  path_image cp
  f contour_integrable_on cp
  contour_integral cp f = - n * contour_integral p_circ f
proof -
  show pathstart cp = p + h p and pathfinish cp = p + h p and valid_path cp
    using n_circ unfolding cp_def by auto
next
  have sphere p (h p)  $\subseteq$  s - insert p pts
    using h[rule_format, of p]  $\langle$ insert p pts  $\subseteq$  s $\rangle$  by force
  moreover have p + complex_of_real (h p)  $\in$  s - insert p pts
    using pg(3) pg(4) by (metis pathfinish_in_path_image subsetCE)
  ultimately show path_image cp  $\subseteq$  s - insert p pts unfolding cp_def
    using n_circ(5) by auto
next
  show winding_number cp p = - n
    unfolding cp_def using winding_number_reversepath n_circ  $\langle$ h p $\rangle$ 
    by (auto simp: valid_path_imp_path)
next
  show winding_number cp p'=0  $\wedge p' \notin$  path_image cp when p'  $\notin$  s - pts for
p'
    unfolding cp_def
    apply (auto)
    apply (subst winding_number_reversepath)
    by (auto simp add: valid_path_imp_path n_circ(7)[OF that] n_circ(1))
next
  show f contour_integrable_on cp unfolding cp_def
    using contour_integrable_reversepath_eq n_circ(1,8) by auto
next
  show contour_integral cp f = - n * contour_integral p_circ f
    unfolding cp_def using contour_integral_reversepath[OF n_circ(1)]
n_circ(9)
    by auto
  qed

```

```

define  $g'$  where  $g' \equiv g +++ pg +++ cp +++ (reversepath\ pg)$ 
have  $contour\_integral\ g' f = (\sum_{p \in pts} winding\_number\ g' p * contour\_integral$ 
( $circlepath\ p\ (h\ p)$ )  $f)$ 
proof ( $rule\ 2.hyps(3)[of\ s - \{p\}\ g', OF\_ \_ \langle finite\ pts \rangle ]$ )
  show  $connected\ (s - \{p\} - pts)$  using  $connected\ by\ (metis\ Diff\_insert2)$ 
  show  $open\ (s - \{p\})$  using  $\langle open\ s \rangle$  by  $auto$ 
  show  $pts \subseteq s - \{p\}$  using  $\langle insert\ p\ pts \subseteq s \rangle \langle p \notin pts \rangle$  by  $blast$ 
  show  $f\ holomorphic\_on\ s - \{p\} - pts$  using  $holo\ \langle p \notin pts \rangle$  by ( $metis$ 
 $Diff\_insert2$ )
  show  $valid\_path\ g'$ 
    unfolding  $g'\_def\ cp\_def$  using  $n\_circ\ valid\ pg\ g\_loop$ 
    by ( $auto\ intro!: valid\_path\_join$ )
  show  $pathfinish\ g' = pathstart\ g'$ 
    unfolding  $g'\_def\ cp\_def$  using  $pg(2)$  by  $simp$ 
  show  $path\_image\ g' \subseteq s - \{p\} - pts$ 
    proof -
      define  $s'$  where  $s' \equiv s - \{p\} - pts$ 
      have  $s':s' = s - insert\ p\ pts$  unfolding  $s'\_def$  by  $auto$ 
      then show  $?thesis$  using  $path\_img\ pg(4)\ cp(4)$ 
        unfolding  $g'\_def$ 
        apply ( $fold\ s'\_def\ s'$ )
        apply ( $intro\ subset\_path\_image\_join$ )
        by  $auto$ 
      qed
    note  $path\_join\_imp[simp]$ 
    show  $\forall z. z \notin s - \{p\} \longrightarrow winding\_number\ g' z = 0$ 
      proof  $clarify$ 
        fix  $z$  assume  $z:z \notin s - \{p\}$ 
        have  $winding\_number\ (g\ +++\ pg\ +++\ cp\ +++\ reversepath\ pg)\ z =$ 
 $winding\_number\ g\ z$ 
           $+ winding\_number\ (pg\ +++\ cp\ +++\ (reversepath\ pg))\ z$ 
        proof ( $rule\ winding\_number\_join$ )
          show  $path\ g$  using  $\langle valid\_path\ g \rangle$  by ( $simp\ add: valid\_path\_imp\_path$ )
          show  $z \notin path\_image\ g$  using  $z\ path\_img$  by  $auto$ 
          show  $path\ (pg\ +++\ cp\ +++\ reversepath\ pg)$  using  $pg(3)\ cp$ 
            by ( $simp\ add: valid\_path\_imp\_path$ )
          next
            have  $path\_image\ (pg\ +++\ cp\ +++\ reversepath\ pg) \subseteq s - insert\ p\ pts$ 
              using  $pg(4)\ cp(4)$  by ( $auto\ simp: subset\_path\_image\_join$ )
            then show  $z \notin path\_image\ (pg\ +++\ cp\ +++\ reversepath\ pg)$  using
 $z\ by\ auto$ 
          next
            show  $pathfinish\ g = pathstart\ (pg\ +++\ cp\ +++\ reversepath\ pg)$  using
 $g\_loop$  by  $auto$ 
          qed
        also have  $... = winding\_number\ g\ z + (winding\_number\ pg\ z$ 
           $+ winding\_number\ (cp\ +++\ (reversepath\ pg))\ z)$ 
        proof ( $subst\ add\_left\_cancel, rule\ winding\_number\_join$ )
          show  $path\ pg$  and  $path\ (cp\ +++\ reversepath\ pg)$ 

```

```

    and pathfinish pg = pathstart (cp +++ reversepath pg)
      by (auto simp add: valid_path_imp_path)
    show z  $\notin$  path_image pg using pg(4) z by blast
    show z  $\notin$  path_image (cp +++ reversepath pg) using z
    by (metis Diff_iff  $\langle z \notin \text{path\_image } pg \rangle$  contra_subsetD cp(4) insertI1
        not_in_path_image_join path_image_reversepath singletonD)
  qed
  also have ... = winding_number g z + (winding_number pg z
    + (winding_number cp z + winding_number (reversepath pg) z))
    apply (auto intro!:winding_number_join simp: valid_path_imp_path)
    apply (metis Diff_iff contra_subsetD cp(4) insertI1 singletonD z)
    by (metis Diff_insert2 Diff_subset contra_subsetD pg(4) z)
  also have ... = winding_number g z + winding_number cp z
    apply (subst winding_number_reversepath)
    apply (auto simp: valid_path_imp_path)
    by (metis Diff_iff contra_subsetD insertI1 pg(4) singletonD z)
    finally have winding_number g' z = winding_number g z + wind-
ing_number cp z
    unfolding g'_def .
    moreover have winding_number g z + winding_number cp z = 0
      using winding z  $\langle n = \text{winding\_number } g \ p \rangle$  by auto
    ultimately show winding_number g' z = 0 unfolding g'_def by auto
  qed
  show  $\forall pa \in s - \{p\}. 0 < h \ pa \wedge (\forall w \in \text{cball } pa \ (h \ pa). w \in s - \{p\} \wedge (w \neq
pa \longrightarrow w \notin \text{pts}))$ 
    using h by fastforce
  qed
  moreover have contour_integral g' f = contour_integral g f
    - winding_number g p * contour_integral p_circ f
  proof -
    have contour_integral g' f = contour_integral g f
      + contour_integral (pg +++ cp +++ reversepath pg) f
    unfolding g'_def
    apply (subst contour_integral_join)
    by (auto simp add:open_Diff[OF  $\langle \text{open } s \rangle$ ,OF finite_imp_closed[OF fin]]
        intro!: contour_integrable_holomorphic_simple[OF holo __ path_img]
        contour_integrable_reversepath)
    also have ... = contour_integral g f + contour_integral pg f
      + contour_integral (cp +++ reversepath pg) f
    apply (subst contour_integral_join)
    by (auto simp add:contour_integrable_reversepath)
    also have ... = contour_integral g f + contour_integral pg f
      + contour_integral cp f + contour_integral (reversepath pg) f
    apply (subst contour_integral_join)
    by (auto simp add:contour_integrable_reversepath)
    also have ... = contour_integral g f + contour_integral cp f
    using contour_integral_reversepath
    by (auto simp add:contour_integrable_reversepath)
  also have ... = contour_integral g f - winding_number g p * contour_integral

```

```

p_circ f
  using ⟨n=winding_number g p⟩ by auto
  finally show ?thesis .
qed
moreover have winding_number g' p' = winding_number g p' when p'∈pts for
p'
proof -
  have [simp]: p' ∉ path_image g p' ∉ path_image pg p' ∉ path_image cp
  using 2.premis(8) that
  apply blast
  apply (metis Diff_iff Diff_insert2 contra_subsetD pg(4) that)
  by (meson DiffD2 cp(4) rev_subsetD subset_insertI that)
  have winding_number g' p' = winding_number g p'
  + winding_number (pg +++ cp +++ reversepath pg) p' unfolding g'_def
  apply (subst winding_number_join)
  apply (simp_all add: valid_path_imp_path)
  apply (intro not_in_path_image_join)
  by auto
  also have ... = winding_number g p' + winding_number pg p'
  + winding_number (cp +++ reversepath pg) p'
  apply (subst winding_number_join)
  apply (simp_all add: valid_path_imp_path)
  apply (intro not_in_path_image_join)
  by auto
  also have ... = winding_number g p' + winding_number pg p' + wind-
ing_number cp p'
  + winding_number (reversepath pg) p'
  apply (subst winding_number_join)
  by (simp_all add: valid_path_imp_path)
  also have ... = winding_number g p' + winding_number cp p'
  apply (subst winding_number_reversepath)
  by (simp_all add: valid_path_imp_path)
  also have ... = winding_number g p' using that by auto
  finally show ?thesis .
qed
ultimately show ?case unfolding p_circ_def
  apply (subst (asm) sum.cong[OF refl,
of pts _ λp. winding_number g p * contour_integral (circlepath p (h p)) f])
  by (auto simp add: sum.insert[OF ⟨finite pts⟩ ⟨p∉pts⟩] algebra_simps)
qed

```

lemma *Cauchy_theorem_singularities:*

```

assumes open s connected s finite pts and
  holo:f holomorphic_on s-pts and
  valid_path g and
  loop:pathfinish g = pathstart g and
  path_image g ⊆ s-pts and
  homo:∀ z. (z ∉ s) ⟶ winding_number g z = 0 and
  avoid:∀ p∈s. h p>0 ∧ (∀ w∈cball p (h p). w∈s ∧ (w≠p ⟶ w ∉ pts))

```

```

shows contour_integral g f = ( $\sum p \in pts. \text{winding\_number } g \ p * \text{contour\_integral}$ 
(circlepath p (h p)) f)
  (is ?L=?R)
proof -
  define circ where circ  $\equiv \lambda p. \text{winding\_number } g \ p * \text{contour\_integral}$  (circlepath
p (h p)) f
  define pts1 where pts1  $\equiv pts \cap s$ 
  define pts2 where pts2  $\equiv pts - pts1$ 
  have pts=pts1  $\cup$  pts2 pts1  $\cap$  pts2 = {} pts2  $\cap$  s={} pts1  $\subseteq$  s
    unfolding pts1_def pts2_def by auto
  have contour_integral g f = ( $\sum p \in pts1. \text{circ } p$ ) unfolding circ_def
    proof (rule Cauchy_theorem_aux[OF  $\langle \text{open } s \rangle$   $\_ \_ \langle pts1 \subseteq s \rangle$   $\_ \_ \langle \text{valid\_path}$ 
g  $\rangle$  loop\_homo])
      have finite pts1 unfolding pts1_def using  $\langle \text{finite } pts \rangle$  by auto
      then show connected (s - pts1)
        using  $\langle \text{open } s \rangle$   $\langle \text{connected } s \rangle$  connected\_open\_delete\_finite[of s] by auto
      next
        show finite pts1 using  $\langle pts = pts1 \cup pts2 \rangle$  assms(3) by auto
        show f holomorphic_on s - pts1 by (metis Diff_Int2 Int_absorb holo
pts1_def)
        show path_image g  $\subseteq$  s - pts1 using assms(7) pts1_def by auto
        show  $\forall p \in s. 0 < h \ p \wedge (\forall w \in \text{cball } p \ (h \ p). w \in s \wedge (w \neq p \longrightarrow w \notin pts1))$ 
          by (simp add: avoid pts1_def)
      qed
    moreover have sum circ pts2=0
      proof -
        have winding_number g p=0 when p  $\in$  pts2 for p
          using  $\langle pts2 \cap s = \{ \} \rangle$  that homo[rule_format,of p] by auto
        thus ?thesis unfolding circ_def
          apply (intro sum.neutral)
          by auto
        qed
      moreover have ?R=sum circ pts1 + sum circ pts2
        unfolding circ_def
        using sum.union_disjoint[OF  $\_ \_ \langle pts1 \cap pts2 = \{ \} \rangle$ ]  $\langle \text{finite } pts \rangle$   $\langle pts = pts1$ 
 $\cup$  pts2  $\rangle$ 
          by blast
        ultimately show ?thesis
          apply (fold circ_def)
          by auto
      qed

```

theorem *Residue_theorem*:

```

fixes s pts::complex set and f::complex  $\Rightarrow$  complex
and g::real  $\Rightarrow$  complex
assumes open s connected s finite pts and
  holo:f holomorphic_on s-pts and
  valid_path g and
  loop:pathfinish g = pathstart g and

```

```

    path_image g  $\subseteq$  s-pts and
    homo: $\forall z. (z \notin s) \longrightarrow \text{winding\_number } g z = 0$ 
  shows  $\text{contour\_integral } g f = 2 * \pi * i * (\sum_{p \in \text{pts.}} \text{winding\_number } g p * \text{residue } f p)$ 
  proof -
    define c where  $c \equiv 2 * \pi * i$ 
    obtain h where avoid: $\forall p \in s. h p > 0 \wedge (\forall w \in \text{cball } p (h p). w \in s \wedge (w \neq p \longrightarrow w \notin \text{pts}))$ 
    using finite_cball_avoid[OF  $\langle \text{open } s \rangle \langle \text{finite pts} \rangle$ ] by metis
    have  $\text{contour\_integral } g f$ 
      =  $(\sum_{p \in \text{pts.}} \text{winding\_number } g p * \text{contour\_integral } (\text{circlepath } p (h p)) f)$ 
    using Cauchy_theorem_singularities[OF  $\text{assms avoid}$ ] .
    also have ... =  $(\sum_{p \in \text{pts.}} c * \text{winding\_number } g p * \text{residue } f p)$ 
    proof (intro sum.cong)
      show  $\text{pts} = \text{pts}$  by simp
    next
      fix x assume  $x \in \text{pts}$ 
      show  $\text{winding\_number } g x * \text{contour\_integral } (\text{circlepath } x (h x)) f$ 
        =  $c * \text{winding\_number } g x * \text{residue } f x$ 
      proof (cases  $x \in s$ )
        case False
          then have  $\text{winding\_number } g x = 0$  using homo by auto
          thus ?thesis by auto
        next
          case True
            have  $\text{contour\_integral } (\text{circlepath } x (h x)) f = c * \text{residue } f x$ 
              using  $\langle x \in \text{pts} \rangle \langle \text{finite pts} \rangle \text{avoid}[\text{rule\_format}, \text{OF True}]$ 
            apply (intro base_residue[of  $s - (\text{pts} - \{x\})$ ], THEN  $\text{contour\_integral\_unique, folded } c\_def$ )
              by (auto intro: holomorphic_on_subset[OF holo] open_Diff[OF  $\langle \text{open } s \rangle$ ]
                finite_imp_closed])
            then show ?thesis by auto
          qed
      qed
    also have ... =  $c * (\sum_{p \in \text{pts.}} \text{winding\_number } g p * \text{residue } f p)$ 
    by (simp add: sum_distrib_left algebra_simps)
    finally show ?thesis unfolding c_def .
  qed

```

6.2 The argument principle

theorem *argument_principle:*

fixes $f::\text{complex} \Rightarrow \text{complex}$ **and** $\text{poles } s::\text{complex set}$

defines $\text{pz} \equiv \{w \in s. f w = 0 \vee w \in \text{poles}\}$ — pz is the set of poles and zeros

assumes $\text{open } s$ **connected } s **and****

$f_holo:f$ holomorphic_on $s - \text{poles}$ **and**

$h_holo:h$ holomorphic_on s **and**

$\text{valid_path } g$ **and**

$\text{loop: pathfinish } g = \text{pathstart } g$ **and**

```

    path_img:path_image g  $\subseteq$  s - pz and
    homo: $\forall z. (z \notin s) \longrightarrow$  winding_number g z = 0 and
    finite:finite pz and
    poles: $\forall p \in s \cap$ poles. is_pole f p
  shows contour_integral g ( $\lambda x. \text{deriv } f x * h x / f x$ ) = 2 * pi * i *
    ( $\sum_{p \in pz. \text{winding\_number } g p * h p * \text{zorder } f p$ )
    (is ?L=?R)
  proof -
    define c where c  $\equiv$  2 * complex_of_real pi * i
    define ff where ff  $\equiv$  ( $\lambda x. \text{deriv } f x * h x / f x$ )
    define cont where cont  $\equiv$   $\lambda ff p e. (ff \text{ has\_contour\_integral } c * \text{zorder } f p * h p$ 
) (circlepath p e)
    define avoid where avoid  $\equiv$   $\lambda p e. \forall w \in \text{cball } p e. w \in s \wedge (w \neq p \longrightarrow w \notin pz)$ 

    have  $\exists e > 0. \text{avoid } p e \wedge (p \in pz \longrightarrow \text{cont } ff p e)$  when p  $\in$  s for p
    proof -
      obtain e1 where e1 > 0 and e1_avoid:avoid p e1
      using finite_cball_avoid[OF <open s> finite] <p  $\in$  s> unfolding avoid_def by
      auto
      have  $\exists e2 > 0. \text{cball } p e2 \subseteq \text{ball } p e1 \wedge \text{cont } ff p e2$  when p  $\in$  pz
      proof -
        define po where po  $\equiv$  zorder f p
        define pp where pp  $\equiv$  zor_poly f p
        define f' where f'  $\equiv$   $\lambda w. pp w * (w - p) \text{powi } po$ 
        define ff' where ff'  $\equiv$  ( $\lambda x. \text{deriv } f' x * h x / f' x$ )
        obtain r where pp p  $\neq$  0 r > 0 and
          r < e1 and
          pp_holo:pp holomorphic_on cball p r and
          pp_po:( $\forall w \in \text{cball } p r - \{p\}. f w = pp w * (w - p) \text{powi } po \wedge pp w \neq 0$ )
      proof -
        have isolated_singularity_at f p
      proof -
        have f holomorphic_on ball p e1 - {p}
          apply (intro holomorphic_on_subset[OF f_holo])
          using e1_avoid <p  $\in$  pz> unfolding avoid_def pz_def by force
        then show ?thesis unfolding isolated_singularity_at_def
          using <e1 > 0> analytic_on_open open_delete by blast
      qed
    moreover have not_essential f p
    proof (cases is_pole f p)
      case True
      then show ?thesis unfolding not_essential_def by auto
    next
      case False
      then have p  $\in$  s - poles using <p  $\in$  s> poles unfolding pz_def by auto
      moreover have open (s - poles)
      proof -
        have closed (s  $\cap$  poles)
          using finite by (simp add: pz_def finite_imp_closed rev_finite_subset

```



```

subset_eq)
  then show ?thesis
    by (metis Diff_Compl Diff_Diff_Int Diff_eq ‹open s› open_Diff)
  qed
  ultimately have isCont f p
    using holomorphic_on_imp_continuous_on[OF f_holo] continu-
ous_on_eq_continuous_at
    by auto
  then show ?thesis unfolding isCont_def not_essential_def by auto
  qed
  moreover have  $\exists_F w$  in at p.  $f w \neq 0$ 
  proof (rule ccontr)
    assume  $\neg (\exists_F w$  in at p.  $f w \neq 0$ )
    then have  $\forall_F w$  in at p.  $f w = 0$  unfolding frequently_def by auto
    then obtain r1 where  $r1 > 0$  and  $r1: \forall w \in \text{ball } p \ r1 - \{p\}. f w = 0$ 
      unfolding eventually_at by (auto simp add: dist_commute)
    obtain r2 where  $r2 > 0$  and  $r2: \text{ball } p \ r2 \subseteq s$ 
      using ‹p ∈ s› ‹open s› openE by blast
    define rr where  $rr = \min \ r1 \ r2$ 

    from r1 r2
    have  $\text{ball } p \ rr - \{p\} \subseteq \{w \in s \cap \text{ball } p \ rr - \{p\}. f w = 0\}$ 
      unfolding rr_def by auto
    moreover have infinite (ball p rr - {p})
      using ‹r1 > 0› ‹r2 > 0› finite_imp_not_open
      unfolding rr_def by fastforce
    ultimately have infinite  $\{w \in s \cap \text{ball } p \ rr - \{p\}. f w = 0\}$  using infi-
nite_super by blast
    then have infinite pz
      unfolding pz_def by (smt (verit) infinite_super Collect_mono_iff DiffE
Int_iff)
    then show False using ‹finite pz› by auto
  qed
  ultimately obtain r where  $pp \ p \neq 0$  and  $r: r > 0$   $pp$  holomorphic_on cball
p r
    ( $\forall w \in \text{cball } p \ r - \{p\}. f w = pp \ w * (w - p) \ \text{powi } po \wedge pp \ w \neq 0$ )
    using zorder_exist[of f p, folded po_def pp_def] by auto
  define r1 where  $r1 = \min \ r \ e1 / 2$ 
  have  $r1 < e1$  unfolding r1_def using ‹e1 > 0› ‹r > 0› by auto
  moreover have  $r1 > 0$   $pp$  holomorphic_on cball p r1
    ( $\forall w \in \text{cball } p \ r1 - \{p\}. f w = pp \ w * (w - p) \ \text{powi } po \wedge pp \ w \neq 0$ )
    unfolding r1_def using ‹e1 > 0› r by auto
  ultimately show ?thesis using that ‹pp p ≠ 0› by auto
  qed

define e2 where  $e2 \equiv r / 2$ 
have  $e2 > 0$  using ‹r > 0› unfolding e2_def by auto
define anal where  $anal \equiv \lambda w. \text{deriv } pp \ w * h \ w / pp \ w$ 
define prin where  $prin \equiv \lambda w. po * h \ w / (w - p)$ 

```

```

have (( $\lambda w.$  prin  $w + anal\ w$ ) has_contour_integral  $c * po * h\ p$ ) (circlepath
 $p\ e2$ )
proof (rule has_contour_integral_add[of _ _ _ 0,simplified])
  have ball  $p\ r \subseteq s$ 
    using  $\langle r < e1 \rangle$  avoid_def ball_subset_cball  $e1\_avoid$  by (simp add:
subset_eq)
  then have cball  $p\ e2 \subseteq s$ 
    using  $\langle r > 0 \rangle$  unfolding  $e2\_def$  by auto
  then have ( $\lambda w.$   $po * h\ w$ ) holomorphic_on cball  $p\ e2$ 
    using  $h\_holo$  by (auto intro!: holomorphic_intros)
  then show (prin has_contour_integral  $c * po * h\ p$ ) (circlepath  $p\ e2$ )
    using Cauchy_integral_circlepath_simple[folded  $c\_def$ , of  $\lambda w.$   $po * h\ w$ ]
 $\langle e2 > 0 \rangle$ 
    unfolding prin_def by (auto simp add: mult.assoc)
  have anal holomorphic_on ball  $p\ r$  unfolding anal_def
    using pp_holo  $h\_holo\ pp\_po\ \langle ball\ p\ r \subseteq s \rangle\ \langle pp \neq 0 \rangle$ 
    by (auto intro!: holomorphic_intros)
  then show (anal has_contour_integral 0) (circlepath  $p\ e2$ )
    using  $e2\_def\ \langle r > 0 \rangle$ 
    by (auto elim!: Cauchy_theorem_disc_simple)
qed
then have cont  $ff'\ p\ e2$  unfolding cont_def  $po\_def$ 
proof (elim has_contour_integral_eq)
  fix  $w$  assume  $w \in path\_image$  (circlepath  $p\ e2$ )
  then have  $w \in ball\ p\ r$  and  $w \neq p$  unfolding  $e2\_def$  using  $\langle r > 0 \rangle$  by auto
  define  $wp$  where  $wp \equiv w - p$ 
  have  $wp \neq 0$  and  $pp\ w \neq 0$ 
    unfolding  $wp\_def$  using  $\langle w \neq p \rangle\ \langle w \in ball\ p\ r \rangle\ pp\_po$  by auto
  moreover have  $der\_f': deriv\ f'\ w = po * pp\ w * (w - p)\ powi\ (po - 1) +$ 
 $deriv\ pp\ w * (w - p)\ powi\ po$ 
  proof (rule DERIV_imp_deriv)
    have ( $pp$  has_field_derivative (deriv  $pp\ w$ )) (at  $w$ )
      using DERIV_deriv_iff_has_field_derivative  $pp\_holo\ \langle w \neq p \rangle$ 
    by (meson open_ball  $\langle w \in ball\ p\ r \rangle\ ball\_subset\_cball\ holomorphic\_derivI$ 
holomorphic_on_subset)
    then show ( $f'$  has_field_derivative of_int  $po * pp\ w * (w - p)\ powi\ (po$ 
 $- 1)$ 
      +  $deriv\ pp\ w * (w - p)\ powi\ po$ ) (at  $w$ )
      unfolding  $f'\_def$  using  $\langle w \neq p \rangle$ 
    by (auto intro!: derivative_eq_intros DERIV_cong[OF has_field_derivative_powr_of_int])
qed
ultimately show prin  $w + anal\ w = ff'\ w$ 
  unfolding  $ff'\_def\ prin\_def\ anal\_def$ 
  apply simp
  apply (unfold  $f'\_def$ )
  apply (fold  $wp\_def$ )
  apply (auto simp add: field_simps)
  by (metis (no_types, lifting) mult.commute power_int_minus_mult)
qed

```

```

then have cont ff p e2 unfolding cont_def
proof (elim has_contour_integral_eq)
  fix w assume w ∈ path_image (circlepath p e2)
  then have w ∈ ball p r and w ≠ p unfolding e2_def using ‹r > 0› by auto
  have deriv f' w = deriv f w
  proof (rule complex_derivative_transform_within_open[where s = ball p r
- {p}])
    show f' holomorphic_on ball p r - {p} unfolding f'_def using pp_holo
      by (auto intro!: holomorphic_intros)
    next
    have ball p e1 - {p} ⊆ s - poles
      using ball_subset_cball e1_avoid[unfolded avoid_def] unfolding pz_def
      by auto
    then have ball p r - {p} ⊆ s - poles
      apply (elim dual_order.trans)
      using ‹r < e1› by auto
    then show f holomorphic_on ball p r - {p} using f_holo
      by auto
    next
    show open (ball p r - {p}) by auto
    show w ∈ ball p r - {p} using ‹w ∈ ball p r› ‹w ≠ p› by auto
    next
    fix x assume x ∈ ball p r - {p}
    then show f' x = f x
      using pp_po unfolding f'_def by auto
    qed
    moreover have f' w = f w
      using ‹w ∈ ball p r› ball_subset_cball subset_iff pp_po ‹w ≠ p›
      unfolding f'_def by auto
    ultimately show ff' w = ff w
      unfolding ff'_def ff_def by simp
    qed
    moreover have cball p e2 ⊆ ball p e1
      using ‹0 < r› ‹r < e1› e2_def by auto
    ultimately show ?thesis using ‹e2 > 0› by auto
    qed
  then obtain e2 where e2: p ∈ pz → e2 > 0 ∧ cball p e2 ⊆ ball p e1 ∧ cont ff
  p e2
  by auto
  define e4 where e4 ≡ if p ∈ pz then e2 else e1
  have e4 > 0 using e2 ‹e1 > 0› unfolding e4_def by auto
  moreover have avoid p e4 using e2 ‹e1 > 0› e1_avoid unfolding e4_def
  avoid_def by auto
  moreover have p ∈ pz → cont ff p e4
  by (auto simp add: e2 e4_def)
  ultimately show ?thesis by auto
  qed
  then obtain get_e where get_e: ∀ p ∈ s. get_e p > 0 ∧ avoid p (get_e p)
  ∧ (p ∈ pz → cont ff p (get_e p))

```

```

  by metis
  define ci where ci  $\equiv$   $\lambda p.$  contour_integral (circlepath p (get_e p)) ff
  define w where w  $\equiv$   $\lambda p.$  winding_number g p
  have contour_integral g ff =  $(\sum p \in pz. w p * ci p)$  unfolding ci_def w_def
  proof (rule Cauchy_theorem_singularities[OF  $\langle$ open s $\rangle$   $\langle$ connected s $\rangle$  finite _
 $\langle$ valid_path g $\rangle$  loop
    path_img homo])
    have open (s - pz) using open_Diff[OF _ finite_imp_closed[OF finite]]  $\langle$ open
s $\rangle$  by auto
    then show ff holomorphic_on s - pz unfolding ff_def using f_holo h_holo
    by (auto intro!: holomorphic_intros simp add:pz_def)
  next
    show  $\forall p \in s. 0 < \text{get\_e } p \wedge (\forall w \in \text{cball } p (\text{get\_e } p). w \in s \wedge (w \neq p \longrightarrow w \notin
pz))$ 
    using get_e using avoid_def by blast
  qed
  also have ... =  $(\sum p \in pz. c * w p * h p * zorder f p)$ 
  proof (rule sum.cong[of pz pz,simplified])
    fix p assume p  $\in$  pz
    show w p * ci p = c * w p * h p * (zorder f p)
    proof (cases p  $\in$  s)
      assume p  $\in$  s
      have ci p = c * h p * (zorder f p) unfolding ci_def
      apply (rule contour_integral_unique)
      using get_e  $\langle$ p  $\in$  s $\rangle$   $\langle$ p  $\in$  pz $\rangle$  unfolding cont_def by (metis mult.assoc
mult.commute)
      thus ?thesis by auto
    next
      assume p  $\notin$  s
      then have w p = 0 using homo unfolding w_def by auto
      then show ?thesis by auto
    qed
  qed
  also have ... = c *  $(\sum p \in pz. w p * h p * zorder f p)$ 
  unfolding sum_distrib_left by (simp add:algebra_simps)
  finally have contour_integral g ff = c *  $(\sum p \in pz. w p * h p * of_int (zorder f
p))$ .
  then show ?thesis unfolding ff_def c_def w_def by simp
qed

```

6.3 Coefficient asymptotics for generating functions

For a formal power series that has a meromorphic continuation on some disc in the context plane, we can use the Residue Theorem to extract precise asymptotic information from the residues at the poles. This can be used to derive the asymptotic behaviour of the coefficients ($a_n \sim \dots$). If the function is meromorphic on the entire complex plane, one can even derive a full asymptotic expansion.

We will first show the relationship between the coefficients and the sum over the residues with an explicit remainder term (the contour integral along the circle used in the Residue theorem).

theorem

fixes $f :: \text{complex} \Rightarrow \text{complex}$ **and** $n :: \text{nat}$ **and** $r :: \text{real}$
defines $g \equiv (\lambda w. f w / w^{\wedge} \text{Suc } n)$ **and** $\gamma \equiv \text{circlepath } 0 r$
assumes $\text{open } A$ $\text{cball } 0 r \subseteq A$ $r > 0$
assumes $f \text{ holomorphic_on } A - S$ $S \subseteq \text{ball } 0 r$ $\text{finite } S$ $0 \notin S$
shows $\text{fps_coeff_conv_residues}$:
 $(\text{deriv } \sim^n) f 0 / \text{fact } n =$
 $\text{contour_integral } \gamma g / (2 * \text{pi} * i) - (\sum z \in S. \text{residue } g z)$ **(is ?thesis1)**
and $\text{fps_coeff_residues_bound}$:
 $(\bigwedge z. \text{norm } z = r \implies z \notin k \implies \text{norm } (f z) \leq C) \implies C \geq 0 \implies \text{finite}$
 $k \implies$
 $\text{norm } ((\text{deriv } \sim^n) f 0 / \text{fact } n + (\sum z \in S. \text{residue } g z)) \leq C / r^{\wedge} n$

proof –

have $\text{holo}: g \text{ holomorphic_on } A - \text{insert } 0 S$
unfolding g_def **using** assms **by** $(\text{auto intro!}: \text{holomorphic_intros})$
have $\text{contour_integral } \gamma g = 2 * \text{pi} * i * (\sum z \in \text{insert } 0 S. \text{winding_number } \gamma z * \text{residue } g z)$
proof $(\text{rule Residue_theorem})$
show $g \text{ holomorphic_on } A - \text{insert } 0 S$ **by fact**
from assms **show** $\forall z. z \notin A \longrightarrow \text{winding_number } \gamma z = 0$
unfolding γ_def **by** $(\text{intro allI impI winding_number_zero_outside}[of _ \text{cball } 0 r]) \text{ auto}$
qed $(\text{insert } \text{assms}, \text{auto simp: } \gamma_def)$
also have $\text{winding_number } \gamma z = 1$ **if** $z \in \text{insert } 0 S$ **for** z
unfolding γ_def **using** assms that **by** $(\text{intro winding_number_circlepath}) \text{ auto}$
hence $(\sum z \in \text{insert } 0 S. \text{winding_number } \gamma z * \text{residue } g z) = (\sum z \in \text{insert } 0 S. \text{residue } g z)$
by $(\text{intro sum.cong}) \text{ simp_all}$
also have $\dots = \text{residue } g 0 + (\sum z \in S. \text{residue } g z)$
using $\langle 0 \notin S \rangle$ **and** $\langle \text{finite } S \rangle$ **by** $(\text{subst sum.insert}) \text{ auto}$
also from $\langle r > 0 \rangle$ **have** $0 \in \text{cball } 0 r$ **by simp**
with assms **have** $0 \in A - S$ **by blast**
with assms **have** $\text{residue } g 0 = (\text{deriv } \sim^n) f 0 / \text{fact } n$
unfolding g_def **by** $(\text{subst residue_holomorphic_over_power}[of } A - S])$
 $(\text{auto simp: finite_imp_closed})$
finally show $?thesis1$
by $(\text{simp add: field_simps})$

assume $C: \bigwedge z. \text{norm } z = r \implies z \notin k \implies \text{norm } (f z) \leq C$ $C \geq 0$ **and** $k: \text{finite } k$

have $(\text{deriv } \sim^n) f 0 / \text{fact } n + (\sum z \in S. \text{residue } g z) = \text{contour_integral } \gamma g / (2 * \text{pi} * i)$
using $\langle ?thesis1 \rangle$ **by** $(\text{simp add: algebra_simps})$
also have $\text{norm } \dots = \text{norm } (\text{contour_integral } \gamma g) / (2 * \text{pi})$
by $(\text{simp add: norm_divide norm_mult})$
also have $\text{norm } (\text{contour_integral } \gamma g) \leq C / r^{\wedge} \text{Suc } n * (2 * \text{pi} * r)$

```

proof (rule has_contour_integral_bound_circlepath_strong)
  from ⟨open A⟩ and ⟨finite S⟩ have open (A - insert 0 S)
    by (blast intro: finite_imp_closed)
  with assms show (g has_contour_integral contour_integral  $\gamma$  g) (circlepath 0
r)
  unfolding  $\gamma\_def$ 
  by (intro has_contour_integral_integral contour_integrable_holomorphic_simple
[OF holo]) auto
  next
  fix z assume z: norm (z - 0) = r z  $\notin$  k
  hence norm (g z) = norm (f z) / r ^ Suc n
    by (simp add: norm_divide g_def norm_mult norm_power)
  also have ...  $\leq$  C / r ^ Suc n
    using k and ⟨r > 0⟩ and z by (intro divide_right_mono C zero_le_power)
  auto
  finally show norm (g z)  $\leq$  C / r ^ Suc n .
  qed (insert C(2) k ⟨r > 0⟩, auto)
  also from ⟨r > 0⟩ have C / r ^ Suc n * (2 * pi * r) / (2 * pi) = C / r ^ n
    by simp
  finally show norm ((deriv  $\sim$  n) f 0 / fact n + ( $\sum$  z $\in$ S. residue g z))  $\leq$  ...
    by - (simp_all add: divide_right_mono)
qed

```

Since the circle is fixed, we can get an upper bound on the values of the generating function on the circle and therefore show that the integral is $O(r^{-n})$.

corollary *fps_coeff_residues_bigo*:

```

fixes f :: complex  $\Rightarrow$  complex and r :: real
assumes open A connected A cball 0 r  $\subseteq$  A r > 0
assumes f holomorphic_on A - S S  $\subseteq$  ball 0 r finite S 0  $\notin$  S
assumes g: eventually ( $\lambda$ n. g n = -( $\sum$  z $\in$ S. residue ( $\lambda$ z. f z / z ^ Suc n) z))
sequentially
  (is eventually ( $\lambda$ n.  $\_ = -?g'$  n)  $\_$ )
shows ( $\lambda$ n. (deriv  $\sim$  n) f 0 / fact n - g n)  $\in$  O( $\lambda$ n. 1 / r ^ n) (is ( $\lambda$ n. ?c n
-  $\_$ )  $\in$  O( $\_$ ))
proof -
from assms have compact (f ' sphere 0 r)
  by (intro compact_continuous_image holomorphic_on_imp_continuous_on
holomorphic_on_subset[OF ⟨f holomorphic_on A - S⟩]) auto
hence bounded (f ' sphere 0 r) by (rule compact_imp_bounded)
then obtain C where C:  $\bigwedge$ z. z  $\in$  sphere 0 r  $\implies$  norm (f z)  $\leq$  C
  by (auto simp: bounded_iff_sphere_def)
have 0  $\leq$  norm (f (of_real r)) by simp
also from C[of of_real r] and ⟨r > 0⟩ have ...  $\leq$  C by simp
finally have C_nonneg: C  $\geq$  0 .

have ( $\lambda$ n. ?c n + ?g' n)  $\in$  O( $\lambda$ n. of_real (1 / r ^ n))
proof (intro bigoI[of  $\_ C$ ] always_eventually_allI )
  fix n :: nat

```

```

from assms and C and C_nonneg have  $\text{norm } (?c\ n + ?g'\ n) \leq C / r \wedge n$ 
  by (intro fps_coeff_residues_bound[where A = A and k = {}]) auto
also have  $\dots = C * \text{norm } (\text{complex\_of\_real } (1 / r \wedge n))$ 
  using  $\langle r > 0 \rangle$  by (simp add: norm_divide norm_power)
finally show  $\text{norm } (?c\ n + ?g'\ n) \leq \dots$  .
qed
also have  $?this \iff (\lambda n. ?c\ n - g\ n) \in O(\lambda n. \text{of\_real } (1 / r \wedge n))$ 
  by (intro landau_o.big.in_cong eventually_mono[OF g]) simp_all
finally show ?thesis .
qed

corollary fps_coeff_residues_bigo':
  fixes f :: complex  $\Rightarrow$  complex and r :: real
  assumes exp: f has_fps_expansion F
  assumes open A connected A cball 0 r  $\subseteq$  A r  $>$  0
  assumes f holomorphic_on A - S  $S \subseteq \text{ball } 0\ r$  finite S  $0 \notin S$ 
  assumes eventually  $(\lambda n. g\ n = -(\sum_{z \in S. \text{residue } (\lambda z. f\ z / z \wedge \text{Suc } n)} z))$ 
sequentially
    (is eventually  $(\lambda n. \_ = -?g'\ n)$   $\_$ )
  shows  $(\lambda n. \text{fps\_nth } F\ n - g\ n) \in O(\lambda n. 1 / r \wedge n)$  (is  $(\lambda n. ?c\ n - \_)$   $\in$ 
 $O(\_)$ )
proof -
  have  $\text{fps\_nth } F = (\lambda n. (\text{deriv } \wedge n) f\ 0 / \text{fact } n)$ 
    using fps_nth_fps_expansion[OF exp] by (intro ext) simp_all
  with fps_coeff_residues_bigo[OF assms(2-)] show ?thesis by simp
qed

```

6.4 Rouché's theorem

```

theorem Rouche_theorem:
  fixes f g :: complex  $\Rightarrow$  complex and s :: complex set
  defines  $fg \equiv (\lambda p. f\ p + g\ p)$ 
  defines  $\text{zeros\_fg} \equiv \{p \in s. fg\ p = 0\}$  and  $\text{zeros\_f} \equiv \{p \in s. f\ p = 0\}$ 
  assumes
    open s and connected s and
    finite  $\text{zeros\_fg}$  and
    finite  $\text{zeros\_f}$  and
    f_holo: f holomorphic_on s and
    g_holo: g holomorphic_on s and
    valid_path  $\gamma$  and
    loop:  $\text{pathfinish } \gamma = \text{pathstart } \gamma$  and
    path_img:  $\text{path\_image } \gamma \subseteq s$  and
    path_less:  $\forall z \in \text{path\_image } \gamma. \text{cmod}(f\ z) > \text{cmod}(g\ z)$  and
    homo:  $\forall z. (z \notin s) \longrightarrow \text{winding\_number } \gamma\ z = 0$ 
  shows  $(\sum_{p \in \text{zeros\_fg}. \text{winding\_number } \gamma\ p * \text{zorder } fg\ p}$ 
     $= (\sum_{p \in \text{zeros\_f}. \text{winding\_number } \gamma\ p * \text{zorder } f\ p})$ 
proof -
  have  $\text{path\_fg}: \text{path\_image } \gamma \subseteq s - \text{zeros\_fg}$ 
proof -

```

```

    have False when  $z \in \text{path\_image } \gamma$  and  $f z + g z = 0$  for  $z$ 
    proof -
      have  $\text{cmod } (f z) > \text{cmod } (g z)$  using  $\langle z \in \text{path\_image } \gamma \rangle \text{ path\_less}$  by auto
      moreover have  $f z = -g z$  using  $\langle f z + g z = 0 \rangle$  by (simp add:
    eq_neg_iff_add_eq_0)
      then have  $\text{cmod } (f z) = \text{cmod } (g z)$  by auto
      ultimately show False by auto
    qed
    then show ?thesis unfolding zeros_fg_def fg_def using path_img by auto
  qed
  have path_f:  $\text{path\_image } \gamma \subseteq s - \text{zeros}_f$ 
  proof -
    have False when  $z \in \text{path\_image } \gamma$  and  $f z = 0$  for  $z$ 
    proof -
      have  $\text{cmod } (g z) < \text{cmod } (f z)$  using  $\langle z \in \text{path\_image } \gamma \rangle \text{ path\_less}$  by auto
      then have  $\text{cmod } (g z) < 0$  using  $\langle f z = 0 \rangle$  by auto
      then show False by auto
    qed
    then show ?thesis unfolding zeros_f_def using path_img by auto
  qed
  define w where  $w \equiv \lambda p. \text{winding\_number } \gamma p$ 
  define c where  $c \equiv 2 * \text{complex\_of\_real } \pi * i$ 
  define h where  $h \equiv \lambda p. g p / f p + 1$ 
  obtain spikes
    where finite_spikes and spikes:  $\forall x \in \{0..1\} - \text{spikes}. \gamma \text{ differentiable at } x$ 
    using  $\langle \text{valid\_path } \gamma \rangle$ 
  by (auto simp: valid_path_def piecewise_C1_differentiable_on_def C1_differentiable_on_eq)
  have h_contour:  $((\lambda x. \text{deriv } h x / h x) \text{ has\_contour\_integral } 0) \gamma$ 
  proof -
    have outside_img:  $0 \in \text{outside } (\text{path\_image } (h \circ \gamma))$ 
    proof -
      have  $h p \in \text{ball } 1 1$  when  $p \in \text{path\_image } \gamma$  for  $p$ 
      proof -
        have  $\text{cmod } (g p / f p) < 1$ 
          by (smt (verit) divide_less_eq_1_pos norm_divide norm_ge_zero
    path_less that)
        then show ?thesis unfolding h_def by (auto simp add: dist_complex_def)
      qed
      then have  $\text{path\_image } (h \circ \gamma) \subseteq \text{ball } 1 1$ 
        by (simp add: image_subset_iff path_image_compose)
      moreover have  $(0::\text{complex}) \notin \text{ball } 1 1$  by (simp add: dist_norm)
      ultimately show ?thesis
        using convex_in_outside[of ball 1 1 0] outside_mono by blast
    qed
    have valid_h:  $\text{valid\_path } (h \circ \gamma)$ 
  proof (rule valid_path_compose_holomorphic[OF  $\langle \text{valid\_path } \gamma \rangle$  __ path_f])
    show  $h \text{ holomorphic\_on } s - \text{zeros}_f$ 
      unfolding h_def using f_holo g_holo
      by (auto intro!: holomorphic_intros simp add: zeros_f_def)
  qed

```



```

next
  show open (s - zeros_f) using ⟨finite zeros_f⟩ ⟨open s⟩ finite_imp_closed
    by auto
qed
have ((λz. 1/z) has_contour_integral 0) (h ∘ γ)
proof -
  have 0 ∉ path_image (h ∘ γ) using outside_img by (simp add: outside_def)
  then have ((λz. 1/z) has_contour_integral c * winding_number (h ∘ γ) 0)
(h ∘ γ)
    using has_contour_integral_winding_number[of h ∘ γ 0,simplified] valid_h
      unfolding c_def by auto
  moreover have winding_number (h ∘ γ) 0 = 0
proof -
  have 0 ∈ outside (path_image (h ∘ γ)) using outside_img .
  moreover have path (h ∘ γ)
    using valid_h by (simp add: valid_path_imp_path)
  moreover have pathfinish (h ∘ γ) = pathstart (h ∘ γ)
    by (simp add: loop_pathfinish_compose_pathstart_compose)
  ultimately show ?thesis using winding_number_zero_in_outside by auto
qed
ultimately show ?thesis by auto
qed
moreover have vector_derivative (h ∘ γ) (at x) = vector_derivative γ (at x)
* deriv h (γ x)
  when x ∈ {0..1} - spikes for x
proof (rule vector_derivative_chain_at_general)
  show γ differentiable at x using that ⟨valid_path γ⟩ spikes by auto
next
  define der where der ≡ λp. (deriv g p * f p - g p * deriv f p)/(f p * f p)
  define t where t ≡ γ x
  have f t ≠ 0 unfolding zeros_f_def t_def
    by (metis DiffD1 image_eqI norm_not_less_zero norm_zero path_defs(4)
path_less that)
  moreover have t ∈ s
    using contra_subsetD path_image_def path_fg t_def that by fastforce
  ultimately have (h has_field_derivative der t) (at t)
    unfolding h_def der_def using g_holo f_holo ⟨open s⟩
    by (auto intro!: holomorphic_derivI derivative_eq_intros)
  then show h field_differentiable at (γ x)
    unfolding t_def field_differentiable_def by blast
qed
then have ((/) 1 has_contour_integral 0) (h ∘ γ)
  = ((λx. deriv h x / h x) has_contour_integral 0) γ
  unfolding has_contour_integral
  apply (intro has_integral_spike_eq[OF negligible_finite, OF ⟨finite spikes⟩])
  by auto
ultimately show ?thesis by auto
qed
then have contour_integral γ (λx. deriv h x / h x) = 0

```

```

using contour_integral_unique by simp
moreover have contour_integral  $\gamma$   $(\lambda x. \text{deriv } fg \ x / fg \ x) = \text{contour\_integral } \gamma$ 
 $(\lambda x. \text{deriv } f \ x / f \ x)$ 
  + contour_integral  $\gamma$   $(\lambda p. \text{deriv } h \ p / h \ p)$ 
proof -
  have  $(\lambda p. \text{deriv } f \ p / f \ p)$  contour_integrable_on  $\gamma$ 
  proof (rule contour_integrable_holomorphic_simple[OF __  $\langle \text{valid\_path } \gamma \rangle$ 
path_f])
    show open  $(s - \text{zeros}_f)$  using finite_imp_closed[OF  $\langle \text{finite zeros}_f \rangle$ ]  $\langle \text{open } s \rangle$ 
    by auto
  then show  $(\lambda p. \text{deriv } f \ p / f \ p)$  holomorphic_on  $s - \text{zeros}_f$ 
    using f_holo
    by (auto intro!: holomorphic_intros simp add:zeros_f_def)
qed
moreover have  $(\lambda p. \text{deriv } h \ p / h \ p)$  contour_integrable_on  $\gamma$ 
  using h_contour
  by (simp add: has_contour_integral_integrable)
ultimately have contour_integral  $\gamma$   $(\lambda x. \text{deriv } f \ x / f \ x + \text{deriv } h \ x / h \ x) =$ 
  contour_integral  $\gamma$   $(\lambda p. \text{deriv } f \ p / f \ p) + \text{contour\_integral } \gamma$ 
 $(\lambda p. \text{deriv } h \ p / h \ p)$ 
  using contour_integral_add[of  $(\lambda p. \text{deriv } f \ p / f \ p)$   $\gamma$   $(\lambda p. \text{deriv } h \ p / h \ p)$ ]
  by auto
moreover have deriv fg p / fg p = deriv f p / f p + deriv h p / h p
  when  $p \in \text{path\_image } \gamma$  for  $p$ 
proof -
  have fg p  $\neq 0$  and f p  $\neq 0$  using path_f path_fg that unfolding zeros_f_def
zeros_fg_def
  by auto
  have h p  $\neq 0$ 
  proof (rule ccontr)
    assume  $\neg h \ p \neq 0$ 
    then have g p / f p = -1 unfolding h_def by (simp add: add_eq_0_iff2)
    then have cmod  $(g \ p / f \ p) = 1$  by auto
    moreover have cmod  $(g \ p / f \ p) < 1$ 
    by (simp add: f p  $\neq 0$ ) norm_divide path_less that)
    ultimately show False by auto
  qed
  have der_fg:deriv fg p = deriv f p + deriv g p unfolding fg_def
  using f_holo g_holo holomorphic_on_imp_differentiable_at[OF __  $\langle \text{open } s \rangle$ ]
path_img that
  by auto
  have der_h:deriv h p =  $(\text{deriv } g \ p * f \ p - g \ p * \text{deriv } f \ p) / (f \ p * f \ p)$ 
proof -
  define der where der  $\equiv \lambda p. (\text{deriv } g \ p * f \ p - g \ p * \text{deriv } f \ p) / (f \ p * f \ p)$ 
  have  $p \in s$  using path_img that by auto
  then have  $(h \ \text{has\_field\_derivative } der \ p)$   $(at \ p)$ 
  unfolding h_def der_def using g_holo f_holo  $\langle \text{open } s \rangle$   $\langle f \ p \neq 0 \rangle$ 
  by (auto intro!: derivative_eq_intros holomorphic_derivI)

```

```

    then show ?thesis unfolding der_def using DERIV_imp_deriv by auto
  qed
  show ?thesis
    apply (simp only: der_fg der_h)
    apply (auto simp add: field_simps ‹h p ≠ 0› ‹f p ≠ 0› ‹fg p ≠ 0›)
    by (auto simp add: field_simps h_def ‹f p ≠ 0› fg_def)
  qed
  then have contour_integral  $\gamma$  ( $\lambda p. \text{deriv } fg \ p / fg \ p$ )
    = contour_integral  $\gamma$  ( $\lambda p. \text{deriv } f \ p / f \ p + \text{deriv } h \ p / h \ p$ )
    by (elim contour_integral_eq)
  ultimately show ?thesis by auto
  qed
  moreover have contour_integral  $\gamma$  ( $\lambda x. \text{deriv } fg \ x / fg \ x$ ) =  $c * (\sum_{p \in \text{zeros\_fg}} w \ p * \text{zorder } fg \ p)$ 
  proof -
    have fg holomorphic_on s unfolding fg_def using f_holo g_holo holomorphic_on_add by auto
    moreover
      have path_image  $\gamma \subseteq s - \{p \in s. fg \ p = 0\}$  using path_fg unfolding zeros_fg_def .
    moreover
      have finite  $\{p \in s. fg \ p = 0\}$  using ‹finite zeros_fg› unfolding zeros_fg_def .
    ultimately show ?thesis
      unfolding c_def zeros_fg_def w_def
      using argument_principle[OF ‹open s› ‹connected s› _ _ ‹valid_path  $\gamma$ › loop _ homo, of _ {}  $\lambda_. 1$ ]
      by simp
    qed
    moreover have contour_integral  $\gamma$  ( $\lambda x. \text{deriv } f \ x / f \ x$ ) =  $c * (\sum_{p \in \text{zeros\_f}} w \ p * \text{zorder } f \ p)$ 
      unfolding c_def zeros_f_def w_def
    proof (rule argument_principle[OF ‹open s› ‹connected s› _ _ ‹valid_path  $\gamma$ › loop _ homo, of _ {}  $\lambda_. 1, \text{simplified}$ ])
      show f holomorphic_on s using f_holo g_holo holomorphic_on_add by auto
      show path_image  $\gamma \subseteq s - \{p \in s. f \ p = 0\}$  using path_f unfolding zeros_f_def .
    .
      show finite  $\{p \in s. f \ p = 0\}$  using ‹finite zeros_f› unfolding zeros_f_def .
    qed
    ultimately have  $c * (\sum_{p \in \text{zeros\_fg}} w \ p * (\text{zorder } fg \ p)) = c * (\sum_{p \in \text{zeros\_f}} w \ p * (\text{zorder } f \ p))$ 
      by auto
    then show ?thesis unfolding c_def using w_def by auto
  qed
end
theory Laurent_Convergence
imports HOL-Computational_Algebra.Formal_Laurent_Series HOL-Library.Landau_Symbols
  Residue_Theorem

```

begin

TODO: Better than $?f \text{ field_differentiable at } (?c * ?z) \implies \text{deriv } (\lambda w. ?f (?c * w)) ?z = ?c * \text{deriv } ?f (?c * ?z)?$

lemma *deriv_compose_linear'*:

assumes $f \text{ field_differentiable at } (c * z+a)$
shows $\text{deriv } (\lambda w. f (c * w+a)) z = c * \text{deriv } f (c * z+a)$
apply (*subst deriv_chain[where f= $\lambda w. c * w+a$, unfolded comp_def]*)
using *assms* **by** (*auto intro:derivative_intros*)

TODO: Better than $[[?f \text{ holomorphic_on } ?T; \text{open } ?S; \text{open } ?T; ?z \in ?S; \bigwedge w. w \in ?S \implies ?u * w \in ?T]] \implies \text{deriv}^{?n} (\lambda w. ?f (?u * w)) ?z = ?u^{?n} * \text{deriv}^{?n} ?f (?u * ?z)?$

lemma *higher_deriv_compose_linear'*:

fixes $z::\text{complex}$
assumes $f: \text{holomorphic_on } T \text{ and } S: \text{open } S \text{ and } T: \text{open } T \text{ and } z: z \in S$
and $fg: \bigwedge w. w \in S \implies u * w+c \in T$
shows $(\text{deriv } \hat{\sim} n) (\lambda w. f (u * w+c)) z = u \hat{\sim} n * (\text{deriv } \hat{\sim} n) f (u * z+c)$
using z
proof (*induction n arbitrary: z*)
case 0 **then show** *?case* **by** *simp*
next
case (*Suc n z*)
have *holo0*: $f \text{ holomorphic_on } (\lambda w. u * w+c) \text{ ' } S$
by (*meson fg f holomorphic_on_subset image_subset_iff*)
have *holo2*: $(\text{deriv } \hat{\sim} n) f \text{ holomorphic_on } (\lambda w. u * w+c) \text{ ' } S$
by (*meson fg holomorphic_higher_deriv holomorphic_on_subset image_subset_iff T*)
have *holo3*: $(\lambda z. u \hat{\sim} n * (\text{deriv } \hat{\sim} n) f (u * z+c)) \text{ holomorphic_on } S$
by (*intro holo2 holomorphic_on_compose [where g= $(\text{deriv } \hat{\sim} n) f$, unfolded o_def] holomorphic_intros*)
have $(\lambda w. u * w+c) \text{ holomorphic_on } S \text{ f holomorphic_on } (\lambda w. u * w+c) \text{ ' } S$
by (*rule holo0 holomorphic_intros*)+
then have *holo1*: $(\lambda w. f (u * w+c)) \text{ holomorphic_on } S$
by (*rule holomorphic_on_compose [where g= f , unfolded o_def]*)
have $\text{deriv } ((\text{deriv } \hat{\sim} n) (\lambda w. f (u * w+c))) z = \text{deriv } (\lambda z. u \hat{\sim} n * (\text{deriv } \hat{\sim} n) f (u * z+c)) z$
proof (*rule complex_derivative_transform_within_open [OF_ holo3 S Suc.premis]*)
show $(\text{deriv } \hat{\sim} n) (\lambda w. f (u * w+c)) \text{ holomorphic_on } S$
by (*rule holomorphic_higher_deriv [OF holo1 S]*)
qed (*simp add: Suc.IH*)
also have $\dots = u \hat{\sim} n * \text{deriv } (\lambda z. (\text{deriv } \hat{\sim} n) f (u * z+c)) z$
proof –
have $(\text{deriv } \hat{\sim} n) f \text{ analytic_on } T$
by (*simp add: analytic_on_open f holomorphic_higher_deriv T*)
then have $(\lambda w. (\text{deriv } \hat{\sim} n) f (u * w+c)) \text{ analytic_on } S$
proof –
have $(\text{deriv } \hat{\sim} n) f \circ (\lambda w. u * w+c) \text{ holomorphic_on } S$

```

      using holomorphic_on_compose[OF __ holo2] ⟨(λw. u * w+c) holomor-
phic_on S⟩
    by simp
  then show ?thesis
    by (simp add: S analytic_on_open o_def)
  qed
  then show ?thesis
    by (intro deriv_cmult analytic_on_imp_differentiable_at [OF __ Suc.prem])
  qed
  also have ... = u * u ^ n * deriv ((deriv ^ n) f) (u * z+c)
  proof -
    have (deriv ^ n) f field_differentiable at (u * z+c)
      using Suc.prem T ffg holomorphic_higher_deriv holomorphic_on_imp_differentiable_at
    by blast
    then show ?thesis
      by (simp add: deriv_compose_linear')
  qed
  finally show ?case
    by simp
  qed

```

```

lemma fps_to_fls_numeral [simp]: fps_to_fls (numeral n) = numeral n
  by (metis fps_to_fls_of_nat of_nat_numeral)

```

```

lemma fls_const_power: fls_const (a ^ b) = fls_const a ^ b
  by (induction b) (auto simp flip: fls_const_mult_const)

```

```

lemma fls_deriv_numeral [simp]: fls_deriv (numeral n) = 0
  by (metis fls_deriv_of_int of_int_numeral)

```

```

lemma fls_const_numeral [simp]: fls_const (numeral n) = numeral n
  by (metis fls_of_nat of_nat_numeral)

```

```

lemma fls_mult_of_int_nth [simp]:
  shows fls_nth (numeral k * f) n = numeral k * fls_nth f n
  and fls_nth (f * numeral k) n = fls_nth f n * numeral k
  by (metis fls_const_numeral fls_mult_const_nth)+

```

```

lemma fls_nth_numeral' [simp]:
  fls_nth (numeral n) 0 = numeral n k ≠ 0 ⇒ fls_nth (numeral n) k = 0
  by (subst fls_const_numeral [symmetric], subst fls_const_nth, simp)+

```

```

lemma fls_subdegree_prod:
  fixes F :: 'a ⇒ 'b :: field_char_0 fls
  assumes ∧x. x ∈ I ⇒ F x ≠ 0
  shows fls_subdegree (∏ x∈I. F x) = (∑ x∈I. fls_subdegree (F x))
  using assms by (induction I rule: infinite_finite_induct) auto

```

```

lemma fls_subdegree_prod':

```

```

fixes  $F :: 'a \Rightarrow 'b :: \text{field\_char\_0}$   $fls$ 
assumes  $\bigwedge x. x \in I \implies fls\_subdegree (F x) \neq 0$ 
shows  $fls\_subdegree (\prod_{x \in I}. F x) = (\sum_{x \in I}. fls\_subdegree (F x))$ 
proof (intro fls_subdegree_prod)
  show  $F x \neq 0$  if  $x \in I$  for  $x$ 
    using assms[OF that] by auto
qed

```

```

instance  $fps :: (\text{semiring\_char\_0}) \text{semiring\_char\_0}$ 
proof
  show inj (of_nat :: nat  $\Rightarrow$  'a fps)
  proof
    fix  $m n :: \text{nat}$ 
    assume  $of\_nat\ m = (of\_nat\ n :: 'a\ fps)$ 
    hence  $fps\_nth (of\_nat\ m)\ 0 = (fps\_nth (of\_nat\ n)\ 0 :: 'a)$ 
      by (simp only:)
    thus  $m = n$ 
      by simp
  qed
qed

```

```

instance  $fls :: (\text{semiring\_char\_0}) \text{semiring\_char\_0}$ 
proof
  show inj (of_nat :: nat  $\Rightarrow$  'a fls)
  proof
    fix  $m n :: \text{nat}$ 
    assume  $of\_nat\ m = (of\_nat\ n :: 'a\ fls)$ 
    hence  $fls\_nth (of\_nat\ m)\ 0 = (fls\_nth (of\_nat\ n)\ 0 :: 'a)$ 
      by (simp only:)
    thus  $m = n$ 
      by (simp add: fls_of_nat_nth)
  qed
qed

```

```

lemma  $fls\_const\_eq\_0\_iff$  [simp]:  $fls\_const\ c = 0 \iff c = 0$ 
  using  $fls\_const\_0\ fls\_const\_nonzero$  by blast

```

```

lemma  $fls\_subdegree\_add\_eq1$ :
  assumes  $f \neq 0$   $fls\_subdegree\ f < fls\_subdegree\ g$ 
  shows  $fls\_subdegree (f + g) = fls\_subdegree\ f$ 
proof (intro antisym)
  from assms have  $*$ :  $fls\_nth (f + g) (fls\_subdegree\ f) \neq 0$ 
    by auto
  from  $*$  show  $fls\_subdegree (f + g) \leq fls\_subdegree\ f$ 
    by (rule fls_subdegree_leI)
  from  $*$  have  $f + g \neq 0$ 
    using  $fls\_nonzeroI$  by blast
  thus  $fls\_subdegree\ f \leq fls\_subdegree (f + g)$ 
    using assms(2)  $fls\_plus\_subdegree$  by force

```

qed

lemma *fls_subdegree_add_eq2*:
 assumes $g \neq 0$ $\text{fls_subdegree } g < \text{fls_subdegree } f$
 shows $\text{fls_subdegree } (f + g) = \text{fls_subdegree } g$
proof (*intro antisym*)
 from *assms* **have** *: $\text{fls_nth } (f + g) (\text{fls_subdegree } g) \neq 0$
 by *auto*
 from * **show** $\text{fls_subdegree } (f + g) \leq \text{fls_subdegree } g$
 by (*rule fls_subdegree_leI*)
 from * **have** $f + g \neq 0$
 using *fls_nonzeroI* **by** *blast*
 thus $\text{fls_subdegree } g \leq \text{fls_subdegree } (f + g)$
 using *assms*(2) *fls_plus_subdegree* **by** *force*
 qed

lemma *fls_subdegree_diff_eq1*:
 assumes $f \neq 0$ $\text{fls_subdegree } f < \text{fls_subdegree } g$
 shows $\text{fls_subdegree } (f - g) = \text{fls_subdegree } f$
 using *fls_subdegree_add_eq1*[*of* $f - g$] *assms* **by** *simp*

lemma *fls_subdegree_diff_eq2*:
 assumes $g \neq 0$ $\text{fls_subdegree } g < \text{fls_subdegree } f$
 shows $\text{fls_subdegree } (f - g) = \text{fls_subdegree } g$
 using *fls_subdegree_add_eq2*[*of* $-g$ f] *assms* **by** *simp*

lemma *nat_minus_fls_subdegree_plus_const_eq*:
 $\text{nat } (-\text{fls_subdegree } (F + \text{fls_const } c)) = \text{nat } (-\text{fls_subdegree } F)$
proof (*cases fls_subdegree F < 0*)
 case *True*
 hence $\text{fls_subdegree } (F + \text{fls_const } c) = \text{fls_subdegree } F$
 by (*intro fls_subdegree_add_eq1*) *auto*
 thus *?thesis*
 by *simp*
 next
 case *False*
 thus *?thesis*
 by (*auto simp: fls_subdegree_ge0I*)
 qed

lemma *at_to_0'*: $\text{NO_MATCH } 0 z \implies \text{at } z = \text{filtermap } (\lambda x. x + z) (\text{at } 0)$
 for $z :: 'a :: \text{real_normed_vector}$
 by (*rule at_to_0*)

lemma *nhds_to_0*: $\text{nhds } (x :: 'a :: \text{real_normed_vector}) = \text{filtermap } ((+) x) (\text{nhds } 0)$
proof –
 have $(\lambda xa. xa - - x) = (+) x$
 by *auto*

thus *?thesis*
using *filtermap_nhds_shift[of -x 0]* **by** *simp*
qed

lemma *nhds_to_0'*: *NO_MATCH 0 x* \implies *nhds* (*x* :: 'a :: *real_normed_vector*)
 $=$ *filtermap* ((+) *x*) (*nhds 0*)
by (*rule nhds_to_0*)

definition *fls_conv_radius* :: *complex fls* \Rightarrow *ereal* **where**
fls_conv_radius f = *fps_conv_radius* (*fls_regpert f*)

definition *eval_fls* :: *complex fls* \Rightarrow *complex* \Rightarrow *complex* **where**
eval_fls F z = *eval_fps* (*fls_base_factor_to_fps F*) *z* * *z* *powi fls_subdegree F*

definition
has_laurent_expansion :: (*complex* \Rightarrow *complex*) \Rightarrow *complex fls* \Rightarrow *bool*
(infixl *has'_laurent'_expansion 60*)
where (*f has_laurent_expansion F*) \longleftrightarrow
fls_conv_radius F > 0 \wedge *eventually* ($\lambda z. \text{eval_fls } F z = f z$) (*at 0*)

lemma *has_laurent_expansion_schematicI*:
f has_laurent_expansion F \implies *F = G* \implies *f has_laurent_expansion G*
by *simp*

lemma *has_laurent_expansion_cong*:
assumes *eventually* ($\lambda x. f x = g x$) (*at 0*) *F = G*
shows (*f has_laurent_expansion F*) \longleftrightarrow (*g has_laurent_expansion G*)
proof –
have *eventually* ($\lambda z. \text{eval_fls } F z = g z$) (*at 0*)
if *eventually* ($\lambda z. \text{eval_fls } F z = f z$) (*at 0*) *eventually* ($\lambda x. f x = g x$) (*at 0*)
for *f g*
using *that* **by** *eventually_elim auto*
from *this[of f g]* *this[of g f]* **show** *?thesis*
using *assms* **by** (*auto simp: eq_commute has_laurent_expansion_def*)
qed

lemma *has_laurent_expansion_cong'*:
assumes *eventually* ($\lambda x. f x = g x$) (*at z*) *F = G* *z = z'*
shows (($\lambda x. f (z + x)$) *has_laurent_expansion F*) \longleftrightarrow (($\lambda x. g (z' + x)$)
has_laurent_expansion G)
by (*intro has_laurent_expansion_cong*)
(use assms in <auto simp: at_to_0' eventually_filtermap add_ac>)

lemma *fls_conv_radius_altdef*:
fls_conv_radius F = *fps_conv_radius* (*fls_base_factor_to_fps F*)
proof –
have *conv_radius* ($\lambda n. \text{fls_nth } F (\text{int } n)$) = *conv_radius* ($\lambda n. \text{fls_nth } F (\text{int } n$
 $+ \text{fls_subdegree } F)$)


```

proof (cases fls_subdegree F  $\geq 0$ )
  case True
    hence conv_radius ( $\lambda n.$  fls_nth F (int n + fls_subdegree F)) =
      conv_radius ( $\lambda n.$  fls_nth F (int (n + nat (fls_subdegree F))))
    by auto
    thus ?thesis
    by (subst (asm) conv_radius_shift) auto
  next
    case False
    hence conv_radius ( $\lambda n.$  fls_nth F (int n)) =
      conv_radius ( $\lambda n.$  fls_nth F (fls_subdegree F + int (n + nat (-fls_subdegree
F))))
    by auto
    thus ?thesis
    by (subst (asm) conv_radius_shift) (auto simp: add_ac)
  qed
thus ?thesis
by (simp add: fls_conv_radius_def fps_conv_radius_def)
qed

```

```

lemma eval_fps_of_nat [simp]: eval_fps (of_nat n) z = of_nat n
and eval_fps_of_int [simp]: eval_fps (of_int m) z = of_int m
by (simp_all flip: fps_of_nat fps_of_int)

```

```

lemma fls_subdegree_numeral [simp]: fls_subdegree (numeral n) = 0
by (metis fls_subdegree_of_nat of_nat_numeral)

```

```

lemma fls_regpart_numeral [simp]: fls_regpart (numeral n) = numeral n
by (metis fls_regpart_of_nat of_nat_numeral)

```

```

lemma fps_conv_radius_of_nat [simp]: fps_conv_radius (of_nat n) =  $\infty$ 
and fps_conv_radius_of_int [simp]: fps_conv_radius (of_int m) =  $\infty$ 
by (simp_all flip: fps_of_nat fps_of_int)

```

```

lemma fps_conv_radius_fls_regpart: fps_conv_radius (fls_regpart F) = fls_conv_radius
F
by (simp add: fls_conv_radius_def)

```

```

lemma fls_conv_radius_0 [simp]: fls_conv_radius 0 =  $\infty$ 
and fls_conv_radius_1 [simp]: fls_conv_radius 1 =  $\infty$ 
and fls_conv_radius_const [simp]: fls_conv_radius (fls_const c) =  $\infty$ 
and fls_conv_radius_numeral [simp]: fls_conv_radius (numeral num) =  $\infty$ 
and fls_conv_radius_of_nat [simp]: fls_conv_radius (of_nat n) =  $\infty$ 
and fls_conv_radius_of_int [simp]: fls_conv_radius (of_int m) =  $\infty$ 
and fls_conv_radius_X [simp]: fls_conv_radius fls_X =  $\infty$ 
and fls_conv_radius_X_inv [simp]: fls_conv_radius fls_X_inv =  $\infty$ 
and fls_conv_radius_X_intpow [simp]: fls_conv_radius (fls_X_intpow m) =
 $\infty$ 
by (simp_all add: fls_conv_radius_def fls_X_intpow_regpart)

```

lemma *fls_conv_radius_shift* [*simp*]: $\text{fls_conv_radius} (\text{fls_shift } n \ F) = \text{fls_conv_radius } F$

unfolding *fls_conv_radius_altdef* **by** (*subst fls_base_factor_to_fps_shift*) (*rule refl*)

lemma *fls_conv_radius_fps_to_fl* [*simp*]: $\text{fls_conv_radius} (\text{fps_to_fls } F) = \text{fps_conv_radius } F$

by (*simp add: fls_conv_radius_def*)

lemma *fls_conv_radius_deriv* [*simp*]: $\text{fls_conv_radius} (\text{fls_deriv } F) \geq \text{fls_conv_radius } F$

proof –

have $\text{fls_conv_radius} (\text{fls_deriv } F) = \text{fps_conv_radius} (\text{fls_regpart} (\text{fls_deriv } F))$

by (*simp add: fls_conv_radius_def*)

also have $\text{fls_regpart} (\text{fls_deriv } F) = \text{fps_deriv} (\text{fls_regpart } F)$

by (*intro fps_ext*) (*auto simp: add_ac*)

also have $\text{fps_conv_radius} \dots \geq \text{fls_conv_radius } F$

by (*simp add: fls_conv_radius_def fps_conv_radius_deriv*)

finally show *?thesis* .

qed

lemma *fls_conv_radius_uminus* [*simp*]: $\text{fls_conv_radius} (-F) = \text{fls_conv_radius } F$

by (*simp add: fls_conv_radius_def*)

lemma *fls_conv_radius_add*: $\text{fls_conv_radius} (F + G) \geq \min (\text{fls_conv_radius } F) (\text{fls_conv_radius } G)$

by (*simp add: fls_conv_radius_def fps_conv_radius_add*)

lemma *fls_conv_radius_diff*: $\text{fls_conv_radius} (F - G) \geq \min (\text{fls_conv_radius } F) (\text{fls_conv_radius } G)$

by (*simp add: fls_conv_radius_def fps_conv_radius_diff*)

lemma *fls_conv_radius_mult*: $\text{fls_conv_radius} (F * G) \geq \min (\text{fls_conv_radius } F) (\text{fls_conv_radius } G)$

proof (*cases F = 0 ∨ G = 0*)

case *False*

hence [*simp*]: $F \neq 0 \ G \neq 0$

by *auto*

have $\text{fls_conv_radius} (F * G) = \text{fps_conv_radius} (\text{fls_regpart} (\text{fls_shift} (\text{fls_subdegree } F + \text{fls_subdegree } G) (F * G)))$

by (*simp add: fls_conv_radius_altdef*)

also have $\text{fls_regpart} (\text{fls_shift} (\text{fls_subdegree } F + \text{fls_subdegree } G) (F * G))$

$=$

$\text{fls_base_factor_to_fps } F * \text{fls_base_factor_to_fps } G$

by (*simp add: fls_times_def*)

also have $\text{fps_conv_radius} \dots \geq \min (\text{fls_conv_radius } F) (\text{fls_conv_radius } G)$

unfolding *fls_conv_radius_altdef* **by** (rule *fps_conv_radius_mult*)
finally show *?thesis* .
qed *auto*

lemma *fps_conv_radius_add_ge*:
 $fps_conv_radius\ F \geq r \implies fps_conv_radius\ G \geq r \implies fps_conv_radius\ (F + G) \geq r$
using *fps_conv_radius_add[of F G]* **by** (*simp add: min_def split: if_splits*)

lemma *fps_conv_radius_diff_ge*:
 $fps_conv_radius\ F \geq r \implies fps_conv_radius\ G \geq r \implies fps_conv_radius\ (F - G) \geq r$
using *fps_conv_radius_diff[of F G]* **by** (*simp add: min_def split: if_splits*)

lemma *fps_conv_radius_mult_ge*:
 $fps_conv_radius\ F \geq r \implies fps_conv_radius\ G \geq r \implies fps_conv_radius\ (F * G) \geq r$
using *fps_conv_radius_mult[of F G]* **by** (*simp add: min_def split: if_splits*)

lemma *fls_conv_radius_add_ge*:
 $fls_conv_radius\ F \geq r \implies fls_conv_radius\ G \geq r \implies fls_conv_radius\ (F + G) \geq r$
using *fls_conv_radius_add[of F G]* **by** (*simp add: min_def split: if_splits*)

lemma *fls_conv_radius_diff_ge*:
 $fls_conv_radius\ F \geq r \implies fls_conv_radius\ G \geq r \implies fls_conv_radius\ (F - G) \geq r$
using *fls_conv_radius_diff[of F G]* **by** (*simp add: min_def split: if_splits*)

lemma *fls_conv_radius_mult_ge*:
 $fls_conv_radius\ F \geq r \implies fls_conv_radius\ G \geq r \implies fls_conv_radius\ (F * G) \geq r$
using *fls_conv_radius_mult[of F G]* **by** (*simp add: min_def split: if_splits*)

lemma *fls_conv_radius_power*: $fls_conv_radius\ (F \wedge n) \geq fls_conv_radius\ F$
by (*induction n*) (*auto intro!: fls_conv_radius_mult_ge*)

lemma *eval_fls_0* [*simp*]: $eval_fls\ 0\ z = 0$
and *eval_fls_1* [*simp*]: $eval_fls\ 1\ z = 1$
and *eval_fls_const* [*simp*]: $eval_fls\ (fls_const\ c)\ z = c$
and *eval_fls_numeral* [*simp*]: $eval_fls\ (numeral\ num)\ z = numeral\ num$
and *eval_fls_of_nat* [*simp*]: $eval_fls\ (of_nat\ n)\ z = of_nat\ n$
and *eval_fls_of_int* [*simp*]: $eval_fls\ (of_int\ m)\ z = of_int\ m$
and *eval_fls_X* [*simp*]: $eval_fls\ fls_X\ z = z$
and *eval_fls_X_intpow* [*simp*]: $eval_fls\ (fls_X_intpow\ m)\ z = z\ powi\ m$
by (*simp_all add: eval_fls_def*)

lemma *eval_fls_at_0*: $eval_fls\ F\ 0 = (if\ fls_subdegree\ F \geq 0\ then\ fls_nth\ F\ 0\ else\ 0)$

by (cases fls_subdegree F = 0)
 (simp_all add: eval_fls_def fls_repart_def eval_fps_at_0)

lemma eval_fps_to_fls:

assumes norm z < fps_conv_radius F

shows eval_fls (fps_to_fls F) z = eval_fps F z

proof (cases F = 0)

case [simp]: False

have eval_fps F z = eval_fps (unit_factor F * normalize F) z

by (metis unit_factor_mult_normalize)

also have ... = eval_fps (unit_factor F * fps_X ^ subdegree F) z

by simp

also have ... = eval_fps (unit_factor F) z * z ^ subdegree F

using assms by (subst eval_fps_mult) auto

also have ... = eval_fls (fps_to_fls F) z

unfolding eval_fls_def fls_base_factor_to_fps_to_fls fls_subdegree_fls_to_fps
 power_int_of_nat ..

finally show ?thesis ..

qed auto

lemma eval_fls_shift:

assumes [simp]: z ≠ 0

shows eval_fls (fls_shift n F) z = eval_fls F z * z powi -n

proof (cases F = 0)

case [simp]: False

show ?thesis

unfolding eval_fls_def

by (subst fls_base_factor_to_fps_shift, subst fls_shift_subdegree[OF ‹F ≠ 0›],
 subst power_int_diff)

(auto simp: power_int_minus divide_simps)

qed auto

lemma eval_fls_add:

assumes ereal (norm z) < fls_conv_radius F ereal (norm z) < fls_conv_radius
 G z ≠ 0

shows eval_fls (F + G) z = eval_fls F z + eval_fls G z

using assms

proof (induction fls_subdegree F fls_subdegree G arbitrary: F G rule: linorder_wlog)

case (sym F G)

show ?case

using sym(1)[of G F] sym(2-) by (simp add: add_ac)

next

case (le F G)

show ?case

proof (cases F = 0 ∨ G = 0)

case False

hence [simp]: F ≠ 0 G ≠ 0

by auto

note [simp] = ‹z ≠ 0›

```

define  $F' G'$  where  $F' = \text{fls\_base\_factor\_to\_fps } F$   $G' = \text{fls\_base\_factor\_to\_fps } G$ 
define  $m n$  where  $m = \text{fls\_subdegree } F$   $n = \text{fls\_subdegree } G$ 
have  $m \leq n$ 
  using  $le$  by ( $auto$   $simp$ :  $m\_n\_def$ )
have  $\text{conv1}$ :  $\text{ereal } (\text{cmod } z) < \text{fps\_conv\_radius } F'$   $\text{ereal } (\text{cmod } z) < \text{fps\_conv\_radius } G'$ 
  using  $assms$   $le$  by ( $simp\_all$   $add$ :  $F'_G'_def$   $\text{fls\_conv\_radius\_altdef}$ )
have  $\text{conv2}$ :  $\text{ereal } (\text{cmod } z) < \text{fps\_conv\_radius } (G' * \text{fps\_X}^{\text{nat } (n - m)})$ 
  using  $\text{conv1}$  by ( $intro$   $less\_le\_trans$ [ $OF\_fps\_conv\_radius\_mult$ ])  $auto$ 
have  $\text{conv3}$ :  $\text{ereal } (\text{cmod } z) < \text{fps\_conv\_radius } (F' + G' * \text{fps\_X}^{\text{nat } (n - m)})$ 
  using  $\text{conv1}$   $\text{conv2}$  by ( $intro$   $less\_le\_trans$ [ $OF\_fps\_conv\_radius\_add$ ])  $auto$ 

have  $\text{eval\_fls } F z + \text{eval\_fls } G z = \text{eval\_fps } F' z * z^{\text{powi } m} + \text{eval\_fps } G' z * z^{\text{powi } n}$ 
  unfolding  $\text{eval\_fls\_def } m\_n\_def$ [ $symmetric$ ]  $F'_G'_def$ [ $symmetric$ ]
  by ( $simp$   $add$ :  $\text{power\_int\_add algebra\_simps}$ )
also have  $\dots = (\text{eval\_fps } F' z + \text{eval\_fps } G' z * z^{\text{powi } (n - m)}) * z^{\text{powi } m}$ 
  by ( $simp$   $add$ :  $\text{algebra\_simps power\_int\_diff}$ )
also have  $\text{eval\_fps } G' z * z^{\text{powi } (n - m)} = \text{eval\_fps } (G' * \text{fps\_X}^{\text{nat } (n - m)}) z$ 
  using  $assms$   $\langle m \leq n \rangle$   $\text{conv1}$  by ( $subst$   $\text{eval\_fps\_mult}$ ) ( $auto$   $simp$ :  $\text{power\_int\_def}$ )
also have  $\text{eval\_fps } F' z + \dots = \text{eval\_fps } (F' + G' * \text{fps\_X}^{\text{nat } (n - m)}) z$ 
  using  $\text{conv1}$   $\text{conv2}$  by ( $subst$   $\text{eval\_fps\_add}$ )  $auto$ 
also have  $\dots = \text{eval\_fls } (\text{fps\_to\_fls } (F' + G' * \text{fps\_X}^{\text{nat } (n - m)})) z$ 
  using  $\text{conv3}$  by ( $subst$   $\text{eval\_fps\_to\_fls}$ )  $auto$ 
also have  $\dots * z^{\text{powi } m} = \text{eval\_fls } (\text{fls\_shift } (-m) (\text{fps\_to\_fls } (F' + G' * \text{fps\_X}^{\text{nat } (n - m)}))) z$ 
  by ( $subst$   $\text{eval\_fls\_shift}$ )  $auto$ 
also have  $\text{fls\_shift } (-m) (\text{fps\_to\_fls } (F' + G' * \text{fps\_X}^{\text{nat } (n - m)})) = F + G$ 
  using  $\langle m \leq n \rangle$ 
  by ( $simp$   $add$ :  $\text{fls\_times\_fps\_to\_fls}$   $\text{fps\_to\_fls\_power}$   $\text{fls\_X\_power\_conv\_shift\_1}$   $\text{fls\_shifted\_times\_simps } F'_G'_def$   $m\_n\_def$ )
finally show  $?thesis$  ..
qed  $auto$ 
qed

```

lemma eval_fls_minus :

```

assumes  $\text{ereal } (\text{norm } z) < \text{fls\_conv\_radius } F$ 
shows  $\text{eval\_fls } (-F) z = -\text{eval\_fls } F z$ 
using  $assms$  by ( $simp$   $add$ :  $\text{eval\_fls\_def}$   $\text{eval\_fps\_minus}$   $\text{fls\_conv\_radius\_altdef}$ )

```

lemma eval_fls_diff :

```

assumes  $\text{ereal } (\text{norm } z) < \text{fls\_conv\_radius } F$   $\text{ereal } (\text{norm } z) < \text{fls\_conv\_radius } G$ 
and [ $simp$ ]:  $z \neq 0$ 

```

```

shows  $eval\_fls (F - G) z = eval\_fls F z - eval\_fls G z$ 
proof -
  have  $eval\_fls (F + (-G)) z = eval\_fls F z - eval\_fls G z$ 
    using assms by (subst eval_fls_add) (auto simp: eval_fls_minus)
  thus ?thesis
    by simp
qed

```

```

lemma eval_fls_mult:
  assumes  $ereal (norm z) < fls\_conv\_radius F$   $ereal (norm z) < fls\_conv\_radius G$ 
   $z \neq 0$ 
  shows  $eval\_fls (F * G) z = eval\_fls F z * eval\_fls G z$ 
proof (cases  $F = 0 \vee G = 0$ )
  case False
  hence [simp]:  $F \neq 0 \ G \neq 0$ 
    by auto
  note [simp] =  $\langle z \neq 0 \rangle$ 
  define  $F' G'$  where  $F' = fls\_base\_factor\_to\_fps F$   $G' = fls\_base\_factor\_to\_fps G$ 
  define  $m n$  where  $m = fls\_subdegree F$   $n = fls\_subdegree G$ 
  have  $eval\_fls F z * eval\_fls G z = (eval\_fps F' z * eval\_fps G' z) * z^{powi (m + n)}$ 
  unfolding eval_fls_def m_n_def [symmetric] F'_G'_def [symmetric]
  by (simp add: power_int_add algebra_simps)
  also have  $\dots = eval\_fps (F' * G') z * z^{powi (m + n)}$ 
    using assms by (subst eval_fps_mult) (auto simp: F'_G'_def fls_conv_radius_altdef)
  also have  $\dots = eval\_fls (F * G) z$ 
    by (simp add: eval_fls_def F'_G'_def m_n_def) (simp add: fls_times_def)
  finally show ?thesis ..
qed auto

```

```

lemma eval_fls_power:
  assumes  $ereal (norm z) < fls\_conv\_radius F$   $z \neq 0$ 
  shows  $eval\_fls (F ^ n) z = eval\_fls F z ^ n$ 
proof (induction n)
  case (Suc n)
  have  $eval\_fls (F ^ Suc n) z = eval\_fls (F * F ^ n) z$ 
    by simp
  also have  $\dots = eval\_fls F z * eval\_fls (F ^ n) z$ 
    using assms by (subst eval_fls_mult) (auto intro!: less_le_trans[OF fls_conv_radius_power])
  finally show ?case
    using Suc by simp
qed auto

```

```

lemma norm_summable_fls:
   $norm z < fls\_conv\_radius f \implies summable (\lambda n. norm (fls_nth f n * z ^ n))$ 
  using norm_summable_fps [of z fls_regpart f] by (simp add: fls_conv_radius_def)

```

```

lemma norm_summable_fls':

```

```

  norm z < fls_conv_radius f  $\implies$  summable ( $\lambda n.$  norm (fls_nth f (n + fls_subdegree
f) * z ^ n))
  using norm_summable_fps[of z fls_base_factor_to_fps f] by (simp add: fls_conv_radius_altdef)

```

lemma summable_fls:

```

  norm z < fls_conv_radius f  $\implies$  summable ( $\lambda n.$  fls_nth f n * z ^ n)
  by (rule summable_norm_cancel[OF norm_summable_fls])

```

theorem sums_eval_fls:

```

  fixes f
  defines n  $\equiv$  fls_subdegree f
  assumes norm z < fls_conv_radius f and z  $\neq$  0  $\vee$  n  $\geq$  0
  shows ( $\lambda k.$  fls_nth f (int k + n) * z powi (int k + n)) sums eval_fls f z
  proof (cases z = 0)
  case [simp]: False
  have ( $\lambda k.$  fps_nth (fls_base_factor_to_fps f) k * z ^ k * z powi n) sums
    (eval_fps (fls_base_factor_to_fps f) z * z powi n)
  using assms(2) by (intro sums_eval_fps sums_mult2) (auto simp: fls_conv_radius_altdef)
  thus ?thesis
    by (simp add: power_int_add n_def eval_fls_def mult_ac)
  next
  case [simp]: True
  with assms have n  $\geq$  0
  by auto
  have ( $\lambda k.$  fls_nth f (int k + n) * z powi (int k + n)) sums
    ( $\sum k \in (\text{if } n \leq 0 \text{ then } \{\text{nat } (-n)\} \text{ else } \{\}) . \text{fls\_nth } f \text{ (int } k + n) * z \text{ powi}$ 
(int k + n))
  by (intro sums_finite) (auto split: if_splits)
  also have ... = eval_fls f z
  using <n  $\geq$  0> by (auto simp: eval_fls_at_0 n_def not_le)
  finally show ?thesis .
  qed

```

lemma holomorphic_on_eval_fls:

```

  fixes f
  defines n  $\equiv$  fls_subdegree f
  assumes A  $\subseteq$  eball 0 (fls_conv_radius f) - (if n  $\geq$  0 then {} else {0})
  shows eval_fls f holomorphic_on A
  proof (cases n  $\geq$  0)
  case True
  have eval_fls f = ( $\lambda z.$  eval_fps (fls_base_factor_to_fps f) z * z ^ nat n)
  using True by (simp add: fun_eq_iff eval_fls_def power_int_def n_def)
  moreover have ... holomorphic_on A
  using True assms(2) by (intro holomorphic_intros) (auto simp: fls_conv_radius_altdef)
  ultimately show ?thesis
    by simp
  next
  case False
  show ?thesis using assms

```

unfolding *eval_fls_def* **by** (*intro holomorphic_intros*) (*auto simp: fls_conv_radius_altdef*)
qed

lemma *holomorphic_on_eval_fls'* [*holomorphic_intros*]:
assumes *g holomorphic_on A*
assumes *g ' A ⊆ eball 0 (fls_conv_radius f) – (if fls_subdegree f ≥ 0 then {} else {0})*
shows $(\lambda x. \text{eval_fls } f (g x)) \text{ holomorphic_on } A$
proof –
have *eval_fls f ∘ g holomorphic_on A*
by (*intro holomorphic_on_compose[OF assms(1) holomorphic_on_eval_fls]*)
(*use assms in auto*)
thus *?thesis*
by (*simp add: o_def*)
qed

lemma *continuous_on_eval_fls*:
fixes *f*
defines $n \equiv \text{fls_subdegree } f$
assumes *A ⊆ eball 0 (fls_conv_radius f) – (if n ≥ 0 then {} else {0})*
shows *continuous_on A (eval_fls f)*
by (*intro holomorphic_on_imp_continuous_on holomorphic_on_eval_fls*)
(*use assms in auto*)

lemma *continuous_on_eval_fls'* [*continuous_intros*]:
fixes *f*
defines $n \equiv \text{fls_subdegree } f$
assumes *g ' A ⊆ eball 0 (fls_conv_radius f) – (if n ≥ 0 then {} else {0})*
assumes *continuous_on A g*
shows *continuous_on A (λx. eval_fls f (g x))*
using *assms(3)*
by (*intro continuous_on_compose2[OF continuous_on_eval_fls __ assms(2)]*)
(*auto simp: n_def*)

lemmas *has_field_derivative_eval_fps'* [*derivative_intros*] =
DERIV_chain2[OF has_field_derivative_eval_fps]

lemma *fps_deriv_fls_regpart*: $\text{fps_deriv } (\text{fls_regpart } F) = \text{fls_regpart } (\text{fls_deriv } F)$
by (*intro fps_ext*) (*auto simp: add_ac*)

lemma *has_field_derivative_eval_fls*:
assumes $z \in \text{eball } 0 (\text{fls_conv_radius } f) - \{0\}$
shows $(\text{eval_fls } f \text{ has_field_derivative } \text{eval_fls } (\text{fls_deriv } f) z) \text{ (at } z \text{ within } A)$
proof –
define *g* **where** $g = \text{fls_base_factor_to_fps } f$
define *n* **where** $n = \text{fls_subdegree } f$
have [*simp*]: $\text{fps_conv_radius } g = \text{fls_conv_radius } f$


```

  by (simp add: fls_conv_radius_altdef g_def)
  have conv1: fps_conv_radius (fps_deriv g * fps_X) ≥ fls_conv_radius f
  by (intro fps_conv_radius_mult_ge order.trans[OF_fps_conv_radius_deriv])
auto
  have conv2: fps_conv_radius (of_int n * g) ≥ fls_conv_radius f
  by (intro fps_conv_radius_mult_ge) auto
  have conv3: fps_conv_radius (fps_deriv g * fps_X + of_int n * g) ≥ fls_conv_radius
f
  by (intro fps_conv_radius_add_ge conv1 conv2)

  have [simp]: fps_conv_radius g = fls_conv_radius f
  by (simp add: g_def fls_conv_radius_altdef)
  have ((λz. eval_fps g z * z powi fls_subdegree f) has_field_derivative
(eval_fps (fps_deriv g) z * z powi n + of_int n * z powi (n - 1) * eval_fps
g z))
  (at z within A)
  using assms by (auto intro!: derivative_eq_intros simp: n_def)
  also have (λz. eval_fps g z * z powi fls_subdegree f) = eval_fls f
  by (simp add: eval_fls_def g_def fun_eq_iff)
  also have eval_fps (fps_deriv g) z * z powi n + of_int n * z powi (n - 1) *
eval_fps g z =
  (z * eval_fps (fps_deriv g) z + of_int n * eval_fps g z) * z powi (n -
1)
  using assms by (auto simp: power_int_diff field_simps)
  also have (z * eval_fps (fps_deriv g) z + of_int n * eval_fps g z) =
  eval_fps (fps_deriv g * fps_X + of_int n * g) z
  using conv1 conv2 assms fps_conv_radius_deriv[of g]
  by (subst eval_fps_add) (auto simp: eval_fps_mult)
  also have ... = eval_fls (fps_to_fls (fps_deriv g * fps_X + of_int n * g)) z
  using conv3 assms by (subst eval_fps_to_fls) auto
  also have ... * z powi (n - 1) = eval_fls (fls_shift (1 - n) (fps_to_fls
(fps_deriv g * fps_X + of_int n * g))) z
  using assms by (subst eval_fls_shift) auto
  also have fls_shift (1 - n) (fps_to_fls (fps_deriv g * fps_X + of_int n * g))
= fls_deriv f
  by (intro fls_eqI) (auto simp: g_def n_def algebra_simps eq_commute[of _
fls_subdegree f])
  finally show ?thesis .
qed

```

lemma *eval_fls_deriv*:

```

  assumes z ∈ eball 0 (fls_conv_radius F) - {0}
  shows eval_fls (fls_deriv F) z = deriv (eval_fls F) z
  by (rule sym, rule DERIV_imp_deriv, rule has_field_derivative_eval_fls, rule
assms)

```

lemma *analytic_on_eval_fls*:

```

  assumes A ⊆ eball 0 (fls_conv_radius f) - (if fls_subdegree f ≥ 0 then {} else
{0})

```

```

  shows eval_fls f analytic_on A
proof (rule analytic_on_subset [OF _ assms])
  show eval_fls f analytic_on eball 0 (fls_conv_radius f) - (if fls_subdegree f ≥
0 then {} else {0})
  using holomorphic_on_eval_fls[OF order.refl]
  by (subst analytic_on_open) auto
qed

```

```

lemma analytic_on_eval_fls' [analytic_intros]:
  assumes g analytic_on A
  assumes g ' A ⊆ eball 0 (fls_conv_radius f) - (if fls_subdegree f ≥ 0 then {}
else {0})
  shows (λx. eval_fls f (g x)) analytic_on A
proof -
  have eval_fls f ∘ g analytic_on A
  by (intro analytic_on_compose[OF assms(1) analytic_on_eval_fls]) (use assms
in auto)
  thus ?thesis
  by (simp add: o_def)
qed

```

```

lemma continuous_eval_fls [continuous_intros]:
  assumes z ∈ eball 0 (fls_conv_radius F) - (if fls_subdegree F ≥ 0 then {} else
{0})
  shows continuous (at z within A) (eval_fls F)
proof -
  have isCont (eval_fls F) z
  using continuous_on_eval_fls[OF order.refl] assms
  by (subst (asm) continuous_on_eq_continuous_at) auto
  thus ?thesis
  using continuous_at_imp_continuous_at_within by blast
qed

```

named_theorems laurent_expansion_intros

```

lemma has_laurent_expansion_imp_asymp_equiv_0:
  assumes F: f has_laurent_expansion F
  defines n ≡ fls_subdegree F
  shows f ~[at 0] (λz. fls_nth F n * z powi n)
proof (cases F = 0)
  case True
  thus ?thesis using assms
  by (auto simp: has_laurent_expansion_def)
next
  case [simp]: False
  define G where G = fls_base_factor_to_fps F

```

```

have fls_conv_radius F > 0
  using F by (auto simp: has_laurent_expansion_def)
hence isCont (eval_fps G) 0
  by (intro continuous_intros) (auto simp: G_def fps_conv_radius_fls_regpart
zero_ereal_def)
hence lim: eval_fps G -0 → eval_fps G 0
  by (meson isContD)
have [simp]: fps_nth G 0 ≠ 0
  by (auto simp: G_def)

have f ~[at 0] eval_fls F
  using F by (intro asymp_equiv_refl_ev) (auto simp: has_laurent_expansion_def
eq_commute)
also have ... = (λz. eval_fps G z * z powi n)
  by (intro ext) (simp_all add: eval_fls_def G_def n_def)
also have ... ~[at 0] (λz. fps_nth G 0 * z powi n) using lim
  by (intro asymp_equiv_intros tendsto_imp_asymp_equiv_const) (auto simp:
eval_fps_at_0)
also have fps_nth G 0 = fls_nth F n
  by (simp add: G_def n_def)
finally show ?thesis
  by simp
qed

```

```

lemma has_laurent_expansion_imp_asymp_equiv:
  assumes F: (λw. f (z + w)) has_laurent_expansion F
  defines n ≡ fls_subdegree F
  shows f ~[at z] (λw. fls_nth F n * (w - z) powi n)
  using has_laurent_expansion_imp_asymp_equiv_0[OF assms(1)] unfolding
n_def
  by (simp add: at_to_0[of z] asymp_equiv_filtermap_iff add_ac)

```

lemmas [tendsto_intros del] = tendsto_power_int

```

lemma has_laurent_expansion_imp_tendsto_0:
  assumes F: f has_laurent_expansion F and fls_subdegree F ≥ 0
  shows f -0 → fls_nth F 0
proof (rule asymp_equiv_tendsto_transfer)
  show (λz. fls_nth F (fls_subdegree F) * z powi fls_subdegree F) ~[at 0] f
    by (rule asymp_equiv_symI, rule has_laurent_expansion_imp_asymp_equiv_0)
fact
  show (λz. fls_nth F (fls_subdegree F) * z powi fls_subdegree F) -0 → fls_nth
F 0
  by (rule tendsto_eq_intros refl | use assms(2) in simp)+
  (use assms(2) in <auto simp: power_int_0_left>If>)
qed

```

lemma has_laurent_expansion_imp_tendsto:

assumes $F: (\lambda w. f (z + w))$ *has_laurent_expansion* F **and** *fls_subdegree* $F \geq 0$

shows $f \xrightarrow{-z} \text{fls_nth } F \ 0$

using *has_laurent_expansion_imp_tendsto_0*[OF *assms*]

by (*simp* *add: at_to_0*[*of z*] *filterlim_filtermap_add_ac*)

lemma *has_laurent_expansion_imp_filterlim_infinity_0*:

assumes $F: f$ *has_laurent_expansion* F **and** *fls_subdegree* $F < 0$

shows *filterlim* f *at_infinity* (*at 0*)

proof (*rule asymp_equiv_at_infinity_transfer*)

have [*simp*]: $F \neq 0$

using *assms*(2) **by** *auto*

show $(\lambda z. \text{fls_nth } F (\text{fls_subdegree } F) * z \text{ powi } \text{fls_subdegree } F) \sim[\text{at } 0] f$

by (*rule asymp_equiv_symI*, *rule has_laurent_expansion_imp_asymp_equiv_0*)

fact

show *filterlim* $(\lambda z. \text{fls_nth } F (\text{fls_subdegree } F) * z \text{ powi } \text{fls_subdegree } F)$ *at_infinity* (*at 0*)

by (*rule tendsto_mult_filterlim_at_infinity_tendsto_const*

filterlim_power_int_neg_at_infinity | *use assms*(2) **in** *simp*) +

(*auto simp: eventually_at_filter*)

qed

lemma *has_laurent_expansion_imp_neg_fls_subdegree*:

assumes $F: f$ *has_laurent_expansion* F

and *infy:filterlim* f *at_infinity* (*at 0*)

shows *fls_subdegree* $F < 0$

proof (*rule ccontr*)

assume *asm*: $\neg \text{fls_subdegree } F < 0$

define *ff* **where** $\text{ff} = (\lambda z. \text{fls_nth } F (\text{fls_subdegree } F) * z \text{ powi } \text{fls_subdegree } F)$

have (*ff* \longrightarrow (*if fls_subdegree* $F = 0$ *then fls_nth* $F \ 0$ *else 0*)) (*at 0*)

using *asm* **unfolding** *ff_def*

by (*auto intro!: tendsto_eq_intros*)

moreover **have** *filterlim* *ff* *at_infinity* (*at 0*)

proof (*rule asymp_equiv_at_infinity_transfer*)

show $f \sim[\text{at } 0] \text{ff}$ **unfolding** *ff_def*

using *has_laurent_expansion_imp_asymp_equiv_0*[OF F] **unfolding** *ff_def*

.

show *filterlim* f *at_infinity* (*at 0*) **by** *fact*

qed

ultimately **show** *False*

using *not_tendsto_and_filterlim_at_infinity*[*of at (0::complex)*] **by** *auto*

qed

lemma *has_laurent_expansion_imp_filterlim_infinity*:

assumes $F: (\lambda w. f (z + w))$ *has_laurent_expansion* F **and** *fls_subdegree* $F < 0$

shows *filterlim* f *at_infinity* (*at z*)

using *has_laurent_expansion_imp_filterlim_infinity_0* [*OF assms*]
by (*simp add: at_to_0* [*of z*] *filterlim_filtermap add_ac*)

lemma *has_laurent_expansion_imp_is_pole_0*:
assumes *F*: *f* *has_laurent_expansion F* **and** *fls_subdegree F < 0*
shows *is_pole f 0*
using *has_laurent_expansion_imp_filterlim_infinity_0* [*OF assms*]
by (*simp add: is_pole_def*)

lemma *is_pole_0_imp_neg_fls_subdegree*:
assumes *F*: *f* *has_laurent_expansion F* **and** *is_pole f 0*
shows *fls_subdegree F < 0*
using *F assms(2)* *has_laurent_expansion_imp_neg_fls_subdegree is_pole_def*
by blast

lemma *has_laurent_expansion_imp_is_pole*:
assumes *F*: $(\lambda x. f(z + x))$ *has_laurent_expansion F* **and** *fls_subdegree F < 0*
shows *is_pole f z*
using *has_laurent_expansion_imp_is_pole_0* [*OF assms*]
by (*simp add: is_pole_shift_0'*)

lemma *is_pole_imp_neg_fls_subdegree*:
assumes *F*: $(\lambda x. f(z + x))$ *has_laurent_expansion F* **and** *is_pole f z*
shows *fls_subdegree F < 0*
apply (*rule is_pole_0_imp_neg_fls_subdegree* [*OF F*])
using *assms(2)* *is_pole_shift_0* **by blast**

lemma *is_pole_fls_subdegree_iff*:
assumes $(\lambda x. f(z + x))$ *has_laurent_expansion F*
shows *is_pole f z* \longleftrightarrow *fls_subdegree F < 0*
using *assms is_pole_imp_neg_fls_subdegree has_laurent_expansion_imp_is_pole*
by auto

lemma
assumes *f* *has_laurent_expansion F*
shows *has_laurent_expansion_isolated_0: isolated_singularity_at f 0*
and *has_laurent_expansion_not_essential_0: not_essential f 0*

proof –

from *assms* **have** *eventually* $(\lambda z. \text{eval_fls } F z = f z)$ (*at 0*)

by (*auto simp: has_laurent_expansion_def*)

then obtain *r* **where** *r*: $r > 0 \wedge z. z \in \text{ball } 0 r - \{0\} \implies \text{eval_fls } F z = f z$

by (*auto simp: eventually_at_filter ball_def eventually_nhds_metric*)

have *fls_conv_radius F > 0*

using *assms* **by** (*auto simp: has_laurent_expansion_def*)

then obtain *R* :: *real* **where** *R*: $R > 0 \wedge R \leq \min r (\text{fls_conv_radius } F)$

using $\langle r > 0 \rangle$ **by** (*metis dual_order.strict_implies_order ereal_dense2 ereal_less(2) min_def*)

```

have eval_fls F holomorphic_on ball 0 R - {0}
  using r R by (intro holomorphic_intros ball_eball_mono Diff_mono) (auto
simp: ereal_le_less)
also have ?this  $\longleftrightarrow$  f holomorphic_on ball 0 R - {0}
  using r R by (intro holomorphic_cong) auto
also have ...  $\longleftrightarrow$  f analytic_on ball 0 R - {0}
  by (subst analytic_on_open) auto
finally show isolated_singularity_at f 0
  unfolding isolated_singularity_at_def using ‹R > 0› by blast

show not_essential f 0
proof (cases fls_subdegree F  $\geq$  0)
  case True
  hence f  $-0 \rightarrow$  fls_nth F 0
    by (intro has_laurent_expansion_imp_tendsto_0[OF assms])
  thus ?thesis
    by (auto simp: not_essential_def)
  next
  case False
  hence is_pole f 0
    by (intro has_laurent_expansion_imp_is_pole_0[OF assms]) auto
  thus ?thesis
    by (auto simp: not_essential_def)
qed
qed

```

lemma

```

assumes ( $\lambda w. f(z + w)$ ) has_laurent_expansion F
shows has_laurent_expansion_isolated: isolated_singularity_at f z
  and has_laurent_expansion_not_essential: not_essential f z
using has_laurent_expansion_isolated_0[OF assms] has_laurent_expansion_not_essential_0[OF
assms]
by (simp_all add: isolated_singularity_at_shift_0 not_essential_shift_0)

```

lemma has_laurent_expansion_fps:

```

assumes f has_fps_expansion F
shows f has_laurent_expansion fps_to_fls F
proof -
  from assms have radius: 0 < fps_conv_radius F and eval:  $\forall_F z$  in nhds 0.
eval_fps F z = f z
  by (auto simp: has_fps_expansion_def)
  from eval have eval':  $\forall_F z$  in at 0. eval_fps F z = f z
  using eventually_at_filter eventually_mono by fastforce
  moreover have eventually ( $\lambda z. z \in$  eball 0 (fps_conv_radius F) - {0}) (at 0)
  using radius by (intro eventually_at_in_open) (auto simp: zero_ereal_def)
  ultimately have eventually ( $\lambda z. eval\_fls$  (fps_to_fls F) z = f z) (at 0)
  by eventually_elim (auto simp: eval_fps_to_fls)
  thus ?thesis using radius
  by (auto simp: has_laurent_expansion_def)

```

qed

lemma *has_laurent_expansion_const* [*simp, intro, laurent_expansion_intros*]:
 ($\lambda_. c$) *has_laurent_expansion_fls_const* c
by (*auto simp: has_laurent_expansion_def*)

lemma *has_laurent_expansion_0* [*simp, intro, laurent_expansion_intros*]:
 ($\lambda_. 0$) *has_laurent_expansion* 0
by (*auto simp: has_laurent_expansion_def*)

lemma *has_fps_expansion_0_iff*: f *has_fps_expansion* $0 \iff$ *eventually* ($\lambda z. f$
 $z = 0$) (*nhds* 0)
by (*auto simp: has_fps_expansion_def*)

lemma *has_laurent_expansion_1* [*simp, intro, laurent_expansion_intros*]:
 ($\lambda_. 1$) *has_laurent_expansion* 1
by (*auto simp: has_laurent_expansion_def*)

lemma *has_laurent_expansion_numeral* [*simp, intro, laurent_expansion_intros*]:
 ($\lambda_. \text{numeral } n$) *has_laurent_expansion* *numeral* n
by (*auto simp: has_laurent_expansion_def*)

lemma *has_laurent_expansion_fps_X_power* [*laurent_expansion_intros*]:
 ($\lambda x. x \wedge n$) *has_laurent_expansion* (*fls_X_intpow* n)
by (*auto simp: has_laurent_expansion_def*)

lemma *has_laurent_expansion_fps_X_power_int* [*laurent_expansion_intros*]:
 ($\lambda x. x \text{ powi } n$) *has_laurent_expansion* (*fls_X_intpow* n)
by (*auto simp: has_laurent_expansion_def*)

lemma *has_laurent_expansion_fps_X* [*laurent_expansion_intros*]:
 ($\lambda x. x$) *has_laurent_expansion_fls_X*
by (*auto simp: has_laurent_expansion_def*)

lemma *has_laurent_expansion_cmult_left* [*laurent_expansion_intros*]:
assumes f *has_laurent_expansion* F

shows ($\lambda x. c * f x$) *has_laurent_expansion_fls_const* $c * F$

proof –

from *assms* **have** *eventually* ($\lambda z. z \in \text{eball } 0 (\text{fls_conv_radius } F) - \{0\}$) (*at* 0)

by (*intro eventually_at_in_open*) (*auto simp: has_laurent_expansion_def*
zero_ereal_def)

moreover from *assms* **have** *eventually* ($\lambda z. \text{eval_fls } F z = f z$) (*at* 0)

by (*auto simp: has_laurent_expansion_def*)

ultimately have *eventually* ($\lambda z. \text{eval_fls } (\text{fls_const } c * F) z = c * f z$) (*at* 0)

by *eventually_elim* (*simp_all* *add: eval_fls_mult*)

with *assms* **show** *?thesis*

by (*auto simp: has_laurent_expansion_def intro!: less_le_trans[OF fls_conv_radius_mult]*)

qed

lemma *has_laurent_expansion_cmult_right* [*laurent_expansion_intros*]:
assumes *f has_laurent_expansion F*
shows $(\lambda x. f x * c)$ *has_laurent_expansion F * fls_const c*
proof –
have $F * fls_const\ c = fls_const\ c * F$
by (*intro fls_eqI*) (*auto simp: mult.commute*)
with *has_laurent_expansion_cmult_left* [*OF assms*] **show** *?thesis*
by (*simp add: mult.commute*)
qed

lemma *has_fps_expansion_scaleR* [*fps_expansion_intros*]:
fixes $F :: 'a :: \{\text{banach, real_normed_div_algebra, comm_ring_1}\}$ *fps*
shows f *has_fps_expansion F* $\implies (\lambda x. c *_{\mathbb{R}} f x)$ *has_fps_expansion fps_const*
*(of_real c) * F*
unfolding *scaleR_conv_of_real* **by** (*intro fps_expansion_intros*)

lemma *has_laurent_expansion_scaleR* [*laurent_expansion_intros*]:
 f *has_laurent_expansion F* $\implies (\lambda x. c *_{\mathbb{R}} f x)$ *has_laurent_expansion fls_const*
*(of_real c) * F*
unfolding *scaleR_conv_of_real* **by** (*intro laurent_expansion_intros*)

lemma *has_laurent_expansion_minus* [*laurent_expansion_intros*]:
assumes f *has_laurent_expansion F*
shows $(\lambda x. - f x)$ *has_laurent_expansion -F*
proof –
from *assms* **have** *eventually* $(\lambda x. x \in \text{eball } 0 (\text{fls_conv_radius } F) - \{0\})$ *(at 0)*
by (*intro eventually_at_in_open*) (*auto simp: has_laurent_expansion_def zero_ereal_def*)
moreover from *assms* **have** *eventually* $(\lambda x. \text{eval_fls } F\ x = f\ x)$ *(at 0)*
by (*auto simp: has_laurent_expansion_def*)
ultimately have *eventually* $(\lambda x. \text{eval_fls } (-F)\ x = -f\ x)$ *(at 0)*
by *eventually_elim* (*auto simp: eval_fls_minus*)
thus *?thesis using assms by* (*auto simp: has_laurent_expansion_def*)
qed

lemma *has_laurent_expansion_add* [*laurent_expansion_intros*]:
assumes f *has_laurent_expansion F* g *has_laurent_expansion G*
shows $(\lambda x. f x + g x)$ *has_laurent_expansion F + G*
proof –
from *assms* **have** $0 < \min (\text{fls_conv_radius } F) (\text{fls_conv_radius } G)$
by (*auto simp: has_laurent_expansion_def*)
also have $\dots \leq \text{fls_conv_radius } (F + G)$
by (*rule fls_conv_radius_add*)
finally have *radius: ... > 0* .

from *assms* **have** *eventually* $(\lambda x. x \in \text{eball } 0 (\text{fls_conv_radius } F) - \{0\})$ *(at 0)*
eventually $(\lambda x. x \in \text{eball } 0 (\text{fls_conv_radius } G) - \{0\})$ *(at 0)*

by (*intro eventually_at_in_open*; *force simp: has_laurent_expansion_def zero_ereal_def*) +
moreover have *eventually* ($\lambda x. \text{eval_fls } F \ x = f \ x$) (*at 0*)
and *eventually* ($\lambda x. \text{eval_fls } G \ x = g \ x$) (*at 0*)
using *assms* **by** (*auto simp: has_laurent_expansion_def*)
ultimately have *eventually* ($\lambda x. \text{eval_fls } (F + G) \ x = f \ x + g \ x$) (*at 0*)
by *eventually_elim* (*auto simp: eval_fls_add*)
with radius show *?thesis* **by** (*auto simp: has_laurent_expansion_def*)
qed

lemma *has_laurent_expansion_diff* [*laurent_expansion_intros*]:
assumes *f has_laurent_expansion F g has_laurent_expansion G*
shows ($\lambda x. f \ x - g \ x$) *has_laurent_expansion F - G*
using *has_laurent_expansion_add*[*of f F \lambda x. - g x - G*] *assms*
by (*simp add: has_laurent_expansion_minus*)

lemma *has_laurent_expansion_mult* [*laurent_expansion_intros*]:
assumes *f has_laurent_expansion F g has_laurent_expansion G*
shows ($\lambda x. f \ x * g \ x$) *has_laurent_expansion F * G*
proof –
from *assms* **have** $0 < \min(\text{fls_conv_radius } F) (\text{fls_conv_radius } G)$
by (*auto simp: has_laurent_expansion_def*)
also have $\dots \leq \text{fls_conv_radius } (F * G)$
by (*rule fls_conv_radius_mult*)
finally have *radius: \dots > 0* .

from *assms* **have** *eventually* ($\lambda x. x \in \text{eball } 0 (\text{fls_conv_radius } F) - \{0\}$) (*at 0*)
and *eventually* ($\lambda x. x \in \text{eball } 0 (\text{fls_conv_radius } G) - \{0\}$) (*at 0*)
by (*intro eventually_at_in_open*; *force simp: has_laurent_expansion_def zero_ereal_def*) +
moreover have *eventually* ($\lambda x. \text{eval_fls } F \ x = f \ x$) (*at 0*)
and *eventually* ($\lambda x. \text{eval_fls } G \ x = g \ x$) (*at 0*)
using *assms* **by** (*auto simp: has_laurent_expansion_def*)
ultimately have *eventually* ($\lambda x. \text{eval_fls } (F * G) \ x = f \ x * g \ x$) (*at 0*)
by *eventually_elim* (*auto simp: eval_fls_mult*)
with radius show *?thesis* **by** (*auto simp: has_laurent_expansion_def*)
qed

lemma *has_fps_expansion_power* [*fps_expansion_intros*]:
fixes $F :: 'a :: \{\text{banach, real_normed_div_algebra, comm_ring_1}\}$ *fps*
shows *f has_fps_expansion F* $\implies (\lambda x. f \ x \wedge m)$ *has_fps_expansion F ^ m*
by (*induction m*) (*auto intro!: fps_expansion_intros*)

lemma *has_laurent_expansion_power* [*laurent_expansion_intros*]:
assumes *f has_laurent_expansion F*
shows ($\lambda x. f \ x \wedge n$) *has_laurent_expansion F ^ n*
by (*induction n*) (*auto intro!: laurent_expansion_intros assms*)

lemma *has_laurent_expansion_sum* [*laurent_expansion_intros*]:
assumes $\bigwedge x. x \in I \implies f \ x$ *has_laurent_expansion F x*

shows $(\lambda y. \sum x \in I. f x y)$ *has_laurent_expansion* $(\sum x \in I. F x)$
using *assms* **by** (*induction I rule: infinite_finite_induct*) (*auto intro!: laurent_expansion_intros*)

lemma *has_laurent_expansion_prod* [*laurent_expansion_intros*]:
assumes $\bigwedge x. x \in I \implies f x$ *has_laurent_expansion* $F x$
shows $(\lambda y. \prod x \in I. f x y)$ *has_laurent_expansion* $(\prod x \in I. F x)$
using *assms* **by** (*induction I rule: infinite_finite_induct*) (*auto intro!: laurent_expansion_intros*)

lemma *has_laurent_expansion_deriv* [*laurent_expansion_intros*]:
assumes f *has_laurent_expansion* F
shows *deriv f* *has_laurent_expansion* *fls_deriv F*
proof –
have *eventually* $(\lambda z. z \in \text{eball } 0 (\text{fls_conv_radius } F) - \{0\})$ (at 0)
using *assms* **by** (*intro eventually_at_in_open*)
(auto simp: has_laurent_expansion_def zero_ereal_def)
moreover from *assms* **have** *eventually* $(\lambda z. \text{eval_fls } F z = f z)$ (at 0)
by (*auto simp: has_laurent_expansion_def*)
then obtain s **where** *open s* $0 \in s$ **and** $s: \bigwedge w. w \in s - \{0\} \implies \text{eval_fls } F w = f w$
by (*auto simp: eventually_nhds eventually_at_filter*)
hence *eventually* $(\lambda w. w \in s - \{0\})$ (at 0)
by (*intro eventually_at_in_open*) *auto*
ultimately have *eventually* $(\lambda z. \text{eval_fls } (\text{fls_deriv } F) z = \text{deriv } f z)$ (at 0)
proof *eventually_elim*
case (*elim z*)
hence $\text{eval_fls } (\text{fls_deriv } F) z = \text{deriv } (\text{eval_fls } F) z$
by (*simp add: eval_fls_deriv*)
also have *eventually* $(\lambda w. w \in s - \{0\})$ (*nhds z*)
using *elim* **and** *open s* **by** (*intro eventually_nhds_in_open*) *auto*
hence *eventually* $(\lambda w. \text{eval_fls } F w = f w)$ (*nhds z*)
by *eventually_elim* (*use s in auto*)
hence $\text{deriv } (\text{eval_fls } F) z = \text{deriv } f z$
by (*intro deriv_cong_ev refl*)
finally show ?*case* .
qed
with *assms* **show** ?*thesis*
by (*auto simp: has_laurent_expansion_def intro!: less_le_trans[OF fls_conv_radius_deriv]*)
qed

lemma *has_laurent_expansion_shift* [*laurent_expansion_intros*]:
assumes f *has_laurent_expansion* F
shows $(\lambda x. f x * x^{\text{pow } n})$ *has_laurent_expansion* $(\text{fls_shift } (-n) F)$
proof –
have *eventually* $(\lambda x. x \in \text{eball } 0 (\text{fls_conv_radius } F) - \{0\})$ (at 0)
using *assms* **by** (*intro eventually_at_in_open*) (*auto simp: has_laurent_expansion_def zero_ereal_def*)
moreover have *eventually* $(\lambda x. \text{eval_fls } F x = f x)$ (at 0)
using *assms* **by** (*auto simp: has_laurent_expansion_def*)
ultimately have *eventually* $(\lambda x. \text{eval_fls } (\text{fls_shift } (-n) F) x = f x * x^{\text{pow } n})$

n) (at 0)
 by eventually_elim (auto simp: eval_fls_shift assms)
 with assms show ?thesis by (auto simp: has_laurent_expansion_def)
 qed

lemma has_laurent_expansion_shift' [laurent_expansion_intros]:
 assumes f has_laurent_expansion F
 shows $(\lambda x. f x * x \text{ powi } (-n))$ has_laurent_expansion (fls_shift $n F$)
 using has_laurent_expansion_shift[OF assms, of $-n$] by simp

lemma has_laurent_expansion_deriv':
 assumes f has_laurent_expansion F
 assumes open A $0 \in A \wedge x. x \in A - \{0\} \implies (f \text{ has_field_derivative } f' x)$ (at x)
 shows f' has_laurent_expansion fls_deriv F
 proof -
 have deriv f has_laurent_expansion fls_deriv F
 by (intro laurent_expansion_intros assms)
 also have ?this \longleftrightarrow ?thesis
 proof (intro has_laurent_expansion_cong_refl)
 have eventually $(\lambda z. z \in A - \{0\})$ (at 0)
 by (intro eventually_at_in_open assms)
 thus eventually $(\lambda z. \text{deriv } f z = f' z)$ (at 0)
 by eventually_elim (auto intro!: DERIV_imp_deriv assms)
 qed
 finally show ?thesis .
 qed

definition laurent_expansion :: (complex \Rightarrow complex) \Rightarrow complex \Rightarrow complex fls
 where

laurent_expansion $f z =$
 (if eventually $(\lambda z. f z = 0)$ (at z) then 0
 else fls_shift $(-zorder f z)$ (fps_to_fls (fps_expansion (zor_poly $f z$) z)))

lemma laurent_expansion_cong:
 assumes eventually $(\lambda w. f w = g w)$ (at z) $z = z'$
 shows laurent_expansion $f z =$ laurent_expansion $g z'$
 unfolding laurent_expansion_def
 using zor_poly_cong[OF assms(1,2)] zorder_cong[OF assms] assms
 by (intro if_cong_refl) (auto elim: eventually_elim2)

theorem not_essential_has_laurent_expansion_0:
 assumes isolated_singularity_at f 0 not_essential f 0
 shows f has_laurent_expansion laurent_expansion f 0
 proof (cases $\exists_F w$ in at 0. $f w \neq 0$)
 case False
 have $(\lambda_. 0)$ has_laurent_expansion 0
 by simp

```

also have ?this  $\longleftrightarrow$  f has_laurent_expansion 0
using False by (intro has_laurent_expansion_cong) (auto simp: frequently_def)
finally show ?thesis
using False by (simp add: laurent_expansion_def frequently_def)
next
case True
define n where n = zorder f 0
obtain r where r: zor_poly f 0 0  $\neq$  0 zor_poly f 0 holomorphic_on cball 0 r r
> 0
       $\forall w \in \text{cball } 0 \ r - \{0\}. f \ w = \text{zor\_poly } f \ 0 \ w * w^{\text{powi } n} \wedge$ 
       $\text{zor\_poly } f \ 0 \ w \neq 0$ 
using zorder_exist[OF assms True] unfolding n_def by auto
have holo: zor_poly f 0 holomorphic_on ball 0 r
by (rule holomorphic_on_subset[OF r(2)]) auto

define F where F = fps_expansion (zor_poly f 0) 0
have F: zor_poly f 0 has_fps_expansion F
unfolding F_def by (rule has_fps_expansion_fps_expansion[OF _ _ holo])
(use <r > 0> in auto)
have ( $\lambda z. \text{zor\_poly } f \ 0 \ z * z^{\text{powi } n}$ ) has_laurent_expansion fls_shift (-n)
(fps_to_fl F)
by (intro laurent_expansion_intros has_laurent_expansion_fps[OF F])
also have ?this  $\longleftrightarrow$  f has_laurent_expansion fls_shift (-n) (fps_to_fl F)
by (intro has_laurent_expansion_cong refl eventually_mono[OF eventually_at_in_open[of
ball 0 r]])
      (use r in <auto simp: complex_powr_of_int>)
finally show ?thesis using True
by (simp add: laurent_expansion_def F_def n_def frequently_def)
qed

lemma not_essential_has_laurent_expansion:
assumes isolated_singularity_at f z not_essential f z
shows ( $\lambda x. f (z + x)$ ) has_laurent_expansion laurent_expansion f z
proof -
from assms(1) have iso: isolated_singularity_at ( $\lambda x. f (z + x)$ ) 0
by (simp add: isolated_singularity_at_shift_0)
moreover from assms(2) have ness: not_essential ( $\lambda x. f (z + x)$ ) 0
by (simp add: not_essential_shift_0)
ultimately have ( $\lambda x. f (z + x)$ ) has_laurent_expansion laurent_expansion ( $\lambda x.$ 
f (z + x)) 0
by (rule not_essential_has_laurent_expansion_0)

also have ... = laurent_expansion f z
proof (cases  $\exists_F w$  in at z. f w  $\neq$  0)
case False
then have  $\forall_F w$  in at z. f w = 0 using not_frequently by force
then have laurent_expansion ( $\lambda x. f (z + x)$ ) 0 = 0
by (smt (verit, best) add commute eventually_at_to_0 eventually_mono
laurent_expansion_def)

```

```

moreover have laurent_expansion f z = 0
  using ⟨ $\forall_F w$  in at z. f w = 0⟩ unfolding laurent_expansion_def by auto
ultimately show ?thesis by auto
next
case True
define df where df=zor_poly ( $\lambda x. f (z + x)$ ) 0
define g where g=( $\lambda u. u - z$ )

have fps_expansion df 0
  = fps_expansion (df o g) z
proof -
  have  $\exists_F w$  in at 0. f (z + w)  $\neq$  0 using True
    by (smt (verit, best) add.commute eventually_at_to_0
      eventually_mono not_frequently)
  from zorder_exist[OF iso ness this, folded df_def]
  obtain r where r>0 and df_holo:df holomorphic_on cball 0 r and df 0  $\neq$ 
0
     $\forall w \in \text{cball } 0 \ r - \{0\}.$ 
    f (z + w) = df w * w powi (zorder ( $\lambda w. f (z + w)$ ) 0)  $\wedge$ 
    df w  $\neq$  0
  by auto
  then have df_nz: $\forall w \in \text{ball } 0 \ r. \text{df } w \neq 0$  by auto

  have (deriv  $\sim$  n) df 0 = (deriv  $\sim$  n) (df o g) z for n
    unfolding comp_def g_def
  proof (subst higher_deriv_compose_linear'[where u=1 and c=-z, simplified])
    show df holomorphic_on ball 0 r
      using df_holo by auto
    show open (ball z r) open (ball 0 r) z  $\in$  ball z r
      using ⟨r>0⟩ by auto
    show  $\bigwedge w. w \in \text{ball } z \ r \implies w - z \in \text{ball } 0 \ r$ 
      by (simp add: dist_norm)
  qed auto
  then show ?thesis
    unfolding fps_expansion_def by auto
qed
also have ... = fps_expansion (zor_poly f z) z
proof (rule fps_expansion_cong)
  have  $\forall_F w$  in nhds z. zor_poly f z w
    = zor_poly ( $\lambda u. f (z + u)$ ) 0 (w - z)
  apply (rule zor_poly_shift)
  using True assms by auto
  then show  $\forall_F w$  in nhds z. (df o g) w = zor_poly f z w
    unfolding df_def g_def comp_def
    by (auto elim:eventually_mono)
qed
finally show ?thesis unfolding df_def
  by (auto simp: laurent_expansion_def at_to_0[of z]
    eventually_filtermap_add_ac zorder_shift')

```

qed
finally show *?thesis* .
qed

lemma *has_fps_expansion_to_laurent*:
 $f \text{ has_fps_expansion } F \iff f \text{ has_laurent_expansion } \text{fps_to_fls } F \wedge f 0 = \text{fps_nth } F 0$
proof *safe*
assume *: $f \text{ has_laurent_expansion } \text{fps_to_fls } F f 0 = \text{fps_nth } F 0$
have *eventually* $(\lambda z. z \in \text{eball } 0 (\text{fps_conv_radius } F)) (\text{nhds } 0)$
using * **by** *(intro eventually_nhds_in_open) (auto simp: has_laurent_expansion_def zero_ereal_def)*
moreover have *eventually* $(\lambda z. z \neq 0 \longrightarrow \text{eval_fls } (\text{fps_to_fls } F) z = f z) (\text{nhds } 0)$
using * **by** *(auto simp: has_laurent_expansion_def eventually_at_filter)*
ultimately have *eventually* $(\lambda z. f z = \text{eval_fps } F z) (\text{nhds } 0)$
by *eventually_elim*
*(auto simp: has_laurent_expansion_def eventually_at_filter eval_fps_at_0 eval_fps_to_fls *(2))*
thus $f \text{ has_fps_expansion } F$
using * **by** *(auto simp: has_fps_expansion_def has_laurent_expansion_def eq_commute)*
next
assume $f \text{ has_fps_expansion } F$
thus $f 0 = \text{fps_nth } F 0$
by *(metis eval_fps_at_0 has_fps_expansion_imp_holomorphic)*
qed *(auto intro: has_laurent_expansion_fps)*

lemma *eval_fps_fls_base_factor* [*simp*]:
assumes $z \neq 0$
shows $\text{eval_fps } (\text{fls_base_factor_to_fps } F) z = \text{eval_fls } F z * z^{\text{powi } -\text{fls_subdegree } F}$
using *assms unfolding eval_fls_def* **by** *(simp add: power_int_minus field_simps)*

lemma *has_fps_expansion_imp_analytic_0*:
assumes $f \text{ has_fps_expansion } F$
shows $f \text{ analytic_on } \{0\}$
by *(meson analytic_at_two assms has_fps_expansion_imp_holomorphic)*

lemma *has_fps_expansion_imp_analytic*:
assumes $(\lambda x. f (z + x)) \text{ has_fps_expansion } F$
shows $f \text{ analytic_on } \{z\}$
proof –
have $(\lambda x. f (z + x)) \text{ analytic_on } \{0\}$
by *(rule has_fps_expansion_imp_analytic_0) fact*
hence $(\lambda x. f (z + x)) \circ (\lambda x. x - z) \text{ analytic_on } \{z\}$
by *(intro analytic_on_compose_gen analytic_intros) auto*
thus *?thesis*
by *(simp add: o_def)*

qed

lemma *is_pole_cong_asymp_equiv*:

assumes $f \sim[at\ z] g\ z = z'$

shows $is_pole\ f\ z = is_pole\ g\ z'$

using *asymp_equiv_at_infinity_transfer*[OF *assms*(1)]

asymp_equiv_at_infinity_transfer[OF *asymp_equiv_symI*[OF *assms*(1)]]

assms(2)

unfolding *is_pole_def* **by** *auto*

lemma *not_is_pole_const* [*simp*]: $\neg is_pole\ (\lambda_::'a::perfect_space.\ c :: complex)$

z

using *not_tendsto_and_filterlim_at_infinity*[of at *z* $\lambda_::'a.\ c\ c$] **by** (*auto simp: is_pole_def*)

lemma *has_laurent_expansion_imp_is_pole_iff*:

assumes $F: (\lambda x. f\ (z + x))\ has_laurent_expansion\ F$

shows $is_pole\ f\ z \longleftrightarrow fls_subdegree\ F < 0$

proof

assume *pole*: $is_pole\ f\ z$

have [*simp*]: $F \neq 0$

proof

assume $F = 0$

hence $is_pole\ f\ z \longleftrightarrow is_pole\ (\lambda_.\ 0 :: complex)\ z$ **using** *assms*

by (*intro is_pole_cong*)

(*auto simp: has_laurent_expansion_def at_to_0*[of *z*] *eventually_filtermap add_ac*)

with *pole* **show** *False*

by *simp*

qed

note *pole*

also **have** $is_pole\ f\ z \longleftrightarrow$

$is_pole\ (\lambda w. fls_nth\ F\ (fls_subdegree\ F) * (w - z)\ powi\ fls_subdegree$

$F)\ z$

using *has_laurent_expansion_imp_asymp_equiv*[OF *F*] **by** (*intro is_pole_cong_asymp_equiv refl*)

also **have** $\dots \longleftrightarrow is_pole\ (\lambda w. (w - z)\ powi\ fls_subdegree\ F)\ z$

by *simp*

finally **have** *pole'*: \dots .

have *False* **if** $fls_subdegree\ F \geq 0$

proof -

have $(\lambda w. (w - z)\ powi\ fls_subdegree\ F)\ holomorphic_on\ UNIV$

using *that* **by** (*intro holomorphic_intros*) *auto*

hence $\neg is_pole\ (\lambda w. (w - z)\ powi\ fls_subdegree\ F)\ z$

by (*meson UNIV_I not_is_pole_holomorphic_open_UNIV*)

with *pole'* **show** *False*

by *simp*

```

qed
thus  $fls\_subdegree\ F < 0$ 
  by force
qed (use has_laurent_expansion_imp_is_pole[OF assms] in auto)

lemma analytic_at_imp_has_fps_expansion_0:
  assumes  $f\ analytic\_on\ \{0\}$ 
  shows  $f\ has\_fps\_expansion\ fps\_expansion\ f\ 0$ 
  using  $assms\ has\_fps\_expansion\_fps\_expansion\ analytic\_at$  by fast

lemma deriv_shift_0:  $deriv\ f\ z = deriv\ (f\ \circ\ (\lambda x. z + x))\ 0$ 
proof -
  have *:  $(f\ \circ\ (+)\ z\ has\_field\_derivative\ D)\ (at\ z')$ 
    if  $(f\ has\_field\_derivative\ D)\ (at\ (z + z'))$  for  $D\ z\ z'$  and  $f\ ::\ 'a\ \Rightarrow\ 'a$ 
  proof -
    have  $(f\ \circ\ (+)\ z\ has\_field\_derivative\ D * 1)\ (at\ z')$ 
      by (rule DERIV_chain that derivative_eq_intros refl)+ auto
    thus ?thesis by simp
  qed
  have  $(\lambda D. (f\ has\_field\_derivative\ D)\ (at\ z)) = (\lambda D. (f\ \circ\ (+)\ z\ has\_field\_derivative\ D)\ (at\ 0))$ 
  using *[of f _ z 0] *[of f \circ (+) z _ - z z] by (intro ext iffI) (auto simp: o_def)
  thus ?thesis
    by (simp add: deriv_def)
qed

lemma deriv_shift_0':  $NO\_MATCH\ 0\ z \Longrightarrow deriv\ f\ z = deriv\ (f\ \circ\ (\lambda x. z + x))\ 0$ 
  by (rule deriv_shift_0)

lemma higher_deriv_shift_0:  $(deriv\ \hat{\sim}\ n)\ f\ z = (deriv\ \hat{\sim}\ n)\ (f\ \circ\ (\lambda x. z + x))\ 0$ 
proof (induction n arbitrary: f)
  case (Suc n)
  have  $(deriv\ \hat{\sim}\ Suc\ n)\ f\ z = (deriv\ \hat{\sim}\ n)\ (deriv\ f)\ z$ 
    by (subst funpow_Suc_right) auto
  also have  $\dots = (deriv\ \hat{\sim}\ n)\ (\lambda x. deriv\ f\ (z + x))\ 0$ 
    by (subst Suc) (auto simp: o_def)
  also have  $\dots = (deriv\ \hat{\sim}\ n)\ (\lambda x. deriv\ (\lambda xa. f\ (z + x + xa))\ 0)\ 0$ 
    by (subst deriv_shift_0) (auto simp: o_def)
  also have  $(\lambda x. deriv\ (\lambda xa. f\ (z + x + xa))\ 0) = deriv\ (\lambda x. f\ (z + x))$ 
    by (rule ext) (simp add: deriv_shift_0' o_def add_ac)
  also have  $(deriv\ \hat{\sim}\ n)\ \dots\ 0 = (deriv\ \hat{\sim}\ Suc\ n)\ (f\ \circ\ (\lambda x. z + x))\ 0$ 
    by (subst funpow_Suc_right) (auto simp: o_def)
  finally show ?case .
qed auto

lemma higher_deriv_shift_0':  $NO\_MATCH\ 0\ z \Longrightarrow (deriv\ \hat{\sim}\ n)\ f\ z = (deriv\ \hat{\sim}\ n)\ (f\ \circ\ (\lambda x. z + x))\ 0$ 
  by (rule higher_deriv_shift_0)

```


lemma *analytic_at_imp_has_fps_expansion*:

assumes f *analytic_on* $\{z\}$

shows $(\lambda x. f (z + x))$ *has_fps_expansion* *fps_expansion* $f z$

proof –

have $f \circ (\lambda x. z + x)$ *analytic_on* $\{0\}$

by (*intro analytic_on_compose_gen*[*OF* *assms*] *analytic_intros*) *auto*

hence $(f \circ (\lambda x. z + x))$ *has_fps_expansion* *fps_expansion* $(f \circ (\lambda x. z + x)) 0$

unfolding *o_def* **by** (*intro analytic_at_imp_has_fps_expansion_0*) *auto*

also have $\dots =$ *fps_expansion* $f z$

by (*simp add: fps_expansion_def higher_deriv_shift_0'*)

finally show *?thesis* **by** (*simp add: add_ac*)

qed

lemma *has_laurent_expansion_zorder_0*:

assumes f *has_laurent_expansion* $F F \neq 0$

shows *zorder* $f 0 =$ *fls_subdegree* F

proof –

define G **where** $G =$ *fls_base_factor_to_fps* F

from *assms* **obtain** A **where** $A: 0 \in A$ *open* $A \wedge x. x \in A - \{0\} \implies$ *eval_fls* $F x = f x$

unfolding *has_laurent_expansion_def eventually_at_filter eventually_nhds*

by *blast*

show *?thesis*

proof (*rule zorder_eqI*)

show *open* $(A \cap$ *eball* 0 $($ *fls_conv_radius* $F)) 0 \in A \cap$ *eball* 0 $($ *fls_conv_radius* $F)$

using *assms* A **by** (*auto simp: has_laurent_expansion_def zero_ereal_def*)

show *eval_fps* G *holomorphic_on* $A \cap$ *eball* 0 $($ *fls_conv_radius* $F)$

by (*intro holomorphic_intros*) (*auto simp: fls_conv_radius_altdef G_def*)

show *eval_fps* $G 0 \neq 0$ **using** $\langle F \neq 0 \rangle$

by (*auto simp: eval_fps_at_0 G_def*)

next

fix $w ::$ *complex* **assume** $w \in A \cap$ *eball* 0 $($ *fls_conv_radius* $F)$ $w \neq 0$

thus $f w =$ *eval_fps* $G w * (w - 0)$ *powi* $($ *fls_subdegree* $F)$

using A **unfolding** G_def

by (*subst eval_fps_fls_base_factor*)

$($ *auto simp: complex_powr_of_int power_int_minus field_simps*)

qed

qed

lemma *has_laurent_expansion_zorder*:

assumes $(\lambda w. f (z + w))$ *has_laurent_expansion* $F F \neq 0$

shows *zorder* $f z =$ *fls_subdegree* F

using *has_laurent_expansion_zorder_0*[*OF* *assms*] **by** (*simp add: zorder_shift' add_ac*)

lemma *has_fps_expansion_zorder_0*:

assumes f has_fps_expansion F $F \neq 0$
shows $\text{zorder } f \ 0 = \text{int } (\text{subdegree } F)$
using *assms* has_laurent_expansion_zorder_0[*of f fps_to_fls F*]
by (*auto simp: has_fps_expansion_to_laurent fls_subdegree_fls_to_fps*)

lemma *has_fps_expansion_zorder*:

assumes $(\lambda w. f(z + w))$ has_fps_expansion F $F \neq 0$
shows $\text{zorder } f \ z = \text{int } (\text{subdegree } F)$
using *has_fps_expansion_zorder_0*[*OF assms*]
by (*simp add: zorder_shift' add_ac*)

lemma *has_fps_expansion_fls_base_factor_to_fps*:

assumes f has_laurent_expansion F
defines $n \equiv \text{fls_subdegree } F$
defines $c \equiv \text{fps_nth } (\text{fls_base_factor_to_fps } F) \ 0$
shows $(\lambda z. \text{if } z = 0 \text{ then } c \text{ else } f \ z * z^{\text{powi } -n})$ has_fps_expansion *fls_base_factor_to_fps* F

proof –

have $(\lambda z. f \ z * z^{\text{powi } -n})$ has_laurent_expansion *fls_shift* $(-(-n))$ F
by (*intro laurent_expansion_intros assms*)
also have *fls_shift* $(-(-n))$ $F = \text{fps_to_fls } (\text{fls_base_factor_to_fps } F)$
by (*simp add: n_def fls_shift_nonneg_subdegree*)
also have $(\lambda z. f \ z * z^{\text{powi } -n})$ has_laurent_expansion *fps_to_fls* $(\text{fls_base_factor_to_fps } F) \longleftrightarrow$
 $(\lambda z. \text{if } z = 0 \text{ then } c \text{ else } f \ z * z^{\text{powi } -n})$ has_laurent_expansion
fps_to_fls $(\text{fls_base_factor_to_fps } F)$
by (*intro has_laurent_expansion_cong*) (*auto simp: eventually_at_filter*)
also have $\dots \longleftrightarrow (\lambda z. \text{if } z = 0 \text{ then } c \text{ else } f \ z * z^{\text{powi } -n})$ has_fps_expansion
fls_base_factor_to_fps F
by (*subst has_fps_expansion_to_laurent*) (*auto simp: c_def*)
finally show *?thesis* .

qed

lemma *zero_has_laurent_expansion_imp_eq_0*:

assumes $(\lambda_. 0)$ has_laurent_expansion F
shows $F = 0$
proof –
have $at(0 :: \text{complex}) \neq \text{bot}$
by *auto*
moreover have $(\lambda z. \text{if } z = 0 \text{ then } \text{fls_nth } F (\text{fls_subdegree } F) \text{ else } 0)$ has_fps_expansion
fls_base_factor_to_fps F (**is** *?f* has_fps_expansion $_$)
using *has_fps_expansion_fls_base_factor_to_fps*[*OF assms*] **by** (*simp cong:*
if_cong)
hence *isCont* *?f* 0
using *has_fps_expansion_imp_continuous* **by** *blast*
hence *?f* $-0 \rightarrow \text{fls_nth } F (\text{fls_subdegree } F)$
by (*auto simp: isCont_def*)
moreover have *?f* $-0 \rightarrow 0 \longleftrightarrow (\lambda_::\text{complex}. 0 :: \text{complex}) -0 \rightarrow 0$
by (*intro filterlim_cong*) (*auto simp: eventually_at_filter*)

hence $?f - 0 \rightarrow 0$
 by *simp*
 ultimately have $\text{fls_nth } F (\text{fls_subdegree } F) = 0$
 by (*rule tendsto_unique*)
 thus *?thesis*
 by (*meson nth_fls_subdegree_nonzero*)
qed

lemma *has_laurent_expansion_unique*:
 assumes *f has_laurent_expansion F f has_laurent_expansion G*
 shows $F = G$
proof –
 from *assms* have $(\lambda x. f x - f x)$ *has_laurent_expansion F - G*
 by (*intro laurent_expansion_intros*)
 hence $(\lambda_. 0)$ *has_laurent_expansion F - G*
 by *simp*
 hence $F - G = 0$
 by (*rule zero_has_laurent_expansion_imp_eq_0*)
 thus *?thesis*
 by *simp*
qed

lemma *laurent_expansion_eqI*:
 assumes $(\lambda x. f (z + x))$ *has_laurent_expansion F*
 shows *laurent_expansion f z = F*
 using *assms has_laurent_expansion_isolated has_laurent_expansion_not_essential*
has_laurent_expansion_unique not_essential_has_laurent_expansion **by**
blast

lemma *laurent_expansion_0_eqI*:
 assumes *f has_laurent_expansion F*
 shows *laurent_expansion f 0 = F*
 using *assms laurent_expansion_eqI[of f 0]* **by** *simp*

lemma *has_laurent_expansion_nonzero_imp_eventually_nonzero*:
 assumes *f has_laurent_expansion F F \neq 0*
 shows *eventually* $(\lambda x. f x \neq 0)$ *(at 0)*
proof (*rule ccontr*)
 assume \neg *eventually* $(\lambda x. f x \neq 0)$ *(at 0)*
 with *assms* have *eventually* $(\lambda x. f x = 0)$ *(at 0)*
 by (*intro not_essential_frequently_0_imp_eventually_0 has_laurent_expansion_isolated*
has_laurent_expansion_not_essential)
 (*auto simp: frequently_def*)
 hence $(f \text{ has_laurent_expansion } 0) \longleftrightarrow ((\lambda_. 0) \text{ has_laurent_expansion } 0)$
 by (*intro has_laurent_expansion_cong*) *auto*
 hence *f has_laurent_expansion 0*
 by *simp*
 with *assms(1)* have $F = 0$
 using *has_laurent_expansion_unique* **by** *blast*

with $\langle F \neq 0 \rangle$ **show** *False*
by *contradiction*
qed

lemma *has_laurent_expansion_eventually_nonzero_iff'*:

assumes *f has_laurent_expansion F*

shows *eventually* $(\lambda x. f x \neq 0)$ *(at 0)* $\longleftrightarrow F \neq 0$

proof

assume $\forall_F x$ *in at 0. f x $\neq 0$*

moreover have $\neg (\forall_F x$ *in at 0. f x $\neq 0$)* **if** $F=0$

proof $-$

have $\forall_F x$ *in at 0. f x = 0*

using *assms that unfolding has_laurent_expansion_def* **by** *simp*

then show *?thesis unfolding not_eventually*

by *(auto elim:eventually_frequentlyE)*

qed

ultimately show $F \neq 0$ **by** *auto*

qed *(simp add:has_laurent_expansion_nonzero_imp_eventually_nonzero[OF assms])*

lemma *has_laurent_expansion_eventually_nonzero_iff*:

assumes $(\lambda w. f (z+w))$ *has_laurent_expansion F*

shows *eventually* $(\lambda x. f x \neq 0)$ *(at z)* $\longleftrightarrow F \neq 0$

apply *(subst eventually_at_to_0)*

apply *(rule has_laurent_expansion_eventually_nonzero_iff')*

using *assms by (simp add:add.commute)*

lemma *has_laurent_expansion_inverse* [*laurent_expansion_intros*]:

assumes *f has_laurent_expansion F*

shows $(\lambda x. inverse (f x))$ *has_laurent_expansion inverse F*

proof *(cases F = 0)*

case *True*

thus *?thesis using assms*

by *(auto simp: has_laurent_expansion_def)*

next

case *False*

define *G* **where** $G = laurent_expansion (\lambda x. inverse (f x)) 0$

from *False* **have** *ev: eventually* $(\lambda z. f z \neq 0)$ *(at 0)*

by *(intro has_laurent_expansion_nonzero_imp_eventually_nonzero[OF assms])*

have $*$: $(\lambda x. inverse (f x))$ *has_laurent_expansion G* **unfolding** *G_def*

by *(intro not_essential_has_laurent_expansion_0 isolated_singularity_at_inverse*

not_essential_inverse

has_laurent_expansion_isolated_0[OF assms] has_laurent_expansion_not_essential_0[OF

assms])

have $(\lambda x. f x * inverse (f x))$ *has_laurent_expansion F * G*

by *(intro laurent_expansion_intros assms *)*

also have *?this $\longleftrightarrow (\lambda x. 1)$ has_laurent_expansion F * G*

by *(intro has_laurent_expansion_cong_refl eventually_mono[OF ev]) auto*

finally have $(\lambda x. 1)$ *has_laurent_expansion F * G* .

moreover have $(\lambda_. 1)$ *has_laurent_expansion 1*
by *simp*
ultimately have $F * G = 1$
using *has_laurent_expansion_unique* **by** *blast*
hence $G = \text{inverse } F$
using *inverse_unique* **by** *blast*
with * show *?thesis*
by *simp*
qed

lemma *has_laurent_expansion_power_int* [*laurent_expansion_intros*]:
 f *has_laurent_expansion F* $\implies (\lambda x. f x \text{ powi } n)$ *has_laurent_expansion (F powi n)*
by (*auto simp: power_int_def intro!: laurent_expansion_intros*)

lemma *has_fps_expansion_0_analytic_continuation*:
assumes f *has_fps_expansion 0* *f holomorphic_on A*
assumes *open A connected A 0 ∈ A x ∈ A*
shows $f x = 0$
proof –
have *eventually* $(\lambda z. z \in A \wedge f z = 0)$ (*nhds 0*) **using** *assms*
by (*intro eventually_conj eventually_nhds_in_open*) (*auto simp: has_fps_expansion_def*)
then obtain B where B : *open B 0 ∈ B $\forall z \in B. z \in A \wedge f z = 0$*
unfolding *eventually_nhds* **by** *blast*
show *?thesis*
proof (*rule analytic_continuation_open*[**where** $f = f$ **and** $g = \lambda_. 0$])
show $B \neq \{\}$
using $\langle \text{open } B \rangle B$ **by** *auto*
show *connected A*
using *assms* **by** *auto*
qed (*use assms B in auto*)
qed

lemma *has_laurent_expansion_0_analytic_continuation*:
assumes f *has_laurent_expansion 0* *f holomorphic_on A - {0}*
assumes *open A connected A 0 ∈ A x ∈ A - {0}*
shows $f x = 0$
proof –
have *eventually* $(\lambda z. z \in A - \{0\} \wedge f z = 0)$ (*at 0*) **using** *assms*
by (*intro eventually_conj eventually_at_in_open*) (*auto simp: has_laurent_expansion_def*)
then obtain B where B : *open B 0 ∈ B $\forall z \in B - \{0\}. z \in A - \{0\} \wedge f z = 0$*
unfolding *eventually_at_filter eventually_nhds* **by** *blast*
show *?thesis*
proof (*rule analytic_continuation_open*[**where** $f = f$ **and** $g = \lambda_. 0$])
show $B - \{0\} \neq \{\}$
using $\langle \text{open } B \rangle \langle 0 \in B \rangle$ **by** (*metis insert_Diff not_open_singleton*)
show *connected (A - {0})*
using *assms* **by** (*intro connected_open_delete*) *auto*

qed (use *assms B in auto*)
qed

lemma *has_fps_expansion_cong*:
assumes *eventually* ($\lambda x. f x = g x$) (*nhds 0*) $F = G$
shows f *has_fps_expansion* $F \longleftrightarrow g$ *has_fps_expansion* G
using *assms(2)* **by** (*auto simp: has_fps_expansion_def elim!: eventually_elim2[OF assms(1)]*)

lemma *zor_poly_has_fps_expansion*:
assumes f *has_laurent_expansion* F $F \neq 0$
shows $\text{zor_poly } f \ 0$ *has_fps_expansion* *fls_base_factor_to_fps* F
proof –
note [*simp*] = $\langle F \neq 0 \rangle$
have *eventually* ($\lambda z. f z \neq 0$) (*at 0*)
by (*rule has_laurent_expansion_nonzero_imp_eventually_nonzero[OF assms]*)
hence *freq*: *frequently* ($\lambda z. f z \neq 0$) (*at 0*)
by (*rule eventually_frequently[rotated]*) *auto*

have *: *isolated_singularity_at* $f \ 0$ *not_essential* $f \ 0$
using *has_laurent_expansion_isolated_0[OF assms(1)] has_laurent_expansion_not_essential_0[OF assms(1)]*
by *auto*

define G **where** $G = \text{fls_base_factor_to_fps } F$
define n **where** $n = \text{zorder } f \ 0$
have n_altdef : $n = \text{fls_subdegree } F$
using *has_laurent_expansion_zorder_0 [OF assms(1)]* **by** (*simp add: n_def*)
obtain r **where** r : $\text{zor_poly } f \ 0 \ 0 \neq 0$ $\text{zor_poly } f \ 0$ *holomorphic_on* $\text{cball } 0 \ r$ $r > 0$

$$\forall w \in \text{cball } 0 \ r - \{0\}. f \ w = \text{zor_poly } f \ 0 \ w * w^{\text{powi } n} \wedge \text{zor_poly } f \ 0 \ w \neq 0$$

using *zorder_exist[OF * freq]* **unfolding** n_def **by** *auto*
obtain r' **where** r' : $r' > 0 \ \forall x \in \text{ball } 0 \ r' - \{0\}. \text{eval_fls } F \ x = f \ x$
using *assms(1)* **unfolding** *has_laurent_expansion_def eventually_at_filter eventually_nhds_metric ball_def*
by (*auto simp: dist_commute*)
have *holo*: $\text{zor_poly } f \ 0$ *holomorphic_on* $\text{ball } 0 \ r$
by (*rule holomorphic_on_subset[OF r(2)]*) *auto*

have ($\lambda z. \text{if } z = 0 \text{ then } \text{fps_nth } G \ 0 \text{ else } f \ z * z^{\text{powi } -n}$) *has_fps_expansion* G
unfolding G_def n_altdef **by** (*intro has_fps_expansion_fls_base_factor_to_fps assms*)

also **have** $?this \longleftrightarrow \text{zor_poly } f \ 0$ *has_fps_expansion* G
proof (*intro has_fps_expansion_cong*)
have *eventually* ($\lambda z. z \in \text{ball } 0 \ (\min \ r \ r')$) (*nhds 0*)
using $\langle r > 0 \rangle \langle r' > 0 \rangle$ **by** (*intro eventually_nhds_in_open*) *auto*
thus $\forall F \ x$ *in nhds 0*. (*if* $x = 0$ *then* $G \ \$ \ 0$ *else* $f \ x * x^{\text{powi } -n} = \text{zor_poly } f \ 0 \ x$)

```

proof eventually_elim
  case (elim w)
  have w: w ∈ ball 0 r w ∈ ball 0 r'
    using elim by auto
  show ?case
  proof (cases w = 0)
    case False
    hence f w = zor_poly f 0 w * w powi n
      using r w by auto
    thus ?thesis using False
      by (simp add: powr_minus complex_powr_of_int power_int_minus)
  next
  case [simp]: True
  obtain R where R: R > 0 R ≤ r R ≤ r' R ≤ fls_conv_radius F
    using ⟨r > 0⟩ ⟨r' > 0⟩ assms(1) unfolding has_laurent_expansion_def
    by (smt (verit, ccfv SIG) ereal_dense2 ereal_less(2) less_ereal.simps(1)
order.strict_implies_order order_trans)
  have eval_fps G 0 = zor_poly f 0 0
  proof (rule analytic_continuation_open[where f = eval_fps G and g =
zor_poly f 0])
    show connected (ball 0 R :: complex set)
      by auto
    have of_real R / 2 ∈ ball 0 R - {0 :: complex}
      using R by auto
    thus ball 0 R - {0 :: complex} ≠ {}
      by blast
    show eval_fps G holomorphic_on ball 0 R
      using R less_le_trans[OF _ R(4)] unfolding G_def
      by (intro holomorphic_intros) (auto simp: fls_conv_radius_altdef)
    show zor_poly f 0 holomorphic_on ball 0 R
      by (rule holomorphic_on_subset[OF holo]) (use R in auto)
    show eval_fps G z = zor_poly f 0 z if z ∈ ball 0 R - {0} for z
      using that r r' R n_altdef unfolding G_def
      by (subst eval_fps_fls_base_factor)
        (auto simp: complex_powr_of_int field_simps power_int_minus
n_def)
    qed (use R in auto)
  hence zor_poly f 0 0 = fps_nth G 0
    by (simp add: eval_fps_at_0)
  thus ?thesis by simp
  qed
qed
qed (use r' in auto)
finally show ?thesis
  by (simp add: G_def)
qed

```

lemma zorder_geI_0:

assumes f analytic_on {0} f holomorphic_on A open A connected A 0 ∈ A z ∈

```

A f z ≠ 0
  assumes  $\bigwedge k. k < n \implies (\text{deriv } \sim k) f 0 = 0$ 
  shows  $\text{zorder } f 0 \geq n$ 
proof -
  define F where  $F = \text{fps\_expansion } f 0$ 
  from assms have f has_fps_expansion F
  unfolding F_def using analytic_at_imp_has_fps_expansion_0 by blast
  hence laurent: f has_laurent_expansion fps_to_fls F and [simp]:  $f 0 = \text{fps\_nth } F 0$ 
  by (simp_all add: has_fps_expansion_to_laurent)

  have [simp]:  $F \neq 0$ 
  proof
    assume [simp]:  $F = 0$ 
    hence  $f z = 0$ 
    proof (cases z = 0)
      case False
      have f has_laurent_expansion 0
      using laurent by simp
      thus ?thesis
    proof (rule has_laurent_expansion_0_analytic_continuation)
      show f holomorphic_on A - {0}
      using assms(2) by (rule holomorphic_on_subset) auto
    qed (use assms False in auto)
  qed auto
  with  $\langle f z \neq 0 \rangle$  show False by contradiction
qed

  have  $\text{zorder } f 0 = \text{int } (\text{subdegree } F)$ 
  using has_laurent_expansion_zorder_0[OF laurent] by (simp add: fls_subdegree_fls_to_fps)
  also have  $\text{subdegree } F \geq n$ 
  using assms by (intro subdegree_geI  $\langle F \neq 0 \rangle$ ) (auto simp: F_def fps_expansion_def)
  hence  $\text{int } (\text{subdegree } F) \geq \text{int } n$ 
  by simp
  finally show ?thesis .
qed

lemma zorder_geI:
  assumes f analytic_on {x} f holomorphic_on A open A connected A  $x \in A$   $z \in A$   $f z \neq 0$ 
  assumes  $\bigwedge k. k < n \implies (\text{deriv } \sim k) f x = 0$ 
  shows  $\text{zorder } f x \geq n$ 
proof -
  have  $\text{zorder } f x = \text{zorder } (f \circ (\lambda u. u + x)) 0$ 
  by (subst zorder_shift) (auto simp: o_def)
  also have  $\dots \geq n$ 
  proof (rule zorder_geI_0)
    show  $(f \circ (\lambda u. u + x)) \text{ analytic\_on } \{0\}$ 
    by (intro analytic_on_compose_gen[OF _ assms(1)] analytic_intros) auto
  qed

```



```

show  $f \circ (\lambda u. u + x)$  holomorphic_on  $((+) (-x)) \text{ ' } A$ 
by (intro holomorphic_on_compose_gen[OF_assms(2)] holomorphic_intros)
auto
show connected  $((+) (-x)) \text{ ' } A$ 
by (intro connected_continuous_image continuous_intros assms)
show open  $((+) (-x)) \text{ ' } A$ 
by (intro open_translation assms)
show  $z - x \in (+) (-x) \text{ ' } A$ 
using  $\langle z \in A \rangle$  by auto
show  $0 \in (+) (-x) \text{ ' } A$ 
using  $\langle x \in A \rangle$  by auto
show  $(f \circ (\lambda u. u + x)) (z - x) \neq 0$ 
using  $\langle f z \neq 0 \rangle$  by auto
next
fix  $k :: \text{nat}$  assume  $k < n$ 
hence  $(\text{deriv } \widehat{\sim} k) f x = 0$ 
using assms by simp
also have  $(\text{deriv } \widehat{\sim} k) f x = (\text{deriv } \widehat{\sim} k) (f \circ (+) x) 0$ 
by (subst higher_deriv_shift_0) auto
finally show  $(\text{deriv } \widehat{\sim} k) (f \circ (\lambda u. u + x)) 0 = 0$ 
by (subst add commute) auto
qed
finally show ?thesis .
qed

```

```

lemma has_laurent_expansion_divide [laurent_expansion_intros]:
assumes  $f$  has_laurent_expansion  $F$  and  $g$  has_laurent_expansion  $G$ 
shows  $(\lambda x. f x / g x)$  has_laurent_expansion  $(F / G)$ 
proof -
have  $(\lambda x. f x * \text{inverse } (g x))$  has_laurent_expansion  $(F * \text{inverse } G)$ 
by (intro laurent_expansion_intros assms)
thus ?thesis
by (simp add: field_simps)
qed

```

```

lemma vector_derivative_translate [simp]:
 $\text{vector\_derivative } ((+) z \circ g) \text{ (at } x \text{ within } A) = \text{vector\_derivative } g \text{ (at } x \text{ within } A)$ 
proof -
have  $((+) z \circ g)$  has_vector_derivative  $g'$  (at  $x$  within  $A$ )
if  $g$  has_vector_derivative  $g'$  (at  $x$  within  $A$ ) for  $g :: \text{real} \Rightarrow 'a$  and  $z g'$ 
unfolding o_def using that by (auto intro!: derivative_eq_intros)
from this[of  $g \_ z$ ] this[of  $\lambda x. z + g x \_ -z$ ] show ?thesis
unfolding vector_derivative_def
by (intro arg_cong[where  $f = \text{Eps}$ ] ext) (auto simp: o_def algebra_simps)
qed

```

```

lemma has_contour_integral_translate:
 $(f \text{ has\_contour\_integral } I) ((+) z \circ g) \longleftrightarrow ((\lambda x. f (x + z)) \text{ has\_contour\_integral } I)$ 

```

I) g
by (simp add: has_contour_integral_def add_ac)

lemma contour_integrable_translate:
 f contour_integrable_on $((+) z \circ g) \longleftrightarrow (\lambda x. f (x + z))$ contour_integrable_on g
by (simp add: contour_integrable_on_def has_contour_integral_translate)

lemma contour_integral_translate:
 $\text{contour_integral } ((+) z \circ g) f = \text{contour_integral } g (\lambda x. f (x + z))$
by (simp add: contour_integral_def contour_integrable_translate has_contour_integral_translate)

lemma residue_shift_0: residue $f z = \text{residue } (\lambda x. f (z + x)) 0$

proof -

define Q **where**

$Q = (\lambda r f z \varepsilon. (f \text{ has_contour_integral_complex_of_real } (2 * \text{pi}) * i * r) (\text{circlepath } z \varepsilon))$

define P **where**

$P = (\lambda r f z. \exists e > 0. \forall \varepsilon > 0. \varepsilon < e \longrightarrow Q r f z \varepsilon)$

have path_eq: circlepath $(z - w) \varepsilon = (+) (-w) \circ \text{circlepath } z \varepsilon$ **for** $z w \varepsilon$

by (simp add: circlepath_def o_def part_circlepath_def algebra_simps)

have *: $P r f z$ **if** $P r (\lambda x. f (x + w)) (z - w)$ **for** $r w f z$

using that **by** (auto simp: P_def Q_def path_eq has_contour_integral_translate)

have (SOME $r. P r f z$) = (SOME $r. P r (\lambda x. f (z + x)) 0$)

using *[of $f z z$] *[of $\lambda x. f (z + x) - z$]

by (intro arg_cong[where $f = Eps$] ext iffI) (simp_all add: add_ac)

thus ?thesis

by (simp add: residue_def P_def Q_def)

qed

lemma residue_shift_0': NO_MATCH $0 z \implies \text{residue } f z = \text{residue } (\lambda x. f (z + x)) 0$

by (rule residue_shift_0)

lemma has_laurent_expansion_residue_0:

assumes f has_laurent_expansion F

shows residue $f 0 = \text{fls_residue } F$

proof (cases fls_subdegree $F \geq 0$)

case True

have residue $f 0 = \text{residue } (\text{eval_fls } F) 0$

using assms **by** (intro residue_cong) (auto simp: has_laurent_expansion_def eq_commute)

also have ... = 0

by (rule residue_holo[OF __ holomorphic_on_eval_flss[OF order_refl]])

(use True assms in <auto simp: has_laurent_expansion_def zero_ereal_def>)

also have ... = fls_residue F

using True **by** simp

finally show ?thesis .

next

```

case False
hence  $F \neq 0$ 
  by auto
have *: zor_poly f 0 has_fps_expansion fls_base_factor_to_fps F
  by (intro zor_poly_has_fps_expansion False assms  $\langle F \neq 0 \rangle$ )

have residue f 0 = (deriv  $\sim$  (nat ( $-zorder$  f 0) - 1)) (zor_poly f 0) 0 / fact
(nat ( $-zorder$  f 0) - 1)
  by (intro residue_pole_order has_laurent_expansion_isolated_0[OF assms]
      has_laurent_expansion_imp_is_pole_0[OF assms]) (use False in auto)
also have ... = fls_residue F
  using has_laurent_expansion_zorder_0[OF assms  $\langle F \neq 0 \rangle$ ] False
  by (subst fps_nth_fps_expansion [OF *, symmetric]) (auto simp: of_nat_diff)
finally show ?thesis .
qed

```

```

lemma has_laurent_expansion_residue:
assumes  $(\lambda x. f(z + x))$  has_laurent_expansion F
shows residue f z = fls_residue F
using has_laurent_expansion_residue_0[OF assms] by (simp add: residue_shift_0')

```

```

lemma eval_fls_has_laurent_expansion [laurent_expansion_intros]:
assumes fls_conv_radius F > 0
shows eval_fls F has_laurent_expansion F
using assms by (auto simp: has_laurent_expansion_def)

```

```

lemma fps_expansion_unique_complex:
fixes F G :: complex fps
assumes f has_fps_expansion F f has_fps_expansion G
shows  $F = G$ 
using assms unfolding fps_eq_iff by (auto simp: fps_eq_iff fps_nth_fps_expansion)

```

```

lemma fps_expansion_eqI:
assumes f has_fps_expansion F
shows fps_expansion f 0 = F
using assms unfolding fps_eq_iff
by (auto simp: fps_eq_iff fps_nth_fps_expansion fps_expansion_def)

```

```

lemma has_fps_expansion_imp_eval_fps_eq:
assumes f has_fps_expansion F norm z < r
assumes f holomorphic_on ball 0 r
shows eval_fps F z = f z

```

```

proof -
have [simp]: fps_expansion f 0 = F
  by (rule fps_expansion_eqI) fact
have *: f holomorphic_on eball 0 (ereal r)
  using assms by simp
from conv_radius_fps_expansion[OF *] have fps_conv_radius F  $\geq$  ereal r
  by simp

```

```

have eval_fps (fps_expansion f 0) z = f (0 + z)
  by (rule eval_fps_expansion'[OF *]) (use assms in auto)
thus ?thesis
  by simp
qed

```

lemma fls_conv_radius_ge:

```

assumes f has_laurent_expansion F
assumes f holomorphic_on eball 0 r - {0}
shows fls_conv_radius F ≥ r
proof -
define n where n = fls_subdegree F
define G where G = fls_base_factor_to_fps F
define g where g = (λz. if z = 0 then fps_nth G 0 else f z * z powi -n)
have G: g has_fps_expansion G
  unfolding G_def g_def n_def
  by (intro has_fps_expansion fls_base_factor_to_fps assms)
have (λz. f z * z powi -n) holomorphic_on eball 0 r - {0}
  by (intro holomorphic_intros assms) auto
also have ?this ↔ g holomorphic_on eball 0 r - {0}
  by (intro holomorphic_cong) (auto simp: g_def)
finally have g analytic_on eball 0 r - {0}
  by (subst analytic_on_open) auto
moreover have g analytic_on {0}
  using G has_fps_expansion_imp_analytic_0 by auto
ultimately have g analytic_on (eball 0 r - {0} ∪ {0})
  by (subst analytic_on_Un) auto
hence g analytic_on eball 0 r
  by (rule analytic_on_subset) auto
hence g holomorphic_on eball 0 r
  by (subst (asm) analytic_on_open) auto
hence fps_conv_radius (fps_expansion g 0) ≥ r
  by (intro conv_radius_fps_expansion)
also have fps_expansion g 0 = G
  using G by (intro fps_expansion_eqI)
finally show ?thesis
  by (simp add: fls_conv_radius_altdef G_def)
qed

```

```

lemma connected_eball [intro]: connected (eball (z :: 'a :: real_normed_vector) r)
  by (cases r) auto

```

lemma eval_fls_eqI:

```

assumes f has_laurent_expansion F f holomorphic_on eball 0 r - {0}
assumes z ∈ eball 0 r - {0}
shows eval_fls F z = f z
proof -
have conv: fls_conv_radius F ≥ r
  by (intro fls_conv_radius_ge[OF assms(1,2)])

```

```

have ( $\lambda z. \text{eval\_fls } F z - f z$ ) has_laurent_expansion  $F - F$ 
using assms by (intro laurent_expansion_intros assms) (auto simp: has_laurent_expansion_def)
hence ( $\lambda z. \text{eval\_fls } F z - f z$ ) has_laurent_expansion 0
by simp
hence  $\text{eval\_fls } F z - f z = 0$ 
proof (rule has_laurent_expansion_0_analytic_continuation)
have  $\text{ereal } 0 \leq \text{ereal } (\text{norm } z)$ 
by simp
also have  $\text{norm } z < r$ 
using assms by auto
finally have  $r > 0$ 
by (simp add: zero_ereal_def)
thus open (eball 0  $r :: \text{complex set}$ ) connected (eball 0  $r :: \text{complex set}$ )
 $0 \in \text{eball } 0 r \ z \in \text{eball } 0 r - \{0\}$ 
using assms by (auto simp: zero_ereal_def)
qed (auto intro!: holomorphic_intros assms less_le_trans[OF _ conv] split:
if_splits)
thus ?thesis by simp
qed

```

lemma *fls_nth_as_contour_integral*:

```

assumes  $F: f$  has_laurent_expansion  $F$ 
assumes holo:  $f$  holomorphic_on ball 0  $r - \{0\}$ 
assumes  $R: 0 < R \ R < r$ 
shows ( $\lambda z. f z * z^{\text{powi } -(n+1)}$ ) has_contour_integral
 $\text{complex\_of\_real } (2 * \text{pi}) * \text{i} * \text{fls\_nth } F n$  (circlepath 0  $R$ )
proof -
define  $I$  where  $I = (\lambda z. f z * z^{\text{powi } -(n+1)})$ 
have ( $I$  has_contour_integral  $\text{complex\_of\_real } (2 * \text{pi}) * \text{i} * \text{residue } I 0$ )
(circlepath 0  $R$ )
proof (rule base_residue)
show open (ball (0::complex)  $r$ )  $0 \in \text{ball } (0::complex) r$ 
using  $R \ F$  by (auto simp: has_laurent_expansion_def zero_ereal_def)
qed (use R in  $\langle \text{auto intro!: holomorphic_intros holomorphic_on_subset[OF holo]$ 
 $\text{simp: } I\_def \text{ split: if_splits} \rangle$ )
also have  $\text{residue } I 0 = \text{fls\_residue } (\text{fls\_shift } (n + 1) F)$ 
unfolding  $I\_def$  by (intro has_laurent_expansion_residue_0 laurent_expansion_intros
 $F$ )
also have  $\dots = \text{fls\_nth } F n$ 
by simp
finally show ?thesis
by (simp add: I_def)
qed

```

lemma *tendsto_0_subdegree_iff_0*:

```

assumes  $F: f$  has_laurent_expansion  $F$  and  $F \neq 0$ 
shows ( $f - 0 \rightarrow 0$ )  $\longleftrightarrow \text{fls\_subdegree } F > 0$ 
proof -
have ?thesis if is_pole  $f 0$ 

```

```

proof -
  have fls_subdegree  $F < 0$ 
    using is_pole_0_imp_neg_fls_subdegree[OF F that] .
  moreover then have  $\neg f \rightarrow 0$ 
    using  $\langle is\_pole\ f\ 0 \rangle F\ at\_neq\_bot$ 
      has_laurent_expansion_imp_filterlim_infinity_0
      not_tendsto_and_filterlim_at_infinity that
    by blast
  ultimately show ?thesis by auto
qed
moreover have ?thesis if  $\neg is\_pole\ f\ 0 \ \exists x. f \rightarrow x$ 
proof -
  have fls_subdegree  $F \geq 0$ 
    using has_laurent_expansion_imp_is_pole_0[OF F] that(1)
    by linarith
  have  $f \rightarrow 0$  if fls_subdegree  $F > 0$ 
    using fls_eq0_below_subdegree[OF that]
    by (metis  $F < 0 \leq fls\_subdegree\ F \rangle$  has_laurent_expansion_imp_tendsto_0)
  moreover have fls_subdegree  $F > 0$  if  $f \rightarrow 0$ 
proof -
  have False if fls_subdegree  $F = 0$ 
proof -
  have  $f \rightarrow fls\_nth\ F\ 0$ 
    using has_laurent_expansion_imp_tendsto_0
      [OF F  $\langle fls\_subdegree\ F \geq 0 \rangle$ ] .
  then have  $fls\_nth\ F\ 0 = 0$  using  $\langle f \rightarrow 0 \rangle$ 
    using LIM_unique by blast
  then have  $F = 0$ 
    using nth_fls_subdegree_zero_iff  $\langle fls\_subdegree\ F = 0 \rangle$ 
    by metis
  with  $\langle F \neq 0 \rangle$  show False by auto
qed
  with  $\langle fls\_subdegree\ F \geq 0 \rangle$ 
  show ?thesis by fastforce
qed
  ultimately show ?thesis by auto
qed
moreover have is_pole  $f\ 0 \vee (\exists x. f \rightarrow x)$ 
proof -
  have not_essential  $f\ 0$ 
    using  $F$  has_laurent_expansion_not_essential_0 by auto
  then show ?thesis unfolding not_essential_def
    by auto
qed
  ultimately show ?thesis by auto
qed

```

lemma *tendsto_0_subdegree_iff*:

assumes $F:(\lambda w. f(z+w))$ *has_laurent_expansion* F **and** $F \neq 0$

```

shows  $(f - z \rightarrow 0) \longleftrightarrow \text{fls\_subdegree } F > 0$ 
apply (subst Lim_at_zero)
apply (rule tendsto_0_subdegree_iff_0)
using assms by auto

```

lemma *is_pole_0_deriv_divide_iff*:

```

assumes  $F: f \text{ has\_laurent\_expansion } F$  and  $F \neq 0$ 
shows  $\text{is\_pole } (\lambda x. \text{deriv } f x / f x) 0 \longleftrightarrow \text{is\_pole } f 0 \vee (f - 0 \rightarrow 0)$ 

```

proof -

```

have  $(\lambda x. \text{deriv } f x / f x) \text{ has\_laurent\_expansion } \text{fls\_deriv } F / F$ 
using  $F$  by (auto intro: laurent_expansion_intros)

```

```

have  $\text{is\_pole } (\lambda x. \text{deriv } f x / f x) 0 \longleftrightarrow$ 
 $\text{fls\_subdegree } (\text{fls\_deriv } F / F) < 0$ 
apply (rule is_pole_fls_subdegree_iff)
using  $F$  by (auto intro: laurent_expansion_intros)

```

```

also have  $\dots \longleftrightarrow \text{is\_pole } f 0 \vee (f - 0 \rightarrow 0)$ 

```

proof (cases $\text{fls_subdegree } F = 0$)

case True

```

then have  $\text{fls\_subdegree } (\text{fls\_deriv } F / F) \geq 0$ 
by (metis diff_zero div_0 ‹ $F \neq 0$ › fls_deriv_subdegree0
fls_divide_subdegree)

```

moreover then have $\neg \text{is_pole } f 0$

```

by (metis  $F$  True is_pole_0_imp_neg_fls_subdegree less_le)

```

moreover have $\neg (f - 0 \rightarrow 0)$

```

using tendsto_0_subdegree_iff_0[OF  $F$  ‹ $F \neq 0$ ›] True by auto

```

ultimately show *?thesis* by auto

next

case False

then have $\text{fls_deriv } F \neq 0$

```

by (metis fls_const_subdegree fls_deriv_eq_0_iff)

```

then have $\text{fls_subdegree } (\text{fls_deriv } F / F) =$

```

 $\text{fls\_subdegree } (\text{fls\_deriv } F) - \text{fls\_subdegree } F$ 

```

```

by (rule fls_divide_subdegree[OF ‹ $F \neq 0$ ›])

```

moreover have $\text{fls_subdegree } (\text{fls_deriv } F) = \text{fls_subdegree } F - 1$

```

using fls_subdegree_deriv[OF False] .

```

ultimately have $\text{fls_subdegree } (\text{fls_deriv } F / F) < 0$ by auto

moreover have $f - 0 \rightarrow 0 = (0 < \text{fls_subdegree } F)$

```

using tendsto_0_subdegree_iff_0[OF  $F$  ‹ $F \neq 0$ ›] .

```

moreover have $\text{is_pole } f 0 = (\text{fls_subdegree } F < 0)$

```

using is_pole_fls_subdegree_iff  $F$  by auto

```

ultimately show *?thesis* using False by auto

qed

finally show *?thesis* .

qed

lemma *is_pole_deriv_divide_iff*:

```

assumes  $F: (\lambda w. f (z+w)) \text{ has\_laurent\_expansion } F$  and  $F \neq 0$ 
shows  $\text{is\_pole } (\lambda x. \text{deriv } f x / f x) z \longleftrightarrow \text{is\_pole } f z \vee (f - z \rightarrow 0)$ 

```

proof –

define $ff\ df$ **where** $ff=(\lambda w. f\ (z+w))$ **and** $df=(\lambda w. deriv\ f\ (z + w))$
have $is_pole\ (\lambda x. deriv\ f\ x / f\ x)\ z$
 $\longleftrightarrow is_pole\ (\lambda x. deriv\ ff\ x / ff\ x)\ 0$
unfolding $ff_def\ df_def$
by ($simp\ add:deriv_shift_0'\ is_pole_shift_0'\ comp_def\ algebra_simps$)
moreover **have** $is_pole\ f\ z \longleftrightarrow is_pole\ ff\ 0$
unfolding ff_def **by** ($auto\ simp:is_pole_shift_0'$)
moreover **have** $(f\ -z \rightarrow 0) \longleftrightarrow (ff\ -0 \rightarrow 0)$
unfolding ff_def **by** ($simp\ add: LIM_offset_zero_iff$)
moreover **have** $is_pole\ (\lambda x. deriv\ ff\ x / ff\ x)\ 0 = (is_pole\ ff\ 0 \vee ff\ -0 \rightarrow 0)$
apply ($rule\ is_pole_0\ deriv_divide_iff$)
using $F\ ff_def\ \langle F \neq 0 \rangle$ **by** $blast+$
ultimately **show** $?thesis$ **by** $auto$

qed

lemma $subdegree_imp_eventually_deriv_nzero_0$:

assumes $F:f$ **has_laurent_expansion** F **and** $fls_subdegree\ F \neq 0$
shows $eventually\ (\lambda z. deriv\ f\ z \neq 0)$ **(at** $0)$

proof –

have $deriv\ f$ **has_laurent_expansion** $fls_deriv\ F$
using $has_laurent_expansion_deriv[OF\ F]$.
moreover **have** $fls_deriv\ F \neq 0$
using $\langle fls_subdegree\ F \neq 0 \rangle$
by ($metis\ fls_const_subdegree\ fls_deriv_eq_0_iff$)
ultimately **show** $?thesis$
using $has_laurent_expansion_eventually_nonzero_iff'$ **by** $blast$

qed

lemma $subdegree_imp_eventually_deriv_nzero$:

assumes $F:(\lambda w. f\ (z+w))$ **has_laurent_expansion** F
and $fls_subdegree\ F \neq 0$
shows $eventually\ (\lambda w. deriv\ f\ w \neq 0)$ **(at** $z)$

proof –

have $\forall_F\ x$ **in** $at\ 0. deriv\ (\lambda w. f\ (z + w))\ x \neq 0$
using $subdegree_imp_eventually_deriv_nzero_0\ asms$ **by** $auto$
then **show** $?thesis$
apply ($subst\ eventually_at_to_0$)
by ($simp\ add:deriv_shift_0'\ comp_def\ algebra_simps$)

qed

lemma $has_fps_expansion_imp_asympt_equiv_0$:

fixes $f :: complex \Rightarrow complex$
assumes $F: f$ **has_fps_expansion** F
defines $n \equiv subdegree\ F$
shows $f \sim[nhds\ 0]$ $(\lambda z. fps_nth\ F\ n * z \wedge n)$

proof –

have F' : f **has_laurent_expansion** $fps_to_fls\ F$
using F **has_laurent_expansion_fps** **by** $blast$


```

have  $f \sim[at\ 0]$   $(\lambda z. \text{fps\_nth } F\ n * z^{\wedge} n)$ 
  using has_laurent_expansion_imp_asymp_equiv_0[OF F']
  by (simp add: fls_subdegree_fls_to_fps_n_def)
moreover have  $f\ 0 = \text{fps\_nth } F\ n * 0^{\wedge} n$ 
  using F by (auto simp: n_def has_fps_expansion_to_laurent_power_0_left)
ultimately show ?thesis
  by (auto simp: asymp_equiv_nhds_iff)
qed

```

```

lemma has_fps_expansion_imp_tendsto_0:
  fixes  $f :: \text{complex} \Rightarrow \text{complex}$ 
  assumes  $f$  has_fps_expansion F
  shows  $(f \longrightarrow \text{fps\_nth } F\ 0)$   $(\text{nhds } 0)$ 
proof (rule asymp_equiv_tendsto_transfer)
  show  $(\lambda z. \text{fps\_nth } F\ (\text{subdegree } F) * z^{\wedge} \text{subdegree } F) \sim[\text{nhds } 0]$   $f$ 
    by (rule asymp_equiv_symI, rule has_fps_expansion_imp_asymp_equiv_0)
fact
  have  $(\lambda z. F\ \$\ \text{subdegree } F * z^{\wedge} \text{subdegree } F) \longrightarrow F\ \$\ 0$   $(\text{at } 0)$ 
    by (rule tendsto_eq_intros refl | simp) + (auto simp: power_0_left)
  thus  $(\lambda z. F\ \$\ \text{subdegree } F * z^{\wedge} \text{subdegree } F) \longrightarrow F\ \$\ 0$   $(\text{nhds } 0)$ 
    by (auto simp: tendsto_nhds_iff power_0_left)
qed

```

```

lemma has_fps_expansion_imp_0_eq_fps_nth_0:
  assumes  $f$  has_fps_expansion F
  shows  $f\ 0 = \text{fps\_nth } F\ 0$ 
proof -
  have eventually  $(\lambda x. f\ x = \text{eval\_fps } F\ x)$   $(\text{nhds } 0)$ 
    using assms by (auto simp: has_fps_expansion_def eq_commute)
  then obtain  $A$  where open A  $0 \in A \forall x \in A. f\ x = \text{eval\_fps } F\ x$ 
    unfolding eventually_nhds by blast
  hence  $f\ 0 = \text{eval\_fps } F\ 0$ 
    by blast
  thus ?thesis
    by (simp add: eval_fps_at_0)
qed

```

```

lemma fls_nth_compose_aux:
  assumes  $f$  has_fps_expansion F
  assumes  $G: g$  has_fps_expansion G  $\text{fps\_nth } G\ 0 = 0$   $\text{fps\_deriv } G \neq 0$ 
  assumes  $(f \circ g)$  has_laurent_expansion H
  shows  $\text{fls\_nth } H\ (\text{int } n) = \text{fps\_nth } (\text{fps\_compose } F\ G)\ n$ 
  using assms(1,5)
proof (induction n arbitrary: f F H rule: less_induct)
  case  $(\text{less } n\ f\ F\ H)$ 
  have [simp]:  $g\ 0 = 0$ 
    using has_fps_expansion_imp_0_eq_fps_nth_0[OF G(1)] G(2) by simp
  have ana_f:  $f$  analytic_on  $\{0\}$ 

```

```

    using less.premis by (meson has_fps_expansion_imp_analytic_0)
  have ana_g: g analytic_on {0}
    using G by (meson has_fps_expansion_imp_analytic_0)
  have (f ∘ g) has_laurent_expansion_fps_to_fls (fps_expansion (f ∘ g) 0)
    by (rule analytic_at_imp_has_fps_expansion_0 analytic_intros has_laurent_expansion_fps
        analytic_on_compose_gen ana_f ana_g)+ auto
  with less.premis have H = fps_to_fls (fps_expansion (f ∘ g) 0)
    using has_laurent_expansion_unique by blast
  also have fls_subdegree ... ≥ 0
    by (simp add: fls_subdegree_fls_to_fps)
  finally have subdeg: fls_subdegree H ≥ 0 .

show ?case
proof (cases n = 0)
  case [simp]: True
    have lim_g: g -0→ 0
      using has_laurent_expansion_imp_tendsto_0[of g fps_to_fls G] G
      by (auto simp: fls_subdegree_fls_to_fps_gt0 has_fps_expansion_to_laurent)
    have lim_f: (f ⟶ fps_nth F 0) (nhds 0)
      by (intro has_fps_expansion_imp_tendsto_0 less.premis)
    have (λx. f (g x)) -0→ fps_nth F 0
      by (rule filterlim_compose[OF lim_f lim_g])
    moreover have (f ∘ g) -0→ fls_nth H 0
      by (intro has_laurent_expansion_imp_tendsto_0 less.premis subdeg)
    ultimately have fps_nth F 0 = fls_nth H 0
      using tendsto_unique by (force simp: o_def)
    thus ?thesis
      by simp
  next
  case n: False
    define GH where GH = (fls_deriv H / fls_deriv (fps_to_fls G))
    define GH' where GH' = fls_regpart GH

    have (λx. deriv (f ∘ g) x / deriv g x) has_laurent_expansion
      fls_deriv H / fls_deriv (fps_to_fls G)
      by (intro laurent_expansion_intros less.premis has_laurent_expansion_fps[of
        _ G] G)
    also have ?this ⟷ (deriv f ∘ g) has_laurent_expansion fls_deriv H /
      fls_deriv (fps_to_fls G)
    proof (rule has_laurent_expansion_cong)
      from ana_f obtain r1 where r1: r1 > 0 f holomorphic_on ball 0 r1
        unfolding analytic_on_def by blast
      from ana_g obtain r2 where r2: r2 > 0 g holomorphic_on ball 0 r2
        unfolding analytic_on_def by blast
      have lim_g: g -0→ 0
        using has_laurent_expansion_imp_tendsto_0[of g fps_to_fls G] G
      by (auto simp: fls_subdegree_fls_to_fps_gt0 has_fps_expansion_to_laurent)
      moreover have open (ball 0 r1) 0 ∈ ball 0 r1
        using r1 by auto
    end

```

```

ultimately have eventually ( $\lambda x. g x \in \text{ball } 0 r1$ ) (at 0)
  unfolding tendsto_def by blast
moreover have eventually ( $\lambda x. \text{deriv } g x \neq 0$ ) (at 0)
using G fps_to_fls_eq_0_iff has_fps_expansion_deriv has_fps_expansion_to_laurent
  has_laurent_expansion_nonzero_imp_eventually_nonzero by blast
moreover have eventually ( $\lambda x. x \in \text{ball } 0 (\min r1 r2) - \{0\}$ ) (at 0)
  by (intro eventually_at_in_open) (use r1 r2 in auto)
ultimately show eventually ( $\lambda x. \text{deriv } (f \circ g) x / \text{deriv } g x = (\text{deriv } f \circ g)$ 
x) (at 0)
  proof eventually_elim
    case (elim x)
    thus ?case using r1 r2
      by (subst deriv_chain)
        (auto simp: field_simps holomorphic_on_def at_within_open[of _ ball
_ _])
  qed
qed auto
finally have GH: (deriv f  $\circ$  g) has_laurent_expansion GH
  unfolding GH_def .

have (deriv f  $\circ$  g) has_laurent_expansion fps_to_fls (fps_expansion (deriv f
 $\circ$  g) 0)
  by (rule analytic_at_imp_has_fps_expansion_0 analytic_intros has_laurent_expansion_fps
analytic_on_compose_gen ana_f ana_g)+ auto
with GH have GH = fps_to_fls (fps_expansion (deriv f  $\circ$  g) 0)
  using has_laurent_expansion_unique by blast
also have fls_subdegree ...  $\geq 0$ 
  by (simp add: fls_subdegree_fls_to_fps)
finally have subdeg': fls_subdegree GH  $\geq 0$  .

have deriv f has_fps_expansion fps_deriv F
  by (intro fps_expansion_intros less.prem)
from this and GH have IH: fls_nth GH (int k) = fps_nth (fps_compose
(fps_deriv F) G) k
  if  $k < n$  for k
  by (intro less.IH that)

have fps_nth (fps_compose (fps_deriv F) G) n = ( $\sum i=0..n. \text{of\_nat } (\text{Suc } i)
* F \$ \text{Suc } i * G ^ i \$ n$ )
  by (simp add: fps_compose_nth)

have fps_nth (fps_compose F G) n =
  fps_nth (fps_deriv (fps_compose F G)) (n - 1) / of_nat n
  using n by (cases n) (auto simp del: of_nat_Suc)
also have fps_deriv (fps_compose F G) = fps_compose (fps_deriv F) G *
fps_deriv G
  using G by (subst fps_compose_deriv) auto
also have fps_nth ... (n - 1) = ( $\sum i=0..n-1. (\text{fps\_deriv } F \text{ oo } G) \$ i *
\text{fps\_deriv } G \$ (n - 1 - i)$ )

```

```

unfolding fps_mult_nth ..
also have ... = ( $\sum i=0..n-1$ . fps_nth GH' i * of_nat (n - i) * G $ (n -
i))
using n by (intro sum.cong) (auto simp: IH Suc_diff_Suc GH'_def)
also have ... = ( $\sum i=0..n$ . fps_nth GH' i * of_nat (n - i) * G $ (n - i))
by (intro sum.mono_neutral_left) auto
also have ... = fps_nth (GH' * Abs_fps ( $\lambda i$ . of_nat i * fps_nth G i)) n
by (simp add: fps_mult_nth mult_ac)
also have Abs_fps ( $\lambda i$ . of_nat i * fps_nth G i) = fps_X * fps_deriv G
by (simp add: fps_mult_fps_X_deriv_shift)
also have fps_nth (GH' * (fps_X * fps_deriv G)) n =
  fps_nth (fps_to_fls (GH' * (fps_X * fps_deriv G))) (int n)
by simp
also have fps_to_fls (GH' * (fps_X * fps_deriv G)) =
  GH * fps_to_fls (fps_deriv G) * fls_X
using subdeg' by (simp add: mult_ac fls_times_fps_to_fls GH'_def)
also have GH * fps_to_fls (fps_deriv G) = fls_deriv H
unfolding GH_def using G by (simp add: fls_deriv_fps_to_fls)
also have fls_deriv H * fls_X = fls_shift (-1) (fls_deriv H)
using fls_X_times_conv_shift(2) by blast
finally show ?thesis
using n by simp
qed
qed

```

lemma has_fps_expansion_compose [fps_expansion_intros]:

```

fixes f g :: complex  $\Rightarrow$  complex
assumes F: f has_fps_expansion F
assumes G: g has_fps_expansion G fps_nth G 0 = 0
shows (f  $\circ$  g) has_fps_expansion fps_compose F G
proof (cases fps_deriv G = 0)
case False
have [simp]: g 0 = 0
using has_fps_expansion_imp_0_eq_fps_nth_0[OF G(1)] G(2) False by
simp
have ana_f: f analytic_on {0}
using F by (meson has_fps_expansion_imp_analytic_0)
have ana_g: g analytic_on {0}
using G by (meson has_fps_expansion_imp_analytic_0)
have fg: (f  $\circ$  g) has_fps_expansion fps_expansion (f  $\circ$  g) 0
by (rule analytic_at_imp_has_fps_expansion_0 analytic_intros
analytic_on_compose_gen ana_f ana_g)+ auto

have fls_nth (fps_to_fls (fps_expansion (f  $\circ$  g) 0)) (int n) = fps_nth (fps_compose
F G) n for n
by (rule fls_nth_compose_aux has_laurent_expansion_fps F G False fg)+
hence fps_expansion (f  $\circ$  g) 0 = fps_compose F G
by (simp add: fps_eq_iff)
thus ?thesis using fg

```

```

  by simp
next
case True
have [simp]: f 0 = fps_nth F 0
  using F by (auto dest: has_fps_expansion_imp_0_eq_fps_nth_0)
from True have fps_nth G n = 0 for n
  using G(2) by (cases n) (auto simp del: of_nat_Suc)
hence [simp]: G = 0
  by (auto simp: fps_eq_iff)
have (λ_. f 0) has_fps_expansion fps_const (f 0)
  by (intro fps_expansion_intros)
also have eventually (λx. g x = 0) (nhds 0)
  using G by (auto simp: has_fps_expansion_def)
hence (λ_. f 0) has_fps_expansion fps_const (f 0) ↔ (f ∘ g) has_fps_expansion
fps_const (f 0)
  by (intro has_fps_expansion_cong) (auto elim!: eventually_mono)
thus ?thesis
  by simp
qed

```

hide_const (open) fls_compose_fps

definition fls_compose_fps :: 'a :: field fls ⇒ 'a fps ⇒ 'a fls **where**
 fls_compose_fps F G =
 fps_to_fls (fps_compose (fls_base_factor_to_fps F) G) * fps_to_fls G powi
 fls_subdegree F

lemma fls_compose_of_nat [simp]: fls_compose (of_nat n :: 'a :: comm_ring_1
 fps) H = of_nat n
and fls_compose_of_int [simp]: fls_compose (of_int i) H = of_int i
unfolding fls_of_nat [symmetric] fls_of_int [symmetric] numeral_fps_const
by (rule fps_const_compose)+

lemmas [simp] = fps_to_fls_of_nat fps_to_fls_of_int

lemma fls_compose_fps_0 [simp]: fls_compose_fps 0 H = 0
and fls_compose_fps_1 [simp]: fls_compose_fps 1 H = 1
and fls_compose_fps_const [simp]: fls_compose_fps (fls_const c) H = fls_const
 c
and fls_compose_fps_of_nat [simp]: fls_compose_fps (of_nat n) H = of_nat
 n
and fls_compose_fps_of_int [simp]: fls_compose_fps (of_int i) H = of_int i
and fls_compose_fps_X [simp]: fls_compose_fps fls_X F = fps_to_fls F
by (simp_all add: fls_compose_fps_def)

lemma fls_compose_fps_0_right:
 fls_compose_fps F 0 = (if fls_subdegree F ≥ 0 then fls_const (fls_nth F 0) else
 0)
by (cases fls_subdegree F = 0) (simp_all add: fls_compose_fps_def)

lemma *fls_compose_fps_shift*:

assumes $H \neq 0$
shows $\text{fls_compose_fps } (\text{fls_shift } n \ F) \ H = \text{fls_compose_fps } F \ H * \text{fps_to_fls } H \ \text{powi } (-n)$
proof (cases $F = 0$)
case *False*
thus ?thesis
using *assms* **by** (simp add: *fls_compose_fps_def power_int_diff power_int_minus field_simps*)
qed *auto*

lemma *fls_compose_fps_to_fls [simp]*:

assumes [simp]: $G \neq 0 \ \text{fps_nth } G \ 0 = 0$
shows $\text{fls_compose_fps } (\text{fps_to_fls } F) \ G = \text{fps_to_fls } (\text{fps_compose } F \ G)$
proof (cases $F = 0$)
case *False*
define n **where** $n = \text{subdegree } F$
define F' **where** $F' = \text{fps_shift } n \ F$
have [simp]: $F' \neq 0 \ \text{subdegree } F' = 0$
using *False* **by** (auto simp: *F'_def n_def*)
have F_eq : $F = F' * \text{fps_X } ^n$
unfolding *F'_def n_def* **using** *subdegree_decompose* **by** *blast*
have $\text{fls_compose_fps } (\text{fps_to_fls } F) \ G =$
 $\text{fps_to_fls } (\text{fps_shift } n \ (\text{fls_regpart } (\text{fps_to_fls } F' * \text{fls_X_intpow } (\text{int } n))) \ \text{oo } G) * \text{fps_to_fls } (G ^n)$
unfolding *F_eq fls_compose_fps_def*
by (simp add: *fls_times_fps_to_fls fls_X_power_conv_shift_1 power_int_add fls_subdegree_fls_to_fps fps_to_fls_power fls_regpart_shift_conv_fps_shift flip: fls_times_both_shifted simp*)
also have $\text{fps_to_fls } F' * \text{fls_X_intpow } (\text{int } n) = \text{fps_to_fls } F$
by (simp add: *F_eq fls_times_fps_to_fls fps_to_fls_power fls_X_power_conv_shift_1*)
also have $\text{fps_to_fls } (\text{fps_shift } n \ (\text{fls_regpart } (\text{fps_to_fls } F))) \ \text{oo } G * \text{fps_to_fls } (G ^n) =$
 $\text{fps_to_fls } ((\text{fps_shift } n \ (\text{fls_regpart } (\text{fps_to_fls } F))) * \text{fps_X } ^n) \ \text{oo } G$
by (simp add: *fls_times_fps_to_fls flip: fps_compose_power add: fps_compose_mult_distrib*)
also have $\text{fps_shift } n \ (\text{fls_regpart } (\text{fps_to_fls } F)) * \text{fps_X } ^n = F$
by (simp add: *F_eq*)
finally show ?thesis .
qed (auto simp: *fls_compose_fps_def*)

lemma *fls_compose_fps_mult*:

assumes [simp]: $H \neq 0 \ \text{fps_nth } H \ 0 = 0$
shows $\text{fls_compose_fps } (F * G) \ H = \text{fls_compose_fps } F \ H * \text{fls_compose_fps } G \ H$
using *assms*
proof (cases $F * G = 0$)
case *False*

```

hence [simp]: F ≠ 0 G ≠ 0
  by auto
define n m where n = fls_subdegree F m = fls_subdegree G
define F' where F' = fls_regpart (fls_shift n F)
define G' where G' = fls_regpart (fls_shift m G)
have F_eq: F = fls_shift (-n) (fps_to_fls F') and G_eq: G = fls_shift (-m)
(fps_to_fls G')
  by (simp_all add: F'_def G'_def n_m_def)
have fls_compose_fps (F * G) H = fls_compose_fps (fls_shift (-(n + m))
(fps_to_fls (F' * G'))) H
  by (simp add: fls_times_fps_to_fls F_eq G_eq fls_shifted_times_simps)
also have ... = fps_to_fls ((F' oo H) * (G' oo H)) * fps_to_fls H powi (m +
n)
  by (simp add: fls_compose_fps_shift fps_compose_mult_distrib)
also have ... = fls_compose_fps F H * fls_compose_fps G H
  by (simp add: F_eq G_eq fls_compose_fps_shift fls_times_fps_to_fls power_int_add)
finally show ?thesis .
qed auto

```

lemma *fls_compose_fps_power*:

```

assumes [simp]: G ≠ 0 fps_nth G 0 = 0
shows fls_compose_fps (F ^ n) G = fls_compose_fps F G ^ n
  by (induction n) (auto simp: fls_compose_fps_mult)

```

lemma *fls_compose_fps_add*:

```

assumes [simp]: H ≠ 0 fps_nth H 0 = 0
shows fls_compose_fps (F + G) H = fls_compose_fps F H + fls_compose_fps
G H
proof (cases F = 0 ∨ G = 0)
  case False
  hence [simp]: F ≠ 0 G ≠ 0
    by auto
  define n where n = min (fls_subdegree F) (fls_subdegree G)
  define F' where F' = fls_regpart (fls_shift n F)
  define G' where G' = fls_regpart (fls_shift n G)
  have F_eq: F = fls_shift (-n) (fps_to_fls F') and G_eq: G = fls_shift (-n)
(fps_to_fls G')
    unfolding n_def by (simp_all add: F'_def G'_def n_def)
  have F + G = fls_shift (-n) (fps_to_fls (F' + G'))
    by (simp add: F_eq G_eq)
  also have fls_compose_fps ... H = fls_compose_fps (fps_to_fls (F' + G')) H
* fps_to_fls H powi n
    by (subst fls_compose_fps_shift) auto
  also have ... = fps_to_fls (fps_compose (F' + G') H) * fps_to_fls H powi n
    by (subst fls_compose_fps_to_fls) auto
  also have ... = fls_compose_fps F H + fls_compose_fps G H
    by (simp add: F_eq G_eq fls_compose_fps_shift fps_compose_add_distrib
algebra_simps)
  finally show ?thesis .

```

qed *auto*

lemma *fls_compose_fps_uminus* [*simp*]: $\text{fls_compose_fps } (-F) H = -\text{fls_compose_fps } F H$

by (*simp add: fls_compose_fps_def fps_compose_uminus*)

lemma *fls_compose_fps_diff*:

assumes [*simp*]: $H \neq 0$ $\text{fps_nth } H 0 = 0$

shows $\text{fls_compose_fps } (F - G) H = \text{fls_compose_fps } F H - \text{fls_compose_fps } G H$

using *fls_compose_fps_add*[*of H F -G*] **by** *simp*

lemma *fps_compose_eq_0_iff*:

fixes $F G :: 'a :: \text{idom } \text{fps}$

assumes $\text{fps_nth } G 0 = 0$

shows $\text{fps_compose } F G = 0 \iff F = 0 \vee (G = 0 \wedge \text{fps_nth } F 0 = 0)$

proof *safe*

assume $*$: $\text{fps_compose } F G = 0$ $F \neq 0$

have $\text{fps_nth } (\text{fps_compose } F G) 0 = \text{fps_nth } F 0$

by *simp*

also have $\text{fps_compose } F G = 0$

by (*simp add: **)

finally show $\text{fps_nth } F 0 = 0$

by *simp*

show $G = 0$

proof (*rule ccontr*)

assume $G \neq 0$

hence $\text{subdegree } G > 0$ **using** *assms*

using *subdegree_eq_0_iff* **by** *blast*

define N **where** $N = \text{subdegree } F * \text{subdegree } G$

have $\text{fps_nth } (\text{fps_compose } F G) N = (\sum i = 0..N. \text{fps_nth } F i * \text{fps_nth } (G \wedge i) N)$

unfolding *fps_compose_def* **by** (*simp add: N_def*)

also have $\dots = (\sum i \in \{\text{subdegree } F\}. \text{fps_nth } F i * \text{fps_nth } (G \wedge i) N)$

proof (*intro sum.mono_neutral_right ballI*)

fix i **assume** $i \in \{0..N\} - \{\text{subdegree } F\}$

show $\text{fps_nth } F i * \text{fps_nth } (G \wedge i) N = 0$

proof (*cases i subdegree F rule: linorder_cases*)

assume $i > \text{subdegree } F$

hence $\text{fps_nth } (G \wedge i) N = 0$

using $i < \text{subdegree } G >$ **by** (*intro fps_pow_nth_below_subdegree*) (*auto simp: N_def*)

thus *?thesis* **by** *simp*

qed (*use i in <auto simp: N_def>*)

qed (*use <subdegree G > 0 in <auto simp: N_def>*)

also have $\dots = \text{fps_nth } F (\text{subdegree } F) * \text{fps_nth } (G \wedge \text{subdegree } F) N$

by *simp*

also have $\dots \neq 0$

using $\langle G \neq 0 \rangle \langle F \neq 0 \rangle$ **by** (*auto simp: N_def*)

finally show *False* using * by auto
qed
qed auto

lemma *fls_compose_fps_eq_0_iff*:
assumes $H \neq 0$ *fps_nth* $H\ 0 = 0$
shows $\text{fls_compose_fps}\ F\ H = 0 \iff F = 0$
using *assms fls_base_factor_to_fps_nonzero*[of *F*]
by (cases $F = 0$) (auto simp: *fls_compose_fps_def fls_compose_eq_0_iff*)

lemma *fls_compose_fps_inverse*:
assumes [*simp*]: $H \neq 0$ *fps_nth* $H\ 0 = 0$
shows $\text{fls_compose_fps}\ (\text{inverse}\ F)\ H = \text{inverse}\ (\text{fls_compose_fps}\ F\ H)$
proof (cases $F = 0$)
case *False*
have $\text{fls_compose_fps}\ (\text{inverse}\ F)\ H * \text{fls_compose_fps}\ F\ H =$
 $\text{fls_compose_fps}\ (\text{inverse}\ F * F)\ H$
by (*subst fls_compose_fps_mult*) auto
also have $\text{inverse}\ F * F = 1$
using *False* by *simp*
finally show ?*thesis*
using *False* by (*simp add: field_simps fls_compose_fps_eq_0_iff*)
qed auto

lemma *fls_compose_fps_divide*:
assumes [*simp*]: $H \neq 0$ *fps_nth* $H\ 0 = 0$
shows $\text{fls_compose_fps}\ (F / G)\ H = \text{fls_compose_fps}\ F\ H / \text{fls_compose_fps}\ G\ H$
using *fls_compose_fps_mult*[of $H\ F\ \text{inverse}\ G$] *fls_compose_fps_inverse*[of $H\ G$]
by (*simp add: field_simps*)

lemma *fls_compose_fps_powi*:
assumes [*simp*]: $H \neq 0$ *fps_nth* $H\ 0 = 0$
shows $\text{fls_compose_fps}\ (F\ \text{powi}\ n)\ H = \text{fls_compose_fps}\ F\ H\ \text{powi}\ n$
by (*simp add: power_int_def fls_compose_fps_power fls_compose_fps_inverse*)

lemma *fls_compose_fps_assoc*:
assumes [*simp*]: $G \neq 0$ *fps_nth* $G\ 0 = 0$ $H \neq 0$ *fps_nth* $H\ 0 = 0$
shows $\text{fls_compose_fps}\ (\text{fls_compose_fps}\ F\ G)\ H = \text{fls_compose_fps}\ F\ (\text{fps_compose}\ G\ H)$
proof (cases $F = 0$)
case [*simp*]: *False*
define *n* where $n = \text{fls_subdegree}\ F$
define *F'* where $F' = \text{fls_regpart}\ (\text{fls_shift}\ n\ F)$
have *F_eq*: $F = \text{fls_shift}\ (-n)\ (\text{fps_to_fls}\ F')$
by (*simp add: F'_def n_def*)
show ?*thesis*
by (*simp add: F_eq fls_compose_fps_shift fls_compose_fps_mult fls_compose_fps_powi*)

```

      fps_compose_eq_0_iff fps_compose_assoc)
qed auto

lemma subdegree_pos_iff: subdegree F > 0  $\longleftrightarrow$  F  $\neq$  0  $\wedge$  fps_nth F 0 = 0
  using subdegree_eq_0_iff[of F] by auto

lemma has_fps_expansion_fps_to_fls:
  assumes f has_laurent_expansion_fps_to_fls F
  shows ( $\lambda z$ . if z = 0 then fps_nth F 0 else f z) has_fps_expansion F
  (is ?f' has_fps_expansion _)
proof -
  have f has_laurent_expansion_fps_to_fls F  $\longleftrightarrow$  ?f' has_laurent_expansion
  fps_to_fls F
  by (intro has_laurent_expansion_cong) (auto simp: eventually_at_filter)
  with assms show ?thesis
  by (auto simp: has_fps_expansion_to_laurent)
qed

lemma has_laurent_expansion_compose [laurent_expansion_intros]:
  fixes f g :: complex  $\Rightarrow$  complex
  assumes F: f has_laurent_expansion F
  assumes G: g has_laurent_expansion_fps_to_fls G fps_nth G 0 = 0 G  $\neq$  0
  shows (f  $\circ$  g) has_laurent_expansion_fls_compose_fps F G
proof -
  from assms have lim_g: g  $\rightarrow$  0
  by (subst tendsto_0_subdegree_iff_0[OF G(1)])
  (auto simp: fls_subdegree_fls_to_fps_subdegree_pos_iff)
  have ev1: eventually ( $\lambda z$ . g z  $\neq$  0) (at 0)
  using  $\langle$ G  $\neq$  0 $\rangle$  G(1) fps_to_fls_eq_0_iff has_laurent_expansion_fps
  has_laurent_expansion_nonzero_imp_eventually_nonzero by blast
  moreover have eventually ( $\lambda z$ . z  $\neq$  0) (at (0 :: complex))
  by (auto simp: eventually_at_filter)
  ultimately have ev: eventually ( $\lambda z$ . z  $\neq$  0  $\wedge$  g z  $\neq$  0) (at 0)
  by eventually_elim blast
  from ev1 and lim_g have lim_g': filterlim g (at 0) (at 0)
  by (auto simp: filterlim_at)
  define g' where g' = ( $\lambda z$ . if z = 0 then fps_nth G 0 else g z)

  show ?thesis
proof (cases F = 0)
  assume [simp]: F = 0
  have eventually ( $\lambda z$ . f z = 0) (at 0)
  using F by (auto simp: has_laurent_expansion_def)
  hence eventually ( $\lambda z$ . f (g z) = 0) (at 0)
  using lim_g' by (rule eventually_compose_filterlim)
  thus ?thesis
  by (auto simp: has_laurent_expansion_def)
next

```

```

assume [simp]:  $F \neq 0$ 
define  $n$  where  $n = \text{fls\_subdegree } F$ 
define  $f'$  where
   $f' = (\lambda z. \text{if } z = 0 \text{ then } \text{fps\_nth } (\text{fls\_base\_factor\_to\_fps } F) 0 \text{ else } f z * z \text{ powi } -n)$ 
have  $((\lambda z. (f' \circ g') z * g z \text{ powi } n)) \text{ has\_laurent\_expansion } \text{fls\_compose\_fps } F G$ 
unfolding  $f'\_def\ n\_def\ \text{fls\_compose\_fps\_def}\ g'\_def$ 
by (intro  $\text{fps\_expansion\_intros}\ \text{laurent\_expansion\_intros}\ \text{has\_fps\_expansion\_fps\_to\_fls}\ \text{has\_fps\_expansion\_fls\_base\_factor\_to\_fps}\ \text{assms}\ \text{has\_laurent\_expansion\_fps}$ )
also have  $?this \longleftrightarrow ?thesis$ 
by (intro  $\text{has\_laurent\_expansion\_cong}\ \text{eventually\_mono}[OF\ \text{ev}]$ )
  (auto simp:  $f'\_def\ \text{power\_int\_minus}\ g'\_def$ )
finally show  $?thesis$  .
qed

```

```

lemma  $\text{has\_laurent\_expansion\_fls\_X\_inv}$  [ $\text{laurent\_expansion\_intros}$ ]:
   $\text{inverse}\ \text{has\_laurent\_expansion\_fls\_X\_inv}$ 
using  $\text{has\_laurent\_expansion\_inverse}[OF\ \text{has\_laurent\_expansion\_fps\_X}]$ 
by (simp add:  $\text{fls\_inverse\_X}$ )

```

```

lemma  $\text{fls\_X\_power\_int}$  [simp]:  $\text{fls\_X}\ \text{powi } n = (\text{fls\_X\_intpow } n :: 'a :: \text{division\_ring}\ \text{fls})$ 
by (auto simp:  $\text{power\_int\_def}\ \text{fls\_X\_power\_conv\_shift\_1}\ \text{fls\_inverse\_X}\ \text{fls\_inverse\_shift}\ \text{simp\_flip:}\ \text{fls\_inverse\_X\_power}$ )

```

```

lemma  $\text{fls\_const\_power\_int}$ :  $\text{fls\_const } (c \text{ powi } n) = \text{fls\_const } (c :: 'a :: \text{division\_ring}) \text{ powi } n$ 
by (auto simp:  $\text{power\_int\_def}\ \text{fls\_const\_power}\ \text{fls\_inverse\_const}$ )

```

```

lemma  $\text{fls\_nth\_fls\_compose\_fps\_linear}$ :
  fixes  $c :: 'a :: \text{field}$ 
  assumes [simp]:  $c \neq 0$ 
  shows  $\text{fls\_nth } (\text{fls\_compose\_fps } F (\text{fps\_const } c * \text{fps\_X}))\ n = \text{fls\_nth } F\ n * c \text{ powi } n$ 
proof -
  {
    assume  $*$ :  $n \geq \text{fls\_subdegree } F$ 
    hence  $c \wedge^{\text{nat}} (n - \text{fls\_subdegree } F) = c \text{ powi int } (\text{nat } (n - \text{fls\_subdegree } F))$ 
      by (simp add:  $\text{power\_int\_def}$ )
    also have  $\dots * c \text{ powi } \text{fls\_subdegree } F = c \text{ powi } (\text{int } (\text{nat } (n - \text{fls\_subdegree } F)) + \text{fls\_subdegree } F)$ 
      using  $*$  by (subst  $\text{power\_int\_add}$ ) auto
    also have  $\dots = c \text{ powi } n$ 
      using  $*$  by simp
    finally have  $c \wedge^{\text{nat}} (n - \text{fls\_subdegree } F) * c \text{ powi } \text{fls\_subdegree } F = c \text{ powi } n$  .
  }
qed

```

thus ?thesis
by (simp add: fls_compose_fps_def fps_compose_linear fls_times_fps_to_fls
power_int_mult_distrib
fls_shifted_times_simps
flip: fls_const_power_int)
qed

lemma zorder_times_analytic:

assumes f analytic_on $\{z\}$ g analytic_on $\{z\}$
assumes eventually $(\lambda z. f z * g z \neq 0)$ (at z)
shows $zorder (\lambda z. f z * g z) z = zorder f z + zorder g z$
proof –
have *: $(\lambda w. f (z + w))$ has_fps_expansion fps_expansion $f z$
 $(\lambda w. g (z + w))$ has_fps_expansion fps_expansion $g z$
 $(\lambda w. f (z + w) * g (z + w))$ has_fps_expansion fps_expansion $f z * g z$
by (intro fps_expansion_intros analytic_at_imp_has_fps_expansion assms)+
have [simp]: fps_expansion $f z \neq 0$
proof
assume fps_expansion $f z = 0$
hence eventually $(\lambda z. f z * g z = 0)$ (at z) **using** *(1)
by (auto simp: has_fps_expansion_0_iff_nhds_to_0' eventually_filtermap
eventually_at_filter
elim: eventually_mono)
with assms(3) **have** eventually $(\lambda z. False)$ (at z)
by eventually_elim auto
thus False **by** simp
qed
have [simp]: fps_expansion $g z \neq 0$
proof
assume fps_expansion $g z = 0$
hence eventually $(\lambda z. f z * g z = 0)$ (at z) **using** *(2)
by (auto simp: has_fps_expansion_0_iff_nhds_to_0' eventually_filtermap
eventually_at_filter
elim: eventually_mono)
with assms(3) **have** eventually $(\lambda z. False)$ (at z)
by eventually_elim auto
thus False **by** simp
qed
from *[THEN has_fps_expansion_zorder] **show** ?thesis
by auto
qed

lemma analytic_on_prod [analytic_intros]:

assumes $\bigwedge x. x \in A \implies f x$ analytic_on B
shows $(\lambda z. \prod x \in A. f x z)$ analytic_on B
using assms **by** (induction A rule: infinite_finite_induct) (auto intro!: analytic_intros)

lemma *zorder_const* [*simp*]: $c \neq 0 \implies \text{zorder } (\lambda_. c) z = 0$
by (*intro zorder_eqI*[**where** $s = UNIV$]) *auto*

lemma *zorder_prod_analytic*:

assumes $\bigwedge x. x \in A \implies f x \text{ analytic_on } \{z\}$

assumes *eventually* $(\lambda z. (\prod x \in A. f x z) \neq 0)$ (*at z*)

shows $\text{zorder } (\lambda z. \prod x \in A. f x z) z = (\sum x \in A. \text{zorder } (f x) z)$

using *assms*

proof (*induction A rule: infinite_finite_induct*)

case (*insert x A*)

have $\text{zorder } (\lambda z. f x z * (\prod x \in A. f x z)) z = \text{zorder } (f x) z + \text{zorder } (\lambda z. \prod x \in A. f x z) z$

using *insert.premis insert.hyps* **by** (*intro zorder_times_analytic analytic_intros*)

auto

also have $\text{zorder } (\lambda z. \prod x \in A. f x z) z = (\sum x \in A. \text{zorder } (f x) z)$

using *insert.premis insert.hyps* **by** (*intro insert.IH*) (*auto elim!: eventually_mono*)

finally show *?case* **using** *insert*

by *simp*

qed *auto*

lemma *zorder_eq_0I*:

assumes *g analytic_on* $\{z\}$ $g z \neq 0$

shows $\text{zorder } g z = 0$

proof –

from *assms* **obtain** *r* **where** $r: r > 0$ *g holomorphic_on ball z r*

unfolding *analytic_on_def* **by** *blast*

thus *?thesis* **using** *assms*

by (*intro zorder_eqI*[*of ball z r _ g*]) *auto*

qed

lemma *zorder_pos_iff*:

assumes *f holomorphic_on A open A z* $z \in A$ *frequently* $(\lambda z. f z \neq 0)$ (*at z*)

shows $\text{zorder } f z > 0 \iff f z = 0$

proof –

have *f analytic_on* $\{z\}$

using *assms analytic_at* **by** *blast*

hence $*$: $(\lambda w. f (z + w)) \text{ has_fps_expansion } \text{fps_expansion } f z$

using *analytic_at_imp_has_fps_expansion* **by** *blast*

have *nz: fps_expansion f z* $\neq 0$

proof

assume *fps_expansion f z* $= 0$

hence *eventually* $(\lambda z. f z = 0)$ (*nhds z*)

using $*$ **by** (*auto simp: has_fps_expansion_def nhds_to_0' eventually_filtermap add_ac*)

hence *eventually* $(\lambda z. f z = 0)$ (*at z*)

by (*auto simp: eventually_at_filter elim: eventually_mono*)

with *assms* **show** *False*

by (*auto simp: frequently_def*)

qed

```

from has_fps_expansion_zorder[OF * this] have eq: zorder f z = int (subdegree
(fps_expansion f z))
  by auto
moreover have subdegree (fps_expansion f z) = 0  $\longleftrightarrow$  fps_nth (fps_expansion
f z) 0  $\neq$  0
  using nz by (auto simp: subdegree_eq_0_iff)
moreover have fps_nth (fps_expansion f z) 0 = f z
  by (auto simp: fps_expansion_def)
ultimately show ?thesis
  by auto
qed

```

lemma zorder_pos_iff':

```

assumes f analytic_on {z} frequently ( $\lambda z. f z \neq 0$ ) (at z)
shows zorder f z > 0  $\longleftrightarrow$  f z = 0
proof -
from assms(1) obtain A where A: open A {z}  $\subseteq$  A f holomorphic_on A
  unfolding analytic_on_holomorphic by auto
with zorder_pos_iff [OF A(3,1), of z] assms show ?thesis
  by auto
qed

```

lemma zorder_ge_0:

```

assumes f analytic_on {z} frequently ( $\lambda z. f z \neq 0$ ) (at z)
shows zorder f z  $\geq$  0
proof -
have *: ( $\lambda w. f (z + w)$ ) has_laurent_expansion_fps_to_fls (fps_expansion f z)
  using assms by (simp add: analytic_at_imp_has_fps_expansion has_laurent_expansion_fps)
from * assms(2) have fps_to_fls (fps_expansion f z)  $\neq$  0
  by (auto simp: has_laurent_expansion_def frequently_def at_to_0' eventu-
ally_filtermap add_ac)
with has_laurent_expansion_zorder[OF *] show ?thesis
  by (simp add: fls_subdegree_fls_to_fps)
qed

```

lemma zorder_eq_0_iff:

```

assumes f analytic_on {z} frequently ( $\lambda w. f w \neq 0$ ) (at z)
shows zorder f z = 0  $\longleftrightarrow$  f z  $\neq$  0
proof
  assume f z  $\neq$  0
  thus zorder f z = 0
    using assms zorder_eq_0I by blast
next
  assume zorder f z = 0
  thus f z  $\neq$  0
    using assms zorder_pos_iff' by fastforce
qed

```

lemma dist_mult_left:

$dist (a * b) (a * c :: 'a :: real_normed_field) = norm a * dist b c$
unfolding $dist_norm$ $right_diff_distrib$ $[symmetric]$ $norm_mult$ **by** $simp$

lemma $dist_mult_right$:

$dist (b * a) (c * a :: 'a :: real_normed_field) = norm a * dist b c$
using $dist_mult_left$ [of $a b c$] **by** ($simp$ $add: mult_ac$)

lemma $zorder_scale$:

assumes f $analytic_on$ $\{a * z\}$ $eventually$ $(\lambda w. f w \neq 0)$ $(at (a * z))$ $a \neq 0$
shows $zorder (\lambda w. f (a * w)) z = zorder f (a * z)$

proof –

from $assms(1)$ **obtain** r **where** $r: r > 0$ f $holomorphic_on$ $ball (a * z) r$
by ($auto$ $simp: analytic_on_def$)

have $*$: $open (ball (a * z) r)$ $connected (ball (a * z) r)$ $a * z \in ball (a * z) r$
using $r \langle a \neq 0 \rangle$ **by** ($auto$ $simp: dist_norm$)

from $assms(2)$ **have** $eventually$ $(\lambda w. f w \neq 0 \wedge w \in ball (a * z) r - \{a * z\})$
 $(at (a * z))$

using $\langle r > 0 \rangle$ **by** ($intro$ $eventually_conj$ $eventually_at_in_open$) $auto$

then obtain $z0$ **where** $f z0 \neq 0 \wedge z0 \in ball (a * z) r - \{a * z\}$

using $eventually_happens$ [of $_$ $at (a * z)$] **by** $force$

hence $**$: $\exists w \in ball (a * z) r. f w \neq 0$

by $blast$

define n **where** $n = nat (zorder f (a * z))$

obtain r' **where** r' :

$(if f (a * z) = 0$ $then$ $0 < zorder f (a * z)$ $else$ $zorder f (a * z) = 0)$

$r' > 0$ $cball (a * z) r' \subseteq ball (a * z) r$ $zor_poly f (a * z)$ $holomorphic_on$
 $cball (a * z) r'$

$\bigwedge w. w \in cball (a * z) r' \implies$

$f w = zor_poly f (a * z) w * (w - a * z) ^ n \wedge zor_poly f (a * z) w \neq 0$

unfolding n_def **using** $zorder_exist_zero$ [OF $r(2) * **$] **by** $blast$

show $?thesis$

proof ($rule$ $zorder_eqI$)

show $open (ball z (r' / norm a))$ $z \in ball z (r' / norm a)$

using $r \langle r' > 0 \rangle \langle a \neq 0 \rangle$ **by** $auto$

have $(*)$ $a \in ball z (r' / cmod a) \subseteq cball (a * z) r'$

proof $safe$

fix w **assume** $w \in ball z (r' / cmod a)$

thus $a * w \in cball (a * z) r'$

using $dist_mult_left$ [of $a z w$] $\langle a \neq 0 \rangle$ **by** ($auto$ $simp: divide_simps$ $mult_ac$)

qed

thus $(\lambda w. a ^ n * (zor_poly f (a * z) \circ (\lambda w. a * w)) w)$ $holomorphic_on$ $ball$
 $z (r' / norm a)$

using $\langle a \neq 0 \rangle$ **by** ($intro$ $holomorphic_on_compose_gen$ [OF $_$ $r'(4)$] $holomorphic_intros$) $auto$

show $a ^ n * (zor_poly f (a * z) \circ (\lambda w. a * w)) z \neq 0$

using $r' \langle a \neq 0 \rangle$ **by** $auto$

show $f (a * w) = a ^ n * (zor_poly f (a * z) \circ (*) a) w * (w - z) powi (zorder$

```

f (a * z))
  if w ∈ ball z (r' / norm a) w ≠ z for w
  proof -
    have f (a * w) = zor_poly f (a * z) (a * w) * (a * (w - z)) ^ n
      using that r'(5)[of a * w] dist_mult_left[of a z w] ‹a ≠ 0› unfolding
ring_distrib
    by (auto simp: divide_simps mult_ac)
    also have ... = a ^ n * zor_poly f (a * z) (a * w) * (w - z) ^ n
      by (subst power_mult_distrib) (auto simp: mult_ac)
    also have (w - z) ^ n = (w - z) powi of_nat n
      by simp
    also have of_nat n = zorder f (a * z)
      using r'(1) by (auto simp: n_def split: if_splits)
    finally show ?thesis
      unfolding o_def n_def .
  qed
qed
qed

```

```

lemma subdegree_fps_compose [simp]:
  fixes F G :: 'a :: idom fps
  assumes [simp]: fps_nth G 0 = 0
  shows subdegree (fps_compose F G) = subdegree F * subdegree G
proof (cases G = 0; cases F = 0)
  assume [simp]: G ≠ 0 F ≠ 0
  define m where m = subdegree F
  define F' where F' = fps_shift m F
  have F_eq: F = F' * fps_X ^ m
    unfolding F'_def by (simp add: fps_shift_times_fps_X_power m_def)
  have [simp]: F' ≠ 0
    using ‹F ≠ 0› unfolding F_eq by auto
  have subdegree (fps_compose F G) = subdegree (fps_compose F' G) + m *
subdegree G
    by (simp add: F_eq fps_compose_mult_distrib fps_compose_eq_0_iff flip:
fps_compose_power)
  also have subdegree (fps_compose F' G) = 0
    by (intro subdegree_eq_0) (auto simp: F'_def m_def)
  finally show ?thesis by (simp add: m_def)
qed auto

```

```

lemma fls_subdegree_power_int [simp]:
  fixes F :: 'a :: field fls
  shows fls_subdegree (F powi n) = n * fls_subdegree F
  by (auto simp: power_int_def fls_subdegree_pow)

```

```

lemma subdegree_fls_compose_fps [simp]:
  fixes G :: 'a :: field fps
  assumes [simp]: fps_nth G 0 = 0
  shows fls_subdegree (fls_compose_fps F G) = fls_subdegree F * subdegree G

```



```

proof (cases  $F = 0$ ; cases  $G = 0$ )
  assume [simp]:  $G \neq 0$   $F \neq 0$ 
  have nz1:  $\text{fls\_base\_factor\_to\_fps } F \neq 0$ 
    using  $\langle F \neq 0 \rangle$   $\text{fls\_base\_factor\_to\_fps\_nonzero}$  by blast
  show ?thesis
    unfolding  $\text{fls\_compose\_fps\_def}$  using nz1
    by (subst  $\text{fls\_subdegree\_mult}$ ) (simp_all add:  $\text{fps\_compose\_eq\_0\_iff\_fls\_subdegree\_fls\_to\_fps}$ )
qed (auto simp:  $\text{fls\_compose\_fps\_0\_right}$ )

```

lemma $\text{zorder_compose_aux}$:

```

assumes  $\text{isolated\_singularity\_at } f \ 0$   $\text{not\_essential } f \ 0$ 
assumes  $G$ :  $g$  has  $\text{fps\_expansion } G$   $G \neq 0$   $g \ 0 = 0$ 
assumes  $\text{eventually } (\lambda w. f \ w \neq 0)$  (at  $0$ )
shows  $\text{zorder } (f \circ g) \ 0 = \text{zorder } f \ 0 * \text{subdegree } G$ 
proof -
  obtain  $F$  where  $F$ :  $f$  has  $\text{laurent\_expansion } F$ 
    using  $\text{not\_essential\_has\_laurent\_expansion\_0}$  [ $OF$   $\text{assms}(1,2)$ ] by blast
  have [simp]:  $\text{fps\_nth } G \ 0 = 0$ 
    using  $G$   $\langle g \ 0 = 0 \rangle$  by (simp add:  $\text{has\_fps\_expansion\_imp\_0\_eq\_fps\_nth\_0}$ )
  note [simp] =  $\langle G \neq 0 \rangle$   $\langle g \ 0 = 0 \rangle$ 
  have [simp]:  $F \neq 0$ 
    using  $\text{has\_laurent\_expansion\_eventually\_nonzero\_iff}$  [ $of \ f \ 0 \ F$ ]  $F$   $\text{assms}$  by
  simp
  have  $FG$ :  $(f \circ g)$  has  $\text{laurent\_expansion } \text{fls\_compose\_fps } F \ G$ 
    by (intro  $\text{has\_laurent\_expansion\_compose}$   $\text{has\_laurent\_expansion\_fps } F \ G$ )
  auto

```

```

  have  $\text{zorder } (f \circ g) \ 0 = \text{fls\_subdegree } (\text{fls\_compose\_fps } F \ G)$ 
    using  $\text{has\_laurent\_expansion\_zorder\_0}$  [ $OF \ FG$ ] by (auto simp:  $\text{fls\_compose\_fps\_eq\_0\_iff}$ )
  also have  $\dots = \text{fls\_subdegree } F * \text{int } (\text{subdegree } G)$ 
    by simp
  also have  $\text{fls\_subdegree } F = \text{zorder } f \ 0$ 
    using  $\text{has\_laurent\_expansion\_zorder\_0}$  [ $OF \ F$ ] by auto
  finally show ?thesis .
qed

```

lemma zorder_compose :

```

assumes  $\text{isolated\_singularity\_at } f \ (g \ z)$   $\text{not\_essential } f \ (g \ z)$ 
assumes  $G$ :  $(\lambda x. g \ (z + x) - g \ z)$  has  $\text{fps\_expansion } G$   $G \neq 0$ 
assumes  $\text{eventually } (\lambda w. f \ w \neq 0)$  (at  $(g \ z)$ )
shows  $\text{zorder } (f \circ g) \ z = \text{zorder } f \ (g \ z) * \text{subdegree } G$ 
proof -
  define  $f'$  where  $f' = (\lambda w. f \ (g \ z + w))$ 
  define  $g'$  where  $g' = (\lambda w. g \ (z + w) - g \ z)$ 
  have  $\text{zorder } f \ (g \ z) = \text{zorder } f' \ 0$ 
    by (simp add:  $f'\_def$   $\text{zorder\_shift' add\_ac}$ )
  have  $\text{zorder } (\lambda x. g \ x - g \ z) \ z = \text{zorder } g' \ 0$ 
    by (simp add:  $g'\_def$   $\text{zorder\_shift' add\_ac}$ )
  have  $\text{zorder } (f \circ g) \ z = \text{zorder } (f' \circ g') \ 0$ 

```

```

  by (simp add: zorder_shift' f'_def g'_def add_ac o_def)
  also have ... = zorder f' 0 * int (subdegree G)
  proof (rule zorder_compose_aux)
    show isolated_singularity_at f' 0 unfolding f'_def
    using assms has_laurent_expansion_isolated_0 not_essential_has_laurent_expansion
  by blast
    show not_essential f' 0 unfolding f'_def
    using assms has_laurent_expansion_not_essential_0 not_essential_has_laurent_expansion
  by blast
  qed (use assms in ⟨auto simp: f'_def g'_def at_to_0' eventually_filtermap
add_ac⟩)
  also have zorder f' 0 = zorder f (g z)
  by (simp add: f'_def zorder_shift' add_ac)
  finally show ?thesis .
qed

```

lemma `fps_to_fls_eq_fls_const_iff` [simp]: $fps_to_fls\ F = fls_const\ c \iff F = fps_const\ c$

```

proof
  assume F = fps_const c
  thus fps_to_fls F = fls_const c
  by simp
next
  assume fps_to_fls F = fls_const c
  thus F = fps_const c
  by (metis fls_regpart_const fls_regpart_fps_trivial)
qed

```

lemma `zorder_compose'`:

```

  assumes isolated_singularity_at f (g z) not_essential f (g z)
  assumes g_analytic_on {z}
  assumes eventually (λw. f w ≠ 0) (at (g z))
  assumes eventually (λw. g w ≠ g z) (at z)
  shows zorder (f ∘ g) z = zorder f (g z) * zorder (λx. g x - g z) z
proof -
  obtain G where G [fps_expansion_intros]: (λx. g (z + x)) has_fps_expansion
G
  using assms analytic_at_imp_has_fps_expansion by blast
  have G': (λx. g (z + x) - g z) has_fps_expansion G - fps_const (g z)
  by (intro fps_expansion_intros)
  hence G'': (λx. g (z + x) - g z) has_laurent_expansion fps_to_fls (G -
fps_const (g z))
  using has_laurent_expansion_fps by blast
  have nz: G - fps_const (g z) ≠ 0
  using has_laurent_expansion_eventually_nonzero_iff[OF G''] assms by auto
  have zorder (f ∘ g) z = zorder f (g z) * subdegree (G - fps_const (g z))
proof (rule zorder_compose)
  show (λx. g (z + x) - g z) has_fps_expansion G - fps_const (g z)
  by (intro fps_expansion_intros)

```

```

qed (use assms nz in auto)
also have int (subdegree (G - fps_const (g z))) = fls_subdegree (fps_to_fls G
- fls_const (g z))
  by (simp flip: fls_subdegree_fls_to_fps)
also have ... = zorder ( $\lambda x. g x - g z$ ) z
  using has_laurent_expansion_zorder [OF G'] nz by auto
finally show ?thesis .
qed

```

lemma analytic_at_cong:

```

assumes eventually ( $\lambda x. f x = g x$ ) (nhds x) x = y
shows f analytic_on {x}  $\longleftrightarrow$  g analytic_on {y}
proof -
  have g analytic_on {x} if f analytic_on {x} eventually ( $\lambda x. f x = g x$ ) (nhds x)
for f g
  proof -
    have ( $\lambda y. f (x + y)$ ) has_fps_expansion fps_expansion f x
    by (rule analytic_at_imp_has_fps_expansion) fact
    also have ?this  $\longleftrightarrow$  ( $\lambda y. g (x + y)$ ) has_fps_expansion fps_expansion f x
    using that by (intro has_fps_expansion_cong refl) (auto simp: nhds_to_0'
eventually_filtermap)
    finally show ?thesis
    by (rule has_fps_expansion_imp_analytic)
  qed
from this[of f g] this[of g f] show ?thesis using assms
  by (auto simp: eq_commute)
qed

```

lemma has_laurent_expansion_sin' [laurent_expansion_intros]:

```

sin has_laurent_expansion fps_to_fls (fps_sin 1)
using has_fps_expansion_sin' has_fps_expansion_to_laurent by blast

```

lemma has_laurent_expansion_cos' [laurent_expansion_intros]:

```

cos has_laurent_expansion fps_to_fls (fps_cos 1)
using has_fps_expansion_cos' has_fps_expansion_to_laurent by blast

```

lemma has_laurent_expansion_sin [laurent_expansion_intros]:

```

( $\lambda z. \sin (c * z)$ ) has_laurent_expansion fps_to_fls (fps_sin c)
by (intro has_laurent_expansion_fps has_fps_expansion_sin)

```

lemma has_laurent_expansion_cos [laurent_expansion_intros]:

```

( $\lambda z. \cos (c * z)$ ) has_laurent_expansion fps_to_fls (fps_cos c)
by (intro has_laurent_expansion_fps has_fps_expansion_cos)

```

lemma has_laurent_expansion_tan' [laurent_expansion_intros]:

```

tan has_laurent_expansion fps_to_fls (fps_tan 1)
using has_fps_expansion_tan' has_fps_expansion_to_laurent by blast

```

```

lemma has_laurent_expansion_tan [laurent_expansion_intros]:
  ( $\lambda z. \tan (c * z)$ ) has_laurent_expansion_fps_to_fls (fps_tan c)
  by (intro has_laurent_expansion_fps has_fps_expansion_tan)

end

```

7 The Great Picard Theorem and its Applications

Ported from HOL Light (cauchy.ml) by L C Paulson, 2017

```

theory Great_Picard
  imports Conformal_Mappings
begin

```

7.1 Schottky's theorem

```

lemma Schottky_lemma0:
  assumes holg: f holomorphic_on S and cons: contractible S and  $a \in S$ 
  and  $f: \bigwedge z. z \in S \implies f z \neq 1 \wedge f z \neq -1$ 
  obtains g where g holomorphic_on S
     $norm(g a) \leq 1 + norm(f a) / 3$ 
     $\bigwedge z. z \in S \implies f z = \cos(of\_real\ pi * g z)$ 
proof –
  obtain g where holg: g holomorphic_on S and  $g: norm(g a) \leq pi + norm(f a)$ 
    and f_eq_cos:  $\bigwedge z. z \in S \implies f z = \cos(g z)$ 
    using contractible_imp_holomorphic_arccos_bounded [OF assms]
    by blast
  show ?thesis
proof
  show ( $\lambda z. g z / pi$ ) holomorphic_on S
    by (auto intro: holomorphic_intros holg)
  have  $3 \leq pi$ 
    using pi_approx by force
  have  $3 * norm(g a) \leq 3 * (pi + norm(f a))$ 
    using g by auto
  also have  $\dots \leq pi * 3 + pi * cmod (f a)$ 
    using  $\langle 3 \leq pi \rangle$  by (simp add: mult_right_mono algebra_simps)
  finally show  $cmod (g a / complex\_of\_real\ pi) \leq 1 + cmod (f a) / 3$ 
    by (simp add: field_simps norm_divide)
  show  $\bigwedge z. z \in S \implies f z = \cos (complex\_of\_real\ pi * (g z / complex\_of\_real\ pi))$ 
    by (simp add: f_eq_cos)
qed
qed

```

```

lemma Schottky_lemma1:
  fixes  $n::nat$ 
  assumes  $0 < n$ 

```

```

  shows  $0 < n + \text{sqrt}(\text{real } n^2 - 1)$ 
proof -
  have  $0 < n * n$ 
  by (simp add: assms)
  then show ?thesis
  by (metis add.commute add.right_neutral add_pos_nonneg assms diff_ge_0_iff_ge
  nat_less_real_le of_nat_0 of_nat_0_less_iff of_nat_power power2_eq_square
  real_sqrt_ge_0_iff)
qed

```

lemma Schottky_lemma2:

```

  fixes  $x::\text{real}$ 
  assumes  $0 \leq x$ 
  obtains  $n$  where  $0 < n \wedge |x - \ln(\text{real } n + \text{sqrt}((\text{real } n)^2 - 1)) / \pi| < 1/2$ 
proof -
  obtain  $n0::\text{nat}$  where  $0 < n0 \wedge \ln(n0 + \text{sqrt}(\text{real } n0^2 - 1)) / \pi \leq x$ 
  proof
    show  $\ln(\text{real } 1 + \text{sqrt}(\text{real } 1^2 - 1)) / \pi \leq x$ 
    by (auto simp: assms)
  qed auto
  moreover
  obtain  $M::\text{nat}$  where  $\bigwedge n. [0 < n; \ln(n + \text{sqrt}(\text{real } n^2 - 1)) / \pi \leq x] \implies n \leq M$ 
  proof
    fix  $n::\text{nat}$ 
    assume  $0 < n \wedge \ln(n + \text{sqrt}((\text{real } n)^2 - 1)) / \pi \leq x$ 
    then have  $\ln(n + \text{sqrt}((\text{real } n)^2 - 1)) \leq x * \pi$ 
    by (simp add: field_split_simps)
    then have  $*: \exp(\ln(n + \text{sqrt}((\text{real } n)^2 - 1))) \leq \exp(x * \pi)$ 
    by blast
    have  $0: 0 \leq \text{sqrt}((\text{real } n)^2 - 1)$ 
    using  $\langle 0 < n \rangle$  by auto
    have  $n + \text{sqrt}((\text{real } n)^2 - 1) = \exp(\ln(n + \text{sqrt}((\text{real } n)^2 - 1)))$ 
    by (simp add: Suc_leI  $\langle 0 < n \rangle$  add_pos_nonneg real_of_nat_ge_one_iff)
    also have  $\dots \leq \exp(x * \pi)$ 
    using  $*$  by blast
    finally have  $\text{real } n \leq \exp(x * \pi)$ 
    using  $0$  by linarith
    then show  $n \leq \text{nat}(\text{ceiling}(\exp(x * \pi)))$ 
    by linarith
  qed
  ultimately obtain  $n$  where
     $0 < n$  and  $le_x: \ln(n + \text{sqrt}(\text{real } n^2 - 1)) / \pi \leq x$ 
    and  $le_n: \bigwedge k. [0 < k; \ln(k + \text{sqrt}(\text{real } k^2 - 1)) / \pi \leq x] \implies k \leq n$ 
  using bounded_Max_nat [of  $\lambda n. 0 < n \wedge \ln(n + \text{sqrt}((\text{real } n)^2 - 1)) / \pi \leq x$ ] by metis
  define  $a$  where  $a \equiv \ln(n + \text{sqrt}(\text{real } n^2 - 1)) / \pi$ 

```

```

define b where  $b \equiv \ln (1 + \text{real } n + \text{sqrt} ((1 + \text{real } n)^2 - 1)) / \pi$ 
have le_xa:  $a \leq x$ 
and le_na:  $\bigwedge k. \llbracket 0 < k; \ln(k + \text{sqrt}(\text{real } k^2 - 1)) / \pi \leq x \rrbracket \implies k \leq n$ 
using le_x le_n by (auto simp: a_def)
moreover have  $x < b$ 
using le_n [of Suc n] by (force simp: b_def)
moreover have  $b - a < 1$ 
proof -
have  $\ln (1 + \text{real } n + \text{sqrt} ((1 + \text{real } n)^2 - 1)) - \ln (\text{real } n + \text{sqrt} ((\text{real } n)^2 - 1)) =$ 
 $\ln ((1 + \text{real } n + \text{sqrt} ((1 + \text{real } n)^2 - 1)) / (\text{real } n + \text{sqrt} ((\text{real } n)^2 - 1)))$ 
by (simp add: <0 < n> Schottky_lemma1 add_pos_nonneg ln_div [symmetric])
also have  $\dots \leq 3$ 
proof (cases n = 1)
case True
have  $\text{sqrt } 3 \leq 2$ 
by (simp add: real_le_sqrt)
then have  $(2 + \text{sqrt } 3) \leq 4$ 
by simp
also have  $\dots \leq \text{exp } 3$ 
using exp_ge_add_one_self [of 3::real] by simp
finally have  $\ln (2 + \text{sqrt } 3) \leq 3$ 
by (metis add_nonneg_nonneg add_pos_nonneg dbl_def dbl_inc_def
dbl_inc_simps(3)
dbl_simps(3) exp_gt_zero ln_exp ln_le_cancel_iff real_sqrt_ge_0_iff
zero_le_one zero_less_one)
then show ?thesis
by (simp add: True)
next
case False with  $\langle 0 < n \rangle$  have  $1 < n \ 2 \leq n$ 
by linarith+
then have  $1: 1 \leq \text{real } n * \text{real } n$ 
by (metis less_imp_le_nat one_le_power power2_eq_square real_of_nat_ge_one_iff)
have  $*$ :  $4 + (m+2) * 2 \leq (m+2) * ((m+2) * 3)$  for  $m::\text{nat}$ 
by simp
have  $4 + n * 2 \leq n * (n * 3)$ 
using  $*$  [of n-2]  $\langle 2 \leq n \rangle$ 
by (metis le_add_diff_inverse2)
then have  $**$ :  $4 + \text{real } n * 2 \leq \text{real } n * (\text{real } n * 3)$ 
by (metis (mono_tags, opaque_lifting) of_nat_le_iff of_nat_add of_nat_mult
of_nat_numeral)
have  $\text{sqrt} ((1 + \text{real } n)^2 - 1) \leq 2 * \text{sqrt} ((\text{real } n)^2 - 1)$ 
by (auto simp: real_le_sqrt power2_eq_square algebra_simps 1 **)
then
have  $((1 + \text{real } n + \text{sqrt} ((1 + \text{real } n)^2 - 1)) / (\text{real } n + \text{sqrt} ((\text{real } n)^2 - 1))) \leq 2$ 
using Schottky_lemma1 <0 < n> by (simp add: field_split_simps)
then have  $\ln ((1 + \text{real } n + \text{sqrt} ((1 + \text{real } n)^2 - 1)) / (\text{real } n + \text{sqrt} ((\text{real } n)^2 - 1))) \leq 2$ 

```

```

n)^2 - 1))) ≤ ln 2
  using Schottky_lemma1 [of n] ‹0 < n›
  by (simp add: field_split_simps add_pos_nonneg)
  also have ... ≤ 3
  using ln_add_one_self_le_self [of 1] by auto
  finally show ?thesis .
qed
also have ... < pi
  using pi_approx by simp
finally show ?thesis
  by (simp add: a_def b_def field_split_simps)
qed
ultimately have |x - a| < 1/2 ∨ |x - b| < 1/2
  by (auto simp: abs_if)
then show thesis
proof
  assume |x - a| < 1/2
  then show ?thesis
    by (rule_tac n=n in that) (auto simp: a_def ‹0 < n›)
next
  assume |x - b| < 1/2
  then show ?thesis
    by (rule_tac n=Suc n in that) (auto simp: b_def ‹0 < n›)
qed
qed

```

```

lemma Schottky_lemma3:
  fixes z::complex
  assumes z ∈ (⋃ m ∈ Ints. ⋃ n ∈ {0<..}. {Complex m (ln(n + sqrt(real n ^ 2
- 1)) / pi)})
    ∪ (⋃ m ∈ Ints. ⋃ n ∈ {0<..}. {Complex m (-ln(n + sqrt(real n ^ 2 -
1)) / pi)})
  shows cos(pi * cos(pi * z)) = 1 ∨ cos(pi * cos(pi * z)) = -1
proof -
  have sqrt2 [simp]: complex_of_real (sqrt x) * complex_of_real (sqrt x) = x if x
≥ 0 for x::real
  by (metis abs_of_nonneg of_real_mult real_sqrt_mult_self that)
  define plusi where plusi (e::complex) ≡ e + inverse e for e
  have 1: ∃ k. plusi (exp (i * (of_int m * complex_of_real pi) - ln (real n + sqrt
((real n)^2 - 1)))) = of_int k * 2
    (is ∃ k. ?Φ k)
    if n > 0 for m n
proof -
  have eqq: e ≠ 0 ⇒ plusi e = n ↔ (inverse e) ^ 2 = n/e - 1 for n e::complex
  by (auto simp: plusi_def field_simps power2_eq_square)
  have [simp]: 1 ≤ real n * real n
  using nat_0_less_mult_iff nat_less_real_le that by force
  consider odd m | even m

```

```

    by blast
  then have  $\exists k. ?\Phi k$ 
  proof cases
    case 1
    then have  $?\Phi (-n)$ 
      using Schottky_lemma1 [OF that]
      by (simp add: eeq) (simp add: power2_eq_square exp_diff exp_Euler
    exp_of_real algebra_simps sin_int_times_real cos_int_times_real)
    then show ?thesis ..
  next
    case 2
    then have  $?\Phi n$ 
      using Schottky_lemma1 [OF that]
      by (simp add: eeq) (simp add: power2_eq_square exp_diff exp_Euler
    exp_of_real algebra_simps)
    then show ?thesis ..
  qed
  then show ?thesis by blast
qed
have 2:  $\exists k. \text{plusi } (\exp (i * (\text{of\_int } m * \text{complex\_of\_real } \pi) +$ 
       $(\ln (\text{real } n + \text{sqrt } ((\text{real } n)^2 - 1)))) = \text{of\_int } k * 2$ 
      (is  $\exists k. ?\Phi k$ )
      if  $n > 0$  for  $m n$ 
proof -
  have eeq:  $e \neq 0 \implies \text{plusi } e = n \iff e^2 - n * e + 1 = 0$  for  $n e :: \text{complex}$ 
    by (auto simp: plusi_def field_simps power2_eq_square)
  have [simp]:  $1 \leq \text{real } n * \text{real } n$ 
    by (metis One_nat_def add.commute nat_less_real_le of_nat_1 of_nat_Suc
  one_le_power power2_eq_square that)
  consider odd  $m \mid$  even  $m$ 
  by blast
  then have  $\exists k. ?\Phi k$ 
  proof cases
    case 1
    then have  $?\Phi (-n)$ 
      using Schottky_lemma1 [OF that]
      by (simp add: eeq) (simp add: power2_eq_square exp_add exp_Euler
    exp_of_real algebra_simps sin_int_times_real cos_int_times_real)
    then show ?thesis ..
  next
    case 2
    then have  $?\Phi n$ 
      using Schottky_lemma1 [OF that]
      by (simp add: eeq) (simp add: power2_eq_square exp_add exp_Euler
    exp_of_real algebra_simps)
    then show ?thesis ..
  qed
  then show ?thesis by blast
qed

```



```

have  $\exists x. \cos(\text{complex\_of\_real } \pi * z) = \text{of\_int } x$ 
  using assms
  apply (auto simp: Ints_def cos_exp_eq exp_minus Complex_eq simp flip: plusi_def)
  apply (auto simp: algebra_simps dest: 1 2)
  done
then have  $\sin(\pi * \cos(\pi * z)) ^ 2 = 0$ 
  by (simp add: Complex_Transcendental.sin_eq_0)
then have  $1 - \cos(\pi * \cos(\pi * z)) ^ 2 = 0$ 
  by (simp add: sin_squared_eq)
then show ?thesis
  using power2_eq_1_iff by auto
qed

```

theorem *Schottky*:

```

assumes holf: f holomorphic_on cball 0 1
  and nof0: norm(f 0) ≤ r
  and not01:  $\bigwedge z. z \in \text{cball } 0 \ 1 \implies \neg(f z = 0 \vee f z = 1)$ 
  and  $0 < t < 1$  norm z ≤ t
  shows  $\text{norm}(f z) \leq \exp(\pi * \exp(\pi * (2 + 2 * r + 12 * t / (1 - t))))$ 
proof -
  obtain h where holf: h holomorphic_on cball 0 1
    and nh0: norm(h 0) ≤ 1 + norm(2 * f 0 - 1) / 3
    and h:  $\bigwedge z. z \in \text{cball } 0 \ 1 \implies 2 * f z - 1 = \cos(\text{of\_real } \pi * h z)$ 
  proof (rule Schottky_lemma0 [of  $\lambda z. 2 * f z - 1$  cball 0 1 0])
    show  $(\lambda z. 2 * f z - 1)$  holomorphic_on cball 0 1
      by (intro holomorphic_intros holf)
    show contractible (cball (0::complex) 1)
      by (auto simp: convex_imp_contractible)
    show  $\bigwedge z. z \in \text{cball } 0 \ 1 \implies 2 * f z - 1 \neq 1 \wedge 2 * f z - 1 \neq -1$ 
      using not01 by force
  qed auto
  obtain g where holg: g holomorphic_on cball 0 1
    and ng0: norm(g 0) ≤ 1 + norm(h 0) / 3
    and g:  $\bigwedge z. z \in \text{cball } 0 \ 1 \implies h z = \cos(\text{of\_real } \pi * g z)$ 
  proof (rule Schottky_lemma0 [OF holf convex_imp_contractible, of 0])
    show  $\bigwedge z. z \in \text{cball } 0 \ 1 \implies h z \neq 1 \wedge h z \neq -1$ 
      using h not01 by fastforce+
  qed auto
  have g0_2_f0: norm(g 0) ≤ 2 + norm(f 0)
  proof -
    have cmod (2 * f 0 - 1) ≤ cmod (2 * f 0) + 1
      by (metis norm_one norm_triangle_ineq4)
    also have  $\dots \leq 6 + 9 * \text{cmod}(f 0)$ 
      by auto
    finally have  $1 + \text{norm}(2 * f 0 - 1) / 3 \leq (2 + \text{norm}(f 0) - 1) * 3$ 
      by (simp add: divide_simps)
    with nh0 have  $\text{norm}(h 0) \leq (2 + \text{norm}(f 0) - 1) * 3$ 

```

```

    by linarith
  then have  $1 + \text{norm}(h \ 0) / 3 \leq 2 + \text{norm}(f \ 0)$ 
    by simp
  with ng0 show ?thesis
    by auto
qed
have  $z \in \text{ball } 0 \ 1$ 
  using assms by auto
have norm_g_12:  $\text{norm}(g \ z - g \ 0) \leq (12 * t) / (1 - t)$ 
proof -
  obtain g' where  $g': \bigwedge x. x \in \text{cball } 0 \ 1 \implies (g \ \text{has\_field\_derivative } g' \ x)$  (at x
  within cball 0 1)
  using holg [unfolded holomorphic_on_def field_differentiable_def] by metis
  have int_g':  $(g' \ \text{has\_contour\_integral } g \ z - g \ 0)$  (linepath 0 z)
  using contour_integral_primitive [OF g' valid_path_linepath, of 0 z]
  using  $\langle z \in \text{ball } 0 \ 1 \rangle$  segment_bound1 by fastforce
  have cmod (g' w)  $\leq 12 / (1 - t)$  if  $w \in \text{closed\_segment } 0 \ z$  for w
proof -
  have w:  $w \in \text{ball } 0 \ 1$ 
    using segment_bound [OF that]  $\langle z \in \text{ball } 0 \ 1 \rangle$  by simp
  have *:  $\llbracket \bigwedge b. (\exists w \in T \cup U. w \in \text{ball } b \ 1); \bigwedge x. x \in D \implies g \ x \notin T \cup U \rrbracket$ 
 $\implies \nexists b. \text{ball } b \ 1 \subseteq g \ ' \ D$  for  $T \ U \ D$ 
    by force
  have ttt:  $1 - t \leq \text{dist } w \ u$  if  $\text{cmod } u = 1$  for u
    using  $\langle \text{norm } z \leq t \rangle$  segment_bound1 [OF  $\langle w \in \text{closed\_segment } 0 \ z \rangle$ ]
norm_triangle_ineq2 [of u w] that
  by (simp add: dist_norm norm_minus_commute)
  have  $\nexists b. \text{ball } b \ 1 \subseteq g \ ' \ \text{cball } 0 \ 1$ 
proof (rule *)
  show  $(\exists w \in (\bigcup m \in \text{Ints}. \bigcup n \in \{0 <.. \}. \{\text{Complex } m \ (\ln(n + \text{sqrt}(\text{real } n$ 
 $\wedge 2 - 1)) / \text{pi}\}) \cup$ 
 $(\bigcup m \in \text{Ints}. \bigcup n \in \{0 <.. \}. \{\text{Complex } m \ (-\ln(n + \text{sqrt}(\text{real } n \wedge 2$ 
 $- 1)) / \text{pi}\}). w \in \text{ball } b \ 1)$  for b
proof -
  obtain m where  $m: m \in \mathbb{Z} \mid \text{Re } b - m \leq 1/2$ 
    by (metis Ints_of_int abs_minus_commute of_int_round_abs_le)
  show ?thesis
proof (cases 0::real Im b rule: le_cases)
  case le
  then obtain n where  $0 < n$  and  $n: |\text{Im } b - \ln(n + \text{sqrt}((\text{real } n)^2 -$ 
 $1)) / \text{pi}| < 1/2$ 
    using Schottky_lemma2 [of Im b] by blast
  have  $\text{dist } b \ (\text{Complex } m \ (\text{Im } b)) \leq 1/2$ 
    by (metis cancel_comm_monoid_add_class.diff_cancel cmod_eq_Re
  complex.sel(1) complex.sel(2) dist_norm m(2) minus_complex.code)
  moreover
  have  $\text{dist} \ (\text{Complex } m \ (\text{Im } b)) \ (\text{Complex } m \ (\ln(n + \text{sqrt}((\text{real } n)^2 -$ 
 $1)) / \text{pi})) < 1/2$ 
    using n by (simp add: complex_norm cmod_eq_Re complex_diff

```

```

dist_norm del: Complex_eq)
  ultimately have dist b (Complex m (ln (real n + sqrt ((real n)2 - 1))
/ pi)) < 1
  by (simp add: dist_triangle_lt [of b Complex m (Im b)] dist_commute)
  with le m <0 < n> show ?thesis
  apply (rule_tac x = Complex m (ln (real n + sqrt ((real n)2 - 1)) /
pi) in bexI)
  by (force simp del: Complex_eq greaterThan_0)+
next
case ge
then obtain n where 0 < n and n: |- Im b - ln (real n + sqrt ((real
n)2 - 1)) / pi < 1/2
  using Schottky_lemma2 [of - Im b] by auto
  have dist b (Complex m (Im b)) ≤ 1/2
  by (metis cancel_comm_monoid_add_class.diff_cancel cmod_eq_Re
complex.sel(1) complex.sel(2) dist_norm m(2) minus_complex.code)
  moreover
  have dist (Complex m (- ln (n + sqrt ((real n)2 - 1)) / pi)) (Complex
m (Im b))
    = |- Im b - ln (real n + sqrt ((real n)2 - 1)) / pi|
  by (simp add: complex_norm dist_norm cmod_eq_Re complex_diff)
  ultimately have dist b (Complex m (- ln (real n + sqrt ((real n)2 -
1)) / pi)) < 1
  using n by (simp add: dist_triangle_lt [of b Complex m (Im b)]
dist_commute)
  with ge m <0 < n> show ?thesis
  by (rule_tac x = Complex m (- ln (real n + sqrt ((real n)2 - 1)) /
pi) in bexI) auto
qed
qed
show g v ∉ (⋃ m ∈ Ints. ⋃ n ∈ {0<..}. {Complex m (ln(n + sqrt(real n ^
2 - 1)) / pi)}) ∪
  (⋃ m ∈ Ints. ⋃ n ∈ {0<..}. {Complex m (-ln(n + sqrt(real n ^ 2
- 1)) / pi)})
  if v ∈ cball 0 1 for v
  using not01 [OF that]
  by (force simp: g [OF that, symmetric] h [OF that, symmetric] dest!:
Schottky_lemma3 [of g v])
qed
then have 12: (1 - t) * cmod (deriv g w) / 12 < 1
  using Bloch_general [OF holg _ ttt, of 1] w by force
have g field_differentiable at w within cball 0 1
  using holg w by (simp add: holomorphic_on_def)
then have g field_differentiable at w within ball 0 1
  using ball_subset_cball field_differentiable_within_subset by blast
with w have der_gw: (g has_field_derivative deriv g w) (at w)
by (simp add: field_differentiable_within_open [of _ ball 0 1] field_differentiable_derivI)
with DERIV_unique [OF der_gw] g' w have deriv g w = g' w
by (metis open_ball at_within_open ball_subset_cball has_field_derivative_subset

```

```

subsetCE)
  then show  $cmod (g' w) \leq 12 / (1 - t)$ 
    using  $g' 12 \langle t < 1 \rangle$  by (simp add: field_simps)
  qed
  then have  $cmod (g z - g 0) \leq 12 / (1 - t) * cmod z$ 
    using has_contour_integral_bound_linepath [OF int_g', of 12/(1 - t)] assms
    by simp
  with  $\langle cmod z \leq t \rangle \langle t < 1 \rangle$  show ?thesis
    by (simp add: field_split_simps)
  qed
  have fz:  $f z = (1 + \cos(\text{of\_real } \pi * h z)) / 2$ 
    using  $h \langle z \in \text{ball } 0 \ 1 \rangle$  by (auto simp: field_simps)
  have  $cmod (f z) \leq \exp (cmod (\text{complex\_of\_real } \pi * h z))$ 
    by (simp add: fz mult.commute norm_cos_plus1_le)
  also have ...  $\leq \exp (\pi * \exp (\pi * (2 + 2 * r + 12 * t / (1 - t))))$ 
  proof (simp add: norm_mult)
    have  $cmod (g z - g 0) \leq 12 * t / (1 - t)$ 
      using norm_g_12  $\langle t < 1 \rangle$  by (simp add: norm_mult)
    then have  $cmod (g z) - cmod (g 0) \leq 12 * t / (1 - t)$ 
      using norm_triangle_ineq2 order_trans by blast
    then have *:  $cmod (g z) \leq 2 + 2 * r + 12 * t / (1 - t)$ 
      using g0_2_f0 norm_ge_zero [of f 0] nof0
      by linarith
    have  $cmod (h z) \leq \exp (cmod (\text{complex\_of\_real } \pi * g z))$ 
      using  $\langle z \in \text{ball } 0 \ 1 \rangle$  by (simp add: g norm_cos_le)
    also have ...  $\leq \exp (\pi * (2 + 2 * r + 12 * t / (1 - t)))$ 
      using  $\langle t < 1 \rangle$  nof0 * by (simp add: norm_mult)
    finally show  $cmod (h z) \leq \exp (\pi * (2 + 2 * r + 12 * t / (1 - t)))$  .
  qed
  finally show ?thesis .
qed

```

7.2 The Little Picard Theorem

theorem Landau_Picard:

obtains R

where $\bigwedge z. 0 < R z$

$\bigwedge f. \llbracket f \text{ holomorphic_on } \text{cball } 0 \ (R(f\ 0));$

$\bigwedge z. \text{norm } z \leq R(f\ 0) \implies f z \neq 0 \wedge f z \neq 1 \rrbracket \implies \text{norm}(\text{deriv } f\ 0)$

< 1

proof -

define R where $R \equiv \lambda z. 3 * \exp(\pi * \exp(\pi * (2 + 2 * cmod z + 12)))$

show ?thesis

proof

show $Rpos: \bigwedge z. 0 < R z$

by (auto simp: R_def)

show $\text{norm}(\text{deriv } f\ 0) < 1$

if $holf: f \text{ holomorphic_on } \text{cball } 0 \ (R(f\ 0))$

and $Rf: \bigwedge z. \text{norm } z \leq R(f\ 0) \implies f z \neq 0 \wedge f z \neq 1$ for f

```

proof –
  let ?r = R(f 0)
  define g where g ≡ f ∘ (λz. of_real ?r * z)
  have 0 < ?r
    using Rpos by blast
  have holg: g holomorphic_on cball 0 1
    unfolding g_def
  proof (intro holomorphic_intros holomorphic_on_compose holomorphic_on_subset
[OF holg])
    show (*) (complex_of_real (R (f 0))) ‘ cball 0 1 ⊆ cball 0 (R (f 0))
      using Rpos by (auto simp: less_imp_le norm_mult)
    qed
  have *: norm(g z) ≤ exp(pi * exp(pi * (2 + 2 * norm (f 0) + 12 * t / (1
– t))))
    if 0 < t t < 1 norm z ≤ t for t z
  proof (rule Schottky [OF holg])
    show cmod (g 0) ≤ cmod (f 0)
      by (simp add: g_def)
    show ∧z. z ∈ cball 0 1 ⇒ ¬ (g z = 0 ∨ g z = 1)
      using Rpos by (simp add: g_def less_imp_le norm_mult Rf)
    qed (auto simp: that)
  have C1: g holomorphic_on ball 0 (1/2)
    by (rule holomorphic_on_subset [OF holg]) auto
  have C2: continuous_on (cball 0 (1/2)) g
    by (meson cball_divide_subset_numeral holg holomorphic_on_imp_continuous_on
holomorphic_on_subset)
  have C3: cmod (g z) ≤ R (f 0) / 3 if cmod (0 – z) = 1/2 for z
  proof –
    have norm(g z) ≤ exp(pi * exp(pi * (2 + 2 * norm (f 0) + 12)))
      using * [of 1/2] that by simp
    also have ... = ?r / 3
      by (simp add: R_def)
    finally show ?thesis .
  qed
  then have cmod_g'_le: cmod (deriv g 0) * 3 ≤ R (f 0) * 2
    using Cauchy_inequality [OF C1 C2 C3, of 1] by simp
  have holg': f holomorphic_on ball 0 (R(f 0))
    by (rule holomorphic_on_subset [OF holg]) auto
  then have fd0: f field_differentiable at 0
    by (rule holomorphic_on_imp_differentiable_at [OF _ open_ball])
      (auto simp: Rpos [of f 0])
  have g_eq: deriv g 0 = of_real ?r * deriv f 0
    unfolding g_def
  by (metis DERIV_imp_deriv DERIV_chain DERIV_cmult_Id fd0 field_differentiable_derivI
mult commute mult_zero_right)
  show ?thesis
    using cmod_g'_le Rpos [of f 0] by (simp add: g_eq norm_mult)
  qed
qed

```

qed

lemma *little_Picard_01*:

assumes *holf*: f holomorphic_on UNIV **and** *f01*: $\bigwedge z. f z \neq 0 \wedge f z \neq 1$

obtains c **where** $f = (\lambda x. c)$

proof –

obtain R

where $Rpos$: $\bigwedge z. 0 < R z$

and R : $\bigwedge h. \llbracket h \text{ holomorphic_on cball } 0 (R(h\ 0));$

$\bigwedge z. \text{norm } z \leq R(h\ 0) \implies h z \neq 0 \wedge h z \neq 1 \rrbracket \implies \text{norm}(\text{deriv}$

$h\ 0) < 1$

using *Landau_Picard* **by** *metis*

have *contf*: continuous_on UNIV f

by (*simp add: holf holomorphic_on_imp_continuous_on*)

show *?thesis*

proof (*cases* $\forall x. \text{deriv } f\ x = 0$)

case *True*

have (f has_field_derivative 0) (at x) **for** x

by (*metis True UNIV_I holf holomorphic_derivI open_UNIV*)

then obtain c **where** $\bigwedge x. f(x) = c$

by (*meson UNIV_I DERIV_zero_connected_constant [OF connected_UNIV open_UNIV finite.emptyI contf]*)

then show *?thesis*

using *that by auto*

next

case *False*

then obtain w **where** w : $\text{deriv } f\ w \neq 0$ **by** *auto*

define fw **where** $fw \equiv (f \circ (\lambda z. w + z / \text{deriv } f\ w))$

have *norm_let1*: $\text{norm}(\text{deriv } fw\ 0) < 1$

proof (*rule R*)

show fw holomorphic_on cball 0 ($R (fw\ 0)$)

unfolding *fw_def*

by (*intro holomorphic_intros holomorphic_on_compose w holomorphic_on_subset [OF holf] subset_UNIV*)

show $fw\ z \neq 0 \wedge fw\ z \neq 1$ **if** $c \text{mod } z \leq R (fw\ 0)$ **for** z

using *f01 by (simp add: fw_def)*

qed

have (fw has_field_derivative $\text{deriv } f\ w * \text{inverse} (\text{deriv } f\ w)$) (at 0)

unfolding *fw_def*

apply (*intro DERIV_chain_derivative_eq_intros w*)+

using *holf holomorphic_derivI by (force simp: field_simps)*+

then show *?thesis*

using *norm_let1 w by (simp add: DERIV_imp_deriv)*

qed

qed

theorem *little_Picard*:

assumes *holf*: f holomorphic_on UNIV

```

    and  $a \neq b$  range  $f \cap \{a, b\} = \{\}$ 
    obtains  $c$  where  $f = (\lambda x. c)$ 
  proof -
    let  $?g = \lambda x. 1/(b - a) * (f x - b) + 1$ 
    obtain  $c$  where  $?g = (\lambda x. c)$ 
    proof (rule little_Picard_01)
      show  $?g$  holomorphic_on UNIV
        by (intro holomorphic_intros holf)
      show  $\bigwedge z. ?g z \neq 0 \wedge ?g z \neq 1$ 
        using assms by (auto simp: field_simps)
    qed auto
    then have  $?g x = c$  for  $x$ 
      by meson
    then have  $f x = c * (b - a) + a$  for  $x$ 
      using assms by (auto simp: field_simps)
    then show  $?thesis$ 
      using that by blast
  qed

```

A couple of little applications of Little Picard

```

lemma holomorphic_periodic_fixpoint:
  assumes holf:  $f$  holomorphic_on UNIV
    and  $p \neq 0$  and per:  $\bigwedge z. f(z + p) = f z$ 
  obtains  $x$  where  $f x = x$ 
  proof -
    have False if non:  $\bigwedge x. f x \neq x$ 
    proof -
      obtain  $c$  where  $(\lambda z. f z - z) = (\lambda z. c)$ 
      proof (rule little_Picard)
        show  $(\lambda z. f z - z)$  holomorphic_on UNIV
          by (simp add: holf holomorphic_on_diff)
        show range  $(\lambda z. f z - z) \cap \{p, 0\} = \{\}$ 
          using assms non by auto (metis add.commute diff_eq_eq)
      qed (auto simp: assms)
      with per show False
        by (metis add.commute add_cancel_left_left <math>p \neq 0</math> diff_add_cancel)
    qed
  then show  $?thesis$ 
    using that by blast
  qed

```

```

lemma holomorphic_involution_point:
  assumes holfU:  $f$  holomorphic_on UNIV and non:  $\bigwedge a. f \neq (\lambda x. a + x)$ 
  obtains  $x$  where  $f(f x) = x$ 
  proof -
    { assume non_ff [simp]:  $\bigwedge x. f(f x) \neq x$ 
      then have non_fp [simp]:  $f z \neq z$  for  $z$ 
        by metis
    }

```

```

have holf: f holomorphic_on X for X
  using assms holomorphic_on_subset by blast
obtain c where c: (λx. (f(f x) - x)/(f x - x)) = (λx. c)
proof (rule little_Picard_01)
  show (λx. (f(f x) - x)/(f x - x)) holomorphic_on UNIV
    using non_fp
  by (intro holomorphic_intros holf holomorphic_on_compose [unfolded o_def,
OF holf]) auto
qed auto
then obtain c ≠ 0 c ≠ 1
by (metis (no_types, opaque_lifting) non_ff diff_zero divide_eq_0_iff right_inverse_eq)
have eq: f(f x) - c * f x = x*(1 - c) for x
  using fun_cong [OF c, of x] by (simp add: field_simps)
have df_times_dff: deriv f z * (deriv f (f z) - c) = 1 * (1 - c) for z
proof (rule DERIV_unique)
  show ((λx. f (f x) - c * f x) has_field_derivative
    deriv f z * (deriv f (f z) - c)) (at z)
  by (rule derivative_eq_intros holomorphic_derivI [OF holfU]
    DERIV_chain [unfolded o_def, where f=f and g=f] | simp add:
algebra_simps)+
  show ((λx. f (f x) - c * f x) has_field_derivative 1 * (1 - c)) (at z)
  by (simp add: eq mult_commute_abs)
qed
{ fix z::complex
  obtain k where k: deriv f ∘ f = (λx. k)
  proof (rule little_Picard)
    show (deriv f ∘ f) holomorphic_on UNIV
      by (meson holfU holomorphic_deriv holomorphic_on_compose holomor-
phic_on_subset open_UNIV subset_UNIV)
    obtain deriv f (f x) ≠ 0 deriv f (f x) ≠ c for x
      using df_times_dff ⟨c ≠ 1⟩ eq_iff_diff_eq_0
    by (metis lambda_one mult_zero_left mult_zero_right)
    then show range (deriv f ∘ f) ∩ {0,c} = {}
      by force
  qed (use ⟨c ≠ 0⟩ in auto)
  have ¬ f constant_on UNIV
    by (meson UNIV_I non_ff constant_on_def)
  with holf open_mapping_thm have open(range f)
    by blast
  obtain l where l: ∧x. f x - k * x = l
  proof (rule DERIV_zero_connected_constant [of UNIV {} λx. f x - k * x],
simp_all)
    have deriv f w - k = 0 for w
      proof (rule analytic_continuation [OF open_UNIV connected_UNIV
subset_UNIV, of λz. deriv f z - k f z range f w])
        show (λz. deriv f z - k) holomorphic_on UNIV
          by (intro holomorphic_intros holf open_UNIV)
        show f z islimpt range f
          by (metis (no_types, lifting) IntI UNIV_I ⟨open (range f)⟩ im-

```



```

age_eqI inf.absorb_iff2 inf_aci(1) islimpt_UNIV islimpt_eq_acc_point open_Int
top_greatest)
  show  $\bigwedge z. z \in \text{range } f \implies \text{deriv } f z - k = 0$ 
    by (metis comp_def diff_self image_iff k)
  qed auto
  moreover
  have (( $\lambda x. f x - k * x$ ) has_field_derivative deriv f x - k) (at x) for x
    by (metis DERIV_cmult_Id Deriv.field_differentiable_diff UNIV_I
field_differentiable_derivI holf holomorphic_on_def)
  ultimately
  show  $\forall x. ((\lambda x. f x - k * x) \text{ has\_field\_derivative } 0)$  (at x)
    by auto
  show continuous_on UNIV ( $\lambda x. f x - k * x$ )
  by (simp add: continuous_on_diff holf holomorphic_on_imp_continuous_on)
  qed (auto simp: connected_UNIV)
  have False
  proof (cases k=1)
  case True
  then have  $\exists x. k * x + l \neq a + x$  for a
    using l non [of a] ext [of f (+) a]
    by (metis add.commute diff_eq_eq)
  with True show ?thesis by auto
  next
  case False
  have  $\bigwedge x. (1 - k) * x \neq f 0$ 
    using l [of 0]
    by (simp add: algebra_simps) (metis diff_add_cancel l mult.commute
non_fp)
  then show False
  by (metis False eq_iff_diff_eq_0 mult.commute nonzero_mult_div_cancel_right
times_divide_eq_right)
  qed
}
}
then show thesis
  using that by blast
qed

```

7.3 The Arzelà–Ascoli theorem

lemma subsequence_diagonalization_lemma:

```

  fixes P :: nat  $\Rightarrow$  (nat  $\Rightarrow$  'a)  $\Rightarrow$  bool
  assumes sub:  $\bigwedge i r. \exists k. \text{strict\_mono } (k :: \text{nat} \Rightarrow \text{nat}) \wedge P i (r \circ k)$ 
  and P_P:  $\bigwedge i r :: \text{nat} \Rightarrow 'a. \bigwedge k1 k2 N. \llbracket P i (r \circ k1); \bigwedge j. N \leq j \implies \exists j'. j \leq j' \wedge k2 j = k1 j' \rrbracket \implies P i (r \circ k2)$ 
  obtains k where strict_mono (k :: nat  $\Rightarrow$  nat)  $\bigwedge i. P i (r \circ k)$ 
  proof -
  obtain kk where  $\bigwedge i r. \text{strict\_mono } (kk i r :: \text{nat} \Rightarrow \text{nat}) \wedge P i (r \circ (kk i r))$ 

```

```

    using sub by metis
  then have sub_kk:  $\bigwedge i r. \text{strict\_mono } (kk\ i\ r)$  and P_kk:  $\bigwedge i r. P\ i\ (r \circ (kk\ i\ r))$ 
  by auto
  define rr where  $rr \equiv \text{rec\_nat } (kk\ 0\ r) (\lambda n x. x \circ kk\ (Suc\ n)\ (r \circ x))$ 
  then have [simp]:  $rr\ 0 = kk\ 0\ r \wedge n. rr\ (Suc\ n) = rr\ n \circ kk\ (Suc\ n)\ (r \circ rr\ n)$ 
  by auto
  show thesis
  proof
    have sub_rr: strict_mono (rr i) for i
    using sub_kk by (induction i) (auto simp: strict_mono_def o_def)
    have P_rr:  $P\ i\ (r \circ rr\ i)$  for i
    using P_kk by (induction i) (auto simp: o_def)
    have  $i \leq i+d \implies rr\ i\ n \leq rr\ (i+d)\ n$  for d i n
    proof (induction d)
      case 0 then show ?case
      by simp
    next
      case (Suc d) then show ?case
      using seq_suble [OF sub_kk] strict_mono_less_eq [OF sub_rr]
      by (simp add: order_subst1)
    qed
    then have  $\bigwedge i j n. i \leq j \implies rr\ i\ n \leq rr\ j\ n$ 
    by (metis le_iff_add)
    show strict_mono ( $\lambda n. rr\ n\ n$ )
    unfolding strict_mono_Suc_iff
    by (simp add: Suc_le_lessD strict_monoD strict_mono_imp_increasing
    sub_kk sub_rr)
    have  $\exists j. i \leq j \wedge rr\ (n+d)\ i = rr\ n\ j$  for d n i
    proof (induction d arbitrary: i)
      case (Suc d)
      then show ?case
      using seq_suble [OF sub_kk] by (simp (meson order_trans))
    qed auto
    then have  $\bigwedge m n i. n \leq m \implies \exists j. i \leq j \wedge rr\ m\ i = rr\ n\ j$ 
    by (metis le_iff_add)
    then show  $P\ i\ (r \circ (\lambda n. rr\ n\ n))$  for i
    by (meson P_rr P_P)
  qed
  qed

lemma function_convergent_subsequence:
  fixes  $f :: [nat, 'a] \Rightarrow 'b::\{real\_normed\_vector, heine\_borel\}$ 
  assumes countable S and M:  $\bigwedge n::nat. \bigwedge x. x \in S \implies \text{norm}(f\ n\ x) \leq M$ 
  obtains k where strict_mono ( $k::nat \Rightarrow nat$ )  $\bigwedge x. x \in S \implies \exists l. (\lambda n. f\ (k\ n)\ x)$ 
   $\longrightarrow l$ 
  proof (cases S = {})
  case True
  then show ?thesis

```

```

    using strict_mono_id that by fastforce
next
case False
with ⟨countable S⟩ obtain σ :: nat ⇒ 'a where σ: S = range σ
  using uncountable_def by blast
obtain k where strict_mono k and k:  $\bigwedge i. \exists l. (\lambda n. (f \circ k) n (\sigma i)) \longrightarrow l$ 
proof (rule subsequence_diagonalization_lemma
  [of  $\lambda i r. \exists l. ((\lambda n. (f \circ r) n (\sigma i)) \longrightarrow l)$  sequentially id])
  show  $\exists k::nat \Rightarrow nat. \text{strict\_mono } k \wedge (\exists l. (\lambda n. (f \circ (r \circ k)) n (\sigma i)) \longrightarrow l)$ 
for i r
  proof -
    have  $f (r n) (\sigma i) \in \text{cball } 0 M$  for n
    by (simp add: σ M)
    then show ?thesis
    using compact_def [of cball (0::'b) M] by (force simp: o_def)
  qed
  show  $\exists l. (\lambda n. (f \circ (r \circ k2)) n (\sigma i)) \longrightarrow l$ 
  if  $\exists l. (\lambda n. (f \circ (r \circ k1)) n (\sigma i)) \longrightarrow l \wedge j. N \leq j \implies \exists j' \geq j. k2 j = k1 j'$ 
  for i N and r k1 k2 :: nat ⇒ nat
  using that
  by (simp add: lim_sequentially) (metis (no_types, opaque_lifting) le_cases
order_trans)
  qed auto
  with σ that show ?thesis
  by force
qed

```

theorem Arzela_Ascoli:

```

fixes  $\mathcal{F} :: [nat, 'a::euclidean\_space] \Rightarrow 'b::\{\text{real\_normed\_vector, heine\_borel}\}$ 
assumes compact S
  and  $M: \bigwedge n x. x \in S \implies \text{norm}(\mathcal{F} n x) \leq M$ 
  and equicont:
     $\bigwedge x e. \llbracket x \in S; 0 < e \rrbracket$ 
       $\implies \exists d. 0 < d \wedge (\forall n y. y \in S \wedge \text{norm}(x - y) < d \longrightarrow \text{norm}(\mathcal{F} n$ 
 $x - \mathcal{F} n y) < e)$ 
  obtains  $g k$  where continuous_on S g strict_mono (k :: nat ⇒ nat)
     $\bigwedge e. 0 < e \implies \exists N. \forall n x. n \geq N \wedge x \in S \longrightarrow \text{norm}(\mathcal{F}(k n) x -$ 
 $g x) < e$ 
proof -
  have UEQ:  $\bigwedge e. 0 < e \implies \exists d. 0 < d \wedge (\forall n. \forall x \in S. \forall x' \in S. \text{dist } x' x < d$ 
 $\longrightarrow \text{dist}(\mathcal{F} n x')(\mathcal{F} n x) < e)$ 
  apply (rule compact_uniformly_equicontinuous [OF ⟨compact S⟩, of range  $\mathcal{F}$ ])
  using equicont by (force simp: dist_commute dist_norm)+
  have continuous_on S g
  if  $\bigwedge e. 0 < e \implies \exists N. \forall n x. n \geq N \wedge x \in S \longrightarrow \text{norm}(\mathcal{F}(r n) x - g x) < e$ 
  for  $g:: 'a \Rightarrow 'b$  and  $r :: nat \Rightarrow nat$ 
proof (rule uniform_limit_theorem [of _  $\mathcal{F} \circ r$ ])
  have continuous_on S ( $\mathcal{F} (r n)$ ) for n

```

```

    using UEQ by (force simp: continuous_on_iff)
  then show  $\forall_F n$  in sequentially. continuous_on  $S ((\mathcal{F} \circ r) n)$ 
    by (simp add: eventually_sequentially)
  show uniform_limit  $S (\mathcal{F} \circ r) g$  sequentially
    using that by (metis (mono_tags, opaque_lifting) comp_apply dist_norm
uniform_limit_sequentially_iff)
  qed auto
  moreover
  obtain  $R$  where countable  $R$   $R \subseteq S$  and  $SR: S \subseteq \text{closure } R$ 
    by (metis separable that)
  obtain  $k$  where strict_mono  $k$  and  $k: \bigwedge x. x \in R \implies \exists l. (\lambda n. \mathcal{F} (k n) x) \longrightarrow l$ 
    using  $\langle R \subseteq S \rangle$  by (force intro: function_convergent_subsequence [OF  $\langle \text{countable } R \rangle M$ ])
  then have Cauchy: Cauchy  $((\lambda n. \mathcal{F} (k n) x))$  if  $x \in R$  for  $x$ 
    using convergent_eq_Cauchy that by blast
  have  $\exists N. \forall m n x. N \leq m \wedge N \leq n \wedge x \in S \longrightarrow \text{dist } ((\mathcal{F} \circ k) m x) ((\mathcal{F} \circ k) n x) < e$ 
    if  $0 < e$  for  $e$ 
  proof -
    obtain  $d$  where  $0 < d$ 
      and  $d: \bigwedge n. \forall x \in S. \forall x' \in S. \text{dist } x' x < d \longrightarrow \text{dist } (\mathcal{F} n x') (\mathcal{F} n x) < e/3$ 
      by (metis UEQ  $\langle 0 < e \rangle$  divide_pos_pos zero_less_numeral)
    obtain  $T$  where  $T \subseteq R$  and finite  $T$  and  $T: S \subseteq (\bigcup c \in T. \text{ball } c d)$ 
    proof (rule compactE_image [OF  $\langle \text{compact } S \rangle$ , of  $R (\lambda x. \text{ball } x d)$ ])
      have  $\text{closure } R \subseteq (\bigcup c \in R. \text{ball } c d)$ 
        using  $\langle 0 < d \rangle$  by (auto simp: closure_approachable)
      with  $SR$  show  $S \subseteq (\bigcup c \in R. \text{ball } c d)$ 
        by auto
    qed auto
    have  $\exists M. \forall m \geq M. \forall n \geq M. \text{dist } (\mathcal{F} (k m) x) (\mathcal{F} (k n) x) < e/3$  if  $x \in R$  for  $x$ 
      using Cauchy  $\langle 0 < e \rangle$  that unfolding Cauchy_def
      by (metis less_divide_eq_numeral1(1) mult_zero_left)
    then obtain  $MF$  where  $MF: \bigwedge x m n. \llbracket x \in R; m \geq MF x; n \geq MF x \rrbracket \implies \text{norm } (\mathcal{F} (k m) x - \mathcal{F} (k n) x) < e/3$ 
      using dist_norm by metis
    have  $\text{dist } ((\mathcal{F} \circ k) m x) ((\mathcal{F} \circ k) n x) < e$ 
      if  $m: \text{Max } (MF ' T) \leq m$  and  $n: \text{Max } (MF ' T) \leq n$   $x \in S$  for  $m n x$ 
    proof -
      obtain  $t$  where  $t \in T$  and  $t: x \in \text{ball } t d$ 
        using  $\langle x \in S \rangle T$  by auto
      have  $\text{norm } (\mathcal{F} (k m) t - \mathcal{F} (k m) x) < e / 3$ 
        by (metis  $\langle R \subseteq S \rangle \langle T \subseteq R \rangle \langle t \in T \rangle d \text{ dist_norm mem_ball subset_iff } t \langle x \in S \rangle$ )
      moreover
      have  $\text{norm } (\mathcal{F} (k n) t - \mathcal{F} (k n) x) < e / 3$ 
        by (metis  $\langle R \subseteq S \rangle \langle T \subseteq R \rangle \langle t \in T \rangle \text{subsetD } d \text{ dist_norm mem_ball } t \langle x \in S \rangle$ )
      moreover

```

```

have norm( $\mathcal{F}$  (k m) t -  $\mathcal{F}$  (k n) t) < e / 3
proof (rule MF)
  show t ∈ R
    using ⟨T ⊆ R⟩ ⟨t ∈ T⟩ by blast
  show MF t ≤ m MF t ≤ n
    by (meson Max_ge ⟨finite T⟩ ⟨t ∈ T⟩ finite_imageI imageI le_trans m
n)+
  qed
ultimately
show ?thesis
  unfolding dist_norm [symmetric] o_def
    by (metis dist_triangle_third dist_commute)
qed
then show ?thesis
  by force
qed
then obtain g where  $\forall e > 0. \exists N. \forall n x. N \leq n \wedge x \in S \longrightarrow \text{norm} ((\mathcal{F} \circ k) n$ 
 $x - g x) < e$ 
  using uniformly_convergent_eq_cauchy [of  $\lambda x. x \in S \mathcal{F} \circ k$ ] by (auto simp
add: dist_norm)
  ultimately show thesis
  by (metis ⟨strict_mono k⟩ comp_apply that)
qed

```

7.3.1 Montel's theorem

a sequence of holomorphic functions uniformly bounded on compact subsets of an open set S has a subsequence that converges to a holomorphic function, and converges *uniformly* on compact subsets of S .

theorem Montel:

```

fixes  $\mathcal{F} :: [\text{nat}, \text{complex}] \Rightarrow \text{complex}$ 
assumes open S
  and  $\mathcal{H}: \bigwedge h. h \in \mathcal{H} \implies h \text{ holomorphic\_on } S$ 
  and bounded:  $\bigwedge K. [\text{compact } K; K \subseteq S] \implies \exists B. \forall h \in \mathcal{H}. \forall z \in K. \text{norm}(h$ 
 $z) \leq B$ 
  and rng_f: range  $\mathcal{F} \subseteq \mathcal{H}$ 
obtains g r
  where g holomorphic_on S strict_mono (r :: nat  $\Rightarrow$  nat)
     $\bigwedge x. x \in S \implies ((\lambda n. \mathcal{F} (r n) x) \longrightarrow g x)$  sequentially
     $\bigwedge K. [\text{compact } K; K \subseteq S] \implies \text{uniform\_limit } K (\mathcal{F} \circ r) g$  sequentially

```

proof –

```

obtain K where comK:  $\bigwedge n. \text{compact}(K n)$  and KS:  $\bigwedge n::\text{nat}. K n \subseteq S$ 
  and subK:  $\bigwedge X. [\text{compact } X; X \subseteq S] \implies \exists N. \forall n \geq N. X \subseteq K n$ 
  using open_Union_compact_subsets [OF ⟨open S⟩] by metis
then have  $\bigwedge i. \exists B. \forall h \in \mathcal{H}. \forall z \in K i. \text{norm}(h z) \leq B$ 
  by (simp add: bounded)
then obtain B where B:  $\bigwedge i h z. [h \in \mathcal{H}; z \in K i] \implies \text{norm}(h z) \leq B$ 
  by metis

```

```

have *:  $\exists r$  g. strict_mono (r::nat $\Rightarrow$ nat)  $\wedge$  ( $\forall e > 0. \exists N. \forall n \geq N. \forall x \in K i.$ 
norm(( $\mathcal{F} \circ r$ ) n x - g x) < e)
  if  $\bigwedge n. \mathcal{F} n \in \mathcal{H}$  for  $\mathcal{F} i$ 
proof -
  obtain g k where continuous_on (K i) g strict_mono (k::nat $\Rightarrow$ nat)
     $\bigwedge e. 0 < e \implies \exists N. \forall n \geq N. \forall x \in K i. \text{norm}(\mathcal{F}(k n) x - g x) < e$ 
  proof (rule Arzela_Ascoli [of K i  $\mathcal{F}$  B i])
    show  $\exists d > 0. \forall n y. y \in K i \wedge \text{cmod} (z - y) < d \longrightarrow \text{cmod} (\mathcal{F} n z - \mathcal{F} n$ 
y) < e
      if  $z: z \in K i$  and  $0 < e$  for  $z e$ 
    proof -
      obtain r where  $0 < r$  and r: cball z r  $\subseteq$  S
        using z KS [of i]  $\langle$ open S $\rangle$  by (force simp: open_contains_cball)
      have cball z (2/3 * r)  $\subseteq$  cball z r
        using  $\langle 0 < r \rangle$  by (simp add: cball_subset_cball_iff)
      then have z23S: cball z (2/3 * r)  $\subseteq$  S
        using r by blast
      obtain M where  $0 < M$  and M:  $\bigwedge n w. \text{dist} z w \leq 2/3 * r \implies \text{norm}(\mathcal{F}$ 
n w)  $\leq$  M
        proof -
          obtain N where N:  $\forall n \geq N. \text{cball} z (2/3 * r) \subseteq K n$ 
            using subK compact_cball [of z (2/3 * r)] z23S by force
          have cmod ( $\mathcal{F} n w$ )  $\leq$  |B N| + 1 if dist z w  $\leq$  2/3 * r for n w
            proof -
              have w  $\in$  K N
                using N mem_cball that by blast
              then have cmod ( $\mathcal{F} n w$ )  $\leq$  B N
                using B  $\langle$  $\bigwedge n. \mathcal{F} n \in \mathcal{H}$  $\rangle$  by blast
              also have ...  $\leq$  |B N| + 1
                by simp
              finally show ?thesis .
            qed
          then show ?thesis
            by (rule_tac M=|B N| + 1 in that) auto
          qed
        have cmod ( $\mathcal{F} n z - \mathcal{F} n y$ ) < e
          if  $y \in K i$  and y_near_z: cmod (z - y) < r/3 cmod (z - y) < e * r
          / (6 * M)
          for  $n y$ 
        proof -
          have ( $(\lambda w. \mathcal{F} n w / (w - \xi))$  has_contour_integral
            ( $(2 * \pi) * i * \text{winding\_number} (\text{circlepath } z (2/3 * r)) \xi * \mathcal{F} n \xi$ 
            (circlepath z (2/3 * r)))
          if dist  $\xi z < (2/3 * r)$  for  $\xi$ 
        proof (rule Cauchy_integral_formula_convex_simple)
          have  $\mathcal{F} n$  holomorphic_on S
            by (simp add:  $\mathcal{H} \langle$  $\bigwedge n. \mathcal{F} n \in \mathcal{H}$  $\rangle$ )
          with z23S show  $\mathcal{F} n$  holomorphic_on cball z (2/3 * r)
            using holomorphic_on_subset by blast

```

```

qed (use that ⟨0 < r⟩ in ⟨auto simp: dist_commute⟩)
then have *: ((λw.  $\mathcal{F} \ n \ w / (w - \xi)$ ) has_contour_integral (2 * pi) * i
*  $\mathcal{F} \ n \ \xi$ )
      (circlepath z (2/3 * r))
if dist  $\xi \ z < (2/3 * r)$  for  $\xi$ 
using that by (simp add: winding_number_circlepath dist_norm)
have y: ((λw.  $\mathcal{F} \ n \ w / (w - y)$ ) has_contour_integral (2 * pi) * i *  $\mathcal{F} \ n \ y$ )
      (circlepath z (2/3 * r))
proof (rule *)
  show dist  $y \ z < 2/3 * r$ 
    using that ⟨0 < r⟩ by (simp only: dist_norm norm_minus_commute)
qed
have z: ((λw.  $\mathcal{F} \ n \ w / (w - z)$ ) has_contour_integral (2 * pi) * i *  $\mathcal{F} \ n \ z$ )
      (circlepath z (2/3 * r))
    using ⟨0 < r⟩ by (force intro!: *)
have le_er: cmod ( $\mathcal{F} \ n \ x / (x - y) - \mathcal{F} \ n \ x / (x - z)$ ) ≤ e / r
    if cmod (x - z) = r/3 + r/3 for x
proof -
  have ¬ (cmod (x - y) < r/3)
    using y_near_z(1) that ⟨M > 0⟩ ⟨r > 0⟩
  by (metis (full_types) norm_diff_triangle_less norm_minus_commute
order_less_irrefl)
  then have r4_le_xy: r/4 ≤ cmod (x - y)
    using ⟨r > 0⟩ by simp
  then have neq: x ≠ y x ≠ z
  using that ⟨r > 0⟩ by (auto simp: field_split_simps norm_minus_commute)
  have leM: cmod ( $\mathcal{F} \ n \ x$ ) ≤ M
    by (simp add: M dist_commute dist_norm that)
  have cmod ( $\mathcal{F} \ n \ x / (x - y) - \mathcal{F} \ n \ x / (x - z)$ ) = cmod ( $\mathcal{F} \ n \ x$ ) *
cmod (1 / (x - y) - 1 / (x - z))
  by (metis (no_types, lifting) divide_inverse mult.left_neutral norm_mult
right_diff_distrib')
  also have ... = cmod ( $\mathcal{F} \ n \ x$ ) * cmod ((y - z) / ((x - y) * (x - z)))
    using neq by (simp add: field_split_simps)
  also have ... = cmod ( $\mathcal{F} \ n \ x$ ) * (cmod (y - z) / (cmod(x - y) * (2/3
* r)))
    by (simp add: norm_mult norm_divide that)
  also have ... ≤ M * (cmod (y - z) / (cmod(x - y) * (2/3 * r)))
    using ⟨r > 0⟩ ⟨M > 0⟩ by (intro mult_mono [OF leM]) auto
  also have ... < M * ((e * r / (6 * M)) / (cmod(x - y) * (2/3 * r)))
    unfolding mult_less_cancel_left
    using y_near_z(2) ⟨M > 0⟩ ⟨r > 0⟩ neq
    by (simp add: field_simps mult_less_0_iff norm_minus_commute)
  also have ... ≤ e/r
    using ⟨e > 0⟩ ⟨r > 0⟩ r4_le_xy by (simp add: field_split_simps)
  finally show ?thesis by simp
qed

```

```

      have  $(2 * pi) * cmod (\mathcal{F} n y - \mathcal{F} n z) = cmod ((2 * pi) * i * \mathcal{F} n y -$ 
 $(2 * pi) * i * \mathcal{F} n z)$ 
      by (simp add: right_diff_distrib [symmetric] norm_mult)
      also have  $cmod ((2 * pi) * i * \mathcal{F} n y - (2 * pi) * i * \mathcal{F} n z) \leq e / r *$ 
 $(2 * pi * (2/3 * r))$ 

    proof (rule has_contour_integral_bound_circlepath [OF has_contour_integral_diff
[OF y z]])
      show  $\bigwedge x. cmod (x - z) = 2/3 * r \implies cmod (\mathcal{F} n x / (x - y) - \mathcal{F} n$ 
 $x / (x - z)) \leq e / r$ 
      using le_er by auto
      qed (use <e > 0> <r > 0> in auto)
      also have  $\dots = (2 * pi) * e * ((2/3))$ 
      using <r > 0> by (simp add: field_split_simps)
      finally have  $cmod (\mathcal{F} n y - \mathcal{F} n z) \leq e * (2/3)$ 
      by simp
      also have  $\dots < e$ 
      using <e > 0> by simp
      finally show ?thesis by (simp add: norm_minus_commute)
    qed
    then show ?thesis
    apply (rule_tac x=min (r/3) ((e * r)/(6 * M)) in exI)
    using <0 < e> <0 < r> <0 < M> by simp
  qed
  show  $\bigwedge n x. x \in K i \implies cmod (\mathcal{F} n x) \leq B i$ 
  using B < $\bigwedge n. \mathcal{F} n \in \mathcal{H}$ > by blast
next
  fix g :: complex  $\Rightarrow$  complex and k :: nat  $\Rightarrow$  nat
  assume *:  $\bigwedge (g :: complex \Rightarrow complex) (k :: nat \Rightarrow nat). continuous\_on (K i) g \implies$ 
  strict_mono k  $\implies$ 
  ( $\bigwedge e. 0 < e \implies \exists N. \forall n \geq N. \forall x \in K i. cmod (\mathcal{F} (k n) x - g x) <$ 
 $e) \implies thesis$ 
  continuous_on (K i) g
  strict_mono k
   $\bigwedge e. 0 < e \implies \exists N. \forall n x. N \leq n \wedge x \in K i \longrightarrow cmod (\mathcal{F} (k n) x - g$ 
 $x) < e$ 
  show ?thesis
  by (rule *(1)[OF *(2,3)], drule *(4)) auto
  qed (use comK in simp_all)
  then show ?thesis
  by auto
  qed
  define  $\Phi$  where  $\Phi \equiv \lambda g i r. \lambda k :: nat \Rightarrow nat. \forall e > 0. \exists N. \forall n \geq N. \forall x \in K i. cmod$ 
 $((\mathcal{F} \circ (r \circ k)) n x - g x) < e$ 
  obtain k :: nat  $\Rightarrow$  nat where strict_mono k and k:  $\bigwedge i. \exists g. \Phi g i id k$ 
  proof (rule subsequence_diagonalization_lemma [where r=id])
    show  $\exists g. \Phi g i id (r \circ k2)$ 
    if ex:  $\exists g. \Phi g i id (r \circ k1)$  and  $\bigwedge j. N \leq j \implies \exists j' \geq j. k2 j = k1 j'$ 
    for i k1 k2 N and r :: nat  $\Rightarrow$  nat

```



```

proof -
  obtain g where  $\Phi$  g i id (r o k1)
    using ex by blast
  then have  $\Phi$  g i id (r o k2)
    using that
    by (simp add:  $\Phi$ _def) (metis (no_types, opaque_lifting) le_trans linear)
  then show ?thesis
    by metis
qed
have  $\exists k$  g. strict_mono (k::nat $\Rightarrow$ nat)  $\wedge$   $\Phi$  g i id (r o k) for i r
  unfolding  $\Phi$ _def o_assoc using rng_f by (force intro!: *)
  then show  $\bigwedge i$  r.  $\exists k$ . strict_mono (k::nat $\Rightarrow$ nat)  $\wedge$  ( $\exists g$ .  $\Phi$  g i id (r o k))
    by force
qed fastforce
have  $\exists l$ .  $\forall e > 0$ .  $\exists N$ .  $\forall n \geq N$ . norm( $\mathcal{F}$  (k n) z - l) < e if z  $\in$  S for z
proof -
  obtain G where G:  $\bigwedge i$  e. e > 0  $\implies$   $\exists M$ .  $\forall n \geq M$ .  $\forall x \in K$  i. cmod (( $\mathcal{F}$  o k) n
x - G i x) < e
    using k unfolding  $\Phi$ _def by (metis id_comp)
  obtain N where  $\bigwedge n$ . n  $\geq N \implies z \in K$  n
    using subK [of {z}] that  $\langle z \in S \rangle$  by auto
  moreover have  $\bigwedge e$ . e > 0  $\implies$   $\exists M$ .  $\forall n \geq M$ .  $\forall x \in K$  N. cmod (( $\mathcal{F}$  o k) n x -
G N x) < e
    using G by auto
  ultimately show ?thesis
    by (metis comp_apply order_refl)
qed
then obtain g where g:  $\bigwedge z$  e. [ $z \in S$ ; e > 0]  $\implies$   $\exists N$ .  $\forall n \geq N$ . norm( $\mathcal{F}$  (k n)
z - g z) < e
  by metis
show ?thesis
proof
  show g_lim:  $\bigwedge x$ . x  $\in$  S  $\implies$  ( $\lambda n$ .  $\mathcal{F}$  (k n) x)  $\longrightarrow$  g x
    by (simp add: lim_sequentially g dist_norm)
  have dg_le_e:  $\exists N$ .  $\forall n \geq N$ .  $\forall x \in T$ . cmod ( $\mathcal{F}$  (k n) x - g x) < e
    if T: compact T T  $\subseteq$  S and 0 < e for T e
  proof -
    obtain N where N:  $\bigwedge n$ . n  $\geq N \implies T \subseteq K$  n
      using subK [OF T] by blast
    obtain h where h:  $\bigwedge e$ . e > 0  $\implies$   $\exists M$ .  $\forall n \geq M$ .  $\forall x \in K$  N. cmod (( $\mathcal{F}$  o k) n x
- h x) < e
      using k unfolding  $\Phi$ _def by (metis id_comp)
    have geq: g w = h w if w  $\in$  T for w
      proof (rule LIMSEQ_unique)
        show ( $\lambda n$ .  $\mathcal{F}$  (k n) w)  $\longrightarrow$  g w
          using  $\langle T \subseteq S \rangle$  g_lim that by blast
        show ( $\lambda n$ .  $\mathcal{F}$  (k n) w)  $\longrightarrow$  h w
          using h N that by (force simp: lim_sequentially dist_norm)
      qed
  qed

```

```

show ?thesis
using T h N ‹0 < e› by (fastforce simp add: geq)
qed
then show  $\bigwedge K. \llbracket \text{compact } K; K \subseteq S \rrbracket$ 
   $\implies \text{uniform\_limit } K (\mathcal{F} \circ k) g \text{ sequentially}$ 
by (simp add: uniform_limit_iff dist_norm eventually_sequentially)
show g holomorphic_on S
proof (rule holomorphic_uniform_sequence [OF ‹open S› H])
show  $\bigwedge n. (\mathcal{F} \circ k) n \in \mathcal{H}$ 
by (simp add: range_subsetD rng_f)
show  $\exists d > 0. \text{cball } z d \subseteq S \wedge \text{uniform\_limit } (\text{cball } z d) (\lambda n. (\mathcal{F} \circ k) n) g$ 
  sequentially
if z ∈ S for z
proof –
obtain d where d: d > 0 cball z d ⊆ S
using ‹open S› ‹z ∈ S› open_contains_cball by blast
then have uniform_limit (cball z d) (F ∘ k) g sequentially
using dg_le_e compact_cball by (auto simp: uniform_limit_iff eventually_sequentially dist_norm)
with d show ?thesis by blast
qed
qed
qed (auto simp: ‹strict_mono k›)
qed

```

7.4 Some simple but useful cases of Hurwitz's theorem

proposition Hurwitz_no_zeros:

assumes S: open S connected S

and holf: $\bigwedge n::\text{nat}. \mathcal{F} n \text{ holomorphic_on } S$

and holg: g holomorphic_on S

and ul_g: $\bigwedge K. \llbracket \text{compact } K; K \subseteq S \rrbracket \implies \text{uniform_limit } K \mathcal{F} g \text{ sequentially}$

and nonconst: $\neg g \text{ constant_on } S$

and nz: $\bigwedge n z. z \in S \implies \mathcal{F} n z \neq 0$

and z0 ∈ S

shows g z0 ≠ 0

proof

assume g0: g z0 = 0

obtain h r m

where 0 < m 0 < r **and** subS: ball z0 r ⊆ S

and holh: h holomorphic_on ball z0 r

and geq: $\bigwedge w. w \in \text{ball } z0 r \implies g w = (w - z0)^m * h w$

and hnz: $\bigwedge w. w \in \text{ball } z0 r \implies h w \neq 0$

by (blast intro: holomorphic_factor_zero_nonconstant [OF holg S ‹z0 ∈ S› g0 nonconst])

then have holf0: F n holomorphic_on ball z0 r **for** n

by (meson holf holomorphic_on_subset)

have *: ((λz. deriv (F n) z / F n z) has_contour_integral 0) (circlepath z0 (r/2)) **for** n

```

proof (rule Cauchy_theorem_disc_simple)
  show ( $\lambda z. \text{deriv } (\mathcal{F} \ n) \ z / \mathcal{F} \ n \ z$ ) holomorphic_on ball z0 r
  by (metis (no_types) ‹open S› holf holomorphic_deriv holomorphic_on_divide
holomorphic_on_subset nz subS)
qed (use ‹0 < r› in auto)
have hol_dg: deriv g holomorphic_on S
  by (simp add: ‹open S› holg holomorphic_deriv)
have continuous_on (sphere z0 (r/2)) (deriv g)
  using ‹0 < r› subS
  by (intro holomorphic_on_imp_continuous_on holomorphic_on_subset [OF
hol_dg]) auto
then have compact (deriv g ‘ (sphere z0 (r/2)))
  by (rule compact_continuous_image [OF __ compact_sphere])
then have bo_dg: bounded (deriv g ‘ (sphere z0 (r/2)))
  using compact_imp_bounded by blast
have continuous_on (sphere z0 (r/2)) (cmod ∘ g)
  using ‹0 < r› subS
  by (intro continuous_intros holomorphic_on_imp_continuous_on holomor-
phic_on_subset [OF holg]) auto
then have compact ((cmod ∘ g) ‘ sphere z0 (r/2))
  by (rule compact_continuous_image [OF __ compact_sphere])
moreover have (cmod ∘ g) ‘ sphere z0 (r/2) ≠ {}
  using ‹0 < r› by auto
ultimately obtain b where b: b ∈ (cmod ∘ g) ‘ sphere z0 (r/2)
   $\wedge t. t \in (cmod \circ g) \text{ ‘ sphere } z0 \ (r/2) \implies b \leq t$ 
  using compact_attains_inf [of (norm ∘ g) ‘ (sphere z0 (r/2))] by blast
have ( $\lambda n. \text{contour\_integral } (\text{circlepath } z0 \ (r/2)) \ (\lambda z. \text{deriv } (\mathcal{F} \ n) \ z / \mathcal{F} \ n \ z)$ )
 $\longrightarrow$ 
  contour_integral (circlepath z0 (r/2)) ( $\lambda z. \text{deriv } g \ z / g \ z$ )
proof (rule contour_integral_uniform_limit_circlepath)
  show  $\forall_F n$  in sequentially. ( $\lambda z. \text{deriv } (\mathcal{F} \ n) \ z / \mathcal{F} \ n \ z$ ) contour_integrable_on
circlepath z0 (r/2)
  using * contour_integrable_on_def eventually_sequentiallyI by meson
  show uniform_limit (sphere z0 (r/2)) ( $\lambda n \ z. \text{deriv } (\mathcal{F} \ n) \ z / \mathcal{F} \ n \ z$ ) ( $\lambda z. \text{deriv } g \ z / g \ z$ )
sequentially
proof (rule uniform_lim_divide [OF __ bo_dg])
  show uniform_limit (sphere z0 (r/2)) ( $\lambda a. \text{deriv } (\mathcal{F} \ a)$ ) (deriv g) sequentially
proof (rule uniform_limitI)
  fix e::real
  assume 0 < e

  show  $\forall_F n$  in sequentially.  $\forall x \in \text{sphere } z0 \ (r/2). \text{dist } (\text{deriv } (\mathcal{F} \ n) \ x)$ 
(deriv g x) < e
proof –
  have dist (deriv (F n) w) (deriv g w) < e
  if e8:  $\bigwedge x. \text{dist } z0 \ x \leq 3 * r / 4 \implies \text{dist } (\mathcal{F} \ n \ x) \ (g \ x) * 8 < r * e$ 
  and w: w ∈ sphere z0 (r/2) for n w
proof –
  have ball w (r/4) ⊆ ball z0 r cball w (r/4) ⊆ ball z0 r

```

```

    using ‹0 < r› w by (simp_all add: ball_subset_ball_iff cball_subset_ball_iff
dist_commute)
  with subS have wr4_sub: ball w (r/4) ⊆ S cball w (r/4) ⊆ S by force+
  moreover
  have (λz. ℱ n z - g z) holomorphic_on S
    by (intro holomorphic_intros holf holg)
  ultimately have hol: (λz. ℱ n z - g z) holomorphic_on ball w (r/4)
    and cont: continuous_on (cball w (r / 4)) (λz. ℱ n z - g z)
  using holomorphic_on_subset by (blast intro: holomorphic_on_imp_continuous_on)+
  have w ∈ S
    using ‹0 < r› wr4_sub by auto
  have dist z0 y ≤ 3 * r / 4 if dist w y < r/4 for y
  proof (rule dist_triangle_le [where z=w])
    show dist z0 w + dist y w ≤ 3 * r / 4
      using w that by (simp add: dist_commute)
  qed
  with e8 have in_ball: ∧y. y ∈ ball w (r/4) ⇒ ℱ n y - g y ∈ ball 0
(r/4 * e/2)
    by (simp add: dist_norm [symmetric])
  have ℱ n field_differentiable at w
    by (metis holomorphic_on_imp_differentiable_at ‹w ∈ S› holf ‹open
S›)
  moreover
  have g field_differentiable at w
    using ‹w ∈ S› ‹open S› holg holomorphic_on_imp_differentiable_at
by auto
  moreover
  have cmod (deriv (λw. ℱ n w - g w) w) * 2 ≤ e
    using Cauchy_higher_deriv_bound [OF hol cont in_ball, of 1] ‹r >
0› by auto
  ultimately have dist (deriv (ℱ n) w) (deriv g w) ≤ e/2
    by (simp add: dist_norm)
  then show ?thesis
    using ‹e > 0› by auto
  qed
  moreover
  have cball z0 (3 * r / 4) ⊆ ball z0 r
    by (simp add: cball_subset_ball_iff ‹0 < r›)
  with subS have uniform_limit (cball z0 (3 * r/4)) ℱ g sequentially
    by (force intro: ul_g)
  then have ∀_F n in sequentially. ∀ x∈cball z0 (3 * r / 4). dist (ℱ n x) (g
x) < r / 4 * e / 2
    using ‹0 < e› ‹0 < r› by (force simp: intro!: uniform_limitD)
  ultimately show ?thesis
    by (force simp add: eventually_sequentially)
  qed
  qed
  show uniform_limit (sphere z0 (r/2)) ℱ g sequentially
  proof (rule uniform_limitI)

```

```

    fix e::real
    assume 0 < e
    have sphere z0 (r/2) ⊆ ball z0 r
      using ‹0 < r› by auto
    with subS have uniform_limit (sphere z0 (r/2))  $\mathcal{F}$  g sequentially
      by (force intro: ul_g)
    then show  $\forall_F n$  in sequentially.  $\forall x \in$  sphere z0 (r/2). dist ( $\mathcal{F}$  n x) (g x)
  < e
    using ‹0 < e› uniform_limit_iff by blast
  qed
  show b > 0  $\wedge$  x. x ∈ sphere z0 (r/2)  $\implies$  b ≤ cmod (g x)
    using b ‹0 < r› by (fastforce simp: geq hnz)+
  qed
  qed (use ‹0 < r› in auto)
  then have (λn. 0)  $\longrightarrow$  contour_integral (circlepath z0 (r/2)) (λz. deriv g z
/ g z)
    by (simp add: contour_integral_unique [OF *])
  then have contour_integral (circlepath z0 (r/2)) (λz. deriv g z / g z) = 0
    by (simp add: LIMSEQ_const_iff)
  moreover
  have contour_integral (circlepath z0 (r/2)) (λz. deriv g z / g z) =
    contour_integral (circlepath z0 (r/2)) (λz. m / (z - z0) + deriv h z / h z)
  proof (rule contour_integral_eq, use ‹0 < r› in simp)
    fix w
    assume w: dist z0 w * 2 = r
    then have w_inb: w ∈ ball z0 r
      using ‹0 < r› by auto
    have h_der: (h has_field_derivative deriv h w) (at w)
      using holh holomorphic_derivI w_inb by blast
    have deriv g w = ((of_nat m * h w + deriv h w * (w - z0)) * (w - z0) ^ m)
/ (w - z0)
      if r = dist z0 w * 2 w ≠ z0
    proof -
      have ((λw. (w - z0) ^ m * h w) has_field_derivative
(m * h w + deriv h w * (w - z0)) * (w - z0) ^ m / (w - z0)) (at w)
        apply (rule derivative_eq_intros h_der refl)+
        using that ‹m > 0› ‹0 < r› apply (simp add: divide_simps distrib_right)
        by (metis Suc_pred mult.commute power_Suc)
      then show ?thesis
    proof (rule DERIV_imp_deriv [OF has_field_derivative_transform_within_open])
      show  $\wedge$ x. x ∈ ball z0 r  $\implies$  (x - z0) ^ m * h x = g x
        by (simp add: hnz geq)
      qed (use that ‹m > 0› ‹0 < r› in auto)
    qed
  with ‹0 < r› ‹0 < m› w w_inb show deriv g w / g w = of_nat m / (w - z0)
+ deriv h w / h w
    by (auto simp: geq field_split_simps hnz)
  qed
  moreover

```

```

have contour_integral (circlepath z0 (r/2)) ( $\lambda z. m / (z - z0) + \text{deriv } h z / h z$ ) =
  2 * of_real pi * i * m + 0
proof (rule contour_integral_unique [OF has_contour_integral_add])
  show ( $(\lambda x. m / (x - z0))$  has_contour_integral 2 * of_real pi * i * m)
(circlepath z0 (r/2))
  by (force simp:  $\langle 0 < r \rangle$  intro: Cauchy_integral_circlepath_simple)
  show ( $(\lambda x. \text{deriv } h x / h x)$  has_contour_integral 0) (circlepath z0 (r/2))
  using hnz holh holomorphic_deriv holomorphic_on_divide  $\langle 0 < r \rangle$ 
  by (fastforce intro!: Cauchy_theorem_disc_simple [of  $z0$  r])
qed
ultimately show False using  $\langle 0 < m \rangle$  by auto
qed

```

corollary *Hurwitz_injective*:

```

assumes S: open S connected S
  and holf:  $\bigwedge n::\text{nat}. \mathcal{F} \ n$  holomorphic_on S
  and holg: g holomorphic_on S
  and ulg:  $\bigwedge K. \llbracket \text{compact } K; K \subseteq S \rrbracket \implies \text{uniform\_limit } K \ \mathcal{F} \ g$  sequentially
  and nonconst:  $\neg$  g constant_on S
  and inj:  $\bigwedge n. \text{inj\_on } (\mathcal{F} \ n) \ S$ 
shows inj_on g S
proof -
  have False if z12: z1  $\in$  S z2  $\in$  S z1  $\neq$  z2 g z2 = g z1 for z1 z2
  proof -
    obtain z0 where z0  $\in$  S and z0: g z0  $\neq$  g z2
    using constant_on_def nonconst by blast
    have ( $\lambda z. g \ z - g \ z1$ ) holomorphic_on S
    by (intro holomorphic_intros holg)
    then obtain r where  $0 < r$  ball z2 r  $\subseteq$  S  $\bigwedge z. \text{dist } z2 \ z < r \wedge z \neq z2 \implies g$ 
z  $\neq$  g z1
    apply (rule isolated_zeros [of  $\lambda z. g \ z - g \ z1$  S z2 z0])
    using S  $\langle z0 \in S \rangle$  z0 z12 by auto
    have g z2 - g z1  $\neq$  0
    proof (rule Hurwitz_no_zeros [of S - {z1}  $\lambda n \ z. \mathcal{F} \ n \ z - \mathcal{F} \ n \ z1$   $\lambda z. g \ z -$ 
g z1])
      show open (S - {z1})
      by (simp add: S open_delete)
      show connected (S - {z1})
      by (simp add: connected_open_delete [OF S])
      show  $\bigwedge n. (\lambda z. \mathcal{F} \ n \ z - \mathcal{F} \ n \ z1)$  holomorphic_on S - {z1}
      by (intro holomorphic_intros holomorphic_on_subset [OF holf]) blast
      show ( $\lambda z. g \ z - g \ z1$ ) holomorphic_on S - {z1}
      by (intro holomorphic_intros holomorphic_on_subset [OF holg]) blast
      show uniform_limit K ( $\lambda n \ z. \mathcal{F} \ n \ z - \mathcal{F} \ n \ z1$ ) ( $\lambda z. g \ z - g \ z1$ ) sequentially
      if compact K K  $\subseteq$  S - {z1} for K
    proof (rule uniform_limitI)
      fix e::real
      assume e > 0

```

```

have uniform_limit K  $\mathcal{F}$  g sequentially
  using that ul_g by fastforce
then have K:  $\forall_F n$  in sequentially.  $\forall x \in K$ .  $\text{dist } (\mathcal{F} n x) (g x) < e/2$ 
  using  $\langle 0 < e \rangle$  by (force simp: intro!: uniform_limitD)
have uniform_limit {z1}  $\mathcal{F}$  g sequentially
  by (simp add: ul_g z12)
then have  $\forall_F n$  in sequentially.  $\forall x \in \{z1\}$ .  $\text{dist } (\mathcal{F} n x) (g x) < e/2$ 
  using  $\langle 0 < e \rangle$  by (force simp: intro!: uniform_limitD)
then have z1:  $\forall_F n$  in sequentially.  $\text{dist } (\mathcal{F} n z1) (g z1) < e/2$ 
  by simp
show  $\forall_F n$  in sequentially.  $\forall x \in K$ .  $\text{dist } (\mathcal{F} n x - \mathcal{F} n z1) (g x - g z1) < e$ 
  apply (rule eventually_mono [OF eventually_conj [OF K z1]])
  by (metis (no_types, opaque_lifting) diff_add_eq diff_diff_eq2 dist_commute
dist_norm dist_triangle_add_half)
qed
show  $\neg (\lambda z. g z - g z1)$  constant_on  $S - \{z1\}$ 
  unfolding constant_on_def
  by (metis Diff_iff  $\langle z0 \in S \rangle$  empty_iff insert_iff right_minus_eq z0 z12)
show  $\bigwedge n z. z \in S - \{z1\} \implies \mathcal{F} n z - \mathcal{F} n z1 \neq 0$ 
  by (metis DiffD1 DiffD2 eq_iff_diff_eq_0 inj inj_onD insertI1  $\langle z1 \in S \rangle$ )
show  $z2 \in S - \{z1\}$ 
  using  $\langle z2 \in S \rangle \langle z1 \neq z2 \rangle$  by auto
qed
with z12 show False by auto
qed
then show ?thesis by (auto simp: inj_on_def)
qed

```

7.5 The Great Picard theorem

lemma GPicard1:

```

assumes S: open S connected S and w  $\in S$   $0 < r$   $Y \subseteq X$ 
  and holX:  $\bigwedge h. h \in X \implies h$  holomorphic_on S
  and X01:  $\bigwedge h z. \llbracket h \in X; z \in S \rrbracket \implies h z \neq 0 \wedge h z \neq 1$ 
  and r:  $\bigwedge h. h \in Y \implies \text{norm}(h w) \leq r$ 
obtains B Z where  $0 < B$  open Z  $w \in Z$   $Z \subseteq S$   $\bigwedge h z. \llbracket h \in Y; z \in Z \rrbracket \implies$ 
norm(h z)  $\leq B$ 
proof -
  obtain e where  $e > 0$  and e: cball w e  $\subseteq S$ 
    using assms open_contains_cball_eq by blast
  show ?thesis
  proof
    show  $0 < \exp(\pi * \exp(\pi * (2 + 2 * r + 12)))$ 
      by simp
    show ball w (e / 2)  $\subseteq S$ 
      using e ball_divide_subset_numeral ball_subset_cball by blast
    show cmod (h z)  $\leq \exp(\pi * \exp(\pi * (2 + 2 * r + 12)))$ 
      if  $h \in Y$   $z \in \text{ball } w (e / 2)$  for h z
    proof -

```

```

have  $h \in X$ 
  using  $\langle Y \subseteq X \rangle \langle h \in Y \rangle$  by blast
have  $hol\_h\_o$ :  $(h \circ (\lambda z. (w + of\_real\ e * z)))$  holomorphic_on cball 0 1
proof (intro holomorphic_intros holomorphic_on_compose)
  have  $h$  holomorphic_on S
    using  $holX \langle h \in X \rangle$  by auto
  then have  $h$  holomorphic_on cball w e
    by (metis e holomorphic_on_subset)
  moreover have  $(\lambda z. w + complex\_of\_real\ e * z)$  ‘cball 0 1  $\subseteq$  cball w e
    using that  $\langle e > 0 \rangle$  by (auto simp: dist_norm norm_mult)
  ultimately show  $h$  holomorphic_on  $(\lambda z. w + complex\_of\_real\ e * z)$  ‘
cball 0 1
    by (rule holomorphic_on_subset)
qed
have  $norm\_le\_r$ :  $cmod ((h \circ (\lambda z. w + complex\_of\_real\ e * z))\ 0) \leq r$ 
  by (auto simp: r  $\langle h \in Y \rangle$ )
have  $le12$ :  $norm (of\_real (inverse\ e) * (z - w)) \leq 1/2$ 
using that  $\langle e > 0 \rangle$  by (simp add: inverse_eq_divide dist_norm norm_minus_commute
norm_divide)
  have  $non01$ :  $h (w + e * z) \neq 0 \wedge h (w + e * z) \neq 1$  if  $z \in cball\ 0\ 1$  for
 $z::complex$ 
proof (rule X01 [OF  $\langle h \in X \rangle$ ])
  have  $w + complex\_of\_real\ e * z \in cball\ w\ e$ 
    using  $\langle 0 < e \rangle$  that by (auto simp: dist_norm norm_mult)
  then show  $w + complex\_of\_real\ e * z \in S$ 
    by (rule subsetD [OF e])
qed
have  $cmod (h\ z) \leq cmod (h (w + of\_real\ e * (inverse\ e * (z - w))))$ 
  using  $\langle 0 < e \rangle$  by (simp add: field_split_simps)
also have  $\dots \leq exp (pi * exp (pi * (14 + 2 * r)))$ 
  using  $r$  [OF  $\langle h \in Y \rangle$ ] Schottky [OF hol_h_o norm_le_r ___ le12]  $non01$ 
by auto
  finally
    show ?thesis by simp
qed
qed (use  $\langle e > 0 \rangle$  in auto)
qed

```

lemma *GPicard2*:

```

assumes  $S \subseteq T$  connected  $T\ S \neq \{\}$  open  $S \wedge x. \llbracket x \text{ islimpt } S; x \in T \rrbracket \implies x \in S$ 
shows  $S = T$ 
by (metis assms open_subset connected_clopen closedin_limpt)

```

lemma *GPicard3*:

```

assumes  $S$ : open  $S$  connected  $S\ w \in S$  and  $Y \subseteq X$ 
and  $holX$ :  $\bigwedge h. h \in X \implies h$  holomorphic_on  $S$ 
and  $X01$ :  $\bigwedge h\ z. \llbracket h \in X; z \in S \rrbracket \implies h\ z \neq 0 \wedge h\ z \neq 1$ 
and  $no\_hw\_le1$ :  $\bigwedge h. h \in Y \implies norm(h\ w) \leq 1$ 

```



```

    and compact K K ⊆ S
  obtains B where  $\bigwedge h z. \llbracket h \in Y; z \in K \rrbracket \implies \text{norm}(h z) \leq B$ 
proof -
  define U where  $U \equiv \{z \in S. \exists B Z. 0 < B \wedge \text{open } Z \wedge z \in Z \wedge Z \subseteq S \wedge$ 
     $(\forall h z'. h \in Y \wedge z' \in Z \implies \text{norm}(h z') \leq B)\}$ 
  then have  $U \subseteq S$  by blast
  have  $U = S$ 
proof (rule GPicard2 [OF  $\langle U \subseteq S \rangle \langle \text{connected } S \rangle$ ])
  show  $U \neq \{\}$ 
  proof -
    obtain B Z where  $0 < B$  open Z  $w \in Z$   $Z \subseteq S$ 
    and  $\bigwedge h z. \llbracket h \in Y; z \in Z \rrbracket \implies \text{norm}(h z) \leq B$ 
    using GPicard1 [OF S zero_less_one  $\langle Y \subseteq X \rangle$  holX] X01 no_hw_le1 by
blast
  then show ?thesis
    unfolding U_def using  $\langle w \in S \rangle$  by blast
  qed
  show open U
    unfolding open_subopen [of U] by (auto simp: U_def)
  fix v
  assume v:  $v \text{ islimpt } U$   $v \in S$ 
  have  $\neg (\forall r > 0. \exists h \in Y. r < \text{cmod } (h v))$ 
  proof
    assume  $\forall r > 0. \exists h \in Y. r < \text{cmod } (h v)$ 
    then have  $\forall n. \exists h \in Y. \text{Suc } n < \text{cmod } (h v)$ 
      by simp
    then obtain F where  $F Y: \bigwedge n. F n \in Y$  and  $ltF: \bigwedge n. \text{Suc } n < \text{cmod } (F$ 
n v)
      by metis
    define G where  $G \equiv \lambda n z. \text{inverse}(F n z)$ 
    have  $\text{hol } G: G n \text{ holomorphic\_on } S$  for n
    proof (simp add: G_def)
      show  $(\lambda z. \text{inverse } (F n z)) \text{ holomorphic\_on } S$ 
        using FY X01  $\langle Y \subseteq X \rangle$  holX by (blast intro: holomorphic_on_inverse)
    qed
    have  $G \text{ not } 0: G n z \neq 0$  and  $G \text{ not } 1: G n z \neq 1$  if  $z \in S$  for n z
      using FY X01  $\langle Y \subseteq X \rangle$  that by (force simp: G_def)+
    have  $G\_le1: \text{cmod } (G n v) \leq 1$  for n
      using less_le_trans linear ltF
      by (fastforce simp add: G_def norm_inverse inverse_le_1_iff)
    define W where  $W \equiv \{h. h \text{ holomorphic\_on } S \wedge (\forall z \in S. h z \neq 0 \wedge h z$ 
 $\neq 1)\}$ 
    obtain B Z where  $0 < B$  open Z  $v \in Z$   $Z \subseteq S$ 
      and  $B: \bigwedge h z. \llbracket h \in \text{range } G; z \in Z \rrbracket \implies \text{norm}(h z) \leq B$ 
    apply (rule GPicard1 [OF  $\langle \text{open } S \rangle \langle \text{connected } S \rangle \langle v \in S \rangle$  zero_less_one,
of range G W])
      using holG Gnot0 Gnot1 G_le1 by (force simp: W_def)+
    then obtain e where  $e > 0$  and  $e: \text{ball } v e \subseteq Z$ 
      by (meson open_contains_ball)

```

```

obtain  $h j$  where  $holh$ :  $h$  holomorphic_on ball  $v e$  and  $strict\_mono j$ 
and  $lim$ :  $\bigwedge x. x \in ball\ v\ e \implies (\lambda n. \mathcal{G}\ (j\ n)\ x) \longrightarrow h\ x$ 
and  $ulim$ :  $\bigwedge K. \llbracket compact\ K; K \subseteq ball\ v\ e \rrbracket$ 
 $\implies uniform\_limit\ K\ (\mathcal{G} \circ j)\ h$  sequentially
proof (rule Montel)
  show  $\bigwedge h. h \in range\ \mathcal{G} \implies h$  holomorphic_on ball  $v e$ 
    by (metis  $\langle Z \subseteq S \rangle e\ hol\mathcal{G}$  holomorphic_on_subset imageE)
  show  $\bigwedge K. \llbracket compact\ K; K \subseteq ball\ v\ e \rrbracket \implies \exists B. \forall h \in range\ \mathcal{G}. \forall z \in K. cmod$ 
 $(h\ z) \leq B$ 
    using  $B\ e$  by blast
qed auto
have  $h\ v = 0$ 
proof (rule LIMSEQ_unique)
  show  $(\lambda n. \mathcal{G}\ (j\ n)\ v) \longrightarrow h\ v$ 
    using  $\langle e > 0 \rangle\ lim$  by simp
  have  $lt\_Fj$ :  $real\ x \leq cmod\ (\mathcal{F}\ (j\ x)\ v)$  for  $x$ 
    by (metis of_nat_Suc ltF  $\langle strict\_mono\ j \rangle$  add.commute less_eq_real_def
less_le_trans nat_le_real less_seq_suble)
  show  $(\lambda n. \mathcal{G}\ (j\ n)\ v) \longrightarrow 0$ 
proof (rule Lim_null_comparison [OF eventually_sequentiallyI lim_inverse_n])
  show  $cmod\ (\mathcal{G}\ (j\ x)\ v) \leq inverse\ (real\ x)$  if  $1 \leq x$  for  $x$ 
    using that by (simp add:  $\mathcal{G\_def}$  norm_inverse_le_norm [OF lt_Fj])
qed
qed
have  $h\ v \neq 0$ 
proof (rule Hurwitz_no_zeros [of ball v e  $\mathcal{G} \circ j\ h$ ])
  show  $\bigwedge n. (\mathcal{G} \circ j)\ n$  holomorphic_on ball  $v e$ 
    using  $\langle Z \subseteq S \rangle e\ hol\mathcal{G}$  by force
  show  $\bigwedge n\ z. z \in ball\ v\ e \implies (\mathcal{G} \circ j)\ n\ z \neq 0$ 
    using  $\mathcal{G}not0\ \langle Z \subseteq S \rangle e$  by fastforce
  show  $\neg h$  constant_on ball  $v e$ 
proof (clarsimp simp: constant_on_def)
  fix  $c$ 
  have False if  $\bigwedge z. dist\ v\ z < e \implies h\ z = c$ 
proof –
  have  $h\ v = c$ 
    by (simp add:  $\langle 0 < e \rangle$  that)
  obtain  $y$  where  $y \in U\ y \neq v$  and  $y$ :  $dist\ y\ v < e$ 
    using  $v\ \langle e > 0 \rangle$  by (auto simp: islimpt_approachable)
  then obtain  $C\ T$  where  $y \in S\ C > 0$  open  $T\ y \in T\ T \subseteq S$ 
and  $\bigwedge h\ z'. \llbracket h \in Y; z' \in T \rrbracket \implies cmod\ (h\ z') \leq C$ 
    using  $\langle y \in U \rangle$  by (auto simp: U_def)
  then have  $le\_C$ :  $\bigwedge n. cmod\ (\mathcal{F}\ n\ y) \leq C$ 
    using FY by blast

have  $\forall_F\ n$  in sequentially.  $dist\ (\mathcal{G}\ (j\ n)\ y)\ (h\ y) < inverse\ C$ 
  using uniform_limitD [OF ulim [of  $\{y\}$ ], of inverse C]  $\langle C > 0 \rangle\ y$ 
  by (simp add: dist_commute)
then obtain  $n$  where  $dist\ (\mathcal{G}\ (j\ n)\ y)\ (h\ y) < inverse\ C$ 

```

```

      by (meson eventually_at_top_linorder order_refl)
    moreover
    have  $h\ y = h\ v$ 
      by (metis  $\langle h\ v = c \rangle$  dist_commute that y)
    ultimately have  $cmod\ (inverse\ (\mathcal{F}\ (j\ n)\ y)) < inverse\ C$ 
      by (simp add:  $\langle h\ v = 0 \rangle$   $\mathcal{G\_def}$ )
    then have  $C < norm\ (\mathcal{F}\ (j\ n)\ y)$ 
      by (metis  $\mathcal{G\_def}$   $\mathcal{G}not0\ \langle y \in S \rangle$  inverse_less_imp_less inverse_zero
norm_inverse zero_less_norm_iff)
    show False
      using  $\langle C < cmod\ (\mathcal{F}\ (j\ n)\ y) \rangle$  le_C not_less by blast
    qed
    then show  $\exists x \in ball\ v\ e.\ h\ x \neq c$  by force
  qed
  show  $h$  holomorphic_on ball v e
    by (simp add: holh)
    show  $\bigwedge K.\ \llbracket compact\ K; K \subseteq ball\ v\ e \rrbracket \implies uniform\_limit\ K\ (\mathcal{G} \circ j)\ h$ 
sequentially
    by (simp add: ulim)
  qed (use  $\langle e > 0 \rangle$  in auto)
  with  $\langle h\ v = 0 \rangle$  show False by blast
qed
then obtain r where  $0 < r$  and  $r: \bigwedge h.\ h \in Y \implies cmod\ (h\ v) \leq r$ 
  by (metis not_le)
moreover
  obtain B Z where  $0 < B$  open Z  $v \in Z$   $Z \subseteq S$   $\bigwedge h\ z.\ \llbracket h \in Y; z \in Z \rrbracket \implies$ 
norm(h z)  $\leq B$ 
  using X01
  by (auto simp: r intro: GPicard1[OF  $\langle open\ S \rangle$   $\langle connected\ S \rangle$   $\langle v \in S \rangle$   $\langle r > 0 \rangle$ 
 $\langle Y \subseteq X \rangle$  holX] X01)
  ultimately show  $v \in U$ 
    using v by (simp add: U_def) meson
  qed
  have  $\bigwedge x.\ x \in K \longrightarrow x \in U$ 
    using  $\langle U = S \rangle$   $\langle K \subseteq S \rangle$  by blast
  then have  $\bigwedge x.\ x \in K \longrightarrow (\exists B\ Z.\ 0 < B \wedge open\ Z \wedge x \in Z \wedge$ 
 $(\forall h\ z'.\ h \in Y \wedge z' \in Z \longrightarrow norm(h\ z') \leq B))$ 
    unfolding U_def by blast
  then obtain F Z where  $F: \bigwedge x.\ x \in K \implies open\ (Z\ x) \wedge x \in Z\ x \wedge$ 
 $(\forall h\ z'.\ h \in Y \wedge z' \in Z\ x \longrightarrow norm(h\ z') \leq F\ x)$ 
    by metis
  then obtain L where  $L \subseteq K$  finite L and  $L: K \subseteq (\bigcup c \in L.\ Z\ c)$ 
    by (auto intro: compactE_image [OF  $\langle compact\ K \rangle$ , of K Z])
  then have *:  $\bigwedge x\ h\ z'.\ \llbracket x \in L; h \in Y \wedge z' \in Z\ x \rrbracket \implies cmod\ (h\ z') \leq F\ x$ 
    using F by blast
  have  $\exists B.\ \forall h\ z.\ h \in Y \wedge z \in K \longrightarrow norm(h\ z) \leq B$ 
  proof (cases L = {})
    case True with L show ?thesis by simp
  next

```

```

case False
then have  $\forall h z. h \in Y \wedge z \in K \longrightarrow (\exists x \in L. \text{cmod } (h z) \leq F x)$ 
  by (metis * L UN_E subset_iff)
with False  $\langle \text{finite } L \rangle$  show ?thesis
  by (rule_tac x = Max (F ` L) in exI) (simp add: linorder_class.Max_ge_iff)
qed
with that show ?thesis by metis
qed

```

lemma *GPicard4*:

```

assumes  $0 < k$  and holf:  $f \text{ holomorphic\_on } (\text{ball } 0 k - \{0\})$ 
  and AE:  $\bigwedge e. [0 < e; e < k] \implies \exists d. 0 < d \wedge d < e \wedge (\forall z \in \text{sphere } 0 d. \text{norm}(f z) \leq B)$ 
obtains  $\varepsilon$  where  $0 < \varepsilon \wedge \varepsilon < k \wedge \bigwedge z. z \in \text{ball } 0 \varepsilon - \{0\} \implies \text{norm}(f z) \leq B$ 
proof -
  obtain  $\varepsilon$  where  $0 < \varepsilon \wedge \varepsilon < k/2$  and  $\varepsilon$ :  $\bigwedge z. \text{norm } z = \varepsilon \implies \text{norm}(f z) \leq B$ 
    using AE [of k/2] 0 < k by auto
  show ?thesis
  proof
    show  $\varepsilon < k$ 
      using  $\langle 0 < k \rangle \langle \varepsilon < k/2 \rangle$  by auto
    show  $\text{cmod } (f \xi) \leq B$  if  $\xi: \xi \in \text{ball } 0 \varepsilon - \{0\}$  for  $\xi$ 
      proof -
        obtain  $d$  where  $0 < d \wedge d < \text{norm } \xi$  and  $d$ :  $\bigwedge z. \text{norm } z = d \implies \text{norm}(f z) \leq B$ 
          using AE [of norm ξ] 0 < k ξ by auto
        have [simp]:  $\text{closure } (\text{cball } 0 \varepsilon - \text{ball } 0 d) = \text{cball } 0 \varepsilon - \text{ball } 0 d$ 
          by (blast intro!: closure_closed)
        have [simp]:  $\text{interior } (\text{cball } 0 \varepsilon - \text{ball } 0 d) = \text{ball } 0 \varepsilon - \text{cball } (0::\text{complex}) d$ 
          using  $\langle 0 < \varepsilon \rangle \langle 0 < d \rangle$  by (simp add: interior_diff)
        have *:  $\text{norm}(f w) \leq B$  if  $w \in \text{cball } 0 \varepsilon - \text{ball } 0 d$  for  $w$ 
          proof (rule maximum_modulus_frontier [of f cball 0 ε - ball 0 d])
            show  $f \text{ holomorphic\_on interior } (\text{cball } 0 \varepsilon - \text{ball } 0 d)$ 
              using  $\langle \varepsilon < k \rangle \langle 0 < d \rangle$  that by (auto intro: holomorphic_on_subset [OF holf])
            show  $\text{continuous\_on } (\text{closure } (\text{cball } 0 \varepsilon - \text{ball } 0 d)) f$ 
              proof (intro holomorphic_on_imp_continuous_on holomorphic_on_subset [OF holf])
                show  $\text{closure } (\text{cball } 0 \varepsilon - \text{ball } 0 d) \subseteq \text{ball } 0 k - \{0\}$ 
                  using  $\langle 0 < d \rangle \langle \varepsilon < k \rangle$  by auto
                qed
            show  $\bigwedge z. z \in \text{frontier } (\text{cball } 0 \varepsilon - \text{ball } 0 d) \implies \text{cmod } (f z) \leq B$ 
              unfolding frontier_def
              using  $\varepsilon d \text{ less\_eq\_real\_def}$  by force
            qed (use that in auto)
          show ?thesis
            using *  $\langle d < \text{cmod } \xi \rangle$  that by auto
          qed
        qed
      qed

```

qed (use <0 < ε> in auto)
qed

lemma GPicard5:

assumes *hol*: f holomorphic_on (ball 0 1 - {0})
and *f01*: $\bigwedge z. z \in \text{ball } 0 \ 1 - \{0\} \implies f z \neq 0 \wedge f z \neq 1$
obtains $e B$ where $0 < e \ e < 1 \ 0 < B$
 $(\forall z \in \text{ball } 0 \ e - \{0\}. \text{norm}(f z) \leq B) \vee$
 $(\forall z \in \text{ball } 0 \ e - \{0\}. \text{norm}(f z) \geq B)$

proof -

have [*simp*]: $1 + \text{of_nat } n \neq (0::\text{complex})$ for n
using *of_nat_eq_0_iff* by fastforce
have [*simp*]: $\text{cmod } (1 + \text{of_nat } n) = 1 + \text{of_nat } n$ for n
by (*metis norm_of_nat of_nat_Suc*)
have *: $(\lambda x::\text{complex}. x / \text{of_nat } (\text{Suc } n)) \text{ ' } (\text{ball } 0 \ 1 - \{0\}) \subseteq \text{ball } 0 \ 1 - \{0\}$
for n

by (*auto simp: norm_divide field_split_simps split: if_split_asm*)

define *h* where $h \equiv \lambda n z::\text{complex}. f (z / (\text{Suc } n))$

have *holh*: $(h \ n)$ holomorphic_on ball 0 1 - {0} for n

unfolding *h_def*

proof (*rule holomorphic_on_compose_gen [unfolded o_def, OF _ holf *]*)

show $(\lambda x. x / \text{of_nat } (\text{Suc } n))$ holomorphic_on ball 0 1 - {0}

by (*intro holomorphic_intros*) auto

qed

have *h01*: $\bigwedge n z. z \in \text{ball } 0 \ 1 - \{0\} \implies h \ n \ z \neq 0 \wedge h \ n \ z \neq 1$

unfolding *h_def*

using * by (*force intro!*: *f01*)

obtain w where $w: w \in \text{ball } 0 \ 1 - \{0::\text{complex}\}$

by (*rule_tac w = 1/2 in that*) auto

consider *infinite* { $n. \text{norm}(h \ n \ w) \leq 1$ } | *infinite* { $n. 1 \leq \text{norm}(h \ n \ w)$ }

by (*metis (mono_tags, lifting) infinite_nat_iff_unbounded_le le_cases mem_Collect_eq*)

then show ?thesis

proof cases

case 1

with *infinite_enumerate* obtain $r :: \text{nat} \Rightarrow \text{nat}$

where *strict_mono* r and $r: \bigwedge n. r \ n \in \{n. \text{norm}(h \ n \ w) \leq 1\}$

by blast

obtain B where $B: \bigwedge j z. \llbracket \text{norm } z = 1/2; j \in \text{range } (h \circ r) \rrbracket \implies \text{norm}(j z) \leq B$

proof (*rule GPicard3 [OF _ _ w, where K = sphere 0 (1/2)]*)

show $\text{range } (h \circ r) \subseteq$

$\{g. g \text{ holomorphic_on ball } 0 \ 1 - \{0\} \wedge (\forall z \in \text{ball } 0 \ 1 - \{0\}. g z \neq 0 \wedge g z \neq 1)\}$

using *h01* by (*auto intro: holomorphic_intros holomorphic_on_compose holh*)

show *connected* (ball 0 1 - {0::complex})

by (*simp add: connected_open_delete*)

qed (use r in auto)

```

    have normf_le_B: cmod(f z) ≤ B if norm z = 1 / (2 * (1 + of_nat (r n)))
  for z n
  proof -
    have *:  $\bigwedge w. \text{norm } w = 1/2 \implies \text{cmod}((f (w / (1 + \text{of\_nat } (r n)))))) \leq B$ 
      using B by (auto simp: h_def o_def)
    have half:  $\text{norm } (z * (1 + \text{of\_nat } (r n))) = 1/2$ 
      by (simp add: norm_mult divide_simps that)
    show ?thesis
      using * [OF half] by simp
  qed
  obtain  $\varepsilon$  where  $0 < \varepsilon \varepsilon < 1 \bigwedge z. z \in \text{ball } 0 \varepsilon - \{0\} \implies \text{cmod}(f z) \leq B$ 
  proof (rule GPicard4 [OF zero_less_one half, of B])
    fix e::real
    assume  $0 < e \varepsilon < 1$ 
    obtain n where  $(1/e - 2) / 2 < \text{real } n$ 
      using reals_Archimedean2 by blast
    also have  $\dots \leq r n$ 
      using <strict_mono r> by (simp add: seq_suble)
    finally have  $(1/e - 2) / 2 < \text{real } (r n)$  .
    with <0 < e> have e:  $e > 1 / (2 + 2 * \text{real } (r n))$ 
      by (simp add: field_simps)
    show  $\exists d > 0. d < e \wedge (\forall z \in \text{sphere } 0 d. \text{cmod } (f z) \leq B)$ 
      apply (rule_tac x=1 / (2 * (1 + of_nat (r n))) in exI)
      using normf_le_B by (simp add: e)
  qed blast
  then have  $\varepsilon: \text{cmod } (f z) \leq |B| + 1$  if  $\text{cmod } z < \varepsilon z \neq 0$  for z
    using that by fastforce
  have  $0 < |B| + 1$ 
    by simp
  then show ?thesis
    using  $\varepsilon$  by (force intro!: that [OF <0 <  $\varepsilon$ > < $\varepsilon < 1$ >])
next
case 2
with infinite_enumerate obtain r :: nat  $\Rightarrow$  nat
  where strict_mono r and r:  $\bigwedge n. r n \in \{n. \text{norm}(h n w) \geq 1\}$ 
  by blast
obtain B where B:  $\bigwedge j z. [\text{norm } z = 1/2; j \in \text{range } (\lambda n. \text{inverse } \circ h (r n))]$ 
 $\implies \text{norm}(j z) \leq B$ 
proof (rule GPicard3 [OF __ w, where K = sphere 0 (1/2)])
  show  $\text{range } (\lambda n. \text{inverse } \circ h (r n)) \subseteq$ 
    {g. g holomorphic_on ball 0 1 - {0}  $\wedge$  ( $\forall z \in \text{ball } 0 1 - \{0\}. g z \neq 0 \wedge$ 
    g z  $\neq 1$ )}
  using h01 by (auto intro!: holomorphic_intros holomorphic_on_compose_gen
  [unfolded o_def, OF __ holh] holomorphic_on_compose)
  show connected (ball 0 1 - {0::complex})
    by (simp add: connected_open_delete)
  show  $\bigwedge j. j \in \text{range } (\lambda n. \text{inverse } \circ h (r n)) \implies \text{cmod } (j w) \leq 1$ 
    using r norm_inverse_le_norm by fastforce
  qed (use r in auto)

```

```

    have norm_if_le_B: cmod(inverse (f z)) ≤ B if norm z = 1 / (2 * (1 +
of_nat (r n))) for z n
  proof -
    have *: inverse (cmod((f (z / (1 + of_nat (r n)))))) ≤ B if norm z = 1/2
for z
    using B [OF that] by (force simp: norm_inverse h_def)
    have half: norm (z * (1 + of_nat (r n))) = 1/2
    by (simp add: norm_mult divide_simps that)
    show ?thesis
    using * [OF half] by (simp add: norm_inverse)
  qed
  have hol_if: (inverse ∘ f) holomorphic_on (ball 0 1 - {0})
    by (metis (no_types, lifting) holf_comp_apply f01 holomorphic_on_inverse
holomorphic_transform)
  obtain ε where 0 < ε ε < 1 and leB: ∧z. z ∈ ball 0 ε - {0} ⇒ cmod((inverse
∘ f) z) ≤ B
  proof (rule GPicard4 [OF zero_less_one hol_if, of B])
    fix e::real
    assume 0 < e e < 1
    obtain n where (1/e - 2) / 2 < real n
    using reals_Archimedean2 by blast
    also have ... ≤ r n
    using ‹strict_mono r› by (simp add: seq_suble)
    finally have (1/e - 2) / 2 < real (r n) .
    with ‹0 < e› have e: e > 1 / (2 + 2 * real (r n))
    by (simp add: field_simps)
    show ∃ d>0. d < e ∧ (∀ z∈sphere 0 d. cmod ((inverse ∘ f) z) ≤ B)
    apply (rule_tac x=1 / (2 * (1 + of_nat (r n))) in exI)
    using norm_if_le_B by (simp add: e)
  qed blast
  have ε: cmod (f z) ≥ inverse B and B > 0 if cmod z < ε z ≠ 0 for z
  proof -
    have inverse (cmod (f z)) ≤ B
    using leB that by (simp add: norm_inverse)
    moreover
    have f z ≠ 0
    using ‹ε < 1› f01 that by auto
    ultimately show cmod (f z) ≥ inverse B
    by (simp add: norm_inverse inverse_le_imp_le)
    show B > 0
    using ‹f z ≠ 0› ‹inverse (cmod (f z)) ≤ B› not_le order.trans by fastforce
  qed
  then have B > 0
  by (metis ‹0 < ε› dense leI order.asym vector_choose_size)
  then have inverse B > 0
  by (simp add: field_split_simps)
  then show ?thesis
  using ε that [OF ‹0 < ε› ‹ε < 1›]
  by (metis Diff_iff dist_0_norm insert_iff mem_ball)

```

qed
qed

lemma GPicard6:

assumes open M $z \in M$ $a \neq 0$ and holf: f holomorphic_on $(M - \{z\})$
 and f0a: $\bigwedge w. w \in M - \{z\} \implies f w \neq 0 \wedge f w \neq a$
 obtains r where $0 < r$ ball z $r \subseteq M$
 bounded($f^{-1}(\text{ball } z \text{ } r - \{z\})$) \vee
 bounded($(\text{inverse} \circ f)^{-1}(\text{ball } z \text{ } r - \{z\})$)
 proof –
 obtain r where $0 < r$ and r : ball z $r \subseteq M$
 using assms openE by blast
 let ?g = $\lambda w. f(z + \text{of_real } r * w) / a$
 obtain e B where $0 < e$ $e < 1$ $0 < B$
 and B : $(\forall z \in \text{ball } 0 \text{ } e - \{0\}. \text{norm}(\text{?g } z) \leq B) \vee (\forall z \in \text{ball } 0 \text{ } e - \{0\}. \text{norm}(\text{?g } z) \geq B)$
 proof (rule GPicard5)
 show ?g holomorphic_on ball 0 $1 - \{0\}$
 proof (intro holomorphic_intros holomorphic_on_compose_gen [unfolded o_def, OF_holf])
 show $(\lambda x. z + \text{complex_of_real } r * x)^{-1}(\text{ball } 0 \text{ } 1 - \{0\}) \subseteq M - \{z\}$
 using $\langle 0 < r \rangle$
 by (auto simp: dist_norm norm_mult subset_eq)
 qed (use $\langle a \neq 0 \rangle$ in auto)
 show $\bigwedge w. w \in \text{ball } 0 \text{ } 1 - \{0\} \implies f(z + \text{of_real } r * w) / a \neq 0 \wedge f(z + \text{of_real } r * w) / a \neq 1$
 using f0a $\langle 0 < r \rangle$ $\langle a \neq 0 \rangle$ r
 by (auto simp: field_split_simps dist_norm norm_mult subset_eq)
 qed
 show ?thesis
 proof
 show $0 < e * r$
 by (simp add: $\langle 0 < e \rangle$ $\langle 0 < r \rangle$)
 have ball z $(e * r) \subseteq \text{ball } z$ r
 by (simp add: $\langle 0 < r \rangle$ $\langle e < 1 \rangle$ order.strict_implies_order subset_ball)
 then show ball z $(e * r) \subseteq M$
 using r by blast
 consider $\bigwedge z. z \in \text{ball } 0 \text{ } e - \{0\} \implies \text{norm}(\text{?g } z) \leq B \mid \bigwedge z. z \in \text{ball } 0 \text{ } e - \{0\} \implies \text{norm}(\text{?g } z) \geq B$
 using B by blast
 then show bounded $(f^{-1}(\text{ball } z \text{ } (e * r) - \{z\})) \vee$
 bounded $((\text{inverse} \circ f)^{-1}(\text{ball } z \text{ } (e * r) - \{z\}))$
 proof cases
 case 1
 have $\llbracket \text{dist } z \text{ } w < e * r; w \neq z \rrbracket \implies \text{cmod}(f w) \leq B * \text{norm } a$ for w
 using $\langle a \neq 0 \rangle$ $\langle 0 < r \rangle$ 1 [of $(w - z) / r$]
 by (simp add: norm_divide dist_norm field_split_simps)
 then show ?thesis


```

    by (force simp: intro!: boundedI)
  next
  case 2
  have [[dist z w < e * r; w ≠ z]] ⇒ cmod (f w) ≥ B * norm a for w
    using ‹a ≠ 0› ‹0 < r› 2 [of (w - z) / r]
    by (simp add: norm_divide dist_norm field_split_simps)
  then have [[dist z w < e * r; w ≠ z]] ⇒ inverse (cmod (f w)) ≤ inverse (B
* norm a) for w
    by (metis ‹0 < B› ‹a ≠ 0› mult_pos_pos norm_inverse norm_inverse_le_norm
zero_less_norm_iff)
    then show ?thesis
      by (force simp: norm_inverse intro!: boundedI)
  qed
qed
qed

```

theorem *great_Picard*:

```

  assumes open M z ∈ M a ≠ b and holf: f holomorphic_on (M - {z})
    and fab: ∧w. w ∈ M - {z} ⇒ f w ≠ a ∧ f w ≠ b
  obtains l where (f → l) (at z) ∨ ((inverse ∘ f) → l) (at z)
proof -
  obtain r where 0 < r and zrM: ball z r ⊆ M
    and r: bounded((λz. f z - a) ‘ (ball z r - {z})) ∨
      bounded((inverse ∘ (λz. f z - a)) ‘ (ball z r - {z}))
  proof (rule GPicard6 [OF ‹open M› ‹z ∈ M›])
    show b - a ≠ 0
      using assms by auto
    show (λz. f z - a) holomorphic_on M - {z}
      by (intro holomorphic_intros holf)
  qed (use fab in auto)
  have holfb: f holomorphic_on ball z r - {z}
    using zrM by (auto intro: holomorphic_on_subset [OF holf])
  have holfb_i: (λz. inverse(f z - a)) holomorphic_on ball z r - {z}
    using fab zrM by (fastforce intro!: holomorphic_intros holfb)
  show ?thesis
    using r
  proof
    assume bounded ((λz. f z - a) ‘ (ball z r - {z}))
    then obtain B where B: ∧w. w ∈ (λz. f z - a) ‘ (ball z r - {z}) ⇒ norm
w ≤ B
      by (force simp: bounded_iff)
    then have ∀x. x ≠ z ∧ dist x z < r → cmod (f x - a) ≤ B
      by (simp add: dist_commute)
    with ‹0 < r› have ∀F w in at z. cmod (f w - a) ≤ B
      by (force simp add: eventually_at)
    moreover have ∧x. cmod (f x - a) ≤ B ⇒ cmod (f x) ≤ B + cmod a
      by (metis add.commute add_le_cancel_right norm_triangle_sub order.trans)
    ultimately have ∃B. ∀F w in at z. cmod (f w) ≤ B

```

```

    by (metis (mono_tags, lifting) eventually_at)
    then obtain g where holg: g holomorphic_on ball z r and gf:  $\bigwedge w. w \in \text{ball } z r - \{z\} \implies g w = f w$ 
    using  $\langle 0 < r \rangle$  holomorphic_on_extend_bounded [OF holfb] by auto
    then have g -z→ g z
    unfolding continuous_at [symmetric]
    using  $\langle 0 < r \rangle$  centre_in_ball field_differentiable_imp_continuous_at
    holomorphic_on_imp_differentiable_at by blast
    then have (f → g z) (at z)
    using Lim_transform_within_open [of g g z z]
    using  $\langle 0 < r \rangle$  centre_in_ball gf by blast
    then show ?thesis
    using that by blast
next
    assume bounded((inverse o (λz. f z - a)) ' (ball z r - {z}))
    then obtain B where B:  $\bigwedge w. w \in (\text{inverse o } (\lambda z. f z - a)) ' (\text{ball } z r - \{z\}) \implies \text{norm } w \leq B$ 
    by (force simp: bounded_iff)
    then have  $\forall x. x \neq z \wedge \text{dist } x z < r \longrightarrow \text{cmod } (\text{inverse } (f x - a)) \leq B$ 
    by (simp add: dist_commute)
    with  $\langle 0 < r \rangle$  have  $\forall_F w$  in at z.  $\text{cmod } (\text{inverse } (f w - a)) \leq B$ 
    by (auto simp add: eventually_at)
    then have  $\exists B. \forall_F z$  in at z.  $\text{cmod } (\text{inverse } (f z - a)) \leq B$ 
    by blast
    then obtain g where holg: g holomorphic_on ball z r and gf:  $\bigwedge w. w \in \text{ball } z r - \{z\} \implies g w = \text{inverse } (f w - a)$ 
    using  $\langle 0 < r \rangle$  holomorphic_on_extend_bounded [OF holfb_i] by auto
    then have gz: g -z→ g z
    unfolding continuous_at [symmetric]
    using  $\langle 0 < r \rangle$  centre_in_ball field_differentiable_imp_continuous_at
    holomorphic_on_imp_differentiable_at by blast
    have gnz:  $\bigwedge w. w \in \text{ball } z r - \{z\} \implies g w \neq 0$ 
    using gf fab zrM by fastforce
    show ?thesis
    proof (cases g z = 0)
    case True
    have *:  $\llbracket g \neq 0; \text{inverse } g = f - a \rrbracket \implies g / (1 + a * g) = \text{inverse } f$  for f
    g::complex
    by (auto simp: field_simps)
    have (inverse o f) -z→ 0
    proof (rule Lim_transform_within_open [of λw. g w / (1 + a * g w) _ _
    UNIV ball z r])
    show (λw. g w / (1 + a * g w)) -z→ 0
    using True by (auto simp: intro!: tendsto_eq_intros gz)
    show  $\bigwedge x. \llbracket x \in \text{ball } z r; x \neq z \rrbracket \implies g x / (1 + a * g x) = (\text{inverse o } f) x$ 
    using * gf gnz by simp
    qed (use  $\langle 0 < r \rangle$  in auto)
    with that show ?thesis by blast
next

```

```

case False
show ?thesis
proof (cases  $1 + a * g z = 0$ )
  case True
  have  $(f \longrightarrow 0)$  (at z)
  proof (rule Lim_transform_within_open [of  $\lambda w. (1 + a * g w) / g w$  ___
    _ ball z r])
    show  $(\lambda w. (1 + a * g w) / g w) -z \rightarrow 0$ 
      by (rule tendsto_eq_intros refl gz  $\langle g z \neq 0 \rangle$  | simp add: True)+
    show  $\bigwedge x. \llbracket x \in \text{ball } z \text{ } r; x \neq z \rrbracket \implies (1 + a * g x) / g x = f x$ 
      using fab fab zrM by (fastforce simp add: gf field_split_simps)
  qed (use  $\langle 0 < r \rangle$  in auto)
  then show ?thesis
    using that by blast
next
case False
have *:  $\llbracket g \neq 0; \text{inverse } g = f - a \rrbracket \implies g / (1 + a * g) = \text{inverse } f$  for f
g::complex
  by (auto simp: field_simps)
have  $(\text{inverse } \circ f) -z \rightarrow g z / (1 + a * g z)$ 
proof (rule Lim_transform_within_open [of  $\lambda w. g w / (1 + a * g w)$  ___
UNIV ball z r])
  show  $(\lambda w. g w / (1 + a * g w)) -z \rightarrow g z / (1 + a * g z)$ 
    using False by (auto simp: False intro!: tendsto_eq_intros gz)
  show  $\bigwedge x. \llbracket x \in \text{ball } z \text{ } r; x \neq z \rrbracket \implies g x / (1 + a * g x) = (\text{inverse } \circ f) x$ 
    using * gf gnz by simp
  qed (use  $\langle 0 < r \rangle$  in auto)
with that show ?thesis by blast
qed
qed
qed
qed

```

corollary great_Picard_alt:

```

assumes M: open M z  $\in$  M and holf: f holomorphic_on (M - {z})
and non:  $\bigwedge l. \neg (f \longrightarrow l)$  (at z)  $\bigwedge l. \neg ((\text{inverse } \circ f) \longrightarrow l)$  (at z)
obtains a where  $-\{a\} \subseteq f^{-1}(M - \{z\})$ 
unfolding subset_iff image_iff
by (metis great_Picard [OF M _ holf] non Compl_iff insertI1)

```

corollary great_Picard_infinite:

```

assumes M: open M z  $\in$  M and holf: f holomorphic_on (M - {z})
and non:  $\bigwedge l. \neg (f \longrightarrow l)$  (at z)  $\bigwedge l. \neg ((\text{inverse } \circ f) \longrightarrow l)$  (at z)
obtains a where  $\bigwedge w. w \neq a \implies \text{infinite } \{x. x \in M - \{z\} \wedge f x = w\}$ 
proof -
  have False if  $a \neq b$  and ab: finite  $\{x. x \in M - \{z\} \wedge f x = a\}$  finite  $\{x. x \in M - \{z\} \wedge f x = b\}$  for a b

```

```

proof -
  have finab: finite {x. x ∈ M - {z} ∧ f x ∈ {a, b}}
  using finite_UnI [OF ab] unfolding mem_Collect_eq insert_iff empty_iff
  by (simp add: conj_disj_distribL)
  obtain r where  $0 < r$  and zrM:  $\text{ball } z \ r \subseteq M$  and r:  $\bigwedge x. [x \in M - \{z\}; f \ x \in \{a, b\}] \implies x \notin \text{ball } z \ r$ 
  proof -
    obtain e where  $e > 0$  and e:  $\text{ball } z \ e \subseteq M$ 
    using assms openE by blast
    show ?thesis
    proof (cases {x ∈ M - {z}. f x ∈ {a, b}} = {})
      case True
      then show ?thesis
        using  $e < e > 0$  that by fastforce
    next
    case False
    let ?r = min e (Min (dist z ‘ {x ∈ M - {z}. f x ∈ {a, b}}))
    show ?thesis
    proof
      show  $0 < ?r$ 
      using min_less_iff_conj Min_gr_iff finab False  $0 < e$  by auto
      have  $\text{ball } z \ ?r \subseteq \text{ball } z \ e$ 
      by (simp add: subset_ball)
      with e show  $\text{ball } z \ ?r \subseteq M$  by blast
      show  $\bigwedge x. [x \in M - \{z\}; f \ x \in \{a, b\}] \implies x \notin \text{ball } z \ ?r$ 
      using min_less_iff_conj Min_gr_iff finab False  $0 < e$  by auto
    qed
  qed
  qed
  have holfb: f holomorphic_on ( $\text{ball } z \ r - \{z\}$ )
  apply (rule holomorphic_on_subset [OF holfb])
  using zrM by auto
  show ?thesis
  apply (rule great_Picard [OF open_ball _  $a \neq b$ ] holfb)
  using  $non \ 0 < r$  r zrM by auto
  qed
with that show thesis
  by meson
qed

theorem Casorati_Weierstrass:
  assumes open M z ∈ M f holomorphic_on (M - {z})
  and  $\bigwedge l. \neg (f \longrightarrow l) \text{ (at } z) \wedge \bigwedge l. \neg ((\text{inverse} \circ f) \longrightarrow l) \text{ (at } z)$ 
  shows  $\text{closure}(f \text{ ‘ } (M - \{z\})) = \text{UNIV}$ 
proof -
  obtain a where  $a: - \{a\} \subseteq f \text{ ‘ } (M - \{z\})$ 
  using great_Picard_alt [OF assms].
  have  $\text{UNIV} = \text{closure}(- \{a\})$ 
  by (simp add: closure_interior)

```

```

also have ...  $\subseteq$  closure( $f^{-1}(M - \{z\})$ )
  by (simp add: a_closure_mono)
finally show ?thesis
  by blast
qed

```

```

end

```

8 Moebius functions, Equivalents of Simply Connected Sets, Riemann Mapping Theorem

```

theory Riemann_Mapping
imports Great_Picard
begin

```

8.1 Moebius functions are biholomorphisms of the unit disc

```

definition Moebius_function :: [real, complex, complex]  $\Rightarrow$  complex where
  Moebius_function  $\equiv \lambda t w z. \exp(i * \text{of\_real } t) * (z - w) / (1 - \text{cnj } w * z)$ 

```

```

lemma Moebius_function_simple:
  Moebius_function 0 w z = (z - w) / (1 - cnj w * z)
by (simp add: Moebius_function_def)

```

```

lemma Moebius_function_eq_zero:
  Moebius_function t w w = 0
by (simp add: Moebius_function_def)

```

```

lemma Moebius_function_of_zero:
  Moebius_function t w 0 = - exp(i * of_real t) * w
by (simp add: Moebius_function_def)

```

```

lemma Moebius_function_norm_lt_1:
  assumes w1: norm w < 1 and z1: norm z < 1
  shows norm (Moebius_function t w z) < 1

```

```

proof -

```

```

  have 1 - cnj w * z  $\neq$  0
    by (metis complex_cnj_cnj_complex_mod_sqrt_Re_mult_cnj_mult.commute
  mult_less_cancel_right1 norm_ge_zero norm_mult norm_one order.asym right_minus_eq
  w1 z1)

```

```

  then have VV: 1 - w * cnj z  $\neq$  0

```

```

    by (metis complex_cnj_cnj_complex_cnj_mult_complex_cnj_one_right_minus_eq)

```

```

  then have 1 - norm (Moebius_function t w z)  $^2 =$ 

```

$$((1 - \text{norm } w^2) / (\text{norm } (1 - \text{cnj } w * z)^2)) * (1 - \text{norm } z^2)$$

```

  apply (cases w)

```

```

  apply (cases z)

```

```

  apply (simp add: Moebius_function_def divide_simps norm_divide norm_mult)

```

```

  apply (simp add: complex_norm_complex_diff_complex_mult_one_complex.code)

```

```

complex_cnj)
  apply (auto simp: algebra_simps power2_eq_square)
  done
  then have  $1 - (\text{cmod } (\text{Moebius\_function } t \ w \ z))^2 = (1 - \text{cmod } (w * w)) / (\text{cmod } (1 - \text{cnj } w * z))^2 * (1 - \text{cmod } (z * z))$ 
  by (simp add: norm_mult power2_eq_square)
  moreover have  $0 < 1 - \text{cmod } (z * z)$ 
  by (metis (no_types) z1 diff_gt_0_iff_gt mult.left_neutral norm_mult_less)
  ultimately have  $0 < 1 - \text{norm } (\text{Moebius\_function } t \ w \ z)^2$ 
  using  $\langle 1 - \text{cnj } w * z \neq 0 \rangle$  w1 norm_mult_less by fastforce
  then show ?thesis
  using linorder_not_less by fastforce
qed

```

```

lemma Moebius_function_holomorphic:
  assumes  $\text{norm } w < 1$ 
  shows  $\text{Moebius\_function } t \ w \ \text{holomorphic\_on\_ball } 0 \ 1$ 
proof -
  have *:  $1 - z * w \neq 0$  if  $\text{norm } z < 1$  for  $z$ 
  proof -
    have  $\text{norm } (1::\text{complex}) \neq \text{norm } (z * w)$ 
    using assms that norm_mult_less by fastforce
    then show ?thesis by auto
  qed
  show ?thesis
  apply (simp add: Moebius_function_def)
  apply (intro holomorphic_intros)
  using assms *
  by (metis complex_cnj_cnj complex_cnj_mult complex_cnj_one complex_mod_cnj
  mem_ball_0 mult.commute right_minus_eq)
qed

```

```

lemma Moebius_function_compose:
  assumes  $\text{meq}: -w1 = w2$  and  $\text{norm } w1 < 1 \ \text{norm } z < 1$ 
  shows  $\text{Moebius\_function } 0 \ w1 \ (\text{Moebius\_function } 0 \ w2 \ z) = z$ 
proof -
  have  $\text{norm } w2 < 1$ 
  using assms by auto
  then have  $-w1 = z$  if  $\text{cnj } w2 * z = 1$ 
  by (metis assms(3) complex_mod_cnj less_irrefl mult.right_neutral norm_mult_less
  norm_one that)
  moreover have  $z=0$  if  $1 - \text{cnj } w2 * z = \text{cnj } w1 * (z - w2)$ 
  proof -
    have  $w2 * \text{cnj } w2 = 1$ 
    using that meq by (auto simp: algebra_simps)
    then show  $z = 0$ 
    by (metis (no_types)  $\langle \text{cmod } w2 < 1 \rangle$  complex_mod_cnj less_irrefl mult.right_neutral
  norm_mult_less norm_one)
  qed

```

```

moreover have  $z - w2 - w1 * (1 - cnj w2 * z) = z * (1 - cnj w2 * z - cnj$ 
 $w1 * (z - w2))$ 
using meq by (fastforce simp: algebra_simps)
ultimately
show ?thesis
by (simp add: Moebius_function_def divide_simps norm_divide norm_mult)
qed

```

lemma *ball_biholomorphism_exists*:

assumes $a \in \text{ball } 0 \ 1$

obtains $f \ g$ **where** $f \ a = 0$

f *holomorphic_on* $\text{ball } 0 \ 1 \ f \ ' \ \text{ball } 0 \ 1 \subseteq \text{ball } 0 \ 1$

g *holomorphic_on* $\text{ball } 0 \ 1 \ g \ ' \ \text{ball } 0 \ 1 \subseteq \text{ball } 0 \ 1$

$\bigwedge z. z \in \text{ball } 0 \ 1 \implies f (g \ z) = z$

$\bigwedge z. z \in \text{ball } 0 \ 1 \implies g (f \ z) = z$

proof

show *Moebius_function 0 a holomorphic_on ball 0 1 Moebius_function 0 (-a)*
holomorphic_on ball 0 1

using *Moebius_function_holomorphic assms mem_ball_0* **by** *auto*

show *Moebius_function 0 a a = 0*

by (*simp add: Moebius_function_eq_zero*)

show *Moebius_function 0 a ' ball 0 1 ⊆ ball 0 1*

Moebius_function 0 (- a) ' ball 0 1 ⊆ ball 0 1

using *Moebius_function_norm_lt_1 assms* **by** *auto*

show *Moebius_function 0 a (Moebius_function 0 (- a) z) = z*

Moebius_function 0 (- a) (Moebius_function 0 a z) = z **if** $z \in \text{ball } 0 \ 1$ **for**

z

using *Moebius_function_compose assms that* **by** *auto*

qed

8.2 A big chain of equivalents of simple connectedness for an open set

lemma *biholomorphic_to_disc_aux*:

assumes *open S connected S 0 ∈ S and S01: S ⊆ ball 0 1*

and *prev: $\bigwedge f. \llbracket f \text{ holomorphic_on } S; \bigwedge z. z \in S \implies f \ z \neq 0; \text{inj_on } f \ S \rrbracket$*

$\implies \exists g. g \text{ holomorphic_on } S \wedge (\forall z \in S. f \ z = (g \ z)^2)$

shows $\exists f \ g. f \text{ holomorphic_on } S \wedge g \text{ holomorphic_on } \text{ball } 0 \ 1 \wedge$

$(\forall z \in S. f \ z \in \text{ball } 0 \ 1 \wedge g(f \ z) = z) \wedge$

$(\forall z \in \text{ball } 0 \ 1. g \ z \in S \wedge f(g \ z) = z)$

proof –

define F **where** $F \equiv \{h. h \text{ holomorphic_on } S \wedge h \ ' \ S \subseteq \text{ball } 0 \ 1 \wedge h \ 0 = 0 \wedge$
inj_on h S\}

have *idF: id ∈ F*

using *S01* **by** (*auto simp: F_def*)

then have $F \neq \{\}$

by *blast*

have *imF_ne: $(\lambda h. \text{norm}(\text{deriv } h \ 0)) \ ' \ F \neq \{\}$*

using *idF* **by** *auto*

```

have holF:  $\bigwedge h. h \in F \implies h \text{ holomorphic\_on } S$ 
  by (auto simp: F_def)
obtain f where  $f \in F$  and normf:  $\bigwedge h. h \in F \implies \text{norm}(\text{deriv } h \ 0) \leq \text{norm}(\text{deriv } f \ 0)$ 
proof -
  obtain r where  $r > 0$  and r:  $\text{ball } 0 \ r \subseteq S$ 
  using  $\langle \text{open } S \rangle \langle 0 \in S \rangle \text{ openE}$  by auto
  have bdd:  $\text{bdd\_above } ((\lambda h. \text{norm}(\text{deriv } h \ 0))) \ ' F$ 
  proof (intro bdd_aboveI exI ballI, clarify)
    show  $\text{norm}(\text{deriv } f \ 0) \leq 1 / r$  if  $f \in F$  for f
    proof -
      have r01:  $(*) (\text{complex\_of\_real } r) \ ' \text{ball } 0 \ 1 \subseteq S$ 
      using that  $\langle r > 0 \rangle$  by (auto simp: norm_mult r [THEN subsetD])
      then have f holomorphic_on  $(*) (\text{complex\_of\_real } r) \ ' \text{ball } 0 \ 1$ 
      using holomorphic_on_subset [OF holF] by (simp add: that)
      then have holf:  $f \circ (\lambda z. (r * z)) \text{ holomorphic\_on } (\text{ball } 0 \ 1)$ 
      by (intro holomorphic_intros holomorphic_on_compose)
      have f0:  $(f \circ (*)) (\text{complex\_of\_real } r) \ 0 = 0$ 
      using F_def that by auto
      have f  $' S \subseteq \text{ball } 0 \ 1$ 
      using F_def that by blast
      with r01 have fr1:  $\bigwedge z. \text{norm } z < 1 \implies \text{norm}((f \circ (*))(\text{of\_real } r)z) < 1$ 
      by force
      have *:  $((\lambda w. f(r * w)) \text{ has\_field\_derivative } \text{deriv } f(r * z) * r) \text{ (at } z)$ 
      if  $z \in \text{ball } 0 \ 1$  for  $z::\text{complex}$ 
      proof (rule DERIV_chain' [where g=f])
        show  $(f \text{ has\_field\_derivative } \text{deriv } f(\text{of\_real } r * z)) \text{ (at } (\text{of\_real } r * z))$ 
        apply (rule holomorphic_derivI [OF holF  $\langle \text{open } S \rangle$ ])
        apply (rule  $\langle f \in F \rangle$ )
        by (meson imageI r01 subset_iff that)
      qed simp
      have df0:  $((\lambda w. f(r * w)) \text{ has\_field\_derivative } \text{deriv } f \ 0 * r) \text{ (at } 0)$ 
      using * [of 0] by simp
      have deq:  $\text{deriv } (\lambda x. f(\text{complex\_of\_real } r * x)) \ 0 = \text{deriv } f \ 0 * \text{complex\_of\_real } r$ 
      using DERIV_imp_deriv df0 by blast
      have norm  $(\text{deriv } (f \circ (*)) (\text{complex\_of\_real } r) \ 0) \leq 1$ 
      by (auto intro: Schwarz_Lemma [OF holf f0 fr1, of 0])
      with  $\langle r > 0 \rangle$  show ?thesis
      by (simp add: deq norm_mult divide_simps o_def)
    qed
  qed
  define l where  $l \equiv \text{SUP } h \in F. \text{norm}(\text{deriv } h \ 0)$ 
  have eql:  $\text{norm}(\text{deriv } f \ 0) = l$  if  $le: l \leq \text{norm}(\text{deriv } f \ 0)$  and  $f \in F$  for f
  apply (rule order_antisym [OF _ le])
  using  $\langle f \in F \rangle$  bdd_cSUP_upper by (fastforce simp: l_def)
  obtain  $\mathcal{F}$  where  $\mathcal{F}in: \bigwedge n. \mathcal{F} \ n \in F$  and  $\mathcal{F}lim: (\lambda n. \text{norm}(\text{deriv } (\mathcal{F} \ n) \ 0)) \longrightarrow l$ 
  proof -

```



```

have  $\exists f. f \in F \wedge |norm (deriv f 0) - l| < 1 / (Suc n)$  for  $n$ 
proof -
  obtain  $f$  where  $f \in F$  and  $f: l < norm (deriv f 0) + 1/(Suc n)$ 
  using  $cSup\_least [OF imF\_ne, of l - 1/(Suc n)]$  by ( $fastforce simp add:$ 
 $l\_def$ )
  then have  $|norm (deriv f 0) - l| < 1 / (Suc n)$ 
  by ( $fastforce simp add: abs\_if not\_less eql$ )
  with  $\langle f \in F \rangle$  show  $?thesis$ 
  by  $blast$ 
qed
then obtain  $\mathcal{F}$  where  $f\mathcal{F}: \bigwedge n. (\mathcal{F} n) \in F$ 
and  $fless: \bigwedge n. |norm (deriv (\mathcal{F} n) 0) - l| < 1 / (Suc n)$ 
by  $metis$ 
have  $(\lambda n. norm (deriv (\mathcal{F} n) 0)) \longrightarrow l$ 
proof ( $rule metric\_LIMSEQ\_I$ )
  fix  $e::real$ 
  assume  $e > 0$ 
  then obtain  $N::nat$  where  $N: e > 1/(Suc N)$ 
  using  $nat\_approx\_posE$  by  $blast$ 
  show  $\exists N. \forall n \geq N. dist (norm (deriv (\mathcal{F} n) 0)) l < e$ 
  proof ( $intro exI allI impI$ )
    fix  $n$  assume  $N \leq n$ 
    have  $dist (norm (deriv (\mathcal{F} n) 0)) l < 1 / (Suc n)$ 
    using  $fless$  by ( $simp add: dist\_norm$ )
    also have  $\dots < e$ 
    using  $N \langle N \leq n \rangle inverse\_of\_nat\_le le\_less\_trans$  by  $blast$ 
    finally show  $dist (norm (deriv (\mathcal{F} n) 0)) l < e$  .
  qed
qed
with  $f\mathcal{F}$  show  $?thesis$ 
using  $that$  by  $blast$ 
qed
have  $\bigwedge K. \llbracket compact K; K \subseteq S \rrbracket \implies \exists B. \forall h \in F. \forall z \in K. norm (h z) \leq B$ 
by ( $rule\_tac x=1$  in  $exI$ ) ( $force simp: F\_def$ )
moreover have  $range \mathcal{F} \subseteq F$ 
using  $\langle \bigwedge n. \mathcal{F} n \in F \rangle$  by  $blast$ 
ultimately obtain  $f$  and  $r :: nat \Rightarrow nat$ 
where  $holF: f$   $holomorphic\_on S$  and  $r: strict\_mono r$ 
and  $limf: \bigwedge x. x \in S \implies (\lambda n. \mathcal{F} (r n) x) \longrightarrow f x$ 
and  $ulimf: \bigwedge K. \llbracket compact K; K \subseteq S \rrbracket \implies uniform\_limit K (\mathcal{F} \circ r) f$ 
 $sequentially$ 
using  $Montel [of S F \mathcal{F}, OF \langle open S \rangle holF]$  by  $auto+$ 
have  $der: \bigwedge n x. x \in S \implies ((\mathcal{F} \circ r) n has\_field\_derivative ((\lambda n. deriv (\mathcal{F} n))$ 
 $\circ r) n x) (at x)$ 
using  $\langle \bigwedge n. \mathcal{F} n \in F \rangle \langle open S \rangle holF holomorphic\_derivI$  by  $fastforce$ 
have  $ulim: \bigwedge x. x \in S \implies \exists d > 0. cball x d \subseteq S \wedge uniform\_limit (cball x d)$ 
 $(\mathcal{F} \circ r) f$   $sequentially$ 
by ( $meson ulimf \langle open S \rangle compact\_cball open\_contains\_cball$ )
obtain  $f' :: complex \Rightarrow complex$  where  $f': (f has\_field\_derivative f' 0) (at 0)$ 

```

```

    and tof'0: ( $\lambda n. ((\lambda n. \text{deriv } (\mathcal{F} \ n)) \circ r) \ n \ 0) \longrightarrow f' \ 0$ 
    using has_complex_derivative_uniform_sequence [OF ‹open S› der ulim] ‹0
 $\in S$ › by metis
  then have derf0:  $\text{deriv } f \ 0 = f' \ 0$ 
    by (simp add: DERIV_imp_deriv)
  have f_field_differentiable (at 0)
    using field_differentiable_def f' by blast
  have ( $\lambda x. (\text{norm } (\text{deriv } (\mathcal{F} \ (r \ x)) \ 0)) \longrightarrow \text{norm } (\text{deriv } f \ 0)$ )
    using isCont_tendsto_compose [OF continuous_norm [OF continuous_ident]
tof'0] derf0 by auto
  with LIMSEQ_subseq_LIMSEQ [OF Flim r] have no_df0:  $\text{norm}(\text{deriv } f \ 0) = l$ 
    by (force simp: o_def intro: tendsto_unique)
  have nonconstf:  $\neg f \ \text{constant\_on } S$ 
  proof -
    have False if  $\bigwedge x. x \in S \implies f \ x = c$  for c
    proof -
      have  $\text{deriv } f \ 0 = 0$ 
      by (metis that ‹open S› ‹0  $\in S$ › DERIV_imp_deriv [OF has_field_derivative_transform_within_open
[OF DERIV_const]])
      with no_df0 have  $l = 0$ 
      by auto
      with eq [OF _ idF] show False by auto
    qed
  then show ?thesis
    by (meson constant_on_def)
  qed
  show ?thesis
  proof
    show  $f \in F$ 
      unfolding F_def
    proof (intro CollectI conjI holf)
      have  $\text{norm}(f \ z) \leq 1$  if  $z \in S$  for z
      proof (intro Lim_norm_ubound [OF _ limf] always_eventually allI that)
        fix n
        have  $\mathcal{F} \ (r \ n) \in F$ 
          by (simp add: Fin)
        then show  $\text{norm } (\mathcal{F} \ (r \ n) \ z) \leq 1$ 
          using that by (auto simp: F_def)
      qed simp
      then have fless1:  $\text{norm}(f \ z) < 1$  if  $z \in S$  for z
      using maximum_modulus_principle [OF holf ‹open S› ‹connected S› ‹open
S›] nonconstf that
        by fastforce
      then show  $f \ ' \ S \subseteq \text{ball } 0 \ 1$ 
        by auto
      have ( $\lambda n. \mathcal{F} \ (r \ n) \ 0) \longrightarrow 0$ 
        using Fin by (auto simp: F_def)
      then show  $f \ 0 = 0$ 

```

```

    using tendsto_unique [OF _ limf ] ⟨0 ∈ S⟩ trivial_limit_sequentially by
blast
show inj_on f S
proof (rule Hurwitz_injective [OF ⟨open S⟩ ⟨connected S⟩ _ holF])
  show  $\bigwedge n. (\mathcal{F} \circ r) n$  holomorphic_on S
  by (simp add: Fin holF)
show  $\bigwedge K. [\text{compact } K; K \subseteq S] \implies$  uniform_limit K  $(\mathcal{F} \circ r) f$  sequentially
  by (metis ulimf)
show  $\neg f$  constant_on S
  using nonconstf by auto
show  $\bigwedge n. inj\_on ((\mathcal{F} \circ r) n) S$ 
  using Fin by (auto simp: F_def)
qed
qed
show  $\bigwedge h. h \in F \implies$  norm (deriv h 0)  $\leq$  norm (deriv f 0)
  by (metis eql le_cases no_df0)
qed
qed
have holF: f holomorphic_on S and injf: inj_on f S and f01: f ' S  $\subseteq$  ball 0 1
  using ⟨f ∈ F⟩ by (auto simp: F_def)
obtain g where holg: g holomorphic_on (f ' S)
  and derg:  $\bigwedge z. z \in S \implies$  deriv f z * deriv g (f z) = 1
  and gf:  $\bigwedge z. z \in S \implies$  g(f z) = z
  using holomorphic_has_inverse [OF holF ⟨open S⟩ injf] by metis
have ball 0 1  $\subseteq$  f ' S
proof
  fix a::complex
  assume a: a ∈ ball 0 1
  have False if  $\bigwedge x. x \in S \implies$  f x  $\neq$  a
  proof -
    obtain h k where h a = 0
      and holh: h holomorphic_on ball 0 1 and h01: h ' ball 0 1  $\subseteq$  ball 0 1
      and holk: k holomorphic_on ball 0 1 and k01: k ' ball 0 1  $\subseteq$  ball 0 1
      and hk:  $\bigwedge z. z \in$  ball 0 1  $\implies$  h (k z) = z
      and kh:  $\bigwedge z. z \in$  ball 0 1  $\implies$  k (h z) = z
      using ball_biholomorphism_exists [OF a] by blast
    have nf1:  $\bigwedge z. z \in S \implies$  norm(f z) < 1
      using ⟨f ∈ F⟩ by (auto simp: F_def)
    have 1: h ∘ f holomorphic_on S
      using F_def ⟨f ∈ F⟩ holh holomorphic_on_compose holomorphic_on_subset
    by blast
    have 2:  $\bigwedge z. z \in S \implies$  (h ∘ f) z  $\neq$  0
      by (metis ⟨h a = 0⟩ a comp_eq_dest_lhs nf1 kh mem_ball_0 that)
    have 3: inj_on (h ∘ f) S
      by (metis (no_types, lifting) F_def ⟨f ∈ F⟩ comp_inj_on inj_on_inverseI
    injf kh mem_Collect_eq subset_inj_on)
    obtain  $\psi$  where hol $\psi$ :  $\psi$  holomorphic_on ((h ∘ f) ' S)
      and  $\psi$ 2:  $\bigwedge z. z \in S \implies$   $\psi(h(f z))$  ^ 2 = h(f z)
    proof (rule exE [OF prev [OF 1 2 3]], safe)

```

```

fix  $\vartheta$ 
assume hol $\vartheta$ :  $\vartheta$  holomorphic_on S and  $\vartheta 2$ :  $(\forall z \in S. (h \circ f) z = (\vartheta z)^2)$ 
show thesis
proof
  show  $(\vartheta \circ g \circ k)$  holomorphic_on  $(h \circ f) ' S$ 
  proof (intro holomorphic_on_compose)
    show  $k$  holomorphic_on  $(h \circ f) ' S$ 
    apply (rule holomorphic_on_subset [OF holk])
    using f01 h01 by force
    show  $g$  holomorphic_on  $k ' (h \circ f) ' S$ 
    apply (rule holomorphic_on_subset [OF holg])
    by (auto simp: kh nf1)
    show  $\vartheta$  holomorphic_on  $g ' k ' (h \circ f) ' S$ 
    apply (rule holomorphic_on_subset [OF hol $\vartheta$ ])
    by (auto simp: gf kh nf1)
  qed
  show  $((\vartheta \circ g \circ k) (h (f z)))^2 = h (f z)$  if  $z \in S$  for  $z$ 
  proof -
    have  $f z \in \text{ball } 0 \ 1$ 
    by (simp add: nf1 that)
    then have  $(\vartheta (g (k (h (f z))))^2 = (\vartheta (g (f z)))^2$ 
    by (metis kh)
    also have  $\dots = h (f z)$ 
    using  $\vartheta 2$  gf that by auto
    finally show ?thesis
    by (simp add: o_def)
  qed
  qed
  have norm $\psi 1$ :  $\text{norm}(\psi (h (f z))) < 1$  if  $z \in S$  for  $z$ 
  proof -
    have  $\text{norm} (\psi (h (f z)) ^ 2) < 1$ 
    by (metis (no_types) that DIM_complex  $\psi 2$  h01 image_subset_iff
    mem_ball_0 nf1)
    then show ?thesis
    by (metis le_less_trans mult_less_cancel_left2 norm_ge_zero norm_power
    not_le power2_eq_square)
  qed
  then have  $\psi 01$ :  $\psi (h (f 0)) \in \text{ball } 0 \ 1$ 
  by (simp add: <0  $\in S$ >)
  obtain  $p \ q$  where  $p 0$ :  $p (\psi (h (f 0))) = 0$ 
  and holp:  $p$  holomorphic_on ball 0 1 and  $p 01$ :  $p ' \text{ball } 0 \ 1 \subseteq \text{ball } 0 \ 1$ 
  and holq:  $q$  holomorphic_on ball 0 1 and  $q 01$ :  $q ' \text{ball } 0 \ 1 \subseteq \text{ball } 0 \ 1$ 
  and pq:  $\bigwedge z. z \in \text{ball } 0 \ 1 \implies p (q z) = z$ 
  and qp:  $\bigwedge z. z \in \text{ball } 0 \ 1 \implies q (p z) = z$ 
  using ball_biholomorphism_exists [OF  $\psi 01$ ] by metis
  have  $p \circ \psi \circ h \circ f \in F$ 
  unfolding F_def
  proof (intro CollectI conjI holp)

```

```

show  $p \circ \psi \circ h \circ f$  holomorphic_on  $S$ 
proof (intro holomorphic_on_compose holf)
  show  $h$  holomorphic_on  $f^{-1} S$ 
    apply (rule holomorphic_on_subset [OF holh])
    using f01 by force
  show  $\psi$  holomorphic_on  $h^{-1} f^{-1} S$ 
    apply (rule holomorphic_on_subset [OF hol $\psi$ ])
    by auto
  show  $p$  holomorphic_on  $\psi^{-1} h^{-1} f^{-1} S$ 
    apply (rule holomorphic_on_subset [OF holp])
    by (auto simp: norm $\psi$ 1)
qed
show  $(p \circ \psi \circ h \circ f)^{-1} S \subseteq \text{ball } 0 \ 1$ 
  apply clarsimp
  by (meson norm $\psi$ 1 p01 image_subset_iff mem_ball_0)
show  $(p \circ \psi \circ h \circ f) \ 0 = 0$ 
  by (simp add:  $\langle p(\psi(h(f\ 0))) = 0 \rangle$ )
show inj_on  $(p \circ \psi \circ h \circ f) S$ 
  unfolding inj_on_def o_def
  by (metis  $\psi$ 2 dist_0_norm gf kh mem_ball nf1 norm $\psi$ 1 qp)
qed
then have le_norm_df0:  $\text{norm}(\text{deriv}(p \circ \psi \circ h \circ f) \ 0) \leq \text{norm}(\text{deriv } f \ 0)$ 
  by (rule normf)
have 1:  $k \circ \text{power2} \circ q$  holomorphic_on ball 0 1
proof (intro holomorphic_on_compose holq)
  show  $\text{power2}$  holomorphic_on  $q^{-1} \text{ball } 0 \ 1$ 
    using holomorphic_on_subset holomorphic_on_power
    by (blast intro: holomorphic_on_ident)
  show  $k$  holomorphic_on  $\text{power2}^{-1} q^{-1} \text{ball } 0 \ 1$ 
    apply (rule holomorphic_on_subset [OF holk])
    using q01 by (auto simp: norm_power abs_square_less_1)
qed
have 2:  $(k \circ \text{power2} \circ q) \ 0 = 0$ 
  using p0 F_def  $\langle f \in F \rangle$   $\psi$ 01  $\psi$ 2  $\langle 0 \in S \rangle$  kh qp by force
have 3:  $\text{norm}((k \circ \text{power2} \circ q) \ z) < 1$  if  $\text{norm } z < 1$  for  $z$ 
proof -
  have  $\text{norm}((\text{power2} \circ q) \ z) < 1$ 
    using that q01 by (force simp: norm_power abs_square_less_1)
  with k01 show ?thesis
    by fastforce
qed
have False if  $c: \forall z. \text{norm } z < 1 \longrightarrow (k \circ \text{power2} \circ q) \ z = c * z$  and  $\text{norm } c = 1$  for  $c$ 
proof -
  have  $c \neq 0$  using that by auto
  have  $\text{norm}(p(1/2)) < 1$   $\text{norm}(p(-1/2)) < 1$ 
    using p01 by force+
  then have  $(k \circ \text{power2} \circ q) \ (p(1/2)) = c * p(1/2)$   $(k \circ \text{power2} \circ q) \ (p(-1/2)) = c * p(-1/2)$ 

```

```

    using c by force+
  then have  $p (1/2) = p (- (1/2))$ 
    by (auto simp:  $\langle c \neq 0 \rangle$  qp o_def)
  then have  $q (p (1/2)) = q (p (- (1/2)))$ 
    by simp
  then have  $1/2 = - (1/2::\text{complex})$ 
    by (auto simp: qp)
  then show False
    by simp
qed
moreover
have False if norm (deriv (k o power2 o q) 0)  $\neq 1$  norm (deriv (k o power2
o q) 0)  $\leq 1$ 
  and le:  $\bigwedge \xi. \text{norm } \xi < 1 \implies \text{norm } ((k \circ \text{power2} \circ q) \xi) \leq \text{norm } \xi$ 
proof -
  have norm (deriv (k o power2 o q) 0)  $< 1$ 
    using that by simp
  moreover have eq:  $\text{deriv } f \ 0 = \text{deriv } (k \circ (\lambda z. z^2) \circ q) \ 0 * \text{deriv } (p \circ$ 
 $\psi \circ h \circ f) \ 0$ 
  proof (intro DERIV_imp_deriv_has_field_derivative_transform_within_open
[OF DERIV_chain])
    show (k o power2 o q has_field_derivative deriv (k o power2 o q) 0) (at
((p o  $\psi$  o h o f) 0))
      using 1 holomorphic_derivI p0 by auto
    show (p o  $\psi$  o h o f has_field_derivative deriv (p o  $\psi$  o h o f) 0) (at 0)
      using  $\langle p \circ \psi \circ h \circ f \in F \rangle \langle \text{open } S \rangle \langle 0 \in S \rangle \text{holF holomorphic\_derivI}$ 
by blast
    show  $\bigwedge x. x \in S \implies (k \circ \text{power2} \circ q \circ (p \circ \psi \circ h \circ f)) \ x = f \ x$ 
      using  $\psi^2 \ f01 \ kh \ \text{norm}\psi1 \ qp$  by auto
    qed (use assms in simp_all)
    ultimately have cmod (deriv (p o  $\psi$  o h o f) 0)  $\leq 0$ 
      using le_norm_df0
  by (metis linorder_not_le mult.commute mult_less_cancel_left2 norm_mult)
  moreover have  $1 \leq \text{norm } (\text{deriv } f \ 0)$ 
    using normf [of id] by (simp add: idF)
  ultimately show False
    by (simp add: eq)
qed
ultimately show ?thesis
  using Schwarz_Lemma [OF 1 2 3] norm_one by blast
qed
then show  $a \in f \ ' \ S$ 
  by blast
qed
then have  $f \ ' \ S = \text{ball } 0 \ 1$ 
  using F_def  $\langle f \in F \rangle$  by blast
then show ?thesis
  apply (rule_tac  $x=f$  in exI)
  apply (rule_tac  $x=g$  in exI)

```

```

    using holf holg derg gf by safe force+
qed

```

```

locale SC_Chain =
  fixes S :: complex set
  assumes openS: open S
begin

```

```

lemma winding_number_zero:
  assumes simply_connected S
  shows connected S  $\wedge$ 
    ( $\forall \gamma z. \text{path } \gamma \wedge \text{path\_image } \gamma \subseteq S \wedge$ 
       $\text{pathfinish } \gamma = \text{pathstart } \gamma \wedge z \notin S \longrightarrow \text{winding\_number } \gamma z = 0$ )
  using assms
  by (auto simp: simply_connected_imp_connected simply_connected_imp_winding_number_zero)

```

```

lemma contour_integral_zero:
  assumes valid_path g path_image g  $\subseteq$  S pathfinish g = pathstart g f holomorphic_on S
     $\wedge \gamma z. \llbracket \text{path } \gamma; \text{path\_image } \gamma \subseteq S; \text{pathfinish } \gamma = \text{pathstart } \gamma; z \notin S \rrbracket \implies$ 
    winding_number  $\gamma z = 0$ 
  shows (f has_contour_integral 0) g
  using assms by (meson Cauchy_theorem_global openS valid_path_imp_path)

```

```

lemma global_primitive:
  assumes connected S and holf: f holomorphic_on S
  and prev:  $\wedge \gamma f. \llbracket \text{valid\_path } \gamma; \text{path\_image } \gamma \subseteq S; \text{pathfinish } \gamma = \text{pathstart } \gamma; f$ 
    holomorphic_on S  $\rrbracket \implies (f \text{ has\_contour\_integral } 0) \gamma$ 
  shows  $\exists h. \forall z \in S. (h \text{ has\_field\_derivative } f z) (\text{at } z)$ 
proof (cases S = {})
case True then show ?thesis
  by simp
next
  case False
  then obtain a where a  $\in$  S
  by blast
  show ?thesis
  proof (intro exI ballI)
    fix x assume x  $\in$  S
    then obtain d where d > 0 and d: cball x d  $\subseteq$  S
    using openS open_contains_cball_eq by blast
    let ?g =  $\lambda z. (\text{SOME } g. \text{polynomial\_function } g \wedge \text{path\_image } g \subseteq S \wedge \text{pathstart}$ 
       $g = a \wedge \text{pathfinish } g = z)$ 
    show (( $\lambda z. \text{contour\_integral } (?g z) f$ ) has_field_derivative f x)
      (at x)
  proof (simp add: has_field_derivative_def has_derivative_at2 bounded_linear_mult_right,
    rule Lim_transform)
    show ( $\lambda y. \text{inverse}(\text{norm}(y - x)) *_R (\text{contour\_integral}(\text{linepath } x y) f - f x$ 

```

```

* (y - x))) -x → 0
  proof (clarsimp simp add: Lim_at)
    fix e::real assume e > 0
    moreover have continuous (at x) f
      using openS ⟨x ∈ S⟩ holf continuous_on_eq_continuous_at holomor-
phic_on_imp_continuous_on by auto
    ultimately obtain d1 where d1 > 0
      and d1: ∧x'. dist x' x < d1 ⇒ dist (f x') (f x) < e/2
      unfolding continuous_at_eps_delta
      by (metis less_divide_eq numeral1(1) mult_zero_left)
    obtain d2 where d2 > 0 and d2: ball x d2 ⊆ S
      using openS ⟨x ∈ S⟩ open_contains_ball_eq by blast
    have inverse (norm (y - x)) * norm (contour_integral (linepath x y) f - f
x * (y - x)) < e
      if 0 < d1 0 < d2 y ≠ x dist y x < d1 dist y x < d2 for y
      proof -
        have f contour_integrable_on linepath x y
          proof (rule contour_integrable_continuous_linepath [OF continuous_on_subset])
            show continuous_on S f
              by (simp add: holf holomorphic_on_imp_continuous_on)
            have closed_segment x y ⊆ ball x d2
              by (meson dist_commute_lessI dist_in_closed_segment le_less_trans
mem_ball_subsetI that(5))
            with d2 show closed_segment x y ⊆ S
              by blast
          qed
        then obtain z where z: (f has_contour_integral z) (linepath x y)
          by (force simp: contour_integrable_on_def)
        have con: ((λw. f x) has_contour_integral f x * (y - x)) (linepath x y)
          using has_contour_integral_const_linepath [of f x y x] by metis
        have norm (z - f x * (y - x)) ≤ (e/2) * norm (y - x)
          proof (rule has_contour_integral_bound_linepath)
            show ((λw. f w - f x) has_contour_integral z - f x * (y - x)) (linepath
x y)
              by (rule has_contour_integral_diff [OF z con])
            show ∧w. w ∈ closed_segment x y ⇒ norm (f w - f x) ≤ e/2
              by (metis d1 dist_norm_less_le_trans not_less not_less_iff_gr_or_eq
segment_bound1 that(4))
          qed (use ⟨e > 0⟩ in auto)
        with ⟨e > 0⟩ have inverse (norm (y - x)) * norm (z - f x * (y - x))
≤ e/2
          by (simp add: field_split_simps)
        also have ... < e
          using ⟨e > 0⟩ by simp
        finally show ?thesis
          by (simp add: contour_integral_unique [OF z])
      qed
    with ⟨d1 > 0⟩ ⟨d2 > 0⟩
    show ∃d>0. ∀z. z ≠ x ∧ dist z x < d →

```



```

      inverse (norm (z - x)) * norm (contour_integral (linepath x z) f -
f x * (z - x)) < e
    by (rule_tac x=min d1 d2 in exI) auto
  qed
next
  have *: (1 / norm (y - x)) *R (contour_integral (?g y) f -
    (contour_integral (?g x) f + f x * (y - x))) =
    (contour_integral (linepath x y) f - f x * (y - x)) /R norm (y - x)
  if 0 < d y ≠ x and yx: dist y x < d for y
  proof -
    have y ∈ S
      by (metis subsetD d dist_commute less_eq_real_def mem_cball yx)
    have gxy: polynomial_function (?g x) ∧ path_image (?g x) ⊆ S ∧ pathstart
      (?g x) = a ∧ pathfinish (?g x) = x
      polynomial_function (?g y) ∧ path_image (?g y) ⊆ S ∧ pathstart
      (?g y) = a ∧ pathfinish (?g y) = y
      using someI_ex [OF connected_open_polynomial_connected [OF openS
      ‹connected S› ‹a ∈ S›]] ‹x ∈ S› ‹y ∈ S›
      by meson+
    then have vp: valid_path (?g x) valid_path (?g y)
      by (simp_all add: valid_path_polynomial_function)
    have f0: (f has_contour_integral 0) ((?g x) +++ linepath x y +++
reversepath (?g y))
      proof (rule prev)
        show valid_path ((?g x) +++ linepath x y +++ reversepath (?g y))
          using gxy vp by (auto simp: valid_path_join)
        have closed_segment x y ⊆ cball x d
          using yx by (auto simp: dist_commute dest!: dist_in_closed_segment)
        then have closed_segment x y ⊆ S
          using d by blast
        then show path_image ((?g x) +++ linepath x y +++ reversepath (?g
y)) ⊆ S
          using gxy by (auto simp: path_image_join)
        qed (use gxy holf in auto)
    then have fintry: f contour_integrable_on linepath x y
      by (metis (no_types, lifting) contour_integrable_joinD1 contour_integrable_joinD2
gxy(2) has_contour_integral_integrable pathfinish_linepath pathstart_reversepath
valid_path_imp_reverse valid_path_join valid_path_linepath vp(2))
    have fintgx: f contour_integrable_on (?g x) f contour_integrable_on (?g y)
      using openS contour_integrable_holomorphic_simple gxy holf vp by blast+
    show ?thesis
      apply (clarsimp simp add: divide_simps)
      using contour_integral_unique [OF f0]
      apply (simp add: fintry gxy contour_integrable_reversepath contour_integral_reversepath
fintgx vp)
      apply (simp add: algebra_simps)
      done
  qed
  show (λz. (1 / norm (z - x)) *R

```

```

      (contour_integral (?g z) f - (contour_integral (?g x) f + f x * (z
- x))) -
      (contour_integral (linepath x z) f - f x * (z - x)) /R norm (z - x))
    -x→ 0
  apply (rule tendsto_eventually)
  apply (simp add: eventually_at)
  apply (rule_tac x=d in exI)
  using ‹d > 0› * by simp
qed
qed
qed

```

lemma *holomorphic_log*:

```

  assumes connected S and holf: f holomorphic_on S and nz:  $\bigwedge z. z \in S \implies f z \neq 0$ 
  and prev:  $\bigwedge f. f \text{ holomorphic\_on } S \implies \exists h. \forall z \in S. (h \text{ has\_field\_derivative } f z)$ 
  (at z)
  shows  $\exists g. g \text{ holomorphic\_on } S \wedge (\forall z \in S. f z = \exp(g z))$ 
proof -
  have  $(\lambda z. \text{deriv } f z / f z) \text{ holomorphic\_on } S$ 
  by (simp add: openS holf holomorphic_deriv holomorphic_on_divide nz)
  then obtain g where  $g: \bigwedge z. z \in S \implies (g \text{ has\_field\_derivative } \text{deriv } f z / f z)$ 
  (at z)
  using prev [of  $\lambda z. \text{deriv } f z / f z$ ] by metis
  have hfd:  $\bigwedge x. x \in S \implies ((\lambda z. \exp (g z) / f z) \text{ has\_field\_derivative } 0)$  (at x)
  apply (rule derivative_eq_intros g| simp)+
  apply (subst DERIV_deriv_iff_field_differentiable)
  using openS holf holomorphic_on_imp_differentiable_at nz apply auto
  done
  obtain c where  $c: \bigwedge x. x \in S \implies \exp (g x) / f x = c$ 
proof (rule DERIV_zero_connected_constant[OF ‹connected› openS finite.emptyI])
  show continuous_on S  $(\lambda z. \exp (g z) / f z)$ 
  by (metis (full_types) openS g continuous_on_divide continuous_on_exp holf
holomorphic_on_imp_continuous_on holomorphic_on_open nz)
  then show  $\forall x \in S - \{ \}. ((\lambda z. \exp (g z) / f z) \text{ has\_field\_derivative } 0)$  (at x)
  using hfd by (blast intro: DERIV_zero_connected_constant [OF ‹connected
S› openS finite.emptyI, of  $\lambda z. \exp (g z) / f z$ ])
qed auto
show ?thesis
proof (intro exI ballI conjI)
  show  $(\lambda z. \text{Ln}(\text{inverse } c) + g z) \text{ holomorphic\_on } S$ 
  apply (intro holomorphic_intros)
  using openS g holomorphic_on_open by blast
  fix z :: complex
  assume z  $\in S$ 
  then have  $\exp (g z) / c = f z$ 
  by (metis c divide_divide_eq_right exp_not_eq_zero nonzero_mult_div_cancel_left)
  moreover have  $1 / c \neq 0$ 
  using ‹z  $\in S$ › c nz by fastforce

```

ultimately show $f z = \exp (Ln (\text{inverse } c) + g z)$
 by (simp add: exp_add inverse_eq_divide)
 qed
 qed

lemma holomorphic_sqrt:

assumes holf: f holomorphic_on S and nz: $\bigwedge z. z \in S \implies f z \neq 0$
 and prev: $\bigwedge f. [f \text{ holomorphic_on } S; \forall z \in S. f z \neq 0] \implies \exists g. g \text{ holomorphic_on } S \wedge (\forall z \in S. f z = \exp(g z))$
 shows $\exists g. g \text{ holomorphic_on } S \wedge (\forall z \in S. f z = (g z)^2)$
 proof –
 obtain g where holg: $g \text{ holomorphic_on } S$ and g : $\bigwedge z. z \in S \implies f z = \exp (g z)$
 using prev [of f] holf nz by metis
 show ?thesis
 proof (intro exI ballI conjI)
 show $(\lambda z. \exp(g z/2)) \text{ holomorphic_on } S$
 by (intro holomorphic_intros) (auto simp: holg)
 show $\bigwedge z. z \in S \implies f z = (\exp (g z/2))^2$
 by (metis (no_types) g exp_double nonzero_mult_div_cancel_left times_divide_eq_right zero_neq_numeral)
 qed
 qed

lemma biholomorphic_to_disc:

assumes connected S and S : $S \neq \{\}$ $S \neq \text{UNIV}$
 and prev: $\bigwedge f. [f \text{ holomorphic_on } S; \forall z \in S. f z \neq 0] \implies \exists g. g \text{ holomorphic_on } S \wedge (\forall z \in S. f z = (g z)^2)$
 shows $\exists f g. f \text{ holomorphic_on } S \wedge g \text{ holomorphic_on ball } 0 \ 1 \wedge$
 $(\forall z \in S. f z \in \text{ball } 0 \ 1 \wedge g(f z) = z) \wedge$
 $(\forall z \in \text{ball } 0 \ 1. g z \in S \wedge f(g z) = z)$
 proof –
 obtain $a b$ where $a \in S$ $b \notin S$
 using S by blast
 then obtain δ where $\delta > 0$ and δ : $\text{ball } a \ \delta \subseteq S$
 using openS openE by blast
 obtain g where holg: $g \text{ holomorphic_on } S$ and eqg: $\bigwedge z. z \in S \implies z - b = (g z)^2$
 proof (rule exE [OF prev [of $\lambda z. z - b$]])
 show $(\lambda z. z - b) \text{ holomorphic_on } S$
 by (intro holomorphic_intros)
 qed (use $\langle b \notin S \rangle$ in auto)
 have $\neg g \text{ constant_on } S$
 proof –
 have $(a + \delta/2) \in \text{ball } a \ \delta$ $a + (\delta/2) \neq a$
 using $\langle \delta > 0 \rangle$ by (simp_all add: dist_norm)
 then show ?thesis
 unfolding constant_on_def
 using eqg [of a] eqg [of $a + \delta/2$] $\langle a \in S \rangle$ δ

```

    by (metis diff_add_cancel subset_eq)
  qed
  then have open (g ` ball a δ)
    using open_mapping_thm [of g S ball a δ, OF holg openS ⟨connected S⟩] δ by
blast
  then obtain r where r > 0 and r: ball (g a) r ⊆ (g ` ball a δ)
    by (metis ⟨0 < δ⟩ centre_in_ball imageI openE)
  have g_not_r: g z ∉ ball (-(g a)) r if z ∈ S for z
  proof
    assume g z ∈ ball (-(g a)) r
    then have - g z ∈ ball (g a) r
      by (metis add.inverse_inverse dist_minus mem_ball)
    with r have - g z ∈ (g ` ball a δ)
      by blast
    then obtain w where w: - g z = g w dist a w < δ
      by auto
    then have w ∈ ball a δ
      by simp
    then have w ∈ S
      using δ by blast
    then have w = z
      by (metis diff_add_cancel eqg power_minus_Bit0 that w(1))
    then have g z = 0
      using ⟨- g z = g w⟩ by auto
    with eqg [OF that] have z = b
      by auto
    with that ⟨b ∉ S⟩ show False
      by simp
  qed
  then have nz: ⋀z. z ∈ S ⇒ g z + g a ≠ 0
    by (metis ⟨0 < r⟩ add commute add_diff_cancel_left' centre_in_ball diff_0)
  let ?f = λz. (r/3) / (g z + g a) - (r/3) / (g a + g a)
  obtain h where holh: h holomorphic_on S and h a = 0 and h01: h ` S ⊆ ball
0 1 and inj_on h S
  proof
    show ?f holomorphic_on S
      by (intro holomorphic_intros holg nz)
    have 3: [norm x ≤ 1/3; norm y ≤ 1/3] ⇒ norm(x - y) < 1 for x y::complex
      using norm_triangle_ineq4 [of x y] by simp
    have norm ((r/3) / (g z + g a) - (r/3) / (g a + g a)) < 1 if z ∈ S for z
      apply (rule 3)
      unfolding norm_divide
      using ⟨r > 0⟩ g_not_r [OF ⟨z ∈ S⟩] g_not_r [OF ⟨a ∈ S⟩]
      by (simp_all add: field_split_simps dist_commute dist_norm)
  then show ?f ` S ⊆ ball 0 1
    by auto
  show inj_on ?f S
    using ⟨r > 0⟩ eqg apply (clarsimp simp: inj_on_def)
    by (metis diff_add_cancel)

```

```

qed auto
obtain k where holk: k holomorphic_on (h ' S)
  and derk:  $\bigwedge z. z \in S \implies \text{deriv } h \ z * \text{deriv } k \ (h \ z) = 1$ 
  and kh:  $\bigwedge z. z \in S \implies k(h \ z) = z$ 
  using holomorphic_has_inverse [OF holh openS <inj_on h S>] by metis

have 1: open (h ' S)
  by (simp add: <inj_on h S> holh openS open_mapping_thm3)
have 2: connected (h ' S)
  by (simp add: connected_continuous_image <connected S> holh holomorphic_on_imp_continuous_on)
have 3:  $0 \in h \ ' \ S$ 
  using <a  $\in S$ > <h a = 0> by auto
have 4:  $\exists g. g \text{ holomorphic\_on } h \ ' \ S \wedge (\forall z \in h \ ' \ S. f \ z = (g \ z)^2)$ 
  if holf: f holomorphic_on h ' S and nz:  $\bigwedge z. z \in h \ ' \ S \implies f \ z \neq 0$  inj_on f (h
' S) for f
  proof -
    obtain g where holg: g holomorphic_on S and egg:  $\bigwedge z. z \in S \implies (f \circ h) \ z$ 
=  $(g \ z)^2$ 
    proof -
      have f  $\circ$  h holomorphic_on S
      by (simp add: holh holomorphic_on_compose holf)
      moreover have  $\forall z \in S. (f \circ h) \ z \neq 0$ 
      by (simp add: nz)
      ultimately show thesis
      using prev that by blast
    qed
  show ?thesis
  proof (intro exI conjI)
    show g  $\circ$  k holomorphic_on h ' S
    proof -
      have k ' h ' S  $\subseteq S$ 
      by (simp add: < $\bigwedge z. z \in S \implies k \ (h \ z) = z$ > image_subset_iff)
      then show ?thesis
      by (meson holg holk holomorphic_on_compose holomorphic_on_subset)
    qed
    show  $\forall z \in h \ ' \ S. f \ z = ((g \circ k) \ z)^2$ 
    using egg kh by auto
  qed
qed
obtain f g where f: f holomorphic_on h ' S and g: g holomorphic_on ball 0 1
  and gf:  $\forall z \in h \ ' \ S. f \ z \in \text{ball } 0 \ 1 \wedge g \ (f \ z) = z$  and fg:  $\forall z \in \text{ball } 0 \ 1. g \ z \in h$ 
' S  $\wedge f \ (g \ z) = z$ 
  using biholomorphic_to_disc_aux [OF 1 2 3 h01 4] by blast
show ?thesis
proof (intro exI conjI)
  show f  $\circ$  h holomorphic_on S
  by (simp add: f holh holomorphic_on_compose)
  show k  $\circ$  g holomorphic_on ball 0 1
  by (metis holomorphic_on_subset image_subset_iff fg holk g holomor-

```

phic_on_compose)

qed (*use fg gf kh in auto*)

qed

lemma *homeomorphic_to_disc:*

assumes $S: S \neq \{\}$

and *prev:* $S = UNIV \vee$

$(\exists f g. f \text{ holomorphic_on } S \wedge g \text{ holomorphic_on ball } 0 \ 1 \wedge$

$(\forall z \in S. f z \in \text{ball } 0 \ 1 \wedge g(f z) = z) \wedge$

$(\forall z \in \text{ball } 0 \ 1. g z \in S \wedge f(g z) = z))$ (**is** $_ \vee ?P$)

shows $S \text{ homeomorphic ball } (0::\text{complex}) \ 1$

using *prev*

proof

assume $S = UNIV$ **then show** *?thesis*

using *homeomorphic_ball01_UNIV homeomorphic_sym* **by** *blast*

next

assume *?P*

then show *?thesis*

unfolding *homeomorphic_minimal*

using *holomorphic_on_imp_continuous_on* **by** *blast*

qed

lemma *homeomorphic_to_disc_imp_simply_connected:*

assumes $S = \{\} \vee S \text{ homeomorphic ball } (0::\text{complex}) \ 1$

shows *simply_connected* S

using *assms homeomorphic_simply_connected_eq convex_imp_simply_connected*
by *auto*

end

proposition

assumes *open* S

shows *simply_connected_eq_winding_number_zero:*

$\text{simply_connected } S \iff$

$\text{connected } S \wedge$

$(\forall g z. \text{path } g \wedge \text{path_image } g \subseteq S \wedge$

$\text{pathfinish } g = \text{pathstart } g \wedge (z \notin S)$

$\implies \text{winding_number } g \ z = 0)$ (**is** *?wn0*)

and *simply_connected_eq_contour_integral_zero:*

$\text{simply_connected } S \iff$

$\text{connected } S \wedge$

$(\forall g f. \text{valid_path } g \wedge \text{path_image } g \subseteq S \wedge$

$\text{pathfinish } g = \text{pathstart } g \wedge f \text{ holomorphic_on } S$

$\implies (f \text{ has_contour_integral } 0) \ g)$ (**is** *?ci0*)

and *simply_connected_eq_global_primitive:*

$\text{simply_connected } S \iff$

$\text{connected } S \wedge$

$(\forall f. f \text{ holomorphic_on } S \implies$

$(\exists h. \forall z. z \in S \implies (h \text{ has_field_derivative } f z) \text{ (at } z)))$ (**is** *?gp*)

```

and simply_connected_eq_holomorphic_log:
  simply_connected S  $\longleftrightarrow$ 
    connected S  $\wedge$ 
      ( $\forall f. f$  holomorphic_on S  $\wedge$  ( $\forall z \in S. f z \neq 0$ )
         $\longrightarrow$  ( $\exists g. g$  holomorphic_on S  $\wedge$  ( $\forall z \in S. f z = \exp(g z)$ ))) (is ?log)
and simply_connected_eq_holomorphic_sqrt:
  simply_connected S  $\longleftrightarrow$ 
    connected S  $\wedge$ 
      ( $\forall f. f$  holomorphic_on S  $\wedge$  ( $\forall z \in S. f z \neq 0$ )
         $\longrightarrow$  ( $\exists g. g$  holomorphic_on S  $\wedge$  ( $\forall z \in S. f z = (g z)^2$ ))) (is ?sqrt)
and simply_connected_eq_biholomorphic_to_disc:
  simply_connected S  $\longleftrightarrow$ 
    S = {}  $\vee$  S = UNIV  $\vee$ 
      ( $\exists f g. f$  holomorphic_on S  $\wedge$   $g$  holomorphic_on ball 0 1  $\wedge$ 
        ( $\forall z \in S. f z \in$  ball 0 1  $\wedge$   $g(f z) = z$ )  $\wedge$ 
        ( $\forall z \in$  ball 0 1.  $g z \in S$   $\wedge$   $f(g z) = z$ ) (is ?bih)
and simply_connected_eq_homeomorphic_to_disc:
  simply_connected S  $\longleftrightarrow$  S = {}  $\vee$  S homeomorphic ball (0::complex) 1
(is ?disc)
proof -
  interpret SC_Chain
  using assms by (simp add: SC_Chain_def)
  have ?wn0  $\wedge$  ?ci0  $\wedge$  ?gp  $\wedge$  ?log  $\wedge$  ?sqrt  $\wedge$  ?bih  $\wedge$  ?disc
proof -
  have *:  $\llbracket \alpha \implies \beta; \beta \implies \gamma; \gamma \implies \delta; \delta \implies \zeta; \zeta \implies \eta; \eta \implies \vartheta; \vartheta \implies \xi; \xi \implies \alpha \rrbracket$ 
     $\implies (\alpha \longleftrightarrow \beta) \wedge (\alpha \longleftrightarrow \gamma) \wedge (\alpha \longleftrightarrow \delta) \wedge (\alpha \longleftrightarrow \zeta) \wedge$ 
      ( $\alpha \longleftrightarrow \eta) \wedge (\alpha \longleftrightarrow \vartheta) \wedge (\alpha \longleftrightarrow \xi)$  for  $\alpha \beta \gamma \delta \zeta \eta \vartheta \xi$ 
  by blast
  show ?thesis
  apply (rule *)
  using winding_number_zero apply metis
  using contour_integral_zero apply metis
  using global_primitive apply metis
  using holomorphic_log apply metis
  using holomorphic_sqrt apply simp
  using biholomorphic_to_disc apply blast
  using homeomorphic_to_disc apply blast
  using homeomorphic_to_disc_imp_simply_connected apply blast
  done
qed
  then show ?wn0 ?ci0 ?gp ?log ?sqrt ?bih ?disc
  by safe
qed

corollary contractible_eq_simply_connected_2d:
  fixes S :: complex set
  shows open S  $\implies$  (contractible S  $\longleftrightarrow$  simply_connected S)
  apply safe

```

```

apply (simp add: contractible_imp_simply_connected)
using convex_imp_contractible_homeomorphic_contractible_eq_simply_connected_eq_homeomorphic_to_disc
by auto

```

8.3 A further chain of equivalences about components of the complement of a simply connected set

(following 1.35 in Burckel'S book)

```

context SC_Chain

```

```

begin

```

```

lemma frontier_properties:

```

```

  assumes simply_connected S

```

```

  shows if bounded S then connected(frontier S)

```

```

    else  $\forall C \in \text{components}(\text{frontier } S). \neg \text{bounded } C$ 

```

```

proof –

```

```

  have S = {}  $\vee$  S homeomorphic ball (0::complex) 1

```

```

    using simply_connected_eq_homeomorphic_to_disc assms open S by blast

```

```

  then show ?thesis

```

```

proof

```

```

  assume S = {}

```

```

  then have bounded S

```

```

    by simp

```

```

  with ⟨S = {}⟩ show ?thesis

```

```

    by simp

```

```

next

```

```

  assume S01: S homeomorphic ball (0::complex) 1

```

```

  then obtain g f

```

```

    where gim: g ‘ S = ball 0 1 and fg:  $\bigwedge x. x \in S \implies f(g x) = x$ 

```

```

    and fim: f ‘ ball 0 1 = S and gf:  $\bigwedge y. \text{cmod } y < 1 \implies g(f y) = y$ 

```

```

    and contg: continuous_on S g and conf: continuous_on (ball 0 1) f

```

```

    by (fastforce simp: homeomorphism_def homeomorphic_def)

```

```

  define D where D  $\equiv \lambda n. \text{ball } (0::\text{complex}) (1 - 1/(\text{of\_nat } n + 2))$ 

```

```

  define A where A  $\equiv \lambda n. \{z::\text{complex}. 1 - 1/(\text{of\_nat } n + 2) < \text{norm } z \wedge \text{norm } z < 1\}$ 

```

```

  define X where X  $\equiv \lambda n::\text{nat}. \text{closure}(f ‘ A n)$ 

```

```

  have D01: D n  $\subseteq$  ball 0 1 for n

```

```

    by (simp add: D_def ball_subset_ball_iff)

```

```

  have A01: A n  $\subseteq$  ball 0 1 for n

```

```

    by (auto simp: A_def)

```

```

  have cloX: closed(X n) for n

```

```

    by (simp add: X_def)

```

```

  have Xsubclo: X n  $\subseteq$  closure S for n

```

```

    unfolding X_def by (metis A01 closure_mono fim image_mono)

```

```

  have connX: connected(X n) for n

```

```

    unfolding X_def

```

```

    apply (rule connected_imp_connected_closure)

```

```

    apply (rule connected_continuous_image)

```

```

    apply (simp add: continuous_on_subset [OF contf A01])

```



```

    using connected_annulus [of _ 0::complex] by (simp add: A_def)
  have nestX:  $X\ n \subseteq X\ m$  if  $m \leq n$  for  $m\ n$ 
  proof -
    have  $1 - 1 / (\text{real } m + 2) \leq 1 - 1 / (\text{real } n + 2)$ 
      using that by (auto simp: field_simps)
    then show ?thesis
      by (auto simp: X_def A_def intro!: closure_mono)
  qed
  have closure  $S - S \subseteq (\bigcap n. X\ n)$ 
  proof
    fix  $x$ 
    assume  $x \in \text{closure } S - S$ 
    then have  $x \in \text{closure } S$   $x \notin S$  by auto
    show  $x \in (\bigcap n. X\ n)$ 
  proof
    fix  $n$ 
    have  $\text{ball } 0\ 1 = \text{closure } (D\ n) \cup A\ n$ 
      by (auto simp: D_def A_def le_less_trans)
    with  $f$  have  $\text{Seq}: S = f\ ' (\text{closure } (D\ n)) \cup f\ ' (A\ n)$ 
      by (simp add: image_Un)
    have continuous_on  $(\text{closure } (D\ n))\ f$ 
      by (simp add: D_def cball_subset_ball_iff continuous_on_subset [OF
contf])
    moreover have compact  $(\text{closure } (D\ n))$ 
      by (simp add: D_def)
    ultimately have  $\text{clo\_fim}: \text{closed } (f\ ' \text{closure } (D\ n))$ 
      using compact_continuous_image compact_imp_closed by blast
    have *:  $(f\ ' \text{cball } 0\ (1 - 1 / (\text{real } n + 2))) \subseteq S$ 
      by (force simp: D_def Seq)
    show  $x \in X\ n$ 
      using  $\langle x \in \text{closure } S \rangle$  unfolding X_def Seq
      using  $\langle x \notin S \rangle$  * D_def clo_fim by auto
  qed
  qed
  moreover have  $(\bigcap n. X\ n) \subseteq \text{closure } S - S$ 
  proof -
    have  $(\bigcap n. X\ n) \subseteq \text{closure } S$ 
  proof -
    have  $(\bigcap n. X\ n) \subseteq X\ 0$ 
      by blast
    also have  $\dots \subseteq \text{closure } S$ 
      apply (simp add: X_def fim [symmetric])
      apply (rule closure_mono)
      by (auto simp: A_def)
    finally show  $(\bigcap n. X\ n) \subseteq \text{closure } S$  .
  qed
  moreover have  $(\bigcap n. X\ n) \cap S \subseteq \{\}$ 
  proof (clarify, clarsimp simp: X_def fim [symmetric])
    fix  $x$  assume  $x$  [rule_format]:  $\forall n. f\ x \in \text{closure } (f\ ' A\ n)$  and  $\text{cmod } x < 1$ 

```

```

then obtain  $n$  where  $n: 1 / (1 - \text{norm } x) < \text{of\_nat } n$ 
  using reals_Archimedean2 by blast
with  $\langle \text{cmod } x < 1 \rangle$  gr0I have  $XX: 1 / \text{of\_nat } n < 1 - \text{norm } x$  and  $n > 0$ 
  by (fastforce simp: field_split_simps algebra_simps)+
have  $f x \in f' (D n)$ 
  using  $n \langle \text{cmod } x < 1 \rangle$  by (auto simp: field_split_simps algebra_simps
D_def)
moreover have  $f' D n \cap \text{closure } (f' A n) = \{\}$ 
proof -
  have op_fDn:  $\text{open}(f' (D n))$ 
  proof (rule invariance_of_domain)
    show continuous_on  $(D n) f$ 
      by (rule continuous_on_subset [OF contf D01])
    show open  $(D n)$ 
      by (simp add: D_def)
    show inj_on  $f (D n)$ 
      unfolding inj_on_def using D01 by (metis gf mem_ball_0 subsetCE)
  qed
have injf: inj_on  $f (\text{ball } 0 1)$ 
  by (metis mem_ball_0 inj_on_def gf)
have  $D n \cup A n \subseteq \text{ball } 0 1$ 
  using D01 A01 by simp
moreover have  $D n \cap A n = \{\}$ 
  by (auto simp: D_def A_def)
ultimately have  $f' D n \cap f' A n = \{\}$ 
  by (metis A01 D01 image_is_empty inj_on_image_Int injf)
then show ?thesis
  by (simp add: open_Int_closure_eq_empty [OF op_fDn])
qed
ultimately show False
  using  $x [of n]$  by blast
qed
ultimately
show  $(\bigcap n. X n) \subseteq \text{closure } S - S$ 
  using closure_subset disjoint_iff_not_equal by blast
qed
ultimately have  $\text{closure } S - S = (\bigcap n. X n)$  by blast
then have frontierS:  $\text{frontier } S = (\bigcap n. X n)$ 
  by (simp add: frontier_def openS interior_open)
show ?thesis
proof (cases bounded S)
  case True
  have bouX: bounded  $(X n)$  for  $n$ 
    apply (simp add: X_def)
    apply (rule bounded_closure)
    by (metis A01 fim image_mono bounded_subset [OF True])
  have compaX: compact  $(X n)$  for  $n$ 
    apply (simp add: compact_eq_bounded_closed bouX)
    apply (auto simp: X_def)

```

```

done
have connected ( $\bigcap n. X n$ )
  by (metis nestX compaX connX connected_nest)
then show ?thesis
  by (simp add: True ‹frontier S = ( $\bigcap n. X n$ )›)
next
case False
have unboundedX:  $\neg$  bounded(X n) for n
proof
  assume bXn: bounded(X n)
  have continuous_on (cball 0 (1 - 1 / (2 + real n))) f
    by (simp add: cball_subset_ball_iff continuous_on_subset [OF contf])
  then have bounded (f ‹ cball 0 (1 - 1 / (2 + real n)))
    by (simp add: compact_imp_bounded [OF compact_continuous_image])
  moreover have bounded (f ‹ A n)
    by (auto simp: X_def closure_subset image_subset_iff bounded_subset
[OF bXn])
  ultimately have bounded (f ‹ (cball 0 (1 - 1/(2 + real n))  $\cup$  A n))
    by (simp add: image_Un)
  then have bounded (f ‹ ball 0 1)
    apply (rule bounded_subset)
    apply (auto simp: A_def algebra_simps)
  done
  then show False
    using False by (simp add: fim [symmetric])
qed
have clo_INTX: closed( $\bigcap$ (range X))
  by (metis cloX closed_INT)
then have lcX: locally compact ( $\bigcap$ (range X))
  by (metis closed_imp_locally_compact)
have False if C: C  $\in$  components (frontier S) and boC: bounded C for C
proof -
  have closed C
    by (metis C closed_components frontier_closed)
  then have compact C
    by (metis boC compact_eq_bounded_closed)
  have Cco: C  $\in$  components ( $\bigcap$ (range X))
    by (metis frontierS C)
  obtain K where C  $\subseteq$  K compact K
    and Ksub: K  $\subseteq$   $\bigcap$ (range X) and clo: closed( $\bigcap$ (range X) - K)
  proof (cases {k. C  $\subseteq$  k  $\wedge$  compact k  $\wedge$  openin (top_of_set ( $\bigcap$ (range X)))
k} = {})
    case True
    then show ?thesis
      using Sura_Bura [OF lcX Cco ‹compact C›] boC
      by (simp add: True)
  next
    case False
    then obtain L where compact L C  $\subseteq$  L and K: openin (top_of_set

```

```

( $\bigcap x. X x$ ) L
  by blast
show ?thesis
proof
  show  $L \subseteq \bigcap (\text{range } X)$ 
    by (metis K openin_imp_subset)
  show closed ( $\bigcap (\text{range } X) - L$ )
    by (metis closedin_diff closedin_self closedin_closed_trans [OF _
clo_INTX] K)
  qed (use <compact L> < $C \subseteq L$ > in auto)
qed
obtain U V where open U and compact (closure U) and open V  $K \subseteq U$ 
  and  $V: \bigcap (\text{range } X) - K \subseteq V$  and  $U \cap V = \{\}$ 
  using separation_normal_compact [OF <compact K> clo] by blast
then have  $U \cap (\bigcap (\text{range } X) - K) = \{\}$ 
  by blast
have (closure U - U)  $\cap (\bigcap n. X n \cap \text{closure } U) \neq \{\}$ 
proof (rule compact_imp_fip)
  show compact (closure U - U)
    by (metis <compact (closure U)> <open U> compact_diff)
  show  $\bigwedge T. T \in \text{range } (\lambda n. X n \cap \text{closure } U) \implies \text{closed } T$ 
    by clarify (metis cloX closed_Int closed_closure)
  show (closure U - U)  $\cap \bigcap \mathcal{F} \neq \{\}$ 
    if finite  $\mathcal{F}$  and  $\mathcal{F}: \mathcal{F} \subseteq \text{range } (\lambda n. X n \cap \text{closure } U)$  for  $\mathcal{F}$ 
  proof
    assume empty: (closure U - U)  $\cap \bigcap \mathcal{F} = \{\}$ 
    obtain J where finite J and  $J: \mathcal{F} = (\lambda n. X n \cap \text{closure } U) \text{ ` } J$ 
      using finite_subset_image [OF <finite  $\mathcal{F}$ >  $\mathcal{F}$ ] by auto
    show False
    proof (cases  $J = \{\}$ )
      case True
        with J_empty have closed U
          by (simp add: closure_subset_eq)
        have  $C \neq \{\}$ 
          using C_in_components_nonempty by blast
        then have  $U \neq \{\}$ 
          using < $K \subseteq U$ > < $C \subseteq K$ > by blast
        moreover have  $U \neq \text{UNIV}$ 
          using <compact (closure U)> by auto
        ultimately show False
          using <open U> <closed U> clopen by blast
      case False
    next
      case False
        define j where  $j \equiv \text{Max } J$ 
        have  $j \in J$ 
          by (simp add: False <finite J> j_def)
        have  $j_{\text{max}}: \bigwedge m. m \in J \implies m \leq j$ 
          by (simp add: j_def <finite J>)
        have  $\bigcap ((\lambda n. X n \cap \text{closure } U) \text{ ` } J) = X j \cap \text{closure } U$ 

```

```

    using False jmax nestX ⟨j ∈ J⟩ by auto
  then have X j ∩ closure U = X j ∩ U
    apply safe
    using DiffI J empty apply auto[1]
    using closure_subset by blast
  then have openin (top_of_set (X j)) (X j ∩ closure U)
    by (simp add: openin_open_Int ⟨open U⟩)
  moreover have closedin (top_of_set (X j)) (X j ∩ closure U)
    by (simp add: closedin_closed_Int)
  moreover have X j ∩ closure U ≠ X j
    by (metis unboundedX ⟨compact (closure U)⟩ bounded_subset
compact_eq_bounded_closed_inf.order_iff)
  moreover have X j ∩ closure U ≠ {}
  proof -
    have C ≠ {}
      using C in_components_nonempty by blast
    moreover have C ⊆ X j ∩ closure U
      using ⟨C ⊆ K⟩ ⟨K ⊆ U⟩ Ksub_closure_subset by blast
    ultimately show ?thesis by blast
  qed
  ultimately show False
    using connX [of j] by (force simp: connected_clopen)
  qed
  qed
  qed
  moreover have (⋂ n. X n ∩ closure U) = (⋂ n. X n) ∩ closure U
    by blast
  moreover have x ∈ U if ⋀ n. x ∈ X n x ∈ closure U for x
  proof -
    have x ∉ V
      using ⟨U ∩ V = {}⟩ ⟨open V⟩ closure_iff_nhds_not_empty that(2) by
blast
    then show ?thesis
      by (metis (no_types) Diff_iff INT_I V ⟨K ⊆ U⟩ contra_subsetD that(1))
    qed
    ultimately show False
      by (auto simp: open_Int_closure_eq_empty [OF ⟨open V⟩, of U])
    qed
    then show ?thesis
      by (auto simp: False)
    qed
  qed
  qed
  qed

```

lemma *unbounded_complement_components:*

assumes $C \in \text{components } (- S)$ **and** $S \in \text{connected } S$
and prev: *if bounded S then connected(frontier S)*
else $\forall C \in \text{components(frontier S)}. \neg \text{bounded } C$

```

shows  $\neg$  bounded  $C$ 
proof (cases bounded  $S$ )
case True
with prev have  $S \neq \text{UNIV}$  and confr: connected(frontier  $S$ )
by auto
obtain  $w$  where  $C_{\text{ccsw}}: C = \text{connected\_component\_set } (- S) w$  and  $w \notin S$ 
using  $C$  by (auto simp: components_def)
show ?thesis
proof (cases  $S = \{\}$ )
case True with  $C$  show ?thesis by auto
next
case False
show ?thesis
proof
assume bounded  $C$ 
then have outside  $C \neq \{\}$ 
using outside_bounded_nonempty by metis
then obtain  $z$  where  $z: \neg$  bounded (connected_component_set  $(- C) z$ )
and  $z \notin C$ 
by (auto simp: outside_def)
have clo_ccs: closed (connected_component_set  $(- S) x$ ) for  $x$ 
by (simp add: closed_Compl closed_connected_component openS)
have connected_component_set  $(- S) w = \text{connected\_component\_set } (- S) z$ 
z
proof (rule joinable_connected_component_eq [OF confr])
show frontier  $S \subseteq - S$ 
using openS by (auto simp: frontier_def interior_open)
have False if connected_component_set  $(- S) w \cap \text{frontier } (- S) = \{\}$ 
proof -
have  $C \cap \text{frontier } S = \{\}$ 
using that by (simp add:  $C_{\text{ccsw}}$ )
then show False
by (metis  $C_{\text{ccsw}}$  ComplI Compl_eq_Compl_iff Diff_subset False  $\langle w \notin S \rangle$  clo_ccs closure_closed compl_bot_eq connected_component_eq UNIV connected_component_eq_empty empty_subsetI frontier_complement frontier_def frontier_not_empty frontier_of_connected_component_subset le_inf_iff subset_antisym)
qed
then show connected_component_set  $(- S) w \cap \text{frontier } S \neq \{\}$ 
by auto
have *:  $\llbracket \text{frontier } C \subseteq C; \text{frontier } C \subseteq F; \text{frontier } C \neq \{\} \rrbracket \implies C \cap F \neq \{\}$ 
for  $C F::\text{complex set}$ 
by blast
have connected_component_set  $(- S) z \cap \text{frontier } (- S) \neq \{\}$ 
proof (rule *)
show frontier (connected_component_set  $(- S) z) \subseteq \text{connected\_component\_set } (- S) z$ 
by (auto simp: closed_Compl closed_connected_component frontier_def openS)
show frontier (connected_component_set  $(- S) z) \subseteq \text{frontier } (- S)$ 

```

```

    using frontier_of_connected_component_subset by fastforce
  have  $\neg$  bounded  $(-S)$ 
    by (simp add: True cobounded_imp_unbounded)
  then have connected_component_set  $(- S) z \neq \{\}$ 
    apply (simp only: connected_component_eq_empty)
    using confr openS  $\langle$ bounded  $C\rangle \langle w \notin S\rangle$ 
    apply (simp add: frontier_def interior_open C_ccsw)
  by (metis ComplI Compl_eq_Diff_UNIV connected_UNIV closed_closure
closure_subset connected_component_eq_self
    connected_diff_open_from_closed subset_UNIV)
  then show frontier (connected_component_set  $(- S) z) \neq \{\}$ 
    apply (simp add: frontier_eq_empty connected_component_eq_UNIV)
    apply (metis False compl_top_eq double_compl)
  done
qed
then show connected_component_set  $(- S) z \cap$  frontier  $S \neq \{\}$ 
  by auto
qed
then show False
  by (metis C_ccsw Compl_iff  $\langle w \notin S\rangle \langle z \notin C\rangle$  connected_component_eq_empty
connected_component_idemp)
qed
qed
next
case False
obtain  $w$  where C_ccsw:  $C =$  connected_component_set  $(- S) w$  and  $w \notin S$ 
  using C by (auto simp: components_def)
have frontier (connected_component_set  $(- S) w) \subseteq$  connected_component_set
 $(- S) w$ 
  by (simp add: closed_Ccompl closed_connected_component frontier_subset_eq
openS)
moreover have frontier (connected_component_set  $(- S) w) \subseteq$  frontier  $S$ 
  using frontier_complement frontier_of_connected_component_subset by blast
moreover have frontier (connected_component_set  $(- S) w) \neq \{\}$ 
  by (metis C C_ccsw False bounded_empty compl_top_eq connected_component_eq_UNIV
double_compl frontier_not_empty in_components_nonempty)
ultimately obtain  $z$  where zin:  $z \in$  frontier  $S$  and  $z: z \in$  connected_component_set
 $(- S) w$ 
  by blast
have *: connected_component_set (frontier  $S) z \in$  components(frontier  $S)$ 
  by (simp add:  $\langle z \in$  frontier  $S\rangle$  componentsI)
with prev False have  $\neg$  bounded (connected_component_set (frontier  $S) z)$ 
  by simp
moreover have connected_component  $(- S) w =$  connected_component  $(- S)$ 
 $z$ 
  using connected_component_eq [OF  $z$ ] by force
ultimately show ?thesis
  by (metis C_ccsw * zin bounded_subset closed_Ccompl closure_closed con-
nected_component_maximal)

```

connected_component_refl connected_connected_component frontier_closures_in_components_subset le_inf_iff mem_Collect_eq openS
qed

lemma *empty_inside:*

assumes *connected S* \wedge *C. C* \in *components* $(- S) \implies \neg$ *bounded C*
shows *inside S = {}*
using *assms* **by** (*auto simp: components_def inside_def*)

lemma *empty_inside_imp_simply_connected:*

\llbracket *connected S; inside S = {}* $\rrbracket \implies$ *simply_connected S*
by (*metis ComplI inside_Un_outside openS outside_mono simply_connected_eq_winding_number_zero subsetCE sup_bot.left_neutral winding_number_zero_in_outside*)

end

proposition

fixes *S :: complex set*

assumes *open S*

shows *simply_connected_eq_frontier_properties:*

simply_connected S \longleftrightarrow

connected S \wedge

(*if bounded S then connected(frontier S)*

else $(\forall C \in$ *components(frontier S). \neg bounded C)) (**is** *?fp*)*

and *simply_connected_eq_unbounded_complement_components:*

simply_connected S \longleftrightarrow

connected S \wedge $(\forall C \in$ *components(- S). \neg bounded C)) (**is** *?ucc*)*

and *simply_connected_eq_empty_inside:*

simply_connected S \longleftrightarrow

connected S \wedge *inside S = {}*) (**is** *?ei*)

proof –

interpret *SC_Chain*

using *assms* **by** (*simp add: SC_Chain_def*)

have *?fp* \wedge *?ucc* \wedge *?ei*

using *empty_inside empty_inside_imp_simply_connected frontier_properties*

unbounded_complement_components winding_number_zero **by** *blast*

then show *?fp* *?ucc* *?ei*

by *blast+*

qed

lemma *simply_connected_iff_simple:*

fixes *S :: complex set*

assumes *open S bounded S*

shows *simply_connected S* \longleftrightarrow *connected S* \wedge *connected(- S)*

apply (*simp add: simply_connected_eq_unbounded_complement_components*
assms, safe)

apply (*metis DIM_complex assms(2) cobounded_has_bounded_component double_compl order_refl*)

by (*meson assms inside_bounded_complement_connected_empty simply_connected_eq_empty_inside*)

simply_connected_eq_unbounded_complement_components)

lemma *subset_simply_connected_imp_inside_subset*:

fixes $A :: \text{complex set}$

assumes *simply_connected* A *open* A $B \subseteq A$

shows *inside* $B \subseteq A$

by (*metis* *assms* *Diff_eq_empty_iff_inside_mono* *subset_empty* *simply_connected_eq_empty_inside*)

8.4 Further equivalences based on continuous logs and sqrts

context *SC_Chain*

begin

lemma *continuous_log*:

fixes $f :: \text{complex} \Rightarrow \text{complex}$

assumes S : *simply_connected* S

and *contf*: *continuous_on* S f **and** *nz*: $\bigwedge z. z \in S \implies f z \neq 0$

shows $\exists g. \text{continuous_on } S g \wedge (\forall z \in S. f z = \exp(g z))$

proof –

consider $S = \{ \} \mid S$ *homeomorphic* *ball* $(0 :: \text{complex}) 1$

using *simply_connected_eq_homeomorphic_to_disc* [*OF* *openS*] S **by** *metis*

then show *?thesis*

proof *cases*

case 1 **then show** *?thesis*

by *simp*

next

case 2

then obtain $h k :: \text{complex} \Rightarrow \text{complex}$

where *kh*: $\bigwedge x. x \in S \implies k(h x) = x$ **and** *him*: $h \text{ ' } S = \text{ball } 0 1$

and *conth*: *continuous_on* S h

and *hk*: $\bigwedge y. y \in \text{ball } 0 1 \implies h(k y) = y$ **and** *kim*: $k \text{ ' } \text{ball } 0 1 = S$

and *contk*: *continuous_on* $(\text{ball } 0 1)$ k

unfolding *homeomorphism_def* *homeomorphic_def* **by** *metis*

obtain g **where** *contg*: *continuous_on* $(\text{ball } 0 1)$ g

and *expg*: $\bigwedge z. z \in \text{ball } 0 1 \implies (f \circ k) z = \exp(g z)$

proof (*rule* *continuous_logarithm_on_ball*)

show *continuous_on* $(\text{ball } 0 1)$ $(f \circ k)$

using *contf* *continuous_on_compose* *contk* *kim* **by** *blast*

show $\bigwedge z. z \in \text{ball } 0 1 \implies (f \circ k) z \neq 0$

using *kim* *nz* **by** *auto*

qed *auto*

then show *?thesis*

by (*metis* *comp_apply* *conth* *continuous_on_compose* *him* *imageI* *kh*)

qed

qed

lemma *continuous_sqrt*:

fixes $f :: \text{complex} \Rightarrow \text{complex}$

assumes *contf*: *continuous_on* S f **and** *nz*: $\bigwedge z. z \in S \implies f z \neq 0$

```

and prev:  $\bigwedge f::\text{complex} \Rightarrow \text{complex}. \llbracket \text{continuous\_on } S f; \bigwedge z. z \in S \Longrightarrow f z \neq 0 \rrbracket$ 
       $\Longrightarrow \exists g. \text{continuous\_on } S g \wedge (\forall z \in S. f z = \text{exp}(g z))$ 
shows  $\exists g. \text{continuous\_on } S g \wedge (\forall z \in S. f z = (g z)^2)$ 
proof -
  obtain g where contg: continuous_on S g and geq:  $\bigwedge z. z \in S \Longrightarrow f z = \text{exp}(g z)$ 
using contf nz prev by metis
show ?thesis
proof (intro exI ballI conjI)
  show continuous_on S  $(\lambda z. \text{exp}(g z/2))$ 
    by (intro continuous_intros) (auto simp: contg)
  show  $\bigwedge z. z \in S \Longrightarrow f z = (\text{exp}(g z/2))^2$ 
    by (metis (no_types, lifting) divide_inverse exp_double_geq mult.left_commute
mult.right_neutral right_inverse zero_neq_numeral)
  qed
qed

lemma continuous_sqrt_impSimplyConnected:
  assumes connected S
  and prev:  $\bigwedge f::\text{complex} \Rightarrow \text{complex}. \llbracket \text{continuous\_on } S f; \forall z \in S. f z \neq 0 \rrbracket$ 
       $\Longrightarrow \exists g. \text{continuous\_on } S g \wedge (\forall z \in S. f z = (g z)^2)$ 
  shows simply_connected S
proof (clarsimp simp add: simply_connected_eq_holomorphic_sqrt [OF openS]
  <connected S>)
  fix f
  assume f holomorphic_on S and nz:  $\forall z \in S. f z \neq 0$ 
  then obtain g where contg: continuous_on S g and geq:  $\bigwedge z. z \in S \Longrightarrow f z =$ 
   $(g z)^2$ 
    by (metis holomorphic_on_imp_continuous_on prev)
  show  $\exists g. g \text{ holomorphic\_on } S \wedge (\forall z \in S. f z = (g z)^2)$ 
proof (intro exI ballI conjI)
  show g holomorphic_on S
proof (clarsimp simp add: holomorphic_on_open [OF openS])
  fix z
  assume  $z \in S$ 
  with nz geq have  $g z \neq 0$ 
    by auto
  obtain  $\delta$  where  $0 < \delta \wedge w. \llbracket w \in S; \text{dist } w z < \delta \rrbracket \Longrightarrow \text{dist } (g w) (g z) <$ 
  cmod  $(g z)$ 
    using contg [unfolded continuous_on_iff] by (metis <g z ≠ 0> <z ∈ S>
zero_less_norm_iff)
  then have  $\delta: \bigwedge w. \llbracket w \in S; w \in \text{ball } z \delta \rrbracket \Longrightarrow g w + g z \neq 0$ 
    apply (clarsimp simp: dist_norm)
    by (metis <g z ≠ 0> add_diff_cancel_left' diff_0_right norm_eq_zero
norm_increases_online norm_minus_commute norm_not_less_zero not_less_iff_gr_or_eq)
  have  $*$ :  $(\lambda x. (f x - f z) / (x - z) / (g x + g z)) -z \rightarrow \text{deriv } f z / (g z + g z)$ 
    apply (intro tendsto_intros)
    using SC_Chain.openS SC_Chain_axioms <f holomorphic_on S> <z ∈ S>

```

```

has_field_derivativeD holomorphic_derivI apply fastforce
  using ⟨z ∈ S⟩ contg continuous_on_eq_continuous_at isCont_def openS
apply blast
  by (simp add: ⟨g z ≠ 0⟩)
then have (g has_field_derivative deriv f z / (g z + g z)) (at z)
  unfolding has_field_derivative_iff
proof (rule Lim_transform_within_open)
  show open (ball z δ ∩ S)
    by (simp add: openS open_Int)
  show z ∈ ball z δ ∩ S
    using ⟨z ∈ S⟩ ⟨0 < δ⟩ by simp
  show  $\bigwedge x. [x \in \text{ball } z \ \delta \cap S; x \neq z]$ 
     $\implies (f x - f z) / (x - z) / (g x + g z) = (g x - g z) / (x - z)$ 
    using δ
    apply (simp add: geq ⟨z ∈ S⟩ divide_simps)
    apply (auto simp: algebra_simps power2_eq_square)
    done
  qed
then show  $\exists f'. (g \text{ has\_field\_derivative } f') (at z) ..$ 
qed
qed (use geq in auto)
qed
end

```

proposition

```

fixes S :: complex set
assumes open S
shows simply_connected_eq_continuous_log:
  simply_connected S  $\longleftrightarrow$ 
  connected S  $\wedge$ 
  ( $\forall f::\text{complex} \implies \text{complex}. \text{continuous\_on } S f \wedge (\forall z \in S. f z \neq 0)$ 
 $\longrightarrow (\exists g. \text{continuous\_on } S g \wedge (\forall z \in S. f z = \exp (g z)))$ ) (is ?log)
and simply_connected_eq_continuous_sqrt:
  simply_connected S  $\longleftrightarrow$ 
  connected S  $\wedge$ 
  ( $\forall f::\text{complex} \implies \text{complex}. \text{continuous\_on } S f \wedge (\forall z \in S. f z \neq 0)$ 
 $\longrightarrow (\exists g. \text{continuous\_on } S g \wedge (\forall z \in S. f z = (g z)^2))$ ) (is ?sqrt)
proof -
interpret SC_Chain
  using assms by (simp add: SC_Chain_def)
  have ?log  $\wedge$  ?sqrt
proof -
have *:  $[\alpha \implies \beta; \beta \implies \gamma; \gamma \implies \alpha]$ 
 $\implies (\alpha \longleftrightarrow \beta) \wedge (\alpha \longleftrightarrow \gamma)$  for α β γ
  by blast
show ?thesis
  apply (rule *)
  apply (simp add: local.continuous_log winding_number_zero)

```

```

    apply (simp add: continuous_sqrt)
    apply (simp add: continuous_sqrt_imp_simply_connected)
  done
qed
  then show ?log ?sqrt
    by safe
qed

```

8.5 More Borsukian results

```

lemma Borsukian_componentwise_eq:
  fixes S :: 'a::euclidean_space set
  assumes S: locally_connected S  $\vee$  compact S
  shows Borsukian S  $\longleftrightarrow$  ( $\forall C \in$  components S. Borsukian C)
proof -
  have *: ANR( $\{-0::\text{complex}\}$ )
    by (simp add: ANR_delete_open_Comp open_imp_ANR)
  show ?thesis
    using cohomotopically_trivial_on_components [OF assms *] by (auto simp:
Borsukian_alt)
qed

```

```

lemma Borsukian_componentwise:
  fixes S :: 'a::euclidean_space set
  assumes locally_connected S  $\vee$  compact S  $\wedge$  C. C  $\in$  components S  $\implies$  Borsukian
C
  shows Borsukian S
  by (metis Borsukian_componentwise_eq assms)

```

```

lemma simply_connected_eq_Borsukian:
  fixes S :: complex set
  shows open S  $\implies$  (simply_connected S  $\longleftrightarrow$  connected S  $\wedge$  Borsukian S)
by (auto simp: simply_connected_eq_continuous_log Borsukian_continuous_logarithm)

```

```

lemma Borsukian_eq_simply_connected:
  fixes S :: complex set
  shows open S  $\implies$  Borsukian S  $\longleftrightarrow$  ( $\forall C \in$  components S. simply_connected C)
by (meson Borsukian_componentwise_eq in_components_connected open_components
open_imp_locally_connected simply_connected_eq_Borsukian)

```

```

lemma Borsukian_separation_open_closed:
  fixes S :: complex set
  assumes S: open S  $\vee$  closed S and bounded S
  shows Borsukian S  $\longleftrightarrow$  connected( $-$  S)
  using S
proof
  assume open S
  show ?thesis
    unfolding Borsukian_eq_simply_connected [OF  $\langle$ open S $\rangle$ ]

```

```

  by (metis ‹open S› ‹bounded S› bounded_subset_in_components_maximal
  separation_by_component_eq open_componentsSimply_connected_iff_simple)
next
  assume closed S
  with ‹bounded S› show ?thesis
  by (simp add: Borsukian_separation_compact_eq_bounded_closed)
qed

```

8.6 Finally, the Riemann Mapping Theorem

theorem *Riemann_mapping_theorem*:

```

open S  $\wedge$  simply_connected S  $\longleftrightarrow$ 
S = {}  $\vee$  S = UNIV  $\vee$ 
( $\exists$  f g. f holomorphic_on S  $\wedge$  g holomorphic_on ball 0 1  $\wedge$ 
  ( $\forall$  z  $\in$  S. f z  $\in$  ball 0 1  $\wedge$  g(f z) = z)  $\wedge$ 
  ( $\forall$  z  $\in$  ball 0 1. g z  $\in$  S  $\wedge$  f(g z) = z))
(is _ = ?rhs)

```

proof –

have simply_connected S \longleftrightarrow ?rhs if open S

by (simp add: simply_connected_eq_biholomorphic_to_disc that)

moreover have open S if ?rhs

proof –

```

{ fix f g
  assume g: g holomorphic_on ball 0 1  $\forall$  z  $\in$  ball 0 1. g z  $\in$  S  $\wedge$  f (g z) = z
  and  $\forall$  z  $\in$  S. cmod (f z) < 1  $\wedge$  g (f z) = z
  then have S = g ‘ (ball 0 1)
  by (force simp:)
  then have open S
  by (metis open_ball g inj_on_def open_mapping_thm3)
}

```

with that show open S by auto

qed

ultimately show ?thesis by metis

qed

8.7 Applications to Winding Numbers

lemma *simply_connected_inside_simple_path*:

fixes p :: real \Rightarrow complex

shows simple_path p \implies simply_connected(inside(path_image p))

using Jordan_inside_outside_connected_simple_path_image_inside_simple_curve_imp_closed
 simply_connected_eq_frontier_properties

by fastforce

lemma *simply_connected_Int*:

fixes S :: complex set

assumes open S open T simply_connected S simply_connected T connected (S
 \cap T)

shows simply_connected (S \cap T)

using assms by (force simp: simply_connected_eq_winding_number_zero open_Int)

8.8 The winding number defines a continuous logarithm for the path itself

lemma *winding_number_as_continuous_log*:

assumes *path p* and $\zeta: \zeta \notin \text{path_image } p$

obtains *q* where *path q*

$$\begin{aligned} & \text{pathfinish } q - \text{pathstart } q = 2 * \text{of_real } \pi i * i * \text{winding_number } p \zeta \\ & \wedge t. t \in \{0..1\} \implies p \ t = \zeta + \exp(q \ t) \end{aligned}$$

proof –

let $?q = \lambda t. 2 * \text{of_real } \pi i * i * \text{winding_number}(\text{subpath } 0 \ t \ p) \ \zeta + \text{Ln}(\text{pathstart } p - \zeta)$

show *?thesis*

proof

have $*$: *continuous* (at *t* within $\{0..1\}$) ($\lambda x. \text{winding_number}(\text{subpath } 0 \ x \ p) \ \zeta$)

if $t: t \in \{0..1\}$ **for** *t*

proof –

let $?B = \text{ball}(p \ t) (\text{norm}(p \ t - \zeta))$

have $p \ t \neq \zeta$

using *path_image_def that* ζ **by** *blast*

then have *simply_connected ?B*

by (*simp add: convex_imp_simply_connected*)

then have $\wedge f::\text{complex} \Rightarrow \text{complex. continuous_on } ?B \ f \wedge (\forall \zeta \in ?B. f \ \zeta \neq 0)$

$$\longrightarrow (\exists g. \text{continuous_on } ?B \ g \wedge (\forall \zeta \in ?B. f \ \zeta = \exp(g \ \zeta)))$$

by (*simp add: simply_connected_eq_continuous_log*)

moreover have *continuous_on ?B* ($\lambda w. w - \zeta$)

by (*intro continuous_intros*)

moreover have ($\forall z \in ?B. z - \zeta \neq 0$)

by (*auto simp: dist_norm*)

ultimately obtain *g* where *contg: continuous_on ?B g*

and *geg: $\wedge z. z \in ?B \implies z - \zeta = \exp(g \ z)$* **by** *blast*

obtain *d* where $0 < d$ **and** *d*:

$$\wedge x. \llbracket x \in \{0..1\}; \text{dist } x \ t < d \rrbracket \implies \text{dist}(p \ x) (p \ t) < \text{cmod}(p \ t - \zeta)$$

using $\langle \text{path } p \rangle \ t$ **unfolding** *path_def continuous_on_iff*

by (*metis $\langle p \ t \neq \zeta \rangle$ right_minus_eq zero_less_norm_iff*)

have ($(\lambda x. \text{winding_number}(\lambda w. \text{subpath } 0 \ x \ p \ w - \zeta) \ 0 - \text{winding_number}(\lambda w. \text{subpath } 0 \ t \ p \ w - \zeta) \ 0) \longrightarrow 0$)
(at *t* within $\{0..1\}$)

proof (*rule Lim_transform_within [OF _ $\langle d > 0 \rangle$]*)

have *continuous* (at *t* within $\{0..1\}$) (*g o p*)

proof (*rule continuous_within_compose*)

show *continuous* (at *t* within $\{0..1\}$) *p*

using $\langle \text{path } p \rangle$ *continuous_on_eq_continuous_within path_def that* **by**

blast

show *continuous* (at $(p \ t)$ within $p \ ^\{0..1\}$) *g*

by (*metis (no_types, lifting) open_ball UNIV_I $\langle p \ t \neq \zeta \rangle$ centre_in_ball*

contg continuous_on_eq_continuous_at continuous_within_topological_right_minus_eq zero_less_norm_iff)

qed

```

      with LIM_zero have (( $\lambda u. (g (subpath\ t\ u\ p\ 1) - g (subpath\ t\ u\ p\ 0))$ )
 $\longrightarrow 0$ ) (at t within {0..1})
      by (auto simp: subpath_def continuous_within_o_def)
      then show (( $\lambda u. (g (subpath\ t\ u\ p\ 1) - g (subpath\ t\ u\ p\ 0)) / (2 * of\_real$ 
 $pi * i)$ )  $\longrightarrow 0$ )
        (at t within {0..1})
        by (simp add: tendsto_divide_zero)
      show (g (subpath t u p 1) - g (subpath t u p 0)) / (2 * of_real pi * i) =
        winding_number ( $\lambda w. subpath\ 0\ u\ p\ w - \zeta$ ) 0 - winding_number ( $\lambda w.$ 
 $subpath\ 0\ t\ p\ w - \zeta$ ) 0
      if  $u \in \{0..1\}$   $0 < dist\ u\ t$   $dist\ u\ t < d$  for u
      proof -
        have closed_segment t u  $\subseteq \{0..1\}$ 
          using closed_segment_eq_real_ivl t that by auto
        then have piB: path_image(subpath t u p)  $\subseteq ?B$ 
          apply (clarsimp simp add: path_image_subpath_gen)
          by (metis subsetD le_less_trans ‹dist u t < d› d dist_commute
 $dist\_in\_closed\_segment$ )
        have *: path (g  $\circ$  subpath t u p)
          apply (rule path_continuous_image)
          using ‹path p› t that apply auto[1]
          using piB contg continuous_on_subset by blast
        have (g (subpath t u p 1) - g (subpath t u p 0)) / (2 * of_real pi * i)
          = winding_number (exp  $\circ$  g  $\circ$  subpath t u p) 0
          using winding_number_compose_exp [OF *]
          by (simp add: pathfinish_def pathstart_def o_assoc)
        also have ... = winding_number ( $\lambda w. subpath\ t\ u\ p\ w - \zeta$ ) 0
        proof (rule winding_number_cong)
          have exp(g y) = y -  $\zeta$  if  $y \in (subpath\ t\ u\ p) \text{ ‘}\{0..1\}$  for y
            by (metis that geq_path_image_def piB subset_eq)
          then show  $\bigwedge x. \llbracket 0 \leq x; x \leq 1 \rrbracket \implies (exp \circ g \circ subpath\ t\ u\ p)\ x = subpath$ 
 $t\ u\ p\ x - \zeta$ 
            by auto
        qed
        qed
        also have ... = winding_number ( $\lambda w. subpath\ 0\ u\ p\ w - \zeta$ ) 0 -
          winding_number ( $\lambda w. subpath\ 0\ t\ p\ w - \zeta$ ) 0
          apply (simp add: winding_number_offset [symmetric])
          using winding_number_subpath_combine [OF ‹path p›  $\zeta$ , of 0 t u] ‹t  $\in$ 
 $\{0..1\}$ › ‹u  $\in \{0..1\}$ ›
          by (simp add: add commute eq_diff_eq)
        finally show ?thesis .
      qed
    qed
  then show ?thesis
    by (subst winding_number_offset) (simp add: continuous_within LIM_zero_iff)
  qed
show path ?q
  unfolding path_def
  by (intro continuous_intros) (simp add: continuous_on_eq_continuous_within

```

```

*)

  have  $\zeta \neq p \ 0$ 
    by (metis  $\zeta$  pathstart_def pathstart_in_path_image)
  then show pathfinish ?q - pathstart ?q = 2 * of_real pi * i * winding_number
  p  $\zeta$ 
    by (simp add: pathfinish_def pathstart_def)
  show p t =  $\zeta + \exp (?q \ t)$  if  $t \in \{0..1\}$  for t
  proof -
    have path (subpath 0 t p)
      using ⟨path p⟩ that by auto
    moreover
    have  $\zeta \notin \text{path\_image (subpath 0 t p)}$ 
    using  $\zeta$  [unfolded path_image_def] that by (auto simp: path_image_subpath)
    ultimately show ?thesis
      using winding_number_exp_2pi [of subpath 0 t p  $\zeta$ ] ⟨ $\zeta \neq p \ 0$ ⟩
      by (auto simp: exp_add algebra_simps pathfinish_def pathstart_def sub-
  path_def)
    qed
  qed
  qed

```

8.9 Winding number equality is the same as path/loop homotopy in $\mathbb{C} - 0$

```

lemma winding_number_homotopic_loops_null_eq:
  assumes path p and  $\zeta: \zeta \notin \text{path\_image } p$ 
  shows winding_number p  $\zeta = 0 \iff (\exists a. \text{homotopic\_loops } (-\{\zeta\}) \ p \ (\lambda t. \ a))$ 
  (is ?lhs = ?rhs)
proof
  assume [simp]: ?lhs
  obtain q where path q
    and qeq: pathfinish q - pathstart q = 2 * of_real pi * i * winding_number
  p  $\zeta$ 
    and peq:  $\bigwedge t. t \in \{0..1\} \implies p \ t = \zeta + \exp(q \ t)$ 
    using winding_number_as_continuous_log [OF assms] by blast
  have *: homotopic_with_canon ( $\lambda r. \text{pathfinish } r = \text{pathstart } r$ )
    {0..1} ( $-\{\zeta\}$ ) (( $\lambda w. \zeta + \exp w$ )  $\circ$  q) (( $\lambda w. \zeta + \exp w$ )  $\circ$  ( $\lambda t. \ 0$ ))
  proof (rule homotopic_with_compose_continuous_left)
    show homotopic_with_canon ( $\lambda f. \text{pathfinish } ((\lambda w. \zeta + \exp w) \circ f) = \text{pathstart}$ 
    (( $\lambda w. \zeta + \exp w$ )  $\circ$  f))
      {0..1} UNIV q ( $\lambda t. \ 0$ )
    proof (rule homotopic_with_mono, simp_all add: pathfinish_def pathstart_def)
      have homotopic_loops UNIV q ( $\lambda t. \ 0$ )
      by (rule homotopic_loops_linear) (use qeq ⟨path q⟩ in ⟨auto simp: path_defs⟩)
      then have homotopic_with ( $\lambda r. \ r \ 1 = r \ 0$ ) (top_of_set {0..1}) euclidean q
      ( $\lambda t. \ 0$ )
      by (simp add: homotopic_loops_def pathfinish_def pathstart_def)
      then show homotopic_with ( $\lambda h. \ \exp (h \ 1) = \exp (h \ 0)$ ) (top_of_set {0..1})

```



```

euclidean q ( $\lambda t. 0$ )
  by (rule homotopic_with_mono) simp
qed
show continuous_on UNIV ( $\lambda w. \zeta + \exp w$ )
  by (rule continuous_intros)+
show ( $\lambda w. \zeta + \exp w$ )  $\in$  UNIV  $\rightarrow -\{\zeta\}$ 
  by auto
qed
then have homotopic_with_canon ( $\lambda r. \text{pathfinish } r = \text{pathstart } r$ )  $\{0..1\} (-\{\zeta\})$ 
p ( $\lambda x. \zeta + 1$ )
  by (rule homotopic_with_eq) (auto simp: o_def peq pathfinish_def pathstart_def)
then have homotopic_loops ( $-\{\zeta\}$ ) p ( $\lambda t. \zeta + 1$ )
  by (simp add: homotopic_loops_def)
then show ?rhs ..
next
assume ?rhs
then obtain a where homotopic_loops ( $-\{\zeta\}$ ) p ( $\lambda t. a$ ) ..
then have winding_number p  $\zeta =$  winding_number ( $\lambda t. a$ )  $\zeta$   $a \neq \zeta$ 
  using winding_number_homotopic_loops homotopic_loops_imp_subset by
(force simp:)+
moreover have winding_number ( $\lambda t. a$ )  $\zeta = 0$ 
  by (metis winding_number_zero_const  $\langle a \neq \zeta \rangle$ )
ultimately show ?lhs by metis
qed

lemma winding_number_homotopic_paths_null_explicit_eq:
  assumes path p and  $\zeta: \zeta \notin \text{path\_image } p$ 
  shows winding_number p  $\zeta = 0 \iff$  homotopic_paths ( $-\{\zeta\}$ ) p (linepath
(pathstart p) (pathstart p))
    (is ?lhs = ?rhs)
proof
  assume ?lhs
  then show ?rhs
    using homotopic_loops_imp_homotopic_paths_null
    by (force simp add: linepath_refl winding_number_homotopic_loops_null_eq
[OF assms])
next
  assume ?rhs
  then show ?lhs
    by (metis  $\zeta$  pathstart_in_path_image winding_number_homotopic_paths wind-
ing_number_trivial)
qed

lemma winding_number_homotopic_paths_null_eq:
  assumes path p and  $\zeta: \zeta \notin \text{path\_image } p$ 
  shows winding_number p  $\zeta = 0 \iff$  ( $\exists a. \text{homotopic\_paths } (-\{\zeta\}) p (\lambda t. a)$ )
    (is ?lhs = ?rhs)
proof
  assume ?lhs

```

```

then show ?rhs
  by (auto simp: winding_number_homotopic_paths_null_explicit_eq [OF assms]
linepath_refl)
next
  assume ?rhs
  then show ?lhs
    by (metis  $\zeta$  homotopic_paths_imp_pathfinish pathfinish_def pathfinish_in_path_image
winding_number_homotopic_paths winding_number_zero_const)
qed

```

proposition *winding_number_homotopic_paths_eq*:

```

assumes path  $p$  and  $\zeta p$ :  $\zeta \notin \text{path\_image } p$ 
  and path  $q$  and  $\zeta q$ :  $\zeta \notin \text{path\_image } q$ 
  and  $qp$ : pathstart  $q = \text{pathstart } p$  pathfinish  $q = \text{pathfinish } p$ 
  shows winding_number  $p \ \zeta = \text{winding\_number } q \ \zeta \iff \text{homotopic\_paths}$ 
( $-\{\zeta\}$ )  $p \ q$ 
  (is ?lhs = ?rhs)

```

proof

```

assume ?lhs
then have winding_number ( $p \ \text{+++ } \text{reversepath } q$ )  $\zeta = 0$ 
  using assms by (simp add: winding_number_join winding_number_reversepath)
moreover
have path ( $p \ \text{+++ } \text{reversepath } q$ )  $\zeta \notin \text{path\_image } (p \ \text{+++ } \text{reversepath } q)$ 
  using assms by (auto simp: not_in_path_image_join)
ultimately obtain  $a$  where homotopic_paths ( $-\{\zeta\}$ ) ( $p \ \text{+++ } \text{reversepath } q$ )
(linepath  $a \ a$ )
  using winding_number_homotopic_paths_null_explicit_eq by blast
then show ?rhs
  using homotopic_paths_imp_pathstart assms
  by (fastforce simp add: dest: homotopic_paths_imp_homotopic_loops homo-
topic_paths_loop_parts)
qed (simp add: winding_number_homotopic_paths)

```

lemma *winding_number_homotopic_loops_eq*:

```

assumes path  $p$  and  $\zeta p$ :  $\zeta \notin \text{path\_image } p$ 
  and path  $q$  and  $\zeta q$ :  $\zeta \notin \text{path\_image } q$ 
  and loops: pathfinish  $p = \text{pathstart } p$  pathfinish  $q = \text{pathstart } q$ 
  shows winding_number  $p \ \zeta = \text{winding\_number } q \ \zeta \iff \text{homotopic\_loops}$ 
( $-\{\zeta\}$ )  $p \ q$ 
  (is ?lhs = ?rhs)

```

proof

```

assume  $L$ : ?lhs
have pathstart  $p \neq \zeta$  pathstart  $q \neq \zeta$ 
  using  $\zeta p \ \zeta q$  by blast+
moreover have path_connected ( $-\{\zeta\}$ )
  by (simp add: path_connected_punctured_universe)
ultimately obtain  $r$  where path  $r$  and rim: path_image  $r \subseteq -\{\zeta\}$ 
  and pas: pathstart  $r = \text{pathstart } p$  and paf: pathfinish  $r =$ 
pathstart  $q$ 

```

```

    by (auto simp: path_connected_def)
  then have pathstart  $r \neq \zeta$  by blast
  have homotopic_loops ( $-\{\zeta\}$ )  $p$  ( $r$   $+++$   $q$   $+++$  reversepath  $r$ )
  proof (rule homotopic_paths_imp_homotopic_loops)
    have path ( $r$   $+++$   $q$   $+++$  reversepath  $r$ )
      by (simp add: <path  $r$ > <path  $q$ > loops paf)
    moreover have  $\zeta \notin \text{path\_image } (r$   $+++$   $q$   $+++$  reversepath  $r$ )
      by (metis  $\zeta q$  not_in_path_image_join path_image_reversepath rim subset_Cmpl_singleton)
    moreover have homotopic_loops ( $-\{\zeta\}$ ) ( $r$   $+++$   $q$   $+++$  reversepath  $r$ )  $q$ 
      using <path  $q$ > <path  $r$ >  $\zeta q$  homotopic_loops_conjugate loops(2) paf rim by blast
    ultimately show homotopic_paths ( $-\{\zeta\}$ )  $p$  ( $r$   $+++$   $q$   $+++$  reversepath  $r$ )
      using loops pathfinish_join pathfinish_reversepath pathstart_join
      by (metis  $L \zeta p$  <path  $p$ > pas winding_number_homotopic_loops winding_number_homotopic_paths_eq)
    qed (use loops pas in auto)
    moreover have homotopic_loops ( $-\{\zeta\}$ ) ( $r$   $+++$   $q$   $+++$  reversepath  $r$ )  $q$ 
      using rim  $\zeta q$  by (auto simp: homotopic_loops_conjugate paf <path  $q$ > <path  $r$ > loops)
    ultimately show ?rhs
      using homotopic_loops_trans by metis
  qed (simp add: winding_number_homotopic_loops)

end
theory Meromorphic
  imports Laurent_Convergence Riemann_Mapping
begin

lemma analytic_at_cong:
  assumes eventually ( $\lambda x. f x = g x$ ) (nhds  $x$ )  $x = y$ 
  shows  $f$  analytic_on  $\{x\} \iff g$  analytic_on  $\{y\}$ 
proof -
  have  $g$  analytic_on  $\{x\}$  if  $f$  analytic_on  $\{x\}$  eventually ( $\lambda x. f x = g x$ ) (nhds  $x$ )
for  $f g$ 
proof -
  have ( $\lambda y. f (x + y)$ ) has_fps_expansion fps_expansion  $f x$ 
    by (rule analytic_at_imp_has_fps_expansion) fact
  also have ?this  $\iff$  ( $\lambda y. g (x + y)$ ) has_fps_expansion fps_expansion  $f x$ 
    using that by (intro has_fps_expansion_cong refl) (auto simp: nhds_to_0' eventually_filtermap)
  finally show ?thesis
    by (rule has_fps_expansion_imp_analytic)
  qed
from this[of  $f g$ ] this[of  $g f$ ] show ?thesis using assms
  by (auto simp: eq_commute)
qed

definition remove_sings :: (complex  $\Rightarrow$  complex)  $\Rightarrow$  complex  $\Rightarrow$  complex where
  remove_sings  $f z =$  (if  $\exists c. f -z \rightarrow c$  then  $\text{Lim } (at z) f$  else 0)

```

```

lemma remove_sings_eqI [intro]:
  assumes  $f -z \rightarrow c$ 
  shows  $\text{remove\_sings } f z = c$ 
  using assms unfolding remove_sings_def by (auto simp: tendsto_Lim)

lemma remove_sings_at_analytic [simp]:
  assumes  $f \text{ analytic\_on } \{z\}$ 
  shows  $\text{remove\_sings } f z = f z$ 
  using assms by (intro remove_sings_eqI) (simp add: analytic_at_imp_isCont isContD)

lemma remove_sings_at_pole [simp]:
  assumes  $\text{is\_pole } f z$ 
  shows  $\text{remove\_sings } f z = 0$ 
  using assms unfolding remove_sings_def is_pole_def
  by (meson at_neq_bot not_tendsto_and_filterlim_at_infinity)

lemma eventually_remove_sings_eq_at:
  assumes  $\text{isolated\_singularity\_at } f z$ 
  shows  $\text{eventually } (\lambda w. \text{remove\_sings } f w = f w) \text{ (at } z)$ 
proof -
  from assms obtain  $r$  where  $r: r > 0$   $f \text{ analytic\_on ball } z r - \{z\}$ 
  by (auto simp: isolated_singularity_at_def)
  hence  $*$ :  $f \text{ analytic\_on } \{w\}$  if  $w \in \text{ball } z r - \{z\}$  for  $w$ 
  using  $r$  that by (auto intro: analytic_on_subset)
  have  $\text{eventually } (\lambda w. w \in \text{ball } z r - \{z\}) \text{ (at } z)$ 
  using  $r$  by (intro eventually_at_in_open) auto
  thus ?thesis
  by eventually_elim (auto simp: remove_sings_at_analytic *)
qed

lemma eventually_remove_sings_eq_nhds:
  assumes  $f \text{ analytic\_on } \{z\}$ 
  shows  $\text{eventually } (\lambda w. \text{remove\_sings } f w = f w) \text{ (nhds } z)$ 
proof -
  from assms obtain  $A$  where  $A: \text{open } A z \in A$   $f \text{ holomorphic\_on } A$ 
  by (auto simp: analytic_at)
  have  $\text{eventually } (\lambda z. z \in A) \text{ (nhds } z)$ 
  by (intro eventually_nhds_in_open A)
  thus ?thesis
proof eventually_elim
  case (elim w)
  from elim have  $f \text{ analytic\_on } \{w\}$ 
  using  $A$  analytic_at by blast
  thus ?case by auto
qed
qed

```

```

lemma remove_sings_compose:
  assumes filtermap g (at z) = at z'
  shows remove_sings (f ∘ g) z = remove_sings f z'
proof (cases  $\exists c. f -z' \rightarrow c$ )
  case True
  then obtain c where  $c: f -z' \rightarrow c$ 
    by auto
  from c have remove_sings f z' = c
    by blast
  moreover from c have remove_sings (f ∘ g) z = c
    using c by (intro remove_sings_eqI) (auto simp: filterlim_def filtermap_compose
assms)
  ultimately show ?thesis
    by simp
next
  case False
  hence  $\neg(\exists c. (f \circ g) -z \rightarrow c)$ 
    by (auto simp: filterlim_def filtermap_compose assms)
  with False show ?thesis
    by (auto simp: remove_sings_def)
qed

```

```

lemma remove_sings_cong:
  assumes eventually  $(\lambda x. f x = g x)$  (at z)  $z = z'$ 
  shows remove_sings f z = remove_sings g z'
proof (cases  $\exists c. f -z \rightarrow c$ )
  case True
  then obtain c where  $c: f -z \rightarrow c$  by blast
  hence remove_sings f z = c
    by blast
  moreover have  $f -z \rightarrow c \iff g -z' \rightarrow c$ 
    using assms by (intro filterlim_cong refl) auto
  with c have remove_sings g z' = c
    by (intro remove_sings_eqI) auto
  ultimately show ?thesis
    by simp
next
  case False
  have  $f -z \rightarrow c \iff g -z' \rightarrow c$  for c
    using assms by (intro filterlim_cong) auto
  with False show ?thesis
    by (auto simp: remove_sings_def)
qed

```

```

lemma deriv_remove_sings_at_analytic [simp]:
  assumes f analytic_on  $\{z\}$ 
  shows deriv (remove_sings f) z = deriv f z
  apply (rule deriv_cong_ev)

```

apply (rule eventually_remove_sings_eq_nhds)
using *assms* **by** *auto*

lemma *isolated_singularity_at_remove_sings* [*simp*, *intro*]:
assumes *isolated_singularity_at* *f z*
shows *isolated_singularity_at* (remove_sings *f*) *z*
using *isolated_singularity_at_cong*[*OF* eventually_remove_sings_eq_at[*OF* *assms*]
refl] *assms*
by *simp*

lemma *not_essential_remove_sings_iff* [*simp*]:
assumes *isolated_singularity_at* *f z*
shows *not_essential* (remove_sings *f*) *z* \longleftrightarrow *not_essential* *f z*
using *not_essential_cong*[*OF* eventually_remove_sings_eq_at[*OF* *assms*(1)]
refl]
by *simp*

lemma *not_essential_remove_sings* [*intro*]:
assumes *isolated_singularity_at* *f z* *not_essential* *f z*
shows *not_essential* (remove_sings *f*) *z*
by (*subst not_essential_remove_sings_iff*) (use *assms* **in** *auto*)

lemma
assumes *isolated_singularity_at* *f z*
shows *is_pole_remove_sings_iff* [*simp*]:
is_pole (remove_sings *f*) *z* \longleftrightarrow *is_pole* *f z*
and *zorder_remove_sings* [*simp*]:
zorder (remove_sings *f*) *z* = *zorder* *f z*
and *zor_poly_remove_sings* [*simp*]:
zor_poly (remove_sings *f*) *z* = *zor_poly* *f z*
and *has_laurent_expansion_remove_sings_iff* [*simp*]:
 $(\lambda w. \text{remove_sings } f (z + w)) \text{ has_laurent_expansion } F \longleftrightarrow$
 $(\lambda w. f (z + w)) \text{ has_laurent_expansion } F$
and *tendsto_remove_sings_iff* [*simp*]:
remove_sings *f* $-z \rightarrow c \longleftrightarrow$ *f* $-z \rightarrow c$
by (*intro is_pole_cong eventually_remove_sings_eq_at refl zorder_cong*
zor_poly_cong has_laurent_expansion_cong' tendsto_cong assms)+

lemma *get_all_poles_from_remove_sings*:
fixes *f*:: *complex* \Rightarrow *complex*
defines *ff* \equiv *remove_sings* *f*
assumes *f_holo*:*f* *holomorphic_on* *s - pts* **and** *finite pts*
pts \subseteq *s* *open* *s* **and** *not_ess*: $\forall x \in \text{pts}. \text{not_essential } f x$
obtains *pts'* **where**
pts' \subseteq *pts* *finite* *pts'* *ff* *holomorphic_on* *s - pts'* $\forall x \in \text{pts}'. \text{is_pole } ff x$
proof –
define *pts'* **where** *pts'* = {*x* \in *pts*. *is_pole* *f x*}

have *pts'* \subseteq *pts* **unfolding** *pts'_def* **by** *auto*

```

then have finite pts' using ‹finite pts›
  using rev_finite_subset by blast
then have open (s - pts') using ‹open s›
  by (simp add: finite_imp_closed open_Diff)

have isolated:isolated_singularity_at f z if z ∈ pts for z
proof (rule isolated_singularity_at_holomorphic)
  show f holomorphic_on (s - (pts - {z})) - {z}
    by (metis Diff_insert f_holo insert_Diff that)
  show open (s - (pts - {z}))
    by (meson assms(3) assms(5) finite_Diff finite_imp_closed open_Diff)
  show z ∈ s - (pts - {z})
    using assms(4) that by auto
qed

have ff holomorphic_on s - pts'
proof (rule no_isolated_singularity')
  show (ff  $\longrightarrow$  ff z) (at z within s - pts') if z ∈ pts - pts' for z
  proof -
    have at z within s - pts' = at z
      apply (rule at_within_open)
      using ‹open (s - pts')› that ‹pts ⊆ s› by auto
    moreover have ff -z → ff z
      unfolding ff_def
    proof (subst tendsto_remove_sings_iff)
      show isolated_singularity_at f z
        apply (rule isolated)
        using that by auto
      have not_essential f z
        using not_ess that by auto
      moreover have  $\neg$ is_pole f z
        using that unfolding pts'_def by auto
      ultimately have  $\exists c. f -z \rightarrow c$ 
        unfolding not_essential_def by auto
      then show f -z → remove_sings f z
        using remove_sings_eqI by blast
    qed
  ultimately show ?thesis by auto
qed

have ff holomorphic_on s - pts
  using f_holo
proof (elim holomorphic_transform)
  fix x assume x ∈ s - pts
  then have f analytic_on {x}
    using assms(3) assms(5) f_holo
    by (meson finite_imp_closed
      holomorphic_on_imp_analytic_at open_Diff)
  from remove_sings_at_analytic[OF this]
  show f x = ff x unfolding ff_def by auto

```

```

qed
then show ff holomorphic_on s - pts' - (pts - pts')
  apply (elim holomorphic_on_subset)
  by blast
show open (s - pts')
  by (simp add: ⟨open (s - pts')⟩)
show finite (pts - pts')
  by (simp add: assms(3))
qed
moreover have  $\forall x \in pts'. is\_pole\ ff\ x$ 
  unfolding pts'_def
  using ff_def is_pole_remove_sings_iff isolated by blast
moreover note ⟨pts'  $\subseteq$  pts⟩ ⟨finite pts'⟩
ultimately show ?thesis using that by auto
qed

```

```

lemma remove_sings_eq_0_iff:
  assumes not_essential f w
  shows remove_sings f w = 0  $\longleftrightarrow$  is_pole f w  $\vee$  f -w  $\rightarrow$  0
proof (cases is_pole f w)
  case True
  then show ?thesis by simp
next
  case False
  then obtain c where c: f -w  $\rightarrow$  c
  using ⟨not_essential f w⟩ unfolding not_essential_def by auto
  then show ?thesis
  using False remove_sings_eqI by auto
qed

```

```

definition meromorphic_on:: [complex  $\Rightarrow$  complex, complex set, complex set]  $\Rightarrow$ 
bool
  ( $\_$  (meromorphic'_on)  $\_$   $\_$  [50,50,50] 50) where
  f meromorphic_on D pts  $\equiv$ 
  open D  $\wedge$  pts  $\subseteq$  D  $\wedge$  ( $\forall z \in pts. isolated\_singularity\_at\ f\ z \wedge not\_essential\ f\ z$ )
 $\wedge$ 
  ( $\forall z \in D. \neg(z\ islimpt\ pts)$ )  $\wedge$  (f holomorphic_on D - pts)

```

```

lemma meromorphic_imp_holomorphic: f meromorphic_on D pts  $\implies$  f holomor-
phic_on (D - pts)
  unfolding meromorphic_on_def by auto

```

```

lemma meromorphic_imp_closedin_pts:
  assumes f meromorphic_on D pts
  shows closedin (top_of_set D) pts
  by (meson assms closedin_limpt meromorphic_on_def)

```

```

lemma meromorphic_imp_open_diff':
  assumes f meromorphic_on D pts pts'  $\subseteq$  pts

```



```

shows open (D - pts')
proof -
  have D - pts' = D - closure pts'
  proof safe
    fix x assume x: x ∈ D x ∈ closure pts' x ∉ pts'
    hence x islimpt pts'
    by (subst islimpt_in_closure) auto
    hence x islimpt pts
    by (rule islimpt_subset) fact
  with assms x show False
  by (auto simp: meromorphic_on_def)
qed (use closure_subset in auto)
then show ?thesis
  using assms meromorphic_on_def by auto
qed

```

```

lemma meromorphic_imp_open_diff: f meromorphic_on D pts ⇒ open (D -
pts)
  by (erule meromorphic_imp_open_diff') auto

```

```

lemma meromorphic_pole_subset:
  assumes merf: f meromorphic_on D pts
  shows {x∈D. is_pole f x} ⊆ pts
  by (smt (verit) Diff_iff assms mem_Collect_eq meromorphic_imp_open_diff
meromorphic_on_def not_is_pole_holomorphic subsetI)

```

named_theorems meromorphic_intros

```

lemma meromorphic_on_subset:
  assumes f meromorphic_on A pts open B B ⊆ A pts' = pts ∩ B
  shows f meromorphic_on B pts'
  unfolding meromorphic_on_def
proof (intro ballI conjI)
  fix z assume z ∈ B
  show ¬z islimpt pts'
  proof
    assume z islimpt pts'
    hence z islimpt pts
    by (rule islimpt_subset) (use ⟨pts' = _⟩ in auto)
  thus False using ⟨z ∈ B⟩ ⟨B ⊆ A⟩ assms(1)
  by (auto simp: meromorphic_on_def)
qed
qed (use assms in ⟨auto simp: meromorphic_on_def⟩)

```

```

lemma meromorphic_on_superset_pts:
  assumes f meromorphic_on A pts pts ⊆ pts' pts' ⊆ A ∀x∈A. ¬x islimpt pts'
  shows f meromorphic_on A pts'
  unfolding meromorphic_on_def
proof (intro conjI ballI impI)

```

```

fix  $z$  assume  $z \in pts'$ 
from  $assms(1)$  have  $holo: f \text{ holomorphic\_on } A - pts$  and  $open\ A$ 
  unfolding  $meromorphic\_on\_def$  by  $blast+$ 
have  $open\ (A - pts)$ 
  by ( $intro\ meromorphic\_imp\_open\_diff[OF\ assms(1)]$ )

show  $isolated\_singularity\_at\ f\ z$ 
proof ( $cases\ z \in pts$ )
  case  $False$ 
  thus  $?thesis$ 
  using  $\langle open\ (A - pts) \rangle\ assms\ \langle z \in pts' \rangle$ 
  by ( $intro\ isolated\_singularity\_at\_holomorphic[of\ _\ A - pts]\ holomorphic\_on\_subset[OF\ holo]$ )
   $auto$ 
qed ( $use\ assms\ in\ \langle auto\ simp: meromorphic\_on\_def \rangle$ )

show  $not\_essential\ f\ z$ 
proof ( $cases\ z \in pts$ )
  case  $False$ 
  thus  $?thesis$ 
  using  $\langle open\ (A - pts) \rangle\ assms\ \langle z \in pts' \rangle$ 
  by ( $intro\ not\_essential\_holomorphic[of\ _\ A - pts]\ holomorphic\_on\_subset[OF\ holo]$ )
   $auto$ 
qed ( $use\ assms\ in\ \langle auto\ simp: meromorphic\_on\_def \rangle$ )
qed ( $use\ assms\ in\ \langle auto\ simp: meromorphic\_on\_def \rangle$ )

lemma  $meromorphic\_on\_no\_singularities: f \text{ meromorphic\_on } A \{\} \longleftrightarrow f \text{ holomorphic\_on } A \wedge open\ A$ 
  by ( $auto\ simp: meromorphic\_on\_def$ )

lemma  $holomorphic\_on\_imp\_meromorphic\_on:$ 
   $f \text{ holomorphic\_on } A \implies pts \subseteq A \implies open\ A \implies \forall x \in A. \neg x \text{ islimpt } pts \implies f \text{ meromorphic\_on } A\ pts$ 
  by ( $rule\ meromorphic\_on\_superset\_pts[where\ pts = \{\}]\$ )
  ( $auto\ simp: meromorphic\_on\_no\_singularities$ )

lemma  $meromorphic\_on\_const$  [ $meromorphic\_intros$ ]:
  assumes  $open\ A\ \forall x \in A. \neg x \text{ islimpt } pts\ pts \subseteq A$ 
  shows  $(\lambda_. c) \text{ meromorphic\_on } A\ pts$ 
  by ( $rule\ holomorphic\_on\_imp\_meromorphic\_on$ ) ( $use\ assms\ in\ auto$ )

lemma  $meromorphic\_on\_ident$  [ $meromorphic\_intros$ ]:
  assumes  $open\ A\ \forall x \in A. \neg x \text{ islimpt } pts\ pts \subseteq A$ 
  shows  $(\lambda x. x) \text{ meromorphic\_on } A\ pts$ 
  by ( $rule\ holomorphic\_on\_imp\_meromorphic\_on$ ) ( $use\ assms\ in\ auto$ )

lemma  $meromorphic\_on\_id$  [ $meromorphic\_intros$ ]:
  assumes  $open\ A\ \forall x \in A. \neg x \text{ islimpt } pts\ pts \subseteq A$ 

```

shows id meromorphic_on A pts
using meromorphic_on_ident assms **unfolding** id_def .

lemma not_essential_add [singularity_intros]:
assumes f_ness : not_essential f z **and** g_ness : not_essential g z
assumes f_iso : isolated_singularity_at f z **and** g_iso : isolated_singularity_at g z
shows not_essential $(\lambda w. f w + g w)$ z
proof –
have $(\lambda w. f (z + w) + g (z + w))$ has_laurent_expansion laurent_expansion f z + laurent_expansion g z
by (intro not_essential_has_laurent_expansion laurent_expansion_intros assms)
hence not_essential $(\lambda w. f (z + w) + g (z + w))$ 0
using has_laurent_expansion_not_essential_0 **by** blast
thus ?thesis
by (simp add: not_essential_shift_0)
qed

lemma meromorphic_on_uminus [meromorphic_intros]:
assumes f meromorphic_on A pts
shows $(\lambda z. -f z)$ meromorphic_on A pts
unfolding meromorphic_on_def
by (use assms in ⟨auto simp: meromorphic_on_def intro!: holomorphic_intros singularity_intros⟩)

lemma meromorphic_on_add [meromorphic_intros]:
assumes f meromorphic_on A pts g meromorphic_on A pts
shows $(\lambda z. f z + g z)$ meromorphic_on A pts
unfolding meromorphic_on_def
by (use assms in ⟨auto simp: meromorphic_on_def intro!: holomorphic_intros singularity_intros⟩)

lemma meromorphic_on_add':
assumes f meromorphic_on A pts1 g meromorphic_on A pts2
shows $(\lambda z. f z + g z)$ meromorphic_on A (pts1 \cup pts2)
proof (rule meromorphic_intros)
show f meromorphic_on A (pts1 \cup pts2)
by (rule meromorphic_on_superset_pts[OF assms(1)])
 (use assms in ⟨auto simp: meromorphic_on_def islimpt_Un⟩)
show g meromorphic_on A (pts1 \cup pts2)
by (rule meromorphic_on_superset_pts[OF assms(2)])
 (use assms in ⟨auto simp: meromorphic_on_def islimpt_Un⟩)
qed

lemma meromorphic_on_add_const [meromorphic_intros]:
assumes f meromorphic_on A pts
shows $(\lambda z. f z + c)$ meromorphic_on A pts
unfolding meromorphic_on_def
by (use assms in ⟨auto simp: meromorphic_on_def intro!: holomorphic_intros

singularity_intros)

lemma *meromorphic_on_minus_const* [*meromorphic_intros*]:
assumes *f meromorphic_on A pts*
shows $(\lambda z. f z - c)$ *meromorphic_on A pts*
using *meromorphic_on_add_const*[*OF assms, of -c*] **by** *simp*

lemma *meromorphic_on_diff* [*meromorphic_intros*]:
assumes *f meromorphic_on A pts g meromorphic_on A pts*
shows $(\lambda z. f z - g z)$ *meromorphic_on A pts*
using *meromorphic_on_add*[*OF assms(1) meromorphic_on_uminus*[*OF assms(2)*]]
by *simp*

lemma *meromorphic_on_diff'*:
assumes *f meromorphic_on A pts1 g meromorphic_on A pts2*
shows $(\lambda z. f z - g z)$ *meromorphic_on A (pts1 \cup pts2)*
proof (*rule meromorphic_intros*)
show *f meromorphic_on A (pts1 \cup pts2)*
by (*rule meromorphic_on_superset_pts*[*OF assms(1)*])
(use assms in <auto simp: meromorphic_on_def islimpt_Un>)
show *g meromorphic_on A (pts1 \cup pts2)*
by (*rule meromorphic_on_superset_pts*[*OF assms(2)*])
(use assms in <auto simp: meromorphic_on_def islimpt_Un>)
qed

lemma *meromorphic_on_mult* [*meromorphic_intros*]:
assumes *f meromorphic_on A pts g meromorphic_on A pts*
shows $(\lambda z. f z * g z)$ *meromorphic_on A pts*
unfolding *meromorphic_on_def*
by (*use assms in <auto simp: meromorphic_on_def intro!: holomorphic_intros singularity_intros>*)

lemma *meromorphic_on_mult'*:
assumes *f meromorphic_on A pts1 g meromorphic_on A pts2*
shows $(\lambda z. f z * g z)$ *meromorphic_on A (pts1 \cup pts2)*
proof (*rule meromorphic_intros*)
show *f meromorphic_on A (pts1 \cup pts2)*
by (*rule meromorphic_on_superset_pts*[*OF assms(1)*])
(use assms in <auto simp: meromorphic_on_def islimpt_Un>)
show *g meromorphic_on A (pts1 \cup pts2)*
by (*rule meromorphic_on_superset_pts*[*OF assms(2)*])
(use assms in <auto simp: meromorphic_on_def islimpt_Un>)
qed

lemma *meromorphic_on_imp_not_essential*:
assumes *f meromorphic_on A pts z \in A*
shows *not_essential f z*

```

proof (cases  $z \in pts$ )
  case False
  thus ?thesis
  using not_essential_holomorphic[of  $f A - pts z$ ] meromorphic_imp_open_diff[OF
  assms(1)] assms
  by (auto simp: meromorphic_on_def)
qed (use assms in ⟨auto simp: meromorphic_on_def⟩)

```

```

lemma meromorphic_imp_analytic:  $f$  meromorphic_on  $D pts \implies f$  analytic_on
( $D - pts$ )
  unfolding meromorphic_on_def
  apply (subst analytic_on_open)
  using meromorphic_imp_open_diff meromorphic_on_id apply blast
  apply auto
  done

```

```

lemma not_islimpt_isolated_zeros:
  assumes mero:  $f$  meromorphic_on  $A pts$  and  $z \in A$ 
  shows  $\neg z$  islimpt  $\{w \in A. \text{isolated\_zero } f w\}$ 
proof
  assume islimpt:  $z$  islimpt  $\{w \in A. \text{isolated\_zero } f w\}$ 
  have holo:  $f$  holomorphic_on  $A - pts$  and open  $A$ 
    using assms by (auto simp: meromorphic_on_def)
  have open': open  $(A - (pts - \{z\}))$ 
    by (intro meromorphic_imp_open_diff'[OF mero]) auto
  then obtain  $r$  where  $r: r > 0$  ball  $z r \subseteq A - (pts - \{z\})$ 
    using meromorphic_imp_open_diff[OF mero] ⟨ $z \in A$ ⟩ openE by blast

  have not_essential  $f z$ 
    using assms by (rule meromorphic_on_imp_not_essential)
  then consider  $c$  where  $f -z \rightarrow c$  | is_pole  $f z$ 
    unfolding not_essential_def by blast
  thus False
proof cases
  assume is_pole  $f z$ 
  hence eventually  $(\lambda w. f w \neq 0)$  (at  $z$ )
    by (rule non_zero_neighbour_pole)
  hence  $\neg z$  islimpt  $\{w. f w = 0\}$ 
    by (simp add: islimpt_conv_frequently_at_frequently_def)
  moreover have  $z$  islimpt  $\{w. f w = 0\}$ 
    using islimpt by (rule islimpt_subset) (auto simp: isolated_zero_def)
  ultimately show False by contradiction
next
  fix  $c$  assume  $c: f -z \rightarrow c$ 
  define  $g$  where  $g = (\lambda w. \text{if } w = z \text{ then } c \text{ else } f w)$ 
  have holo':  $g$  holomorphic_on  $A - (pts - \{z\})$  unfolding g_def
    by (intro removable_singularity holomorphic_on_subset[OF holo] open'  $c$ )
  auto

```

```

have eq_zero:  $g\ w = 0$  if  $w \in \text{ball } z\ r$  for  $w$ 
proof (rule analytic_continuation[where  $f = g$ ])
  show open (ball  $z\ r$ ) connected (ball  $z\ r$ )  $\{w \in \text{ball } z\ r. \text{isolated\_zero } f\ w\} \subseteq$ 
ball  $z\ r$ 
  by auto
  have  $z \text{ islimpt } \{w \in A. \text{isolated\_zero } f\ w\} \cap \text{ball } z\ r$ 
  using  $\text{islimpt } \langle r > 0 \rangle$  by (intro islimpt_Int_eventually_eventually_at_in_open)
auto
  also have  $\dots = \{w \in \text{ball } z\ r. \text{isolated\_zero } f\ w\}$ 
  using  $r$  by auto
  finally show  $z \text{ islimpt } \{w \in \text{ball } z\ r. \text{isolated\_zero } f\ w\}$ 
  by simp
next
fix  $w$  assume  $w: w \in \{w \in \text{ball } z\ r. \text{isolated\_zero } f\ w\}$ 
show  $g\ w = 0$ 
proof (cases  $w = z$ )
  case False
  thus ?thesis using  $w$  by (auto simp:  $g\_def$  isolated_zero_def)
next
  case True
  have  $z \text{ islimpt } \{z. f\ z = 0\}$ 
  using  $\text{islimpt}$  by (rule islimpt_subset) (auto simp: isolated_zero_def)
  thus ?thesis
  using  $w$  by (simp add: isolated_zero_altdef True)
qed
qed (use  $r$  that in  $\langle \text{auto intro!}: \text{holomorphic\_on\_subset}[OF\ \text{holo}] \text{ simp: isolated\_zero\_def} \rangle$ )

have infinite ( $\{w \in A. \text{isolated\_zero } f\ w\} \cap \text{ball } z\ r$ )
  using  $\text{islimpt } \langle r > 0 \rangle$  unfolding  $\text{islimpt\_eq\_infinite\_ball}$  by blast
hence  $\{w \in A. \text{isolated\_zero } f\ w\} \cap \text{ball } z\ r \neq \{\}$ 
  by force
then obtain  $z0$  where  $z0: z0 \in A \text{ isolated\_zero } f\ z0\ z0 \in \text{ball } z\ r$ 
  by blast
have  $\forall_F\ y$  in at  $z0. y \in \text{ball } z\ r - (\text{if } z = z0 \text{ then } \{\} \text{ else } \{z\}) - \{z0\}$ 
  using  $r\ z0$  by (intro eventually_at_in_open) auto
hence eventually  $(\lambda w. f\ w = 0)$  (at  $z0$ )
proof eventually_elim
  case (elim  $w$ )
  show ?case
  using  $\text{eq\_zero}[of\ w]$  elim by (auto simp:  $g\_def$  split: if_splits)
qed
hence eventually  $(\lambda w. f\ w = 0)$  (at  $z0$ )
  by (auto simp:  $g\_def$  eventually_at_filter elim!: eventually_mono split:
if_splits)
moreover from  $z0$  have eventually  $(\lambda w. f\ w \neq 0)$  (at  $z0$ )
  by (simp add: isolated_zero_def)
ultimately have eventually  $(\lambda_. \text{False})$  (at  $z0$ )
  by eventually_elim auto

```

```

    thus False
      by simp
  qed
qed

lemma closedin_isolated_zeros:
  assumes f meromorphic_on A pts
  shows closedin (top_of_set A) {z ∈ A. isolated_zero f z}
  unfolding closedin_limpt using not_islimpt_isolated_zeros[OF assms] by auto

lemma meromorphic_on_deriv':
  assumes f meromorphic_on A pts open A
  assumes  $\bigwedge x. x \in A - pts \implies (f \text{ has\_field\_derivative } f' x) (at x)$ 
  shows f' meromorphic_on A pts
  unfolding meromorphic_on_def
proof (intro conjI ballI)
  have open (A - pts)
    by (intro meromorphic_imp_open_diff[OF assms(1)])
  thus f' holomorphic_on A - pts
    by (rule derivative_is_holomorphic) (use assms in auto)
next
  fix z assume z ∈ pts
  hence z ∈ A
    using assms(1) by (auto simp: meromorphic_on_def)
  from ⟨z ∈ pts⟩ obtain r where r: r > 0 f analytic_on ball z r - {z}
    using assms(1) by (auto simp: meromorphic_on_def isolated_singularity_at_def)

  have open (ball z r ∩ (A - (pts - {z})))
    by (intro open_Int assms meromorphic_imp_open_diff'[OF assms(1)]) auto
  then obtain r' where r': r' > 0 ball z r' ⊆ ball z r ∩ (A - (pts - {z}))
    using r ⟨z ∈ A⟩ by (subst (asm) open_contains_ball) fastforce

  have open (ball z r' - {z})
    by auto
  hence f' holomorphic_on ball z r' - {z}
    by (rule derivative_is_holomorphic[of _ f]) (use r' in ⟨auto intro!: assms(3)⟩)
  moreover have open (ball z r' - {z})
    by auto
  ultimately show isolated_singularity_at f' z
    unfolding isolated_singularity_at_def using ⟨r' > 0⟩
    by (auto simp: analytic_on_open intro!: exI[of _ r'])
next
  fix z assume z: z ∈ pts
  hence z': not_essential f z z ∈ A
    using assms by (auto simp: meromorphic_on_def)
  from z'(1) show not_essential f' z
proof (rule not_essential_deriv')
  show z ∈ A - (pts - {z})
    using ⟨z ∈ A⟩ by blast

```

```

  show open (A - (pts - {z}))
    by (intro meromorphic_imp_open_diff'[OF assms(1)]) auto
  qed (use assms in auto)
qed (use assms in ‹auto simp: meromorphic_on_def›)

```

```

lemma meromorphic_on_deriv [meromorphic_intros]:
  assumes f meromorphic_on A pts open A
  shows deriv f meromorphic_on A pts
proof (intro meromorphic_on_deriv'[OF assms(1)])
  have *: open (A - pts)
    by (intro meromorphic_imp_open_diff'[OF assms(1)])
  show (f has_field_derivative deriv f x) (at x) if x ∈ A - pts for x
    using assms(1) by (intro holomorphic_derivI[OF _ * that]) (auto simp: meromorphic_on_def)
  qed fact

```

```

lemma meromorphic_on_imp_analytic_at:
  assumes f meromorphic_on A pts z ∈ A - pts
  shows f analytic_on {z}
  using assms by (metis analytic_at meromorphic_imp_open_diff meromorphic_on_def)

```

```

lemma meromorphic_compact_finite_pts:
  assumes f meromorphic_on D pts compact S S ⊆ D
  shows finite (S ∩ pts)
proof -
  { assume infinite (S ∩ pts)
    then obtain z where z ∈ S and z: z islimpt (S ∩ pts)
      using assms by (metis compact_eq_Bolzano_Weierstrass inf_le1)
    then have False
      using assms by (meson in_mono inf_le2 islimpt_subset meromorphic_on_def)
  }
  then show ?thesis by metis
qed

```

```

lemma meromorphic_imp_countable:
  assumes f meromorphic_on D pts
  shows countable pts
proof -
  obtain K :: nat ⇒ complex set where K: D = (⋃ n. K n) ∧ n. compact (K n)
    using assms unfolding meromorphic_on_def by (metis open_Union_compact_subsets)
  then have pts = (⋃ n. K n ∩ pts)
    using assms meromorphic_on_def by auto
  moreover have ∧ n. finite (K n ∩ pts)
    by (metis K(1) K(2) UN_I assms image_iff meromorphic_compact_finite_pts rangeI subset_eq)
  ultimately show ?thesis
    by (metis countableI_type countable_UN countable_finite)
qed

```



```

lemma meromorphic_imp_connected_diff':
  assumes f meromorphic_on D pts connected D pts'  $\subseteq$  pts
  shows connected (D - pts')
proof (rule connected_open_diff_countable)
  show countable pts'
    by (rule countable_subset [OF assms(3)]) (use assms(1) in  $\langle$ auto simp: meromorphic_imp_countable $\rangle$ )
qed (use assms in  $\langle$ auto simp: meromorphic_on_def $\rangle$ )

lemma meromorphic_imp_connected_diff:
  assumes f meromorphic_on D pts connected D
  shows connected (D - pts)
  using meromorphic_imp_connected_diff'[OF assms order.refl] .

lemma meromorphic_on_compose [meromorphic_intros]:
  assumes f: f meromorphic_on A pts and g: g holomorphic_on B
  assumes open B and g ' B  $\subseteq$  A
  shows  $(\lambda x. f (g x))$  meromorphic_on B (isolated_points_of (g - ' pts  $\cap$  B))
  unfolding meromorphic_on_def
proof (intro ballI conjI)
  fix z assume z: z  $\in$  isolated_points_of (g - ' pts  $\cap$  B)
  hence z': z  $\in$  B g z  $\in$  pts
    using isolated_points_of_subset by blast+
  have g': g analytic_on {z}
    using g z'  $\langle$ open B $\rangle$  analytic_at by blast

  show isolated_singularity_at  $(\lambda x. f (g x))$  z
    by (rule isolated_singularity_at_compose[OF _ g']) (use f z' in  $\langle$ auto simp: meromorphic_on_def $\rangle$ )
  show not_essential  $(\lambda x. f (g x))$  z
    by (rule not_essential_compose[OF _ g']) (use f z' in  $\langle$ auto simp: meromorphic_on_def $\rangle$ )
next
  fix z assume z: z  $\in$  B
  hence g z  $\in$  A
    using assms by auto
  hence  $\neg$ g z islimpt pts
    using f by (auto simp: meromorphic_on_def)
  hence ev: eventually  $(\lambda w. w \notin$  pts) (at (g z))
    by (auto simp: islimpt_conv_frequently_at frequently_def)
  have g': g analytic_on {z}
    by (rule holomorphic_on_imp_analytic_at[OF g]) (use assms z in auto)

have eventually  $(\lambda w. w \notin$  isolated_points_of (g - ' pts  $\cap$  B)) (at z)
proof (cases isolated_zero  $(\lambda w. g w - g z)$  z)
  case True
    have eventually  $(\lambda w. w \notin$  pts) (at (g z))
      using ev by (auto simp: islimpt_conv_frequently_at frequently_def)

```

```

moreover have  $g - z \rightarrow g z$ 
  using analytic_at_imp_isCont[OF  $g'$ ] isContD by blast
hence lim: filterlim  $g$  (at ( $g z$ )) (at  $z$ )
  using True by (auto simp: filterlim_at isolated_zero_def)
have eventually ( $\lambda w. g w \notin pts$ ) (at  $z$ )
  using ev lim by (rule eventually_compose_filterlim)
thus ?thesis
  by eventually_elim (auto simp: isolated_points_of_def)
next
case False
have eventually ( $\lambda w. g w - g z = 0$ ) (nhds  $z$ )
  using False by (rule non_isolated_zero) (auto intro!: analytic_intros g')
hence eventually ( $\lambda w. g w = g z \wedge w \in B$ ) (nhds  $z$ )
  using eventually_nhds_in_open[OF  $\langle open B \rangle \langle z \in B \rangle$ ]
  by eventually_elim auto
then obtain  $X$  where  $X$ : open  $X$   $z \in X$   $X \subseteq B$   $\forall x \in X. g x = g z$ 
  unfolding eventually_nhds by blast

have  $z0 \notin isolated\_points\_of (g - ' pts \cap B)$  if  $z0 \in X$  for  $z0$ 
proof (cases  $g z \in pts$ )
  case False
  with that have  $g z0 \notin pts$ 
    using  $X$  by metis
  thus ?thesis
    by (auto simp: isolated_points_of_def)
next
case True
have eventually ( $\lambda w. w \in X$ ) (at  $z0$ )
  by (intro eventually_at_in_open') fact+
hence eventually ( $\lambda w. w \in g - ' pts \cap B$ ) (at  $z0$ )
  by eventually_elim (use X True in fastforce)
hence frequently ( $\lambda w. w \in g - ' pts \cap B$ ) (at  $z0$ )
  by (meson at_neq_bot eventually_frequently)
thus  $z0 \notin isolated\_points\_of (g - ' pts \cap B)$ 
  unfolding isolated_points_of_def by (auto simp: frequently_def)
qed
moreover have eventually ( $\lambda x. x \in X$ ) (at  $z$ )
  by (intro eventually_at_in_open') fact+
ultimately show ?thesis
  by (auto elim!: eventually_mono)
qed
thus  $\neg z$  islimpt isolated_points_of ( $g - ' pts \cap B$ )
  by (auto simp: islimpt_conv_frequently_at frequently_def)
next
have  $f \circ g$  analytic_on ( $\bigcup z \in B - isolated\_points\_of (g - ' pts \cap B). \{z\}$ )
  unfolding analytic_on_UN
proof
  fix  $z$  assume  $z$ :  $z \in B - isolated\_points\_of (g - ' pts \cap B)$ 
  hence  $z \in B$  by blast

```

```

have g': g analytic_on {z}
  by (rule holomorphic_on_imp_analytic_at[OF g]) (use assms z in auto)
show f ∘ g analytic_on {z}
proof (cases g z ∈ pts)
  case False
  show ?thesis
  proof (rule analytic_on_compose)
    show f analytic_on g -' {z} using False z assms
    by (auto intro!: meromorphic_on_imp_analytic_at[OF f])
  qed fact
next
case True
show ?thesis
proof (cases isolated_zero (λw. g w - g z) z)
  case False
  hence eventually (λw. g w - g z = 0) (nhds z)
    by (rule non_isolated_zero) (auto intro!: analytic_intros g')
  hence f ∘ g analytic_on {z} ↔ (λ_. f (g z)) analytic_on {z}
    by (intro analytic_at_cong) (auto elim!: eventually_mono)
  thus ?thesis
    by simp
next
case True
  hence ev: eventually (λw. g w ≠ g z) (at z)
    by (auto simp: isolated_zero_def)

  have ¬g z islimpt pts
    using ⟨g z ∈ pts⟩ f by (auto simp: meromorphic_on_def)
  hence eventually (λw. w ∉ pts) (at (g z))
    by (auto simp: islimpt_conv_frequently_at frequently_def)
  moreover have g -z → g z
    using analytic_at_imp_isCont[OF g'] isContD by blast
  with ev have filterlim g (at (g z)) (at z)
    by (auto simp: filterlim_at)
  ultimately have eventually (λw. g w ∉ pts) (at z)
    using eventually_compose_filterlim by blast
  hence z ∈ isolated_points_of (g -' pts ∩ B)
    using ⟨g z ∈ pts⟩ ⟨z ∈ B⟩
    by (auto simp: isolated_points_of_def elim!: eventually_mono)
  with z show ?thesis by simp
qed
qed
qed
also have ... = B - isolated_points_of (g -' pts ∩ B)
  by blast
finally show (λx. f (g x)) holomorphic_on B - isolated_points_of (g -' pts ∩ B)
  unfolding o_def using analytic_imp_holomorphic by blast
qed (auto simp: isolated_points_of_def ⟨open B⟩)

```

```

lemma meromorphic_on_compose':
  assumes f: f meromorphic_on A pts and g: g holomorphic_on B
  assumes open B and g ' B  $\subseteq$  A and pts' = (isolated_points_of (g - ' pts  $\cap$  B))
  shows ( $\lambda x. f (g x)$ ) meromorphic_on B pts'
  using meromorphic_on_compose[OF assms(1-4)] assms(5) by simp

lemma meromorphic_on_inverse': inverse meromorphic_on UNIV 0
  unfolding meromorphic_on_def
  by (auto intro!: holomorphic_intros singularity_intros not_essential_inverse
      isolated_singularity_at_inverse simp: islimpt_finite)

lemma meromorphic_on_inverse [meromorphic_intros]:
  assumes mero: f meromorphic_on A pts
  shows ( $\lambda z. inverse (f z)$ ) meromorphic_on A (pts  $\cup$  {z $\in$ A. isolated_zero f z})
proof -
  have open A
    using mero by (auto simp: meromorphic_on_def)
  have open': open (A - pts)
    by (intro meromorphic_imp_open_diff[OF mero])
  have holo: f holomorphic_on A - pts
    using assms by (auto simp: meromorphic_on_def)
  have ana: f analytic_on A - pts
    using open' holo by (simp add: analytic_on_open)

  show ?thesis
    unfolding meromorphic_on_def
  proof (intro conjI ballI)
    fix z assume z: z  $\in$  pts  $\cup$  {z $\in$ A. isolated_zero f z}
    have isolated_singularity_at f z  $\wedge$  not_essential f z
    proof (cases z  $\in$  pts)
      case False
        have f holomorphic_on A - pts - {z}
          by (intro holomorphic_on_subset[OF holo]) auto
        hence isolated_singularity_at f z
          by (rule isolated_singularity_at_holomorphic)
          (use z False in (auto intro!: meromorphic_imp_open_diff[OF mero]))
        moreover have not_essential f z
          using z False
        by (intro not_essential_holomorphic[OF holo] meromorphic_imp_open_diff[OF
mero]) auto
      ultimately show ?thesis by blast
    qed (use assms in (auto simp: meromorphic_on_def))
    thus isolated_singularity_at ( $\lambda z. inverse (f z)$ ) z not_essential ( $\lambda z. inverse (f z)$ ) z
    by (auto intro!: isolated_singularity_at_inverse not_essential_inverse)
  next
    fix z assume z  $\in$  A
    hence  $\neg$  z islimpt {z $\in$ A. isolated_zero f z}

```

```

    by (rule not_islimpt_isolated_zeros[OF mero])
  thus  $\neg z \text{ islimpt } pts \cup \{z \in A. \text{ isolated\_zero } f z\}$  using  $\langle z \in A \rangle$ 
    using mero by (auto simp: islimpt_Un meromorphic_on_def)
next
  show  $pts \cup \{z \in A. \text{ isolated\_zero } f z\} \subseteq A$ 
    using mero by (auto simp: meromorphic_on_def)
next
  have  $(\lambda z. \text{ inverse } (f z)) \text{ analytic\_on } (\bigcup_{w \in A} - (pts \cup \{z \in A. \text{ isolated\_zero } f z\}) . \{w\})$ 
    unfolding analytic_on_UN
  proof (intro ballI)
    fix w assume  $w: w \in A - (pts \cup \{z \in A. \text{ isolated\_zero } f z\})$ 
    show  $(\lambda z. \text{ inverse } (f z)) \text{ analytic\_on } \{w\}$ 
      proof (cases  $f w = 0$ )
        case False
          thus ?thesis using w
            by (intro analytic_intros analytic_on_subset[OF ana]) auto
        case True
          have eventually  $(\lambda w. f w = 0)$  (nhds w)
            using True w by (intro non_isolated_zero analytic_on_subset[OF ana])
          auto
          hence  $(\lambda z. \text{ inverse } (f z)) \text{ analytic\_on } \{w\} \longleftrightarrow (\lambda_. 0) \text{ analytic\_on } \{w\}$ 
            using w by (intro analytic_at_cong_refl) auto
          thus ?thesis
            by simp
      qed
    qed
  also have  $\dots = A - (pts \cup \{z \in A. \text{ isolated\_zero } f z\})$ 
    by blast
  finally have  $(\lambda z. \text{ inverse } (f z)) \text{ analytic\_on } \dots$  .
  moreover have open  $(A - (pts \cup \{z \in A. \text{ isolated\_zero } f z\}))$ 
    using closedin_isolated_zeros[OF mero] open'  $\langle \text{open } A \rangle$ 
  by (metis (no_types, lifting) Diff_Diff_Int Diff_Un closedin_closed open_Diff
open_Int)
  ultimately show  $(\lambda z. \text{ inverse } (f z)) \text{ holomorphic\_on } A - (pts \cup \{z \in A. \text{ isolated\_zero } f z\})$ 
    by (subst (asm) analytic_on_open) auto
  qed (use assms in  $\langle \text{auto simp: meromorphic_on_def islimpt_Un} \rangle$ 
    intro!: holomorphic_intros singularity_intros)
qed

lemma meromorphic_on_inverse'' [meromorphic_intros]:
  assumes  $f \text{ meromorphic\_on } A \text{ pts } \{z \in A. f z = 0\} \subseteq \text{pts}$ 
  shows  $(\lambda z. \text{ inverse } (f z)) \text{ meromorphic\_on } A \text{ pts}$ 
proof -
  have  $(\lambda z. \text{ inverse } (f z)) \text{ meromorphic\_on } A (pts \cup \{z \in A. \text{ isolated\_zero } f z\})$ 
    by (intro meromorphic_on_inverse assms)
  also have  $(pts \cup \{z \in A. \text{ isolated\_zero } f z\}) = \text{pts}$ 

```

```

    using assms(2) by (auto simp: isolated_zero_def)
    finally show ?thesis .
qed

```

```

lemma meromorphic_on_divide [meromorphic_intros]:
  assumes f meromorphic_on A pts and g meromorphic_on A pts
  shows  $(\lambda z. f z / g z)$  meromorphic_on A (pts  $\cup$  {z $\in$ A. isolated_zero g z})
proof -
  have mero1:  $(\lambda z. \text{inverse } (g z))$  meromorphic_on A (pts  $\cup$  {z $\in$ A. isolated_zero g z})
    by (intro meromorphic_intros assms)
  have sparse:  $\forall x \in A. \neg x \text{ islimpt } pts \cup \{z \in A. \text{isolated\_zero } g z\}$  and pts  $\subseteq$  A
    using mero1 by (auto simp: meromorphic_on_def)
  have mero2: f meromorphic_on A (pts  $\cup$  {z $\in$ A. isolated_zero g z})
    by (rule meromorphic_on_superset_pts[OF assms(1)]) (use sparse  $\langle pts \subseteq A \rangle$ )
in auto)
  have  $(\lambda z. f z * \text{inverse } (g z))$  meromorphic_on A (pts  $\cup$  {z $\in$ A. isolated_zero g z})
    by (intro meromorphic_on_mult mero1 mero2)
  thus ?thesis
    by (simp add: field_simps)
qed

```

```

lemma meromorphic_on_divide' [meromorphic_intros]:
  assumes f meromorphic_on A pts g meromorphic_on A pts  $\{z \in A. g z = 0\} \subseteq pts$ 
  shows  $(\lambda z. f z / g z)$  meromorphic_on A pts
proof -
  have  $(\lambda z. f z * \text{inverse } (g z))$  meromorphic_on A pts
    by (intro meromorphic_intros assms)
  thus ?thesis
    by (simp add: field_simps)
qed

```

```

lemma meromorphic_on_cmult_left [meromorphic_intros]:
  assumes f meromorphic_on A pts
  shows  $(\lambda x. c * f x)$  meromorphic_on A pts
  using assms by (intro meromorphic_intros) (auto simp: meromorphic_on_def)

```

```

lemma meromorphic_on_cmult_right [meromorphic_intros]:
  assumes f meromorphic_on A pts
  shows  $(\lambda x. f x * c)$  meromorphic_on A pts
  using assms by (intro meromorphic_intros) (auto simp: meromorphic_on_def)

```

```

lemma meromorphic_on_scaleR [meromorphic_intros]:
  assumes f meromorphic_on A pts
  shows  $(\lambda x. c *_R f x)$  meromorphic_on A pts
  using assms unfolding scaleR_conv_of_real
  by (intro meromorphic_intros) (auto simp: meromorphic_on_def)

```

```

lemma meromorphic_on_sum [meromorphic_intros]:
  assumes  $\bigwedge y. y \in I \implies f y \text{ meromorphic\_on } A \text{ pts}$ 
  assumes  $I \neq \{\}$   $\vee \text{ open } A \wedge \text{ pts} \subseteq A \wedge (\forall x \in A. \neg x \text{ islimpt pts})$ 
  shows  $(\lambda x. \sum_{y \in I} f y x) \text{ meromorphic\_on } A \text{ pts}$ 
proof -
  have *:  $\text{open } A \wedge \text{ pts} \subseteq A \wedge (\forall x \in A. \neg x \text{ islimpt pts})$ 
    using assms(2)
  proof
    assume  $I \neq \{\}$ 
    then obtain  $x$  where  $x \in I$ 
      by blast
    from assms(1)[OF this] show ?thesis
      by (auto simp: meromorphic_on_def)
  qed auto
  show ?thesis
    using assms(1)
    by (induction I rule: infinite_finite_induct) (use * in  $\langle \text{auto intro!: meromorphic\_intros} \rangle$ )
qed

```

```

lemma meromorphic_on_prod [meromorphic_intros]:
  assumes  $\bigwedge y. y \in I \implies f y \text{ meromorphic\_on } A \text{ pts}$ 
  assumes  $I \neq \{\}$   $\vee \text{ open } A \wedge \text{ pts} \subseteq A \wedge (\forall x \in A. \neg x \text{ islimpt pts})$ 
  shows  $(\lambda x. \prod_{y \in I} f y x) \text{ meromorphic\_on } A \text{ pts}$ 
proof -
  have *:  $\text{open } A \wedge \text{ pts} \subseteq A \wedge (\forall x \in A. \neg x \text{ islimpt pts})$ 
    using assms(2)
  proof
    assume  $I \neq \{\}$ 
    then obtain  $x$  where  $x \in I$ 
      by blast
    from assms(1)[OF this] show ?thesis
      by (auto simp: meromorphic_on_def)
  qed auto
  show ?thesis
    using assms(1)
    by (induction I rule: infinite_finite_induct) (use * in  $\langle \text{auto intro!: meromorphic\_intros} \rangle$ )
qed

```

```

lemma meromorphic_on_power [meromorphic_intros]:
  assumes  $f \text{ meromorphic\_on } A \text{ pts}$ 
  shows  $(\lambda x. f x \wedge^n) \text{ meromorphic\_on } A \text{ pts}$ 
proof -
  have  $(\lambda x. \prod_{i \in \{..<n\}} f x) \text{ meromorphic\_on } A \text{ pts}$ 
    by (intro meromorphic_intros assms(1)) (use assms in  $\langle \text{auto simp: meromorphic\_on\_def} \rangle$ )
  thus ?thesis

```

by simp
qed

lemma meromorphic_on_power_int [meromorphic_intros]:

assumes f meromorphic_on A pts

shows $(\lambda z. f z \text{ powi } n)$ meromorphic_on A $(pts \cup \{z \in A. \text{isolated_zero } f z\})$

proof –

have $inv: (\lambda x. \text{inverse } (f x))$ meromorphic_on A $(pts \cup \{z \in A. \text{isolated_zero } f z\})$

by (intro meromorphic_intros $assms$)

have $*$: f meromorphic_on A $(pts \cup \{z \in A. \text{isolated_zero } f z\})$

by (intro meromorphic_on_superset_pts [OF $assms(1)$])

(use inv in $\langle auto \text{ simp: meromorphic_on_def} \rangle$)

show ?thesis

proof (cases $n \geq 0$)

case True

have $(\lambda x. f x \wedge \text{nat } n)$ meromorphic_on A $(pts \cup \{z \in A. \text{isolated_zero } f z\})$

by (intro meromorphic_intros $*$)

thus ?thesis

using True by (simp add: power_int_def)

next

case False

have $(\lambda x. \text{inverse } (f x) \wedge \text{nat } (-n))$ meromorphic_on A $(pts \cup \{z \in A. \text{isolated_zero } f z\})$

by (intro meromorphic_intros $assms$)

thus ?thesis

using False by (simp add: power_int_def)

qed

qed

lemma meromorphic_on_power_int' [meromorphic_intros]:

assumes f meromorphic_on A pts $n \geq 0 \vee (\forall z \in A. \text{isolated_zero } f z \longrightarrow z \in pts)$

shows $(\lambda z. f z \text{ powi } n)$ meromorphic_on A pts

proof (cases $n \geq 0$)

case True

have $(\lambda z. f z \wedge \text{nat } n)$ meromorphic_on A pts

by (intro meromorphic_intros $assms$)

thus ?thesis

using True by (simp add: power_int_def)

next

case False

have $(\lambda z. f z \text{ powi } n)$ meromorphic_on A $(pts \cup \{z \in A. \text{isolated_zero } f z\})$

by (rule meromorphic_on_power_int) fact

also from $assms(2)$ False have $pts \cup \{z \in A. \text{isolated_zero } f z\} = pts$

by auto

finally show ?thesis .

qed


```

lemma has_laurent_expansion_on_imp_meromorphic_on:
  assumes open A
  assumes laurent:  $\bigwedge z. z \in A \implies \exists F. (\lambda w. f (z + w)) \text{ has\_laurent\_expansion } F$ 
  shows f meromorphic_on A  $\{z \in A. \neg f \text{ analytic\_on } \{z\}\}$ 
  unfolding meromorphic_on_def
proof (intro conjI ballI)
  fix z assume  $z \in \{z \in A. \neg f \text{ analytic\_on } \{z\}\}$ 
  then obtain F where F:  $(\lambda w. f (z + w)) \text{ has\_laurent\_expansion } F$ 
  using laurent[of z] by blast
  from F show not_essential f z isolated_singularity_at f z
  using has_laurent_expansion_not_essential has_laurent_expansion_isolated
by blast+
next
  fix z assume  $z \in A$ 
  obtain F where F:  $(\lambda w. f (z + w)) \text{ has\_laurent\_expansion } F$ 
  using laurent[of z]  $\langle z \in A \rangle$  by blast
  from F have isolated_singularity_at f z
  using has_laurent_expansion_isolated z by blast
  then obtain r where  $r: r > 0 \wedge f \text{ analytic\_on ball } z \ r - \{z\}$ 
  unfolding isolated_singularity_at_def by blast
  have f analytic_on  $\{w\}$  if  $w \in \text{ball } z \ r - \{z\}$  for w
  by (rule analytic_on_subset[OF r(2)]) (use that in auto)
  hence eventually  $(\lambda w. f \text{ analytic\_on } \{w\})$  (at z)
  using eventually_at_in_open[of ball z r z]  $\langle r > 0 \rangle$  by (auto elim!: eventually_mono)
  hence  $\neg z \text{ islimpt } \{w. \neg f \text{ analytic\_on } \{w\}\}$ 
  by (auto simp: islimpt_conv_frequently_at_frequently_def)
  thus  $\neg z \text{ islimpt } \{w \in A. \neg f \text{ analytic\_on } \{w\}\}$ 
  using islimpt_subset[of z]  $\{w \in A. \neg f \text{ analytic\_on } \{w\}\}$   $\{w. \neg f \text{ analytic\_on } \{w\}\}$  by blast
next
  have f analytic_on  $A - \{w \in A. \neg f \text{ analytic\_on } \{w\}\}$ 
  by (subst analytic_on_analytic_at) auto
  thus f holomorphic_on  $A - \{w \in A. \neg f \text{ analytic\_on } \{w\}\}$ 
  by (meson analytic_imp_holomorphic)
qed (use assms in auto)

lemma meromorphic_on_imp_has_laurent_expansion:
  assumes f meromorphic_on A pts  $z \in A$ 
  shows  $(\lambda w. f (z + w)) \text{ has\_laurent\_expansion } \text{laurent\_expansion } f \ z$ 
proof (cases z \in pts)
  case True
  thus ?thesis
  using assms by (intro not_essential_has_laurent_expansion) (auto simp: meromorphic_on_def)
next
  case False
  have f holomorphic_on  $(A - \text{pts})$ 
  using assms by (auto simp: meromorphic_on_def)

```

```

moreover have  $z \in A - pts$  open ( $A - pts$ )
using assms(2) False by (auto intro!: meromorphic_imp_open_diff[OF assms(1)])
ultimately have  $f$  analytic_on  $\{z\}$ 
  unfolding analytic_at by blast
thus ?thesis
  using isolated_singularity_at_analytic not_essential_analytic
    not_essential_has_laurent_expansion by blast
qed

```

lemma

```

assumes isolated_singularity_at  $f$   $z$   $f -z \rightarrow c$ 
shows eventually_remove_sings_eq_nhds':
  eventually ( $\lambda w.$  remove_sings  $f$   $w =$  (if  $w = z$  then  $c$  else  $f$   $w$ )) (nhds  $z$ )
  and remove_sings_analytic_at_singularity: remove_sings  $f$  analytic_on  $\{z\}$ 
proof -
  have eventually ( $\lambda w.$   $w \neq z$ ) (at  $z$ )
    by (auto simp: eventually_at_filter)
  hence eventually ( $\lambda w.$  remove_sings  $f$   $w =$  (if  $w = z$  then  $c$  else  $f$   $w$ )) (at  $z$ )
    using eventually_remove_sings_eq_at[OF assms(1)]
    by eventually_elim auto
  moreover have remove_sings  $f$   $z = c$ 
    using assms by auto
  ultimately show ev: eventually ( $\lambda w.$  remove_sings  $f$   $w =$  (if  $w = z$  then  $c$  else
     $f$   $w$ )) (nhds  $z$ )
    by (simp add: eventually_at_filter)

  have ( $\lambda w.$  if  $w = z$  then  $c$  else  $f$   $w$ ) analytic_on  $\{z\}$ 
    by (intro removable_singularity' assms)
  also have ?this  $\longleftrightarrow$  remove_sings  $f$  analytic_on  $\{z\}$ 
    using ev by (intro analytic_at_cong) (auto simp: eq_commute)
  finally show ... .

```

qed

lemma *remove_sings_meromorphic_on*:

```

assumes  $f$  meromorphic_on  $A$   $pts \bigwedge z. z \in pts - pts' \implies \neg is\_pole$   $f$   $z$   $pts' \subseteq$ 
 $pts$ 
shows remove_sings  $f$  meromorphic_on  $A$   $pts'$ 
unfolding meromorphic_on_def
proof safe
have remove_sings  $f$  analytic_on  $\{z\}$  if  $z \in A - pts'$  for  $z$ 
proof (cases  $z \in pts$ )
  case False
    hence  $*$ :  $f$  analytic_on  $\{z\}$ 
      using assms meromorphic_imp_open_diff[OF assms(1)] that
      by (force simp: meromorphic_on_def analytic_at)
    have remove_sings  $f$  analytic_on  $\{z\} \longleftrightarrow$   $f$  analytic_on  $\{z\}$ 
      by (intro analytic_at_cong eventually_remove_sings_eq_nhds  $*$  refl)
    thus ?thesis using  $*$  by simp
  next

```

```

case True
have isol: isolated_singularity_at f z
  using True using assms by (auto simp: meromorphic_on_def)
from assms(1) have not_essential f z
  using True by (auto simp: meromorphic_on_def)
with assms(2) True that obtain c where f  $-z \rightarrow c$ 
  by (auto simp: not_essential_def)
thus remove_sings f analytic_on {z}
  by (intro remove_sings_analytic_at_singularity isol)
qed
hence remove_sings f analytic_on A - pts'
  by (subst analytic_on_analytic_at) auto
thus remove_sings f holomorphic_on A - pts'
  using meromorphic_imp_open_diff'[OF assms(1,3)] by (subst (asm) analytic_on_open)
qed (use assms islimpt_subset[OF assms(3)] in  $\langle$ auto simp: meromorphic_on_def $\rangle$ )

```

```

lemma remove_sings_holomorphic_on:
  assumes f meromorphic_on A pts  $\wedge z. z \in pts \implies \neg is\_pole$  f z
  shows remove_sings f holomorphic_on A
  using remove_sings_meromorphic_on[OF assms(1), of {}] assms(2)
  by (auto simp: meromorphic_on_no_singularities)

```

```

lemma meromorphic_on_Ex_iff:
  ( $\exists pts. f$  meromorphic_on A pts)  $\longleftrightarrow$ 
  open A  $\wedge (\forall z \in A. \exists F. (\lambda w. f$  (z + w)) has_laurent_expansion F)
proof safe
  fix pts assume *: f meromorphic_on A pts
  from * show open A
  by (auto simp: meromorphic_on_def)
  show  $\exists F. (\lambda w. f$  (z + w)) has_laurent_expansion F if z  $\in A$  for z
  using that *
  by (intro exI[of _ Laurent_expansion f z] meromorphic_on_imp_has_laurent_expansion)
qed (blast intro!: has_laurent_expansion_on_imp_meromorphic_on)

```

```

lemma is_pole_inverse_holomorphic_pts:
  fixes pts::complex set and f::complex  $\Rightarrow$  complex
  defines g  $\equiv \lambda x. (if$  x $\in pts$  then 0 else inverse (f x))
  assumes mer: f meromorphic_on D pts
  and non_z:  $\bigwedge z. z \in D - pts \implies f$  z  $\neq 0$ 
  and all_poles:  $\forall x. is\_pole$  f x  $\longleftrightarrow x \in pts$ 
  shows g holomorphic_on D
proof -
  have open D and f_holo: f holomorphic_on (D - pts)
  using mer by (auto simp: meromorphic_on_def)
  have  $\exists r. r > 0 \wedge f$  analytic_on ball z r - {z}
   $\wedge (\forall x \in ball$  z r - {z}. f x  $\neq 0$ ) if z $\in pts$  for z
proof -
  have isolated_singularity_at f z is_pole f z

```

```

    using mer meromorphic_on_def that all_poles by blast+
  then obtain r1 where r1>0 and fan: f analytic_on ball z r1 - {z}
    by (meson isolated_singularity_at_def)
  obtain r2 where r2>0  $\forall x \in \text{ball } z \text{ } r2 - \{z\}. f x \neq 0$ 
    using non_zero_neighbour_pole[OF ‹is_pole f z›]
  unfolding eventually_at by (metis Diff_iff UNIV_I dist_commute insertI1
mem_ball)
  define r where r = min r1 r2
  have r>0 by (simp add: ‹0 < r2› ‹r1>0› r_def)
  moreover have f analytic_on ball z r - {z}
    using r_def by (force intro: analytic_on_subset [OF fan])
  moreover have  $\forall x \in \text{ball } z \text{ } r - \{z\}. f x \neq 0$ 
    by (simp add: ‹ $\forall x \in \text{ball } z \text{ } r2 - \{z\}. f x \neq 0$ › r_def)
  ultimately show ?thesis by auto
qed
then obtain get_r where r_pos:get_r z>0
  and r_ana:f analytic_on ball z (get_r z) - {z}
  and r_nz: $\forall x \in \text{ball } z \text{ } (\text{get\_r } z) - \{z\}. f x \neq 0$ 
  if z∈pts for z
  by metis
define p_balls where p_balls  $\equiv \bigcup z \in \text{pts}. \text{ball } z \text{ } (\text{get\_r } z)$ 
have g_ball:g holomorphic_on ball z (get_r z) if z∈pts for z
proof -
  have ( $\lambda x. \text{if } x = z \text{ then } 0 \text{ else inverse } (f x)$ ) holomorphic_on ball z (get_r z)
  proof (rule is_pole_inverse_holomorphic)
    show f holomorphic_on ball z (get_r z) - {z}
      using analytic_imp_holomorphic r_ana that by blast
    show is_pole f z
      using mer meromorphic_on_def that all_poles by force
    show  $\forall x \in \text{ball } z \text{ } (\text{get\_r } z) - \{z\}. f x \neq 0$ 
      using r_nz that by metis
  qed auto
  then show ?thesis unfolding g_def
    by (smt (verit, ccfv_SIG) Diff_iff Elementary_Metric_Spaces.open_ball
all_poles analytic_imp_holomorphic empty_iff
holomorphic_transform_insert_iff not_is_pole_holomorphic
open_delete r_ana that)
qed
then have g holomorphic_on p_balls
proof -
  have g analytic_on p_balls
    unfolding p_balls_def analytic_on_UN
    using g_ball by (simp add: analytic_on_open)
  moreover have open p_balls using p_balls_def by blast
  ultimately show ?thesis
    by (simp add: analytic_imp_holomorphic)
qed
moreover have g holomorphic_on D-pts
proof -

```

```

    have ( $\lambda z. \text{inverse } (f z)$ ) holomorphic_on  $D - \text{pts}$ 
      using f_holo holomorphic_on_inverse non_z by blast
    then show ?thesis
      by (metis DiffD2 g_def holomorphic_transform)
  qed
  moreover have open p_balls
    using p_balls_def by blast
  ultimately have g holomorphic_on ( $p\_balls \cup (D - \text{pts})$ )
    by (simp add: holomorphic_on_Un meromorphic_imp_open_diff[OF mer])
  moreover have  $D \subseteq p\_balls \cup (D - \text{pts})$ 
    unfolding p_balls_def using  $\langle \lambda z. z \in \text{pts} \implies 0 < \text{get\_r } z \rangle$  by force
  ultimately show g holomorphic_on D by (meson holomorphic_on_subset)
qed

lemma meromorphic_imp_analytic_on:
  assumes f meromorphic_on D pts
  shows f analytic_on ( $D - \text{pts}$ )
  by (metis assms analytic_on_open meromorphic_imp_open_diff meromorphic_on_def)

lemma meromorphic_imp_constant_on:
  assumes merf: f meromorphic_on D pts
    and f constant_on ( $D - \text{pts}$ )
    and  $\forall x \in \text{pts}. \text{is\_pole } f x$ 
  shows f constant_on D
proof -
  obtain c where  $c: \lambda z. z \in D - \text{pts} \implies f z = c$ 
    by (meson assms constant_on_def)

  have  $f z = c$  if  $z \in D$  for  $z$ 
  proof (cases is_pole f z)
    case True
    then obtain r0 where  $r0 > 0$  and r0: f analytic_on ball z r0 - {z} and
      pol: is_pole f z
      using merf unfolding meromorphic_on_def isolated_singularity_at_def
      by (metis  $\langle z \in D \rangle$  insert_Diff insert_Diff_if insert_iff merf
        meromorphic_imp_open_diff not_is_pole_holomorphic)
    have open D
      using merf meromorphic_on_def by auto
    then obtain r where  $r > 0$   $\text{ball } z r \subseteq D$   $r \leq r0$ 
      by (smt (verit, best)  $\langle 0 < r0 \rangle \langle z \in D \rangle$  openE order_subst2 subset_ball)
    have r: f analytic_on ball z r - {z}
      by (meson Diff_mono  $\langle r \leq r0 \rangle$  analytic_on_subset order_refl r0 subset_ball)
    have  $\text{ball } z r - \{z\} \subseteq -\text{pts}$ 
      using merf r unfolding meromorphic_on_def
      by (meson Compl Elementary_Metric_Spaces.open_ball
        analytic_imp_holomorphic assms(3) not_is_pole_holomorphic open_delete
        subsetI)
    with  $\langle \text{ball } z r \subseteq D \rangle$  have  $\text{ball } z r - \{z\} \subseteq D - \text{pts}$ 
      by fastforce

```

```

with  $c$  have  $c': \bigwedge u. u \in \text{ball } z \ r - \{z\} \implies f \ u = c$ 
  by blast
have False if  $\forall_F x \text{ in } \text{at } z. \text{cmod } c + 1 \leq \text{cmod } (f \ x)$ 
proof -
  have  $\forall_F x \text{ in } \text{at } z \text{ within } \text{ball } z \ r - \{z\}. \text{cmod } c + 1 \leq \text{cmod } (f \ x)$ 
  by (smt (verit, best) Diff_UNIV Diff_eq_empty_iff eventually_at_topological
insert_subset that)
  with  $\langle r > 0 \rangle$  show ?thesis
  apply (simp add: c' eventually_at_filter topological_space_class.eventually_nhds
open_dist)
  by (metis dist_commute min_less_iff_conj perfect_choose_dist)
qed
with pol show ?thesis
  by (auto simp: is_pole_def filterlim_at_infinity_conv_norm_at_top filter-
lim_at_top)
next
  case False
  then show ?thesis by (meson DiffI assms(3) c that)
qed
then show ?thesis
  by (simp add: constant_on_def)
qed

```

lemma *meromorphic_isolated*:

assumes *merf*: f *meromorphic_on* D *pts* **and** $p \in \text{pts}$
obtains r **where** $r > 0$ $\text{ball } p \ r \subseteq D$ $\text{ball } p \ r \cap \text{pts} = \{p\}$

proof -

have $\forall z \in D. \exists e > 0. \text{finite } (\text{pts} \cap \text{ball } z \ e)$
using *merf unfolding meromorphic_on_def islimpt_eq_infinite_ball*
by *auto*

then obtain $r0$ **where** $r0: r0 > 0$ $\text{finite } (\text{pts} \cap \text{ball } p \ r0)$
by (*metis assms(2) in_mono merf meromorphic_on_def*)
moreover define pts' **where** $\text{pts}' = \text{pts} \cap \text{ball } p \ r0 - \{p\}$
ultimately have $\text{finite } \text{pts}'$
by *simp*

define $r1$ **where** $r1 = (\text{if } \text{pts}' = \{\} \text{ then } r0 \text{ else}$
 $\min (\text{Min } \{\text{dist } p' \ p \mid p'. p' \in \text{pts}'\} / 2) \ r0)$

have $r1 > 0 \wedge \text{pts} \cap \text{ball } p \ r1 - \{p\} = \{\}$

proof (*cases pts' = {}*)

case *True*

then show *?thesis*

using *pts'_def r0(1) r1_def* **by** *presburger*

next

case *False*

define S **where** $S = \{\text{dist } p' \ p \mid p'. p' \in \text{pts}'\}$

have *nempty*: $S \neq \{\}$

```

    using False S_def by blast
  have finite:finite S
    using ⟨finite pts'⟩ S_def by simp

  have r1>0
  proof -
    have r1=min (Min S/2) r0
      using False unfolding S_def r1_def by auto
    moreover have Min S∈S
      using ⟨S≠{⟩ ⟨finite S⟩ Min_in by auto
    then have Min S>0 unfolding S_def
      using pts'_def by force
    ultimately show ?thesis using ⟨r0>0⟩ by auto
  qed
  moreover have pts ∩ ball p r1 - {p} = {}
  proof (rule ccontr)
    assume pts ∩ ball p r1 - {p} ≠ {}
    then obtain p' where p'∈pts ∩ ball p r1 - {p} by blast
    moreover have r1≤r0 using r1_def by auto
    ultimately have p'∈pts' unfolding pts'_def
      by auto
    then have dist p' p≥Min S
      using S_def eq_Min_iff local.finite by blast
    moreover have dist p' p < Min S
      using ⟨p'∈pts ∩ ball p r1 - {p}⟩ False unfolding r1_def
      apply (fold S_def)
      by (smt (verit, ccfv_threshold) DiffD1 Int_iff dist_commute
          dist_triangle_half_l mem_ball)
    ultimately show False by auto
  qed
  ultimately show ?thesis by auto
qed
then have r1>0 and r1_pts:pts ∩ ball p r1 - {p} = {} by auto

obtain r2 where r2>0 ball p r2 ⊆ D
  by (metis assms(2) merf meromorphic_on_def openE subset_eq)
define r where r=min r1 r2
have r > 0 unfolding r_def
  by (simp add: ⟨0 < r1⟩ ⟨0 < r2⟩)
moreover have ball p r ⊆ D
  using ⟨ball p r2 ⊆ D⟩ r_def by auto
moreover have ball p r ∩ pts = {p}
  using assms(2) ⟨r>0⟩ r1_pts
  unfolding r_def by auto
ultimately show ?thesis using that by auto
qed

lemma meromorphic_pts_closure:
  assumes merf: f meromorphic_on D pts

```

```

shows  $pts \subseteq \text{closure } (D - pts)$ 
proof -
  have  $p \text{ islimpt } (D - pts)$  if  $p \in pts$  for  $p$ 
  proof -
    obtain  $r$  where  $r > 0$   $\text{ball } p \ r \subseteq D$   $\text{ball } p \ r \cap pts = \{p\}$ 
      using  $\text{meromorphic\_isolated}[OF \text{ merf } \langle p \in pts \rangle]$  by auto
    from  $\langle r > 0 \rangle$ 
    have  $p \text{ islimpt } \text{ball } p \ r - \{p\}$ 
      by ( $\text{meson open\_ball ball\_subset\_cball in\_mono islimpt\_ball}$ 
         $\text{islimpt\_punctured le\_less open\_contains\_ball\_eq}$ )
    moreover have  $\text{ball } p \ r - \{p\} \subseteq D - pts$ 
      using  $\langle \text{ball } p \ r \cap pts = \{p\} \rangle \langle \text{ball } p \ r \subseteq D \rangle$  by fastforce
    ultimately show  $?thesis$ 
      using  $\text{islimpt\_subset}$  by auto
  qed
  then show  $?thesis$  by ( $\text{simp add: islimpt\_in\_closure subset\_eq}$ )
qed

lemma  $nconst\_imp\_nzero\_neighbour$ :
  assumes  $\text{merf: } f \text{ meromorphic\_on } D \ pts$ 
    and  $f\_nconst: \neg(\forall w \in D - pts. f \ w = 0)$ 
    and  $z \in D$  and  $\text{connected } D$ 
  shows  $(\forall_F w \text{ in at } z. f \ w \neq 0 \wedge w \in D - pts)$ 
  proof -
    obtain  $\beta$  where  $\beta: \beta \in D - pts$   $f \ \beta \neq 0$ 
      using  $f\_nconst$  by auto

  have  $?thesis$  if  $z \notin pts$ 
  proof -
    have  $\forall_F w \text{ in at } z. f \ w \neq 0 \wedge w \in D - pts$ 
      apply ( $\text{rule non\_zero\_neighbour\_alt}[of f D - pts z \ \beta]$ )
      subgoal using  $\text{merf meromorphic\_on\_def}$  by blast
      subgoal using  $\text{merf meromorphic\_imp\_open\_diff}$  by auto
      subgoal using  $\text{assms}(4)$   $\text{merf meromorphic\_imp\_connected\_diff}$  by blast
      subgoal by ( $\text{simp add: assms}(3)$ ) that)
      using  $\beta$  by auto
    then show  $?thesis$  by ( $\text{auto elim: eventually\_mono}$ )
  qed
  moreover have  $?thesis$  if  $z \in pts \neg f \ -z \rightarrow 0$ 
  proof -
    have  $\forall_F w \text{ in at } z. w \in D - pts$ 
      using  $\text{merf}[unfolding \text{ meromorphic\_on\_def islimpt\_iff\_eventually}] \langle z \in D \rangle$ 
      using  $\text{eventually\_at\_in\_open' eventually\_elim2}$  by fastforce
    moreover have  $\forall_F w \text{ in at } z. f \ w \neq 0$ 
    proof (cases  $\text{is\_pole } f \ z$ )
      case True
        then show  $?thesis$  using  $\text{non\_zero\_neighbour\_pole}$  by auto
    next
      case False

```



```

moreover have not_essential f z
  using merf meromorphic_on_def that(1) by fastforce
ultimately obtain c where  $c \neq 0$  f -z → c
  by (metis  $\langle \neg f -z \rightarrow 0 \rangle$  not_essential_def)
then show ?thesis
  using tendsto_imp_eventually_ne by auto
qed
ultimately show ?thesis by eventually_elim auto
qed
moreover have ?thesis if  $z \in \text{pts}$  f -z → 0
proof -
  define ff where  $ff = (\lambda x. \text{if } x = z \text{ then } 0 \text{ else } f x)$ 
  define A where  $A = D - \{z\}$ 

  have f holomorphic_on A - {z}
    by (metis A_def Diff_insert analytic_imp_holomorphic
      insert_Diff merf meromorphic_imp_analytic_on that(1))
  moreover have open A
    using A_def merf meromorphic_imp_open_diff' by force
  ultimately have ff holomorphic_on A
    using  $\langle f -z \rightarrow 0 \rangle$  unfolding ff_def
    by (rule removable_singularity)
  moreover have connected A
proof -
  have connected (D - pts)
    using assms(4) merf meromorphic_imp_connected_diff by auto
  moreover have  $D - \text{pts} \subseteq A$ 
    unfolding A_def by auto
  moreover have  $A \subseteq \text{closure } (D - \text{pts})$  unfolding A_def
    by (smt (verit, ccfv_SIG) Diff_empty Diff_insert
      closure_subset insert_Diff_single insert_absorb
      insert_subset merf meromorphic_pts_closure that(1))
  ultimately show ?thesis using connected_intermediate_closure
    by auto
qed
moreover have  $z \in A$  using A_def assms(3) by blast
moreover have  $ff z = 0$  unfolding ff_def by auto
moreover have  $\beta \in A$  using A_def  $\beta(1)$  by blast
moreover have  $ff \beta \neq 0$  using  $\beta(1)$   $\beta(2)$  ff_def that(1) by auto
ultimately obtain r where  $0 < r$ 
   $\text{ball } z r \subseteq A \wedge x. x \in \text{ball } z r - \{z\} \implies ff x \neq 0$ 
  using  $\langle \text{open } A \rangle$  isolated_zeros[of ff A z  $\beta$ ] by auto
then show ?thesis unfolding eventually_at ff_def
  by (intro exI[of _ r]) (auto simp: A_def dist_commute ball_def)
qed
ultimately show ?thesis by auto
qed

```

lemma *nconst_imp_nzero_neighbour'*:

```

assumes merf: f meromorphic_on D pts
and f_nconst:  $\neg(\forall w \in D - pts. f w = 0)$ 
and z  $\in D$  and connected D
shows  $\forall_F w$  in at z. f w  $\neq 0$ 
using nconst_imp_nzero_neighbour[OF assms]
by (auto elim:eventually_mono)

lemma meromorphic_compact_finite_zeros:
assumes merf: f meromorphic_on D pts
and compact S S  $\subseteq D$  connected D
and f_nconst:  $\neg(\forall w \in D - pts. f w = 0)$ 
shows finite ( $\{x \in S. f x = 0\}$ )
proof -
have finite ( $\{x \in S. f x = 0 \wedge x \notin pts\}$ )
proof (rule ccontr)
assume infinite  $\{x \in S. f x = 0 \wedge x \notin pts\}$ 
then obtain z where z  $\in S$  and z_lim: z islimpt  $\{x \in S. f x = 0$ 
 $\wedge x \notin pts\}$ 
using  $\langle compact S \rangle$  unfolding compact_eq_Bolzano_Weierstrass
by auto

from z_lim
have  $\exists_F x$  in at z. f x = 0  $\wedge x \in S \wedge x \notin pts$ 
unfolding islimpt_iff_eventually_not_eventually by simp
moreover have  $\forall_F w$  in at z. f w  $\neq 0 \wedge w \in D - pts$ 
using nconst_imp_nzero_neighbour[OF merf f_nconst _  $\langle connected D \rangle$ ]
 $\langle z \in S \rangle \langle S \subseteq D \rangle$ 
by auto
ultimately have  $\exists_F x$  in at z. False
by (simp add: eventually_mono frequently_def)
then show False by auto
qed
moreover have finite ( $S \cap pts$ )
using meromorphic_compact_finite_pts[OF merf  $\langle compact S \rangle \langle S \subseteq D \rangle$ ].
ultimately have finite ( $\{x \in S. f x = 0 \wedge x \notin pts\} \cup (S \cap pts)$ )
unfolding finite_Un by auto
then show ?thesis by (elim rev_finite_subset) auto
qed

lemma meromorphic_onI [intro?]:
assumes open A pts  $\subseteq A$ 
assumes f holomorphic_on A - pts  $\wedge z. z \in A \implies \neg z$  islimpt pts
assumes  $\wedge z. z \in pts \implies$  isolated_singularity_at f z
assumes  $\wedge z. z \in pts \implies$  not_essential f z
shows f meromorphic_on A pts
using assms unfolding meromorphic_on_def by blast

lemma Polygamma_plus_of_nat:
assumes  $\forall k < m. z \neq -of\_nat k$ 

```

```

shows Polygamma n (z + of_nat m) =
        Polygamma n z + (-1) ^ n * fact n * (∑ k<m. 1 / (z + of_nat k) ^
Suc n)
using assms
proof (induction m)
  case (Suc m)
  have Polygamma n (z + of_nat (Suc m)) = Polygamma n (z + of_nat m + 1)
    by (simp add: add_ac)
  also have ... = Polygamma n (z + of_nat m) + (-1) ^ n * fact n * (1 / ((z
+ of_nat m) ^ Suc n))
    using Suc.prem1 by (subst Polygamma_plus1) (auto simp: add_eq_0_iff2)
  also have Polygamma n (z + of_nat m) =
        Polygamma n z + (-1) ^ n * (∑ k<m. 1 / (z + of_nat k) ^ Suc n)
    * fact n
    using Suc.prem1 by (subst Suc.IH) auto
  finally show ?case
    by (simp add: algebra_simps)
qed auto

```

```

lemma tendsto_Gamma [tendsto_intros]:
  assumes (f ⟶ c) F c ∉ ℤ≤0
  shows ((λz. Gamma (f z)) ⟶ Gamma c) F
  by (intro isCont_tendsto_compose[OF _ assms(1)] continuous_intros assms)

```

```

lemma tendsto_Polygamma [tendsto_intros]:
  fixes f :: _ ⇒ 'a :: {real_normed_field, euclidean_space}
  assumes (f ⟶ c) F c ∉ ℤ≤0
  shows ((λz. Polygamma n (f z)) ⟶ Polygamma n c) F
  by (intro isCont_tendsto_compose[OF _ assms(1)] continuous_intros assms)

```

```

lemma analytic_on_Gamma' [analytic_intros]:
  assumes f analytic_on A ∀ x∈A. f x ∉ ℤ≤0
  shows (λz. Gamma (f z)) analytic_on A
  using analytic_on_compose_gen[OF assms(1) analytic_Gamma[of f 'A]] assms(2)
  by (auto simp: o_def)

```

```

lemma analytic_on_Polygamma' [analytic_intros]:
  assumes f analytic_on A ∀ x∈A. f x ∉ ℤ≤0
  shows (λz. Polygamma n (f z)) analytic_on A
  using analytic_on_compose_gen[OF assms(1) analytic_on_Polygamma[of f '
A n]] assms(2)
  by (auto simp: o_def)

```

```

lemma
  shows is_pole_Polygamma: is_pole (Polygamma n) (-of_nat m :: complex)
  and zorder_Polygamma: zorder (Polygamma n) (-of_nat m) = -int (Suc n)
  and residue_Polygamma: residue (Polygamma n) (-of_nat m) = (if n = 0
then -1 else 0)
proof -

```

```

define  $g1 :: \text{complex} \Rightarrow \text{complex}$  where
   $g1 = (\lambda z. \text{Polygamma } n (z + \text{of\_nat } (\text{Suc } m)) +$ 
     $(-1) \wedge \text{Suc } n * \text{fact } n * (\sum k < m. 1 / (z + \text{of\_nat } k) \wedge \text{Suc } n))$ 
define  $g :: \text{complex} \Rightarrow \text{complex}$  where
   $g = (\lambda z. g1 z + (-1) \wedge \text{Suc } n * \text{fact } n / (z + \text{of\_nat } m) \wedge \text{Suc } n)$ 
define  $F$  where  $F = \text{fps\_to\_fls } (\text{fps\_expansion } g1 (-\text{of\_nat } m)) + \text{fls\_const}$ 
 $((-1) \wedge \text{Suc } n * \text{fact } n) / \text{fls\_X} \wedge \text{Suc } n$ 
have  $F\_altdef: F = \text{fps\_to\_fls } (\text{fps\_expansion } g1 (-\text{of\_nat } m)) + \text{fls\_shift}$ 
 $(n+1) (\text{fls\_const } ((-1) \wedge \text{Suc } n * \text{fact } n))$ 
by (simp add: F_def del: power_Suc)

have  $\neg(-\text{of\_nat } m) \text{ islimpt } (\mathbb{Z}_{\leq 0} :: \text{complex set})$ 
by (intro discrete_imp_not_islimpt[where e = 1])
(auto elim!: nonpos_Ints_cases simp: dist_of_int)
hence eventually  $(\lambda z :: \text{complex}. z \notin \mathbb{Z}_{\leq 0})$  (at  $(-\text{of\_nat } m)$ )
by (auto simp: islimpt_conv_frequently_at frequently_def)
hence  $ev: \text{eventually } (\lambda z. \text{Polygamma } n z = g z)$  (at  $(-\text{of\_nat } m)$ )
proof eventually_elim
case (elim z)
hence  $*$ :  $\forall k < \text{Suc } m. z \neq -\text{of\_nat } k$ 
by auto
thus ?case
using Polygamma_plus_of_nat[of Suc m z n, OF *]
by (auto simp: g_def g1_def algebra_simps)
qed

have  $(\lambda w. g (-\text{of\_nat } m + w)) \text{ has\_laurent\_expansion } F$ 
unfolding  $g\_def F\_def$ 
by (intro laurent_expansion_intros has_laurent_expansion_fps analytic_at_imp_has_fps_expansion)
(auto simp: g1_def intro!: laurent_expansion_intros analytic_intros)
also have ?this  $\longleftrightarrow (\lambda w. \text{Polygamma } n (-\text{of\_nat } m + w)) \text{ has\_laurent\_expansion}$ 
 $F$ 
using  $ev$  by (intro has_laurent_expansion_cong refl)
(simp_all add: eq_commute at_to_0' eventually_filtermap)
finally have  $*$ :  $(\lambda w. \text{Polygamma } n (-\text{of\_nat } m + w)) \text{ has\_laurent\_expansion}$ 
 $F$  .

have  $\text{subdegree: fls\_subdegree } F = -\text{int } (\text{Suc } n)$  unfolding  $F\_def$ 
by (subst fls_subdegree_add_eq2) (simp_all add: fls_subdegree_fls_to_fps)
 $\text{fls\_divide\_subdegree}$ )
have [simp]:  $F \neq 0$ 
using  $\text{subdegree}$  by auto

show is_pole  $(\text{Polygamma } n) (-\text{of\_nat } m :: \text{complex})$ 
using  $*$  by (rule has_laurent_expansion_imp_is_pole) (auto simp: subdegree)
show  $\text{zorder } (\text{Polygamma } n) (-\text{of\_nat } m :: \text{complex}) = -\text{int } (\text{Suc } n)$ 
by (subst has_laurent_expansion_zorder[OF *]) (auto simp: subdegree)
show  $\text{residue } (\text{Polygamma } n) (-\text{of\_nat } m :: \text{complex}) = (\text{if } n = 0 \text{ then } -1 \text{ else } 0)$ 

```

by (subst has_laurent_expansion_residue[OF *]) (auto simp: F_altdef)
qed

lemma Gamma_meromorphic_on [meromorphic_intros]: Gamma meromorphic_on UNIV $\mathbb{Z}_{\leq 0}$

proof

show $\neg z$ islimpt $\mathbb{Z}_{\leq 0}$ for $z :: \text{complex}$

by (intro discrete_imp_not_islimpt[of 1]) (auto elim!: nonpos_Ints_cases simp: dist_of_int)

next

fix $z :: \text{complex}$ assume $z: z \in \mathbb{Z}_{\leq 0}$

then obtain n where $n: z = -\text{of_nat } n$

by (elim nonpos_Ints_cases[^])

show not_essential Gamma z

by (auto simp: n intro!: is_pole_imp_not_essential is_pole_Gamma)

have *: open $(-\mathbb{Z}_{\leq 0} - \{z\})$

by (intro open_Cmpl discrete_imp_closed[of 1]) (auto elim!: nonpos_Ints_cases simp: dist_of_int)

have Gamma holomorphic_on $(-\mathbb{Z}_{\leq 0} - \{z\}) - \{z\}$

by (intro holomorphic_intros) auto

thus isolated_singularity_at Gamma z

by (rule isolated_singularity_at_holomorphic) (use $z * \text{in auto}$)

qed (auto intro!: holomorphic_intros)

lemma Polygamma_meromorphic_on [meromorphic_intros]: Polygamma n meromorphic_on UNIV $\mathbb{Z}_{\leq 0}$

proof

show $\neg z$ islimpt $\mathbb{Z}_{\leq 0}$ for $z :: \text{complex}$

by (intro discrete_imp_not_islimpt[of 1]) (auto elim!: nonpos_Ints_cases simp: dist_of_int)

next

fix $z :: \text{complex}$ assume $z: z \in \mathbb{Z}_{\leq 0}$

then obtain m where $n: z = -\text{of_nat } m$

by (elim nonpos_Ints_cases[^])

show not_essential (Polygamma n) z

by (auto simp: n intro!: is_pole_imp_not_essential is_pole_Polygamma)

have *: open $(-\mathbb{Z}_{\leq 0} - \{z\})$

by (intro open_Cmpl discrete_imp_closed[of 1]) (auto elim!: nonpos_Ints_cases simp: dist_of_int)

have Polygamma n holomorphic_on $(-\mathbb{Z}_{\leq 0} - \{z\}) - \{z\}$

by (intro holomorphic_intros) auto

thus isolated_singularity_at (Polygamma n) z

by (rule isolated_singularity_at_holomorphic) (use $z * \text{in auto}$)

qed (auto intro!: holomorphic_intros)

theorem argument_principle':

fixes $f :: \text{complex} \Rightarrow \text{complex}$ and poles $s :: \text{complex set}$

— pz is the set of non-essential singularities and zeros

```

defines pz  $\equiv$  {w $\in$ s. f w = 0  $\vee$  w  $\in$  poles}
assumes open s and
  connected s and
  f_holo:f holomorphic_on s-poles and
  h_holo:h holomorphic_on s and
  valid_path g and
  loop:pathfinish g = pathstart g and
  path_img:path_image g  $\subseteq$  s - pz and
  homo: $\forall$  z. (z  $\notin$  s)  $\longrightarrow$  winding_number g z = 0 and
  finite:finite pz and
  poles: $\forall$  p $\in$ s $\cap$ poles. not_essential f p
shows contour_integral g ( $\lambda$ x. deriv f x * h x / f x) = 2 * pi * i *
  ( $\sum$  p $\in$ pz. winding_number g p * h p * zorder f p)
proof -
  define ff where ff = remove_sings f

  have finite':finite (s  $\cap$  poles)
    using finite unfolding pz_def by (auto elim:rev_finite_subset)

  have isolated:isolated_singularity_at f z if z $\in$ s for z
  proof (rule isolated_singularity_at_holomorphic)
    show f holomorphic_on (s-(poles-{z})) - {z}
      by (metis Diff_empty Diff_insert Diff_insert0 Diff_subset
        f_holo holomorphic_on_subset insert_Diff)
    show open (s - (poles - {z}))
      by (metis Diff_Diff_Int Int_Diff assms(2) finite' finite_Diff
        finite_imp_closed inf.idem open_Diff)
    show z  $\in$  s - (poles - {z})
      using assms(4) that by auto
  qed

  have not_ess:not_essential f w if w $\in$ s for w
    by (metis Diff_Diff_Int Diff_iff Int_Diff Int_absorb assms(2)
      f_holo finite' finite_imp_closed not_essential_holomorphic
      open_Diff poles that)

  have nzero: $\forall$  x in at w. f x  $\neq$  0 if w $\in$ s for w
  proof (rule ccontr)
    assume  $\neg$  ( $\forall$  x in at w. f x  $\neq$  0)
    then have  $\exists$  x in at w. f x = 0
      unfolding not_eventually by simp
    moreover have  $\forall$  x in at w. x $\in$ s
      by (simp add: assms(2) eventually_at_in_open' that)
    ultimately have  $\exists$  x in at w. x $\in$ {w $\in$ s. f w = 0}
      apply (elim frequently_rev_mp)
      by (auto elim:eventually_mono)
    from frequently_at_imp_islimpt[OF this]
    have w islimpt {w  $\in$  s. f w = 0} .
    then have infinite({w  $\in$  s. f w = 0}  $\cap$  ball w 1)

```

```

    unfolding islimpt_eq_infinite_ball by auto
  then have infinite({w ∈ s. f w = 0})
    by auto
  then have infinite pz unfolding pz_def
    by (smt (verit) Collect_mono_iff rev_finite_subset)
  then show False using finite by auto
qed

obtain pts' where pts':pts' ⊆ s ∩ poles
  finite pts' ff holomorphic_on s - pts' ∀ x∈pts'. is_pole ff x
  apply (elim get_all_poles_from_remove_sings
    [of f, folded ff_def, rotated -1])
  subgoal using f_holo by fastforce
  using ⟨open s⟩ poles finite' by auto

have pts'_sub_pz:{w ∈ s. ff w = 0 ∨ w ∈ pts'} ⊆ pz
proof -
  have w∈poles if w∈s w∈pts' for w
    by (meson in_mono le_infE pts'(1) that(2))
  moreover have f w=0 if w∈s w∉poles ff w=0 for w
  proof -
    have ¬ is_pole f w
      by (metis DiffI Diff_Diff_Int Diff_subset assms(2) f_holo
        finite' finite_imp_closed inf.absorb_iff2
        not_is_pole_holomorphic_open_Diff that(1) that(2))
    then have f -w → 0
      using remove_sings_eq_0_iff[OF not_ess[OF ⟨w∈s⟩]] ⟨ff w=0⟩
      unfolding ff_def by auto
    moreover have f analytic_on {w}
      using that(1,2) finite' f_holo assms(2)
      by (metis Diff_Diff_Int Diff_empty Diff_iff Diff_subset
        double_diff finite_imp_closed
        holomorphic_on_imp_analytic_at open_Diff)
    ultimately show ?thesis
      using ff_def remove_sings_at_analytic that(3) by presburger
  qed
  ultimately show ?thesis unfolding pz_def by auto
qed

have contour_integral g (λx. deriv f x * h x / f x)
  = contour_integral g (λx. deriv ff x * h x / ff x)
proof (rule contour_integral_eq)
  fix x assume x ∈ path_image g
  have f analytic_on {x}
  proof (rule holomorphic_on_imp_analytic_at[of _ s-poles])
    from finite'
    show open (s - poles)
      using ⟨open s⟩

```

```

    by (metis Diff_Compl Diff_Diff_Int Diff_eq finite_imp_closed
        open_Diff)
  show  $x \in s - \text{poles}$ 
    using path_img  $\langle x \in \text{path\_image } g \rangle$  unfolding pz_def by auto
qed (use f_holo in simp)
then show  $\text{deriv } f x * h x / f x = \text{deriv } ff x * h x / ff x$ 
  unfolding ff_def by auto
qed
also have ... =  $\text{complex\_of\_real } (2 * \pi) * i *$ 
   $(\sum_{p \in \{w \in s. ff w = 0 \vee w \in pts'\}} \text{winding\_number } g p * h p * \text{of\_int } (zorder ff p))$ 
proof (rule argument_principle[OF  $\langle \text{open } s \rangle \langle \text{connected } s \rangle$ , of ff pts' h g])
  show  $\text{path\_image } g \subseteq s - \{w \in s. ff w = 0 \vee w \in pts'\}$ 
    using path_img pts'_sub_pz by auto
  show  $\text{finite } \{w \in s. ff w = 0 \vee w \in pts'\}$ 
    using pts'_sub_pz finite
    using rev_finite_subset by blast
qed (use pts'_assms in auto)
also have ... =  $2 * \pi * i *$ 
   $(\sum_{p \in pz. \text{winding\_number } g p * h p * zorder f p}$ 
proof -
  have  $(\sum_{p \in \{w \in s. ff w = 0 \vee w \in pts'\}} \text{winding\_number } g p * h p * \text{of\_int } (zorder ff p)) =$ 
   $(\sum_{p \in pz. \text{winding\_number } g p * h p * \text{of\_int } (zorder f p))$ 
proof (rule sum.mono_neutral_cong_left)
  have  $zorder f w = 0$ 
    if  $w \in s \wedge f w = 0 \vee w \in \text{poles} \wedge ff w \neq 0 \wedge w \notin pts'$ 
    for w
  proof -
    define F where  $F = \text{laurent\_expansion } f w$ 
    have  $\text{has\_l} : (\lambda x. f (w + x)) \text{ has\_laurent\_expansion } F$ 
      unfolding F_def
      apply (rule not_essential_has_laurent_expansion)
      using isolated_not_ess  $\langle w \in s \rangle$  by auto
    from  $\text{has\_laurent\_expansion\_eventually\_nonzero\_iff}$ [OF this]
    have  $F \neq 0$ 
      using nzero  $\langle w \in s \rangle$  by auto
    from  $\text{tendsto\_0\_subdegree\_iff}$ [OF has_l this]
    have  $f -w \rightarrow 0 = (0 < fls\_subdegree F)$  .
    moreover have  $\neg (\text{is\_pole } f w \vee f -w \rightarrow 0)$ 
      using  $\text{remove\_sings\_eq\_0\_iff}$ [OF not_ess[OF  $\langle w \in s \rangle$ ]]  $\langle ff w \neq 0 \rangle$ 
      unfolding ff_def by auto
    moreover have  $\text{is\_pole } f w = (fls\_subdegree F < 0)$ 
      using  $\text{is\_pole\_fls\_subdegree\_iff}$ [OF has_l] .
    ultimately have  $fls\_subdegree F = 0$  by auto
    then show ?thesis
      using  $\text{has\_laurent\_expansion\_zorder}$ [OF has_l  $\langle F \neq 0 \rangle$ ] by auto
  qed
then show  $\forall i \in pz - \{w \in s. ff w = 0 \vee w \in pts'\}$ .

```



```

    winding_number g i * h i * of_int (zorder f i) = 0
  unfolding pz_def by auto
  show  $\bigwedge x. x \in \{w \in s. \text{ff } w = 0 \vee w \in \text{pts}'\} \implies$ 
    winding_number g x * h x * of_int (zorder ff x) =
    winding_number g x * h x * of_int (zorder f x)
  using isolated zorder_remove_sings[of f, folded ff_def] by auto
  qed (use pts'_sub_pz finite in auto)
  then show ?thesis by auto
qed
finally show ?thesis .
qed

```

lemma *meromorphic_on_imp_isolated_singularity*:

assumes *f meromorphic_on D pts z ∈ D*

shows *isolated_singularity_at f z*

by (*meson DiffI assms(1) assms(2) holomorphic_on_imp_analytic_at_isolated_singularity_at_analytic*

meromorphic_imp_open_diff meromorphic_on_def)

lemma *meromorphic_imp_not_is_pole*:

assumes *f meromorphic_on D pts z ∈ D - pts*

shows *¬is_pole f z*

proof -

from *assms* **have** *f analytic_on {z}*

using *meromorphic_on_imp_analytic_at* **by** *blast*

thus *?thesis*

using *analytic_at_not_is_pole_holomorphic* **by** *blast*

qed

lemma *meromorphic_all_poles_iff_empty* [*simp*]: *f meromorphic_on pts pts \longleftrightarrow pts = {}*

by (*auto simp: meromorphic_on_def holomorphic_on_def open_imp_islimpt*)

lemma *meromorphic_imp_nonsingular_point_exists*:

assumes *f meromorphic_on A pts A \neq {}*

obtains *x* **where** *x ∈ A - pts*

proof -

have *A \neq pts*

using *assms* **by** *auto*

moreover **have** *pts \subseteq A*

using *assms* **by** (*auto simp: meromorphic_on_def*)

ultimately **show** *?thesis*

using *that* **by** *blast*

qed

lemma *meromorphic_frequently_const_imp_const*:

assumes *f meromorphic_on A pts connected A*

assumes *frequently ($\lambda w. f w = c$) (at z)*

assumes *z ∈ A - pts*

```

assumes  $w \in A - pts$ 
shows  $f w = c$ 
proof -
  have  $f w - c = 0$ 
  proof (rule analytic_continuation[where  $f = \lambda z. f z - c$ ])
    show  $(\lambda z. f z - c)$  holomorphic_on  $(A - pts)$ 
      by (intro holomorphic_intros meromorphic_imp_holomorphic[OF assms(1)])
    show [intro]: open  $(A - pts)$ 
      using assms meromorphic_imp_open_diff by blast
    show connected  $(A - pts)$ 
      using assms meromorphic_imp_connected_diff by blast
    show  $\{z \in A - pts. f z = c\} \subseteq A - pts$ 
      by blast
    have eventually  $(\lambda z. z \in A - pts)$  (at  $z$ )
      using assms by (intro eventually_at_in_open') auto
    hence frequently  $(\lambda z. f z = c \wedge z \in A - pts)$  (at  $z$ )
      by (intro frequently_eventually_frequently assms)
    thus  $z$  islimpt  $\{z \in A - pts. f z = c\}$ 
      by (simp add: islimpt_conv_frequently_at_conj_commute)
  qed (use assms in auto)
thus ?thesis
  by simp
qed

```

lemma meromorphic_imp_eventually_neq:

```

assumes  $f$  meromorphic_on  $A$   $pts$  connected  $A$   $\neg f$  constant_on  $A - pts$ 
assumes  $z \in A - pts$ 
shows eventually  $(\lambda z. f z \neq c)$  (at  $z$ )
proof (rule ccontr)
  assume  $\neg$ eventually  $(\lambda z. f z \neq c)$  (at  $z$ )
  hence *: frequently  $(\lambda z. f z = c)$  (at  $z$ )
    by (auto simp: frequently_def)
  have  $\forall w \in A - pts. f w = c$ 
    using meromorphic_frequently_const_imp_const [OF assms(1,2) * assms(4)]
by blast
  hence  $f$  constant_on  $A - pts$ 
    by (auto simp: constant_on_def)
  thus False
    using assms(3) by contradiction
qed

```

lemma meromorphic_frequently_const_imp_const':

```

assumes  $f$  meromorphic_on  $A$   $pts$  connected  $A$   $\forall w \in pts. is\_pole f w$ 
assumes frequently  $(\lambda w. f w = c)$  (at  $z$ )
assumes  $z \in A$ 
assumes  $w \in A$ 
shows  $f w = c$ 
proof -
  have  $\neg is\_pole f z$ 

```

```

    using frequently_const_imp_not_is_pole[OF assms(4)] .
  with assms have z: z ∈ A - pts
  by auto
  have *: f w = c if w ∈ A - pts for w
  using that meromorphic_frequently_const_imp_const [OF assms(1,2,4) z] by
  auto
  have ¬is_pole f u if u ∈ A for u
  proof -
    have is_pole f u ↔ is_pole (λ_. c) u
    proof (rule is_pole_cong)
      have eventually (λw. w ∈ A - (pts - {u}) - {u}) (at u)
      by (intro eventually_at_in_open meromorphic_imp_open_diff' [OF assms(1)])
      (use that in auto)
      thus eventually (λw. f w = c) (at u)
      by eventually_elim (use * in auto)
    qed auto
  thus ?thesis
  by auto
qed
moreover have pts ⊆ A
  using assms(1) by (simp add: meromorphic_on_def)
ultimately have pts = {}
  using assms(3) by auto
with * and ⟨w ∈ A⟩ show ?thesis
  by blast
qed

```

```

lemma meromorphic_imp_eventually_neq':
  assumes f meromorphic_on A pts connected A ∀ w ∈ pts. is_pole f w ¬f constant_on A
  assumes z ∈ A
  shows eventually (λz. f z ≠ c) (at z)
proof (rule ccontr)
  assume ¬eventually (λz. f z ≠ c) (at z)
  hence *: frequently (λz. f z = c) (at z)
  by (auto simp: frequently_def)
  have ∀ w ∈ A. f w = c
  using meromorphic_frequently_const_imp_const' [OF assms(1,2,3) * assms(5)]
  by blast
  hence f constant_on A
  by (auto simp: constant_on_def)
  thus False
  using assms(4) by contradiction
qed

```

```

lemma zorder_eq_0_iff_meromorphic:
  assumes f meromorphic_on A pts ∀ z ∈ pts. is_pole f z z ∈ A
  assumes eventually (λx. f x ≠ 0) (at z)
  shows zorder f z = 0 ↔ ¬is_pole f z ∧ f z ≠ 0

```

```

proof (cases  $z \in pts$ )
  case True
    from assms obtain  $F$  where  $F$ :  $(\lambda x. f (z + x))$  has_laurent_expansion  $F$ 
      by (metis True meromorphic_on_def not_essential_has_laurent_expansion)
    from  $F$  and assms(4) have [simp]:  $F \neq 0$ 
      using has_laurent_expansion_eventually_nonzero_iff by blast
    show ?thesis using True assms(2)
      using is_pole_fls_subdegree_iff [OF  $F$ ] has_laurent_expansion_zorder [OF
 $F$ ]
      by auto
  next
    case False
      have ana:  $f$  analytic_on  $\{z\}$ 
        using meromorphic_on_imp_analytic_at False assms by blast
      hence  $\neg$ is_pole  $f z$ 
        using analytic_at_not_is_pole_holomorphic by blast
      moreover have frequently  $(\lambda w. f w \neq 0)$  (at  $z$ )
        using assms(4) by (intro eventually_frequently) auto
      ultimately show ?thesis using zorder_eq_0_iff[OF ana] False
        by auto
qed

```

```

lemma zorder_pos_iff_meromorphic:
  assumes  $f$  meromorphic_on  $A$  pts  $\forall z \in pts. is\_pole\ f\ z\ z \in A$ 
  assumes eventually  $(\lambda x. f x \neq 0)$  (at  $z$ )
  shows  $zorder\ f\ z > 0 \iff \neg is\_pole\ f\ z \wedge f\ z = 0$ 
proof (cases  $z \in pts$ )
  case True
    from assms obtain  $F$  where  $F$ :  $(\lambda x. f (z + x))$  has_laurent_expansion  $F$ 
      by (metis True meromorphic_on_def not_essential_has_laurent_expansion)
    from  $F$  and assms(4) have [simp]:  $F \neq 0$ 
      using has_laurent_expansion_eventually_nonzero_iff by blast
    show ?thesis using True assms(2)
      using is_pole_fls_subdegree_iff [OF  $F$ ] has_laurent_expansion_zorder [OF
 $F$ ]
      by auto
  next
    case False
      have ana:  $f$  analytic_on  $\{z\}$ 
        using meromorphic_on_imp_analytic_at False assms by blast
      hence  $\neg$ is_pole  $f z$ 
        using analytic_at_not_is_pole_holomorphic by blast
      moreover have frequently  $(\lambda w. f w \neq 0)$  (at  $z$ )
        using assms(4) by (intro eventually_frequently) auto
      ultimately show ?thesis using zorder_pos_iff'[OF ana] False
        by auto
qed

```

```

lemma zorder_neg_iff_meromorphic:

```

```

assumes  $f$  meromorphic_on  $A$  pts  $\forall z \in \text{pts. is\_pole } f z z \in A$ 
assumes eventually  $(\lambda x. f x \neq 0)$  (at  $z$ )
shows  $\text{zorder } f z < 0 \iff \text{is\_pole } f z$ 
proof -
  have frequently  $(\lambda x. f x \neq 0)$  (at  $z$ )
    using  $\text{assms}$  by (intro eventually_frequently) auto
  moreover from  $\text{assms}$  have isolated_singularity_at  $f z$  not_essential  $f z$ 
    using meromorphic_on_imp_isolated_singularity meromorphic_on_imp_not_essential
by blast+
  ultimately show ?thesis
    using isolated_pole_imp_neg_zorder neg_zorder_imp_is_pole by blast
qed

```

```

lemma meromorphic_on_imp_discrete:
  assumes  $\text{mero}: f$  meromorphic_on  $S$  pts and connected  $S$ 
    and  $\text{nconst}: \neg (\forall w \in S - \text{pts. } f w = c)$ 
  shows discrete  $\{x \in S. f x = c\}$ 
proof -
  define  $g$  where  $g = (\lambda x. f x - c)$ 
  have  $\forall_F w$  in at  $z. g w \neq 0$  if  $z \in S$  for  $z$ 
  proof (rule nconst_imp_nzero_neighbour'[of  $g$   $S$  pts  $z$ ])
    show  $g$  meromorphic_on  $S$  pts using  $\text{mero}$  unfolding  $g\_def$ 
    by (auto intro: meromorphic_intros)
  show  $\neg (\forall w \in S - \text{pts. } g w = 0)$  using  $\text{nconst}$  unfolding  $g\_def$  by auto
qed fact+
  then show ?thesis
    unfolding discrete_altdef  $g\_def$ 
    using eventually_mono by fastforce
qed

```

```

lemma meromorphic_isolated_in:
  assumes  $\text{merf}: f$  meromorphic_on  $D$  pts  $p \in \text{pts}$ 
  shows  $p$  isolated_in pts
  by (meson  $\text{assms}$  isolated_in_islimpt_iff meromorphic_on_def subsetD)

```

```

lemma remove_sings_constant_on:
  assumes  $\text{merf}: f$  meromorphic_on  $D$  pts and connected  $D$ 
    and  $\text{const}: f$  constant_on  $(D - \text{pts})$ 
  shows (remove_sings  $f$ ) constant_on  $D$ 
proof -
  have remove_sings_const: remove_sings  $f$  constant_on  $D - \text{pts}$ 
    using const
  by (metis constant_onE merf meromorphic_on_imp_analytic_at remove_sings_at_analytic)

  have ?thesis if  $D = \{\}$ 
    using that unfolding constant_on_def by auto
  moreover have ?thesis if  $D \neq \{\}$   $\{x \in \text{pts. is\_pole } f x\} = \{\}$ 
proof -
  obtain  $\xi$  where  $\xi \in (D - \text{pts})$   $\xi$  islimpt  $(D - \text{pts})$ 

```

```

proof –
  have open (D – pts)
    using meromorphic_imp_open_diff[OF merf] .
  moreover have (D – pts) ≠ {} using ‹D≠{}›
    by (metis Diff_empty_closure_empty merf
      meromorphic_pts_closure_subset_empty)
  ultimately show ?thesis using open_imp_islimpt that by auto
qed
moreover have remove_sings f holomorphic_on D
  using remove_sings_holomorphic_on[OF merf] that by auto
moreover note remove_sings_const
moreover have open D
  using assms(1) meromorphic_on_def by blast
ultimately show ?thesis
  using Conformal_Mappings.analytic_continuation'
    [of remove_sings f D D–pts ξ] ‹connected D›
  by auto
qed
moreover have ?thesis if D≠{} {x∈pts. is_pole f x} ≠ {}
proof –
  define PP where PP={x∈D. is_pole f x}
  have remove_sings f meromorphic_on D PP
    using merf unfolding PP_def
  apply (elim remove_sings_meromorphic_on)
  subgoal using assms(1) meromorphic_on_def by force
  subgoal using meromorphic_pole_subset merf by auto
  done
moreover have remove_sings f constant_on D – PP
proof –
  obtain ξ where ξ ∈ f ‘ (D – pts)
    by (metis Diff_empty Diff_eq_empty_iff ‹D ≠ {}› assms(1)
      closure_empty_ex_in_conv_imageI meromorphic_pts_closure)
  have ξ:∀ x∈D – pts. f x = ξ
    by (metis ‹ξ ∈ f ‘ (D – pts)› assms(3) constant_on_def image_iff)

  have remove_sings f x = ξ if x∈D – PP for x
  proof (cases x∈pts)
  case True
  then have x isolated_in pts
    using meromorphic_isolated_in[OF merf] by auto
  then obtain T0 where T0:open T0 T0 ∩ pts = {x}
    unfolding isolated_in_def by auto
  obtain T1 where T1:open T1 x∈T1 T1 ⊆ D
    using merf unfolding meromorphic_on_def
    using True by blast
  define T2 where T2 = T1 ∩ T0
  have open T2 x∈T2 T2 – {x} ⊆ D – pts
    using T0 T1 unfolding T2_def by auto
  then have ∀ w∈T2. w≠x → f w =ξ

```

```

    using  $\xi$  by auto
  then have  $\forall_F x$  in at  $x$ .  $f x = \xi$ 
    unfolding eventually_at_topological
    using  $\langle \text{open } T2 \rangle \langle x \in T2 \rangle$  by auto
  then have  $f -x \rightarrow \xi$ 
    using tendsto_eventually by auto
  then show ?thesis by blast
next
case False
then show ?thesis
  using  $\langle \forall x \in D - \text{pts}. f x = \xi \rangle$  assms(1)
  meromorphic_on_imp_analytic_at that by auto
qed

  then show ?thesis unfolding constant_on_def by auto
qed

moreover have is_pole (remove_sings f)  $x$  if  $x \in PP$  for  $x$ 
proof -
  have isolated_singularity_at f  $x$ 
    by (metis (mono_tags, lifting) DiffI PP_def assms(1)
        isolated_singularity_at_analytic mem_Collect_eq
        meromorphic_on_def meromorphic_on_imp_analytic_at that)
  then show ?thesis using that unfolding PP_def by simp
qed
ultimately show ?thesis
  using meromorphic_imp_constant_on
  [of remove_sings f D PP]
  by auto
qed
ultimately show ?thesis by auto
qed

lemma meromorphic_eq_meromorphic_extend:
  assumes  $f$  meromorphic_on  $A$  pts1  $g$  meromorphic_on  $A$  pts1  $\neg z$  islimpt pts2
  assumes  $\bigwedge z. z \in A - \text{pts2} \implies f z = g z$  pts1  $\subseteq$  pts2  $z \in A - \text{pts1}$ 
  shows  $f z = g z$ 
proof -
  have  $g$  analytic_on  $\{z\}$ 
    using assms by (intro meromorphic_on_imp_analytic_at[OF assms(2)]) auto
  hence  $g -z \rightarrow g z$ 
    using analytic_at_imp_isCont isContD by blast
  also have ?this  $\iff f -z \rightarrow g z$ 
proof (intro filterlim_cong)
  have eventually  $(\lambda w. w \notin \text{pts2})$  (at  $z$ )
    using assms by (auto simp: islimpt_conv_frequently_at frequently_def)
  moreover have eventually  $(\lambda w. w \in A)$  (at  $z$ )
    using assms by (intro eventually_at_in_open') (auto simp: meromor-
    phic_on_def)

```

```

ultimately show  $\forall_F x$  in at  $z$ .  $g x = f x$ 
  by eventually_elim (use assms in auto)
qed auto
finally have  $f -z \rightarrow g z$  .
moreover have  $f$  analytic_on  $\{z\}$ 
  using assms by (intro meromorphic_on_imp_analytic_at[OF assms(1)]) auto
hence  $f -z \rightarrow f z$ 
  using analytic_at_imp_isCont isContD by blast
ultimately show ?thesis
  using tendsto_unique by force
qed

```

lemma meromorphic_constant_on_extend:

```

assumes  $f$  constant_on  $A - pts1$   $f$  meromorphic_on  $A$   $pts1$   $f$  meromorphic_on
 $A$   $pts2$   $pts2 \subseteq pts1$ 
shows  $f$  constant_on  $A - pts2$ 
proof -
  from assms(1) obtain  $c$  where  $c: \bigwedge z. z \in A - pts1 \implies f z = c$ 
  unfolding constant_on_def by auto
  have  $f z = c$  if  $z \in A - pts2$  for  $z$ 
  using assms(3)
  proof (rule meromorphic_eq_meromorphic_extend[where  $z = z$ ])
    show  $(\lambda a. c)$  meromorphic_on  $A$   $pts2$ 
      by (intro meromorphic_on_const) (use assms in <auto simp: meromor-
        phic_on_def>)
    show  $\neg z$  islimpt  $pts1$ 
      using that assms by (auto simp: meromorphic_on_def)
  qed (use assms  $c$  that in auto)
  thus ?thesis
    by (auto simp: constant_on_def)
qed

```

lemma meromorphic_remove_sings_constant_on_imp_constant_on:

```

assumes  $f$  meromorphic_on  $A$   $pts$ 
assumes remove_sings  $f$  constant_on  $A$ 
shows  $f$  constant_on  $A - pts$ 
proof -
  from assms(2) obtain  $c$  where  $c: \bigwedge z. z \in A \implies$  remove_sings  $f z = c$ 
  by (auto simp: constant_on_def)
  have  $f z = c$  if  $z \in A - pts$  for  $z$ 
  using meromorphic_on_imp_analytic_at[OF assms(1) that]  $c$ [of  $z$ ] that
  by auto
  thus ?thesis
    by (auto simp: constant_on_def)
qed

```


definition *singularities_on* :: complex set \Rightarrow (complex \Rightarrow complex) \Rightarrow complex set
where

singularities_on A f =
 $\{z \in A. \text{isolated_singularity_at } f z \wedge \text{not_essential } f z \wedge \neg f \text{ analytic_on } \{z\}\}$

lemma *singularities_on_subset*: *singularities_on* A f \subseteq A
by (auto simp: *singularities_on_def*)

lemma *pole_in_singularities_on*:

assumes f meromorphic_on A pts z \in A is_pole f z
shows z \in *singularities_on* A f
unfolding *singularities_on_def* *not_essential_def* **using** *assms*
using *analytic_at_imp_no_pole* meromorphic_on_imp_isolated_singularity **by**
force

lemma *meromorphic_on_subset_pts*:

assumes f meromorphic_on A pts pts' \subseteq pts f analytic_on pts - pts'
shows f meromorphic_on A pts'

proof

show open A pts' \subseteq A
using *assms* **by** (auto simp: *meromorphic_on_def*)
show *isolated_singularity_at* f z *not_essential* f z **if** z \in pts' **for** z
using *assms* **that** **by** (auto simp: *meromorphic_on_def*)
show $\neg z$ *islimpt* pts' **if** z \in A **for** z
using *assms* **that** *islimpt_subset* **unfolding** *meromorphic_on_def* **by** *blast*
have f analytic_on A - pts
using *assms*(1) *meromorphic_imp_analytic* **by** *blast*
with *assms* **have** f analytic_on (A - pts) \cup (pts - pts')
by (*subst analytic_on_Un*) auto
also **have** (A - pts) \cup (pts - pts') = A - pts'
using *assms* **by** (auto simp: *meromorphic_on_def*)
finally **show** f holomorphic_on A - pts'
using *analytic_imp_holomorphic* **by** *blast*

qed

lemma *meromorphic_on_imp_superset_singularities_on*:

assumes f meromorphic_on A pts
shows *singularities_on* A f \subseteq pts

proof

fix z **assume** z \in *singularities_on* A f
hence z \in A \neg f analytic_on {z}
by (auto simp: *singularities_on_def*)
with *assms* **show** z \in pts
by (*meson DiffI meromorphic_on_imp_analytic_at*)

qed

lemma *meromorphic_on_singularities_on*:

assumes f meromorphic_on A pts

```

shows  $f$  meromorphic_on  $A$  (singularities_on  $A$   $f$ )
using assms meromorphic_on_imp_superset_singularities_on[OF assms]
proof (rule meromorphic_on_subset_pts)
  have  $f$  analytic_on  $\{z\}$  if  $z \in pts - singularities\_on\ A$  for  $z$ 
    using that assms by (auto simp: singularities_on_def meromorphic_on_def)
  thus  $f$  analytic_on  $pts - singularities\_on\ A$   $f$ 
    using analytic_on_analytic_at by blast
qed

theorem Residue_theorem_inside:
assumes  $f$ :  $f$  meromorphic_on  $s$  pts
  simply_connected  $s$ 
assumes  $g$ : valid_path  $g$ 
  pathfinish  $g = pathstart\ g$ 
  path_image  $g \subseteq s - pts$ 
defines  $pts1 \equiv pts \cap inside\ (path\_image\ g)$ 
shows finite  $pts1$ 
  and contour_integral  $g\ f = 2 * pi * i * (\sum_{p \in pts1}. winding\_number\ g\ p * residue\ f\ p)$ 
proof -
  note [dest] = valid_path_imp_path
  have  $cl\_g$  [intro]: closed (path_image  $g$ )
    using  $g$  by (auto intro!: closed_path_image)
  have open  $s$ 
    using  $f(1)$  by (auto simp: meromorphic_on_def)
  define  $pts2$  where  $pts2 = pts - pts1$ 

  define  $A$  where  $A = path\_image\ g \cup inside\ (path\_image\ g)$ 
  have closed  $A$ 
    unfolding  $A\_def$  using  $g$  by (intro closed_path_image_Un_inside) auto
  moreover have bounded  $A$ 
    unfolding  $A\_def$  using  $g$  by (auto intro!: bounded_path_image_bounded_inside)
  ultimately have 1: compact  $A$ 
    using compact_eq_bounded_closed by blast
  have 2: open ( $s - pts2$ )
    using  $f$  by (auto intro!: meromorphic_imp_open_diff' [OF  $f(1)$ ] simp:  $pts2\_def$ )
  have 3:  $A \subseteq s - pts2$ 
    unfolding  $A\_def\ pts2\_def\ pts1\_def$ 
    using  $f(2)\ g(3)$  2 subset_simply_connected_imp_inside_subset [of  $s\ path\_image\ g$ ]
    by auto

  obtain  $\varepsilon$  where  $\varepsilon: \varepsilon > 0$  ( $\bigcup_{x \in A}. ball\ x\ \varepsilon \subseteq s - pts2$ )
    using compact_subset_open_imp_ball_epsilon_subset [OF 1 2 3] by blast
  define  $B$  where  $B = (\bigcup_{x \in A}. ball\ x\ \varepsilon)$ 

  have finite ( $A \cap pts$ )
    using 1 3 by (intro meromorphic_compact_finite_pts [OF  $f(1)$ ]) auto
  also have  $A \cap pts = pts1$ 

```

```

  unfolding pts1_def using g by (auto simp: A_def)
  finally show fin: finite pts1 .

  show contour_integral g f = 2 * pi * i * (∑ p∈pts1. winding_number g p *
  residue f p)
  proof (rule Residue_theorem)
    show open B
      by (auto simp: B_def)
    next
    have connected A
      unfolding A_def using g
      by (intro connected_with_inside closed_path_image connected_path_image)
    auto
    hence connected (A ∪ B)
      unfolding B_def using g ⟨ε > 0⟩ f(2)
      by (intro connected_Un_UN connected_path_image valid_path_imp_path)
      (auto simp: simply_connected_imp_connected)
    also have A ∪ B = B
      using ε(1) by (auto simp: B_def)
    finally show connected B .
  next
  have f holomorphic_on (s - pts)
    by (intro meromorphic_imp_holomorphic f)
  moreover have B - pts1 ⊆ s - pts
    using ε unfolding B_def by (auto simp: pts1_def pts2_def)
  ultimately show f holomorphic_on (B - pts1)
    by (rule holomorphic_on_subset)
  next
  have path_image g ⊆ A - pts1
    using g unfolding pts1_def by (auto simp: A_def)
  also have ... ⊆ B - pts1
    unfolding B_def using ε(1) by auto
  finally show path_image g ⊆ B - pts1 .
  next
  show ∀ z. z ∉ B ⟶ winding_number g z = 0
  proof safe
    fix z assume z: z ∉ B
    hence z ∉ A
      using ε(1) by (auto simp: B_def)
    hence z ∈ outside (path_image g)
      unfolding A_def by (simp add: union_with_inside)
    thus winding_number g z = 0
      using g by (intro winding_number_zero_in_outside) auto
  qed
  qed (use g fin in auto)
qed

```

theorem Residue_theorem':
 assumes $f: f$ meromorphic_on s pts

```

      simply_connected s
assumes g: valid_path g
          pathfinish g = pathstart g
          path_image g  $\subseteq$  s - pts
assumes pts': finite pts'
          pts'  $\subseteq$  s
           $\bigwedge z. z \in pts - pts' \implies \text{winding\_number } g \ z = 0$ 
shows contour_integral g f =  $2 * \pi * i * (\sum p \in pts'. \text{winding\_number } g \ p * \text{residue } f \ p)$ 
proof -
  note [dest] = valid_path_imp_path
  define pts1 where pts1 = pts  $\cap$  inside (path_image g)

  have contour_integral g f =  $2 * \pi * i * (\sum p \in pts1. \text{winding\_number } g \ p * \text{residue } f \ p)$ 
  unfolding pts1_def by (intro Residue_theorem_inside[OF f g])
  also have  $(\sum p \in pts1. \text{winding\_number } g \ p * \text{residue } f \ p) =$ 
     $(\sum p \in pts'. \text{winding\_number } g \ p * \text{residue } f \ p)$ 
  proof (intro sum.mono_neutral_cong_refl)
    show finite pts1
    unfolding pts1_def by (intro Residue_theorem_inside[OF f g])
    show finite pts'
    by fact
  next
  fix z assume z: z  $\in$  pts' - pts1
  show winding_number g z * residue f z = 0
  proof (cases z  $\in$  pts)
    case True
    with z have z  $\notin$  path_image g  $\cup$  inside (path_image g)
    using g(3) by (auto simp: pts1_def)
    hence z  $\in$  outside (path_image g)
    by (simp add: union_with_inside)
    hence winding_number g z = 0
    using g by (intro winding_number_zero_in_outside) auto
    thus ?thesis
    by simp
  next
  case False
  with z pts' have z  $\in$  s - pts
  by auto
  with f(1) have f analytic_on {z}
  by (intro meromorphic_on_imp_analytic_at)
  hence residue f z = 0
  using analytic_at_residue_holo by blast
  thus ?thesis
  by simp
qed
next
fix z assume z: z  $\in$  pts1 - pts'

```

```
    hence winding_number g z = 0
      using pts' by (auto simp: pts1_def)
    thus winding_number g z * residue f z = 0
      by simp
  qed
  finally show ?thesis .
qed

end
theory Complex_Analysis
imports
  Residue_Theorem
  Meromorphic
begin

end
```

References

[1]